The University of Texas at Austin Department of Electrical and Computer Engineering

EE381K: Large Scale Optimization — Fall 2015

PROBLEM SET SEVEN

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Due: Thursday, November 5, 2015.

Reading Assignments

1. Reading: Boyd & Vandenberghe: Chapters 7 & 8.

Matlab and Computational Assignments. Please provide a printout of the Matlab code you wrote to generate the solutions to the problems below.

- 1. Sums of Squares and SDP. Construct a 4th (or higher) degree non-convex univariate polynomial that is bounded from below (just make sure the leading term is even degree, with positive coefficient), and show that there is a starting point from which gradient descent will get trapped in a local minimum. Then formulate the problem as a sum of squares, and solve using SDP via CVX.
- 2. (?) Try solving the above SDP via an interior point method. You can try a generic approach as discussed in class, or something more carefully tailored to SDPs.
- 3. **MaxCut**. The MAXCUT problem is as follows: Given a graph G = (V, E) with nonnegative edge weights $W = \{w_{ij}\}$, find a partition of the vertices $V = V_1 \cup V_2$ (where $V_1 \cap V_2 = \emptyset$) so that the sum of the edge weights from V_1 to V_2 is maximized.

This can be formulated as a binary integer programming problem as follows:

$$\max: \sum_{i,j} \frac{w_{ij}}{2} (1 - x_i x_j)$$
s.t.: $x_i \in \{-1, 1\}.$

Written problem 1 below asks you to show that this is equivalent to the following rank-constrained SDP:

max:
$$\sum_{i,j} \frac{w_{ij}}{2} (1 - X_{ij})$$
s.t.:
$$X_{ii} = 1$$

$$X \succeq 0$$

$$\operatorname{rank}(X) = 1.$$

This is still a non-convex problem, because of the rank constraint. The SDP relaxation of this comes from dropping the rank one constraint, and hence obtaining:

$$\max: \sum_{i,j} \frac{w_{ij}}{2} (1 - X_{ij})$$
s.t.:
$$X_{ii} = 1$$

$$X \succeq 0$$

which is a convex optimization problem. Use CVX to formulate and solve MaxCut on:

- (a) Any two planar graphs of your choice (with a reasonable number of nodes).
- (b) The Petersen Graph: http://en.wikipedia.org/wiki/Petersen_graph.

Written Problems

- 1. Show that the rank-constrained SDP is equivalent to the binary variable formulation of Max-Cut.
- Find the dual of the following convex quadratically constrained quadratic program, and show that the dual objective is concave and hence the problem is convex (the dual is always convex I am asking you to check this explicitly for this particular case). In the problem below, the matrices P₀ and P_i are positive definite.

minimize
$$\frac{1}{2}\mathbf{x}^{\top}P_0\mathbf{x} + \mathbf{p}_0^{\top}\mathbf{x} + r_0$$

subject to
$$\frac{1}{2}\mathbf{x}^{\top}P_i\mathbf{x} + \mathbf{p}_i^{\top}\mathbf{x} + r_i \le 0, i = 1, \dots, m.$$
 (1)

- 3. Problem 7.12 form Boyd & Vandenberghe.
- 4. Problem 7.13 from Boyd & Vandenberghe.
- 5. Problem 8.8 from Boyd & Vandenberghe.
- 6. Problem 8.9 from Boyd & Vandenberghe.
- 7. Problem 8.23 from Boyd & Vandenberghe.
- 8. Problem 8.24 from Boyd & Vandenberghe.
- 9. Problem 8.25 from Boyd & Vandenberghe.
- 10. (?) 8.2 from Boyd & Vandenberghe.
- 11. (?) 8.13 from Boyd & Vandenberghe.