

ARQUITECTURA DE LAS COMPUTADORAS
Ejercicios de CIRCUITOS LÓGICOS para dar en clases

- 1) Dada la siguiente tabla de verdad, obtener la 1er. forma canónica algebraica y la numérica, el circuito y la 2da. forma canónica numérica y algebraica de la función:

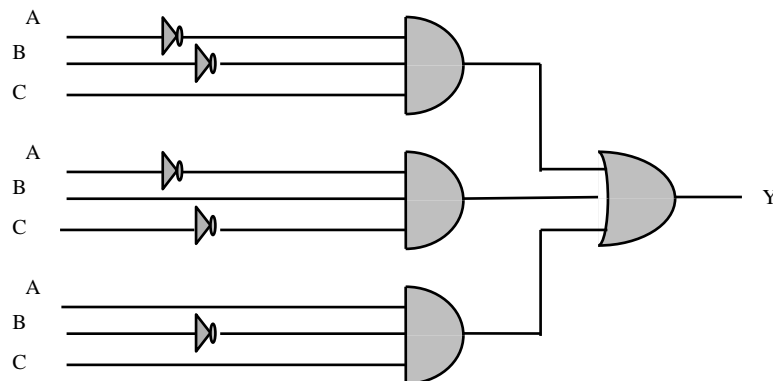
A	B	C	Y
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	0

RESPUESTA: $Y = A' \cdot B' \cdot C + A' \cdot B \cdot C' + A \cdot B' \cdot C$

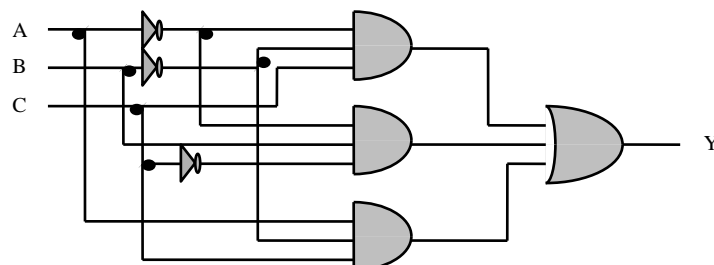
$$Y = Y(A,B,C) = \sum_3 m_i (1,2,5)$$

$$Y = \prod_3 M_i (0,3,4,6,7)$$

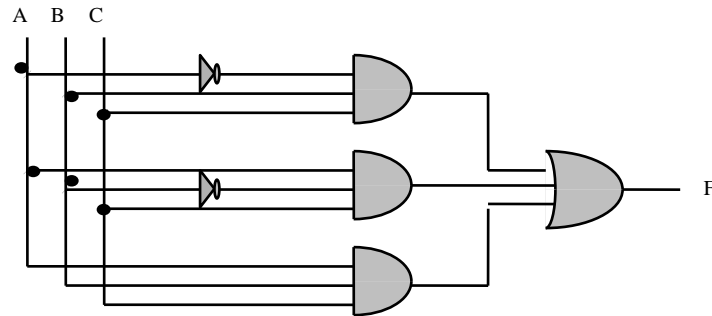
$$Y = (A+B+C) \cdot (A+B'+C') \cdot (A'+B+C) \cdot (A'+B'+C) \cdot (A'+B'+C')$$



Otra forma de hacer el circuito:



- 2) Dado el siguiente circuito, obtener la 1er. forma canónica algebraica y la numérica, la 2da. forma canónica numérica y algebraica y la tabla de verdad de la función respectiva:



RESPUESTA: $F = A' \cdot B \cdot C + A \cdot B' \cdot C + A \cdot B \cdot C$

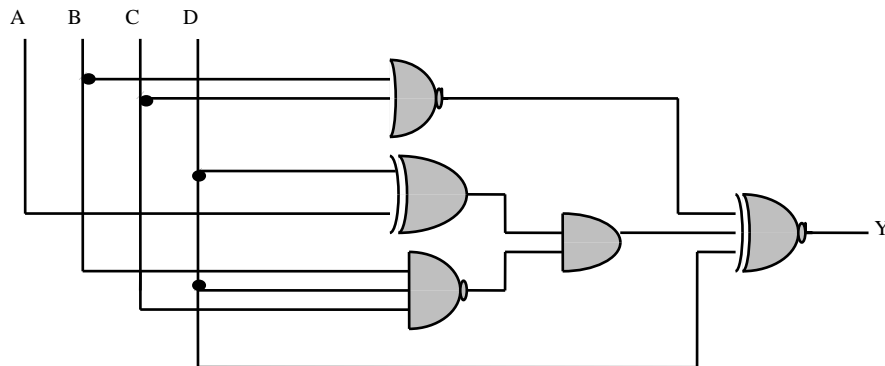
$$F = \sum_3 m_i (3,5,7)$$

$$F = \prod_3 M_i (0,1,2,4,6)$$

$$F = (A+B+C) \cdot (A+B+C') \cdot (A+B'+C) \cdot (A'+B+C) \cdot (A'+B'+C)$$

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

- 3) Escribir la expresión algebraica de la función del siguiente circuito:

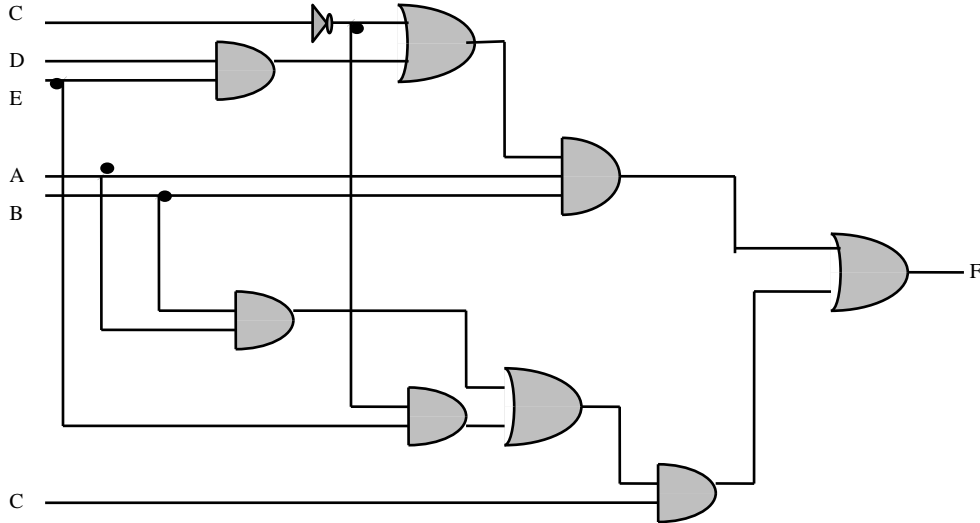


RESPUESTA: $Y = \{ (B + C)' \oplus [(D \oplus A) \cdot (B \cdot C \cdot D)'] \oplus D \}'$

