Algorithms and Datastructures

Levenshtein distance, Dynamic programming

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Structure

Introduction

Edit distance:

- Measurement for similarity of two words / strings
- Algorithm for efficient calculation
- General principle: dynamic programming

Introduction

Motivation: Error tolerant string comparison

BioInfSearch

ejafjatlajökuk

eyjafjallajökull
eyjafjallajökull movie
eyjafjallajälull trailer

Search!

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Wikipedia.org:

"Der Eyjafjallajökull ([ˈeɪjaˌfjatlaˌjœːkʏtl̥])[3], auf Deutsch Eyjafjöll-Gletscher, ist der sechstgrößte Gletscher Islands.

Er liegt an der äußersten Südküste, westlich des Gletschers Mýrdalsjökull in der Gemeinde Rangárþing eystra, die größte Höhe beträgt 1651 m. Unter dem Gletscher befindet sich der Vulkan Eyjafjöll mit eigener Magmakammer, der seit der Besiedelung von Island in den Jahren 920, 1612 (oder 1613), 1821 bis 1823 und zuletzt im Jahr 2010 aktiv war."

A lot of applications where similar string are searched:

- Duplicates in databases:

```
Hein Blöd   27568 Bremerhaven
Hein Bloed  27568 Bremerhafen
Hein Doof   27478 Cuxhaven
```

- Product search:

```
memory stik
```

- Websearch:

```
eyjaföllajaküll
uniwersität verien 2017
```

- Bioinformatics: Similarity of DNA-sequences

Search of similar proteins:

- BLAST (**B**asic **L**ocal **A**lignment **S**earch **T**ool)
- Alignment $\widehat{=}$ Edit distance
- Changed life-science completely
- Cited 63437 times on Google Scholar (Sep. 2017)

Edit distance

Definition of edit distance: (*Levenshtein-distance*)

- Let x, y be two strings
- Edit distance $\mathrm{ED}(x, y)$ of x and y:
  The minimal number of operations to transform x into y
  - `Insert` a character
  - `Replace` a character with another
  - `Delete` a character

Edit distance

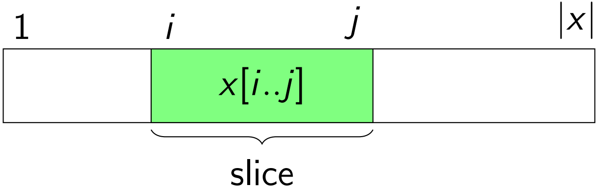
Example

$$\begin{array}{ll}
\mathsf{1\,2\,3\,4\,5} & \\
\mathsf{DOOF} & \\
\downarrow & \texttt{replace(1, B)} \\
\mathsf{BOOF} & \\
\downarrow & \texttt{replace(2, L)} \\
\mathsf{BLOF} & \\
\downarrow & \texttt{insert(4, E)} \\
\mathsf{BLOEF} & \\
\downarrow & \texttt{replace(5, D)} \\
\mathsf{BLOED} & \\
 & \underbrace{\hphantom{\texttt{replace(5, D)}}}_{\text{ED}=4}
\end{array}
\qquad\qquad
\begin{array}{ll}
\mathsf{1\,2\,3\,4\,5} & \\
\mathsf{BLOED} & \\
\downarrow & \texttt{replace(5, F)} \\
\mathsf{BLOEF} & \\
\downarrow & \texttt{delete(4)} \\
\mathsf{BLOF} & \\
\downarrow & \texttt{replace(2, O)} \\
\mathsf{BOOF} & \\
\downarrow & \texttt{replace(1, D)} \\
\mathsf{DOOF} & \\
 & \underbrace{\hphantom{\texttt{replace(1, D)}}}_{\text{ED}=4}
\end{array}$$

Edit distance

Notation:

- ε is the empty string
- $|x|$ is the length of the string x (number of characters)
- $x[i..j]$ is the slice of x from i to j where $1 \leq i \leq j \leq |x|$



