Algorithms and Datastructures Graphs, Depth-/Breadth-first Search, Graph-Connectivity

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Bioinformatics Group / Department of Computer Science

Algorithms and Datastructures, January 2017

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Feedback

Exercises

Lecture

Graphs

Introduction

Implementation

Application example

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Feedback from the exercises

The upcoming exercise sheet 12 and 13 will be merged together (finding largest connected component + Dijkstra)

Some people were asking for more solution sheets for the exercises

We are working on it.

Feedback from the lecture

Code in the lecture will be a litte bit different from exercise sheet.

One person asked for additional explanations regarding proofs.

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Graphs - Overview:

▶ Besides arrays, lists and trees the most common datastructure (Trees are a special type of graph)

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- Representation of graphs in the computer

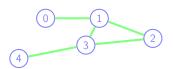
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- Representation of graphs in the computer
- Breadth first search (BFS)

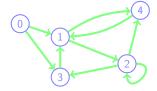
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- Representation of graphs in the computer
- Breadth first search (BFS)
- Depth first search (DFS)
- Connected components of a graph

Introduction

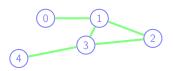
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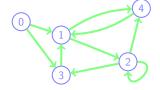




Introduction

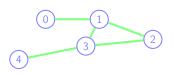
Terminology:

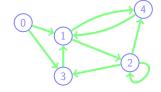




▶ Each Graph G = (V, E) consists of:

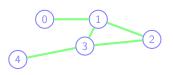
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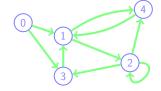




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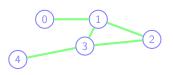
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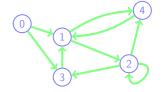




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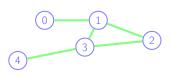
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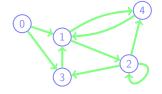




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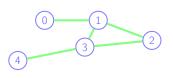
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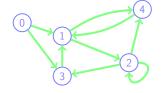




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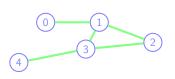
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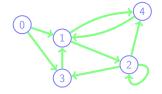




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Introduction





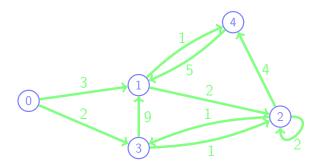
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- ▶ Self-loops are also possible: e = (u, u) or $e = \{u, u\}$

Introduction

Weighted graph:

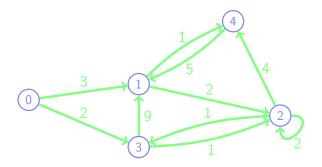
Introduction

Weighted graph:



Introduction

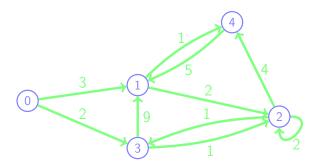
Weighted graph:



▶ Each edge is marked with a real number named weight

Introduction

Weighted graph:



- ► Each edge is marked with a real number named weight
- ► The weight is also named length or cost of the edge depending on the application

Example: Road network

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► Intersections: vertices

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Example: Road network

► Intersections: vertices

► Roads: edges

Travel time: costs of the edges Wengarten C. Freiburg im Principal Court Oberal Wenter State Coorpor

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Figure: Map of Freiburg © OpenStreetMap

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Graphs Implementation

How to represent this graph computationally?

Implementation

How to represent this graph computationally?

► Two classic variants

Implementation

How to represent this graph computationally?

- Two classic variants
 - 1. Adjacency matrix with space consumption $\Theta(|V|^2)$

Implementation

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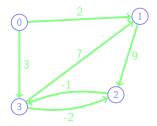


Figure: Weighted graph with

$$|V| = 4$$
, $|E| = 6$

Implementation

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- Two classic variants
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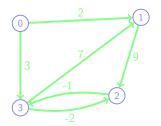


Figure: Weighted graph with |V| = 4, |E| = 6

		ena-vertice			
		0	1	2	3
<u>S</u>	0		2		3
vertice	1			9	
٦.	2				-1
start	3		7	-2	

Figure: Adjacency matrix

Implementation

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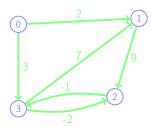


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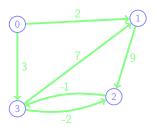


