

Kombinatorika 2 - zapiski s predavanj prof. Konvalinke

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Seznam uporabljenih kratic

kratica	izraz
NSTE	naslednje trditve so ekvivalentne
orf	običajna rodovna funkcija
erf	eksponentna rodovna funkcija
fp	formalni polinom
fpv	formalna potenčna vrsta
dum	delno urejena množica

Poglavje 1

Osnove

1.1 Kako štejemo?

S končna množica, $|S| = ?$

Pogosto $S_n, n \in \mathbb{N}$.

Preštevalno zaporedje $|S_0|, |S_1|, |S_2| \dots$

Kaj je odgovor?

(1) Formula.

$$[n] = \{1, 2 \dots n\}.$$

$$S_n = 2^{[n]} = P([n]).$$

$$|S_n| = 2^n.$$

$$S_n = \{\text{permutacije } n \text{ elementov}\}.$$

$$|S_n| = n! = 1 \cdot 2 \cdots n \text{ „}n \text{ fakulteta“ „}n \text{ factorial“}.$$

$$S_n = \{\text{kompozicije } n \text{ s členi } 1 \text{ ali } 2\}, \text{ npr. } 5 = 1+2+1.$$

$$|S_5| = 8.$$

$$1, 1, 2, 3, 5, 8 \dots$$

$$|S_n| = F_n - \text{Fibonaccijevo zaporedje}.$$

(2) Asimptotska formula.

$$|S_n| \sim a_n \text{ (to pomeni } \lim_{n \rightarrow \infty} \frac{a_n}{|S_n|} = 1).$$

$n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n$ - Stirlingova formula.

$$F_n \sim \frac{1}{\sqrt{5}} \left(\frac{1+\sqrt{5}}{2}\right)^{n+1}.$$

(3) Z rekurzijo.

$$S_n = 2^{[n]}.$$

$$a_n = |S_n|, a_n = 2a_{n-1}; \quad n \geq 1, \quad a_0 = 1.$$

$$S_n = \{\text{kompozicije s členi 1 ali 2}\}.$$

$$S_n = F_n, F_n = F_{n-1} + F_{n-2}; \quad n \geq 2, \quad F_0 = F_1 = 1.$$

F_{n-1} - kompozicije, ki se končajo z 1, F_{n-2} - končajo z 2.

(4) Z rodovno funkcijo (generating function).

$(a_n)_{n \in \mathbb{N}}$ zaporedje.

$$a_0 + a_1x + a_2x^2 + \dots = \sum_{n=0}^{\infty} a_n x^n = \sum_n a_n x^n \text{ običajna (ordinary)}$$

rodovna funkcija - ORF.

$$a_n = 2^n, \quad \sum_{n=0}^{\infty} 2^n x^n = \frac{1}{1-2x}.$$

$$\sum_n F_n x^n = \frac{1}{1-x-x^2}.$$

$$\sum_n n! x^n //.$$

$\sum_n \frac{a_n}{n!} x^n$ eksponentna rodovna funkcija.

$$\sum_n 2^n \frac{x^n}{n!} = e^{2x}.$$

$$\sum_n \frac{n!}{n!} x^n = \frac{1}{1-x}.$$

(4) je najboljši način, da poznamo zaporedje.

- Rodovna funkcija je velikokrat „lepa“, tudi če ni lepe formule za zaporedje.

$i_n \dots \#$ involucij z n elementi ($\pi^2 = \text{id}$).

ni enostavnejše formule za i_n .

$$\sum_{n=0}^{\infty} \frac{i_n}{n!} x^n = e^{x + \frac{x^2}{2}}$$

- Do rodovne funkcije lahko pogosto pridemo neposredno s kombinatoričnim premislekom.

Involucija = permutacija s cikli dolžine 1 ali 2.

$$\sum F_n x^n = \frac{1}{1-x-x^2}; \quad x - \text{cikli dolžine 1, } x^2 - \text{cikli dolžine 2.}$$

- V rodovni funkciji so „skrite“ (1)-(3).

1.2 Osnovne Kombinatorične strukture

$$\mathbb{N} = \{0, 1, 2, \dots\}.$$

$$[n] = \{1, 2, \dots, n\}.$$

$$2^A = P(A) = \{B \subseteq A\}.$$

$$\binom{A}{k} = \{B \subseteq A : |B| = k\} \text{ „A nad } k\text{“ (angl. „A choose } k\text{“).}$$

$$\binom{[4]}{2} = \{\{1, 2\}, \{1, 3\}, \dots, \{3, 4\}\}.$$

$$Y^X = \{f : X \rightarrow Y\}.$$

Statistika na množici S je preslikava $S \rightarrow \mathbb{N}$.

$$S = 2^A.$$

Moč je statistika.

S končna množica, f statistika na S .

Pogosto gledamo polinom $\sum_{s \in S} x^{f(s)}$ (enumeration).

$$| \cdot | \text{ na } 2^{[3]} : 1 + 3x + 3x^2 + x^3 = (1 + x)^3.$$

$$S_n = \{\text{permutacije } [n]\} = \{f : [n] \rightarrow [n] : f \text{ bijektivna}\}.$$

$$\begin{pmatrix} 1 & 2 & 3 \\ 2 & 1 & 3 \end{pmatrix} - \text{dvovrstična notacija.}$$

2 1 3 - enovrstična notacija.

(1 2)(3) - produkt disjunktnih ciklov.

$$i, \pi(i), \pi^2(i) \dots$$

$$\text{Gotovo } \exists j_1 < j_2 : \pi^{j_1}(i) = \pi^{j_2}(i) \implies i = \pi^j(i); j > 0.$$

$$(i \ \pi(i) \dots \pi^{j-1}(i)) \text{ cikel.}$$

$$38241765 = (1 \ 3 \ 2 \ 8 \ 5)(4)(6 \ 7) = (4)(2 \ 8 \ 5 \ 1 \ 3)(7 \ 6).$$

Množenje permutacij: kompozicije.

Nekomutativno za $n > 2$.

Disjunktni cikli komutirajo.

Zapis: enoličen do vrstnega reda ciklov in ciklične ureditve ciklov.

Cikel dolžine 1 = negibna točka.

Cikel dolžine 2 = transpozicija.

$(S_n \cdot)$ simetrična grupa.

$$e = id = 1 \ 2 \dots n.$$

π^{-1} inverz (kot preslikava).

$$3\ 8\ 2\ 4\ 1\ 7\ 6\ 5^{-1} = 5\ 3\ 1\ 4\ 8\ 7\ 6\ 2.$$

$$3\ 1\ 4\ 2 \cdot 4\ 2\ 3\ 1 = 2\ 1\ 4\ 3 - \text{množimo z desne.}$$

Statistika: $\#$ ciklov $= c(\pi)$ (štejemo tudi cikle dolžine 1).

$$n = 3 : x^3 + 3x^2 + 2x = x(x+1)(x+2).$$

$$\sum_{\pi \in S_n} x^{c(\pi)} = \sum_k |\{\pi \in S_n : c(\pi) = n\}| x^k.$$

$|\{\pi \in S_n : c(\pi) = n\}| =: c(n, k)$ - Stirlingovo število 1. vrste.

$$\sum_{B \subseteq [n]} x^{|B|} = \sum_k \binom{[n]}{k} x^k.$$

$|\binom{[n]}{k}| =: \binom{n}{k}$ - binomski koeficient.

Inverzija $\pi \in S_n$ je (i, j) , da je za $i < j$ $\pi_i > \pi_j$.

$$\text{inv}(\pi) = \# \text{ inverzij } \pi.$$

$$\text{inv}(4\ 1\ 6\ 2\ 5\ 3) = 7.$$

$$0 \leq \text{inv}(\pi) \leq \binom{n}{2}.$$

Signatura permutacije: $(-1)^{\text{inv}(\pi)}$.

$sg\pi = 1$ - soda permutacija: produkt sodo mnogo transpozicij.

$sg\pi = -1$ - liha permutacija: produkt liho mnogo transpozicij.

$$\det A = \sum_{\pi \in S_n} (-1)^{\text{inv}(\pi)} a_{1, \pi(1)} \cdots a_{n, \pi(n)}.$$

Izraz brez $(-1)^{\text{inv}(\pi)}$: permanenta.

$$n = 3 :$$

$$1 + 2x + 2x^2 + x^3 = 1 + x^2 + x^3 + x + x^2 + x^3 = (1+x)(1+x^2).$$

$$\sum_{\pi \in S_n} x^{\text{inv}(\pi)} = 1 \cdot (1+x)(1+x^2) \cdots (1+x^{n-1}) - \text{kasneje.}$$

$\#$ permutacij v S_n s k inverzijami: ni standardne oznake.

spust/padec (descent) $i : \pi_i > \pi_{i+1}$.

$$\text{des}(4\ 1\ 6\ 2\ 5\ 3) = 3.$$

$$0 \leq \text{des}(\pi) \leq n - 1.$$

$\#$ permutacij v S_n s $k - 1$ spusti $= A(n, k)$ - Eulersko število ($k - 1$ iz zgodovinskih razlogov).

$$\sum_k A(n, k) x^k = \sum_{\pi \in S_n} x^{1+\text{des}(\pi)} = A_n(x) - \text{eulerski polinom.}$$

$$n = 3 :$$

$$x + 4x^2 + x^3.$$

razdelitev/razbitje (angl. set partition) A je $\{B_1, B_2 \dots B_n\}$, davelja :

$$- B_i \neq \emptyset \ i = 1 \dots k,$$

$$- B_i \cap B_j = \emptyset \quad 1 \leq i < j \leq k,$$

$$- \cup_{i=1}^k B_i = A.$$

B_i : bloki razdelitve,

blokov,

razdelitev $[n]$ s k bloki = $S(n,k)$ - Stirlingovo število druge vrste.

$$A = [3] \quad \{\{1\}, \{2\}, \{3\}\}, \{\{1,2\}, \{3\}\} \dots \{\{1,2,3\}\}.$$

$$x + 3x^2 + x^3.$$

$$S(4,2) = 4 + 3 = 7.$$

Kompozicija # n je $\lambda = (\lambda_1 \dots \lambda_l)$, $\lambda_i > 0$ člen kompozicije, $\lambda_i \in \mathbb{N}$,

$$\sum_{i=1}^l \lambda_i = n.$$

$l(\lambda)$ # členov - dolžina.

$\lambda \models n$ - λ je kompozicija n .

Razčlenitev # n je $\lambda = (\lambda_1 \dots \lambda_l)$, $\lambda_i > 0$, $\lambda_i \in \mathbb{N}$.

$$\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_l, \sum_{i=1}^l \lambda_i = n$$

(angl. integer partition).

$p(n)$ - # razčlenitev n .

$p_k(n)$ - # razčlenitev n s k členi.

$$n = 4 :$$

4, 31, 22, 13, 211, 121, 112, 1111 - 8 kompozicij.

4, 31, 22, 221, 1111 - 5 razčlenitev.

$$p(4) = 5, p_2(4) = 2.$$

$B(n) = \sum_k S(n,k)$ - # razčlenitev $[n]$, Bellovo število.

$$B(3) = 5.$$

$L(n,k)$ - razdelitev $[n]$ na k linearno urejenih blokov.

$$L(4,2) = 4 \cdot 6 + 3 \cdot 2 \cdot 2 = 36 - \text{Lahovo število.}$$

$E_n = \#$ alternirajočih permutacij v S_n - Eulerjevo število (Euler number).

$$\pi_1 > \pi_2 < \pi_3 > \pi_4 \dots$$

Primerjaj: eulerska števila (eulerian number).

$$1, 1, 1, 2, 5.$$

Poti:

npr. poti od $(0,0)$ do (n,m) s korakom $(1,0)$ (vzhod) in $(0,1)$ (sever);
 npr. poti od $(0,0)$ do $(2n,0)$ s korakoma $(1,1)$ in $(1,-1)$;
 npr. poti od $(0,0)$ do $(2n,0)$ s korakoma $(1,1)$ in $(1,-1)$, nikoli pod x osjo - Dyckove poti;
 $c_n = \#$ Dyckovih poti dolžine n (konec v $(2n,0)$) - Catalanova števila.
 $1, 1, 2, 5, 14, 42, \dots$
 Drevesa (povezani aciklični grafi).
 $\#$ označenih dreves na n vozliščih.
 Cayleyev izrek: n^{n-2} .
 Ravninska drevesa.
 (Vrstni red pomembnosti).
 Dvojiška drevesa: vsako vozlišče ima 2 ali 0 naslednikov.

1.3 Osnovna načela preštevanja

Načelo vsote: $A \cap B = \emptyset \implies |A \cup B| = |A| + |B|$.

$i \neq j : A_i \cap A_j = \emptyset \implies |\cup_{i=1}^n A_i| = \sum_{i=1}^n |A_i|$.

Načelo produkta: $|A \times B| = |A| \cdot |B|, |\prod_{i=1}^n A_i| = \prod_{i=1}^n |A_i|$.

Kombinatorično:

2 možnosti, izberemo eno ali drugo (ne pa obe) $\implies \#$ načinov je vsota $\#$ načinov,

dvakrat izbiramo, izbiri sta neodvisni $\implies \#$ načinov je produkt $\#$ načinov.

Trditev 1.3.1. $|2^A| = 2^{|A|}$.

Dokaz 1.3.2. Za vsak element se odločimo, ali ga damo v podmnožico ali ne. 2 izbiri, izbiramo $|A|$ -krat, izbire so neodvisne $2 \cdot 2 \cdots 2 = 2^{|A|}$.

$\phi : 2^A \rightarrow \{0,1\}^{|A|}, A = \{a_1, a_2 \dots a_n\}$.

$\phi(B) = (\epsilon_1 \dots \epsilon_n), \epsilon_i = \begin{cases} 1 & a_i \in B \\ 0 & \text{sicer} \end{cases}$

$$\psi : \{0,1\}^{|A|} \rightarrow 2^A.$$

$$\psi(\epsilon_1 \dots \epsilon_n) = \{a_i : \epsilon_i = 1\}.$$

$$\psi \circ \phi, \phi \circ \psi \text{ identiteti.}$$

$$|\{0,1\}^{|A|}| = 2^{|A|}.$$



Trditev 1.3.3.

$$1. \quad |K^N| = |K|^{|N|}.$$

$$2. \quad |\{f \in K^n \text{ injektivna}\}| = |K|(|K| - 1) \dots (|K| - |N| + 1).$$

$$3. \quad |S_n| = n(n-1) \dots 1 = n!.$$

oznake:

$$n^{\underline{k}} = n(n-1) \dots (n-k+1): n \text{ na } k \text{ padajoče.}$$

$$n^{\overline{k}} = n(n+1) \dots (n+k-1): n \text{ na } k \text{ naraščajoče.}$$

Opomba. Pri 2. in 3. smo uporabili varianto načela produkta: izbire sicer niso neodvisne, je pa neodvisno število izbir.

Dirichletov princip (pigeon-hole principle):

$$\phi : X \rightarrow Y \text{ injektivna} \implies |X| \leq |Y|.$$

Če damo n kroglic v k škatel, $n > k$, sta v vsaj eni škatli vsaj 2 kroglici.

Primer.

(1) n ljudi, med njimi sta dva, ki poznata enako mnogo ljudi.

$$X = \text{ljudje}, f = \# \text{ znanstev.}$$

n kroglic, n škatel, ampak škatli 0 in $n-1$ ne moreta biti obe neprazni.

(2) $X \subseteq [2n], |X| = n+1.$

$$\text{Obstajata } x, y \in X, x \neq y, x|y.$$

$$x = 2^k \cdot l, k \geq 0, k \text{ lih.}$$

$$Y = \{i \in [2n] \text{ liho}\}.$$

$$x \mapsto l.$$

1.4 Binomski koeficienti

$\binom{n}{k} = \left| \binom{[n]}{k} \right|$ = število k -elementnih podmnoživ v $[n]$ = število izbir k elementov izmed n elementov.

$$\binom{4}{2} = 6, \binom{5}{0} = 1, \binom{8}{-2} = 0, \binom{8}{9} = 0.$$

$$\binom{n}{0} = 1, \binom{n}{n} = 1, \binom{n}{1} = n.$$

$$\binom{n}{n-k} = \binom{n}{k}.$$

$$\phi : \binom{[n]}{n-k} \rightarrow \binom{[n]}{k}.$$

$$\phi(A) = A^c.$$

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}.$$

$$\binom{n-1}{k-1}: \text{izberemo } n.$$

$$\binom{n-1}{k}: \text{ne izberemo } n\text{-ja.}$$

Pascalov trikotnik:

$$n = 0$$

$$n = 1 \quad \quad \quad 1$$

$$n = 2 \quad \quad \quad 1 \quad 1$$

$$n = 3 \quad \quad \quad 1 \quad 2 \quad 1$$

$$n = 4 \quad \quad \quad 1 \quad 3 \quad 3 \quad 1$$

$$n = 5 \quad \quad \quad 1 \quad 4 \quad 6 \quad 4 \quad 1$$

$$1 \quad 5 \quad 10 \quad 10 \quad 5 \quad 1$$

Trditev 1.4.1. $\binom{n}{k} = \frac{n^k}{k!} = \begin{cases} \frac{n!}{n!(n-k)!} & 0 \leq k \leq n \\ 0 & k > n \end{cases}$

Dokaz 1.4.2. Izberemo 1 element na n načinov, 2 na $n - 1 \dots \implies n^k$ načinov, vsak izbor smo šteli $k!$ -krat.

Ali: preštejemo urejene izbire k različnih elementov iz $[n]$;

$$n^k = \binom{n}{k} \cdot k!.$$

$$\binom{n}{k}: \text{najprej izberemo } k \text{ elementov.}$$

k : nato jih uredimo. ■

Izrek 1.4.3 (Binomski izrek). $(a + b)^n = \sum_{k=0}^n \binom{n}{k} a^k b^{n-k}$;
 $a, b \in K$ komutativni kolobar, $n \in \mathbb{N}$.

Dokaz 1.4.4.

D1. Indukcija po n :

$$n = 0: 1 = 1$$

$$n - 1 \rightarrow n:$$

$$\begin{aligned} (a+b)^n &= (a+b)^{n-1}(a+b) = \\ &\stackrel{\text{IP}}{=} \sum_{k=0}^{n-1} \binom{n-1}{k} a^k b^{n-1-k} (a+b) = \\ &= \sum_{k=0}^{n-1} \binom{n-1}{k} a^{k+1} b^{n-1-k} + \sum_{k=0}^{n-1} \binom{n-1}{k} a^k b^{n-k} = \\ &= \sum_{k=1}^{n-1} \binom{n-1}{k-1} a^k b^{n-k} + \sum_{k=1}^{n-1} \binom{n-1}{k} a^k b^{n-k} = \\ &= \sum_{k=0}^n \binom{n-1}{k-1} a^k b^{n-k} + \sum_k \binom{n-1}{k} a^k b^{n-k} = \\ &= \sum_{k=0}^n \binom{n}{k} a^k b^{n-k}. \end{aligned}$$

$$\text{D2. } (a+b)^n = \sum_k \binom{n}{k} a^k b^{n-k} \text{ DN.}$$

$$\begin{aligned} \text{D3. } (a+b) \dots (a+b) &= \sum_{\text{izbira } a \text{ ali } b} \text{produkt izbranih} = \\ &= \sum_k \binom{n}{k} a^k b^{n-k}. \end{aligned}$$

a izberemo k -krat.

Izberemo k oklepajev, pri katerih izberemo a .



$$\begin{aligned} \binom{10}{3} &= \frac{10 \cdot 9 \cdot 8}{3 \cdot 2} = 120. \\ \binom{12}{10} &= \binom{12}{2} = \frac{12 \cdot 11}{2} = 66. \end{aligned}$$

Izbori: n kroglic, k izberemo.

	s ponavljanjem	brez ponavljanja	
vrstni red pomemben	n^k	$n^{\underline{k}}$	variacije
ni pomemben	$\binom{n+k-1}{k}$	$\binom{n}{k}$	kombinacije

$$1 \leq i_1 \leq i_2 \leq \dots \leq i_k \leq n.$$

$$j_1 = i_1, j_2 = i_2 + 1 \dots j_k = i_k + k - 1.$$

$$1 \leq j_1 < j_2 < \dots < j_k \leq n + k - 1.$$

Trditev 1.4.5. Število kompozicij n je 2^{n-1} ($n \geq 1$), število kompozicij s k členi je $\binom{n-1}{k-1}$ ($n \geq 1$).

Dokaz 1.4.6. n kroglic $\circ | \circ \circ \circ | \circ \circ : 6 = 1 + 3 + 2.$

$k - 1$ pregrad, $n - 1$ mest za pregrade. ■

Kompozicije: $2^{n-1}, \binom{n-1}{k-1}.$

Šibka kompozicija: $(\lambda_1 \dots \lambda_l); \lambda_i \geq 0, \lambda_1 + \dots + \lambda_l = n.$

$3 : 12, 3, 21, 102, 300, 0102 \dots$

Število šibkih kompozicij n s k členi.

$n + k - 1$ objektov, premešamo na $\binom{n+k-1}{k-1}$ oz. $\binom{n+k-1}{n}$ načinov.

Še en dokaz:

$$\lambda_1 + \dots + \lambda_l = n, \lambda_i \geq 0.$$

$$\mu_i = \lambda_i + 1 \mu_i \geq 1.$$

$$\mu_1 + \dots + \mu_l = n + k \implies \binom{n+k-1}{n-1}.$$

Primerjaj z: kombinacije s ponavljanjem.

n kroglic, k -krat izbiram.

λ_i : kolikokrat izberemo i -to kroglico.

$$\lambda_1 + \dots + \lambda_n = k, \lambda_i \geq 0.$$

Šibke kompozicije k z n členi: $\binom{k+n-1}{k}.$

Trditev 1.4.7.

$$L(n, k) = \frac{n!}{k!} \binom{n-1}{k-1}.$$

Dokaz 1.4.8. Koliko je urejenih razdelitev na linearno urejene bloke:

$$k! \cdot L(n, k) = n! \cdot \binom{n-1}{k-1}.$$

Tukaj predstavljajo

- $L(n, k)$: urejene bloke,

- $k!$: njihov vrstni red,
- $n!$: permutacije,
- $\binom{n-1}{k-1}$: šibke kompozicije.

Poti iz $(0,0)$ v (n,m) , premikamo se gor ali desno.

n -krat gor, m -krat desno: $\binom{n+m}{m}$ možnosti.

Poti iz $(0,0)$ v $(2n,0)$, desno-gor ali desno-dol.

n -krat gor, n -krat dol: $\binom{2n}{n}$.

Dyckove poti: isto kot prej, se ne spustimo pod x -os.

Pot je slaba, če gre pod x -os:

Od tam naprej, kjer 1. doseže $y = -1$, prezrcalimo pot preko $y = -1$.

Konča se v $y = -2$.

Število slabih poti = število poti od $(0,0)$ do $(2n, -2)$.

