MATEMÁTICA DISCRETA

GRAFO DIRIGIDO

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1. Consigna

- 1. Especificar un grafo dirigido y no ponderado.
- 2. Realizar y probar la implementación de un grafo dirigido y no ponderado con lista de aristas. Calcular el orden de cada operación.
- 3. Para probar el buen funcionamiento de la clase anterior:
 - (I) Hacer un método que permita cargar datos a un grafo (este método debe permita cargar los valores de los nodos y las aristas).
 - (II) Hacer un método que genere un grafo en forma aleatoria.
 - (III) Hacer un método que muestre por pantalla al grafo (puede ser un listado con los valores de los nodos y otro con las aristas o bien el dibujo del grafo).
- 4. Implementar los tres algoritmos que permiten recorrer un grafo (búsqueda plana, BFS y DFS)
- 5. Modificando adecuadamente los métodos anteriores, escribir y probar algoritmos que:
 - (I) Obtengan la cantidad de vértices fuentes y la cantidad de sumideros.
 - (II) Verificar si es débilmente conexo.
 - (III) Dados dos vértices verificar si existe un camino de longitud 2 entre ambos.
 - (IV) Implementar el algoritmo de Warshall.

2. Especificación

2.1. DirectedGraph

Description: Represents a directed non-weighted graph.

Digraph: Max Order \rightarrow Digraph

Order: O(1)

Description: Creates a new empty digraph with the given max order limit. If no max order

is given, the graph has unlimited max order.

Precondition: Max Order > 0 Postcondition: A Digraph is created.

Classification: Constructor

 $\mathbf{add_vertex} \colon \mathsf{Digraph} \ \mathsf{X} \ \mathsf{Data} \to \mathsf{void}$

Order: O(1)

Description: Adds a new vertex to the graph and stores the given data.

Precondition:

■ Digraph must exist.

■ Data must exist.

Postcondition: New vertex added to the given graph, storing the given data.

Classification: Modifier

 $\mathbf{remove_vertex}$: Digraph X Key \rightarrow Data

Order: O(n)

Description: Finds the vertex referenced by the given key in the digraph and removes it.

Returns the data stored in the vertex.

Precondition:

■ Digraph must exist.

• Key must reference an existing vertex in the digraph.

Postcondition: Vertex removed, and its data returned

Classification: Modifier

 $\mathbf{key_of} \colon \mathbf{Digraph} \ \mathbf{X} \ \mathbf{Data} \to \mathbf{Key}$

Order: O(n)

Description: Finds the key referencing the vertex storing the given data in the graph.

Precondition:

■ Digraph must exist.

■ The data must be stored in the digraph.

Postcondition: Key referencing the vertex storing the given data is returned.

 add_edge : Digraph X From Key X To Key \rightarrow void

Order: O(n)

Description: Adds a new edge connecting the vertices referenced by the given keys. The edge goes from the vertex referenced by From Key to the vertex referenced by To Key.

Precondition:

• Digraph must exist.

• From Key and To Key must reference vertices in the digraph.

Postcondition: New edge added to the graph, that goes from the vertex referenced by From

Key to the vertex referenced by To Key.

Classification: Modifier

remove_edge: Digraph X From Key X To Key \rightarrow void

Order: O(n)

Description: Finds the edge connecting the vertices referenced by From Key and To Key and removes it.

Precondition:

■ Digraph must exist.

• From Key and To Key must reference a existing vertices in the digraph.

 An edge going from the vertex referenced by From Key to the vertex referenced by To Key must exist.

Postcondition: Edge removed Classification: Modifier

 $\mathbf{get_vertex} \colon \mathbf{Digraph} \ \mathbf{X} \ \mathbf{Key} \to \mathbf{Data}$

Order: O(1)

Description: Finds the vertex referenced by Key in the digraph and returns the data stored

in it

Precondition:

Digraph must exist.

• Key must reference an existing vertex in the digraph.

Postcondition: Return the data stored in the vertex referenced by the given key.

 $\mathbf{get_adjacency_list}$: Digraph X Key \to Key List

Order: O(n)

Description: Returns the list of all vertices adjacent to the one referenced by the given key

in the graph. **Precondition:**

• Digraph must exist.

• Key must reference an existing vertex in the digraph.

Postcondition: Return a list of keys referencing all adjacent vertices to the one referenced

by the given key.

