Algorithms for geometric problems.

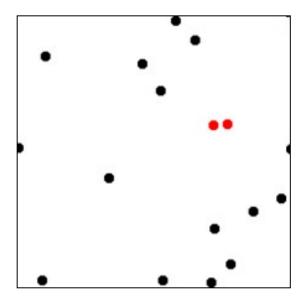
• Relevant for Games, Computer Graphics, Robotics and GIS.

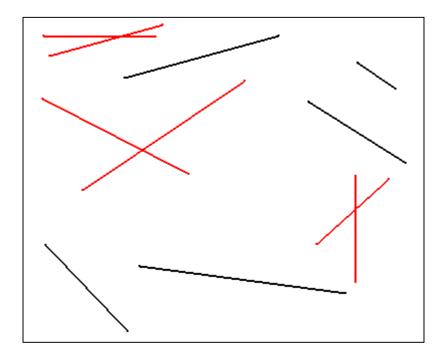
 We focus on "combinatorial computational geometry", that is, the objects under study are basic geometrical objects, such as points, lines segments, polygons and polyhedra.

#### **Readings:**

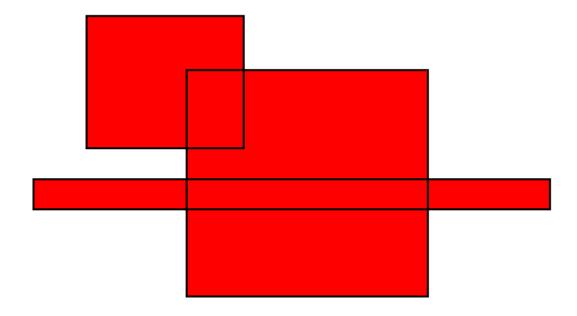
- S. Skiena, M. Revilla, Programming Challenges, Chapter 13
- S. Skiena, The Algorithm Design Manual, Chapter 17
- T. Cormen et al., Introduction to Algorithms, Chapter 33
- David Goldberg. "What Every Computer Scientist Should Know About Floating-Point Arithmetic". ACM Computing Surveys, 23 (1): 5–48, 1991 (link)

Closest pair problem



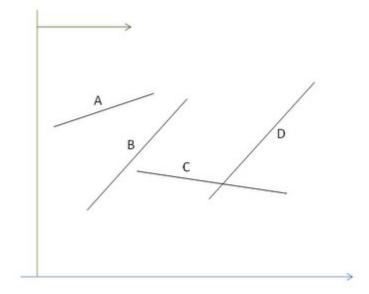


• Area of the union of rectangles problem

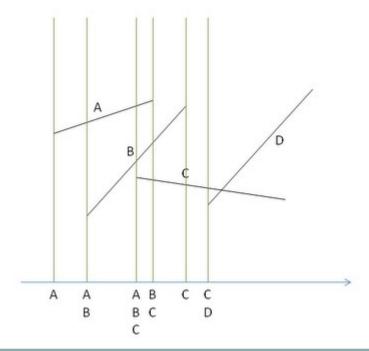


- Special strategies for computational geometry:
  - Line Sweep
    - A line sweeps the plane, stopping at some points.
    - Perform geometric operations at each stop.
    - Assumes that the objects are sorted in one dimension.
    - The solution is obtained once the line passed all objects.
    - For 3D, it is called plane sweep.

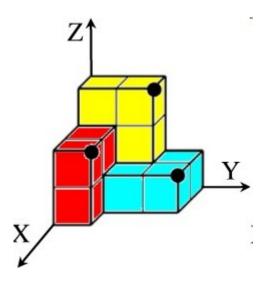
- Special strategies for computational geometry:
  - Line Sweep



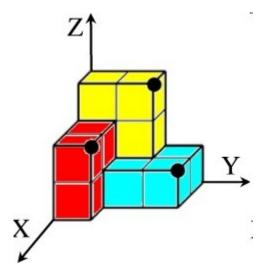
- Special strategies for computational geometry:
  - Line Sweep

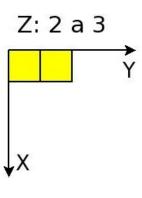


- Special strategies for computational geometry:
  - Plane Sweep

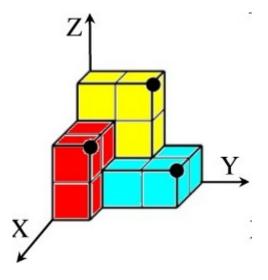


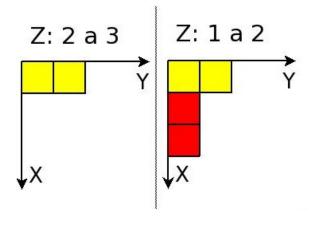
- Special strategies for computational geometry:
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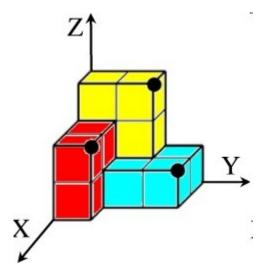