4) Hacer el circuito de la siguiente función:

$$F = (C' + D.E) . A . B + (A.B + C'.E) . C$$

RESPUESTA:



5) Dada la siguiente función, escribir la tabla de verdad de la misma:

$$F_3 = A . B + C$$

RESPUESTA:

A	B	C	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

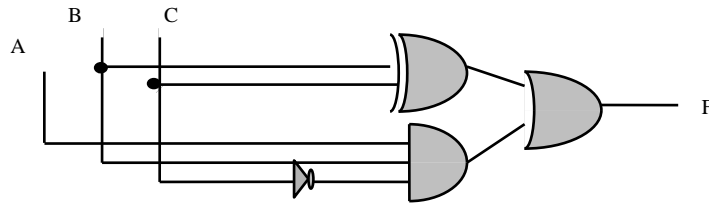
6) Dada la siguiente función, escribir su 1er. f. canónica:

$$F_4 = A . B . C . D' + A . B . C'$$

RESPUESTA: $F = \sum_4 m_i (14,12,13)$

$$F = A.B.C.D' + A.B.C'.D' + A.B.C'.D$$

7) Escribir la expresión algebraica y la 1er. f. canónica de la función cuyo circuito es el siguiente:



RESPUESTA: $F = (B \oplus C) + A \cdot B \cdot C' = B' \cdot C + B \cdot C' + A \cdot B \cdot C'$

$\therefore F = \sum_3 m_i (1, 5, 2, 6)$

$F = A' \cdot B' \cdot C + A \cdot B' \cdot C + A' \cdot B \cdot C' + A \cdot B \cdot C'$

8) Escribir la 2da. forma canónica numérica y la algebraica de la siguiente función:

$F = (A \cdot B \cdot C)' + (A \oplus B)$

RESPUESTA:

A	B	C	$(A \cdot B \cdot C)'$	$A \oplus B$	F
0	0	0	1	0	1
0	0	1	1	0	1
0	1	0	1	1	1
0	1	1	1	1	1
1	0	0	1	1	1
1	0	1	1	1	1
1	1	0	1	0	1
1	1	1	0	0	0

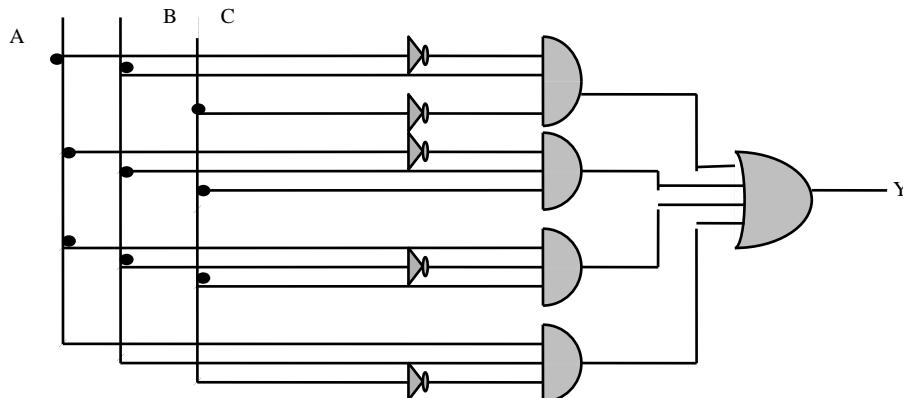
$F = \prod_3 M_i (7)$

$F = \bar{A} + \bar{B} + \bar{C}$

9) Dada la siguiente función, dibujar el circuito correspondiente:

$Y = \sum_3 m_i (2, 3, 5, 6)$

RESPUESTA:



10) Dada la siguiente función, determinar la 1er. f. c. numérica y la 2da. algebraica:

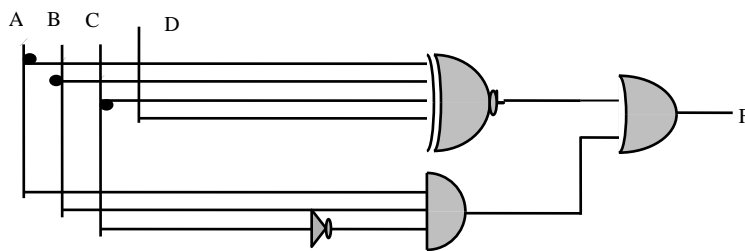
$$F(A, B, C) = A$$

RESPUESTA: $F = \sum_3 m_i (4, 5, 6, 7)$

$$\rightarrow F = \prod_3 M_i (0, 1, 2, 3)$$

$$\rightarrow F = (A+B+C) \cdot (A+B+\bar{C}) \cdot (A+\bar{B}+C) \cdot (A+\bar{B}+\bar{C})$$

11) Dado el siguiente circuito, determinar la 1er. f. c. numérica de la función:



RTA.: a) $F = (A \oplus B \oplus C \oplus D)' + A \cdot B \cdot \bar{C}$

$F = 1$ cuando uno, o los dos sumandos valen "1"

Analizando:

$$F_1 = A \cdot B \cdot \bar{C} = A \cdot B \cdot \bar{C} \cdot D + A \cdot B \cdot \bar{C} \cdot \bar{D} \rightarrow F_1 = \sum_4 m_i (13, 12)$$

$F_2 = (A \oplus B \oplus C \oplus D)'$ marca paridad par \rightarrow busco cuándo hay paridad par:

A	B	C	D		
1	1	1	1	\rightarrow 15	0 1 1 0 \rightarrow 6
1	1	0	0	\rightarrow 12	0 1 0 1 \rightarrow 5
1	0	1	0	\rightarrow 10	0 0 1 1 \rightarrow 3
1	0	0	1	\rightarrow 9	0 0 0 0 \rightarrow 0

$$\therefore F_2 = \sum_4 m_i (0, 3, 5, 6, 9, 10, 12, 15)$$

$$\Rightarrow F = F_1 + F_2 = \sum_4 m_i (0, 3, 5, 6, 9, 10, 12, 13, 15)$$

b) O bien, puede resolverse haciendo la tabla de verdad.