Teh je $\binom{2n}{n-1}$: $(n-1)$ -krat gor, $(n+1)$ -krat dol.

$$C_n = \text{število Dyckovih poti dožine } n = \binom{2n}{n} - \binom{2n}{n-1} \\ = \frac{(2n!)}{n!n!} - \frac{(2n)!}{(n-1)!(n+1)!} = \binom{2n}{n} \left(1 - \frac{n}{n+1}\right) = \frac{1}{n+1} \binom{2n}{n}.$$

Multinomski koeficienti:

$$\alpha_1 \times 1, \alpha_2 \times 2 \dots \alpha_k \times k : 11..12..2..k.$$

Na koliko načinov lahko premešamo:

$$\binom{\alpha_1 + \dots + \alpha_k}{\alpha_1} \binom{\alpha_2 + \dots + \alpha_k}{\alpha_2} \dots \binom{\alpha_k}{\alpha_k} = \frac{(\alpha_1 + \dots + \alpha_k)!}{\alpha_1! \dots \alpha_k!}.$$

Definiramo

$$\binom{\alpha_1 + \dots + \alpha_k}{\alpha_1, \alpha_2 \dots \alpha_k} := \frac{(\alpha_1 + \dots + \alpha_k)!}{\alpha_1! \dots \alpha_k!}. \quad (1.1)$$

Izrazu 1.1 pravimo multinomski simbol.

Figure v 1. vrsti pri šahu: $\frac{8!}{1!1!2!2!} = 7!$.

i -jem damo indekse $\alpha_1 \dots \alpha_k : 1_1 \dots 1_{\alpha_1} 2_1 \dots k_{\alpha_k}$

Premešamo na $(\alpha_1 + \dots + \alpha_k)!$ načinov.

Eno permutacijo dobimo $(\alpha_1! \dots \alpha_k!)$ -krat.

Multimnožica M je množica, v kateri se elementi lahko ponavljajo.

$$M = \{1, 1, 1, 2, 2, 3, 3, 3, 3\} = \{1^3, 2^2, 3^4\}.$$

Število permutacij multimnožice je multinomski simbol.

Formalno je multimnožica (S, f) , kjer je S množica, $f : S \rightarrow \mathbb{N}$ šteje kolikokrat se posamezen element ponovi.

1.5 Dvanajstera pot

n kroglic, k škatel; na koliko načinov lahko damo kroglice v škatle.

$N \setminus K$	vse	injekcije	surjekcije	
L L	k^n	k^n	$k!S(n, k)$	„kompozicije“
N L	$\binom{n+k-1}{k-1}$	$\binom{k}{n}$	$\binom{n-1}{k-1}$	
L N	$\sum_i S(n, i)$	$\begin{cases} 1 & k \geq n \\ 0 & \text{sicer} \end{cases}$	$S(n, k)$	razdelitve
N N	$\overline{p_k(n)}$	$\begin{cases} 1 & k \geq n \\ 0 & \text{sicer} \end{cases}$	$p_k(n)$	razčlenitve

Vpeljemo ekvivalenčne relacije

- $f \sim_N g : \exists \pi \in S_n : f \circ \pi = g$
- $f \sim_K g : \exists \sigma \in S_k : \sigma \circ f = g$
- $f \sim_{N,k} g : \exists \pi \in S_n, \sigma \in S_k : \sigma \circ f \circ \pi = g.$

1.6 Rekurzije

$$\binom{n}{k} = \binom{n-1}{k-1} + \binom{n-1}{k}.$$

$$c(n, k) = c(n-1, k-1) + (n-1)c(n-1, k);$$

$c(n-1, k-1)$: n negibna, $(n-1)$: za kateri element vstavimo.

$$S(n, k) = S(n-1, k-1) + kS(n-1, k);$$

$S(n-1, k-1)$: n v svojem bloku, k : v kateri blok vstavimo.

$$L(n, k) = L(n-1, k-1) + (n+k-1)L(n-1, k);$$

$L(n-1, k-1)$: n v svojem bloku, $(n+k-1)$: kam vstavimo.

$$B(n+1) = \sum_{k=0}^n \binom{n}{k} B(n-k) = \sum_{k=0}^n \binom{n}{k} B(k);$$

odstranimo blok, v katerem je $n+1$, k : število elementov v bloku skupaj

$z\ n + 1, \binom{n}{k}$: kateri elementi v bloku skupaj $z\ n + 1, B(n - k)$: razdelimo ostale.

$$p_k(n) = p_{k-1}(n - 1) + p_k(n - k);$$

$p_{k-1}(n - 1)$: $\lambda_l = 1$, $p_k(n - k)$: $\lambda_l \geq 2$ (odstranimo 1. stolpec v Ferrersovem diagramu).

$A(n, k) = (n + 1 - k)A(n - 1, k - 1) + kA(n - 1, k)$. odstranimo n, k : n damo na konec ali za spust, $(n + 1 - k)$: n damo na začetek ali za vzpon. V S_n velja še: število spustov + število vzponov = $n - 1$.

$$2E_{n+1} = \sum_{k=0}^n \binom{n}{k} E_k E_{n-k} \quad n \geq 1;$$

k : koliko elementov je pred $n + 1$, število obratno alternirajočih = število alternirajočih ($i \rightarrow n + 1 - i$), E_k : pred $n + 1$, E_{n-k} : za $n + 1$, štejemo in alternirajoče in obratno alternirajoče permutacije.

$$C_{n+1} = \sum_{k=0}^n C_k C_{n-k};$$

k : ko 1. pridemo v $y = 0$: pred in za tem sta Dyckovi poti.

$$p(n) = p(n - 1) + p(n - 2) - p(n - 5) - p(n - 7) + p(n - 12) + p(n - 15) - \dots$$

Eulerjev petkotniški izrek (dokaz kasneje) (pentagonal).

1.7 Načelo vključitev in izključitev (NVI)

(Principle of inclusion and exclusion).

$$|A \cup B| = |A| + |B| - |A \cap B|.$$

$$|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|.$$

Izrek 1.7.1 (NVI).

$$\begin{aligned}
|\cup_{i=1}^n A_i| &= |A_1| + |A_2| + \cdots + |A_n| \\
&\quad - |A_1 \cap A_2| - \cdots - |A_{n-1} \cap A_n| \\
&\quad + |A_1 \cap A_2 \cap A_3| + \cdots + |A_{n-2} \cap A_{n-1} \cap A_n| \\
&\quad - \cdots \\
&= \sum_{i=1}^n (-1)^{i-1} \sum_{1 \leq j_1 < \cdots < j_k \leq n} |A_{j_1} \cap \cdots \cap A_{j_k}| \\
&= \sum_{\emptyset \neq S \subseteq [n]} (-1)^{|S|-1} |A_S|,
\end{aligned}$$

kjer je $A_S := \cap_{i \in S} A_i$.

Dokaz 1.7.2.

$x \in \cup_{i=1}^n A_i$.

Trdimo, da x prispeva 1 k vsoti na desni.

Recimo, da je x v natanko m množicah A_i ($1 \leq m \leq n$):

$$\begin{aligned}
&m - \binom{m}{2} + \binom{m}{3} - \cdots + (-1)^m \binom{m}{m} \\
&= 1 - \left(\binom{m}{0} - \binom{m}{1} + \binom{m}{2} - \cdots + (-1)^{m-1} \binom{m}{m} \right) \\
&= 1 - (1 - 1)^m = 1.
\end{aligned}$$

Trditev 1.7.3 (NVI, 2. verzija).

$$|\cap_{i=1}^n A_i^C| = \sum_{S \subseteq [n]} |A_S|.$$

Dokaz 1.7.4.

$$\begin{aligned}
|\cap_{i=1}^n A_i^C| &= |(\cup_{i=1}^n A_i)^C| \\
&= |A| - |\cup_{i=1}^n A_i| \\
&= |A| + \sum_{\emptyset \neq S \subseteq [n]} (-1)^{|S|} |A_S| \\
&= \sum_{S \subseteq [n]} |A_S|,
\end{aligned}$$

kjer velja še $A_\emptyset = A$.

Primer.

(1) Koliko je k -elementnih antiverig v B_n ?

$B_n = (2^{[n]}, \subseteq)$ Boolova algebra, antiveriga - množica neprimerljivih elementov.

$k=1$: 2^n (vsi elementi).

$k=2$:

$$S = \{(A, B) : A, B \subseteq [n]\}$$

$$S_1 = \{(A, B) : A \subseteq B\}$$

$$S_2 = \{(A, B) : B \subseteq A\}$$

$$|S_1^C \cap S_2^C| = |S| - |S_1| - |S_2| + |S_1 \cap S_2| = 4^n - 2 \cdot 3^n + 2^n;$$

4^n : vse možnosti $x \in, \notin A, B$, 3^n : vse razen $x \in A, \notin B \dots$

$$\implies \frac{1}{2}(4^n - 2 \cdot 3^n + 2^n).$$

$k=3$:

$$S = \{(A, B, C) : A, B, C \in 2^{[n]}\}$$

$$S_1 : A \subseteq B, S_2 : B \subseteq A, S_3 : A \subseteq C, S_4 : C \subseteq A$$

$$S_5 : B \subseteq C, S_6 : C \subseteq B.$$

$$|\cap_{i=1}^6 S_i^C| = 8^n - 6 \cdot 6^n + 3 \cdot 4^n + 6 \cdot 5^n - 6 \cdot 4^n - \dots$$

$$6^n : S_i, 4^n : \text{npr. } S_1 \cap S_2, 5^n : \text{npr. } S_1 \cap S_3, 4^n : \text{npr. } S_1 \cap S_4.$$

(2) i_n : število premestitev v S_n = število permutacij v S_n brez negibne

točke (dearangement).

$$\begin{aligned}
 A &= S_n \\
 A_i &= \{\pi \in S_n : \pi_i = i\} \\
 |A_I| &= (n - |I|)! \\
 i_n &= \sum_{I \subseteq [n]} (-1)^{|I|} (n - |I|)! \\
 &= \sum_{k=0}^n \binom{n}{k} (-1)^k (n - k)! \\
 &= n! \sum_{k=0}^n \frac{(-1)^k}{k!}.
 \end{aligned}$$

$$P(\text{število premestitev}) = \sum_{k=0}^n \frac{(-1)^k}{k!} \xrightarrow{n \rightarrow \infty} e^{-1}.$$

(3) Število surjekcij iz $[n]$ v $[k]$.

$$\begin{aligned}
 A &= [k]^{[n]} \\
 A_i &= ([k] \setminus \{i\})^{[n]} \\
 \left| \cap_{i=1}^n A_i^C \right| &= \sum_{I \subseteq [n]} (-1)^{|I|} (k - |I|)^n \\
 &= \sum_{k=1}^n \binom{k}{i} (-1)^i (k - i)^n \\
 &\stackrel{i=k-i}{=} \sum_{j=1}^k (-1)^{k-j} \binom{k}{j} j^n \\
 &= k! S(n, k);
 \end{aligned}$$

surjekcija je urejena razdelitev;

$$S(n, k) = \sum_{j=0}^n \frac{(-1)^{k-j} j^n}{j! (k-j)!}.$$

(4) Eulerjev petkotniški izrek:

$$p(n) = p(n-1) + p(n-2) - p(n-5) - \dots$$

$$A = \{\text{razčlenitve } n\}$$

$$A_i = \{\text{razčlenitve } n, \text{ ki vsebujejo } i \text{ za člen}\} \quad i = 1, 2, \dots, n$$

$$|A_i| = p(n - i)$$

$$|A_i \cap A_j| = p(n - k - j)$$

$$|A_I| = p(n - \sum_{i \in I} i)$$

$$\begin{aligned} p(n) &= p(n - 1) + p(n - 2) + p(n - 3) + \dots \\ &\quad - p(n - 1 - 2) - p(n - 1 - 3) - p(n - 2 - 3) - \dots \\ &\quad + p(n - 1 - 2 - 3) - \dots \\ &= p(n - 1) + p(n - 2) - p(n - 5) - p(n - 7) + \dots \end{aligned}$$

Franklinova bijekcija:

$$p(n) = \sum_{m=1}^{\infty} (\alpha(m) - \beta(m))p(n - m); \quad m - \text{razčlenitve z različnimi členi,}$$

$$\alpha(m) = \text{število razčlenitev } m \text{ z liho mnogo različnimi členi,}$$

$$\beta(m) = \text{število razčlenitev } m \text{ z sodo mnogo različnimi členi,}$$

Bijekcija

$$\begin{aligned} \Phi : \{\text{razčlenitev } m \text{ z liho mnogo različnimi členi}\} \setminus \{\dots\} \\ \rightarrow \{\text{razčlenitev } m \text{ z sodo mnogo različnimi členi}\} \setminus \{\dots\}. \end{aligned}$$

$$f(\lambda) = \max\{i : \lambda_i = \lambda_1 - i + 1\} - \text{bok},$$

$$g(\lambda) = \lambda_{l(\lambda)} - \text{najmanjši člen},$$

$$\text{a) } f(\lambda) \geq g(\lambda): \text{ min} \rightarrow \text{bok},$$

$$\text{b) } f(\lambda) < g(\lambda): \text{ bok} \rightarrow \text{min},$$

$$\text{a) ne dela (število členov se ohrani),}$$

$$\text{b) ne dela (2 člena enako dolga),}$$

$$\text{a) ne dela, ko:}$$

$$f(\lambda) = g(\lambda) = l(\lambda)$$

$$m = k + (k + 1) + \dots + (2k - 1) = \frac{2k(2k-1)}{2} - \frac{k(k-1)}{2} = \frac{k(3k-1)}{2}$$

$$(\alpha(m) - \beta(m)) = (-1)^{k-1} \quad (k \text{ lih ali sod}).$$

$$\text{b) ne dela, ko:}$$

$$f(\lambda) = g(\lambda) - 1 = l(\lambda)$$

$$m = (k+1) + (k+2) + \dots + (2k) = \dots = \frac{k(3k+1)}{2}$$

$$(\alpha(m) - \beta(m)) = (-1)^{k-1}.$$

Eulerjev petkotniški izrek:

$$p(n) = \sum_{k=1}^{\infty} (-1)^{k-1} \left(p\left(n - \frac{k(3k-1)}{2}\right) + p\left(n - \frac{k(3k+1)}{2}\right) \right)$$

$$\text{oz. } \sum_{k \in \mathbb{Z}} (-1)^k p\left(n - \frac{k(3k+1)}{2}\right) = 0.$$

Tukaj smo upoštevali ko vstavimo $-k$: $\frac{-k(-3k-1)}{2} = \frac{k(3k+1)}{2}$ in $p(0) = 0$.

Izrek 1.7.5 („NVI“).

$f, g : B_n \rightarrow K$, K komutativni kolobar.

$$f(T) = \sum_{S \subseteq T} g(S) (\forall T \in B_n) \iff g(T) = \sum_{S \subseteq T} (-1)^{|T \setminus S|} f(S) (\forall T \in B_n).$$

Zgled.

$$des(\pi) = |\{i : \pi(i) > \pi(i+1)\}|$$

$$D(\pi) = \{i : \pi(i) > \pi(i+1)\}$$

$$D(1\ 4\ 2\ 6\ 5\ 3) = \{2, 4, 5\}$$

$$f_n(T) = |\{\pi \in S_n : D(\pi) = T\}|$$

$$\text{npr. } n = 8, T = \{1, 5\}$$

$$g_n(T) = |\{\pi \in S_n : D(\pi) \subseteq T\}|$$

$$T = \{t_1, t_2 \dots t_k\}$$

$$g_n(T) = \binom{n}{t_1} \binom{n-t_1}{t_2-t_1} \binom{n-t_1-\dots-t_{k-1}}{t_k} = \binom{n}{t_1, t_2-t_1, \dots, t_k-t_{k-1}, n-t_k}$$

$_ < _ < _ < \underline{t_i} \leq _ :$ zaradi \subseteq : tam lahko spust ali pa ne.

// če lastnosti točno določene: težko $(f_n(T))$, če „vsebovano“ $(g_n(T))$: lažje

$$g_n(T) = \sum_{S \subseteq T} f_n(S)$$

$$\begin{aligned}
f_n(T) &= \sum_{S \subseteq T} (-1)^{|T \setminus S|} g_n(S) \\
&= \sum_{S \subseteq T} (-1)^{|T \setminus S|} \binom{n}{s_1, s_2 - s_1, \dots, n - s_k} \\
&\stackrel{\text{vaje}}{=} \det \left[\binom{n - t_i}{t_{j+1} - t_j} \right]_{i,j=0}^{|T|}.
\end{aligned}$$

npr. $n = 8$, $T = \{1, 5\}$, $t_0 = 0$, $t_{|T|} = n + 1 = 9$

$$f_8(\{1, 5\}) = \begin{vmatrix} \binom{8}{1} & \binom{8}{5} & \binom{8}{8} \\ \binom{7}{0} & \binom{7}{1} & \binom{7}{7} \\ \binom{3}{-4} & \binom{3}{0} & \binom{3}{3} \end{vmatrix} = 217.$$

Dokaz 1.7.6.

(\implies):

$$\begin{aligned}
\sum_{S \subseteq T} (-1)^{|T \setminus S|} f(S) &= \sum_{S \subseteq T} (-1)^{|T \setminus S|} f(S) \sum_{U \subseteq S} g(U) \\
&= \sum_{U \subseteq T} \left(\sum_{U \subseteq S \subseteq T} (-1)^{|T \setminus S|} \right) g(U) \\
&\stackrel{k=|S \setminus U|}{=} \sum_{U \subseteq T} \sum_{k=0}^{|U|} \binom{|T \setminus U|}{k} (-1)^{|T \setminus U| - k} g(U) \\
&= g(T).
\end{aligned}$$

Na notranji vsoti uporabimo binomski izrek za -1 in 1 :

$$(1 - 1)^{|T \setminus S|} = \begin{cases} 1 : U = T \\ 0 : U \subset T \end{cases}$$

1.8 Polinomske enкости

$$(1 + x)^n = \sum_{k=0}^n \binom{n}{k} x^k$$

Izrek 1.8.1.

- (a) $\sum_k c(n,k)x^k = x^{\bar{n}}$
- (b) $\sum_k (-1)^{n-k} c(n,k)x^k = x^{\underline{n}}$
- (c) $\sum_k S(n,k)x^{\underline{k}} = x^{\bar{n}}$
- (d) $\sum_k (-1)^{n-k} S(n,k)x^{\bar{k}} = x^{\underline{n}}$
- (e) $\sum_k L(n,k)x^{\underline{k}} = x^{\bar{n}}$
- (f) $\sum_k (-1)^{n-k} L(n,k)x^{\bar{k}} = x^{\underline{n}}$

Opomba. $K[x] = \{\text{polinomi v } x\}$ vektorski prostor (celo algebra), K komutativen obseg.

$\{x^{\bar{n}}\}, \{x^{\underline{n}}\}, \{x^{\bar{n}}\}$ naravne baze.

Dokaz 1.8.2.

(a) Indukcija (na vajah drugače):

$$n = 0: 1=1$$

$$n - 1 \rightarrow n:$$

$$\begin{aligned} x^{\bar{n}} &= x^{\overline{n-1}}(x + n - 1) \stackrel{\text{IP}}{=} (x + n - 1) \sum_k c(n-1, k)x^k \\ &= \sum_k c(n-1, k-1)x^k + (n-1) \sum_k c(n-1, k)x^k = \sum_k c(n, k)x^k, \end{aligned}$$

(b) $x \rightarrow -x$ v (a),

(c) Preslikava = razdelitev + injekcija,

število preslikav iz $[n]$ v $[k] = \sum_k S(n, k)x^{\underline{k}}$, kjer predstavljajo

- k : število blokov,
- $S(n, k)$: razdelimo $[n]$ na k blokov,
- $x^{\underline{k}}$: injekcija $[k] \rightarrow [x]$.

Dokazali smo za $x \in \mathbb{N} \implies$ polinoma sta enaka (ujemanje v ∞ točkah).

(e) Z indukcijo DN.

$$\pi = 4 \ 2 \ 5 \ 1 \ 6 \ 3$$

$$\text{inv}(\pi) = 7$$

$$I(\pi) = \{(1,2), (1,4), (1,6) \dots\}$$

$TI(\pi) = (a_1 \dots a_n); a_k = \{(i,j) : \pi_i > \pi_j = k\}$ („desna stran“) - tabela inverzij.

$$TI(\pi) = (3,1,3,0,0,0)$$

$0 \leq a_i \leq n - i$, a_i : koliko levo od i večjih od i .

Trditev 1.8.3.

$TI : S_n \rightarrow [0, n-1] \times [0, n-2] \times \dots \times [0, 0]$ je bijekcija.

Posledica 1.8.4.

$$\sum_{\pi \in S_n} q^{\text{inv}(\pi)} = \underline{n!} = (1+q)(1+q+q^2) \dots (1+q+\dots+q^{n-1}).$$

$$\pi = 4 \ 1 \ 7 \ 3 \ 9 \ 6 \ 2 \ 8 \ 5,$$

$$TI(\pi) = (1, 5, 2, 0, 4, 2, 0, 1, 0),$$

$$\text{inverz: } 9 \rightarrow 9 \ 8 \rightarrow 7 \ 9 \ 8 \rightarrow 7 \ 9 \ 6 \ 8 \rightarrow 7 \ 9 \ 6 \ 8 \ 5 \rightarrow 4 \ 7 \ 9 \ 6 \ 8 \ 5$$

$$\rightarrow 4 \ 7 \ 3 \ 9 \ 6 \ 8 \ 5 \rightarrow 4 \ 7 \ 3 \ 9 \ 6 \ 2 \ 8 \ 5 \rightarrow 4 \ 1 \ 7 \ 3 \ 9 \ 6 \ 2 \ 8 \ 5.$$

Dokaz 1.8.5. trditve.

Skonstruiramo inverz:

$$(a_1 \dots a_n), \ 0 \leq a_i \leq n - i.$$

Vpisujemo $n, n-1 \dots 1$: i pišemo za a_i elementi.

Dokaz 1.8.6. posledice.