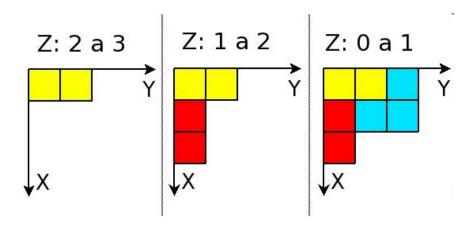
- Special strategies for computational geometry:
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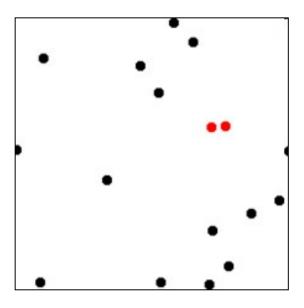


- Special strategies for computational geometry:
  - Plane Sweep





Closest pair problem (given n points in 2D)



- Closest pair problem
- Brute force: O(n²)
  - Compute all the distances between the *n* points and pick the pair with the smallest distance.

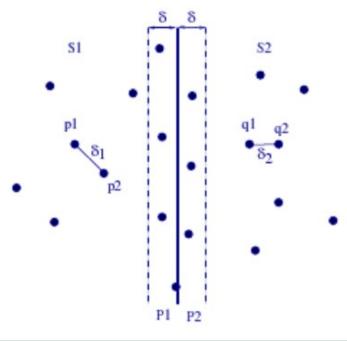
- Closest pair problem
- Divide-and-Conquer: O(n log n)
  - **Step 1**: Sort points along the x-coordinate.
  - Step 2: Divide: Split the points into 2 equal subsets, and solve recursively, the left and right subsets ( $\delta_1$  and  $\delta_2$ ).
  - Step 3: Merge: Find the minimum distance between points in the left and in the right subset ( $\delta_c$ ). The minimum distance is min( $\delta_1$ , $\delta_2$ ,  $\delta_c$ ). How to do it efficiently?

- Closest pair problem
- Divide-and-Conquer: O(n log n)

#### Step 3:

How to compute the minimum distance between the points in P1 and P2?

$$\delta = \min(\delta_1, \delta_2)$$

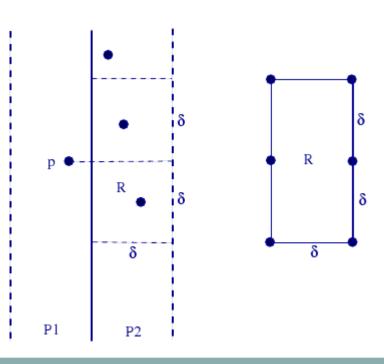


- Closest pair problem
- Divide-and-Conquer: O(n log n)

#### Step 3:

For each point p, check neighbors at a distance at most  $\delta$  of p

There can be at most 6 neighbors at a distance  $\delta$  of p.

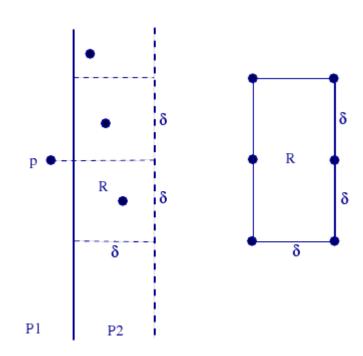


- Closest pair problem
- Divide-and-Conquer: O(n log n)

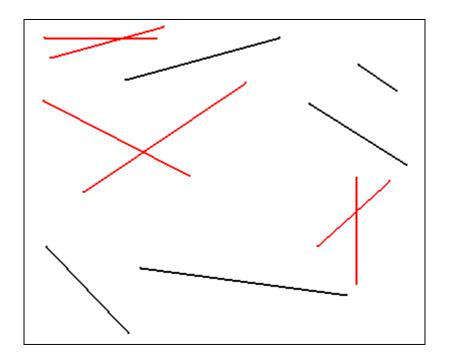
Step 3 in O(n):

- Project points in P1 and P2 on the vertical line
- For each point, compute the minimum distance from the 6 neighbors