12) Escribir la 2da. f. c. numérica y algebraica de la siguiente función lógica:

$$Y = (A \oplus B) \cdot C + (A \cdot B + B \cdot C)' \cdot A$$

$$\text{RTA.: } Y = \underbrace{(A \oplus B) \cdot C}_{Y_1} + \underbrace{(A \cdot B + B \cdot C)' \cdot A}_{Y_2}$$

A	B	C	(A ⊕ B) // Y ₁	A · B // B · C	(A · B + B · C)' // Y ₂	Y
0	0	0	0	0	1	0
0	0	1	0	0	1	0
0	1	0	1	0	1	0
0	1	1	1	1	0	1
1	0	0	1	0	1	1
1	0	1	1	0	1	1
1	1	0	0	1	0	0
1	1	1	0	1	0	0

$$Y = \prod_3 M_i(0, 1, 2, 6, 7)$$

$$Y = (A+B+C) \cdot (A+B+\bar{C}) \cdot (A+\bar{B}+C) \cdot (\bar{A}+\bar{B}+C) \cdot (\bar{A}+\bar{B}+\bar{C})$$

13) Dada $F(A,B,C) = A' \cdot B \cdot C + A \cdot B \cdot C' + A' \cdot B \cdot C' + A \cdot B' \cdot C' + A \cdot B \cdot C$, encontrar la 2da. Forma Canónica y el circuito que ésta representa.

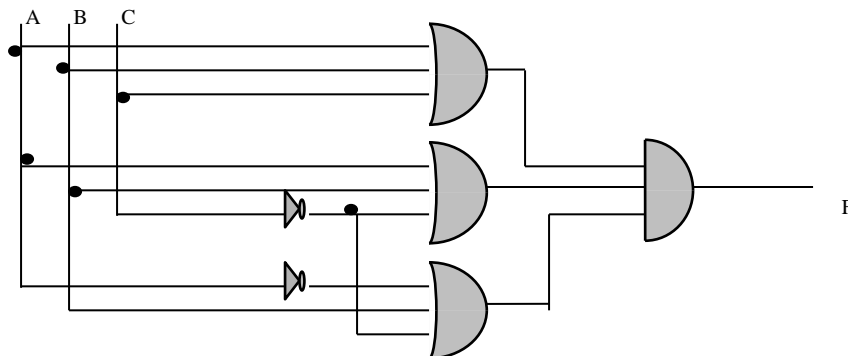
$$\text{RTA.: } F = \sum_3 m_i(3, 6, 2, 4, 7) = \sum_3 m_i(2, 3, 4, 6, 7)$$

$$\rightarrow F = \prod_3 M_i(0, 1, 5)$$

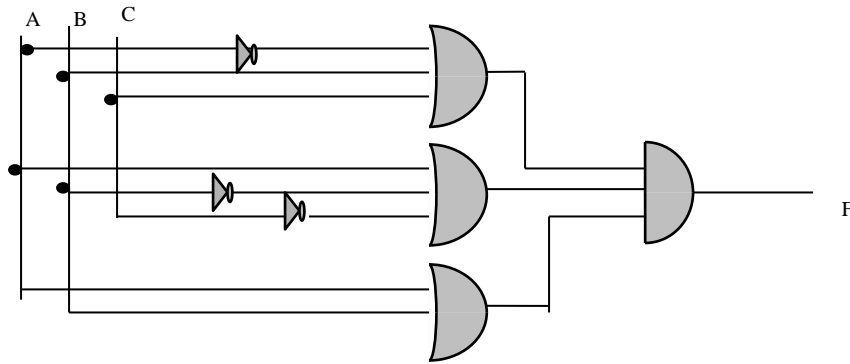
2a F.C.N.

$$F = (A+B+C) \cdot (A+B+C') \cdot (A'+B+C')$$

2a F.C.A.



14) Dado el circuito, encontrar la función lógica respectiva y la 1er. F.C. y la T. De V.



$$\text{RTA.: } F = (A' + B + C) \cdot (A + B' + C') \cdot (A + B)$$

$$\rightarrow F = \Pi_3 M_i (4, 3, 1, 0)$$

$$\rightarrow F = \Sigma_3 m_i (2, 5, 6, 7)$$

1er.F.C.N.

$$F = A' \cdot B \cdot C' + A \cdot B' \cdot C + A \cdot B \cdot C' + A \cdot B \cdot C$$

1er.F.C.A.

A	B	C	F
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

MINIMIZACIÓNMINIMIZACIÓN POR ÁLGEBRA DE BOOLE:

Simplificar, hasta llegar a la mínima expresión, utilizando las leyes del álgebra de Boole (pueden existir o no, varios caminos diferentes para resolver cada ejercicio, aquí mostraremos una sola forma para resolverlos) y comprobar que verifique la tabla de verdad:

$$1) F = \sum_4 m_i (5,7,13,15) = \overline{A}.B.C'.D + \overline{A}.B.C.D + \overline{A}.B.C'.D + \overline{A}.B.C.D =$$

$$= \overline{A}.B.D.(C'+C) + \overline{A}.B.D.(C'+C) = B.D$$

Para demostrar que la función del resultado es equivalente a la función dada como enunciado, puede realizarse la tabla de verdad, o la parte de la misma que nos interesa:

A	B	C	D	decimal
0	1	0	1	5
0	1	1	1	7
1	1	0	1	13
1	1	1	1	15

$$2) F = \sum_3 m_i (0,1,2,3,4,5) = \overline{A}.B'.C' + \overline{A}.B'.C + \overline{A}.B.C' + \overline{A}.B.C + \overline{A}.B'.C' + \overline{A}.B'.C =$$

$$= \overline{A}.B'.(C'+C) + \overline{A}.B.(C'+C) + \overline{A}.B'.(C'+C) = \overline{A}.B' + \overline{A}.B + \overline{A}.B' =$$

$$I) = \overline{A}.B' + \overline{A}.B + \overline{A}.B' + \overline{A}.B' = \overline{A}.(B'+B) + \overline{A}.B'.(A'+A) = \overline{A} + B' = (A.B)'$$

$$II) = \overline{A}.B' + \overline{A}.B + \overline{A}.B' = \overline{A}.(B'+B) + \overline{A}.B' = \overline{A} + \overline{A}.B' = \overline{A} + B' = (A.B)'$$

$$III) = \overline{A}.B' + \overline{A}.B + \overline{A}.B' = B' . (A' + A) + \overline{A} . B = B' + \overline{A} . B = B' + A' = (B . A)'$$

3)

A	B	C	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	0

$$F = \overline{A}.B'.C + \overline{A}.B'.C' + \overline{A}.B'.C + \overline{A}.B.C' = B'.C.(A'+A) + \overline{A}.C'.(B'+B) = B'.C + \overline{A}.C'$$

O bien, trabajando por los "0":
$$F = (A+B+C).(A+B'+C).(A+B'+C').(A'+B'+C') =$$

$$= (A+C) . (B'+C') = (A+C) . (B . C)'$$

4) $F = \Pi_3 M_i(0,1,4,7) = (A+B+C).(A+B+C').(A'+B+C).(A'+B'+C') =$

$$= (A+B+C.C') . (A.A'+B+C) . (A'+B'+C') = (A+B) . (B+C).(A'+B'+C') =$$

(I) $= (B+A.C) . (B'+(A.C)') = B.(A.C)' + (A.C).B' = B \oplus (A.C)$

(II) $= (B + A.C) . (A . B . C)'$

5) $Y = A'.B'.C'.D' + A'.B.C'.D' + A'.B.A'.B.C'.D + B =$

$$= A'.B'.C'.D' + B . (A'.C'.D' + A' + A'.C'.D + 1) = A'.B'.C'.D' + B = B'.(A+C+D)' + B =$$

$$= B + (A + C + D)'$$

6) $Y = A.B.C' + A.B'.C.D + A.B'.C.D' + A'.B'.C.D + A'.C =$

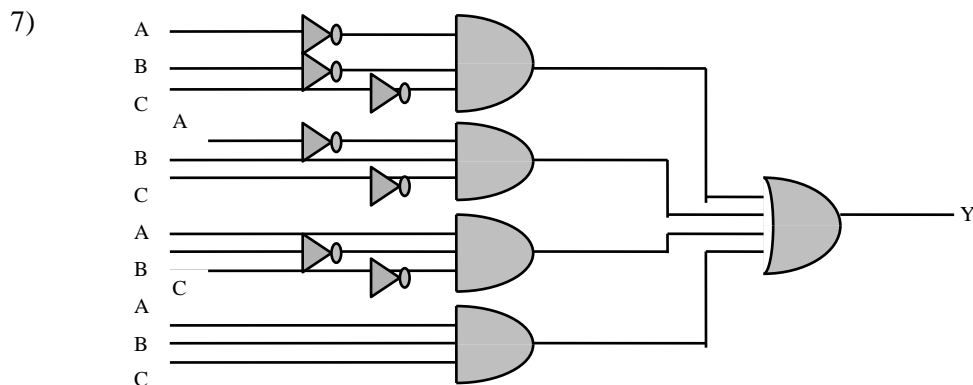
$$= A.B.C' + A.B'.C.(D+D') + A'.C.(B'.D+1) = A.B.C' + A.B'.C + A'.C =$$

(I) $= A.B.C' + C.(A.B' + A') = A.B.C' + C.(B' + A') = A.B.C' + C.(B.A)' = C \oplus B.A$

(II) $= A . (B . C' + B' . C) + A' . C = A . (B \oplus C) + A' . C$

(III) $= A . (B . C' + B' . C) + C . (A . B' + A') = A . (B \oplus C) + C . (A' + B') =$

$$= A . (B \oplus C) + C . (A . B)'$$



$$Y = \overline{A}.B'.C' + \overline{A}.B.C' + \overline{A}.B'.C + A.B.C = \overline{A}.C'.(B'+B) + B'.C'.(\overline{A}+A) + A.B.C =$$

$$= \overline{A}.C' + B'.C' + A.B.C = C'.(\overline{A}+B') + C.A.B = C'(A.B)' + C.A.B = (C \oplus (A.B))'$$

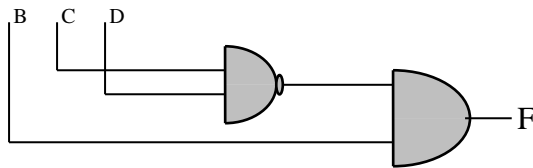
8) $= \Sigma_4 m_i (4,5,6,12,13,14)$, y dibujar el circuito mínimo.

$$F = \overline{A}.B.C'.D' + \overline{A}.B.C'.D + \overline{A}.B.C.D' + \overline{A}.B.C.D + A.B.C'.D' + A.B.C'.D + A.B.C.D' + A.B.C.D =$$

$$= \overline{A}.B.C'.(D'+D) + B.C.D'.(\overline{A}+A) + A.B.C'.(D'+D) =$$

$$= \overline{A}.B.C' + B.C.D' + A.B.C' = B.C.D' + B.C'.(\overline{A}+A) = B.C.D' + B.C' =$$

$$= B.(C.D' + C') = B.(D'+C') = B.(D.C)'$$



9) $S = \overline{A}.B'.C'.D' + \overline{A}.B'.C'.D + \overline{A}.B.C.D' + \overline{A}.B.C.D + A.B.C =$

$$= \overline{A}.B'.C'.(D'+D) + \overline{A}.B.C.(D'+D) + A.B.C = \overline{A}.B'.C' + \overline{A}.B.C + A.B.C =$$

$$(I) = \overline{A}.B'.C' + B.C.(A'+A) = (A+B+C)' + B.C$$

$$(II) = \overline{A}.(B'.C' + B.C) + A.B.C = \overline{A}.(B \oplus C)' + A.B.C = [A + (B \oplus C)]' + A.B.C$$

