

**Verjetnostne metode v
računalništvu - zapiski s
predavanj prof. Marca**

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Kazalo

1	Introduction	1
1.1	Probability	1
1.2	Random variables	2
2	Quicksort, min-cut	4
2.1	Quicksort	4
	Literatura	7

Poglavje 1

Introduction

1.1 Probability

(Ω, F, P_r) :

- $\emptyset \in F$,
- $A \in F \implies A^c \in F$,
- $A_1, A_2 \dots \in F \implies \cup_{i=1}^{\infty} A_i \in F$.

$P_r(A) \geq 0$,

$P_r(\cup_{i=1}^{\infty} A_i) = \sum_{i=1}^{\infty} P_r(A_i)$ if A_i disjoint,

$P_r(\cup_{i=1}^{\infty} A_i) \leq \sum_{i=1}^{\infty} P_r(A_i)$,

$\Omega = \{\omega_1, \omega_2 \dots\}$ - countable case.

$$\begin{pmatrix} \omega_1 & \omega_2 & \dots \\ p_1 & p_2 & \dots \end{pmatrix}$$

Primer.

`Alg():`

`while True:`

`B = sample as random from {0,1} # 1 with probability p`

`if B = 1:`

return

$$\Omega = \{1, 01, 001, 0001 \dots\}$$

$$\begin{pmatrix} 1 & 01 & 001 & 0001 & \dots \\ p & (1-p)p & (1-p)^2p & (1-p)^3p & \dots \end{pmatrix}.$$

1.2 Random variables

$X : \Omega \rightarrow \mathbb{Z}$.

$E[X] = \sum_{c \in \mathbb{Z}} c \cdot P_r(X = c)$ expected value of X .

Properties:

- $E[f(X)] = \sum_{c \in \mathbb{Z}} f(c) \cdot P_r(X = c)$,
- $E[aX + bY] = aE[X] + bE[Y]$,
- $E[X \cdot Y] = E[X] \cdot E[Y]$ if X, Y independent,
- $P_r(X \geq a) \leq \frac{E[X]}{a} \forall a > 0, X \geq 0$ Markov inequality.

Primer. (Continuing from before).

X = number of trials before return.

$X : \Omega \rightarrow \mathbb{Z}$.

$X : 1 \rightarrow 1, 01 \rightarrow 2, 001 \rightarrow 3 \dots$

$$\begin{pmatrix} 1 & 2 & 3 & 4 & \dots \\ p & (1-p)p & (1-p)^2p & (1-p)^3p & \dots \end{pmatrix} - \text{geometric distribution.}$$

Trditev 1.2.1. $E[X] = \frac{1}{p}$.

Dokaz 1.2.2. $X = \sum_{i=1}^{\infty} X_i$.

$$X_i = \begin{cases} 1 & \text{if trial } i \text{ is executed} \\ 0 & \text{else} \end{cases}$$

$$\begin{aligned} E[X] &= E\left[\sum_{i=1}^{\infty} X_i\right] = \sum_{i=1}^{\infty} E[X_i] = \\ &= \sum_{i=1}^{\infty} (1-p)^{i-1} = \frac{i=0}{\infty} (1-p)^i = \frac{1}{1-(1-p)} = \frac{1}{p}. \end{aligned}$$

$$E[X] = \frac{1}{p}.$$

$$P_r(X \geq 100 \cdot \frac{1}{p}) \leq \frac{E[X]}{\frac{1}{p}} = \frac{1}{100}.$$

Definicija 1.2.3. $H_n = 1 + \frac{1}{2} + \frac{1}{3} + \dots + \frac{1}{n} = \sum_{i=1}^{\infty} \frac{1}{i}$.

Izrek 1.2.4. $H_n \leq 1 + \ln(n)$.

