

1 Preliminaries and Notation

One such algorithm is the alternating direction method of multipliers [?], (ADMM). This algorithm solves problems of the form

$$\begin{aligned} \arg \min_x f(x) + g(z) \\ \text{s.t. } Ux + Vz = c \end{aligned} \quad (1.0.1)$$

where f and g are assumed to be closed convex functions with range in \mathbb{R} , $U \in \mathbb{R}^{p \times n}$ and $V \in \mathbb{R}^{p \times m}$ are matrices (not assumed to have full rank), and $c \in \mathbb{R}^p$.

ADMM consists of iteratively minimising the augmented Lagrangian

$$L_p(x, z, \eta) = f(x) + g(z) + \eta^T (Ux + Vz - c) + \frac{\rho}{2} \|Ux + Vz - c\|_2^2$$

(η is a Lagrange multiplier), and ρ is a parameter we can choose to make $g(z)$ smooth [?], with the following iterations:

$$x^{k+1} := \arg \min_x L_\rho(x, z^k, \eta^k) \quad (1.0.2)$$

$$z^{k+1} := \arg \min_z L_\rho(x^{k+1}, z, \eta^k) \quad (1.0.3)$$

$$\eta^{k+1} := \eta^k + \rho (Ux^{k+1} + Vz^{k+1} - c) \quad (1.0.4)$$

The alternating minimisation works because of the decomposability of the objective function: the x minimisation step is independent of the z minimisation step and vice versa.

We will work frequently with the *dual* of problem (1.0.1) which is

$$D(\lambda) = \arg \max_{\lambda} -F^*(A^T x) + \langle \lambda, c \rangle - G^*(B^T \lambda) \quad (1.0.5)$$

where f^* is the convex conjugate:

$$F^*(p) = \sup_p \langle x, p \rangle - f(x) \quad (1.0.6)$$

For example the convex conjugate of the l_1 norm $g(x) = |x|$ is:

$$f^*(x^*) = \begin{cases} 0 & \text{if } |x| \leq 1 \\ \infty & \text{if } |x| > 1 \end{cases} \quad (1.0.7)$$

One way to measure how far the iterations of ADMM are from the optimal solution is to define the primal and dual residuals:

$$r_k = c - Ax_k - Bz_k \quad (1.0.8)$$

$$d_k = \rho A^T B (x_k - x_{k-1}) \quad (1.0.9)$$

2 Convergence of ADMM

Theorem 2.1. *Consider the ADMM iterations defined by (1.0.1). Supposing that both f, g are strongly convex and that:*

$$\rho^3 \leq \frac{\sigma_f \sigma_g}{\tau(A^T A) \tau(B^T B)} \quad (2.0.10)$$

Then the sequence of iterates $\{\eta_k\}$ satisfy:

$$D(\eta^*) - D(\eta) \leq \frac{\|\eta^* - \eta_1\|}{2\rho(k-1)} \quad (2.0.11)$$

Before we prove theorem 2.1 we need a couple of lemmas.

Definition 2.2.

$$\Psi(\lambda) = A\nabla F^*(A^T\lambda)$$

$$\Phi(\lambda) = B\nabla G^*(B^T\lambda)$$

Note that the derivative of $D(\lambda)$ is $b - \Phi - \Psi$, and maximising the dual problem is equivalent to finding a solution to $c \in \Phi(\lambda^*) + \Psi(\lambda^*)$. Since f, g are both strongly convex, then Ψ, Φ are both Lipchitz continuous with $L_\Psi \leq \frac{\tau(A^T A)}{\sigma_f}$ and $L_\Phi \leq \frac{\tau(B^T B)}{\sigma_g}$

Lemma 2.3. Let $\lambda, z \in \mathbb{R}^n$, and define

$$\lambda^{1/2} = \lambda + \rho(c - Ax^{k+1} - Bz^k) \quad (2.0.12)$$

$$\lambda^+ = \lambda + (c - Ax^{k+1} - Bz^{k+1}) \quad (2.0.13)$$

Then we have

$$Ax^{k+1} = A\nabla F^*(A^T\lambda^{1/2}) := \Psi(\lambda^{1/2}) \quad (2.0.14)$$

$$Bz^{k+1} = B\nabla G^*(B^T\lambda^+) := \Phi(\lambda^+) \quad (2.0.15)$$

Lemma 2.4. Suppose that

$$\rho^3 \leq \frac{\sigma_f \sigma_g}{\tau(A^T A) \tau(B^T B)} \quad (2.0.16)$$

and that $Bz = \Phi(\lambda)$. Then for any $\gamma \in \mathbb{R}^n$:

$$D(\lambda^+) - D(\gamma) \geq \rho^{-1} \langle \gamma - \lambda, \lambda - \lambda^+ \rangle + \frac{1}{2\rho} \|\lambda - \lambda^+\|^2 \quad (2.0.17)$$

Proof

Theorem (2.1). □