10) $Y = \overline{A}.B'.C'.D' + \overline{A}.B + \overline{A}.B'.C.D + B' + \overline{A}.B.C'.D' =$

$$= \overline{A}.B.(1 + C'.D') + B'.(\overline{A}.C'.D' + \overline{A}.C.D + 1) = \overline{A}.B + B' = \overline{A} + B' = (A.B)'$$

11) $F = \Sigma_4 m_i (0, 1, 3, 4, 5, 7, 11, 15) = \overline{A}.B'.C'.D' + \overline{A}.B'.C'.D + \overline{A}.B'.C.D + \overline{A}.B.C'.D' +$

$$+ \overline{A}.B.C'.D + \overline{A}.B.C.D + \overline{A}.B'.C.D + A.B.C.D =$$

$$= \overline{A}.B'.C'.(D'+D) + \overline{A}.B.C'.(D'+D) + B'.C.D.(A'+A) + B.C.D.(A'+A) =$$

$$= \overline{A}.B'.C' + \overline{A}.B.C' + B'.C.D + B.C.D = \overline{A}.C'.(B'+B) + C.D.(B'+B) =$$

$$= \overline{A}.C' + C.D = (A+C)' + C.D$$

$$12) F = \prod_4 M_i (0, 2, 5, 7, 13, 15) =$$

$$= (A+B+C+D).(A+B+C'+D).(A+B'+C+D').(A+B'+C'+D').(A'+B'+C+D').(A'+B'+C'+D') =$$

$$= (A+B+D+C.C').(A+B'+D'+C.C').(A'+B'+D'+C.C') =$$

$$= (A+B+D).(A+B'+D').(A'+B'+D') = (A+B+D).(A.A'+B'+D') =$$

$$= (A+B+D).(B'+D') =$$

$$(I) = (BD)' . (A+B+D)$$

$$(II) = A.B'+A.D'+B.D'+D.B' = (\text{por 9 a}) = A.B' + B.D' + D.B' =$$

$$(a) = A.B' + (B \oplus D)$$

$$(b) = B' . (A + D) + (B \oplus D)$$

$$(c) = B' . (A + D) + B . D'$$

$$(III) = A.B' + A.D' + B.D' + D.B' =$$

$$(a) = A . (B'+D') + (B \oplus D) = A . (B.D)' + (B \oplus D)$$

$$(b) = D' + (A + B) + B' . (A + D)$$

$$13) Y = \underline{A'.B'.C'.D' + A'.B'.C'.D + A'.B' + A'.B'.C.D} + B =$$

$$= A'.B'.(C'.D' + C'.D + 1 + C.D) + B = A'.B' + B = B + A'$$

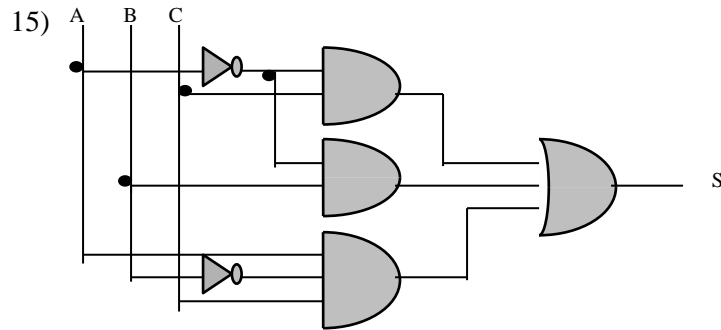
$$14) F = \underline{A'.B'.C'.D + A'.B.C'.D + A'.B + A'.B.D' + A'.B.C' + D} =$$

$$= A'.B.(C'.D + 1 + D' + C') + D.(A'.B'.C' + 1) = A'.B + D$$

O bien:

$$F = \underline{A'.B'.C'.D + A'.B.C'.D + A'.B + A'.B.D' + A'.B.C' + D} =$$

$$= A'.B.(1 + D' + C') + D.(A'.B'.C' + A'.B.C' + 1) = A'.B + D$$



$$S = A'.C + A'.B + A.B'.C = C.(A' + A.B') + A'.B = C.A' + C.B' + A'.B =$$

$$(I)(\text{por 9a}) = A'.B + B'.C$$

$$(II) = A'.B + C.(A' + B') = A'.B + C.(A.B)'$$

$$16) Y = \underline{A'.B'.C'.D} + \underline{A'.B'.C.D} + \underline{A'.B.C'.D} + \underline{A'.B.C.D} + A'.B.C =$$

$$= A'.C'.D.(B' + B) + A'.B.C + A'.C.D.(B' + B) = A'.C'.D + A'.B.C + A'.C.D =$$

$$= A'.D.(C' + C) + A'.B.C = A'.D + A'.B.C = A'.(D + B.C)$$

17)

A	B	C	F
0	0	0	1
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

$$F = \underline{A.B.C} + \underline{A.B'.C} + \underline{A.B.C'} + \underline{A'.B.C} + \underline{A'.B'.C'} =$$

$$= A.B.(C + C') + A.C.(B + B') + B.C.(A + A') + A'.B'.C' =$$

$$= A.B + A.C + B.C + A'.B'.C' = A.B + C.(A + B) + C'.(A + B)' =$$

$$= A.B + (C \oplus (A + B))'$$

O bien:

$$F = (A + B + C') \cdot (A + B' + C) \cdot (A' + B + C) = [A + (B + C') \cdot (B' + C)] \cdot (A' + B + C) =$$

$$= [A + B \cdot C + B' \cdot C'] \cdot (A' + B + C) = [A + (B \oplus C)'] \cdot (A' + B + C)$$

O también:

$$F = (A + B + C') \cdot (A + B' + C) \cdot (A' + B + C) =$$

$$= [A + (B + C') \cdot (B' + C)] \cdot [C + (A + B') \cdot (A' + B)] =$$

$$= [A + B \cdot C + B' \cdot C'] \cdot [C + A \cdot B + B' \cdot A'] = [A + (B \oplus C)'] \cdot [C + (A \oplus B)']$$