Classification: Analyzer

 $\mathbf{edge_exists} \colon \mathbf{Digraph} \ \mathbf{X} \ \mathbf{From} \ \mathbf{Key} \ \mathbf{X} \ \mathbf{To} \ \mathbf{Key} \to \mathbf{Boolean}$

Order: O(n)

Description: Verifies whether an edge from the vertex referenced by From Key to the vertex

referenced by To Key exists.

Precondition:

■ Digraph must exist.

• From Key and To Key must reference a existing vertices in the digraph.

Postcondition:

■ True if the edge exists.

• False if the edge does not exist.

Classification: Analyzer

 $\mathbf{random_digraph}$: Vertex Amount X Edge Amount \rightarrow Digraph

Order: O(n)

Description: Creates a random digraph with the given amount of vertices and edges.

Precondition:

• Vertex Amount > 0

■ Edge Amount > 0

Postcondition: A random digraph with the given amount of vertices and edges is created.

Classification: Constructor.

2.2. Graph Related Algorithms

Description: Algorithms related to graph theory

 $plain_search$: Graph \rightarrow Key List

Order: O(n)

Description: Given a graph returns a list of keys referencing all its vertices in no particular

order.

Precondition: Graph must exist

Postcondition: Return a list of keys referencing all vertices in the graph.

Classification: Analyzer

 $\mathbf{dfs}: \mathbf{Graph} \to \mathbf{Key} \ \mathbf{List}$

Order: $O(n^2)$

Description: Given a graph returns a list of keys referencing all its vertices using Depth

First Search, starting from the vertex first inserted in the graph.

Precondition: Graph must exist

Postcondition: Return a list of keys referencing all vertices in the graph using Depth First

Search.

Classification: Analyzer

bfs: Graph \rightarrow Key List

Order: $O(n^2)$

Description: Given a graph returns a list of keys referencing all its vertices using Breath

First Search, starting from the vertex first inserted in the graph.

Precondition: Graph must exist

Postcondition: Return a list of keys referencing all vertices in the graph using Breath First

Search.

Classification: Analyzer

 $number_of_sources: Graph \rightarrow Number$

Order: O(n)

Description: Given a graph returns the number of source vertices.

Precondition: Graph must exist

Postcondition: Return the amount of source vertices in the graph

 $\mathbf{number_of_sinks} \colon \mathsf{Graph} \to \mathsf{Number}$

Order: O(n)

Description: Given a graph returns the number of sink vertices.

Precondition: Graph must exist

Postcondition: Return the amount of sink vertices in the graph

Classification: Analyzer

 $\mathbf{is_weakly_connected} \colon \mathbf{Graph} \to \mathbf{Boolean}$

Order: $O(n^2)$

Description: Given a graph, checks whether it is weakly connected.

Precondition: Graph must exist

Postcondition:

• True if the graph is weakly connected.

■ False if the graph is not weakly connected.

 ${\bf Classification:} \ {\bf Analyzer}$

path_exists: Graph x From Key x To Key X Length X Transitive Closure \to Boolean **Order:** $O(n^2)$

Description: Checks if a path between the vertices referenced by From Key and To Key exist.

- A length or a transitive closure can be given, not both.
- If a length is given, the path can be no longer than the length.
- If a Transitive Closure is given the path is checked in it, achieving O(1) performance.

Precondition:

- Digraph must exist.
- From Key and To Key must reference a existing vertices in the digraph.
- If given, length > 0.

Postcondition:

- True if a path exists.
- False if a path does not exist.

Classification: Analyzer

warshall: Graph \rightarrow Transitive Closure

Order: $O(n^3)$

Description: Given a graph calculates its transitive closure using Warshall's.

Precondition: Graph must exist

Postcondition: The transitive closure must have the capability of calculating whether there

is a path between two vertices in O(1) time.

 $\mathbf{save} \colon \mathsf{Graph} \ \mathsf{X} \ \mathsf{Name} \ \mathsf{X} \ \mathsf{View} \ \mathsf{X} \ \mathsf{Format} \to \mathsf{void}$

Order: O(n)

Description: Given a graph, stores its graphic representation in the given format with the given name. A view of the file is given if View is True.

Precondition:

■ Graph must exist

- Name must be a valid file name
- If given, view must be True or False (default False)
- If given, format must be a valid file format for the Graphviz program (default 'pdf').

Postcondition:

- A file containing DOT code for Graphviz and an image in the given format are created.
- If View is True, the created image is displayed.