$\sum_{\pi \in S_n} q^{\text{inv}(\pi)} = n!_q = \underline{n!} = \underline{n(n-1)} \dots 1$ - q fakulteta, $\underline{i} = 1 + q + \dots + q^{i-1}$ - polinom, q-naravno število (q-integer).

$$\begin{aligned} D &= (1 + q + \dots + q^{n-1})(1 + q + \dots + q^{n-2}) \dots 1 \\ &= \sum_{0 \leq a_i \leq n-i} q^{a_1} q^{a_2} \dots q^{a_n} \\ &\stackrel{\text{trditev}}{=} \sum_{\pi \in S_n} q^{\text{inv}(\pi)}. \end{aligned}$$

Opomba. $\text{maj}(\pi) = \sum_{i \text{ spust } \pi} i$ oz. $\sum_{i \in D(\pi)} i$ - majorski indeks

$$\text{maj}(4 \ 2 \ 5 \ 1 \ 3) = 1 + 3 = 4$$

$$\sum_{\pi \in S_n} q^{\text{maj}(\pi)} = \sum_{\pi \in S_n} q^{\text{inv}(\pi)} = \underline{n}!$$

Definicija 1.8.7 (q-binomski koeficient).

$$\binom{\underline{n}}{\underline{k}} = \binom{n}{k}_q = \frac{\underline{n}!}{\underline{k}! \underline{(n-k)}!}.$$

$$\begin{aligned} \binom{\underline{n}}{\underline{0}} &= \binom{n}{n} = 1 \\ \binom{\underline{n}}{\underline{1}} &= \underline{n} \binom{\underline{4}}{\underline{2}} = \frac{(1+q+q^2+q^3)(1+q+q^2)(1+q)}{(1+q)(1+q)} = (1+q^2)(1+q+q^2) \quad q = 1 : \binom{\underline{n}}{\underline{k}} = \binom{n}{k}. \end{aligned}$$

Trditev 1.8.8.

$$\binom{\underline{n}}{\underline{k}} = q^{n-k} \binom{\underline{n-1}}{\underline{k-1}} + \binom{\underline{n-1}}{\underline{k}} = \binom{\underline{n-1}}{\underline{k-1}} + q^k \binom{\underline{n-1}}{\underline{k}}.$$

Dokaz 1.8.9.

$$\begin{aligned} & q^{n-1} \frac{(\underline{n-1})!}{(\underline{k-1})! \underline{(n-k)}!} + \frac{(\underline{n-1})!}{(\underline{k})! \underline{(n-1-k)}!} \\ &= \frac{\underline{n}!}{\underline{k}! \underline{(n-k)}!} (q^{n-k} \underline{k}! + \underline{n-k}) \\ &= \frac{\underline{n}!}{\underline{k}! \underline{(n-k)}!} \\ &= \binom{\underline{n}}{\underline{k}}, \end{aligned}$$

kjer je

$$q^{n-k} \underline{k}! + \underline{n-k} = q^{n-k} + \dots + q^n + 1 + \dots + q^{n-k-1} = 1 + q + \dots + q^n.$$

Posledica 1.8.10. $\binom{\underline{n}}{\underline{k}}$ je polinom v q .

Trditev 1.8.11.

$$\prod_{i=1}^n (1 + q^{i-1}x) = \sum_{k=0}^n \binom{\underline{n}}{\underline{k}} x^k.$$

Dokaz 1.8.12. Indukcija:

$$n = 0 : 1 = 1$$

$$n - 1 \rightarrow n:$$

$$\begin{aligned} \prod_{i=1}^n (1 + q^{i-1}x) &= \left(\sum_{k=0}^n \binom{n-1}{k} x^k \right) \cdot (1 + q^{n-1}x) \\ &= \sum_k q^{\binom{k}{2}} \binom{n-1}{k} x^k + \sum_k q^{\binom{k-1}{2} + n-1} \binom{n-1}{k-1} x^k \\ &= \sum_k q^{\binom{k}{2}} \left(\binom{n-1}{k} + q^{\binom{k-1}{2} + n-1 - \binom{k}{2}} \binom{n-1}{k-1} \right) x^k. \end{aligned}$$

$$\text{Upoštevali smo } \binom{k-1}{2} - \binom{k}{2} = -\binom{k-1}{1}.$$

\mathbb{Z}_p, p praštevilo končen obseg.

Izrek 1.8.13. Obseg moči $n \in \mathbb{N}$ obstaja $\iff n = p^k$ p praštevilo. Obseg je do izomorfizma natančno določen.

\mathbb{F}_q - oznaka.

Izrek 1.8.14. V \mathbb{F}_q^n je $\binom{n}{k}$ k -dimenzionalnih podprostorov.

$$\text{Primer. } q = 4, n = 4, k = 2 : (1 + 4^2) + (1 + 4 + 4^2) = 38.$$

Dokaz 1.8.15. Spomnimo se: $[n]$ ima $\binom{n}{k}$ k -podmnožic, štejemo urejene k -terice različnih števil: $k! \binom{n}{k} = n^{\underline{k}}$.

Štejemo k -terice linearno neodvisnih vektorjev v \mathbb{F}_q^n :

$$(q^k - 1)(q^k - q) \dots (q^k - q^{k-1})X = (q^n - 1)(q^n - q) \dots (q^n - q^{n-1});$$

$q^k - q^i$: vsi v podprostoru brez linearnih kombinacij že vzetih,

$q^n - q^i$: vsi brez linearnih kombinacij že vzetih.

X : število izbir podprostora.

$$X = \frac{q^{\binom{k}{2}} (q-1)^k n(n-1) \dots (n-k+1)}{q^{\binom{k}{2}} (q-1)^k k!} = \binom{n}{k}.$$

Definicija 1.8.16 (q-multinomski koeficient).

$$\begin{aligned} \binom{a_1 + \dots + a_k}{\underline{a_1}, \underline{a_2} \dots \underline{a_k}} &= \frac{(a_1 + \dots + a_k)!}{\underline{a_1}! \dots \underline{a_k}!} \\ &= \binom{a_1 + \dots + a_k}{\underline{a_1}} \binom{a_2 + \dots + a_k}{\underline{a_2}} \dots \binom{a_k}{\underline{a_k}}. \end{aligned}$$

\Rightarrow je polinom (produkt polinomov).

$x_1 \dots x_n$ permutacija multimnožice $\{1^{a_1}, 2^{a_2} \dots n^{a_n}\}$

inverzija: $(i, j) : i < j, x_i > x_j$

inv : število inverzij

$inv(1 \ 2 \ 1 \ 1 \ 2 \ 3) = 2$.

Izrek 1.8.17. $M = \{1^{a_1}, 2^{a_2} \dots n^{a_n}\}$

$$\sum_{\pi \in S(M)} q^{inv(\pi)} = \binom{a_1 + \dots + a_n}{\underline{a_1} \dots \underline{a_n}}.$$

Primer.

$$q = 1 : |S(M)| = \binom{a_1 + \dots + a_n}{a_1 \dots a_n}$$

$a_1 = \dots = a_n = 1 : \sum_{\pi \in S_n} q^{inv(\pi)} = n!$ - posplošitev formul za multinomske koeficiente in Stirlingova števila 1. vrste.

Dokaz 1.8.18.

$$\begin{aligned} \sum_{\pi \in S(M)} q^{inv(\pi)} \underline{a_1}! \dots \underline{a_n}! &= \underline{(a_1 + \dots + a_n)!} \\ \sum_{\pi_0 \in S(M)} q^{inv(\pi_0)} \cdot \sum_{\pi_1 \in S_{a_1}} q^{inv(\pi_1)} \dots \sum_{\pi_n \in S_{a_n}} q^{inv(\pi_n)} &= \sum_{\pi \in S_{a_1 + \dots + a_n}} q^{inv(\pi)}. \end{aligned}$$

Iščemo bijekcijo

$$\begin{aligned} \Phi : (\pi_0 \pi_1 \dots \pi_n) &\rightarrow \pi \\ S(M) S_{a_1} \dots S_{a_n} &\mapsto S_{a_1 + \dots + a_n}. \end{aligned}$$

$$M = \{1^4, 2^2, 3^3\}$$

$$(1 \ 2 \ 2 \ 1 \ 3 \ 1 \ 3 \ 3 \ 1, 2 \ 4 \ 1 \ 3, 2 \ 1, 1 \ 3 \ 2)$$

$\mapsto 2\ 6\ 5\ 4\ 7\ 1\ 9\ 8\ 3$.

V π_0 enke spremenimo v $1 \dots a_1$ v vrstnem redu, ki ga določa π_1 , v π_0 dvojke spremenimo v $a_1 + 1 \dots a_2$ v vrstnem redu, ki ga določa π_2 , itn.

$$\text{inv}(\pi_0) + \dots + \text{inv}(\pi_n) = \text{inv}(\Phi(\pi_0 \dots \pi_n)).$$

Vsaka inverzija $\Phi(\pi_0 \dots \pi_n)$ prihaja bodisi od inverzije π_i bodisi od inverzije π_0 (glede na „indeks“ v π_0) \implies vsota enaka.

Poglavje 2

Formalne potenčne vrste

2.1 Uvod

$$\sum_k c(n,k)x^k = x^{\overline{n}}$$

$\sum_n S(n,k)x^n$ neskončna vsota.

V analizi: potenčne vrste:

$$F(x) = \sum_{n=0}^{\infty} a_n x^n.$$

Konvergira za $|x| < R$ - konvergenčni polmer:

$$R = \frac{1}{\limsup_{n \rightarrow \infty} \sqrt[n]{|a_n|}} \stackrel{\text{če obstaja}}{=} \lim_{n \rightarrow \infty} \left| \frac{a_n}{a_{n+1}} \right| \in [0, \infty].$$

Primer. $\sum_{n=0}^{\infty} x^n : R = 1$

$$\sum_{n=0}^{\infty} \frac{x^n}{n!} : R = \lim_{n \rightarrow \infty} \frac{\frac{1}{n!}}{\frac{1}{(n+1)!}} = \frac{(n+1)!}{n!} = \infty$$

$$\sum_{n=0}^{\infty} n! x^n : R = 0$$

$\sum_{n=0}^{\infty} n!^2 x^n = \sum_{n=0}^{\infty} n! x^n$ - definirana samo v 0, obe z vrednostjo 1 tam.

$$F(x) = \begin{cases} e^{-\frac{1}{x^2}} x \neq 0 \\ 0 \quad x = 0 \end{cases} : \mathbb{R} \rightarrow \mathbb{R}$$

$$F^{(n)}(0) = 0 \quad \forall n \geq 0 \implies F(x) = 0 + 0x + 0x^2 + \dots$$

Potenčne vrste niso „najboljše“ za študij zaporedij.

2.2 Formalne potenčne vrste

K komutativni obseg s karakteristiko $0 : 1 + 1 + \dots + 1 \neq 0 \ \forall n \geq 1$.

$\mathbb{Q}, \mathbb{R}, \mathbb{C}$

$\frac{1}{n!}$ je definirano

$K[[x]] = \{(a_n)_n : a_n \in K\} = K^{\mathbb{N}}$ - množica formalnih potenčnih vrst (FPV)

= zaporedje

$K[x] = \{(a_n)_n : a_n \in K, a_n = 0 \ \forall n \geq n_0\}$ - množica polinomov.

V $K[[x]]$ vpeljemo operacije:

$$(a_n)_n + (b_n)_n = (a_n + b_n)_n,$$

$$\lambda(a_n)_n = (\lambda a_n)_n,$$

$$((a_n)_n \cdot (b_n)_n) = (c_n)_n; \ c_n = \sum_{k=0}^n a_k b_{n-k} \text{ - konvolucijsko množenje.}$$

$K[[x]]$ algebra formalnih potenčnih vrst: komutativna, $(1, 0, 0, 0, \dots)$ enota za množenje: $\sum_{k=0}^n a_k \cdot \delta_{n-k,0} = a_n$.

Oznake:

$(a_n)_n \leftrightarrow \sum_n a_n x^n$: ni vsota (samo oznaka), x je ločilo (ni spremenljivka, ne „vstavljamo“),

$$(a_0 + a_1 x + \dots)(b_0 + b_1 x + \dots) = a_0 b_0 + (a_1 b_0 + a_0 b_1)x + \dots,$$

$$1 + 0x + 0x^2 + \dots = 1,$$

$$[x^n]F(x) := a_n \text{ - „koeficient pred } x^n\text{“,}$$

$$F(0) := [x^0]F(x).$$

Trditev 2.2.1. $F(x)$ ima inverz $\iff F(0) \neq 0$.

Dokaz 2.2.2.

(\implies) :

$$F(x)G(x) = 1$$

$$F(0)G(0) = 1 \implies F(0) \neq 0$$

$(\Longleftarrow) :$

$$\begin{aligned}
 F(x) &= a_0 + a_1x + a_2x^2 + \dots, a_0 \neq 0 \\
 G(x) &= b_0 + b_1x + b_2x^2 + \dots \\
 F(x)G(x) &= 1 \\
 a_0b_0 &= 1 \implies b_0 = \frac{1}{a_0} \\
 a_0b_1 + a_1b_0 &= 0 \implies b_1 = \frac{-a_1b_0}{a_0} \\
 a_0b_2 + a_1b_1 + a_2b_0 &= 0 \implies b_2 = \frac{-a_1b_1 - a_2b_0}{a_0} \\
 &\vdots
 \end{aligned}$$

Opomba. K komutativen kolobar s karakteristiko 0.

$F(x)$ ima inverz $\iff F(0)$ ima inverz v K .

$$v(F(x)) = \begin{cases} \min n : [x^n]F(x) \neq 0 & F(x) \neq 0 \\ \infty & F(x) = 0 \end{cases} \text{ - valuacija.}$$

$$v(F(x)G(x)) = v(F(x)) + v(G(x)) \quad (\implies \text{ni deliteljev nič})$$

$$v(F(x) + G(x)) \geq \min\{v(F(x)), v(G(x))\}$$

$$v(\lambda F(x)) = \begin{cases} v(F(x)) & \lambda \neq 0 \\ \infty & \lambda = 0 \end{cases}$$

$$d(F(x), G(x)) = 2^{-v(F(x)-G(x))} \text{ - metrika}$$

$$d(F(x), G(x)) = 2^{-k} \iff [x^n]F(x) = [x^n]G(x) \quad \forall n \leq k.$$

Trditev 2.2.3. $(K[[x]], d)$ je poln metrični prostor.

Dokaz 2.2.4.

$$\begin{aligned}
d &\geq 0, d = 0 \iff F = G \\
d(F(x), G(x)) &= d(G(x), F(x)) \\
d(F(x), H(x)) &= 2^{-v(F(x)-H(x))} \\
&= 2^{-v(F(x)-G(x)+G(x)-H(x))} \\
&\leq \max\{2^{-v(F(x)-G(x))}, 2^{-v(G(x)-H(x))}\} \\
&= \max\{d(F(x), G(x)), d(G(x), H(x))\} \\
&\leq d(F(x), G(x)) + d(G(x), H(x)).
\end{aligned}$$

$F_m(x) = \sum_n a_n^{(m)} x^n$ Cauchyjevo zaporedje

$$\forall k \exists M : M_1, M_2 \geq M \implies d(F_{M_1}(x), F_{M_2}(x)) < 2^{-k}$$

$$\text{oz. } [x^n]F_{M_1}(x) = [x^n]F_{M_2}(x) \quad \forall n \leq k.$$

Torej za vsak $[x^n]F_n(x)$ konstantni od nekod naprej in enaki npr. a_n .

$$F(x) = \sum_n a_n x^n \text{ je limita } (F_n(x))_m.$$

Primer.

$$(\sum_n x^n)(1-x) = 1$$

$$c_n = 1 \cdot (-1) + 1 \cdot 1 = 0 \quad \forall n \geq 1$$

$$c_0 = 1. \text{ Torej } \sum_n x^n = \frac{1}{1-x} \implies 1-x \text{ inverz od } \sum_n x^n.$$

$$\lim_{N \rightarrow \infty} \sum_{n=0}^N x^n = \frac{1}{1-x}.$$

Opomba. $(F_m(x))_m$ konvergira v $K[[x]]$, če je $([x^n]F_m(x))_m$ od nekod naprej konstantno, npr a_n ; v tem primeru je $\lim_{m \rightarrow \infty} F_m(x) = \sum_n a_n x^n$.

Odvajanje:

$$F'(x) = \lim_{h \rightarrow 0} \frac{F(x+h)-F(x)}{h}.$$

Za $K[[x]]$:

$$[x^n]F'(x) := (n+1)[x^{n+1}]F(x)$$

$$(\sum_n a_n x^n)' = F(x)'G(x) + F(x)G(x)'$$

Dokaz: DN.

$$\left(\frac{F(x)}{G(x)}\right)' = \frac{F(x)'G(x) - F(x)G(x)'}{G(x)^2}; \quad G(0) \neq 0$$

Primer.

$$F'(x) = F(x)$$

$$(n+1)a_{n+1} = a_n$$

$$na_n = a_{n-1}$$

a_0 poljubno

$$a_n = \frac{a_0}{n!}.$$

$$e^{\lambda x} := \sum_n \frac{\lambda^n}{n!} x^n$$

$$e^{\lambda x} \cdot e^{\mu x} = e^{(\lambda+\mu)x}$$

$$L = \sum_{k=0}^n \frac{\lambda^k}{k!} \frac{\mu^{n-k}}{(n-k)!} \stackrel{?}{=} \frac{(\lambda+\mu)^n}{n!} = D.$$

Binomski izrek v K : enakost velja.

$$F'(x) = \frac{1}{1-x}, \quad F(0) = 0$$

$$(n+1)a_{n+1} = 1$$

$$a_n = \frac{a_0}{n}$$

$$\log \frac{1}{1-x} := \sum_{n=1}^{\infty} \frac{x^n}{n!}$$

$$e^{\log \frac{1}{1-x}} \stackrel{?}{=} \frac{1}{1-x}.$$

Najprej definicija kompozituma, dokaz enakosti kasneje.

Bolj splošno:

$$F(0) = 1$$

$$\log(F(x)G(x)) = \log F(x) + \log G(x): \text{DN.}$$

Binomska vrsta:

$$\lambda \in K, n \in \mathbb{N}, \binom{\lambda}{n} := \frac{\lambda^n}{n!} \text{ posplošen binomski koeficient.}$$

$$B_\lambda(x) = \sum_{n=0}^{\infty} \binom{\lambda}{n} x^n$$

$$n \in \mathbb{N}: B_n(x) = \sum_{k=0}^{\infty} \binom{n}{k} x^k = (1+x)^n.$$

Trditev 2.2.5.

$$B_\lambda(x) \cdot B_\mu(x) = B_{\lambda+\mu}(x).$$

Dokaz 2.2.6.

$$D = \frac{(\lambda+\mu)^n}{n!} = \sum_{k=0}^n \frac{\lambda^k}{k!} \frac{\mu^{n-k}}{(n-k)!} = L$$

$$\sum_{k=0}^n \binom{n}{k} \lambda^k \mu^{n-k} = (\lambda + \mu)^n.$$

Indukcija: DN.

$$B_\lambda(x) := (1+x)^\lambda$$

$$n \in \mathbb{N} : B_n(x) \cdot B_{-n}(x) = 1$$

$$(1+x)^{-n} = \frac{1}{(1+x)^n}$$

$$(1+x)^{-n} = \sum_k \binom{-n}{k} x^k$$

$$\begin{aligned} \binom{-n}{k} &= \frac{(-n)(-n-1)\dots(-n-k+1)}{k!} \\ &= \frac{(-1)^k(n+k-1)\dots n}{k!} \cdot \frac{(n-1)!}{(n-1)!} \\ &= (-1)^k \binom{n+k-1}{k-1} \end{aligned}$$

$$\begin{aligned} (1-x)^{-k} &= \frac{1}{1-x} \dots \frac{1}{1-x} \\ &= \sum_{n=0}^{\infty} \left(\sum_{n_i \geq 0, \sum n_i = k} 1 \right) x^n \\ &= \sum_n (\text{število šibkih kompozicij } n \text{ s } k \text{ členi}) x^n \\ &= \sum_n \binom{n+k-1}{k-1} x^n \end{aligned}$$

$$F(x)G(x)H(x) = \sum_{n=0}^{\infty} \left(\sum_{n_1, n_2, n_3 \geq 0, n_1+n_2+n_3=n} a_{n_1} b_{n_2} c_{n_3} \right) x^n$$

$$\binom{-1}{n} = (-1)^n \binom{n}{0} = (-1)^n$$

$$(1+x)^{\frac{1}{2}} = 1 + \sum_{n=1}^{\infty} \frac{(-1)^{n-1}}{2^{2n-1}} \binom{2n-2}{n-1} x^n$$

$$\begin{aligned} \binom{\frac{1}{2}}{n} &= \frac{\frac{1}{2} \left(-\frac{1}{2} \right) \cdot \left(\frac{1}{2} - n + 1 \right)}{n!} \\ &= \frac{(-1)^{n-1} (2n-3)!!}{2^n \cdot n!} \cdot \frac{(2n-2)!!}{(2n-2)!!} \\ &= \frac{(-1)^{n-1} (2n-2)!}{2^n \cdot n! \cdot 2^{n-1} \cdot (n-1)!} \\ &= \frac{(-1)^{n-1}}{2^{2n-1} n} \binom{2n-2}{n-1} \quad n \geq 1. \end{aligned}$$

2.3 Kompozitum

$$F(x) = \sum_n a_n x^n$$

$$G(x) = \sum_n b_n x^n$$

$$F \circ G(x) = F(G(x)) = ?$$

$$(F \circ G)(x) = a_0 + a_1 G(x) + a_2 G^2(x) + \dots = \lim_{N \rightarrow \infty} \sum_{n=0}^N a_n G^n(x).$$

Kdaj ta limita obstaja?

Trditev 2.3.1. $(F_n(x))_n$.

$$\lim_{N \rightarrow \infty} F_n(x) \text{ obstaja} \iff \lim_{n \rightarrow \infty} v(F_n(x)) = \infty.$$

Dokaz 2.3.2.

(\implies) :

$$\left(\sum_{n=0}^N F_n(x) \right)_N \text{ je Cauchyjevo :}$$

$$\forall x \exists N_0 \forall N, M \geq N_0 : d \left(\sum_{n=0}^N F_n(x), \sum_{m=0}^M F_m(x) \right) \leq 2^{-k}$$

$$M = N - 1 : v(F_N(x)) \geq k.$$

(\impliedby) :

$$\forall k \exists N_0 \forall N \geq N_0 : v(F_N(x)) \geq k \text{ (predpostavka)}$$

$$\begin{aligned} N > M \geq N_0 : d \left(\sum_{n=0}^N F_n(x), \sum_{m=0}^M F_m(x) \right) \\ &= 2^{-v(F_{M+1}(x) + \dots + F_N(x))} \\ &\leq \max \{ 2^{-v(F_{M+1}(x))} \dots 2^{-v(F_N(x))} \} \\ &\leq 2^{-k}. \end{aligned}$$

$$F \circ G(x) \text{ obstaja} \iff \lim_{n \rightarrow \infty} v(a_n G^n(x)) = \infty$$

$$\iff v(G(x)) > 0 \text{ ali } a_n = 0 \text{ od nekod naprej}$$

$$\iff F \text{ polinom ali } G(0) = 0.$$

$$\text{Velja } v(a_n G^n(x)) = \begin{cases} n \cdot v(G(x)) & a_n \neq 0 \\ \infty & a_n = 0 \end{cases}$$

Primer.

$$F(x) = x^2 - 3x + 1$$

$$G(x) = e^x$$

$$(F \circ G)(x) = e^{2x} - 3e^x + 1 - \text{ok}$$

$$F(x) = G(x) = e^x - \text{ni ok}$$

$$F(x) = e^x$$

$$G(x) = e^x - 1$$

$$e^{e^x - 1} - \text{ok.}$$

Opomba.

$$F(x) = \sum_n a_n x^n$$

$$G(x) = \sum_n b_n x^n \quad b_0 = 0$$

$$a_0 + a_1(b_1x + b_2x^2 + \dots) + a_2(b_1x + b_2x^2 + \dots)^2 + \dots$$

Za izračun koeficienta pri x^5 izračunamo končno vsoto.

$$\text{Enota za kompozitum: } x = 0 + 1 \cdot x + 0 \cdot x^2 + \dots$$

$$F \circ x = a_0 + a_1x + a_2x^2 + \dots = F = x \circ F = 1 \cdot (a_0 + a_1x + \dots)$$

Izrek 2.3.3.