O: $F = (A + B + C') \cdot (A + B' + C) \cdot (A' + B + C) = (A + B + C') \cdot [C + (A + B') \cdot (A' + B)] =$

$$= (A + B + C') \cdot [C + A \cdot B + B' \cdot A'] = (A + B + C') \cdot [C + (A \oplus B)']$$

18) $F = B \cdot C \cdot D + A' \cdot B' \cdot C' \cdot D + A' \cdot B' \cdot C' \cdot D' + B' \cdot C \cdot D =$

$$= C \cdot D \cdot (B + B') + A' \cdot B' \cdot C' \cdot (D + D') = C \cdot D + A' \cdot B' \cdot C' = C \cdot D + (A + B + C)'$$

19) $F = A' \cdot B' \cdot C' + A \cdot B \cdot C \cdot D + B \cdot C' \cdot D + A' \cdot C \cdot D + A \cdot B' \cdot D + B \cdot C' \cdot D =$

$$= A' \cdot B' \cdot C' + B \cdot D \cdot (A \cdot C + C') + C \cdot D \cdot (A' + A \cdot B) + A \cdot D \cdot (B' + B \cdot C) =$$

$$= A' \cdot B' \cdot C' + B \cdot D \cdot C' + B \cdot D \cdot A + C \cdot D \cdot A' + C \cdot D \cdot B + A \cdot D \cdot B' + A \cdot D \cdot C =$$

$$= B \cdot D \cdot (C' + C) + A \cdot D \cdot (B' + B) + C \cdot D \cdot (A' + A) + A' \cdot B' \cdot C' = B \cdot D + A \cdot D + C \cdot D + A' \cdot B' \cdot C' =$$

$$= D \cdot (B + A + C) + A' \cdot B' \cdot C' = D \cdot (B + A + C) + (A + B + C)' = D + (A + B + C)'$$

20) Dada la siguiente tabla, construir la tabla de verdad completa. Desarrollar la función por minterms y por maxterms. Simplificarlas. Dibujar circuito mínimo.

A	B	C	F
0	0	0	1
0	0	1	1
1	0	0	1

RTA.:

A	B	C	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

$$F = \overline{A} \cdot \overline{B} \cdot C + \overline{A} \cdot B \cdot \overline{C} + A \cdot B \cdot \overline{C} =$$

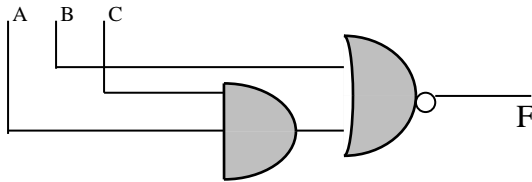
$$I) = \overline{A} \cdot \overline{B} \cdot (C + C) + B \cdot \overline{C} \cdot (\overline{A} + A) = \overline{A} \cdot \overline{B} + B \cdot \overline{C} = \overline{B} \cdot (\overline{A} + C) = \overline{B} \cdot (A \cdot C)' = (B + A \cdot C)'$$

$$II) = \overline{A} \cdot \overline{B} \cdot (C + C) + A \cdot B \cdot \overline{C} = \overline{A} \cdot \overline{B} + A \cdot B \cdot \overline{C} = \overline{B} \cdot (\overline{A} + AC) = \overline{B} \cdot (\overline{A} + C) = \overline{B} \cdot (A \cdot C)' = (B + A \cdot C)'$$

$$F = (\overline{A + B' + C}) \cdot (\overline{A + B' + C}) \cdot (\overline{A' + B + C}) \cdot (\overline{A' + B' + C}) \cdot (\overline{A' + B' + C}) =$$

$$= (\overline{A + B' + C}) \cdot (\overline{A' + C + B \cdot B'}) \cdot (\overline{A' + B' + C}) = (\overline{A + B'}) \cdot (\overline{A' + C}) \cdot (\overline{A' + B'}) =$$

$$= \overline{B} \cdot (\overline{A' + C}) = \overline{B} \cdot (AC)' = (B + A \cdot C)'$$



MINIMIZACIÓN POR MAPA DE KARNAUGH:

Ejercicios: Repetiremos algunos de los ejercicios dados en minimización por Álgebra de Boole.

1) $F = \sum_4 m(5,7,13,15) = A'.B.C'.D + A'.B.C.D + A.B.C'.D + A.B.C.D$

		A B			
C	D	<div><div></div><div></div></div>			
		0 0	0 1	1 1	1 0
	0 0				
	0 1		1	1	
	1 1		1	1	
	1 0				

$$F = B . D$$

2) $F = \sum_3 m(0,1,2,3,4,5) = A'.B'.C' + A'.B'.C + A'.B.C' + A'.B.C + A.B'.C' + A.B'.C$

		A B			
C		0 0	0 1	1 1	1 0
0		1	1		1
1		1	1		1

$$F = A' + B' = (A . B)'$$

3)

A	B	C	F
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	0

		A B			
C		0 0	0 1	1 1	1 0
0				1	1
1		1			1

$$F = A . C' + B' . C$$

4) $F = \prod_3 M_i(0, 1, 4, 7)$

		A B					
C		0 0	0 1	1 1	1 0		
		0	1	0	1		
	0	0			0		
	1	0		0			

$$\begin{aligned}
 F &= (A+B).(B+C).(A'+B'+C') = \\
 &= (B+A.C).(B'+(A.C)') = \\
 &= B.(A.C)' + (A.C).B' = B \oplus (A.C)
 \end{aligned}$$

O BIEN:

		A B					
C		0 0	0 1	1 1	1 0		
		0	1	0	1		
	0		1	1			
	1		1		1		

$$\begin{aligned}
 F &= A'.B + B.C' + A.B'.C = \\
 &= B.(A'+C') + C.A.B' = B.(A.C)' + A.C.B' = \\
 &= B \oplus (A.C)
 \end{aligned}$$