(All points must be sorted in y)

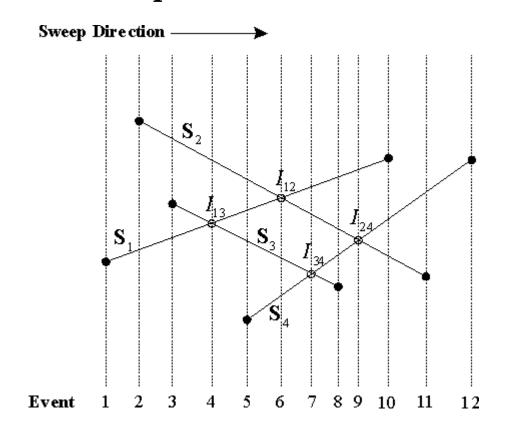


• Line intersection problem (given *n* lines and *k* intersections)



- Line intersection problem
- Brute force: O(n²)
  - Detect intersection between every pair of lines.

Line intersection problem



Bentley-Ottmann algorithm by sweep-line

#### Line intersection problem

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Sweep line: O((n+k) log n)
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Step 1: Sort 2n points in x-coordinate and let  $Q = \{p_1,...,p_{2n}\}$ 

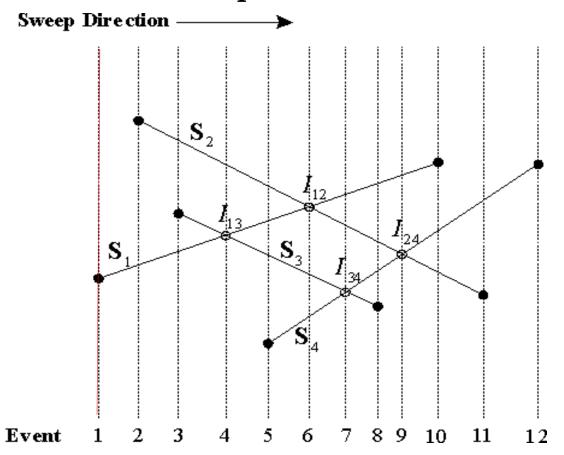
Step 2: From left to right, for each point *p* in *Q*:

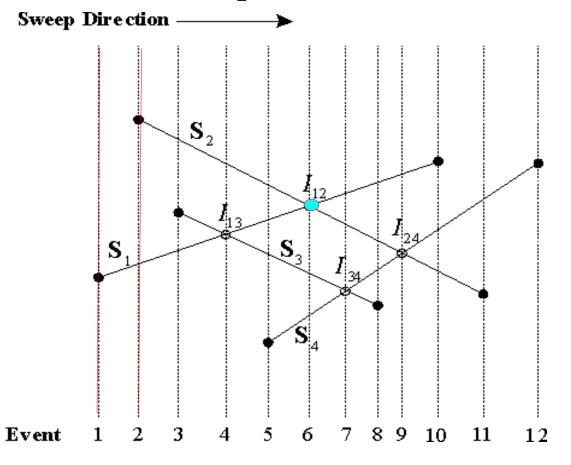
Step 2.1: If *p* is left-endpoint of *l*, add intersection between *l* and its neighboring lines to Q.

Step 2.2: If *p* is right-endpoint of *l*, add intersection between its neighboring lines to *Q* (to the right of *p*).

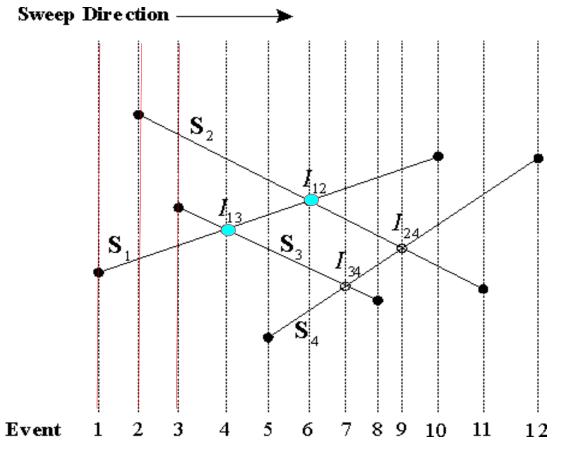
Step 2.3: If p is an intersection between  $l_1$  and  $l_2$ , add intersection between  $l_1$  and  $l_2$  and their neighboring lines to Q.