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<u>o</u>	1, 2	3, 3
verti	2, 9	
.5 (2)	3, -1	
start (3)	1, 7	2, -2

Figure: Adjacency list

Implementation

Graph: Arrangement

Implementation

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► Graph is fully defined through the adjacency matrix / list

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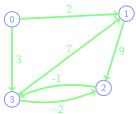


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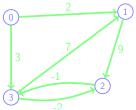


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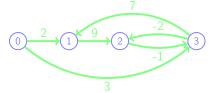


Figure: Same graph ordered by number - outer planar graph

Implementation - Python

```
class Graph:
    def __init__(self):
        self.vertices = []
        self.edges = []
    def addVertice(self, vert):
        self.vertices.append(vert)
    def addEdge(self, fromVert, toVert):
        self.edges.append((fromVert, toVert))
```

Implementation - Python

Degrees (Valency)

Degree of a vertex: Directed graph: G = (V, E)

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Figure: Vertex with in- / outdegree of 3 / 2 $\,$

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$$\deg^+(u) = |\{(v, u) : (v, u) \in E\}|$$

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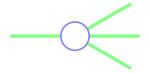


Figure: Vertex with degree of 4

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Paths

Paths

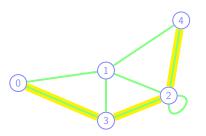


Figure: Undirected path of length 3 P = (0, 3, 2, 4)

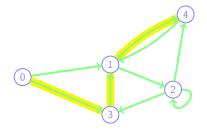


Figure: Directed path of length 3 P = (0, 3, 1, 4)

Paths

Paths in a graph: G = (V, E)

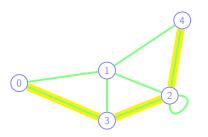


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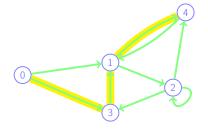


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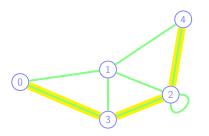


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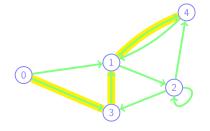


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- ▶ A path of G is a sequence of edges $u_1, u_2, ..., u_i \in V$ with
 - ▶ Undirected graph: $\{u_1, u_2\}, \{u_2, u_3\}, \dots, \{u_{i-1}, u_i\} \in E$
 - ▶ Directed graph: $(u_1, u_2), (u_2, u_3), \dots, (u_{i-1}, u_i) \in E$

Paths

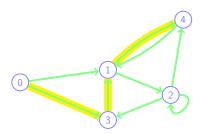


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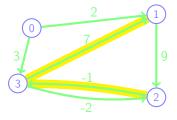


Figure: Weighted path with cost 6 P = (2,3,1)

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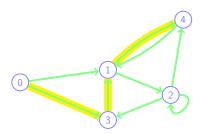


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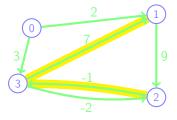


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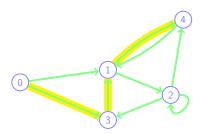


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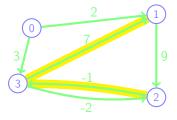


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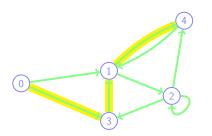


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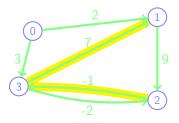


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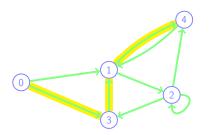


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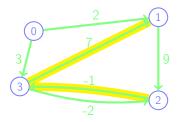


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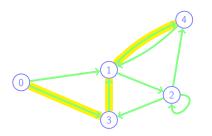


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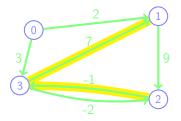


Figure: Weighted path with cost 6 P = (2, 3, 1)

- ► The length of a path is: (also costs of a path)
 - ▶ Without weights: number of edges taken
 - With weights: sum of weigths of edges taken

Paths

Shortest path in a graph: G = (V, E)

Paths

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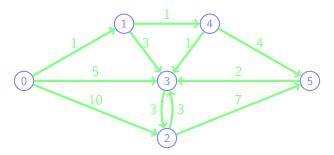


Figure: Shortest path from 0 to 2 with cost / distance d(0,2) = ?