OTRA RESPUESTA, para el primer caso:

$$F = (A+B).(B+C).(A'+B'+C') = (B + A.C).(A . B . C)'$$

5) $Y = A'.B'.C'.D' + A'.B.C'.D' + A'.B.A'.B.C'.D + B$

		A B					
C D		0 0	0 1	1 1	1 0		
		0	1	0	1		
	0 0	1	1	1			
	0 1		1	1			
	1 1		1	1			
	1 0		1	1			

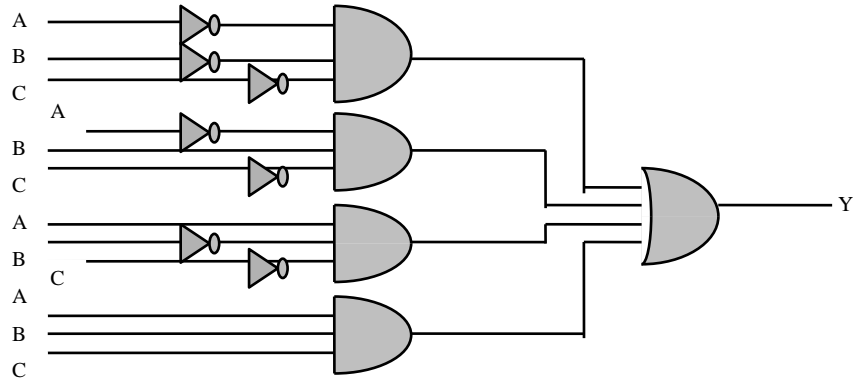
$$Y = B + A'.C'.D' = B + (A+C+D)'$$

6) $Y = A.B.C' + A.B'.C.D + A.B'.C.D' + A'.B'.C.D + A'.C$

		A B					
C D		0 0	0 1	1 1	1 0		
		0	1	0	1		
	0 0			1			
	0 1			1			
	1 1	1	1		1		
	1 0	1	1		1		

$$\begin{aligned}
 Y &= A'.C + B'.C + A.B.C' = \\
 &= C.(A.B)' + C'.A.B = \\
 &= C \oplus (A.B)
 \end{aligned}$$

7)



$$Y = A'.B'.C' + A'.B.C' + A.B'.C' + A.B.C$$

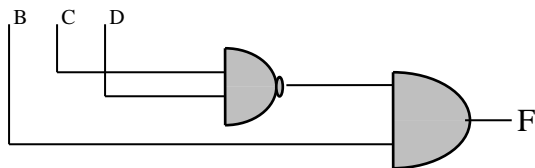
		A B							
C		0 0	0 1	1 1	1 0				
		0	1	1	0				
	0	1	1		1				
	1			1					

$$Y = B'.C' + A'.C' + A.B.C = (C \oplus (A.B))'$$

8) $F = \sum_4 m_i (4, 5, 6, 12, 13, 14)$, y dibujar el circuito mínimo.

		A B							
C D		0 0	0 1	1 1	1 0				
		0	1	1	0				
	0 0		1	1					
	0 1		1	1					
	1 1								
	1 0		1	1					

$$F = B.C' + B.D' = B.(C' + D') = B.(C.D)'$$



9) $F = A'.B'.C'.D' + A'.B'.C'.D + A'.B.C.D' + A'.B.C.D + A.B.C$

		A B					
C	D	0	0	0	1	1	1
		0	0	1	1	0	0
0	0	1					
0	1	1					
1	1			1	1		
1	0			1	1		

$$F = B.C + A'.B'.C'$$

$$= BC + (A+B+C)'$$

10) $Y = A'.B'.C'.D' + A'.B + A'.B'.C.D + B' + A'.B.C'.D'$

		A B					
C	D	0	0	0	1	1	1
		0	0	1	1	0	0
0	0	1		1			1
0	1	1		1			1
1	1	1		1			1
1	0	1		1			1

$$Y = B' + A' = (B . A)'$$

11) $F = \sum_4 m_i (0,1,3,4,5,7,11,15)$

		A B					
C	D	0	0	0	1	1	1
		0	0	1	1	0	0
0	0	1		1			
0	1	1		1			
1	1	1		1	1		1
1	0						

$$F = C . D + A' . C' = (A + C)' + C . D$$

12) $F = \prod_4 M_i (0,2,5,7,13,15)$

		A B					
C	D	0	0	0	1	1	1
		0	0	1	1	0	0
0	0	0					
0	1			0	0		
1	1			0	0		
1	0	0					

$$F = (B'+D') . (A+B+D)$$

$$= (B . D)' . (A+B+D)$$

O bien: $F = A.B' + A.D' + B.D' + D.B' = (\text{por 9a}) = A.B' + B.D' + D.B' = A.B' + (B \oplus D)$

Otra forma de resolverlo:

		A B					
C D		0 0	0 1	1 1	1 0		
		0 0	0 1	1 1	1 0		
	0 0		1	1	1		
	0 1	1			1		
	1 1	1			1		
	1 0		1	1	1		

$F = A.B' + B.D' + B'.D = (I) = A.B' + (B \oplus D)$

$(II) = B'.(A + D) + B.D'$

También podría ser:

		A B					
C D		0 0	0 1	1 1	1 0		
		0 0	0 1	1 1	1 0		
	0 0		1	1	1		
	0 1	1			1		
	1 1	1			1		
	1 0		1	1	1		

$F = B.D' + B'.D + A.D' = (I) = A.D' + (B \oplus D)$

$(II) = D'.(A + B) + B'.D$

13) $Y = A'.B'.C'.D' + A'.B'.C'.D + A'.B' + A'.B'.C.D + B$

		A B					
C D		0 0	0 1	1 1	1 0		
		0 0	0 1	1 1	1 0		
	0 0	1	1	1			
	0 1	1	1	1			
	1 1	1	1	1			
	1 0	1	1	1			

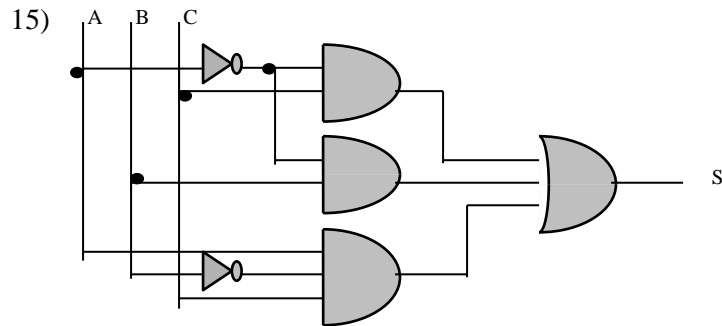
$Y = A' + B$

Trabajando con "0" $\Rightarrow Y = A' + B$

14) $F = A' \cdot B' \cdot C' \cdot D + A' \cdot B \cdot C' \cdot D + A' \cdot B + A' \cdot B \cdot D' + A' \cdot B \cdot C' + D$