Edit distance

Trivial facts:

- $\mathrm{ED}(x, y) = \mathrm{ED}(y, x)$
- $\mathrm{ED}(x, \epsilon) = |x|$
- $\mathrm{ED}(x, y) \geq \mathrm{abs}(|x| - |y|)$ $\quad\quad \mathrm{abs}(x) = \begin{cases} x & \text{if } x \geq 0 \\ -x & \text{else} \end{cases}$
- $\mathrm{ED}(x, y) \leq \mathrm{ED}(x[1..n-1], y[1..m-1]) + 1 \;\; n = |x|, \;\; m = |y|$

Edit distance

Solving examples

Solutions based on examples:

- From `VERIEN` to `FERIEN`?
- From `MEXIKO` to `AMERIKA`?
- From `AAEBEAABEAREEAEBA` to `RBEAAEEBAAAEBBAEAE`?
- Searching biggest substrings can yield the solution but doesn't have to

Recursive approach:

- Dividing in two halves? Not a good idea:

  $\mathrm{ED}(GRAU, RAUM) = 2$ but
  $\mathrm{ED}(GR, RA) + \mathrm{ED}(AU, UM) = 4$

- Finding "smaller" sub problems?
  Let's try it!

Edit distance

Terminology:

- Let x, y be two strings
- Let $\sigma_1, \ldots, \sigma_k$ be a sequence of k operations where $k = \mathrm{ED}(x, y)$ for $x \rightarrow y$ (transform x into y)
  (We do not know this sequence but we assume it exists)

Edit distance

Terminology:

- We only consider monotonous sequences:
  The positition of σ_{i+1} is $\geq$ the position of σ_i where we only allow the positions to be equal on a `delete` operation

1 2 3 4 5		1 2 3 4 5 6 7	
DOOF		SAUDOOF	
$\downarrow$	`replace(1, B)`	$\downarrow$	`delete(1)`
BOOF		AUDOOF	
$\downarrow$	`replace(2, L)`	$\downarrow$	`delete(1)`
BLOF		UDOOF	
$\downarrow$	insert(4, E)	$\downarrow$	`delete(1)`
BLOEF		DOOF	
$\downarrow$	`replace(5, D)`	$\downarrow$	`insert(4, O)`
BLOED		DOOOF	

Edit distance

Terminology:

- **Lemma:** For any x and y with $k = \mathrm{ED}(x, y)$ exists a monotonous sequence of k operations for $x \rightarrow y$
- **Intuition:** The order of our sequence is not relevant (Therefore we can also sort them monotonously)

1	2	3	4	5
D	O	O	F	
B	L	O	E	D

1	2	3	4	5	6	7
S	A	U	D	O	O	F
D	O	O	O	F		

Edit distance

Recursive approach

Consider the last operation:

- Solve blue part recursively

```
DOOF
↓↓↓↓
BLOE
    ↓insert
BLOED
```

Figure: Case 1a

```
  DOOF
↓↓↓↓↓↓
BLOEDF
     ↓delete
BLOED
```

Figure: Case 1b

```
 DOOF
↓↓↓↓↓
BLOEF
    ↓replace
BLOED
```

Figure: Case 1c

Edit distance

Recursive approach

Consider the last operation:

- Solve blue part recursively

W I N T E R
↓ ↓ ↓ ↓ ↓ ↓
S O M M E R
↓`nothing`
S O M M E R

Figure: Case 2

Display of solution:

- Alignment
- Example:

_	_	_	B	L	O	E	D
S	A	U	B	L	O	E	D

Edit distance

Dynamic programming

Dynamic programming:

- Instances of Bellman's principle of optimality:
  - Shortest paths
  - Edit distance

Figure: Richard Bellman (1920 - 1984)

- Optimal solutions consist of optimal partial solutions
  - Shortest paths: Each partial path has to be optimal
  - Edit distance: Each partial alignment has to be optimal

