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Lic. en Cs. de la Computación

# ANÁLISIS DE LENGUAJES DE PROGRAMACIÓN

## Trabajo Práctico 4

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# Ejercicio 1

a) Demostración de State como mónada:

```
newtype State a = State {runState :: Env -> (a, Env)}

instance Monad State where
    return x = State (\s -> (x, s))
    m >>= f = State (\s -> let (v, s') = runState m s
                              in runState (f v) s')

-- Monad 1
return y >>= f =
-- {Definicion de return}
State (\s -> (y, s)) >>= f =
-- {Definicion de >>=}
State (\s -> let (v, s') = runState (State (\s -> (y, s))) s
              in runState (f v) s') =
-- {Definicion de runState}
State (\s -> let (v, s') = (\s -> (y, s)) s
              in runState (f v) s') =
-- {Aplicacion de funcion anonima}
State (\s -> let (v, s') = (y, s)
              in runState (f v) s') =
-- {Sustitucion de let}
State (\s -> runState (f y) s) =
-- {Eta-reduccion}
State (runState (f y)) =
-- {Sea x = State z. Luego, State (runState x) = State z = x}
-- {Entonces, la composicion de State y runState es la identidad}
f y

-- Monad 2
m >>= return =
-- {Definicion de >>=}
State (\s -> let (v, s') = runState m s
              in runState (return v) s') =
-- {Definicion de return}
State (\s -> let (v, s') = runState m s
              in runState (State (\s -> (v, s))) s') =
-- {Definicion de runState}
State (\s -> let (v, s') = runState m s
              in (\s -> (v, s)) s') =
-- {Aplicacion de funcion anonima}
State (\s -> let (v, s') = runState m s
              in (v, s')) =
-- {Sustitucion de let}
State (\s -> runState m s) =
-- {Eta-reduccion}
State (runState m) =
-- {La composicion de State y runState es la identidad}
m

-- Monad 3
(m >>= f) >>= g =
-- {Definicion de >>=}
(State (\s -> let (v, s') = runState m s
              in runState (f v) s')) >>= g =
-- {Definicion de >>=}
State (\s -> let (v, s') = runState (State (\s -> let (v, s') = runState m s
                                                    in runState (f v) s')) s
              in runState (g v) s') =
```

```

— {Definicion de runState}
State (\s -> let (v, s') = (\s -> let (v, s') = runState m s
                                   in runState (f v) s')) s
    in runState (g v) s') =
— {Aplicacion de funcion anonima}
State (\s -> let (v, s') = (let (v, s') = runState m s
                             in runState (f v) s')
    in runState (g v) s') =
— {Se reescribe el let de la siguiente forma:
— let a = (let b = c in f(b)) in g(a)
— let b = c
—     a = f(b)
— in g(a)}
State (\s -> let (v, s') = runState m s
                (w, s'') = runState (f v) s'
    in runState (g w) s'')

m >>= (\x -> f x >>= g) =
— {Definicion de >>=}
State (\s -> let (v, s') = runState m s
    in runState ((\x -> f x >>= g) v) s') =
— {Aplicacion de funcion anonima}
State (\s -> let (v, s') = runState m s
    in runState (f v >>= g) s') =
— {Definicion de >>=}
State (\s -> let (v, s') = runState m s
    in runState (State (\s -> let (v, s') = runState (f v) s
                             in runState (g v) s')) s') =
— {Definicion de runState}
State (\s -> let (v, s') = runState m s
    in (\s -> let (v, s') = runState (f v) s
        in runState (g v) s') s') =
— {Aplicacion de funcion anonima}
State (\s -> let (v, s') = runState m s
    in (let (v, s') = runState (f v) s'
        in runState (g v) s')) =
— {Se reescribe el let de la siguiente forma:
— let a = b in (let c = f(a) in g(c))
— let a = b
—     c = f(a)
— in g(c)}
State (\s -> let (v, s') = runState m s
                (w, s'') = runState (f v) s'
    in runState (g w) s'')

— {Ambos terminos son iguales a una misma expresion}
— {Por transitividad, (m >>= f) >>= g = m >>= (\x -> f x >>= g)}