- Q is a priority queue that keeps the line with the minimum x-coordinate
- Use a balanced binary tree to maintain the order of the segments at each step in the y-coordinate

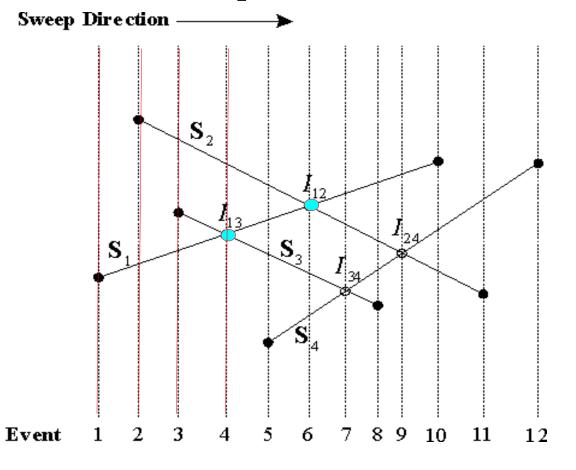




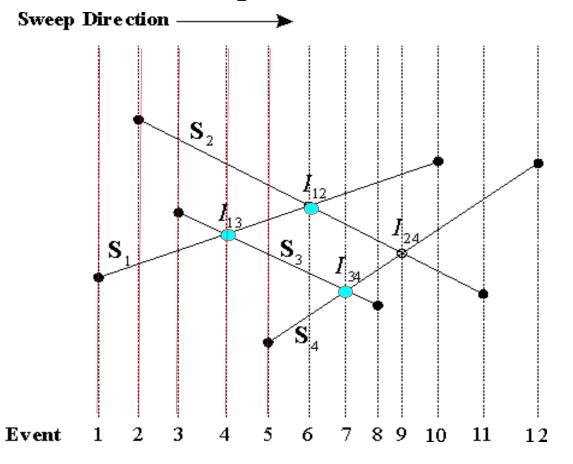
- Check intersection of S1 and S2
- Add intersection I12



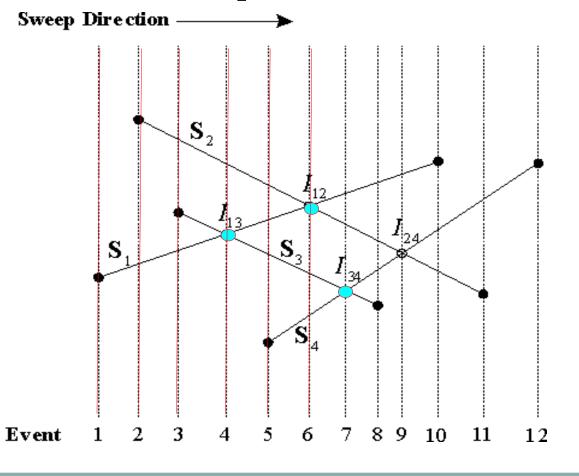
- Check intersection of S3 and S1 and of S3 and S2
- Add intersection I13



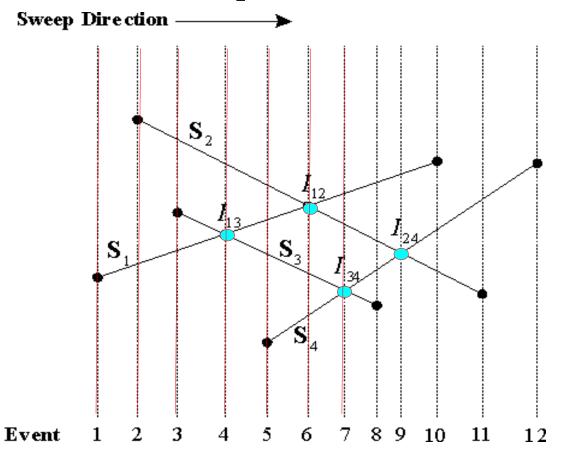
- Swap position of S1 and S3
- Check intersection of S1 and S2



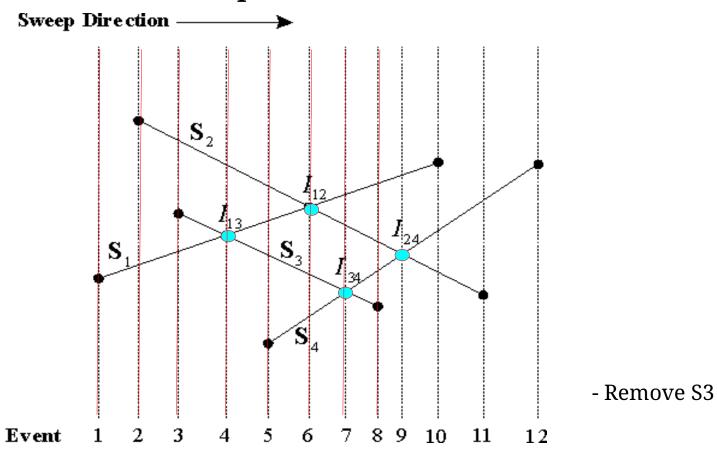
- Check intersection of S4 and S3
- Add intersection I34