$F \in K[[x]]$ ima inverz za kompozitum $\iff F(x) = a_0 + a_1x$; $a_1 \neq 0$ ali $v(F(x)) = 1$.

Primer.

$$x - x^2 \text{ ima inverz,}$$

$$e^x - 1 \text{ ima inverz,}$$

$$x^2 \text{ nima inverza.}$$

$$F^{<-1>} - \text{inverz za kompozitum.}$$

Dokaz 2.3.4.

$(\implies):$

$$F(x) = \sum_n a_n x^n$$

$$G(x) = \sum_n b_n x^n \text{ inverz od } F$$

$$a_0 = 0 \stackrel{?}{\iff} b_0 = 0$$

$$(\Longleftarrow) : F \circ G = a_0 + a_1(b_1x + \dots) + a_2(\dots)^2 + \dots$$

$$[x^0]F(G(x)) = a_0 = [x^0]x = 0$$

$(\implies) : \text{isto?}$

$$1. a_0 \neq 0, b_0 \neq 0$$

$$\implies F, G \text{ polinoma, } \deg(F \circ G) = \deg(F) \cdot \deg(G) = 1$$

$$\implies \deg(F) = \deg(G) = 1$$

$$2. a_0 = b_0 = 0$$

$$v(F \circ G) = v(F) \cdot v(G) = 1$$

$$\implies v(F) = v(G) = 1$$

$$\implies F(x) = a_1x + a_2x^2 + \dots \quad a_1 \neq 0.$$

(\Leftarrow):

$$F(x) = a_0 + a_1x \quad a_1 \neq 0$$

$$a_0 + a_1y = x \implies y = \dots$$

$$F^{<-1>}(x) = -\frac{a_0}{a_1} + \frac{x}{a_1}$$

$$F(x) = a_1x + a_2x^2 + \dots \quad a_1 \neq 0$$

$$\text{levi inverz: } G_1(x) = b_0 + b_1x + \dots$$

$$G_1 \circ F = x$$

$$b_0 + b_1(a_1x + \dots) + b_2(a_1x + \dots)^2 + \dots = x$$

$$[x^0] : b_0 = 0$$

$$[x^1] : a_1b_1 = 0 \implies b_1 = \frac{1}{a_1}$$

$$[x^2] : b_1a_2 + b_1a_1^2 = 0 \implies b_2 = -\frac{b_1a_2}{a_1^2}$$

$$[x^3] : b_1a_3 + 2b_2a_1a_2 + b_3a_1^3 = 0 \implies b_3 = \dots \frac{\ddots}{a_1^3}$$

$$[x^n] : \dots + b_na_1^n = 0 \quad n \geq 1$$

$$b_n = \dots \text{ rekurzivno}$$

$$\text{desni inverz: } G_2(x) = c_0 + c_1x + \dots, \quad c_0 = 0$$

$$F \circ G_2 = x$$

$$a_1(c_1x + \dots) + a_2(c_1x + \dots)^2 + \dots = x$$

$$[x^0] : 0 = 0$$

$$[x^1] : a_1c_1 = 1 \implies c_1 = \frac{1}{a_1}$$

$$[x^2] : a_1c_2 + a_2c_1^2 = 0 \implies c_2 = -\frac{a_2c_1^2}{a_1}$$

$$[x^3] : a_1c_3 + 2a_2c_1c_2 + a_3c_1^3 = 0 \implies c_3 = \frac{\ddots}{a_1}$$

$$[x^n] : a_1c_n + \dots = 0 \implies c_n = \frac{\ddots}{a_1}.$$



$$(G_1 \circ F) \circ G_2 = G_2$$

$$G_1 \circ (F \circ G_2) = G_1.$$

Iz asociativnosti (ki je nismo dokazali) sledi $G_1 = G_2 = F^{<-1>}$.

Trditev 2.3.5.

$$F_n(0) = 0$$

$$\lim_{N \rightarrow \infty} \prod_{n=1}^N (1 + F_n(x)) \text{ obstaja} \iff \lim_{n \rightarrow \infty} v(F_n(x)) = \infty.$$

Dokaz DN.

Primer.

$$(1+x)(1+x)(1+x)\dots - \text{ni ok,}$$

$$(1+x)(1+x^2)(1+x^3)\dots - \text{ok.}$$

Opomba.

$$K[[x]]$$

$$K[[x,y]] = K^{\mathbb{N} \times \mathbb{N}}$$

$\sum a_{n,m} x^n y^m$ bivariantna potenčna vrsta.

$$\sum_{k,m} \binom{n}{k} x^k y^m = \sum_m (1+x)^n y^m = \frac{1}{1-(1+x)y}.$$

$$K[[x_1, x_2 \dots]]$$

$$x_1 x_2^2 x_3 + x_2 x_3 + \dots - \text{ok}$$

$$x_1 x_2 x_3 x_4 \dots - \text{ni ok.}$$

2.4 Reševanje linearnih rekurzivnih enačb s konstantnimi koeficienti

Primer.

$$(1) \quad a_n = 2a_{n-1} + 1 \quad n \geq 1, a_0 = 1$$

$$1, 3, 7, 15 \dots$$

$F(x) = \sum_n a_n x^n$ rodovna funkcija (angl. generating function) zapo-

redja.

$$\begin{aligned} F(x) - 1 &= \sum_{n=1}^{\infty} a_n x^n = \sum_{n=1}^{\infty} (2a_{n-1} + 1)x^n = 2xF(x) + \frac{x}{1-x} \\ F(x)(1-2x) &= 1 + \frac{x}{1-x} = \frac{1}{1-x} \\ F(x) &= \frac{1}{(1-x)(1-2x)}. \end{aligned}$$

Ekvivalentno:

$$\begin{aligned} a_n &= 2a_{n-1} + 1 \quad / \cdot x^n \sum_{n=1}^{\infty} \\ F(x) - 1 &= \frac{x}{1-x} + 2xF(x) \\ F(x) &= \frac{1}{(1-x)(1-2x)} = \frac{A}{1-x} + \frac{B}{1-2x} = \frac{A(1-2x) + B(1-x)}{(1-x)(1-2x)} \\ / \cdot (1-x), x=1 \\ \frac{1}{-1} &= A \implies A = -1 \\ / \cdot (1-2x), x=\frac{1}{2} \\ B &= 2 \end{aligned}$$

$$a_n = -1 + 2^{n+1}.$$

$$(2) \quad F_n = F_{n-1} + F_{n-2} \quad n \geq 2, F_0 = F_1 = 1 \quad / \cdot x^n \sum_{n=2}^{\infty}$$

$$\begin{aligned} F(x) &= \sum_n F_n x^n \\ F(x) - 1 - x &= x(F(x) - 1) + x^2 F(x) \\ F(x) &= \frac{1}{1-x-x^2} = \frac{1}{(1-y_1x)(1-y_2x)}. \end{aligned}$$

Ničli $1-x-x^2$ sta $\frac{1}{y_1}, \frac{1}{y_2}$

$$y_1, y_2 \text{ sta ničli } y^2 - y - 1 \text{ (obrnjen polinom), torej } x_1, x_2 = \frac{-1 \pm \sqrt{5}}{2}.$$

V splošnem:

$$p(x) = c_0 + c_1x + \dots + c_dx^d; \quad c_d \neq 0$$

ima ničle $\lambda_1 \dots \lambda_d$, ima

$p^{\text{obr}}(x) = c_0 x^d + c_1 x^{d-1} + \dots + c_d$ (obrtni polinom) ničle $\frac{1}{\lambda_1} \dots \frac{1}{\lambda_d}$:

$$\begin{aligned} p^{\text{obr}}\left(\frac{1}{\lambda_i}\right) &= c_0 \cdot \frac{1}{\lambda_i^d} + c_1 \cdot \frac{1}{\lambda_i^{d-1}} + \dots + c_d \\ &= \frac{c_0 + c_1 \lambda_i + \dots + c_d \lambda_i^d}{\lambda_i^d} = 0 \end{aligned}$$

$$\begin{aligned} F(x) &= \frac{1}{1-x-x^2} \\ &= \frac{1}{(1-y_1x)(1-y_2x)} \\ &= \frac{\frac{1}{1-\frac{y_2}{y_1}}}{1-y_1x} + \frac{\frac{1}{1-\frac{y_1}{y_2}}}{1-y_2x} \\ &= \frac{1}{y_1-y_2} \left(\frac{y_1}{1-y_1x} - \frac{y_2}{1-y_2x} \right) \\ y_1 - y_2 &= 5 \\ \Rightarrow F_n &= \frac{1}{\sqrt{5}} \left(\left(\frac{1+\sqrt{5}}{2} \right)^{n-1} - \left(\frac{1-\sqrt{5}}{2} \right)^{n+1} \right). \end{aligned}$$

Izrek 2.4.1. NSTE (naslednje trditve so ekvivalentne) za $(a_n)_n, a_n \in \mathbb{C}$:

- (1) $c_d a_n + c_{d-1} a_{n-1} + \dots + c_n a_{n-d} = 0, \quad n \geq d, \quad c_0, c_d \neq 0,$
- (2) $F(x) = \sum_n a_n x^n = \frac{P(x)}{c_d + \dots + c_0 x^d}, \quad \deg P < d,$
- (3) $a_n = \sum_{i=1}^k p_i(n) \lambda_i^n, \quad \lambda_1 \dots \lambda_k$ ničle $c_d y^d + \dots + c_0$ (karakteristični polinom) s kratnostmi $\alpha_1 \dots \alpha_k, \quad \deg p_i < \alpha_i.$

Dokaz 2.4.2.

(1) \implies (2):

$$\begin{aligned}
 c_d a_n + c_{d-1} a_{n-1} + \cdots + c_n a_{n-d} &= 0 \quad / \cdot x^n \sum_{n=d}^{\infty} \\
 c_d (F(x) - a_0 - \cdots - a_{d-1} x^{d-1}) \\
 + c_{d-1} (F(x) - a_0 - \cdots - a_{d-2} x^{d-2}) \\
 + \cdots + c_0 x^d F(x) &= 0 \\
 F(x) = (c_d + c_{d-1} x + c_{d-2} x^2 + \cdots + c_0 x^d) &= P(x) \quad \deg P < d.
 \end{aligned}$$

(2) \implies (1):

$$\begin{aligned}
 (c_d + c_{d-1} x + \cdots + c_0 x^d) \cdot \sum_n a_n x^n &= P(x) \\
 n \geq d: [x^n]: c_d a_n + \cdots + c_0 a_{n-d} &= 0.
 \end{aligned}$$

(2) \implies (3):

$$\begin{aligned}
 \sum_n a_n x^n &= \frac{P(x)}{c_d (1 - \lambda_1 x)^{\alpha_1} \cdots (1 - \lambda_m x)^{\lambda_m}} \\
 &\stackrel{\text{parc}}{=} \sum_{i=1}^k \sum_{j=1}^{\alpha_i} \frac{A_{ij}}{(1 - \lambda_i x)^j} \\
 \frac{1}{(1-x)^d} &= \sum_n \binom{n+d-1}{d-1} x^n \\
 a_n &= \sum_{i=1}^k \left(\sum_{j=1}^{\alpha_i} A_{ij} \cdot \binom{n+j-1}{j-1} \right) \lambda_i^n, \\
 \binom{n+j-1}{j-1} &\text{binom v } n \text{ stopnje } j-1 < \alpha_i.
 \end{aligned}$$

(3) \implies (2): podobno: $p_i(n)$ zapišemo v bazi $\binom{n+j-1}{j-1}$.

Primer.

$$a_n - 7a_{n-1} + 18a_{n-2} - 12a_{n-3} = 0, \quad a_0, a_1, a_2 \text{ dani.}$$

$$y^3 - 7y^2 + 18y - 12 = (y-2)^2(y-3)$$

$$\implies a_n = 2^n(An + B) + 3^n \cdot C.$$

A, B, C dobimo iz a_0, a_1, a_2 (vstavimo, dobimo sistem).

Opomba.

$\sum_n a_n x^n = \frac{P(x)}{Q(x)}$, $\deg P \geq \deg Q \iff c_d a_n + \dots + c_n a_{n-d} = 0$ za $n \geq N$ (dovolj velik).

Opomba.

$c_d a_n + \dots + c_0 a_{n-d} = r(n) \cdot \lambda^n$, $\deg r = \alpha$.

Homogena + partikularna

$\sum_n r(n) \lambda^n x^n = \frac{R(x)}{(1-\lambda x)^\alpha}$.

Če λ^{α_i} –kratna ničla karakterističnega polinoma: $\sum_{j=1}^{\alpha+\alpha_i} \dots$

Nastavek: $n^{\alpha_i} q(n) \lambda^n$, $\deg q = \alpha_i - 1$.

Primer.

$a_n - 4a_{n-1} + 4a_{n-2} = n \cdot 2^n$, $n \geq 2$.

Partikularna: $n^2 \cdot (An + B)2^n$.

2.5 Nadaljevanje uporabe običajnih rodovnih funkcij

$F(x) = \sum_n a_n x^n$

$F(x) \xleftrightarrow{\text{orf}} (a_n)_n$

$F'(x) \xleftrightarrow{\text{orf}} ((n+1)a_{n+1})_n$

$xF'(x) \xleftrightarrow{\text{orf}} (na_n)_n$

$DF(x) := F'(x)$, D : operator odvajanja.

$(xD)^2 F(x) \xleftrightarrow{\text{orf}} (n^2 a_n)_n$

$p(xD)F(x) \xleftrightarrow{\text{orf}} (p(n)a_n)_n$, p polinom.

Primer.

$\sum_j j^2$

$\frac{1}{1-x} \xleftrightarrow{\text{orf}} (1)_n$

$(xD)^2 \frac{1}{1-x} \xleftrightarrow{\text{orf}} \left(\sum_{j=0}^n a_j \right)_n$

$x \cdot \left(\frac{x}{(1-x)^2} \right)' = \dots = \frac{x(1+x)}{(1-x)^3}$ - samo členi. $F(x) \xleftrightarrow{\text{orf}} (a_n)_n$

$F(x) \cdot \frac{1}{1-x} \xleftrightarrow{\text{orf}} \left(\sum_{j=0}^n a_j \right)_n$ - konvolucija z $(1)_n$.

$$\begin{aligned}
[X^n] \left(F(x) \cdot \frac{1}{1-x} \right) &= [x^n] \left(\frac{x^2}{(1-x)^4} + \frac{x^2}{(1-x)^4} \right) \\
&= \binom{n+2}{3} + \binom{n+1}{3} \\
&= \frac{n(n+1)(2n+1)}{6}.
\end{aligned}$$

$$F(x) \cdot G(x) = \sum_n a_n x^n \cdot \sum_n b_n x^n = \sum_n \left(\sum_{k=0}^n a_k b_{n-k} \right) x^n.$$

Naj bo 1. del struktura A $((a_n)_n)$ preštevalno zaporedje,

naj bo 2. del struktura B $((b_n)_n)$ preštevalno zaporedje:

$$\sum_{k=0}^n a_k b_{n-k}.$$

Primer.

- (1) m kroglic, rdeče, črne, zelene, zelenih kroglic sodo in so na koncu.

1, 2, 5, 10 ...

A : rdeče / črne kroglice: $2^n \rightarrow \frac{1}{1-2x}$

B : sodo mnogo zelenih kroglic: $1, 0, 1, 0, 1 \dots \rightarrow \frac{1}{1-x^2}$

$$\frac{1}{1-2x} \cdot \frac{1}{1-x^2} = \frac{\frac{4}{3}}{1-2x} + \frac{-\frac{1}{2}}{1-x} + \frac{\frac{1}{6}}{1+x}$$

$$a_n = \frac{4}{3} \cdot 2^n - \frac{1}{2} + \frac{1}{6}(-1)^n.$$

- (2) Kompozicije s k členi

A : neničelno število: $0, 1, 1, 1, 1 \dots \rightarrow \frac{x}{1-x}$

$$\left(\frac{x}{1-x} \right)^k = \sum_n \binom{n+k-1}{k-1} x^{n+k} = \sum_n \binom{n-1}{k-1} x^n,$$

šibke kompozicije:

$$\left(\frac{1}{1-x} \right)^k,$$

kompozicije z lihimi členi: $0, 1, 0, 1, 0, 1 \dots \rightarrow \frac{x}{1-x^2}$

$$\left(\frac{x}{1-x^2} \right)^k.$$

- (3) $S(n, k)$

$$n = 7, k = 3 : \{ \{1, 4, 5\}, \{2, 7\}, \{3, 6\} \}$$

$$\sum_n S(n, k) x^n = ?$$

Vrstni red določimo: 1 v 1. bloku, v 2. bloku najmanjše število, ki ni v

1. bloku ...

$\rightarrow 1\ 2\ 3\ 1\ 1\ 3\ 2$ (primer od prej).

Dobimo: zaporedje n števil v $[k]$, vsa od 1 do k se pojavijo, 1. pojavitev i je pred 1. pojavitvijo $i + 1$

$1\ (1 \dots 1)2(1/2 \dots 1/2)3(\dots) \dots$

$$x \cdot \frac{1}{1-x} \cdot x \cdot \frac{1}{1-2x} \dots$$

$$\sum_n S(n, k) x^n = \frac{x^k}{(1-x)(1-2x) \dots (1-kx)}.$$

Ekvivalentno: $(1 - kx) \sum_n S(n, k) x^n = \sum_n S(n - 1, k - 1) x^n$

$$[x^n] : S(n, k) - kS(n - 1, k) = S(n - 1, k - 1)$$

$$\frac{x^k}{(1-x) \dots (1-kx)} = \frac{(-1)^k}{k!} + \sum_{j=1}^k \frac{A_j}{1-jx} \stackrel{DN}{=} \dots$$

(4) Razčlenitve

$\overline{p}_k(n)$ $\stackrel{\text{konjugiranje}}{=}$ število razčlenitev n s členi $\leq k$

$$\frac{1}{1-x} \cdot \frac{1}{1-x^2} \cdots \frac{1}{1-x^k}$$

$$= \sum_n \overline{p}_k(n) x^n$$

$$= (1 + x + x^2 + \dots)(1 + x^2 + x^4 + \dots)(1 + x^3 + \dots) \dots (1 + x^k + \dots)$$

$$[x^n] : x^n = x^{m_1} \cdot x^{2m_2} \dots x^{km_k}$$

$$n = m_1 + 2m_2 + \dots + km_k$$

$$k \dots k \dots 32 \dots 21 \dots 1$$

$$\sum_n p_n(n) x^n = \lim_{k \rightarrow \infty} \sum_n \overline{p}_k(n)$$

$$= \lim_{n \rightarrow \infty} \frac{1}{\prod_{j=1}^n (1 - x^j)}$$

$$= \prod_{i=1}^{\infty} \frac{1}{1 - x^i}.$$

$d(n)$: število razčlenitev n z različnimi členi

$$\sum_n d(n) x^n = \prod_{i=1}^{\infty} (1 - x^i) \quad (0 \text{ ali } 1\text{-krat vedno})$$

$o(n)$ = število razčlenitev n z lihimi členi

$$\sum_n o(n) x^n = \prod_{i=0}^{\infty} \frac{1}{1 - x^{2i+1}}$$

$$\prod_i (1+x^i) \cdot \frac{1-x^i}{1+x^i} = \prod_i \frac{1-1^{2i}}{1-x} = \prod_i \frac{1}{1-x^{2i+1}}$$

$$\implies o(n) = d(n).$$

DN: bijekcija.

(5) c_n : Dyckove poti dolžine n

$$c_{n+1} = \prod_{k=0}^n c_k \cdot c_{n-k} \quad / \cdot x^{n+1} \sum_n$$

$$F(x) - 1 = x \cdot \sum_{n=0}^{\infty} (\sum_{k=0}^n c_k c_{n-k}) x^n = x \cdot F^2(x)$$

$$F(x) = 1 + xF^2(x):$$

1: prazna, $xF^2(x)$: dolžine n , $2n$ korakov

Motzkinova pot: v smeri $(1,1), (1, -1), (1,0)$

$$M(x) = 1 + xM(x) + x^2M^2(x):$$

1: prazna, $xM(x)$: naravnost, $x^2M^2(x)$: desno-gor

$$xF^2 - F + 1 = 0$$

$$F = \frac{-1 \pm \sqrt{1-4x}}{2x}$$

$$\sqrt{1-4x} = 1 - \sum_{n=1}^{\infty} \frac{1}{n} \binom{2n-2}{n-1} \cdot \frac{(-1)^n}{2^{2n-1}} (-4x)^n = 1 - \sum_{n=1}^{\infty} \frac{2}{n} \binom{2n-2}{n-1} x^n$$

$$\frac{1+\sqrt{1-4x}}{2x} - \text{ne, ker } \frac{2+\dots}{2x}$$

$$\frac{1-\sqrt{1-4x}}{2x} = \sum_{n=0}^{\infty} \frac{1}{n+1} \binom{2n}{n} x^n.$$

Druga utemeljitev:

$$4x^2F^2 - 4xF + 4x = 0$$

$$(2xF - (1 - \sqrt{1-4x}))(2xF - (1 + \sqrt{1-4x})) = 0 \text{ v } K[[x]].$$

$$2xF - (1 + \sqrt{1-4x}) \neq 0 \text{ (konstantni koeficient nima 0)}$$

$$\implies 2xF = 1 - \sqrt{1-4x}.$$

$F^k(x)$: razdelimo na k delov, vsakemu damo strukturo F .

$\sum_{k=0}^{\infty} F^k(x) = \frac{1}{1-F(x)}$: razdelimo na poljubno mnogo delov, vsakemu F .

Primer.

(1) Kompozicije n .

$$\frac{1}{1-\frac{x}{1-x}} = \frac{1-x}{1-2x} = \begin{cases} 2^{n-1} & n > 0 \\ 0 & n = 0 \end{cases}$$

kompozicije s členi 1 in 2

$$\frac{1}{1-(x+x^2)}.$$

(2) $2 \times n$ plošča, domine 2×1 .

Primitivni tlakovanji

$$\frac{1}{1-x-x^2}$$

Domini 1×1 in 2×1

$n = 1$: 1 možnost,

$n = 2$: 3,

$n = 3$: 2,

$n = 4$: 2,

\vdots

$$\frac{1}{1-(2x+3x^2+2x^3+\dots)} = \frac{1}{1-x^2-\frac{2x}{1-x}} = \frac{1-x}{1-3x-x^2+x^3}.$$

(3) Primitivna Dyckova pot: se ne dotakne x osi.

$$F(x) = \frac{1}{1-xF(x)},$$

$$M(x) = \frac{1}{1-x-x^2F(x)}.$$

Levi faktor Dyckove poti: $L(x) = \frac{F(x^2)}{1-x-x^2F(x)} = \dots = \frac{2}{1-2x+\sqrt{1-4x^2}}$

$F(x^2)$: Dyckova pot (na začetku), $xF(x^2)$: korak + Dyckova pot.

DN: $L_n = \binom{n}{\lfloor \frac{n}{2} \rfloor}$, namig: $\frac{1}{\sqrt{1-4x}} = ?$

$(F \circ G)(x) = a_0 + a_1G(x) + a_2G^2(x) + \dots$: razdelimo na poljubno delov, vsakemu delu damo strukturo G , delom da strukturo F .