```

b) Implementación del evaluador para la mónada State:

```
— Evalua un programa en el estado nulo
eval :: Comm -> Env
eval p = snd (runState (evalComm p) initState)

— Evalua un comando en un estado dado
evalComm :: MonadState m => Comm -> m ()
evalComm Skip = return ()
evalComm (Let var ie) = do n <- evalIntExp ie
                        update var n
evalComm (Seq c1 c2) = do evalComm c1
                        evalComm c2
evalComm (Cond be c1 c2) = do b <- evalBoolExp be
                        if b then (evalComm c1)
                        else (evalComm c2)
evalComm (While be c) = evalComm (Cond be (Seq c (While be c)) Skip)

— Resuelve la evaluacion de dos expresiones enteras, aplicadas a un operador.
evalIntOp :: MonadState m => IntExp -> IntExp -> (Int -> Int -> a) -> m a
evalIntOp ie1 ie2 operator = do n1 <- evalIntExp ie1
                        n2 <- evalIntExp ie2
                        return (operator n1 n2)

— Evalua una expresion entera, sin efectos laterales
evalIntExp :: MonadState m => IntExp -> m Int
evalIntExp (Const i) = return i
evalIntExp (Var x) = lookfor x
evalIntExp (UMinus ie) = do n <- evalIntExp ie
                        return (-n)
evalIntExp (Plus ie1 ie2) = evalIntOp ie1 ie2 (+)
evalIntExp (Minus ie1 ie2) = evalIntOp ie1 ie2 (-)
evalIntExp (Times ie1 ie2) = evalIntOp ie1 ie2 (*)
evalIntExp (Div ie1 ie2) = evalIntOp ie1 ie2 div

— Resuelve la evaluacion de dos expresiones booleanas, aplicadas a un operador.
evalBoolOp :: MonadState m => BoolExp -> BoolExp -> (Bool -> Bool -> a) -> m a
evalBoolOp be1 be2 operator = do n1 <- evalBoolExp be1
                        n2 <- evalBoolExp be2
                        return (operator n1 n2)

— Evalua una expresion booleana, sin efectos laterales
evalBoolExp :: MonadState m => BoolExp -> m Bool
evalBoolExp BTrue = return True
evalBoolExp BFalse = return False
evalBoolExp (Eq ie1 ie2) = evalIntOp ie1 ie2 (==)
evalBoolExp (Lt ie1 ie2) = evalIntOp ie1 ie2 (<)
evalBoolExp (Gt ie1 ie2) = evalIntOp ie1 ie2 (>)
evalBoolExp (And be1 be2) = evalBoolOp be1 be2 (&&)
evalBoolExp (Or be1 be2) = evalBoolOp be1 be2 (||)
evalBoolExp (Not be) = do b <- evalBoolExp be
                        return (not b)
```

## Ejercicio 2

a) Instancia de Monad para StateError:

```
instance Monad StateError where
  return x = StateError (\s -> Just (x, s))
  m >>= f = StateError (\s -> case runStateError m s of
    Nothing -> Nothing
    Just (v, s') -> runStateError (f v) s')
```

b) Instancia de MonadError para StateError:

```
instance MonadError StateError where
  throw = StateError (\s -> Nothing)
```

c) Instancia de MonadState para StateError:

```
instance MonadState StateError where
  lookfor v = StateError (\s -> maybe (Nothing) (\u -> Just (u, s)) (lookfor ' v s))
    where lookfor ' v [] = Nothing
          lookfor ' v ((u,j):ss) | v == u      = Just j
                                | otherwise = lookfor ' v ss
  update v i = StateError (\s -> Just ((), update ' v i s))
    where update ' v i [] = [(v, i)]
          update ' v i ((u, j):ss) | v == u      = (v, i):ss
                                | otherwise = (u, j):(update ' v i ss)
```

