PAPER

Compactness of Finite Union of Regular Patterns and Regular Patterns without Adjacent Variables

Naoto TAKETA†, ***Nonmember***, **Tomoyuki UCHIDA**†, **Takayoshi SHOUDAI**††, **Satoshi MATSUMOTO**†††, **Yusuke SUZUKI**†, ***and*** **Tetsuhiro MIYAHARA**†, ***Members***

SUMMARY A regular pattern is a string consisting of constant symbols and distinct variable symbols. The language $L(p)$ of a regular pattern p is the set of all constant strings obtained by replacing all variable symbols in the regular pattern p with constant strings. $\mathcal{RP}^k$ denotes the class of all sets consisting at most k ($k \geq 2$) regular patterns. Sato et al. (Proc. ALT'98, 1998) showed that the finite set $S_2(P)$ of symbol strings is a characteristic set of $L(P) = \bigcup_{p \in P} L(p)$, where $S_2(P)$ is obtained from $P \in \mathcal{RP}^k$ by substituting variables with symbol strings of at most length 2. They also showed that $\mathcal{RP}^k$ has compactness with respect to containment, if the number of constant symbols is greater than or equal to $2k-1$. In this paper, we check the their results and correct the error of the proof of their theorem. Further, we consider the set $\mathcal{RP}_{NAV}$ of all non-adjacent regular patterns, which are regular patterns without adjacent variables, and show that the set $S_2(P)$ obtained from a set P in the class $\mathcal{RP}^k_{NAV}$ of at most k ($k \geq 1$) non-adjacent regular patterns is a characteristic set of $L(P)$. Further we show that $\mathcal{RP}^k_{NAV}$ has compactness with respect to containment if the number of constant symbols is greater than or equal to $k+2$. Thus we show that we can design an efficient learning algorithm of a finite union of pattern languages of non-adjacent regular patterns with the number of constant symbols which is smaller than the case of regular patterns.
key words: *Regular Pattern Language, Compactness w.r.t. Containment, Non-adjacent Regular Patterns Language*

1. Introduction

A pattern is a string consisting of constant symbols and variable symbols [1], [2]. For example, we consider constant symbols a, b, c and variable symbols x, y, then $axbxcy$ is a pattern. $\mathcal{P}$ denotes the set of all patterns. For a pattern $p \in \mathcal{P}$, *the pattern language generated by* p, denoted by $L(p)$, or simply called a pattern language, is the set of all strings obtained by replacing all variable symbols with constant symbol strings, where the same variable symbol is replaced by the same constant string. For example the pattern language $L(axbxcy)$ generated by the above pattern $axbxcy$ denotes $\{aubucw \mid$ u and w are constant strings that are not $\varepsilon\}$. A pattern where each variable symbol appears at most once is called a *regular pattern*. For example, a pattern $axbxcy$ is not a regular pattern, but a pattern $axbzcy$ with variable symbols x, y, z is a regular pattern. $\mathcal{RP}$ denotes the set of all regular patterns.

If a pattern $p \in \mathcal{P}$ is obtained from a pattern $q \in \mathcal{P}$ by replacing variable symbols in q with patterns, we say that q is a *generalization* of p and denote this by $p \preceq q$. For example, a pattern $q = axz$ is a generalization of a pattern $p = axbxcy$, because p is obtained from q by replacing the variable z in q with a pattern $bxcy$. So we write $p \preceq q$. For patterns $p, q \in \mathcal{P}$, it is obvious that $p \preceq q$ implies $L(p) \subseteq L(q)$. But, the converse, that is, the statement that $L(p) \subseteq L(q)$ implies $p \preceq q$ does not always hold. With respect to this statement, Mukouchi [3] showed that if the number of constant symbols is greater than or equal to 3, for any regular pattern $p, q \in \mathcal{RP}$, $L(p) \subseteq L(q)$ implies $p \preceq q$.