Dokaz 1.2.5.

$$H_n = 1 + \sum_{i=2}^n \frac{1}{i} \stackrel{\text{integral}}{\leq} 1 + \int_1^n \frac{dx}{x} = 1 + \ln(x)|_1^n = 1 + \ln(n)$$

Poglavje 2

Quicksort, min-cut

2.1 Quicksort

Input: set (no equal element) (unordered list) $S \in \mathbb{R}$
(or whatever you can compare linearly)

Output: ordered list

Code:

```
def Quicksort(S):  
    if |S| = 0 or 1:  
        return S  
    else:  
        a = uniformly at random from S  
         $S^- = \{b \in S \mid b < a\}$   
         $S^+ = \{b \in S \mid a < b\}$   
        return Quicksort( $S^-$ ), a, Quicksort( $S^+$ )
```

$C(n)$ - random variable, the number of comparisons in evaluation of Quicksort
with $|S| = n$

Izrek 2.1.1. $E[C(n)] = O(N \log(n))$

Dokaz 2.1.2. $C(0) = C(1) = 0$

$$\begin{aligned}
E[C(n)] &= n - 1 + \sum_{i=1}^n (E[C(i-1)] + E[C(n-i)]) \cdot P_r(a \text{ is } i\text{-it element}) \leq \\
&\leq n + \frac{2}{n} \sum_{i=1}^{n-1} E[C(i)]
\end{aligned}$$

Induction:

$n = 1 : \checkmark$

$n - 1 \rightarrow n$:

$$\begin{aligned}
E[C(n)] &\leq n + \frac{2}{n} \sum_{i=1}^n E[C(i)] \leq \\
&\leq n + \frac{2}{n} \sum_{i=1}^n 5i \log i \leq \\
&\leq n + \frac{2}{n} \sum_{i=1}^{\lfloor \frac{n}{2} \rfloor} 5i \log i + \frac{2}{n} \sum_{i=1+\lfloor \frac{n}{2} \rfloor}^{n-1} 5i \log i \leq \\
&\leq n + \frac{2}{n} \sum_{i=1}^{\lfloor \frac{n}{2} \rfloor} 5i \log \frac{n}{2} + \frac{2}{n} \sum_{i=1+\lfloor \frac{n}{2} \rfloor}^{n-1} 5i \log n \leq \\
&(\log \frac{n}{2} = \log n - 1) \\
&\leq n + \frac{2}{n} \left(\sum_{i=1}^n 5i \log n - \sum_{i=1}^{\frac{n}{2}} 5i \right) = \\
&= n + \frac{10}{n} \left(\frac{n(n-1)}{2} \log n - \frac{\frac{n}{2}(\frac{n}{2}+1)}{2} \right) \leq \\
&\leq n + 5(n-1) \log n - n < \\
&< 5n \log n
\end{aligned}$$

$$P(C(n) \geq b \cdot 5n \log n) \stackrel{\text{Markov}}{\leq} \frac{1}{b}$$

Dokaz 2.1.3.

2:

Let $S_1, S_2 \dots S_n$ sorted elements of S

Define random variable $X_{ij} = \begin{cases} 1 & \text{if } S_i \text{ and } S_j \text{ are compared} \\ 0 & \text{else} \end{cases}$

$$C(n) = \sum_{1 \leq i < j \leq n} E[X_{ij}]$$

$$E[X_{ij}] = P(S_i \text{ and } X_j \text{ compared})$$

S_{ij} - the last set including S_i and S_j

$$E[X_{ij}] = \frac{2}{|S_{ij}|} \leq \frac{2}{j-i+1}$$

$$|S_{ij}| \geq j - i + 1$$

S_{ij} has everything in between

$$\begin{aligned} \Rightarrow E[C(n)] &\leq \sum_{1 \leq i < j \leq n} \frac{2}{j-i+1} = \\ &\stackrel{k=j-i+1}{=} \sum_{i=1}^{n-1} \sum_{k=2}^{n-i+1} \frac{2}{k} \leq \\ &\leq 2 \cdot n \cdot H_n \leq \\ &\leq 2n(1 + \log n) \end{aligned}$$

Literatura