

RZHP 1. domača naloga

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Naloga 1

a)

$$n(\log_4(m+1) * n + 2 * 1) = n^2 * \log_4(m+1) + 2n \Rightarrow \theta(n^2 * \log_4(m+1)) \quad (1)$$

b)

Najprej n definiramo kot razdaljo od e (end) in s (start)

$$n = j = e - s \quad (2)$$

Potem definiramo novi meji pri rekurzivnih klicih

$$n = (end) - (start) = (s + floor(j/6)) - (s) = floor(j/6) = \lfloor \frac{j}{6} \rfloor = \lfloor \frac{n}{6} \rfloor \quad (3)$$

$$n = (end) - (start) = (e) - (s + floor(j/6) + 1) = (e - s) - \lfloor \frac{j}{6} \rfloor + 1 = n - \lfloor \frac{n}{6} \rfloor + 1 = \lceil \frac{5n}{6} \rceil + 1 \quad (4)$$

+1, ki ostane, ostani zaradi tega, ker je med vključno dvema številoma (n_1, n_2): $n_2 - n_1 + 1$ števil. Poleg dveh rekurzivnih klicov moramo v vsakem klicu funkcije skozi for zanko veliko $j = n$. Dobimo končno enačbo:

$$T(n) = n + T(\lfloor \frac{n}{6} \rfloor) + T(\lceil \frac{5n}{6} \rceil) \quad (5)$$

Sedaj s pomočjo Akra Bazzi teorema dobimo tight bound

$$T(n) = T(\lceil \frac{5n}{6} \rceil) + T(\lfloor \frac{n}{6} \rfloor) + n =$$

$$T(\frac{5n}{6} + \lceil \frac{5n}{6} \rceil - \frac{5n}{6}) + T(\frac{n}{6} + \lfloor \frac{n}{6} \rfloor - \frac{n}{6}) + n \quad (6)$$

$$a_1 = 1 \quad b_1 = \frac{5}{6} \quad h_1 = \lceil \frac{5n}{6} \rceil - \frac{5n}{6} = O(\frac{n}{\log^2 n})$$

$$a_2 = 1 \quad b_2 = \frac{5}{6} \quad h_2 = \lfloor \frac{n}{6} \rfloor - \frac{n}{6} = O(\frac{n}{\log^2 n}) \quad (7)$$

Dokaza za obe zgornji meji

$$\lceil \frac{5n}{6} \rceil - \frac{5n}{6} \leq c \frac{n}{\log^2 n} \Rightarrow 1 \leq c \frac{n}{\log^2 n}$$

$$\lfloor \frac{n}{6} \rfloor - \frac{n}{6} \leq c \frac{n}{\log^2 n} \Rightarrow 0 \leq c \frac{n}{\log^2 n} \quad (8)$$

Izračunamo p

$$1 * (\frac{5}{6})^p + 1 * (\frac{1}{6})^p = 1 \Rightarrow p = 1 \quad (9)$$

Končno

$$T(n) = \theta(n^p(1 + \int_1^n \frac{f(u)}{u^{p+1}} du)) = \theta(n^1(1 + \int_1^n \frac{u}{u^2} du)) =$$

$$\theta(n(1 + \int_1^n \frac{1}{u} du)) = \theta(n(1 + \ln u|_1^n)) = \theta(n(1 + \ln n - \ln 1)) = \quad (10)$$

$$\theta(n \log n)$$

c)

Funkcijo lahko zapišemo v obliki

$$T(n) = T(n-1) + T(n-1) + 1 = 2T(n-1) + 1, \quad T(0) = 1 \quad (11)$$

$$T(n) = 1 + 2 + 4 + \dots + 2^n = \sum_{i=0}^n 2^i = 2^{n+1} \Rightarrow \theta(2^n) \quad (12)$$