We denote by $\mathcal{RP}^+$ the class of all non-empty finite sets of regular patterns and by $\mathcal{RP}^k$ the class of at most k ($k \geq 2$) regular patterns. For a set of regular patterns $P \in \mathcal{RP}^k$ we define $L(P) = \bigcup_{p \in P} L(p)$ and consider the class $\mathcal{RPL}^k$ of regular pattern languages of $\mathcal{RP}^k$, where $\mathcal{RPL}^k = \{L(P) \mid P \in \mathcal{RP}^k\}$. Let $P, Q \in \mathcal{RP}^k$ and $Q = \{q_1, \ldots, q_k\}$. We denote by $P \sqsubseteq Q$ that for any regular pattern $p \in P$ there exists a regular pattern q_i such that $p \preceq q_i$ holds. From definition, it is obvious that $P \sqsubseteq Q$ implies $L(P) \subseteq L(Q)$. Then Sato et al. [4] shows that if $k \geq 3$ and the number of constant symbols is $2k-1$ then the finite set $S_2(P)$ of constant symbols obtained from $P \in \mathcal{RP}^k$ by substituting variable symbols with constant strings of at most 2 length is a characteristic set of $L(P)$, that is, for any regular pattern language $L' \in \mathcal{RPL}^k$, $S_2(P) \subseteq L'$ implies $L(P) \subseteq L'$. Thus they shows that the following three statements: (i) $S_2(P) \subseteq L(Q)$, (ii) $P \sqsubseteq Q$ and (iii) $L(P) \subseteq L(Q)$ are equivalent. But the Lemma 14 in [4], which is used in this results, contains an error. In this paper we correct this lemma and give a correct proof showing the equivalence of the three statements shown in [4]. Sato et al. [4] shows that $\mathcal{RP}^k$ has compactness with respect to containment if the number of constant symbols is greater than or equal to $2k-1$. On the contrary to this result, we show that the set $S_2(P)$ obtained from a set P in the class $\mathcal{RP}^k_{NAV}$ of at most k ($k \geq 1$) regular patterns having non-adjacent variables is a characteristic set of $L(P)$. Further, we show that if the number of constant symbols is greater than or equal to $k+2$ then $\mathcal{RP}^k_{NAV}$ has compactness with respect to containment. In Table 1 we summarize the all results in this paper.

Table 1 The conditions of the number of constant symbols with respect to the compactness of inclusion

k	2	$\geqq 3$
$\mathcal{RP}^k$	$\geqq 4$	$\geqq 2k-1$
$\mathcal{RP}^k_{NAV}$	$\geqq k+2$	

Manuscript received January 1, 2015.
Manuscript revised January 1, 2015.
†Graduate School of Information Sciences, Hiroshima City University
††Department of Computer Science and Engineering, Fukuoka Institute of Technology
†††Faculty of Science, Tokai University
DOI: 10.1587/transinf.E0.D.1

Mukouchi [5] examined the decision problem of determining whether a containment relation exists between the languages generated by two given patterns. The inductive inference of formal languages—specifically, pattern languages [2] and unions of pattern languages [6], [7] from positive data has been extensively investigated. Arimura et al. [8] introduced a formal framework for the efficient generalization of unions of pattern languages, presenting a polynomial-time algorithm to identify the minimal set of patterns whose union encompasses a given set of positive examples. In a subsequent study, Arimura et al. [9] proposed the concept of strong compactness of containment for unions of regular pattern languages. Day et al. [10] established that pattern languages are, in general, not closed under standard language operations such as union, intersection, and complement. Matsumoto et al. [11] developed an efficient query learning algorithm for regular pattern languages that requires only a single positive example and a linear number of membership queries. More recently, Takeda et al. [12] proposed a query learning algorithm that utilizes a deep learning model trained on a set of strings as an oracle, enabling the learned features to be visualized as regular patterns. Subsequent research extended the study of regular patterns to Elementary Formal Systems (EFS) [13], thereby broadening the theoretical foundation of pattern languages. This extension inspired further work on tree patterns [14], [15] for generating tree languages, as well as on the development of Formal Graph Systems [16]. These advancements have facilitated the formalization and efficient learning of increasingly complex structured data beyond strings, fostering applications in domains such as grammatical inference and graph-based learning.

This paper is organized as follows. In Sect.2 as preparations, we give definitions of pattern languages, regular pattern languages and compactness, and then introduce the results of Sato et al.[4]. In Sect.3, we show that $S_2(P)$ is a characteristic set of $L(P)$ in $\mathcal{RPL}^k$ and $\mathcal{RP}^k$ has compactness with respect to containment. In Sect.4, we propose regular patterns having non-adjacent variables, show that $S_2(P)$ obtained from a set P in $\mathcal{RP}^k_{NAV}$ is a characteristic set of $L(P)$, and $\mathcal{RP}^k_{NAV}$ has compactness with respect to containment.

2. Preliminaries

2.1 Basic definitions and notations

Let Σ be a non-empty finite set of constant symbols. Let X be an infinite set of variable symbols such that $\Sigma \cap X = \emptyset$ holds. Then, a *string* on $\Sigma \cup X$ is a sequence of symbols in $\Sigma \cup X$. Particularly, the string having no symbol is called the *empty string* and is denoted by ε. We denote by $(\Sigma \cup X)^*$ the set of all strings on $\Sigma \cup X$ and by $(\Sigma \cup X)^+$ the set of all strings on $\Sigma \cup X$ except ε, i.e., $(\Sigma \cup X)^+ = (\Sigma \cup X)^* \setminus \{\varepsilon\}$.