_	_	_	B	L	O	E	D	E	R
S	A	U	B	L	O	E	D	_	_

- Always solvable through dynamic programming (Caching of optimal partial solutions)

Edit distance

Case analysis:

- We consider the last operation σ_k
  - $\sigma_1, \ldots, \sigma_{k-1}$: $x \to z$ and σ_k: $z \to y$
    Example:

$$x = \mathrm{DOOF},\ z = \mathrm{SAUBLOEF},\ y = \mathrm{SAUBLOED}$$

  - Let $n = |x|,\ m = |y|,\ m' = |z|$
  - We note $m' \in \{m-1, m, m+1\}$ why?

Edit distance

Case analysis:

- Case 1: σ_k does something at the outer end:
  - Case 1a: $\sigma_k = \texttt{insert}(m'+1, y[m])$ [then $m' = m-1$]
  - Case 1b: $\sigma_k = \texttt{delete}(m')$ [then $m' = m+1$]
  - Case 1c: $\sigma_k = \texttt{replace}(m', y[m])$ [then $m' = m$]
- Case 2: σ_k does nothing at the outer end:
  - Then $z[m'] = y[m]$ and $x[n'] = z[m']$ and with that $\sigma_1, \ldots, \sigma_{k-1}$: $x[1..n-1] \to y[1..m-1]$ and $x[n] = y[m]$

Edit distance

Case analysis:

- Case 1a (`insert`): $\sigma_1, \ldots, \sigma_{k-1}$: $x \rightarrow y[1..m-1]$
- Case 1b (`delete`): $\sigma_1, \ldots, \sigma_{k-1}$: $x[1..n-1] \rightarrow y$
- Case 1c (`replace`): $\sigma_1, \ldots, \sigma_{k-1}$: $x[1..n-1] \rightarrow y[1..m-1]$
- Case 2 (`nothing`): $\sigma_1, \ldots, \sigma_k$: $x[1..n-1] \rightarrow y[1..m-1]$

This results in the recursive formula:

- For $|x| > 0$ and $|y| > 0$ is $\mathrm{ED}(x, y)$ the minimum of
  - $\mathrm{ED}(x, y[1..m-1]) + 1$ and
  - $\mathrm{ED}(x[1..n-1], y) + 1$ and
  - $\mathrm{ED}(x[1..n-1], y[1..m-1]) + 1$ if $x[n] \neq y[m]$
  - $\mathrm{ED}(x[1..n-1], y[1..m-1]) + 0$ if $x[n] = y[m]$
- For $|x| = 0$ is $\mathrm{ED}(x, y) = |y|$
- For $|y| = 0$ is $\mathrm{ED}(x, y) = |x|$

Edit distance

Implementation - Python

```python
def edit_distance(x, y):
    if len(x) == 0:
        return len(y)
    if len(y) == 0:
        return len(x)

    ed1 = edit_distance(x, y[:-1]) + 1
    ed2 = edit_distance(x[:-1], y) + 1
    ed3 = edit_distance(x[:-1], y[:-1])
    if x[-1] != y[-1]:
        ed3 += 1

    return min(ed1, ed2, ed3)
```

Edit distance

Runtime analysis

Recursive program:

- The algorithm results in the following recursive formular:

$$
\begin{aligned}
T(n,m) &= T(n-1,m) + T(n,m-1) + T(n-1,m-1) + 1 \\
&\geq T(n-1,m-1) + T(n-1,m-1) + T(n-1,m-1) \\
&= 3 \cdot T(n-1,m-1)
\end{aligned}
$$