Primer.

Število kompozicij s sodo mnogo lihimi členi.

$n = 0 : 1$

$n = 1 : 0$

$n = 2 : 1$

$n = 3 : 0$

$n = 4 : 3$

$n = 5 : 0$

$n = 6 : 8$

$n = 7 : 0$

$n = 8 : 21$

$G(x) = \frac{x}{1-x^2}$ - lihi

$F(x) = \frac{1}{1-x^2}$ - sodo mnogo.

$$\begin{aligned}
 (F \circ G)(x) &= \frac{1}{1 - \left(\frac{x}{1-x^2}\right)^2} \\
 &= \frac{(1-x^2)^2}{(1-x-x^2)(1+x-x^2)} \\
 &= \dots \\
 &= 1 + \frac{x}{2} \left(\frac{1}{1-x-x^2} - \frac{1}{1+x-x^2} \right) \\
 &= \sum_{n \text{ lih}} F_n x^n
 \end{aligned}$$

kjer se, ko razpišemo $\left(\frac{1}{1-x-x^2} - \frac{1}{1+x-x^2}\right)$ sodi odštejejo, lihi štejejo 2-krat, to delimo z 2.

Primer (Dobri Will Hunting).

(1) Matrika sosednosti: $A = \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 2 & 1 \\ 0 & 2 & 0 & 0 \\ 1 & 1 & 0 & 0 \end{bmatrix}.$

(2) Matrika, ki opisuje sprehode dolžine 3: $A^3 = \begin{bmatrix} 2 & 7 & 2 & 3 \\ 7 & 2 & 12 & 7 \\ 2 & 12 & 0 & 2 \\ 3 & 7 & 2 & 2 \end{bmatrix}.$

(3) Poišči rodovno funkcijo za sprehode $i \rightarrow j$

$$\sum_{k=0}^{\infty} A^k x^k = (I - Ax)^{-1} = \frac{1}{\det(I - Ax)} [\dots]$$

(4) $1 \rightarrow 3$:

$$\frac{2x^2 + 2x^3}{1 - 7x^2 - 2x^3 + 4x^4}.$$

2.6 Uporaba eksponentnih rodovnih funkcij

$$F(x) = \sum_n \frac{a_n}{n!} x^n$$

$$F(x) \xleftrightarrow{\text{erf}} (a_n)_n$$

$$\left[\frac{x^n}{n!} \right] F(x) = a_n$$

$$\left[\frac{x^n}{n!} \right] F(x) = n! [x^n] F(x)$$

$$F'(x) \xleftrightarrow{\text{erf}} (a_{n+1})_n$$

$$xF'(x) \xleftrightarrow{\text{erf}} (n \cdot a_n)_n$$

$$p(xD)F(x) \xleftrightarrow{\text{erf}} (p(n)a_n)_n.$$

Primer.

$$(1) \quad F_{n+2} = F_{n+1} + F_n; \quad n \geq 0$$

$$F(x) = \sum_n \frac{F_n}{n!} x^n$$

$$F''(x) - F'(x) - F(x) = 0$$

$$\lambda^2 - \lambda - 1 = 0 \implies \lambda_{1,2} = \frac{1 \pm \sqrt{5}}{2}$$

$$F(x) = Ae^{\frac{1+\sqrt{5}}{2}x} + Be^{\frac{1-\sqrt{5}}{2}x}$$

$$F_n = \left[\frac{x^n}{n!} \right] F(x) = A \left(\frac{1+\sqrt{5}}{2} \right)^n + B \left(\frac{1-\sqrt{5}}{2} \right)^n.$$

$$(2) \quad i_n: \text{število involucij v } S_n \text{ } (\pi^2 = id).$$

$$i_n = i_{n-1} + (n-1)i_{n-2}; \quad n \geq 2:$$

$$i_{n-1}: n \text{ fiksna točka}$$

$$i_{n-2}: n \text{ v transpoziciji z enim od } n-1 \text{ ostalih.}$$

$$I(x) = \sum_n \frac{i_n}{n!} x^n$$

$$I'' - I' - (xI' + I) = 0$$

$$I'' - (x+1)I' - I = 0$$

$$(I' - (x+1)I')' = 0$$

$$I' - (x+1)I = c$$

$$x=0: 1-1=0=c$$

$$I' = (x+1)I$$

$$\int \frac{dI}{I} = \int (x+1)dx$$

$$\ln I = \frac{x^2}{2} + x + \log D$$

$$I = D e^{x + \frac{x^2}{2}} \xrightarrow{x=0} D = 1$$

$$I(x) = e^{x + \frac{x^2}{2}}.$$

$$F(x) = \sum_n \frac{a_n}{n!} x^n$$

$$G(x) = \sum_n \frac{b_n}{n!} x^n$$

$$F(x)G(x) = \sum_n \left(\sum_{k=0}^n \frac{a_k}{k!} \frac{b_{n-k}}{(n-k)!} \right) x^n = \sum_n \left(\sum_{k=0}^n \binom{n}{k} a_k b_{n-k} \right) \frac{x^n}{n!}$$

$$\sum_{k=0}^n \binom{n}{k} a_k b_{n-k}: \text{ binomska konvolucija.}$$

orf: neoznačene strukture,

erf: označene strukture.

Primer.

d_n : premestitve v S_n (dearangement) - permutacije brez negibne točke.

$$D(x) = \sum_n \frac{d_n}{n!} x^n.$$

Permutacija = premestitev + množica negibnih točk.

$$(1\ 5\ 2)\ (3)\ (4\ 8\ 7)\ (6)$$

$$\frac{1}{1-x} = D(x) \cdot e^x$$

$$D(x) = \frac{e^{-x}}{1-x}$$

$$e^{-x} = \sum_n \frac{(-1)^n}{n!} x^n$$

$$\frac{e^{-x}}{1-x} = \sum_n \left(\sum_{k=0}^n \frac{(-1)^k}{k!} \right) x^n$$

$$d_n = n! \sum_{k=0}^n \frac{(-1)^k}{k!}.$$

$$\begin{aligned} F(x)G(x) &= \sum_n \left(\sum_{k=0}^n \binom{n}{k} a_k b_{n-k} \right) \frac{x^n}{n!} \\ &= \sum_n \left(\sum_{(S_1, S_2), S_1 \cap S_2 = \emptyset, S_1 \cup S_2 = [n]} a_{|S_1|} b_{|S_2|} \right) \frac{x^n}{n!} \end{aligned}$$

$$F^k(x) = \sum_n \left(\sum_{(i_1 \dots i_k), i_j \geq 0, i_1 + \dots + i_k = n} \binom{n}{i_1 \dots i_k} a_{i_1} \dots a_{i_k} \right) \frac{x^n}{n!}.$$

Predpostavimo $F(0) = 0!!$

$$\begin{aligned} F^k(x) &= \sum_n \left(\sum_{(S_1 \dots S_k), S_i \neq \emptyset, S_i \cap S_j = \emptyset, S_1 \cup \dots \cup S_k = n} a_{|S_1|} \dots a_{|S_k|} \right) \frac{x^n}{n!} \\ &= k! \sum_n \left(\sum_{(B_1 \dots B_k) \text{ razdelitev } [n]} a_{|B_1|} \dots a_{|B_k|} \right) \frac{x^n}{n!}. \end{aligned}$$

Izrek 2.6.1.

$F(0) = 0$.

$\frac{1}{k!} F^k(x)$ je erf za strukturo: izberemo razdelitev in vsakemu bloku damo strukturo F .

Primer.

$$\sum_n S(n, k) \frac{x^n}{n!} = \frac{1}{k!} (e^k - 1)^k$$

F : neprazna množica: $0, 1, 1 \dots \xrightarrow{\text{erf}} e^x - 1$.

Binomski izrek $(e^x - 1)^k = e^{-kx} - \dots$ nam da formulo za $S(n, k)$.

$$\sum_n c(n, k) \frac{x^n}{n!} = \frac{1}{k!} \left(\log \frac{1}{1-x} \right)^k$$

F : cikel: $a_n = (n-1)!$ za $n \geq 1 \xrightarrow{\text{erf}} \log \frac{1}{1-x}$

$$\sum_n L(n, k) \frac{x^n}{n!} = \frac{1}{k!} \left(\frac{x}{1-x} \right)^k$$

F : neprazna linearno urejena množica: $a_n = (n)!$ za $n \geq 1 \xrightarrow{\text{erf}} \log \frac{1}{1-x}$.

Izrek 2.6.2 (Eksponentna formula).

$F(0) = 0$.

$e^{F(x)}$ je erf za strukturo: izberemo razdelitev, vsakemu (bloku) damo strukturo F .

Dokaz 2.6.3. $\sum_{k=0}^{\infty} \frac{1}{k!} F^k(x) = e^{F(x)}$.

Primer.

(1) Permutacija = množica disjunktnih ciklov.

$$\frac{1}{1-x} = e^{\log \frac{1}{1-x}}.$$

DN: direktno.

(2) Involucija = množica ciklov dolžine 1 in 2: $(0,1,1,0,0,\dots)$

$$\sum_n \frac{i_n}{n!} = e^{x + \frac{x^2}{2}}$$

$$a_n = |\{\pi \in S_n : \pi^6 = id\}|$$

$$\sum_n \frac{a_n}{n!} x^n = e^{x + \frac{x^2}{2} + \frac{x^3}{3} + \frac{x^6}{6}}$$

$$\sum_n \frac{d_n}{n!} x^n = e^{\sum_{n \geq 2} \frac{x^n}{n}} = e^{\log \frac{1}{1-x} - x} = \frac{e^{-x}}{1-x}.$$

(3) $\sum_n \frac{B(n)}{n!} x^n = e^{e^x - 1}.$

(4) a_n : število 2-regularnih grafov ($deg v = 2 \forall v \in V(G)$),

F : moč množice neusmerjenih ciklov dolžine ≥ 3 : $a_n = \frac{(n-1)!}{2}$; $n \geq 3$

$$\sum_n \frac{a_n}{n!} x^n = e^{\sum_{n \geq 3} \frac{(n-1)!}{2} \frac{x^n}{n}} = e^{\frac{1}{2} \left(\log \frac{1}{1-x} - x - \frac{x^2}{2} \right)} = \frac{e^{-\frac{x}{2} - \frac{x^2}{4}}}{\sqrt{1-x}}.$$

Kompozitum:

$$(F \circ G)(x) = \sum_k \frac{a_k}{k!} G^k(x).$$

Izrek 2.6.4 (O kompoziciji).

$$F(x), G(x), F(0) = 0.$$

Potem je $(F \circ G)(x)$ erf za strukturo: množico razdelimo na bloke, vsakemu bloku damo strukturo G , množici blokov damo strukturo F .

Primer.

(1) $B(\tilde{n})$: urejena Bellova števila = število urejenih razdelitev množice $[n]$.

$$B(\tilde{2}) = 3 : \{1,2\}; \{1\},\{2\}; \{2\},\{1\}$$

$$B(\tilde{n}) = \sum_k S(n,k).$$

$B(\tilde{n})$: število vseh surjekcij iz $[n]$.

$$\sum_n \frac{B(\tilde{n})}{n!} x^n = \frac{1}{1-(e^x-1)} = \frac{1}{2-e^x}$$

$$G(x) = e^x - 1$$

$$F(x) = \frac{1}{1-x}.$$

(2) Permutacije z lihim številom ciklov

$$\sum_n a_n \frac{x^n}{n!} = \frac{e^{\log \frac{1}{1-x}} - e^{-\log \frac{1}{1-x}}}{2} = \frac{1}{2} \left(\frac{1}{1-x} - (1-x) \right).$$

$$G(x) = \log \frac{1}{1-x}$$

$$F(x) = \frac{e^x - e^{-x}}{2} \quad (F(x) - F(-x) : \text{lihi})$$

$$a_n = \begin{cases} 0 & n = 0 \\ 1 & n = 1 \\ \frac{n}{2} & n \geq 2 \end{cases}$$

orf	erf
$F(x)G(x)$	$F(x)G(x)$
$F^k(x)$	$\frac{1}{k!}F^k(x)$
$\frac{1}{1-F(x)}, F(0) = 0$	$e^{F(x)}$
$F \circ G$	$F \circ G$

2.7 Algebraične rodovne funkcije

$K[x]$ polinomi,

$K[[x]]$ formalni polimon (fp?),

$K(x)$ racionalne funkcije (polje ulomkov za $K[x]$),

$\frac{1}{x} \in K(x)$, $\frac{1}{x} \notin K[[x]]$,

$K(x) \cap K[[x]]$ racionalna rodovna funkcija.

Za taka zaporedja imamo linearne rekurzije.

$$F(x) = \sum_n a_n x^n$$

$$xF^2 - F + 1 = 0$$

$c_{n+1} = \sum_{k=0}^n c_k c_{n-k}$ kvadratična rekurzija.

Ali je $F(x) \in K(x)$?

$$F(x) = \frac{P(x)}{Q(x)}$$

$$xP^2 = PQ - Q^2 = Q(P - Q)$$

L : $\deg P \cdot 2 + 1$ - liha stopnja,

$$D : \begin{cases} \deg P < \deg Q \implies Q(P - Q) \text{ sode stopnje} \\ \deg P \geq \deg Q \implies \deg Q(P - Q) \leq 2 \cdot \deg P \end{cases}$$

Definicija 2.7.1.

$F(x) \in K[[x]]$ je algebraična reda d , če

$$Q_d(x)F^d(x) + Q_{d-1}(x)F^{d-1}(x) + \dots + Q_0(x) = 0 \text{ za } Q_0 \cdot Q_d \in K[X], Q_0, Q_d \neq 0,$$

ne obstaja taka enačba stopnje $< d$.

Algebraična reda $d =$ racionalna fpv (formalna potenčna vrsta).

$F(x) = \sum_n F_n x^n$, $M(x) = \sum_n M_n x^n$ algebraični reda 2.

$Q_d(x)F^d(x) + \dots + Q_0(x) = 0$ za $Q_0, Q_d \neq 0$

$C_n : xF(x)^2 - F(x) + 1 = 0$

$M_n : x^2 F(x)^2 + xF(x) + 1 = 0$.

S -drevo:

$S \subseteq \{1, 2, 3, \dots\}$.

Drevo s korenem, vsak element je list ali pa je število naslednikov v S .

$\{2, 3\}$ -drevo

a_n : število S -dreves z n vozlišči,

b_n : število S -dreves z n listi.

$U(x) = \sum_n a_n x^n$

$V(t) = \sum_n b_n t^n$.

$S = \{2, 3\}$

$U(x) = x + xU^2(x) + xU^3(x)$:

x : 1 vozlišče.

$V(t) = t + v^2(t) + v^3(t)$:

koren ne prispeva k številu listov.

$U(x) = x + \sum_{k \in S} xU^k(x)$

$V(t) = t + \sum_{k \in S} tV^k(t)$, $1 \notin S$.

S končna $\implies S$ algebraična.

Če S neskončna, sta U in V vseeno lahko algebraični.

Primer.

- $S = \{2\}$ - dvojiška drevesa.

$$v = t + v^2$$

$$v^2 - v + t = 0 \implies v = \frac{1 - \sqrt{1 - 4t}}{2} = \sum_{n=1}^{\infty} C_{n-1} t^n$$

C_n : število dvojiških dreves z $n + 1$ listi.

- $S = \{k\}$
 $v = t + v^k$ - Lagrangeeva inverzija (kasneje).
- $S = \{1, 2, 3, 4 \dots\}$
 $U = x + x \sum_{k=1}^{\infty} U^k = x + x \frac{U}{1-U}$
 $U - U^2 = x - xU + xU = x$
 $U^2 - U + x = 0 \implies U = \frac{1-\sqrt{1-4x}}{2} = \sum_{n=1}^{\infty} C_{n-1} x^n$
 C_n : število ravninskih dreves z $n + 1$ vozlišči.

Izkaže se: U, V algebraični $\iff S$ se za končno množico razlikuje od končne unije aritmetičnih zaporedij.

Trditev 2.7.2.

$K_{alg}[[x]] = \{F[x] \in K[[x]] \text{ algebraična}\}$ je podalgebra $K[[x]]$.

$$xF^2 - F + 1 = 0$$

$$F^2 + 2xF F' = 0$$

$$F' = \frac{F^2}{1-2xF} \stackrel{?}{=} a + bF; \quad a, b \in K(x)$$

$$F^2 = a + bF - 2a x F - 2b x F^2$$

$$(1 - 2bx)F^2 + (2ax - b)F - ax = 0$$

$$(1 - 2bx + (2ab - x))F - 1 - 2bx - ax = 0$$

\rightarrow : 2 enačbi, 2 neznanki.

$$a = \frac{1}{x(1-4x)}$$

$$b = \frac{2x-1}{x(1-4x)}$$

$$F' - \frac{1}{x(1-4x)} - \frac{2x-1}{x(1-4x)} F = 0$$

$$x(1-4x)F' - 1 - (2x-1)F = 0$$

$$F' = \sum_n n C_n x^{n+1}$$

$$[x^n] : n C_n - 4(n-1)C_{n-1} + 2C_{n-1} + C_n \text{ za } n > 1$$

$$C_n = \frac{2(n-1)}{n+1} C_{n-1} \implies \dots C_n = \frac{1}{n+1} \binom{2n}{n}.$$

Definicija 2.7.3.

$F(x) \in K[[x]]$ je D -končna, če je

$R_n(x)F^{(d)}(x) + \dots + R_1F'(x) + R_0 = 0$ za $R_i(x) \in K[x]$.

Ekvivalentno: vektorski prostor nad $K(x)$, generiran z $F, F', F'' \dots$ je končno razsežen.

Definicija 2.7.4.

$(a_n)_n$ je P -rekurzivna, če je $p_d(n)a_n + \dots + p_0(n)a_{n-d} = 0$ za $n \geq d$.

Trditev 2.7.5.

$F(x) = \sum_n a_n x^n$ je D -končna $\iff (a_n)_n$ je P -rekurzivna.

Torej: za P -rekurzivno zaporedje lahko člene hitro izračunamo.

Zgled.

$F(x) = \sum_n C_n x^n$ je D -končna,

e^x je D -končna: $F' - F = 0$,

e^x ni algebraična.

Izrek 2.7.6.

$F(x)$ algebraična $\implies D$ -končna.

Dokaz 2.7.7. (skica):

$$Q_d(x)F^d(x) + \dots + Q_0(x) = 0 \quad /'$$

$$Q_d(x)'F^d(x) + dQ_d(x)F^{d-1}(x)F'(x) + \dots + Q_0'(x) = 0$$

$$F'(x) \in K(x, F(x))$$

Iz algebre:

K obseg, u v večjem obsegu;

(i) v algebraičnem: $K[u] = K(u)$ končno razsežen VP,

(ii) v transcendentnem: $K[u] \subseteq K[x]$ („ u spremenljivka“).

$$K = K[x]$$

$$u = F(x)$$

$$K[u] = K(x, F(x)).$$

Torej: $K(x, F(x))$ je končno razsežen VP nad $K(x)$, torej so $1, F, F' \dots$ linearno neodvisni $\implies F$ je D -končna.

2.8 Eulerjeva in eulerska števila

E_n : število alternirajočih permutacij v S_n .

$$E_3 = 2 \text{ (231), (132)}$$

$$2E_{n+1} = \sum_{k=0}^n \binom{n}{k} E_k E_{n-k} + \delta_{n0}$$

$$E(x) = \sum_n \frac{E_n}{n!} x^n$$

$$2F' = F^2 + 1$$

$$\int \frac{2dF}{F^2+1} = \int dx$$

$$2 \arctan F = x + 2c$$

$$F = \tan\left(\frac{x}{2} + c\right)$$

$$F(0) = 1 = \tan c \implies c = \frac{\pi}{4}$$

$$\tan\left(\frac{x}{2} + \frac{\pi}{4}\right) = \frac{\tan \frac{x}{2} + 1}{1 - \tan \frac{x}{2}} = \frac{\sin \frac{x}{2} + \cos \frac{x}{2}}{\cos \frac{x}{2} - \sin \frac{x}{2}} = \frac{1 + \sin x}{\cos x}.$$

Izrek 2.8.1.

$$\sum_n \frac{E_n}{n!} x^n = \frac{1 + \sin x}{\cos x} \text{ oz.}$$

$$\frac{1}{\cos x} = \sum_n \text{sod} \frac{E_n}{n!} x^n$$

$$\frac{1}{\sin x} = \sum_n \text{lih} \frac{E_n}{n!} x^n$$

Opomba.

Bernoullijeva števila.

$$B_n = \begin{cases} 1 & n = 0 \\ \frac{1}{2} & n = 1 \\ 0 & n > 1, n \text{ lih} \\ \frac{(-1)^{\frac{n}{2}+1} E_{n-1}}{2^n (2^n - 1)} & n > 0, n \text{ sod} \end{cases}$$

$$\sum_n B_n \frac{x^n}{n!} = \frac{x e^x}{e^x - 1}$$

$$\sum_{n=1}^{\infty} \frac{1}{n^2} = \frac{\pi^2}{6}$$

$$\sum_{n=1}^{\infty} \frac{1}{n^4} = \frac{\pi^4}{90}$$

$$\sum_{n=1}^{\infty} \frac{1}{n^{2k}} = \frac{2^{2n} (-1)^{k+1} \pi^{2k}}{2 \cdot (2k)!} = \frac{E_{2k-1} \pi^{2k}}{2(2k-1)!(2^{2k}-1)} = \zeta(2k)$$

Riemmanova funkcija ζ :

$$\zeta(s) = \sum_{n=1}^{\infty} \frac{1}{n^s} \text{ za } \operatorname{Re} s > 1.$$

Z analitičnim nadaljevanjem lahko ζ definiramo na $\mathbb{C} \setminus \{1\}$.

$\zeta(-n) = \frac{B_{n+1}}{n+1}$ - soda negativna števila so ničle - trivialne ničle.

Riemmanova hipoteza:

$\operatorname{Re} z = \frac{1}{2}$ za vsako netrivialno ničlo z funkcije ζ .

$$\zeta(-1) = -\frac{B_2}{2} = -\frac{1}{12}$$

$$„\sum_{n=1}^{\infty} n = -\frac{1}{12}“$$

Faulhaberjeva formula:

$$\sum_{i=1}^n i = \frac{n(n+1)}{2} = \frac{n^2}{2} + \frac{n}{2}$$

$$\sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6} = \frac{n^3}{3} + \frac{n^2}{2} + \frac{n}{6}$$

$$\sum_{i=1}^n i^3 = \frac{n^2(n+1)^2}{4} = \frac{n^4}{4} + \frac{n^3}{2} + \frac{n^2}{4}$$

$$\sum_{i=1}^n i^k = \frac{1}{k+1} \sum_{l=0}^k \binom{k+1}{l} B_l n^{k+1-l}$$

$$= \frac{n^{k+1}}{k+1} + \frac{n^k}{2} + \sum_{l=1}^{\lfloor \frac{k}{2} \rfloor} \frac{(-1)^{l+1} \binom{k}{2l-1}}{2^{2l}(2^{2l}-1)} E_{2l-1} n^{k-1-2l}$$

$$= \log n + \gamma + \frac{1}{2n} - \sum_{k=1}^{\infty} \frac{B_{2k}}{2k \cdot n^{2k}}$$

$$= \log n + \gamma + \frac{1}{2n} - \sum_{k=1}^{\infty} \frac{(-1)^k E_{2k-1}}{2^{2k}(2^{2k}-1)n^{2k}},$$

kjer je $H_n = 1 + \frac{1}{2} + \dots + \frac{1}{n}$ n -to harmoično število.