A *pattern* on $\Sigma \cup X$ is a string in $(\Sigma \cup X)^*$. Note that the empty string ε is a pattern on $\Sigma \cup X$. A pattern p is said to be *regular* if each variable symbol appears at most once in p. The length of p, denote by $|p|$, is the number of symbols in p. Note that $|\varepsilon| = 0$ holds. The set of all patterns and regular patterns on $\Sigma \cup X$ are denoted by $\mathcal{P}$ and $\mathcal{RP}$, respectively. For a set S, we denote by $\sharp S$ the number of elements in S. Let p, q be strings. If p and q are equal as strings, we denote it by $p = q$. We denote by $p \cdot q$ the string obtained from p and q by concatenating q after p. Note that for a string p and the empty string ε, $p \cdot \varepsilon = \varepsilon \cdot p = p$.

A substitution θ is a mapping from $(\Sigma \cup X)^*$ to $(\Sigma \cup X)^*$ such that (1) θ is a homomorphism with respect to string concatenation, i.e., $\theta(p \cdot q) = \theta(p) \cdot \theta(q)$ holds for patterns p and q, (2) $\theta(\varepsilon) = \varepsilon$ holds, (3) for each constant symbol $a \in \Sigma$, $\theta(a) = a$ holds, and (4) for each variable symbol $x \in X$, $|\theta(x)| \geq 1$ holds. Let $x_1, \ldots, x_n$ are variable symbols and $p_1, \ldots, p_n$ non-empty patterns. The notation $\{x_1 := p_1, \ldots, x_n := p_n\}$ denotes a substitution that replaces each variable symbol x_i with a non-empty pattern p_i for each $i \in \{1, \ldots, n\}$. For a pattern p and a substitution $\theta = \{x_1 := p_1, \ldots, x_n := p_n\}$, we denote by $p\theta$ a new pattern obtained from p by replacing variable symbols $x_1, \ldots, x_n$ in p with patterns $p_1, \ldots, p_n$ according to θ, respectively.

For a pattern p and q, the pattern q is a *generalization* of p, or p is an *instance* of q, denoted by $p \preceq q$, if there exists a substitution θ such that $p = q\theta$ holds. If $p \preceq q$ and $p \succeq q$ hold, we denote it by $p \equiv q$. The notation $p \equiv q$ means that p and q are equal as strings except for variable symbols. For a pattern p, the *pattern language* of p, denoted by $L(p)$, is the set $\{w \in \Sigma^* \mid w \preceq p\}$. For patterns p and q, it is clear that $L(p) = L(q)$ if $p \equiv q$, and $L(p) \subseteq L(q)$ if $p \preceq q$. Note that $L(\varepsilon) = \{\varepsilon\}$. In particular, if p is a regular pattern, we say that $L(p)$ is a *regular pattern language*. The set of all pattern languages and regular patterns languages are denoted by $\mathcal{PL}$ and $\mathcal{RPL}$, respectively.

Lemma 1 (Mukouchi(Theorem 6.1, [3])): Suppose $\sharp\Sigma \geq 3$. Let p and q be regular patterns. Then $p \preceq q$ if and only if $L(p) \subseteq L(q)$.

Next, we consider unions of pattern languages. The class of all non-empty finite subsets of $\mathcal{P}$ is denoted by $\mathcal{P}^+$, i.e., $\mathcal{P}^+ = \{P \subseteq \mathcal{P} \mid 0 < \sharp P < \infty\}$. For a positive integer k i.e., $k > 0$, the class of non-empty sets consisting of at most k patterns, i.e., $\mathcal{P}^k = \{P \subseteq \mathcal{P} \mid 0 < \sharp P \leq k\}$. For a set P of patterns, the pattern language of P, denoted by $L(P)$, is the set $\bigcup_{p \in P} L(p)$. We denote by $\mathcal{PL}^k$ the class of unions of at most k pattern languages, i.e., $\mathcal{PL}^k = \{L(P) \mid P \in \mathcal{P}^k\}$. In a similar way, we also define $\mathcal{RP}^+$, $\mathcal{RP}^k$ and $\mathcal{RPL}^k$. For P, Q in $\mathcal{P}^+$, the notation $P \sqsubseteq Q$ means that for any $p \in P$ there is a pattern $q \in Q$ such that $p \preceq q$ holds. It is clear that $P \sqsubseteq Q$ implies $L(P) \subseteq L(Q)$. However, the converse is not valid in general.

at most k

2.2 Characteristic sets

Definition 1: Let C be a class of languages, L a language in C and S a non-empty finite subset of L. We say that S is a *characteristic* set of L within C if for any $L' \in C$, $S \subseteq L'$ implies $L \subseteq L'$.

