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Fall 2013, CSE 383, Numerical Linear Algebra, Midterm
 Tuesday, October 31,
 This is a closed-book exam.

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1. [10 points,] Let $A \in \mathbb{R}^{m \times m}$ be unitary and upper triangular. Using induction with base case $m = 1$, show that A is diagonal with nonzero entries in the diagonal.

A unitary
 A upper triangular

Base case $k=1$ $A = [a_{11}] = [\pm 1]$

If true for $k=m-1$ then $A = \begin{bmatrix} a_{11} & & \\ & \ddots & \\ & & a_{m-1,m-1} \end{bmatrix}$ where
 $a_{11} \dots a_{m-1,m-1}$
 $a_{ii} = \pm 1$

✓

$k=m$ $A = \begin{bmatrix} \text{unitary and diagonal} & & \\ & A_{(m-1) \times (m-1)} & \\ a_{m1} & \dots & a_{mm} \end{bmatrix}$

$\Rightarrow a_{1m} = \dots = a_{m(m-1)} = 0$
 $a_{m1} = \dots = a_{(m-1)m} = 0$
 $a_{mm} = \pm 1$
 (b.c. $a_i^* a_m = a_2^* a_m = \dots = a_{m-1}^* a_m = 0$)

2. [5 points,] Let $A \in \mathbb{R}^{m \times m}$ and let $\|A\|$ be a vector-induced matrix norm. Let I be the $m \times m$ identity matrix. Show that if $\|A\| < 1$ then $I - A$ is invertible.

$\|A\| = \max_{\|x\|=1} \|Ax\|$

$0 \leq \|A\| < 1$

$\|I - A\| \leq \|I\| + \|A\| = \|I\| + \|A\| = 1 + \|A\| \geq 1$

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Since $\|I - A\| \geq 1$, $I - A$ is invertible

$\|I\| + \|A\| \geq 1$
 but $\|I - A\| \leq \|I\| + \|A\|$

3. [20 points,] Let

$$A = \begin{bmatrix} 1/2 & -1/2 \\ 1/2 & 1/2 \\ 1/2 & -1/2 \\ 1/2 & 1/2 \end{bmatrix} \begin{bmatrix} 2 & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 1/\sqrt{2} & -1/\sqrt{2} & 0 \\ 1/\sqrt{2} & 1/\sqrt{2} & 0 \end{bmatrix}$$

truncated

(a) Give a set of orthonormal vectors that span $\text{range}(A)$; (b) Give a set of orthonormal vectors that span $\text{rowspace}(A)$; (c) give a set of orthonormal vectors that span $\text{null}(A)$; (d) What is $\|A\|_2$?; (e) What is the 2-norm condition number of the linear map defined by $y = Ax : x \in \text{rowspace}(A)$?

(a) $\left\{ \begin{pmatrix} 1/\sqrt{2} \\ -1/\sqrt{2} \\ 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 1/\sqrt{2} \\ 1/\sqrt{2} \\ 0 \\ 0 \end{pmatrix} \right\}$

(b) $\left\{ \begin{pmatrix} 1/2 \\ 1/2 \\ 1/2 \\ 1/2 \end{pmatrix}, \begin{pmatrix} -1/2 \\ 1/2 \\ -1/2 \\ 1/2 \end{pmatrix} \right\}$

(c) $\left\{ \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix} \right\}$

(d) $\|A\|_2 = 2$ (max sing. value)

(e) $\kappa_2(A) = \frac{2}{1} = 2$ ($\frac{\sigma_{\max}}{\sigma_{\min}}$)

4. [5 points,] Let $u \in \mathbb{R}^n$ such that $\|u\|_2 = 1$ be given. Let $H = I - 2uu^T$. Let $w \in \mathbb{R}^n$ be given. Give a formula for finding a vector v such that $Hv = w$. Explain.

$$(I - 2uu^T)v = w$$

$$H = (I - 2uu^T)$$

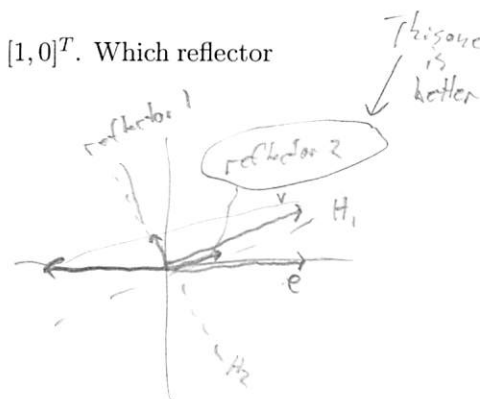
$$v = H^* w$$

↑
since H is unitary

5. [5 points,] Let $v = [1, \epsilon_M]^T$. Find the two reflectors that align v with $[1, 0]^T$. Which reflector is numerically stable? Explain.

$$u_1 = \frac{v - e_1}{\|v - e_1\|_2}$$

$$u_2 = \frac{v + e_1}{\|v + e_1\|_2} = \frac{v + e_1}{\|v + e_1\|_2}$$



We use the reflector that reflects further in order to avoid rounding error due to subtraction.