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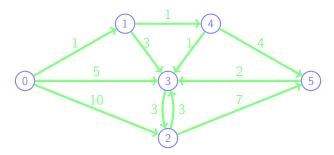


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Paths

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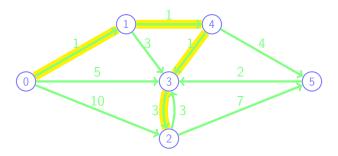


Figure: Shortest path from 0 to 2 with cost / distance d(0,2) = 6P = (0,1,4,3,2)

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Paths

Diameter of a graph: G = (V, E)

Paths

Diameter of a graph:
$$G = (V, E)$$
 $d = \max_{u,v \in V} d(u,v)$

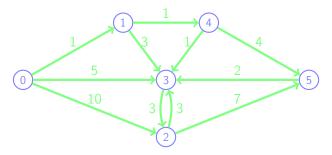


Figure: Diameter of graph is d = ?

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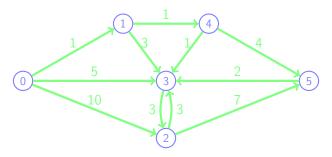


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► The diameter of a graph is the length / the costs of the longest shortest path

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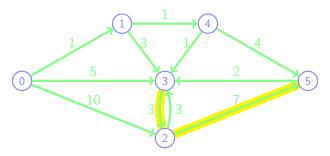


Figure: Diameter of graph is d = 10, P = (3, 2, 5)

► The diameter of a graph is the length / the costs of the longest shortest path

Connected Components

Connected Components

Connected components: G = (V, E)

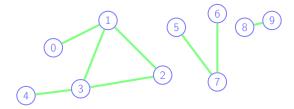


Figure: Three connected components

Undirected graph:

Connected Components

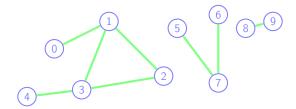


Figure: Three connected components

- Undirected graph:
 - All connected components are a partition of V

$$V = V_1 \cup \cdots \cup V_k$$

Connected Components

Connected components: G = (V, E)

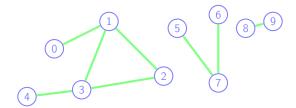


Figure: Three connected components

- Undirected graph:
 - ▶ All connected components are a partition of *V*

$$V = V_1 \cup \cdots \cup V_k$$

► Two vertices *u*, *v* are in the same connected component if a path between *u* and *v* exists

Connected Components

Connected Components

Connected components: G = (V, E)

Directed graph:

Connected Components

- Directed graph:
 - ► Named strongly connected components

Connected Components

- Directed graph:
 - Named strongly connected components
 - Direction of edge has to be regarded

Connected Components

- Directed graph:
 - Named strongly connected components
 - Direction of edge has to be regarded
 - Not part of this lecture

Connected Components - Graph Exploration

Connected Components - Graph Exploration

Graph Exploration: (Informal definition)

▶ Let G = (V, E) be a graph and $s \in V$ a start vertex

Connected Components - Graph Exploration

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- We visit each reachable vertex connected to s

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- ▶ Depth-first search: in sequence of the largest distance to s

Connected Components - Graph Exploration

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- We visit each reachable vertex connected to s
- Breadth-first search: in sequence of the smallest distance to s
- ▶ Depth-first search: in sequence of the largest distance to *s*
- Not a problem on its own but is often used as subroutine of other algorithms

Connected Components - Breadth-First Search

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Idea:

1. We start with all vertices unmarked and mark visited vertices

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- 5. Iteratively mark reachable vertices for all levels
- 6. All connected nodes are now marked and in the same connected component as the start vertex s

Connected Components - Breadth-First Search

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Connected Components - Breadth-First Search

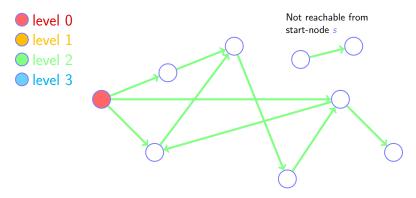


Figure: spanning tree of a breadth-first search

Connected Components - Breadth-First Search

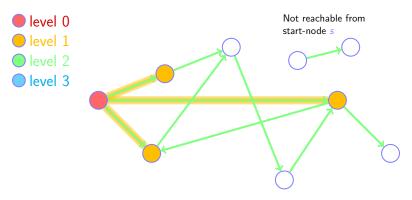


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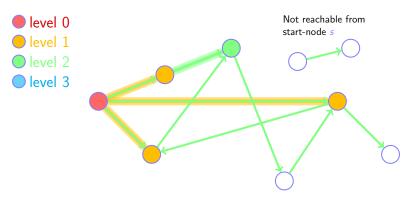