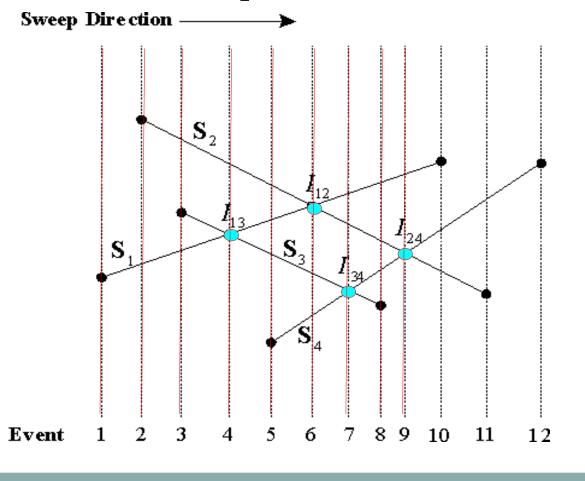


- Swap position of S1 and S2
- Check intersection of S2 and S3



- Swap position of S3 and S4
- Check intersection of S2 and S4
- Add intersection I24

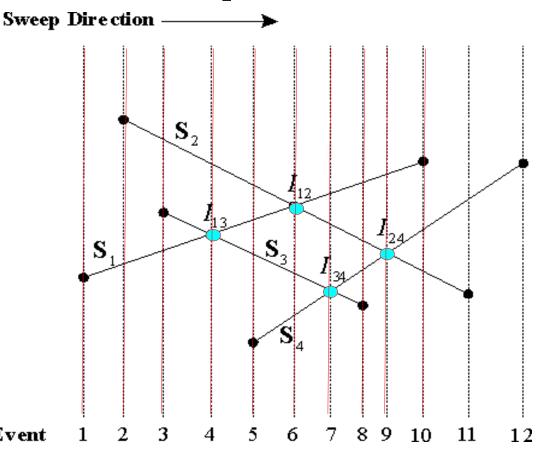




- Swap position of S2 and S4
- Check intersection of S1 and S4

Line intersection problem

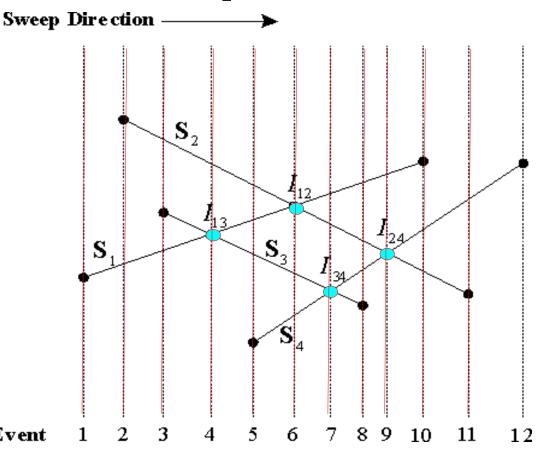
**Event** 



-Remove S1

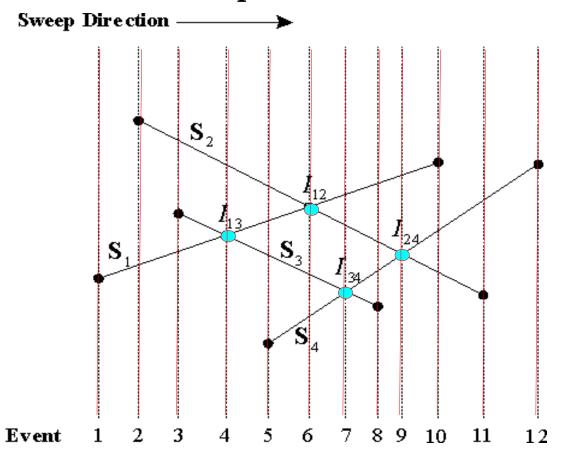
Line intersection problem

**Event** 



-Remove S2

Line intersection problem



-Remove S3