# Algorithmic Strategies 2024/25 Week 9 – Graph Algorithms

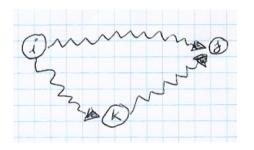


Universidade de Coimbra

- Floyd-Warshall Algorithm is a dynamic programming approach that finds the shortest path between every pair of vertices in  $O(|V|^3)$
- It can detect negative cycles
- It finds the shortest path with negative weights if there is no negative cycle
- It can also be used to solve reachability questions: whether vertex i can reach vertex j.

Let D(i,j,k) be the solution to the subproblem of finding the shortest path from i to j using the vertices from  $\{1,\ldots,k\}$  as intermediate vertices.

$$D(i,j,k) = egin{cases} w(i,j) & \text{if } k=0 \ \min egin{cases} D(i,j,k-1) \ D(i,k,k-1) + D(k,j,k-1) \end{cases}$$
 otherwise



#### Floyd-Warshall Algorithm

```
Function FW(G, D)

for i = 1 to |V| do

D[i, i, 0] = 0

for each arc (i, j) \in A do

D[i, j, 0] = w(i, j)

for k = 1 to |V| do

for i = 1 to |V| do

for j = 1 to |V| do

D[i, j, k] = \min(D[i, j, k - 1], D[i, k, k - 1] + D[k, j, k - 1])

return D
```

The cells of matrix D are initialized to  $\infty$ .

#### Floyd-Warshall Algorithm

```
Function FW(G, D)

for i=1 to |V| do

D[i,i]=0

for each arc (i,j) \in A do

D[i,j]=w(i,j)

for k=1 to |V| do

for i=1 to |V| do

for j=1 to |V| do

D[i,j]=\min(D[i,j],D[i,k]+D[k,j])

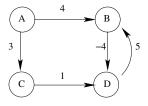
return D
```

Only a two-dimensional matrix is required

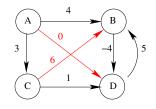
#### Floyd-Warshall Algorithm

```
Function FW(G, D)
  for i = 1 to |V| do
     D[i, i] = 0
  for each arc (i, j) \in A do
     D[i,j] = w(i,j)
  for k = 1 to |V| do
    for i = 1 to |V| do
       for i = 1 to |V| do
          D[i,j] = \min(D[i,j], D[i,k] + D[k,j])
  for i = 1 to |V| do
    if D[i, i] < 0 then
                                            {check for negative cycle}
       return negative cycle found
  return D
```

## Find all-pairs shortest path



## Find all-pairs shortest path

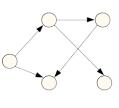


	i	j	k	D[i,j]
Π	Α	D	В	0
	Α	D	C	4
	С	В	D	6

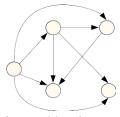
The transitive closure of a graph G = (V, A) is a graph  $G^* = (V, A^*)$ , where

 $A^* = \{(i,j) : \text{ there is a path from vertex } i \text{ to vertex } j \text{ in G } \}$ 

Given a transitive closure, it is possible to know, in constant time, whether one vertice is reachable from another.



 $\mathsf{A} \; \mathsf{graph}$ 



Its transitive closure

#### Floyd-Warshall Algorithm for transitive closure

```
Function FW(G, D)

for each arc (i,j) \in A do

D[i,j] = \text{true}

for k = 1 to |V| do

for i = 1 to |V| do

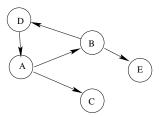
for j = 1 to |V| do

D[i,j] = D[i,j] \lor (D[i,k] \land D[k,j])

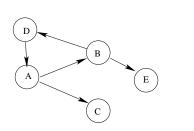
return D
```

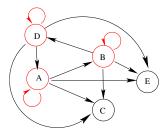
The cells of matrix D are initialized to false.

Compilers routinely create so-called "call graphs" during the compilation process. A call graph represents which functions in a program call each other, and may be used to identify functions that are never called, for example. Call graphs may also be used to identify which functions in a program originate recursion. Your task is to write a program that identifies the recursive functions in a given call graph, i.e., functions that call themselves, either directly or through other functions.



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Check the main diagonal of the adjacency matrix of the transitive closure