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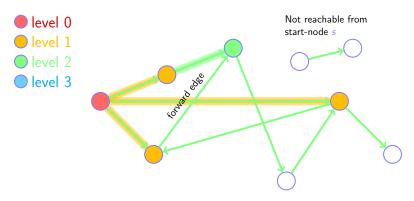


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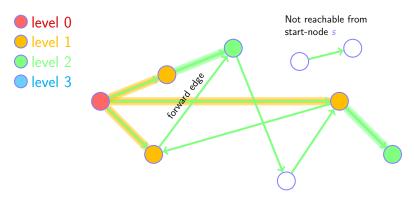


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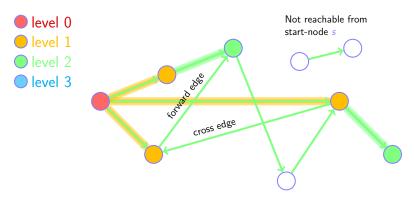


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Connected Components - Breadth-First Search

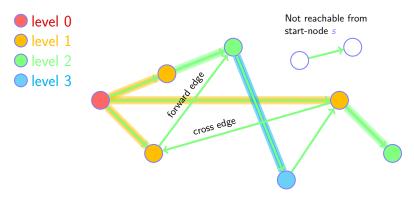


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Connected Components - Breadth-First Search

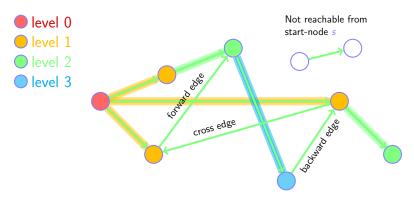


Figure: spanning tree of a breadth-first search

Connected Components - Depth-First Search

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Connected Components - Depth-First Search

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Connected Components - Depth-First Search

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- 1. We start with all vertices unmarked and mark visited vertices
- 2. Mark the start vertex s
- 3. Pick an unmarked connected vertex and start a recursive depth-first search with the vertex as start vertex (continue on step 2)

Connected Components - Depth-First Search

Idea:

- 1. We start with all vertices unmarked and mark visited vertices
- 2. Mark the start vertex s
- Pick an unmarked connected vertex and start a recursive depth-first search with the vertex as start vertex (continue on step 2)
- 4. If no unmarked connected vertex exists go one vertex back (reduce the recursion level by one)

Connected Components - Depth-First Search

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Depth-first search:

► Search starts with long paths (searching with depth)

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- Search starts with long paths (searching with depth)
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- ▶ If the graph is acyclic we get a topological sorting
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 - ▶ The numbers increase with path from the start vertex

Connected Components - Depth-First Search

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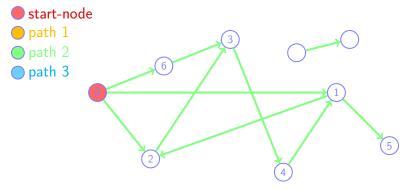


Figure: spanning tree of a depth-first search

Connected Components - Depth-First Search

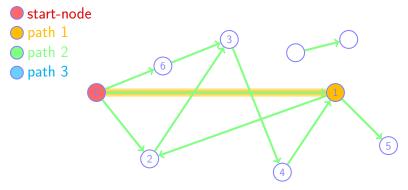


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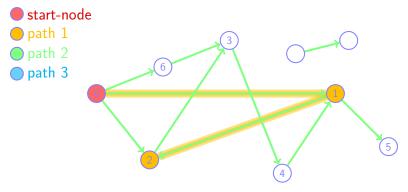


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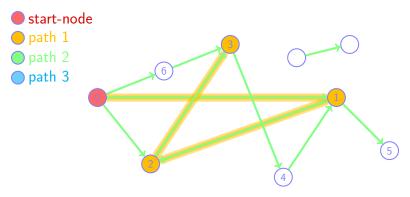


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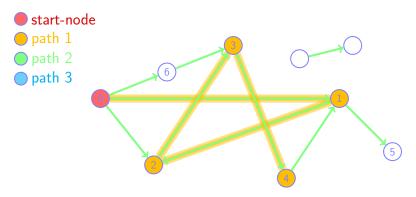