Naloga 2

a)

$$T(n) = 4T\left(\frac{n}{9}\right) + \sqrt[2]{n}$$

$$a = 4 \quad b = 9 \quad d = \frac{1}{2} \quad (13)$$

Uporabimo Master theorem

$$a > b^d \Rightarrow 4 > 9^{\frac{1}{2}} = 3 \Rightarrow T(n) = \theta(n^{\log_9 4}) \quad (14)$$

b)

$$T(n) = 5T\left(\frac{n}{4}\right) + \log(n) \quad (15)$$

$$a = 5 \quad b = 4 \quad f(n) = \log(n) \quad (16)$$

Uporabimo Master theorem

$$f(n) = O(n^{\log_4 5 - \epsilon}) \Rightarrow T(n) = \theta(n^{\log_4 5}) \quad (17)$$

Naloga 3

a)

$$T(n) = 2T\left(\frac{n}{4}\right) + 3T\left(\frac{n}{6}\right) + \theta(n \log n)$$

$$f(n) = n \log n \quad k = 2 \quad a_1 = 2 \quad b_1 = \frac{1}{4} \quad (18)$$

$$a_2 = 3 \quad b_2 = \frac{1}{6}$$

$$\sum_{i=1}^k a_i b_i^p = 1 \Rightarrow 2\left(\frac{1}{4}\right)^p + 3\left(\frac{1}{6}\right)^p = 1 \Rightarrow p = 1 \quad (19)$$

$$T(n) = \theta(n^p(1 + \int_1^n \frac{f(u)}{u^{p+1}} du)) = \theta(n^1(1 + \int_1^n \frac{u \log u}{u^2} du)) = \theta(n(1 + \int_1^n \frac{\log u}{u} du)) \quad (20)$$

$$x = \log u \Rightarrow dx = \frac{1}{u} du \Rightarrow du = u dx \quad (21)$$

$$\begin{aligned} T(n) &= \theta(n(1 + \int_1^n \frac{\log u}{u} du)) = \theta(n(1 + \int_1^n \frac{x}{u} u dx)) = \\ \theta(n(1 + \frac{x^2}{2} \Big|_1^n)) &= \theta(n(1 + \frac{\log^2 u}{2} \Big|_1^n)) = \theta(n(1 + \frac{\log^2 n}{2} - \frac{\log^2 1}{2})) = \\ &\Rightarrow \theta(n \log^2 n) \end{aligned} \quad (22)$$

b)

$$\begin{aligned} T(n) &= T(\frac{n}{9}) + T(\frac{n}{4}) + 2T(\frac{n}{36}) + \sqrt{n^3} \\ f(n) &= \sqrt{n^3} \quad a_1 = 1 \quad b_1 = \frac{1}{9} \\ &\quad a_2 = 1 \quad b_2 = \frac{1}{4} \\ &\quad a_3 = 2 \quad b_3 = \frac{1}{36} \end{aligned} \quad (23)$$

Za izračun vrednosti spremenljivke p sem uporabil wolfram alpha

$$\sum_{i=1}^k a_i b_i^p = 1 \Rightarrow 1(\frac{1}{9})^p + 1(\frac{1}{4})^p + 2(\frac{1}{36})^p = 1 \Rightarrow p \approx 0.5695215 \quad (24)$$

$$\begin{aligned} T(n) &= \theta(n^p(1 + \int_1^n \frac{f(u)}{u^{p+1}} du)) = \theta(n^p(1 + \int_1^n \frac{\sqrt{u^3}}{u^{p+1}} du)) = \\ \theta(n^p(1 + \int_1^n u^{\frac{3}{2}-(p+1)} du)) &= \theta(n^p(1 + \int_1^n u^{-0.0695215} du)) = \\ \theta(n^p(1 + \frac{u^{0.9304785}}{0.9304785} \Big|_1^n)) &= \theta(n^p(1 + \frac{n^{0.9304785}}{0.9304785} - \frac{1^{0.9304785}}{0.9304785})) = \\ \theta(n^{0.5695215} \frac{n^{0.9304785}}{0.9304785}) &= \theta(n^{1.5}) = \theta(\sqrt{n^3}) \end{aligned} \quad (25)$$

Naloga 4

a)

