

MATEMÁTICA DISCRETA

GRAFO DIRIGIDO

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1. Consigna

1. Especificar un grafo dirigido y no ponderado.
2. Realizar y probar la implementación de un grafo dirigido y no ponderado con lista de aristas. Calcular el orden de cada operación.
3. Para probar el buen funcionamiento de la clase anterior:
 - (I) Hacer un método que permita cargar datos a un grafo (este método debe permita cargar los valores de los nodos y las aristas).
 - (II) Hacer un método que genere un grafo en forma aleatoria.
 - (III) Hacer un método que muestre por pantalla al grafo (puede ser un listado con los valores de los nodos y otro con las aristas o bien el dibujo del grafo).
4. Implementar los tres algoritmos que permiten recorrer un grafo (búsqueda plana, BFS y DFS)
5. Modificando adecuadamente los métodos anteriores, escribir y probar algoritmos que:
 - (I) Obtengan la cantidad de vértices fuentes y la cantidad de sumideros.
 - (II) Verificar si es débilmente conexo.
 - (III) Dados dos vértices verificar si existe un camino de longitud 2 entre ambos.
 - (IV) Implementar el algoritmo de Warshall.

2. Especificación

2.1. DirectedGraph

Description: Represents a directed non-weighted graph.

Digraph: Max Order \rightarrow Digraph

Order: $O(1)$

Description: Creates a new empty digraph with the given max order limit. If no max order is given, the graph has unlimited max order.

Precondition: Max Order > 0

Postcondition: A Digraph is created.

Classification: Constructor

add_vertex: Digraph X Data \rightarrow void

Order: $O(1)$

Description: Adds a new vertex to the graph and stores the given data.

Precondition:

- Digraph must exist.
- Data must exist.

Postcondition: New vertex added to the given graph, storing the given data.

Classification: Modifier

remove_vertex: Digraph X Key \rightarrow Data

Order: $O(n)$

Description: Finds the vertex referenced by the given key in the digraph and removes it. Returns the data stored in the vertex.

Precondition:

- Digraph must exist.
- Key must reference an existing vertex in the digraph.

Postcondition: Vertex removed, and its data returned

Classification: Modifier

key_of: Digraph X Data \rightarrow Key

Order: $O(n)$

Description: Finds the key referencing the vertex storing the given data in the graph.

Precondition:

- Digraph must exist.
- The data must be stored in the digraph.

Postcondition: Key referencing the vertex storing the given data is returned.

Classification: Analyzer

add_edge: Digraph X From Key X To Key \rightarrow void

Order: $O(n)$

Description: Adds a new edge connecting the vertices referenced by the given keys. The edge goes from the vertex referenced by From Key to the vertex referenced by To Key.

Precondition:

- Digraph must exist.
- From Key and To Key must reference vertices in the digraph.

Postcondition: New edge added to the graph, that goes from the vertex referenced by From Key to the vertex referenced by To Key.

Classification: Modifier

remove_edge: Digraph X From Key X To Key \rightarrow void

Order: $O(n)$

Description: Finds the edge connecting the vertices referenced by From Key and To Key and removes it.

Precondition:

- Digraph must exist.
- From Key and To Key must reference a existing vertices in the digraph.
- An edge going from the vertex referenced by From Key to the vertex referenced by To Key must exist.

Postcondition: Edge removed

Classification: Modifier

get_vertex: Digraph X Key \rightarrow Data

Order: $O(1)$

Description: Finds the vertex referenced by Key in the digraph and returns the data stored in it

Precondition:

- Digraph must exist.
- Key must reference an existing vertex in the digraph.

Postcondition: Return the data stored in the vertex referenced by the given key.

Classification: Analyzer

get_adjacency_list: Digraph X Key \rightarrow Key List

Order: $O(n)$

Description: Returns the list of all vertices adjacent to the one referenced by the given key in the graph.

Precondition:

- Digraph must exist.
- Key must reference an existing vertex in the digraph.