$A(n, k)$: število permutacij v S_n z $k-1$ spusti.

$$A(n, k) = (n+1-k)A(n-1, k-1) + kA(n-1, k)$$

$$\frac{1}{1-x} = 1 + x + x^2 + \dots \quad /' / \cdot x$$

$$\frac{x}{(1-x)^2} = x + 2x^2 + 3x^3 + \dots \quad /' / \cdot x$$

$$\frac{x+x^2}{(1-x)^3} = x + 4x^2 + 9x^3 + \dots \quad /' / \cdot x$$

$$\frac{x+4x^2+x^3}{(1-x)^4} = x + 8x^2 + 27x^3 + \dots$$

$$A_n(x) = \sum_k A(n, k)x^k \text{ eulerski polinom.}$$

Izrek 2.8.2.

$$\sum_m m^n x^m = \frac{A_n(x)}{(1-x)^{n+1}}.$$

Dokaz 2.8.3.

Indukcija:

$$n=0: \frac{1}{1-x} = \frac{1}{1-x}$$

$$n-1 \rightarrow n:$$

$$\sum_m m^{n-1} x^m = \frac{A_{n-1}(x)}{(1-x)^n} \quad /' / \cdot x$$

$$x \cdot \sum_m m^{n-1} x^{m-1} = \frac{A'_{n-1}(x)(1-x)^n + A_{n-1}(x)n(1-x)^{n-1}}{(1-x)^{2n}} \stackrel{?}{=} \frac{A_n(x)}{(1+x)^{n+1}}$$

$$[x^k]: (k+1)A(n-1, k-1) - kA(n-1, k) + nA(n-1, k) = A(n, k) \quad \checkmark$$

$$A_{n-1}(x) = \sum_k A(n-1, k)x^k$$

$$A'_{n-1}(x) = \sum_k kA(n-1, k)x^{k-1}.$$

Izrek 2.8.4.

$\sum_{n,k} A(n, k)x^k \frac{y^n}{n!} = \frac{1-x}{1-xe^{xy(1-y)}}$ - mešana rodovna funkcija (običajna v x , eksponentna v y).

Dokaz 2.8.5.

$$\begin{aligned} & \sum_{n,k} A(n, k)x^k \frac{y^n}{n!} \\ &= (1-x) \left(\sum_k \frac{A_n(x)}{(1-x)^{n+1}} \cdot \frac{y^n}{n!} (1-x)^n \right) \\ &= (1-x) \sum_n \left(\sum_m m^n x^m \right) \frac{y^n (1-x)^n}{n!} \\ &= (1-x) \sum_m \left(\sum_n \frac{m^n y^n (1-x)^n}{n!} \right) x^m \\ &= (1-x) \sum_m e^{xy(1-x)} x^m \\ &= \frac{1-x}{1-e^{xy(1-x)}}. \end{aligned}$$

2.9 Izračun povprečij in variance

Koliko elementov ima v povprečju podmnožica $[n]$?

$$\frac{\sum_{T \subseteq [n]} |T|}{2^n} = \frac{\sum_n k \binom{n}{k}}{2^n} = \frac{n \cdot 2^{n-1}}{2^n} = \frac{n}{2}$$

$$(1+x)^n = \sum_k \binom{n}{k} x^k \quad /'$$

$$n(1+x)^{n-1} = \sum_k k \binom{n}{k} x^{k-1}$$

$x = 1$:

$$n \cdot 2^{n-1} = \sum_k k \binom{n}{k}.$$

S končna množica.

$$F(x) = \sum_{a \in S} x^{f(a)}$$

$$F(1) = |S|$$

$$F'(x) = \sum_{a \in S} f(a) \cdot x^{f(a)-1}$$

$$F'(1) = \sum_{a \in S} f(a)$$

$$\mu = \frac{F'(1)}{F(1)} = (\log' F)(1)$$

$$F(x) = (1+x)^n$$

$$\log F(x) = n \log(1+x)$$

$$\log' F(x) = \frac{n}{1+x}$$

$$(\log' F)(1) = \frac{n}{2}$$

$$\sigma^2 = E(x^2) - \mu^2$$

$$E(x^2) = \frac{\sum_n f^2(s)}{|S|}$$

$$F'(x) + xF''(x) = (xF'(x))' = \sum_{a \in S} f^2(a)x^{f(a)-1}$$

$$x = 1:$$

$$\sigma^2 = \frac{F'(1)+F''(1)}{F(1)} - \frac{F'(1)^2}{F(1)^2} = \frac{F'(1)}{F(1)} + \frac{F''(1)F(1)-F'(1)^2}{F(1)^2}.$$

$$\text{Torej}$$

$$\mu = (\log' F)(1)$$

$$\sigma^2 = (\log' F)(1) + (\log'' F)(1)$$

$$F(x) = (1+x)^n$$

$$\mu = \frac{n}{2}$$

$$\log' F(x) = \frac{n}{1+x}$$

$$\log'' F(x) = -\frac{n}{(1-x)^2}$$

$$\sigma^2 = \frac{n}{2} - \frac{n}{4} = \frac{n}{4}$$

$$\frac{n}{2} \pm \frac{\sqrt{n}}{2}.$$

$$\text{Koliko ciklov ima v povprečju permutacija v } S_n?$$

$$\sum_{\pi \in S_n} x^{f(\pi)} = \sum_k c(n,k)x^k = x^{\bar{n}} = F(x)$$

$$\log F(x) = \log x + \log(x+1) + \dots + \log(x+n-1)$$

$$\log' F(x) = \frac{1}{x} + \dots + \frac{1}{x+n-1}$$

$$\mu = H_n = \log n + \gamma + o(1)$$

$$\log'' F(x) = -\frac{1}{x^2} - \dots - \frac{1}{(x+n-1)^2}$$

$$\sigma^2 = H_n - \sum_{i=1}^n i^2 = \log n + \gamma - \frac{\pi^2}{6} + o(1)$$

$$\log n \pm \sqrt{\log n}.$$

2.10 Lagrangeeva inverzija

$K[x]$ algebra polinomov,

$K(x)$ obseg racionalnih funkcij (obseg ulomkov $K[x]$),

$K[[x]]$ algebra formalnih potenčnih vrst,

$K((x)) = \{\sum_{n \geq n_0} a_n x^n; n_0 \in \mathbb{Z}, a_i \in K\}$ obseg formalnih Laurentovih vrst

(obseg ulomkov $K[[x]]$).

$$\frac{F(x)}{G(x)} = \frac{F(x)}{x^m H(x)}, \frac{F(x)}{H(x)} \in K[[x]], H(0) \neq 0.$$

Seštevanje, množenje, odvod, kompozitum, valuacija ($\in \mathbb{Z}$).

$$\text{res}F(x) = [x^{-1}]F(x) \text{ residuum.}$$

Lema 2.10.1. $\text{res}F(x) = 0 \leftrightarrow F(x) = G'(x)$ za $K((x))$.

Dokaz 2.10.2.

(\Leftarrow)

$$F(x) = \left(\sum_{n \geq n_0} b_n x^n \right) = \left(\sum_{n \geq n_0} n b_n x^{n-1} \right)$$

$$[x^{-1}]F(x) = 0 \cdot b_0 = 0.$$

(\Rightarrow)

$$F(x) = \sum_{n \geq n_0} a_n x^n$$

$$G(x) = \sum_{n \geq n_0} \frac{a_{n-1} x^n}{n}$$

$$a_{-1} = 0.$$

Lema 2.10.3.

$$F(x) \in K((x)), F(x) \neq 0, \text{res}_{F(x)}^{F'(x)} = v(F(x)).$$

Dokaz 2.10.4.

$$F(x) = x^{n_0} G(x)$$

$$n_0 = v(F(x))$$

$$G(x) \in K[[x]], G(0) \neq 0$$

$$\frac{F'(x)}{F(x)} = \frac{n_0 x^{n_0-1} G(x) + x^{n_0} x^{n_0} G'(x)}{x^{n_0} G(x)} = \frac{n_0}{x} + \frac{G'(x)}{G(x)}$$

$$\frac{G'(x)}{G(x)} \in K[[x]].$$

Lagrangeeva inverzija (1. verzija):

$$F \in K[[x]]$$

$$v(F(x)) = 1$$

$$n \cdot [x^n] (F^{<-1>}(x))^k = k \cdot [x^{-k}] F^{-n}(x);$$

$$F^{-n}(x) \in K((x)).$$

$$\text{Torej: } n \cdot [x^n] F^{<-1>}(x) = \text{res} F^{-1}(x).$$

$$\textbf{Dokaz 2.10.5. } (F^{<-1>}(x))^k = \sum_{m \geq k} c_m x^m$$

$$x \leftrightarrow F(x)$$

$$x^k = \sum_{m \geq k} c_m (F(x))^m \quad /'$$

$$kx^{k-1} = \sum_{m \geq k} m c_m F^{m-1}(x) F'(x) \quad / : F^n(x)$$

$$\frac{kx^{k-1}}{F^n(x)} = \sum_{m \geq k} m c_m F^{m-n-1}(x) F'(x) \quad / \text{res}$$

$$[x^{-1}] \frac{kx^{k-1}}{F^n(x)} = [x^{-k}] \frac{k}{F^n(x)}$$

$$F^{m-n-1}(x) F'(x) = \frac{(F^{m-n}(x))'}{m-n}; \quad m \neq n$$

$$\text{res} \left(F^{m-n-1}(x) F'(x) \right) = 0 \text{ \u0107e } m \neq n \text{ in 1 sicer (lemi)}$$

$$\rightarrow n \cdot a_n \cdot 1 \text{ (leva stran).} \quad \blacksquare$$

Primer.

$$F(x) = x - x^2$$

$$F^{<-1>}(x) = ?$$

$$n[x^n] F^{<-1>}(x) = [x^{-1}] \left(\frac{1}{1-x^2} \right)^n = [x^{-n}] \frac{x^{-n}}{(1-x)^n}$$

$$\frac{1}{(1-x)^n} = \sum_m \binom{m+n-1}{n-1} x^m$$

$$[x^n] F^{<-1>}(x) = \frac{1}{n} \binom{2n-2}{n-1} = C_{n-1}.$$

\u0160e ena razlaga:

$$y - y^2 = x$$

$$y^2 - y + x = 0 \implies y = \frac{1 \pm \sqrt{1-4x}}{2} \implies y = x \sum_n C_n x^n.$$

Lagrangeeva inverzija (2. verzija)

$$F(x) = xG(F(x))$$

$$F(x) \in K[[x]]$$

$$G(x) \in K[[x]], G(0) \neq 0, v(F) = 1$$

$$[x^k] F(x)^k = k[x^{n-k}] G(x)^n.$$

Dokaz 2.10.6.

$$f(x) := \frac{x}{G(x)}, v(f) = 1$$

$$f(F(x)) = \frac{F(x)}{G(F(x))} = 1 \rightarrow \text{ima levi inverz, tudi desni.}$$

$$n[x^n] F(x)^k = k[x^n] (f^{<-1>}(x))^k$$

$$= k[x^{-k}]f^{-k}(x) = k[x^{-k}]x^{-n}G^n(x).$$

Primer.

$$(a) \ S = \{k\}$$

$$k = 3$$

a_n : število k -dreves na n vozliščih.

$$v(x) = \sum_n a_n x^n$$

$$V(x) = x + xV^k(x) = x(1 + V^k(x))$$

$$G(x) = (1 + x)^n$$

$$n[x^n]V(x)[x^{n-1}](1 + x^k)^n = k[x^{n-1}]\sum_{i=0}^n \binom{n}{i} x^{k_i};$$

$$n = ki + 1, i \in \mathbb{N}, a_n = a_{ki+1} = \frac{1}{n} \cdots = \frac{1}{ki+1} \binom{ki+1}{i}.$$

$$(b) \text{ Vpeta drevesa v } K_n.$$

r_n : število vpetih dreves s korenem v K_n .

$$R(x) = \sum_n \frac{r_n}{n!} x^n \text{ (vozlišča so označena).}$$

Označimo drevo s korenem = koren + množica blokov, ki jim damo strukturo označenega drevesa s korenem.

$$R(x) = xe^{R(x)}$$

$$G(x) = e^x$$

$$n[x^n]R(x) = [x^{n-1}]e^{nx}$$

$$e^n = \sum_k \frac{n^k x^k}{n!}$$

$$\frac{nr_n}{n!} = \frac{n^{n-1}}{(n-1)!}$$

$$r_n = n^{n-1}$$

Število vpetih dreves v K_n je n^{n-2} .

2.11 Asimptotika koeficientov

$$K = \mathbb{C}$$

$$F(x) = \sum_n a_n x^n$$

$F(x) \in \mathbb{C}[[x]]$ ima pozitiven konvergenčni polmer

$$R = \frac{1}{\limsup_{n \rightarrow \infty} \sqrt[n]{|a_n|}}.$$

F je holomorfna v okolici 0.

Za $\forall \epsilon > 0$:

- $|a_n| < \frac{1}{R} + \epsilon$ za $\forall n \geq n_0$,
- $|a_n| > \frac{1}{R} - \epsilon$ za neskončno mnogo n .

Npr. $F(z) = \frac{1}{1-z} = 1 + z + z^2 + \dots$

$R = 1$,

$|a_n| < (1 + \epsilon)^n$ za $\forall n$,

$|a_n| > (1 - \epsilon)^n$ za vse sode n .

$R = \infty \implies F(z)$ cela funkcija.

$R < \infty \implies F(z)$ ima singularnost v z_0 , $|z_0| = R$.

Definicija 2.11.1. f ima v z_0 pol reda r , če ima $f(z)(z - z_0)^r$ odpravljivo singularnost v z_0 , $\lim_{z \rightarrow z_0} f(z)(z - z_0)^r \neq 0$.

Funkcija je meromorfna, če so vse singularnosti poli in množica polov nima stekališč (oz. je diskretna).

$$f(z)(z - z_0)^r = b_0 + b_1(z - z_0) + b_2(z - z_0)^2 + \dots \quad / : (z - z_0)^n$$

V kombinatoriki: $1 - \frac{z}{z_0}$, $b_i \mapsto b_{i-r}$

$$f(z) = b_{-r} + b_{-r+1} \left(1 - \frac{z}{z_0}\right) + \dots + b_{-1} \left(1 - \frac{z}{z_0}\right)^{-1} + b_0 + b_1 \left(1 - \frac{z}{z_0}\right) + \dots$$

Glavni del (angl. principal part):

$$PP_{f,z_0}(z) = b_{-r} \left(1 - \frac{z}{z_0}\right)^r + \dots + b_{-1} \left(1 - \frac{z}{z_0}\right)^{-1}.$$

Če je z_0 edina singularnost na $|z| = R$:

$f(z) - PP_{f,z_0}(z)$ ima konvergenčni polmer $R' > R$.

$$[z^n]PP_{f,z_0}(z) = \left(\sum_{i=1}^r b_{-i} \binom{n+i-1}{i-1}\right) z_0^n \sim \frac{b_{-r} n^{r-1}}{z_0^n (r-1)!}.$$

$$\forall \epsilon > 0 : [z^n] |f(z) - PP_{f,z_0}(z)| < \left(\frac{1}{R'} + \epsilon\right)^n \text{ za } n \geq n_0.$$

$$\frac{1}{R'} + \epsilon < \frac{1}{R}$$

$$\lim_{n \rightarrow \infty} \frac{\left(\frac{1}{R'} + \epsilon\right)^n}{\left(\frac{1}{R}\right)^n} = 0.$$

Izrek 2.11.2.

$F(z) \in \mathbb{C}[[x]]$, $R \in (0, \infty)$, z_0 edina singularnost na $|z_0| = R$, z_0 je pol reda r . Potem je

$$[z^n]F(z) \sim \frac{b_{-r}n^{r-1}}{z_0^n(r-1)!}, \text{ kjer je}$$

$$b_{-r} = \lim_{z \rightarrow z_0} f(z) \left(1 - \frac{z}{z_0}\right)^r.$$

Primer.

$$(1) \quad f(z) = \frac{1}{(1-z)(1-2z)}$$

$$R = \frac{1}{2}, z_0 = \frac{1}{2}, r = 1$$

$$\lim_{z \rightarrow \frac{1}{2}} \frac{1}{(1-z)(1-2z)} (1-2z) = 2 = b_{-1}$$

$$a_n \sim \frac{2}{\left(\frac{1}{2}\right)^n} = 2^{n+1}.$$

$$(2) \quad d_n: \text{ število premestitev v } S_n$$

$$\sum_n \frac{d_n}{n!} z^n = \frac{e^{-z}}{1-z}$$

$$z_0 = 1, r = 1$$

$$b_{-1} = \lim_{z \rightarrow 1} \frac{e^{-z}}{1-z} (1-z) = e^{-1}$$

$$\frac{d_n}{n!} \sim \frac{e^{-1}}{1 \cdot 1} = \frac{1}{e}$$

$$d_n \sim \frac{n!}{e}.$$

Koliko dober je za približek?

$\frac{e^{-z}}{1-z} - \frac{e^{-1}}{1-z}$ je cela funkcija.

$$[z^n] (\text{cela funkcija}) < \left(\frac{1}{R} + \epsilon\right)^n = \epsilon^n \text{ za } n \geq n_0.$$

Koeficienti celih funkcij hitro padajo proti 0.

Ker je $z_0 = 1$ edini pol in ker je enostaven, je $\frac{b_{-1}}{z_0^n}$ odličen približek.

$$d_n = \left\lfloor \frac{n!}{e} \right\rfloor.$$

$$(3) \quad \tilde{B}(n): \text{ urejena Bellova števila}$$

$$\tilde{B}(n) = \sum_k k! S(n, k)$$

$$\sum_n \tilde{B}(n) \frac{z^n}{n!} = \frac{1}{1-(e^z-1)} = \frac{1}{2-e^z}.$$

Poli so $\log 2 + 2k\pi i$, $k \in \mathbb{Z}$

$$z_0 = \log 2, r = 1$$

$$b_{-1} = \lim_{z \rightarrow \log 2} \frac{1 - \frac{z}{\log 2}}{2 - e^z} \stackrel{L'H}{=} \lim_{z \rightarrow \log 2} \frac{-\frac{1}{\log 2}}{-e^z} = \frac{1}{2 \log 2} = \frac{1}{\log 4}$$

$$\tilde{B}(n) \sim \frac{n!}{2(\log 2)^{n+1}}$$

$$\tilde{B}(20) = 267 \dots 115 \text{ (23 števk)}$$

$$\left\lfloor \frac{20!}{2(\log 2)^{21}} \right\rfloor = 267 \dots 088$$

$$\frac{\log 2}{\log 2 + 2\pi i} \doteq 0.11.$$

(4) n hiš.

1. družina se vseli v naključno hišo,

2. družina se vseli v naključno naslednjo hišo,

a_n : pričakovano število zasedenih hiš, $\frac{n}{3} < a_n < \frac{n}{2}$?

$$a_0 = 0, a_1 = 1, a_2 = 1, a_3 = \frac{1}{3} \cdot 1 + \frac{2}{3} \cdot 2 = \frac{5}{3}.$$

$$a_n = \frac{1}{n} \sum_{i=1}^n (a_{i-2} + a_{n-i-1} + 1) \quad / \cdot n$$

$$na_n = n + 2(a_0 + a_1 + \dots + a_{n-2})$$

$$F(x) = \sum_n a_n x^n$$

$$xF'(x) + 2xF(x) + 2F(x) = \frac{x}{(1-x)^2} + \frac{2F(x)}{1-x} - \text{linearna DE 1. reda.}$$

$$F(x) = \frac{1-e^{-2x}}{2(1-x)^2}$$

$$z_0 = 1, r = 2$$

$$b_{-2} = \lim_{z \rightarrow 1} \frac{1-e^{-2z}}{2(1-z)^2} (1-z)^2 = \frac{1-e^{-2}}{2 \cdot 1!}$$

$$a_n \sim \left(\frac{1-e^{-2}}{2}\right)^n$$

$$\frac{1-e^{-2}}{2} \doteq 0.423 \in \left(\frac{1}{3}, \frac{1}{2}\right).$$

Kaj pa, če imamo več singularnosti na $|z| = R$?

$z_1 \dots z_k$ poli redov $r_1 \dots r_k$

$$[z^n]f(z) = \sum_{i=1}^k \frac{b_{-r_i} n^{r_i-1}}{z_i^n (r_i-1)!} + O\left(\left(\frac{1}{R'}\right)^n\right), R' > R.$$

Primer.

$$r(x) = \frac{1}{1-z} + \frac{1}{1+z} + \frac{1}{1-z^2}$$

$$a_n = 1 + (-1)^n + \frac{1}{2^n} \asymp 1 + (-1)^n.$$

V praksi štejejo le najvišji poli.

Primer.

$$(a) \sum_n \overline{p}_k(n) x^n = \prod_{i=1}^k \frac{1}{1-x^i}.$$

Racionalna funkcija, poli

1 reda k , -1 reda $\lfloor \frac{k}{2} \rfloor$, $e^{\pm \frac{2\pi i}{3}}$ reda $\lfloor \frac{k}{3} \rfloor \dots$

1 ima najvišji red.

$$z_0 = 1, r = k$$

$$b_{-k} = \lim_{z \rightarrow 1} \prod_{i=1}^k \frac{1}{1-z^i} (1-z)^k = \lim_{z \rightarrow 1} \prod_{i=1}^k \frac{1}{1+z+\dots+z^{i-1}} = \frac{1}{k!}$$

$$\overline{p_k}(n) \sim \frac{n^{k-1}}{k!(k-1)!}$$

$$\sum_k p_k(n) x^k = x^k \prod_{i=1}^k \frac{1}{1-x^i}$$

$$p_k(n) \sim \frac{n^{k-1}}{k!(k-1)!}.$$

(Šibke) kompozicije n s k členi

$$\binom{n+k-1}{k-1} \sim \frac{n^{k-1}}{(k-1)!}$$

$$\binom{n-1}{k-1} \sim \frac{n^{k-1}}{(k-1)!}$$

$$\sum_n p(n) x^n = \prod_{i=1}^{\infty} \frac{1}{1-x^i} \text{ - ni racionalna funkcija.}$$

Singularnosti so bistvene, množica singularnosti ima stekališča.