Seštejemo vsak nivo posebej, zgornja meja je podana z dolžino rekurzije, kjer se pri klicu deli z 2, saj je tam neuravnoteženo drevo najdaljše

$$T(n) = n^3 + 2n^3 + 4n^3 + \dots + 2^{\log_2 n} n^3 = \sum_{i=0}^{\log_2 n} 2^i n^3 = n^3 \sum_{i=0}^{\log_2 n} 2^i = n^4 \quad (26)$$

d)

$$\begin{aligned} T(n) &= 8T(\lfloor \frac{n}{2} \rfloor) + 27T(\lfloor \frac{n}{3} \rfloor) + n^3 = \\ &8T(\frac{n}{2} + \lfloor \frac{n}{2} \rfloor - \frac{n}{2}) + 27T(\frac{n}{3} + \lfloor \frac{n}{3} \rfloor - \frac{n}{3}) + n^3 \\ a_1 &= 8 \quad b_1 = \frac{1}{2} \quad h_1 = \lfloor \frac{n}{2} \rfloor - \frac{n}{2} = O(\frac{n}{\log^2 n}) \\ a_2 &= 27 \quad b_2 = \frac{1}{3} \quad h_2 = \lfloor \frac{n}{3} \rfloor - \frac{n}{3} = O(\frac{n}{\log^2 n}) \end{aligned} \quad (27)$$

Dokaza za obe zgornji meji

$$\begin{aligned} \lfloor \frac{n}{2} \rfloor - \frac{n}{2} &\leq c \frac{n}{\log^2 n} \Rightarrow 0 \leq c \frac{n}{\log^2 n} \\ \lfloor \frac{n}{3} \rfloor - \frac{n}{3} &\leq c \frac{n}{\log^2 n} \Rightarrow 0 \leq c \frac{n}{\log^2 n} \end{aligned} \quad (28)$$

$$\sum_{i=1}^k a_i b_i^p = 1 \Rightarrow 8(\frac{1}{2})^p + 27(\frac{1}{3})^p = 1 \Rightarrow p = 3 \quad (29)$$

$$\begin{aligned} T(n) &= \theta(n^p(1 + \int_1^n \frac{f(u)}{u^{p+1}} du)) = \theta(n^3(1 + \int_1^n \frac{u^3}{u^{3+1}} du)) = \\ &\theta(n^3(1 + \int_1^n u^{-1} du)) = \theta(n^3(1 + \ln u|_1^n)) = \theta(n^3(1 + \ln n - \ln 1)) \\ &\Rightarrow \theta(n^3 \log n) \end{aligned} \quad (30)$$

Naloga 5

a)

$$\begin{aligned}
T(n) &= T(n-1) + n^2 + n - 1 \\
T(n+1) - T(n) &= (n+1)^2 + (n+1) - 1 \\
(E-1)T(n) &= n^2 + 2n + 1 + n \\
(E-1)T(n) &= n^2 + 3n + 1
\end{aligned} \tag{31}$$

polinom na desni ustreza funkciji $(a_0 + a_1n + a_2n^2)a^n$,
kjer je $a = 1$, anihilira pa jo $(E-1)^3 \Rightarrow (E-1)^3(E-1)T(n) = 0$

$$\begin{aligned}
(E-1)^4T(n) &= 0 \quad a = 1 \quad d = 0 \\
T(n) &= \left(\sum_{i=0}^{d-1} \alpha_i n^i \right) a^n = \alpha_0 + \alpha_1 n + \alpha_2 n^2 + \alpha_3 n^3 \\
T(0) &= 1 = \alpha_0 + \alpha_1 0 + \alpha_2 0 + \alpha_3 0 \Rightarrow \alpha_0 = 1 \\
T(1) &= T(0) + n^2 + n - 1 = 1 + 1^2 + 1 - 1 = 2 \\
T(2) &= T(1) + n^2 + n - 1 = 2 + 2^2 + 2 - 1 = 7 \\
T(3) &= T(2) + n^2 + n - 1 = 7 + 3^2 + 3 - 1 = 18
\end{aligned} \tag{32}$$

...

$$\begin{aligned}
\alpha_0 &= 1 \quad \alpha_1 = -\frac{1}{3} \quad \alpha_2 = 1 \quad \alpha_3 = \frac{1}{3} \\
T(n) &= \alpha_0 + \alpha_1 n + \alpha_2 n^2 + \alpha_3 n^3 = 1 - \frac{n}{3} + n^2 + \frac{n^3}{3}
\end{aligned} \tag{33}$$