Postcondition: Return a list of keys referencing all adjacent vertices to the one referenced by the given key.

Classification: Analyzer

edge_exists: Digraph X From Key X To Key \rightarrow Boolean

Order: $O(n)$

Description: Verifies whether an edge from the vertex referenced by From Key to the vertex referenced by To Key exists.

Precondition:

- Digraph must exist.
- From Key and To Key must reference a existing vertices in the digraph.

Postcondition:

- True if the edge exists.
- False if the edge does not exist.

Classification: Analyzer

random_digraph: Vertex Amount X Edge Amount \rightarrow Digraph

Order: $O(n)$

Description: Creates a random digraph with the given amount of vertices and edges.

Precondition:

- Vertex Amount > 0
- Edge Amount > 0

Postcondition: A random digraph with the given amount of vertices and edges is created.

Classification: Constructor.

2.2. Graph Related Algorithms

Description: Algorithms related to graph theory

plain_search: Graph \rightarrow Key List

Order: $O(n)$

Description: Given a graph returns a list of keys referencing all its vertices in no particular order.

Precondition: Graph must exist

Postcondition: Return a list of keys referencing all vertices in the graph.

Classification: Analyzer

dfs: Graph \rightarrow Key List

Order: $O(n^2)$

Description: Given a graph returns a list of keys referencing all its vertices using Depth First Search, starting from the vertex first inserted in the graph.

Precondition: Graph must exist

Postcondition: Return a list of keys referencing all vertices in the graph using Depth First Search.

Classification: Analyzer

bfs: Graph \rightarrow Key List

Order: $O(n^2)$

Description: Given a graph returns a list of keys referencing all its vertices using Breath First Search, starting from the vertex first inserted in the graph.

Precondition: Graph must exist

Postcondition: Return a list of keys referencing all vertices in the graph using Breath First Search.

Classification: Analyzer

number_of_sources: Graph \rightarrow Number

Order: $O(n)$

Description: Given a graph returns the number of source vertices.

Precondition: Graph must exist

Postcondition: Return the amount of source vertices in the graph

Classification: Analyzer

number_of_sinks: Graph \rightarrow Number

Order: $O(n)$

Description: Given a graph returns the number of sink vertices.

Precondition: Graph must exist

Postcondition: Return the amount of sink vertices in the graph

Classification: Analyzer

is_weakly_connected: Graph \rightarrow Boolean

Order: $O(n^2)$

Description: Given a graph, checks whether it is weakly connected.

Precondition: Graph must exist

Postcondition:

- True if the graph is weakly connected.
- False if the graph is not weakly connected.

Classification: Analyzer

warshall: Graph \rightarrow Transitive Closure

Order: $O(n^3)$

Description: Given a graph calculates its transitive closure using Warshall's.

Precondition: Graph must exist

Postcondition: The transitive closure must have the capability of calculating whether there is a path between two vertices in $O(1)$ time.

Classification: Analyzer

save: Graph X Name X View X Format \rightarrow void

Order: $O(n)$

Description: Given a graph, stores its graphic representation in the given format with the given name. A view of the file is given if View is True.

Precondition:

- Graph must exist
- Name must be a valid file name
- If given, view must be True or False (default False)
- If given, format must be a valid file format for the Graphviz program (default 'pdf').

Postcondition:

- A file containing DOT code for Graphviz and an image in the given format are created.
- If View is True, the created image is displayed.

Classification: Analyzer

path_exists: Graph x From Key x To Key X Length X Transitive Closure \rightarrow Boolean

Order: $O(n^2)$

Description: Checks if a path between the vertices referenced by From Key and To Key exist.

- A length or a transitive closure can be given, not both.
- If a length is given, the path can be no longer than the length.
- If a Transitive Closure is given the path is checked in it, achieving $O(1)$ performance.

Precondition:

- Digraph must exist.
- From Key and To Key must reference a existing vertices in the digraph.
- If given, length > 0 .

Postcondition:

- True if a path exists.
- False if a path does not exist.

Classification: Analyzer