		A B					
C	D	0 0	0 1	1 1	1 0		
0	0		1				
0	1	1	1	1	1		
1	1	1	1	1	1		
1	0		1				

$F = A' \cdot B + D$



$S = A' \cdot C + A' \cdot B + A \cdot B' \cdot C$

		A B					
C		0 0	0 1	1 1	1 0		
0			1				
1		1	1		1		

$S = A' \cdot B + B' \cdot C$

16) $Y = A' \cdot B' \cdot C' \cdot D + A' \cdot B' \cdot C \cdot D + A' \cdot B \cdot C' \cdot D + A' \cdot B \cdot C \cdot D + A' \cdot B \cdot C$

		A B					
C	D	0 0	0 1	1 1	1 0		
0	0						
0	1	1	1				
1	1	1	1				
1	0		1				

$Y = A' \cdot D + A' \cdot B \cdot C = A' \cdot (D + B \cdot C)$

17)

A	B	C	F
0	0	0	1
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	1

	A	B	
C	0	0	1
0	1		1
1		1	1

$$F = A' \cdot B' \cdot C' + A \cdot B + B \cdot C + A \cdot C$$

$$F = A \cdot B + C \cdot (A + B) + C' \cdot (A + B)' = A \cdot B + (C \oplus (A + B))'$$

$$\begin{aligned} \text{O bien: } F &= A' \cdot B' \cdot C' + A \cdot B + B \cdot C + A \cdot C = B' \cdot A' \cdot C' + B \cdot (A + C) + A \cdot C = \\ &= B' \cdot A' \cdot C' + B \cdot (A' \cdot C')' + A \cdot C = B \oplus (A' \cdot C') + A \cdot C = \\ &= [B \oplus (A + C)'] + A \cdot C \end{aligned}$$

Trabajando con los "0":

	A	B	
C	0	0	1
0		0	0
1	0		

$$F = (A + B' + C) \cdot (A' + B + C) \cdot (A + B + C')$$

No se logró minimizar. Hay que trabajarlo con lógica de Boole.

$$\begin{aligned} \text{(I)} \quad &= [A + (B + C')] \cdot (B' + C) \cdot (A' + B + C) = \\ &= [A + B \cdot C + B' \cdot C'] \cdot (A' + B + C) = [A + (B \oplus C)'] \cdot (A' + B + C) \end{aligned}$$

$$\begin{aligned} \text{(II)} \quad &= (A + B + C') \cdot [C + (A + B')] \cdot (A' + B) = \\ &= (A + B + C') \cdot [C + A \cdot B + B' \cdot A'] = (A + B + C') \cdot [C + (A \oplus B)'] \end{aligned}$$

$$\begin{aligned} \text{(III)} \quad &= [A + (B + C')] \cdot (B' + C) \cdot [C + (A + B')] = \\ &= [A + B \cdot C + B' \cdot C'] \cdot [C + A \cdot B + B' \cdot A'] = \\ &= [A + (B \oplus C)'] \cdot [C + (A \oplus B)'] \end{aligned}$$

$$18) F = B \cdot C \cdot D + A' \cdot B' \cdot C' \cdot D + A' \cdot B' \cdot C' \cdot D' + B' \cdot C \cdot D$$

		A B			
C	D	0 0	0 1	1 1	1 0
		1			
0	1	1			
1	1	1	1	1	1
1	0				

$$\begin{aligned} F &= C \cdot D + A' \cdot B' \cdot C' \\ &= C \cdot D + (A + B + C)' \end{aligned}$$

$$19) F = A' \cdot B' \cdot C' + A \cdot B \cdot C \cdot D + B \cdot C' \cdot D + A' \cdot C \cdot D + A \cdot B' \cdot D + B \cdot C' \cdot D$$

		A B			
C	D	0 0	0 1	1 1	1 0
		1			
0	1	1	1	1	1
1	1	1	1	1	1
1	0				

$$\begin{aligned} F &= D + A' \cdot B' \cdot C' \\ &= D + (A + B + C)' \end{aligned}$$

Trabajando con los "0":

		A B			
C	D	0 0	0 1	1 1	1 0
			0	0	0
0	1				
1	1				
1	0	0	0	0	0

$$F = (C' + D) \cdot (A' + D) \cdot (B' + D) = D + (C' \cdot A' \cdot B') = D + (C + A + B)'$$

20) Dada la siguiente tabla, construir la tabla de verdad completa. Minimizarla. Dibujar circuito mínimo:

A	B	C	F
0	0	0	1
0	0	1	1
1	0	0	1

RTA.:

A	B	C	F
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	0

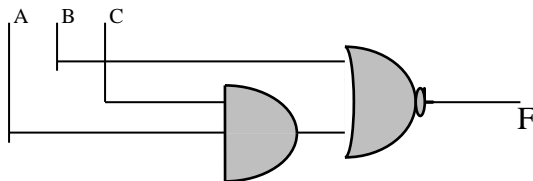
	A	B					
C	0	0	0	1	1	1	0
0	1						1
1	1						

$$\begin{aligned}
 F &= A'.B' + B'.C' = \\
 &= B'.(A' + C') = B'.(A.C)' = \\
 &= (B + (A.C))'
 \end{aligned}$$

Trabajando con los ceros:

	A	B					
C	0	0	0	1	1	1	0
0			0		0		
1			0		0		0

$$F = B'.(A' + C') = B'.(A.C)' = [B + (A.C)]'$$



21) $F = A'.B'.C'.D' + A.B'.C'.D' + A'.B'.C.D' + A.B'.C.D'$

		A	B				
C	D						
		0 0	0 1	1 1	1 0		
0	0	1				1	
0	1						
1	1						
1	0	1				1	

$$F = B'.D' = (B + D)'$$