d) Implementación del evaluador utilizando la mónada StateError:

```
— Evalua un programa en el estado nulo
eval :: Comm -> Env
eval c = maybe initState snd (runStateError (evalComm c) initState)

— Evalua un comando en un estado dado
evalComm :: (MonadState m, MonadError m) => Comm -> m ()
evalComm Skip = return ()
evalComm (Let var ie) = do n <- evalIntExp ie
                        update var n
evalComm (Seq c1 c2) = do evalComm c1
                        evalComm c2
evalComm (Cond be c1 c2) = do b <- evalBoolExp be
                        if b then (evalComm c1)
                        else (evalComm c2)
evalComm (While be c) = evalComm (Cond be (Seq c (While be c)) Skip)

— Resuelve la evaluacion de dos expresiones enteras, aplicadas a un operador.
evalIntOp :: (MonadState m, MonadError m) => IntExp -> IntExp -> (Int -> Int -> a)
-> m a
evalIntOp ie1 ie2 operator = do n1 <- evalIntExp ie1
                        n2 <- evalIntExp ie2
                        return (operator n1 n2)

— Evalua una expresion entera, sin efectos laterales
evalIntExp :: (MonadState m, MonadError m) => IntExp -> m Int
evalIntExp (Const i) = return i
evalIntExp (Var x) = lookfor x
evalIntExp (UMinus ie) = do n <- evalIntExp ie
                        return (-n)
evalIntExp (Plus ie1 ie2) = evalIntOp ie1 ie2 (+)
evalIntExp (Minus ie1 ie2) = evalIntOp ie1 ie2 (-)
evalIntExp (Times ie1 ie2) = evalIntOp ie1 ie2 (*)
evalIntExp (Div ie1 ie2) = do n1 <- evalIntExp ie1
                        n2 <- evalIntExp ie2
                        if n2 == 0 then throw
                        else return (div n1 n2)

— Resuelve la evaluacion de dos expresiones booleanas, aplicadas a un operador.
evalBoolOp :: (MonadState m, MonadError m) => BoolExp -> BoolExp -> (Bool -> Bool -> a)
-> m a
evalBoolOp be1 be2 operator = do n1 <- evalBoolExp be1
                        n2 <- evalBoolExp be2
                        return (operator n1 n2)

— Evalua una expresion booleana, sin efectos laterales
evalBoolExp :: (MonadState m, MonadError m) => BoolExp -> m Bool
evalBoolExp BTrue = return True
evalBoolExp BFalse = return False
evalBoolExp (Eq ie1 ie2) = evalIntOp ie1 ie2 (==)
evalBoolExp (Lt ie1 ie2) = evalIntOp ie1 ie2 (<)
evalBoolExp (Gt ie1 ie2) = evalIntOp ie1 ie2 (>)
evalBoolExp (And be1 be2) = evalBoolOp be1 be2 (&&)
evalBoolExp (Or be1 be2) = evalBoolOp be1 be2 (||)
evalBoolExp (Not be) = do b <- evalBoolExp be
                        return (not b)
```

## Ejercicio 3

a) Mónada propuesta:

```
— Monada estado
newtype StateErrorTick a =
    StateErrorTick {runStateErrorTick :: Env -> Maybe (a, Int, Env)}
```

b) Clase MonadTick:

```
class Monad m => MonadTick m where
    — Crea un contador en 1
    tick :: m ()
```

c) Instancia de MonadTick:

```
instance MonadTick StateErrorTick where
    tick = StateErrorTick (\s -> Just ((), 1, s))
```

d) Instancia de MonadError:

```
instance MonadError StateErrorTick where
    throw = StateErrorTick (\s -> Nothing)
```