Lema 2.11.3.

$$\alpha \in \mathbb{R}.$$

$$\lim_{x \rightarrow \infty} \frac{\Gamma(x+\alpha)}{x^\alpha \Gamma(x)} = 1.$$

$$\Gamma(x) = \int_0^x t^{x-1} e^{-t} dt$$

$$\Gamma(x+1) = x\Gamma(x)$$

$$\Gamma(n) = (n-1)! \quad n = 1, 2, 3, \dots$$

Γ lahko razširimo na $\mathbb{C} \setminus \{0, -1, -2, \dots\}$.

$$n! \sim \sqrt{2\pi n} \left(\frac{n}{e}\right)^n \text{ Stirlingova formula.}$$

$$\lim_{x \rightarrow \infty} \frac{\Gamma(x+\alpha)}{\sqrt{2\pi x} \left(\frac{x}{e}\right)^x} = 1.$$

Dokaz 2.11.4.

$$\begin{aligned} \lim_{x \rightarrow \infty} \frac{\Gamma(x+\alpha)}{x^\alpha \Gamma(x)} &= \lim_{x \rightarrow \infty} \frac{\sqrt{2\pi(x+\alpha-1)} \left(\frac{x-\alpha-1}{e}\right)^{x+\alpha-1}}{x^\alpha \cdot \sqrt{2\pi(x-1)} \left(\frac{x-\alpha}{e}\right)^{x-1}} \\ &= \lim_{x \rightarrow \infty} \frac{1}{e^\alpha} \left(\left(1 + \frac{\alpha}{x-1}\right)^{\frac{x-1}{\alpha}} \right)^\alpha \\ &= \frac{e^\alpha}{e^\alpha} \\ &= 1. \end{aligned}$$

Lema 2.11.5.

$$\beta \in \mathbb{R} \setminus \mathbb{N}.$$

$$\binom{\beta}{n} \sim \frac{(-1)^n}{\Gamma(-\beta)n^{\beta+1}}.$$

Dokaz 2.11.6.

$$\begin{aligned} & \lim_{n \rightarrow \infty} \frac{\beta(\beta-1)\dots(\beta-n+1)\Gamma(-\beta)}{n!(-1)^n} \\ &= \lim_{n \rightarrow \infty} \frac{n^{\beta+1}\Gamma(-\beta+n)}{\Gamma(n+1)} \\ &\stackrel{\text{lema}}{=} 1; \end{aligned}$$

$$x = n - \beta, \alpha = \beta + 1.$$

$$z_0 \in \mathbb{R}$$

$$f(z) = \left(1 - \frac{z}{z_0}\right)^\beta g(z)$$

$$\beta \in \mathbb{Z} \setminus \mathbb{N}: \text{ pol,}$$

$$\beta \notin \mathbb{Z} \setminus \mathbb{N}: \text{ algebraična singularnost.}$$

$$\text{Tipično: } \beta = \frac{1}{2}, \text{ npr. } f(z) = \sqrt{1-z}.$$

$$g \text{ analitična v } 0 \text{ s polmerom } > |z_0|.$$

$$\begin{aligned} f(z) &= \left(1 - \frac{z}{z_0}\right)^\beta \left(b_0 + b_1 \left(1 - \frac{z}{z_0}\right) + \dots\right) \\ &= b_0 \left(1 - \frac{z}{z_0}\right)^\beta + b_1 \left(1 - \frac{z}{z_0}\right)^{\beta+1} + \dots \end{aligned}$$

$$[z^n]f(z) = b_0 \binom{\beta}{n} \frac{(-1)^n}{z_0^n} + b_1 \binom{\beta}{n} \frac{(-1)^n}{z_0^2} + \dots$$

$$b_0 \binom{\beta}{n} \frac{(-1)^n}{z_0^n} \sim b_0 \cdot \frac{1}{\Gamma(-\beta)n^{\beta+1}z_0^n},$$

$$b_1 \binom{\beta+1}{n} \frac{(-1)^n}{z_0^n} \sim b_1 \cdot \frac{1}{\Gamma(-\beta-1)n^{\beta+2}z_0^n}.$$

$$\frac{1}{n^{\beta+1}} > \frac{1}{n^{\beta+2}} \rightarrow \text{majhno.}$$

Izrek 2.11.7.

$f(z) = \left(1 - \frac{z}{z_0}\right)^\beta g(z)$, $z_0 \in \mathbb{R}$, $\beta \in \mathbb{R} \setminus \mathbb{N}$, $g(z_0) \neq 0$, g holomorfná s konvergenčnim polmerom $> |z_0|$. Potem je

$$[z^n]f(z) \sim \frac{g(z_0)}{\Gamma(-\beta)n^{\beta+1}z_0^n}.$$

$$\text{V posebnem: } b = -r : \frac{b_{-r}n^{r-1}}{\Gamma(r)z_0^n}.$$

Primer.

$$(1) \quad F(x) = \sum_n C_n x^n$$

$$F(x) = 1 + xF^2(x)$$

$$F(x) = \frac{1 - \sqrt{1-4x}}{2x}$$

$$xF(x) = \frac{1}{2} + \frac{1}{2}\sqrt{1-4x}$$

$$x_0 = \frac{1}{4}, \beta = \frac{1}{2}, g(x) = -\frac{1}{2}$$

$$C_{n-1} \sim \frac{-\frac{1}{2}}{\Gamma(-\frac{1}{2})n^{\frac{3}{2}}(\frac{1}{4})^n}$$

$$\Gamma\left(\frac{1}{2}\right) = \sqrt{\pi}$$

$$\Gamma\left(-\frac{1}{2}\right) = -\frac{1}{2}\Gamma\left(-\frac{1}{2}\right)$$

$$\Gamma\left(-\frac{1}{2}\right) = -2\sqrt{\pi}$$

$$C_{n-1} \sim \frac{-\frac{1}{2}4^n}{-2\sqrt{\pi}n^{\frac{3}{2}}} = \frac{4^{n-1}}{\sqrt{\pi}n^{\frac{3}{2}}}.$$

D.N. Dokažite to formulo iz $C_n = \frac{1}{n+1} \binom{2n}{n}$ in Stirlingovo formulo.

$$(2) \quad M(k) = \sum_n M_n x^n$$

$$M(x) = 1 + xM(x) + x^2M^2(x)$$

$$x^2M^2 + (x-1)M + 1 = 0$$

$$M(x) = \frac{1-x-\sqrt{1-2x-3x^2}}{2x^2}$$

$$x^2M = \frac{1-x}{2} - \frac{1}{2}\sqrt{(1-3x)(1+x)}$$

$$x_0 = \frac{1}{3}, \beta = \frac{1}{2}, g(x) = -\frac{1}{2}\sqrt{1+x}$$

$$M_{n-2} \sim \frac{-\frac{1}{2}\sqrt{\frac{4}{3}}}{-2\sqrt{\pi}n^{\frac{3}{2}}(\frac{1}{3})^n}$$

$$M_n \sim \frac{3^{\frac{3}{2}} \cdot 3^n}{2\sqrt{\pi}n^{\frac{3}{2}}}.$$

Kaj pa, če je $f(n)$ cela?

Izrek 2.11.8 (Haymanova metoda). Naj bo $f(z)$ dopustna funkcija (brez definicije), npr. $f(z) = e^{P(z)}$, P polinom, $[z^n]f(z) > 0$ od nekega n naprej (npr. e^z , $e^{z+\frac{z^2}{2}}$, ne pa e^{z^2}).

$$\beta(z) := \frac{zf'(z)}{f(z)}.$$

Potem ima enačba $\beta(z) = n$ natanko eno pozitivno rešitev z_n .

$$[z^n]f(z) \sim \frac{f(z_n)}{z_0^n \sqrt{2\pi z_n} \beta'(z_n)}.$$

Primer.

$$(1) \quad f(z) = e^z$$

$$\beta(z) = \frac{ze^z}{e^z} = z$$

$$z_n = n$$

$$[z^n]f(z) \sim \frac{e^n}{n^n \sqrt{2\pi n}} - \text{Stirlingova formula.}$$

$$(2) \quad f(z) = e^{z+\frac{z^2}{2}}$$

$$\beta(z) = \frac{z \cdot e^{z+\frac{z^2}{2}} (1+z)}{e^{z+\frac{z^2}{2}}} = z^2 + z$$

$$z^2 + z + n = 0$$

$$z_n = \frac{-1+\sqrt{1+4n}}{2}$$

$$\frac{i_n}{n!} \sim \frac{e^{\left(\frac{-1+\sqrt{1+4n}}{2}\right)^2 + \frac{-1+\sqrt{1+4n}}{2}}}{\left(\frac{-1+\sqrt{1+4n}}{2}\right)^n \sqrt{2\pi \frac{-1+\sqrt{1+4n}}{2}} \sqrt{1+4n}} \sim \dots$$

Poglavje 3

Incidenčne algebre in Möbiusova inverzija

3.1 Motivacija

$$f, g : \mathbb{N} \rightarrow \mathbb{R}$$

$$g(n) = f(0) + f(1) + \cdots + f(n) \quad n \in \mathbb{N}$$

$$f(n) = g(n) - g(n-1)$$

$$(g(x) = \int_0^x f(t)dt, \quad g'(x) = f(x)).$$

$$f, g : \mathbb{N} \setminus \{0\} \rightarrow \mathbb{R}$$

$$g(n) = \sum_{d|n} f(d)$$

$$f(n) = \sum_{d|n} \mu\left(\frac{n}{d}\right) g(d) \text{ klasična Möbiusova inverzija, } \mu \text{ klasična Möbiusova}$$

funkcija, $\mu(n) \in \{-1, 0, 1\}$.

$$f, g : 2^{[n]} \rightarrow \mathbb{R}$$

$$g(T) = \sum_{S \subseteq T} f(S)$$

$$f(T) = \sum_{S \subseteq T} (-1)^{|T \setminus S|} g(S) \text{ - NVI.}$$

3.2 Delno urejene množice

(P, \leq) je delno urejena množica (dum) (angl. partially ordered set oz. poset);

refleksivnost: $x \leq x$, antisimetričnost: $x \leq y, y \leq x \implies x = y$, tranzitivnost: $x \leq y, y \leq z \implies x \leq z$.

Primer.

$$(1) ([n], \leq) = \underline{n} = \mathbf{n} \\ (\mathbb{N}, \leq).$$

$$(2) (D_n, |) = D_n \text{ delitelji } n \\ (\mathbb{N} \setminus \{0\}, |) = D.$$

$$(3) (2^{[n]}, \subseteq) = B_n \text{ Boolova algebra.}$$

$$(4) (\{\text{razdelitve } [n]\}, \leq) \\ \leq: \text{ biti finejša } \pi \leq \sigma: \text{ vsak blok v } \pi \text{ je vsebovan v bloku v } \sigma \\ 14 - 2 - 378 - 56 \leq 12456 - 378.$$

$$(5) (\text{podprostor } \mathbb{F}_q^n, \subseteq) = L_n(q).$$

$$x \geq y \leftrightarrow y \leq x$$

$$x < y \leftrightarrow x \leq y, x \neq y$$

$$x < \cdot y \leftrightarrow x < y, \nexists z : x < z < y$$

x predhodnik y , y predhodnik x

$$(\mathbb{N}, \leq): i < \cdot i + 1$$

$$B_n: A \subset \cdot A \cup \{i\}; i \notin A$$

$$D: r \mid \cdot s \leftrightarrow \frac{s}{r} \text{ praštevílo}$$

$$L_n(q): U < \cdot V \leftrightarrow U \subseteq V, \dim V - \dim U = 1$$

\mathbb{R} : nikoli ne velja $x < \cdot y$.

Hassejev diagram:

graf,

$$V = P,$$

$$xy \in E \iff x < \cdot y \text{ ali } y < \cdot x$$

$$x < \cdot y \implies x \text{ pod } y.$$

Hassejev diagram B_n je hiperkocka.

$$x \text{ maksimalen element, če velja } y \geq x \implies y = x \text{ (oz } \nexists y : y > x)$$

$$x \text{ minimalen element, če velja } y \leq x \implies y = x \text{ (oz } \nexists y : y < x).$$

$$P \text{ končna dum} \implies P \text{ ima maksimalen element.}$$

$$x \text{ največji element: } y \leq x \forall y \in P.$$

Nima največjega elementa.

$$x, y \text{ največja} \implies x \leq y, y \leq x \implies x = y.$$

$$\hat{0} : \text{najmanjši element (če } \exists),$$

$$\hat{1} : \text{največji element (če } \exists).$$

P, Q dum.

$$\varphi : P \rightarrow Q \text{ homomorfizem, če } x \leq_P y \implies \varphi(x) \leq_Q \varphi(y).$$

$$\varphi : P \rightarrow Q \text{ izomorfizem, če je bijektiven homomorfizem in je inverz tudi homomorfizem, oz. } \varphi \text{ bijekcija, } x \leq_P y \iff \varphi(x) \leq_Q \varphi(y).$$

Bijektivni homomorfizem, ni izomorfizem.

$$P \cong Q \text{ (} P, Q \text{ izomorfna), če obstaja izomorfizem } \varphi : P \rightarrow Q.$$

$$B_3 \cong D_{30}.$$

P, Q dum.

$$P \times Q \text{ (množica } P \times Q), (x, y) \leq (x', y'), \text{ če } x \leq_P x', y \leq_Q y', x, x' \in P, y, y' \in Q - \text{kartezični produkt.}$$

$$P \sqcup Q = P \times \{0\} \cup Q \times \{1\}.$$

$$P + Q \text{ (množica } P \sqcup Q), x \leq y \text{ če } (x, y \in P, x \leq_P y) \text{ ali } (x, y \in Q, x \leq_Q y) - \text{disjunktna unija.}$$

$$P \oplus Q \text{ (množica } P \sqcup Q), x \leq y \text{ če } (x, y \in P, x \leq_P y) \text{ ali } (x, y \in Q, x \leq_Q y) \text{ ali } (x \in P, y \in Q) - \text{disjunktna vsota.}$$

$$\underline{1} \oplus \cdots \oplus \underline{1} \cong \underline{n}$$

$$\underline{2} \times \cdots \times \underline{2} \cong B_n$$

$$\varphi : \underline{2}^n \rightarrow B_n$$

$$\varphi(\epsilon_1 \dots \epsilon_n) = \{i : \epsilon_i = 2\}$$

$$D_n \cong [\underline{0}, \underline{\alpha_1}] \times \cdots \times [\underline{0}, \underline{\alpha_k}]$$

$n = p_1^{\alpha_1} \dots p_k^{\alpha_k}$, $\alpha_i \geq 1$, delitelji $p_1^{\beta_1} \dots p_k^{\beta_k}$, $0 \leq \beta_i \leq \alpha_i$.

Če je n produkt k različnih praštevil, je $D_n \cong B_k$.

Veriga je podmnožica P , če sta poljubna elementa primerljiva ($x \leq y$ ali $y \leq x$).

V B_8 : $\{\emptyset, \{1, 5\}, \{1, 2, 5, 7, 8\}\}$,

v D_{12} : $\{2, 6, 12\}$.

$x_0 < x_1 < \dots < x_k$ veriga dolžine k ,

$x_0 \leq x_1 \leq \dots < x_k$ multiveriga dolžine k .

Antiveriga je podmnožica P , v kateri nobena različna elementa nista primerljiva.

$\binom{[n]}{k}$ antiveriga v B_n ,

\P antiveriga v D .

Stopničasta dum (angl. graded) je P z rangom, t.j.

$\rho : P \rightarrow \mathbb{N}$, če

$$x < y \implies \rho(x) < \rho(y)$$

$$x < \cdot y \implies \rho(y) = \rho(x) + 1.$$

V \mathbb{N} : $\rho = id$,

v B_n : $\rho(A) = |A|$,

v D_n : $\rho(p_1^{\alpha_1} \dots p_k^{\alpha_k}) = \alpha_1 + \dots + \alpha_k$,

ni stopničasta.

Definicija 3.2.1. P je lokalno končna, če je za

$\forall x \leq y : [x, y] := \{z : x \leq z \leq y\}$ končna.

Npr. vsaka končna dum je lokalno končna.

\mathbb{N}, D sta lokalno končni.

3.3 Incidenčna algebra

P lokalno končna dum.

$$Int(P) := \{[x, y] : x \leq y\}$$

$$I(P, K) := \{f : Int(P) \rightarrow K\} \text{ incidenčna algebra.}$$

$$x \leq y : f([x, y]) = f(x, y) \text{ (krajšamo).}$$

$$(f + g)(x, y) = f(x, y) + g(x, y)$$

$$(\lambda f)(x, y) = \lambda \cdot f(x, y)$$

$$(f \cdot g)(x, y) = \sum_{x \leq z \leq y} f(x, z) \cdot g(z, y) - \text{pomembno!}$$

$$\begin{aligned} (f \cdot g) \cdot h(x, y) &= \sum_{x \leq z \leq y} (f \cdot g)(x, z) \cdot h(z, y) \\ &= \sum_{x \leq z \leq y} \left(\sum_{x \leq q \leq z} f(x, q) g(q, z) \right) h(z, y) \\ &= \sum_{x \leq w \leq z \leq y} f(x, w) g(w, z) h(z, y) \\ &= \dots = f \cdot (g \cdot h)(x, y). \end{aligned}$$

(Nekomutativna algebra.)

$$P = \underline{n}.$$

$I(\underline{n}, k) \cong$ algebra zgornje trikotnih matrik nad K .

$$f(i, j) \rightarrow [f(i, j) \text{ če } i \leq j, 0 \text{ sicer}]_{i, j=1}^n$$

$$1 \leq i \leq j \leq n$$

$$(A \cdot B)_{ij} = \sum_{k=1}^n A_{ik} B_{kj} = \sum_{k=i}^j A_{ik} B_{kj}$$

$$\underline{1}(x, y) = \delta_{xy} = \begin{cases} 1 : x = y \\ 0 : x < y \end{cases} \text{ enota za množenje.}$$

$$f : \underline{1}(x, y) = \sum_{x \leq z \leq y} f(x, y) \cdot \underline{1}(z, y) = f(x, y), \text{ ker } \underline{1}(z, y) = 0, \text{ razen za } z = y.$$

$$\underline{1} \cdot f = f.$$

Trditev 3.3.1. $f \in I(\underline{n}, K)$ je obrnljiv $\iff f(x, x) \neq 0$ za $\forall x \in P$.

Dokaz 3.3.2.

(\Rightarrow):

$$\begin{aligned} f \cdot g &= \underline{1} \\ (f \cdot g)(x, x) &= \sum_{x \leq z \leq x} f(x, z) g(z, x) = f(x, y) \cdot g(x, y) \\ &= \underline{1}(x, x) = 1 \\ &\implies f(x, x) \neq 0. \end{aligned}$$

(\Leftarrow):

\exists desni inverz:

$$f \cdot g = \underline{1}$$

$$(f \cdot g)(x, x) = 1 = f(x, x) \cdot g(x, x)$$

$$g(x, x) = \frac{1}{f(x, x)}.$$

Skonstruiramo rekurzivno glede na $||[x, y]||$:

$$|[x, y]| = 1 : \checkmark$$

Imamo $g(x', y')$ za $|[x', y']| < |[x, y]|$

$$g(x, y) = \frac{\sum \dots}{f(x, x)}.$$

Podobno za levi inverz, enaka.

$$\zeta(x, y) = 1 \text{ za } x \leq y$$

$$\zeta^2(x, y) = \sum_{x \leq z \leq y} \zeta(x, z) \zeta(z, y) = |[x, y]|$$

$$\zeta^3(x, y) = \sum_{x \leq w \leq z \leq y} \zeta(x, w) \zeta(w, z) \zeta(z, y) = \text{število multiverig dolžine 3 med } x \text{ in } y$$

$$\zeta^k(x, y) = \text{število multiverig dolžine } k \text{ med } x \text{ in } y.$$

$$(\zeta - 1)(x, y) = \begin{cases} 1 : x < y \\ 0 : x = y \end{cases}$$

$$(\zeta - 1)^2 = |(x, y)| - \text{dolžina odprtega intervala.}$$

$$(\zeta - 1)^k = \text{število (multi?)verig dolžine } k \text{ med } x \text{ in } y = 0 \text{ od nekega } k \text{ naprej.}$$

$$\underline{1} + (\zeta - 1) + (\zeta - 1)^2 + \dots \text{ je dobro definirana (končnost).}$$

$$(1 + (\zeta - 1) + \dots)(x, y) = \text{število verig med } x \text{ in } y.$$

$$(1 + (\zeta - 1) + \dots)(1 - (\zeta - 1)) = 1$$

$$(2 - \zeta)^{-1}(x, y) = \text{število verig med } x \text{ in } y.$$

$$\begin{bmatrix} 1 & 1 & \dots & 1 \\ 0 & 1 & \dots & \\ \vdots & & \ddots & \\ & & & 1 \end{bmatrix}^{-1} = \begin{bmatrix} 1 & 1 & 2 & 4 & \dots & 2^{n-1} \\ & & \vdots & & & \vdots \\ & & & & & 2 \\ & & & & & 1 \\ & & & & & 1 \end{bmatrix}.$$

Število verig med i in j je 2^{j-i-1} za $j \geq i + 1$.

3.4 Möbius funkcija in Möbiusova inverzija

$\mu := \zeta^{-1}$: inverz obstaja, ker je $\zeta(x, x) \neq 0$.

4

$$\zeta \cdot \mu = \underline{1}$$

$$x = y : \zeta(x, x) \cdot \mu(x, x) = 1 \implies \mu(x, x) = 1$$

$$x < y : \sum_{x \leq z \leq y} \zeta(x, z) \cdot \mu(z, y) = 0$$

$$\mu(x, y) = - \sum_{x < z \leq y} \mu(z, y)$$

$$\mu \cdot \zeta = \underline{1}$$

$$\sum_{x \leq z \leq y} \mu(x, z) = 0$$

$$\mu(x, y) = - \sum_{x \leq z < y} \mu(x, z)$$

4:

$$\mu(i, i) = 1$$

$$\mu(i, i+1) = -\mu(i, i) = -1$$

$$\mu(i, i+2) = -\mu(i, i) - \mu(i, i+1) = 0$$

$$\mu(i, i+3) = -\mu(i, i) - \mu(i, i+1) - \mu(i, i+2) = 0$$

$$v \text{ n in } (\mathbb{N}, \leq): \mu(x, y) = \begin{cases} 1 : i = j \\ -1 : j = i + 1 \\ 0 : j - i \geq 2 \end{cases}$$

$$\mu(a, a) = \mu(b, b) = \dots = 1$$

$$\mu(a, b) = \mu(b, c) = \mu(c, e) = \mu(a, d) = \mu(d, e) = -1$$

$$\mu(a, b) = \mu(b, e) = 0$$

$$\mu(a, e) = 1.$$

Izrek 3.4.1 (Möbiusova inverzija). P dum, za $\forall x \in P$ $\{z \in P : z \leq x\}$ je končna ($\implies P$ je lokalno končna.)

$f, g : P \rightarrow K$

$$g(y) = \sum_{x \leq y} f(x) \iff f(y) = \sum_{x \leq y} \mu(x, y)g(x).$$

(Dobro definirano, ker je vsota končna.)

Dokaz 3.4.2.

(\implies):

$$\begin{aligned} \sum_{x \leq y} \mu(x, y)g(x) &= \sum_{x \leq y} \mu(x, y) \sum_{z \leq x} f(z) \\ &= \sum_{z \leq y} \sum_{z \leq x \leq y} \mu(x, y)f(z) = f(y); \\ \text{ker } \sum_{z \leq x \leq y} \mu(x, y) &= \delta_{z, y}. \end{aligned}$$

(\Leftarrow): podobno.