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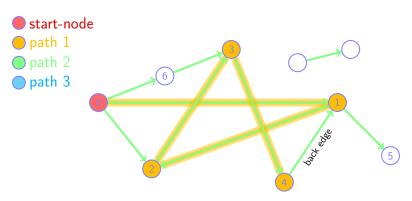


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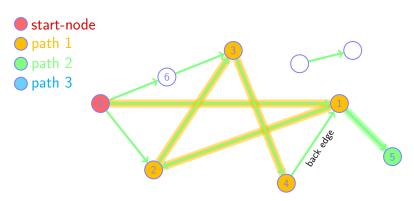


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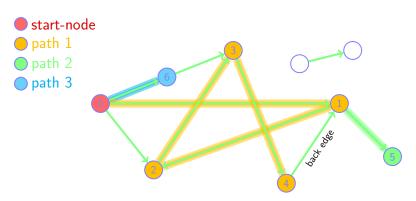


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Connected Components - Depth-First Search

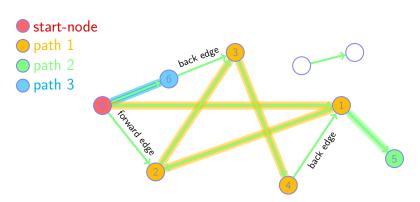


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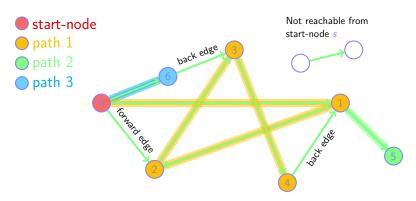


Figure: spanning tree of a depth-first search

Why is this called Breadth - and Depth First Search?

Connected Components - Breadth-/Depth-First Search

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Runtime complexity:

Constant costs for each visited vertex and edge

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Connected Components - Breadth-/Depth-First Search

- Constant costs for each visited vertex and edge
- We get a runtime complexity of $\Theta(|V'| + |E'|)$
- Let V' and E' be the reachable vertices and edges
- ► All vertices of V' are in the same connected component as our start vertex s
- This can only be improved by a constant factor

Structure

Feedback

Exercises

Lecture

Graphs

Introduction

Implementation

Application example

Image processing

Image processing

Connected component labeling

Image processing

- Connected component labeling
- Counting of objects in an image

Image processing

- Connected component labeling
- ► Counting of objects in an image

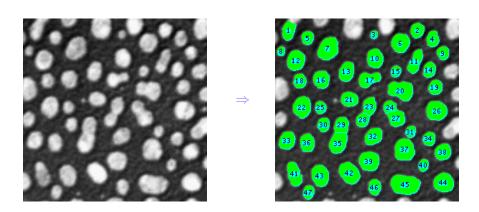


Image processing

What's object, what's background?

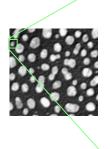


Image processing

Convert to black white using threshold:

value = 255 if value > 100 else 0

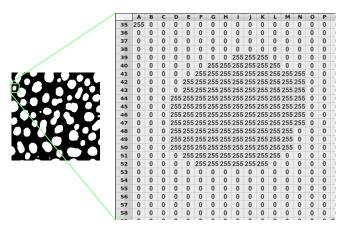


Image processing

Interpret image as graph:

Image processing

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Image processing

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- Edges between adjacent pixels (normally 4 or 8 neighbors)
- Edges are not saved externally, algorithm works directly on array
- Breadth- / depth-first search find all connected components (particles)

Image processing

Image processing

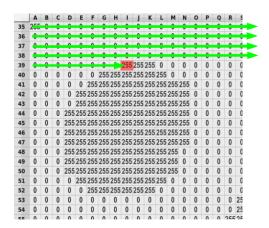
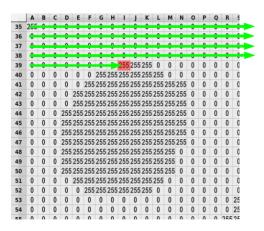


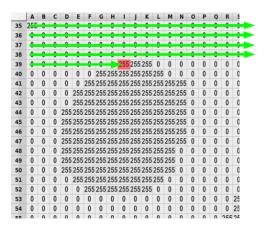
Image processing

Find connected components:



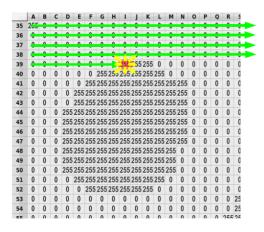
 Search pixel-by-pixel for non-zero intensity

Image processing



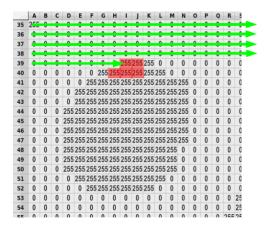
- Search pixel-by-pixel for non-zero intensity
- ► Label found pixel as component 1

Image processing



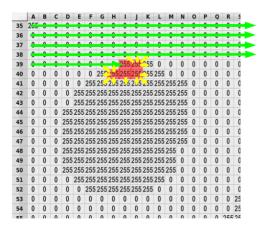
- ► Search pixel-by-pixel for non-zero intensity
- ► Label found pixel as component 1
- Check neighbors of all new labeled pixels

Image processing



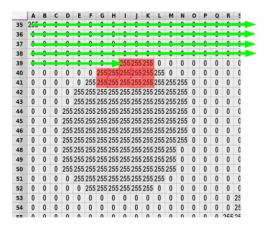
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Image processing



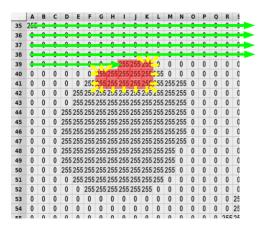
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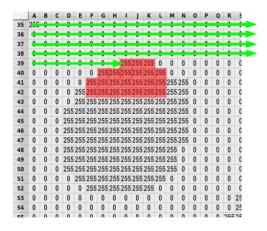
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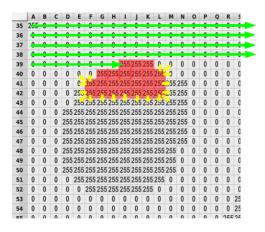
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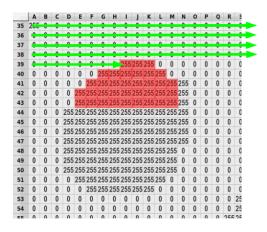
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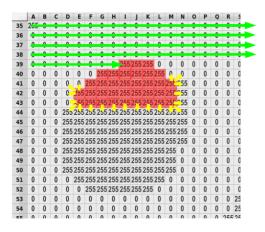
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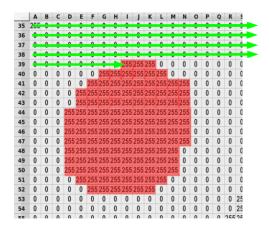
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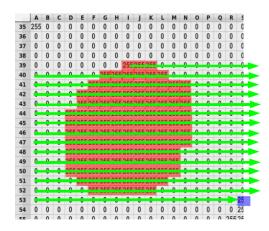
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Image processing



- Search pixel-by-pixel for non-zero intensity
- ► Label found pixel as component 2
- **.** . .

Image processing

Result of connected component labeling:

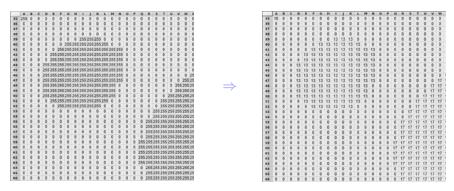


Figure: Result: particle indices instead of intensities

Further Literature

General

- [CRL01] Thomas H. Cormen, Ronald L. Rivest, and Charles E. Leiserson. Introduction to Algorithms. MIT Press, Cambridge, Mass, 2001.
- [MS08] Kurt Mehlhorn and Peter Sanders.
 Algorithms and data structures, 2008.
 https://people.mpi-inf.mpg.de/~mehlhorn/
 ftp/Mehlhorn-Sanders-Toolbox.pdf.

Further Literature

Graph-Search

```
[Wika] Breadth-first search
          https://en.wikipedia.org/wiki/
          Breadth-first_search
[Wikb] Depth-first search
          https:
          //en.wikipedia.org/wiki/Depth-first_search
```

Graph-Connectivity

```
[Wik] Connectivity (graph theory)
    https://en.wikipedia.org/wiki/Connectivity_
        (graph_theory)
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