e) Instancia de MonadState:

```
instance MonadState StateErrorTick where
    lookfor v = StateErrorTick (\s -> maybe (Nothing) (\u -> Just (u, 0, s)) (lookfor' v s))
        where lookfor' v [] = Nothing
              lookfor' v ((u,j):ss) | v == u = Just j
                                   | otherwise = lookfor' v ss
    update v i = StateErrorTick (\s -> Just ((), 0, update' v i s))
        where update' v i [] = [(v, i)]
              update' v i ((u, j):ss) | v == u = (v, i):ss
                                   | otherwise = (u, j):(update' v i ss)
```

f) Implementación del evaluador utilizando la mónada StateErrorTick:

```

— Evalua un programa en el estado nulo
eval :: Comm -> (Int, Env)
eval c = maybe ((-1), [])
          (\(v, count, s) -> (count, s))
          (runStateErrorTick (evalComm c) initState)

— Evalua un comando en un estado dado
evalComm :: (MonadState m, MonadError m, MonadTick m) => Comm -> m ()
evalComm Skip = return ()
evalComm (Let var ie) = do n <- evalIntExp ie
                          update var n
evalComm (Seq c1 c2) = do evalComm c1
                          evalComm c2
evalComm (Cond be c1 c2) = do b <- evalBoolExp be
                              if b then (evalComm c1)
                              else (evalComm c2)
evalComm (While be c) = evalComm (Cond be (Seq c (While be c))) Skip

— Resuelve la evaluacion de dos expresiones enteras, aplicadas a un operador.
evalIntOp :: (MonadState m, MonadError m, MonadTick m) => IntExp -> IntExp ->
  (Int -> Int -> a) -> m a
evalIntOp ie1 ie2 operator = do n1 <- evalIntExp ie1
                                n2 <- evalIntExp ie2
                                return (operator n1 n2)

— Evalua una expresion entera, sin efectos laterales.
evalIntExp :: (MonadState m, MonadError m, MonadTick m) => IntExp -> m Int
evalIntExp (Const i) = return i
evalIntExp (Var x) = lookfor x
evalIntExp (UMinus ie) = do n <- evalIntExp ie
                          return (-n)
evalIntExp (Plus ie1 ie2) = tick >> evalIntOp ie1 ie2 (+)
evalIntExp (Minus ie1 ie2) = tick >> evalIntOp ie1 ie2 (-)
evalIntExp (Times ie1 ie2) = tick >> evalIntOp ie1 ie2 (*)
evalIntExp (Div ie1 ie2) = do n1 <- evalIntExp ie1
                              n2 <- evalIntExp ie2
                              if n2 == 0 then throw
                              else tick >> return (div n1 n2)

— Resuelve la evaluacion de dos expresiones booleanas, aplicadas a un operador.
evalBoolOp :: (MonadState m, MonadError m, MonadTick m) => BoolExp -> BoolExp ->
  (Bool -> Bool -> a) -> m a
evalBoolOp be1 be2 operator = do n1 <- evalBoolExp be1
                                n2 <- evalBoolExp be2
                                return (operator n1 n2)

— Evalua una expresion booleana, sin efectos laterales
evalBoolExp :: (MonadState m, MonadError m, MonadTick m) => BoolExp -> m Bool
evalBoolExp BTrue = return True
evalBoolExp BFalse = return False
evalBoolExp (Eq ie1 ie2) = evalIntOp ie1 ie2 (==)
evalBoolExp (Lt ie1 ie2) = evalIntOp ie1 ie2 (<)
evalBoolExp (Gt ie1 ie2) = evalIntOp ie1 ie2 (>)
evalBoolExp (And be1 be2) = evalBoolOp be1 be2 (&&)
evalBoolExp (Or be1 be2) = evalBoolOp be1 be2 (||)
evalBoolExp (Not be) = do b <- evalBoolExp be
                        return (not b)

```