Primer.

$P = \underline{n}$

$$\begin{aligned} g(j) = \sum_{i \leq j} f(i) &\iff f(j) = \sum_{i=1}^j \mu(i, j)g(i) = g(j) - g(j-1) \text{ za } j \geq 2, \\ f(1) &= g(1). \end{aligned}$$

Kako izračunati μ za $B_n, D_n, M_n, L_n(q)$?

Trditev 3.4.3. P, Q lokalno končni $\implies P \times Q$ lokalno končen.

$$\mu_{P \times Q}((x, y), (x', y')) = \mu_P(x, y) \cdot \mu_Q(x', y').$$

Dokaz 3.4.4.

$$\begin{aligned}
& (\zeta_{P \times Q}(\mu_P, \mu_Q))((x, y), (x', y')) \\
&= \sum_{(x, y) \leq (x'', y'') \leq (x', y')} \mu_P(x'', x') \mu_Q(y'', y') \\
&= \sum_{x \leq x'' \leq x'} \sum_{y \leq y'' \leq y'} \mu_P(x'', x') \cdot \mu_Q(y'', y') \\
&= \left(\sum_{x \leq x'' \leq x'} \mu_P(x'', x') \right) \cdot \left(\sum_{y \leq y'' \leq y'} \mu_Q(y'', y') \right) \\
&= \delta_{x, x'} \cdot \delta_{y, y'} \\
&= \delta_{(x, y), (x', y')}.
\end{aligned}$$

Primer.

$$(1) \ B_n = \underline{2} \times \cdots \times \underline{2}$$

$$\mu(S, T) = \mu((\epsilon_1 \dots \epsilon_n), (\varphi_1 \dots \varphi_n)) = \mu_{\underline{2}}(\epsilon_1, \varphi_1) \dots \mu_{\underline{2}}(\epsilon_n, \varphi_n) = (-1)^{|T \setminus S|}$$

$$S \subseteq T$$

$$f, g : 2^{[n]} \rightarrow K$$

$$g(T) = \sum_{S \subseteq T} f(S) \iff f(T) = \sum_{S \subseteq T} (-1)^{|T \setminus S|} g(S): \text{NVI.}$$

$$(2) \ D_n = [0, \alpha_1] \times \cdots \times [0, \alpha_k]$$

$$n = p_1^{\alpha_1} \dots p_k^{\alpha_k}$$

$$\mu(r, s) = \mu((\beta_1 \dots \beta_k), (\gamma_1 \dots \gamma_k))$$

$$= \mu(\beta_1, \gamma_1) \dots \mu(\beta_k, \gamma_k)$$

$$= \begin{cases} (-1)^l : \frac{s}{r} \text{ produkt } l \text{ razliĉnih praštevil} \\ 0 : p^2 \mid \frac{s}{r}, \text{ ppraštevil} \end{cases} = \mu\left(\frac{s}{r}\right)$$

$$r = p_1^{\beta_1} \dots p_k^{\beta_k}$$

$$s = p_1^{\gamma_1} \dots p_k^{\gamma_k}$$

$$0 \leq \beta_i \leq \gamma_i \leq \alpha_i$$

$$r = p_1^{\gamma_1 - \beta_1} \dots p_k^{\gamma_k - \beta_k}$$

$$\mu(n) = \begin{cases} (-1)^k : n \text{ produkt } k \text{ razliĉnih praštevil} \\ 0 : p^2 \mid n \text{ praštevil} \end{cases}$$

$$f, g : \mathbb{N} \setminus \{0\} \rightarrow K$$

$$g(n) = \sum_{d|n} f(d) \iff f(n) = \sum_{d|n} \mu(d, n) g(d) = \sum_{d|n} \mu\left(\frac{n}{d}\right) g(d).$$

P

$$I(P, K) = \{f : \text{Int}(P) \rightarrow K\}$$

$$f \cdot g(x, y) = \sum_{x \leq z \leq y} f(x, z) g(z, y)$$

$\zeta, \mu.$

Izrek 3.4.5.

P dum, $\{y \leq x\}$ končen $\forall x \in P$,

$f, g : P \rightarrow K.$

$$f(x) = \sum_{y \leq x} g(y) \iff g(x) = \sum_{y \leq x} \mu(y, x) f(y).$$

Izrek 3.4.6.

P dum, $\{y \geq x\}$ končen $\forall x \in P$,

$f, g : P \rightarrow K.$

$$f(x) = \sum_{y \geq x} g(y) \iff g(x) = \sum_{y \geq x} \mu(x, y) f(y).$$

$$B_n : \mu(S, T) = (-1)^{|T \setminus S|}$$

$$B_n \cong \underline{2} \times \cdots \times \underline{2}$$

$$\mu_{P \times Q} = \mu_P \cdot \mu_Q$$

$$D_n : \mu(r, s) = \begin{cases} (-1)^k : \frac{s}{r} \text{ produkt } k \text{ razliĉnih praštevil} \\ 0 : p^2 \Big| \frac{s}{r} \end{cases}.$$

3.5 Mreže

Definicija 3.5.1. $x \leq y$:

y zgornja meja za x ,

x spodnja meja za y .

P je mreža (angl. lattice?), če imata poljubna elementa najmanjšo zgornjo mejo in največjo spodnjo mejo.

$x \vee y$ spoj (angl. join),

$x \wedge y$ stik (angl. meet).

$$x \wedge y \leq x, y \leq x \vee y$$

$$x, y \leq z \implies x \vee y \leq z$$

$$z \leq x, y \implies z \leq x \wedge y.$$

Primer.

- 3 zgornje meje za x, y , noben ni \leq od ostalih, ni mreža.
- $\underline{n}, \mathbb{N}$: $i \vee j = \max\{i, j\}$, $i \wedge j = \min\{i, j\}$.
- B_n : $T \vee S = T \cup S$, $T \wedge S = T \cap S$.
- D_n, D : $r \vee s = l(r, s)$, $r \wedge s = D(r, s)$.
- $L_n(q)$: $U \vee V = U + V$, $U \wedge V = U \cap V$.
- Π_n
 $\pi = 135 - 246, \sigma = 123 - 46 - 5$
 $\pi \wedge \sigma = \{\text{neprazni preseki bloka } \pi \text{ in bloka } \sigma\}$
 $\pi \vee \sigma = \{\text{povezane komponente grafa, } V = [n], i \sim j: i \text{ in } j \text{ v istem bloku } \pi \text{ ali } \sigma\}$
 $\pi \vee \sigma = 123456$.

P končna mreža \implies ima največji in najmanjši element.

Največji: spoj vseh elementov = $\hat{1}$,

najmanjši: stik vseh elementov = $\hat{0}$.

$\forall x < y$:

$$\sum_{x \leq z \leq y} \mu(x, z) = 0 \implies \mu(x, y) = -\sum_{x \leq z < y} \mu(x, z) \text{ ali}$$

$$\sum_{x \leq z \leq y} \mu(z, y) = 0 \implies \mu(x, y) = -\sum_{x < z \leq y} \mu(z, y).$$

Izrek 3.5.2.

P končna mreža,

$a \neq \hat{1}$.

$$\mu(\hat{0}, \hat{1}) = -\sum_{x \neq \hat{0}, x \wedge a = \hat{0}} \mu(x, \wedge 1).$$

Opomba. Vedno: $\mu(\hat{0}, \hat{1}) = -\sum_{x \neq \hat{0}} \mu(x, \hat{1})$.

Torej izrek nam omogoča, da $\mu(\hat{0}, \hat{1})$ izračunamo preko vsote z manj členi.

Tipično $a < \cdot \hat{1}$.

Dokaz 3.5.3.

$$\begin{aligned} \sum_{x \wedge a = \hat{0}} &= \sum_{x \in P} \mu(x, \hat{1}) \cdot 1(\hat{0}, x \wedge a) \\ &= \sum_{x \in P} \mu(x, \hat{1}) \sum_{y \leq x \wedge a} \mu(\hat{0}, y) \\ &\stackrel{(*)}{=} \sum_{x \in P} \mu(x, \hat{1}) \sum_{y \leq x, y \leq a} \mu(\hat{0}, y) \\ &= \sum_{y \leq a} \left(\sum_{x \geq y} \mu(x, \hat{1}) \right) \mu(\hat{0}, y) = 0; \end{aligned}$$

$\ker \sum_{x \geq y} \mu(x, \hat{1}) = 1(y, \hat{1}) = 0$, $\ker y \leq a \neq \hat{1}$,

(*) : $y \leq x \wedge a \implies y \leq x \wedge y \leq a$.

Primer.

(a) B_n

$$\mu_n = \mu(0, [n])$$

$$[S, T] \cong B_{|T \setminus S|}$$

$$[\{n\}, [n]] \cong B_{n-1}$$

$$A = [n-1]$$

$$\mu_n = \sum_{T \neq \emptyset, T \cap [n-1] = \emptyset} \mu(T, [n]) = -\mu(\{n\}, [n]) = -\mu_{n-1}$$

$$\implies \mu_n = (-1)^n$$

$$\mu(S, T) = (-1)^{|T \setminus S|}.$$

(b) D_n

$$n = p_1^{\alpha_1} \dots p_k^{\alpha_k}$$

$$a = p_1^{\alpha_1-1} \dots p_k^{\alpha_k}$$

$$\mu(1, n) = -\sum_{d|n, d \neq 1, D(d, a)=1} \mu(d, n) = \begin{cases} 0 : \alpha_1 \geq 2 \text{ (takega } d \text{ ni)} \\ -\mu(p_1, n) : \alpha_1 = 1 \text{ (} d = p_1 \text{)} \end{cases}$$

$-\mu(p_1, n) = -\mu(1, p_2^{\alpha_2} \dots p_n^{\alpha_n})$:
 rekurzivno, $= 0$ če $\alpha_i \geq 2$, $(-1)^k$ sicer.

(c) $L_n(q)$

$$\mu_n = \mu(0, \Pi_q^n)$$

$$[U, V] \cong L_{\dim V - \dim U}(q)$$

$$A = \Pi_q^{n-1} \times \{0\}$$

$$\mu_n = -\sum_{U \neq 0, U \cap A = 0} \mu(U, \Pi_q^n) = -q^{n-1} \mu_{n-1}.$$

Linearna algebra: $\dim(U \cap A) + \dim(U + A) = \dim(U) + \dim(A)$:

$$\dim(A) = n - 1, \dim(U \cap A) = 0, \dim(U) \geq 1, \dim(U + A) \geq 0$$

$$n \geq \dim(U \cap A), \dim(U) + \dim(A) \geq n$$

$\implies \dim(U) = 1, U = \text{Lin}\{u\}$; zadnja komponenta $\neq 0$, BŠS 1.

q^{n-1} : q možnosti za vsako od $n - 1$ preostalih komponent.

$$\mu_n = (-1)^n q^{\binom{n}{2}}$$

$$\mu(U, V) = (-1)^{\dim V - \dim U} q^{\binom{\dim V - \dim U}{2}}.$$

(d) Π_n

$$\mu := \mu(1 - 2 - 3 \dots - n, 123 \dots n)$$

$$\alpha = 12 \dots (n - 1) - n$$

$$\mu_n = -\sum_{\pi \neq 1-2 \dots n, \pi \wedge \alpha = 1-2 \dots n} \mu(\pi, 12 \dots n) = -(n - 1) \mu_{n-1}$$

$$\pi = 1 - 2 - \dots - (i - 1) - (in) - (i + 1) - \dots - (n - 1)$$

$$[\pi, 12 \dots n] \cong \Pi_{n-1}$$

$$\mu_n = (-1)^{n-1} (n - 1)! \text{ (do } \mu_1, \text{ ne } \mu_0)$$

$$[\pi, \sigma] \cong \pi_{\alpha_1} \times \dots \times \pi_{\alpha_k},$$

kjer i -ti blok σ razpade na a_i blokov v π za $i = 1, 2 \dots k$.

$$\pi = 12 - 3 - 4 - 568 - 7$$

$$\sigma = 1247 - 56 - 8 - 3$$

$$a_1 = 3, a_2 = 2, a_3 = 1$$

$$\Pi_3 \times \Pi_2 \times \Pi_1$$

$$\mu(\pi, \sigma) = (-1)^{a_1} (a_1 - 1)! \cdot (-1)^{a_2} (a_2 - 1)! \cdot (-1)^{a_3} (a_3 - 1)!.$$

3.6 Reducirane incidenčne algebre in Dirichletove rodovne funkcije

Primer.

- $\underline{n}, \mathbb{N}$

$$\mu(i, j) = \begin{cases} 1 : i = j \\ -1 : j = i + 1 \\ 0 : j - i > 1 \end{cases} \quad \text{- odvisen od } j - i.$$
- $B_n, B = \cup_{n=0}^{\infty} B_n = \{\text{končne podmnožice } \{1, 2, 3 \dots\}\}$
 $\mu(S, T) = (-1)^{|T \setminus S|}$ - odvisen od $|T \setminus S|$.
- $L_n(q), L_q = \cup_{n=0}^{\infty} L_n(q)$ (dodamo $\times \{0\}^i$ na konce?)
 $\mu(U, T) = (-1)^{\dim V - \dim U} \dots$ - odvisen od $\dim V - \dim U$.
- D_n, D
 $\mu(r, s)$ - odvisen od $\frac{s}{r}$.

Vedno: $\mu(x, y) = \mu(x', y')$, če je $[x, y] \cong [x', y']$.

(Primer zgoraj za $\mathbb{N}, B, L(q)$.)

V D : $[1, 14] \cong [1, 15] \cong B_2$, vendar $\frac{14}{1} \neq \frac{15}{1}$.

Izrek 3.6.1.

P lokalno končna dum.

$$I_{\cong}(P, K) = \{f : \text{Int}(P) \rightarrow K : [x, y] \cong [x', y'] \implies f(x, y) = f(x', y')\}.$$

(npr. za $P = \underline{n}$ zgornje trikotne matrike, ki so konstantne na diagonalah(ah?))

$(1, \mu, \zeta)$.

Potem velja $f, g \in I_{\cong}(P, I), \lambda \in K \implies f + g, \lambda \cdot f, f \cdot g \in I_{\cong}(P, K)$,

$f \in I_{\cong}(P, K)$ obrnljiv $\implies f^{-1} \in I_{\cong}(P, K)$,

$I_{\cong}(P, K)$ reducirana incidenčna algebra.

Dokaz 3.6.2.

$$[x, y] \cong [x', y']$$

$$(f + g)(x, y) = f(x, y) + g(x, y) = f(x', y') + g(x', y') = (f + g)(x', y'),$$

$\lambda \cdot f$: podobno.

$$(f \cdot g)(x, y) = \sum_{x \leq z \leq y} f(x, z) \cdot g(z, y)$$

$$(f \cdot g)(x', y') = \sum_{x' \leq z' \leq y'} f(x', z') \cdot g(z', y')$$

$\phi : [x, y] \rightarrow [x', y']$ izomorfizem

$$[\phi(z), \phi(w)] \cong [z, w]$$

$$f(x, z) = f(x', z'), g(z, y) = g(z', y')$$

$$f^{-1}(x, y) = f^{-1}(x', y') \text{ z indukcijo po } |[x, y]|.$$

$$|[x, y]| = 1$$

$$x = x', y = y'$$

$$f^{-1}(x, y) = \frac{1}{f(x, y)} = \frac{1}{f(x', y')} = f^{-1}(x', y')$$

$$|[x, y]| > 1$$

$$\sum_{x \leq z \leq y} f(x, z) f^{-1}(z, y) = \sum_{x < z \leq y} f(x, z) f^{-1}(z, y) + f(x, x) f^{-1}(x, y) = 0$$

$$\sum_{x' \leq z' \leq y'} f(x', z') f^{-1}(z', y') = \sum_{x' < z' \leq y'} f(x', z') f^{-1}(z', y') + f(x', x') f^{-1}(x', y') = 0;$$

$$f(x, z) = f(x', z'), f(x, x) = f(x', x'), f^{-1}(z, y) \stackrel{IP}{=} f^{-1}(z', y')$$

$$\implies f^{-1}(x, y) = f^{-1}(x', y').$$

$\tau = \{\text{množica ekvivalenčnih razredov za } \cong\}$: množica tipov.

$$\mathbb{N} : \tau \equiv \mathbb{N}$$

$$B : \tau \equiv \mathbb{N}$$

$$L(q) : \tau \equiv \mathbb{N}$$

$[x, y]$ tipa α .

$$f, g \in I_{\cong}(P, K), f \cdot g(x, y) = \sum_{x \leq z \leq y} f(x, z) g(z, y)$$

$$(f \cdot g)(\alpha) = \sum_{\beta, \gamma} \binom{\alpha}{\beta, \gamma} f(\beta) g(\gamma)$$

$(f \cdot g)$ odvisen samo od tipa.

$\binom{\alpha}{\beta, \gamma} := \text{število elementov } z \in [x, y]; [x, y] \text{ tipa } \alpha, \text{ da je } [x, z] \text{ tipa } \beta, [z, y] \text{ tipa } \gamma.$

γ .

Torej: $I_{\cong}(P, K)$ je izomorfna algebri preslikav $\tau \rightarrow K$ s produktom

$$(f \cdot g)(\alpha) = \sum_{\beta, \gamma} \binom{\alpha}{\beta, \gamma} f(\beta) g(\gamma).$$

\mathbb{N}

$$\binom{n}{i, j} = \begin{cases} 1 : i + j = n \\ 0 : \text{ sicer} \end{cases}$$

$$f \cdot g(n) = \sum_{k=0}^n f(k) g(n - k)$$

$$I_{\cong}(\mathbb{N}, K) \cong K[[x]]$$

$$f \rightarrow \sum_n f(n) x^n$$

B

$$\binom{n}{i, j} = \begin{cases} \binom{n}{i} : i + j = n \\ 0 : \text{ sicer} \end{cases}$$

$$f \cdot g(n) = \sum_{k=0}^n \binom{n}{k} f(k) g(n - k)$$

$$I_{\cong}(B, K) \cong K[[x]]$$

$$f \rightarrow \sum_n \frac{f(n)}{n!} x^n$$

L_q

$$\binom{n}{i, j} = \begin{cases} \binom{n}{i}_q : i + j = n \\ 0 : \text{ sicer} \end{cases}$$

$$f \cdot g(n) = \sum_{k=0}^n \binom{n}{k}_q f(k) g(n - k)$$

$$I_{\cong}(L(q), K) \cong K[[x]]$$

$$f \rightarrow \sum_n \frac{f(n)}{\underline{n}!} x^n$$

\mathbb{N}

$$\zeta \rightarrow \frac{1}{1-x}$$

$$\mu \rightarrow \left(\frac{1}{1-x} \right)^{-1} = 1 - x, \text{ torej } \mu(0) = 1, \mu(1) = -1, \mu(2) = \mu(3) = \dots = 0$$

$$\zeta^k \rightarrow \left(\frac{1}{1-x} \right)^k = \sum_n \binom{n+k-1}{k-1} x^n$$

$\zeta^k(n)$: število multiverig dolžine k med 0 in n

$$0 \leq i_1 \leq \dots \leq i_{k-1} \leq n.$$

Kombinacije s ponavljanjem: $\binom{(n+1)+(k-1)-1}{k-1} = \binom{n+k-1}{k-1}$

$$(\zeta - 1)^k \rightarrow \left(\frac{x}{1-x}\right)^k = \sum_k \binom{n-1}{k-1} x^n$$

$$0 < i_1 < \dots < i_{k-1} < n$$

$$\binom{n-1}{k-1}$$

$$(2 - \zeta)^{-1} \rightarrow \left(2 - \frac{1}{1-x}\right)^{-1} = \left(\frac{2-2x-1}{1-x}\right)^{-1} = \frac{1-x}{1-2x} = 1 + \sum_{n=1}^{\infty} 2^{n-1} x^n$$

$(2 - \zeta)^{-1}(n)$: število vseh verig med 0 in n :

$$0 < i_1 < \dots < i_{k-1} < n$$

2^{n-1} , $n \geq 1$: izberem ali ne.

B

$$\zeta \rightarrow e^x$$

$$\mu \rightarrow e^{-x}, \text{ torej } \mu(n) = (-1)^n$$

$$\zeta^k \rightarrow e^{kx} = \sum_n \frac{k^n}{n!} x^n$$

$\zeta^k(n)$: število multiverig $\emptyset \subseteq A_1 \subseteq \dots \subseteq A_{k-1} \subseteq [n]$.

Za $\forall j = 1, 2 \dots n$ izberemo A_i , v katerem se j prvič pojavi; k izbir, n -krat izbiramo $\rightarrow k^n$

$$(\zeta - 1)^k \rightarrow (e^x - 1)^k = \sum_n \frac{k! S(n, k)}{n!} x^n$$

$(\zeta - 1)^k(n)$: število verig $\emptyset \subseteq A_1 \subset \dots \subset A_{k-1} \subseteq [n]$

$(A_1, A_2 \setminus A_1, A_3 \setminus A_2 \dots)$ urejena razdelitev na k blokov.

Spomnimo se: $\mu(r, s) = \mu(r', s')$, če je $\frac{s}{r} = \frac{s'}{r'}$.

$[r, s] \sim [r', s']$, če je $\frac{s}{r} = \frac{s'}{r'}$.

$I_{\sim}(D, K) = \{f : Int(D) \rightarrow K : [r, s] \sim [r', s'] \implies f(r, s) = f(r', s')\}$ je

tudi podlagebra (dokaz podoben).

$$\tau \equiv \mathbb{N} \setminus \{0\}$$

$$\binom{n}{i,j} = \begin{cases} 1 : i \cdot j = n \\ 0 : \text{sicer} \end{cases}$$

$$f * g(n) = \sum_{i,j} \binom{n}{i,j} f(i)g(j) = \sum_{d|n} f(d)g\left(\frac{n}{d}\right) \text{ Dirichletova konvolucija.}$$

Dirichletove rodovne funkcije:

$$\left\{ \sum_{n=1}^{\infty} \frac{a_n}{n^s}; a_i \in K \right\}$$

$$\sum_{n=1}^{\infty} \frac{a_n}{n^s} \cdot \sum_{n=1}^{\infty} \frac{b_n}{n^s} = \sum_{n=1}^{\infty} \frac{\sum_{d|n} a_d b_{\frac{n}{d}}}{n^s}$$

$$f \rightarrow \sum_n \frac{f(n)}{n^s} \text{ izomorfizem algeber.}$$

$$\zeta \rightarrow \zeta(s) \text{ (Riemmanova) funkcija } \zeta.$$

če $\sum_n \frac{a_n}{n^s}$ in $\sum_n \frac{b_n}{n^s}$ konvergirata:

$$\left(\frac{a_1}{1^s} + \frac{a_2}{2^s} + \frac{a_3}{3^s} + \dots \right) \cdot \left(\frac{b_1}{1^s} + \frac{b_2}{2^s} + \frac{b_3}{3^s} + \dots \right)$$

$$\left[\frac{1}{6^s} \right] : a_1 b_6 + a_2 b_3 + a_3 b_2 + a_6 b_1 \text{ (množenje kot dejanske funkcije).}$$