6. [20 points,] Let $v \in \mathbb{R}^n$. Suggest an algorithm for constructing an orthonormal basis for $\mathbb{R}^n \setminus \text{span}\{v\}$. Give an estimate for the work complexity of your algorithm.

$$v = v_1 + \dots + v_n$$

projects v onto q_1

$$r = v - q_1 q_1^* v - q_2 q_2^* v - \dots - q_n q_n^* v$$

v is a single vector and so it spans 1 dimension
 $\mathbb{R}^n - v$ should be of dimension \mathbb{R}^{n-1} ✓

$$q_n = \frac{v}{\|v\|}$$

Build $n-1$ more vectors orthogonal to q_n

how? that's
the question.

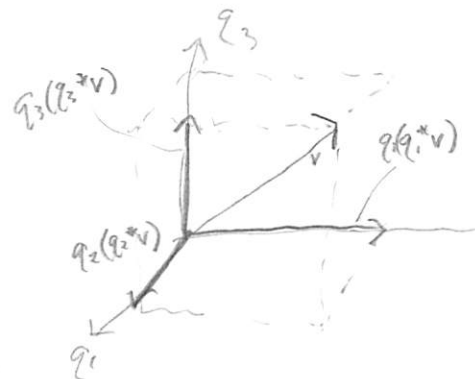
Then we have

$$Q = \begin{bmatrix} | & & | \\ q_1 & \dots & q_n \\ | & & | \end{bmatrix}$$

→ This spans \mathbb{R}^n

Then remove q_n so we're left with
 which spans $\mathbb{R}^n - v$

$$\tilde{Q} = \begin{bmatrix} | & & | \\ q_1 & \dots & q_{n-1} \\ | & & | \end{bmatrix}$$



7. [10 points,] Assume you performing a numerical stability analysis for a function $y = f(x)$. Let $[f(x)]$ denote the numerical evaluation of f using IEEE-754-compliant hardware. Assume that you've managed to show that for every x there exists a y such that $\|[f(x)] - f(y)\| = \mathcal{O}(\epsilon_M)$ with $\|x - y\| = \mathcal{O}(\epsilon_M)$. Using this fact, show that the algorithm for computing f numerically stable.

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Def. numerical stability $\forall x \exists y: \frac{\|y - x\|}{\|x\|} = \mathcal{O}(\epsilon_n)$ and $\frac{\|[f(x)] - f(y)\|}{f(x)} = \mathcal{O}(\epsilon_n)$

This is definition of forward stable

$$\forall x \exists y: \|x - y\| = \mathcal{O}(\epsilon_n) \text{ and } \|[f(x)] - f(y)\| = \mathcal{O}(\epsilon_n)$$

8. [5 points,] What is the smallest positive integer that is not a double precision floating point number (i.e., cannot be represented exactly in IEEE-754)?

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$$(1 = 1.00 \dots \times 2^0)$$

$$(2 = 1.00 \dots \times 2^1)$$

9. [10 points,] Let $A \in \mathbb{R}^{m \times n}$ be a full rank matrix with $m < n$. Show how you can use the Householder QR factorization of A^T to compute the least squares solution to $Ax = b$ that is equal to $x = V\Sigma^{-1}U^Tb$, the solution we obtain using the reduced SVD of A . (Without using the actual SVD of A or any other matrix.) Explain.

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10. [10 points,] Let $A \in \mathbb{A}^{m \times n}$ be a full-rank matrix with $m > n$. Let Q and R be the (reduced) QR factorization of A computed using the modified Gram-Schmidt (MGS) orthogonalization. Since Q is not orthonormal, we cannot use this factorization to solve $Ax = b$. However, consider using the MGS-factorization of the augmented matrix $[A \quad b]$

$$[A \quad b] = [Q \quad q] \begin{bmatrix} R & w \\ 0 & \beta \end{bmatrix}$$

Noticing

$$Ax - b = [A \quad b] \begin{bmatrix} x \\ -1 \end{bmatrix}$$

show that this factorization can be used to compute the least squares solution of $Ax = b$ by setting $x = R^{-1}w$. Is this algorithm numerically stable?