- This results in $T(n,n) \geq 3^n$

$\Rightarrow$ The runtime is at least exponential

Edit distance

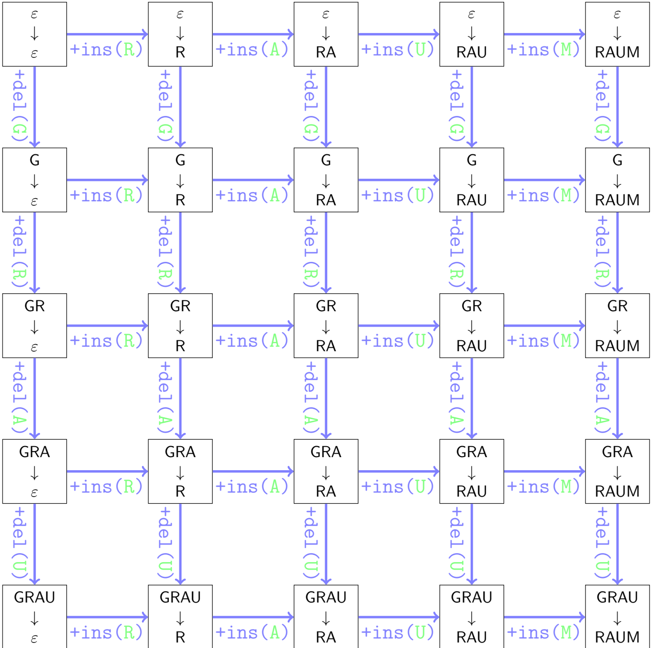
Dynamic programming:

- We create a table with all possible combination of substrings and save calculated entries
- This results in a runtime and space consumption of $O(n \cdot m)$

Visualization on the next slide:

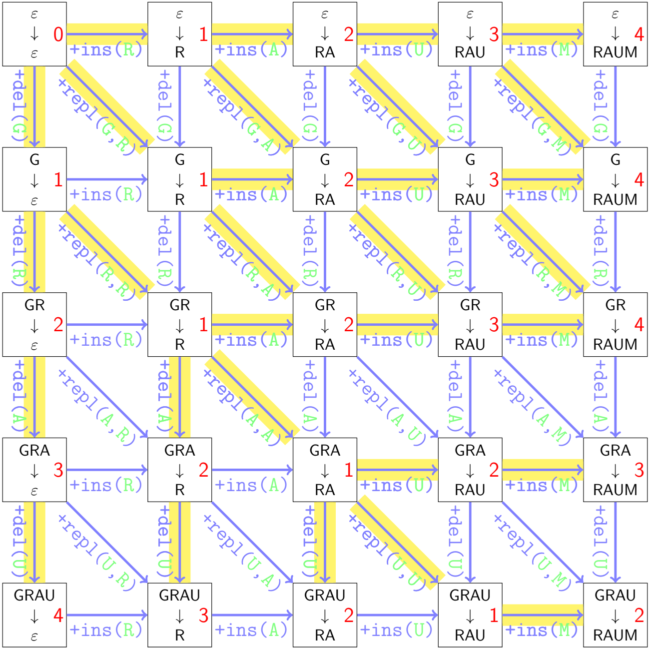
- Operations always refer to the last position (indices are omitted)
- We also display the replaced character on a `replace` operation to visualize operations without costs $\Rightarrow$ `repl(A, A)`

Edit Distance



Fast algorithm:
We can determine the edit distance for all combination of partial strings from the top left to bottom right.

Edit Distance

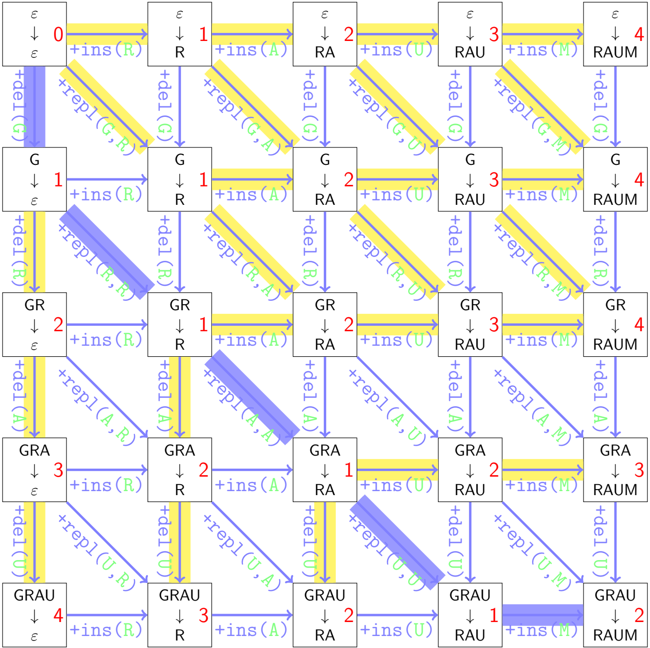


Edit distance

How to get the sequence of operations?

- We save at each recursion the most efficient previous entry (the highlighted arrows in our image)
- There can be more than one arrows to the three previous entries
- If we follow the highlighted path from (n, m) to $(1, 1)$ we get the optimum operations to transform x into y
  - If we can follow more than one path there exist more than one ideal sequence

Edit Distance



Edit distance

General principle:

- Recursive computation of ...
  - ... the same reoccuring partial problems
  - ... a limited number of partial problems
- Computation of the solutions for all partial problems
- In a order that unsolved partial problems consist of already solved partial problems
- The "path" to our solution normally gets computed while searching the best solution
- Dijkstra algorithm is basically dynamic programming!

Edit distance

Additional applications (I)

Additional applications:

- *Edit distance*: global alignment with $O(n^2)$ space and time consumption
- But: Model for deletition unrealistic
  - In evolution larger pieces are more likely
  - `delete` operation: first gap expensive (e.g. 2), remainding are cheaper (e.g. 0.5)

_	_	_	B	L	O	E	D
S	A	U	B	L	O	E	D

  - Solution in $O(n^3)$ time or $O(n^2)$ affine

Edit distance

Additional applications (II)

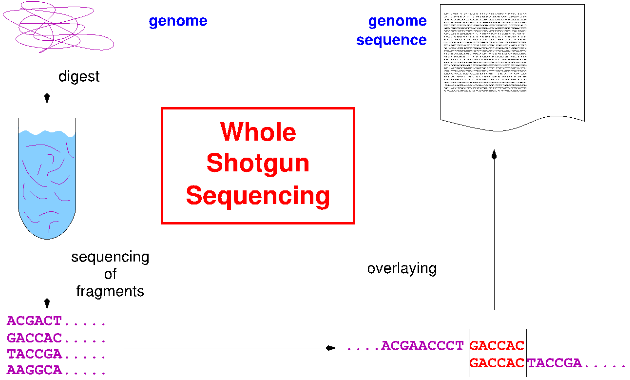
$O(n^2)$ space consumption might be problematic:

Hirschberg algorithm:

- Divide-and-conquer approach
- $O(n)$ space and $O(n^2)$ time consumption

Edit distance

Additional applications (III)



- Sequencing: $O(n^2)$ is too much
- Index: suffixtree, suffixarray, burrow-wheeler-transform

Further Literature

- **General**

[CRL01] Thomas H. Cormen, Ronald L. Rivest, and Charles E. Leiserson.
Introduction to Algorithms.
MIT Press, Cambridge, Mass, 2001.

[MS08] Kurt Mehlhorn and Peter Sanders.
Algorithms and data structures, 2008.
https://people.mpi-inf.mpg.de/~mehlhorn/ftp/Mehlhorn-Sanders-Toolbox.pdf.

Further Literature

- **Dynamic programming**

  [Wik] Dynamic programming
  https://en.wikipedia.org/wiki/Dynamic_programming

- **Edit distance**

  [Wik] Levenshtein distance
  https://en.wikipedia.org/wiki/Levenshtein_distance