Let n be a positive integer and p a regular pattern. We denote by $S_n(p)$ the set of all strings in Σ^* obtained by replacing all variable symbols in p with strings in Σ^+ of length at most n. Moreover, for a positive integer n and a set $P \in \mathcal{RP}^+$, let $S_n(P) = \bigcup_{p \in P} S_n(p)$. It is clear that $S_n(P) \subseteq S_{n+1}(P) \subseteq L(P)$ for any positive integer n.

Theorem 1 (Sato et al.(Theorem 8, [4])): Let k be a positive integer and $P \in \mathcal{RP}^k$. Then, there exists a positive integer n such that $S_n(P)$ is a characteristic set of $L(P)$ within $\mathcal{RPL}^k$.

Theorem 2 (Sato et al.(Lemma 9, [4])): Let p, q, p_1, p_2, q_1, q_2, q_3 be regular patterns in $\mathcal{RP}$ and x a variable symbol such that $p = p_1 x p_2$ and $q = q_1 q_2 q_3$ hold. Then $p \preceq q$ if the following three conditions (i), (ii) and (iii) are holds:

(i) $p_1 \preceq q_1 q_2$, (ii) $p_2 \preceq q_2 q_3$,
(iii) q_2 contains at least one variable symbol.

a,b,c
は定数の字

Lemma 2 (Sato et al.(Lemma 10, [4])): Suppose $\sharp\Sigma \geq 3$. Let p_1, p_2, q be regular patterns in $\mathcal{RP}$ and x a variable symbol. Let a, b and c be mutually distinct constant symbols in Σ. If $p_1 a p_2 \preceq q$, $p_1 b p_2 \preceq q$ and $p_1 c p_2 \preceq q$, then $p_1 x p_2 \preceq q$ holds.

Lemma 3 (Sato et al.(Lemma 13, [4])): Suppose $\sharp\Sigma \geq 3$. Let p_1, p_2, q_1, q_2 be regular patterns in $\mathcal{RP}$ and x a variable symbol. Let a, b be constant symbols in Σ with $a \neq b$ and w a string on Σ^*. Let $p = p_1 AwxwB p_2$ and $q = q_1 AwB q_2$ be regular patterns that satisfy the following three conditions:

(i) $p_1 Aw \preceq q_1$,
(ii) $wB p_2 \preceq q_2$,
(iii) $(A, B) \in \{(a, b), (b, a)\}$.

Then, we have that $p\{x := a\} \preceq q$ and $p\{x := b\} \preceq q$ hold but $p \not\preceq q$.

From Lemma 2, the following theorem holds.

Theorem 3 (Sato et al.(Theorem 10, [4])): Let k be a positive integer. Suppose $\sharp\Sigma \geq 2k+1$. For $P \in \mathcal{RP}^+$ and $Q \in \mathcal{RP}^k$, the following (i), (ii) and (iii) are equivalent:

(i) $S_1(P) \subseteq L(Q)$, (ii) $P \sqsubseteq Q$, (iii) $L(P) \subseteq L(Q)$.

The following Example 1 in $\boxed{\textbf{??}}$ shows that Theorem 3 does not hold if $\sharp\Sigma \leq 2k$ holds.

Example 1: Let k be a positive integer and $\Sigma = \{a_1, \ldots, a_k, b_1, \ldots, b_k\}$. Let $w_1, \ldots, w_k$ be regular patterns in $\mathcal{RP}$ such that $w_k = \varepsilon$ and for $i = 1, 2, \ldots, k-1$, $w_i = w_{i+1} b_{i+1} a_{i+1} w_{i+1}$ hold. Let p, $q_1, \ldots, q_k$ be regular patterns in $\mathcal{RP}$ such that $p = x_1 a_1 w_1 x w_1 b_1 x_2$ and for $i = 1, 2, \ldots, k$, $q_i = x_1 a_i w_i b_i x_2$ hold. Let Q be a set $\{q_1, \ldots, q_k\}$ in $\mathcal{RP}^k$. For $i = 1$, we have $p\{x := a_1\} = (x_1 a_1 w_1) a_1 (w_1 b_1 x_2) = q_1\{x_1 := x_1 a_1 w_1\} \preceq q_1$. For $i \geq 2$, from the definition of w_i, we easily see that $w_1 = (w_i b_i) w^{(i)} = w'(i)(a_i w_i)$ for some strings $w^{(i)}$ and $w'^{(i)}$. Then, for each $i \geq 2$,

$$\begin{aligned}
p\{x := a_i\} &= (x_1 a_1 w_1) a_i (w_1 b_1 x_2) \\
&= (x_1 a_1 w_1) a_i (w_i b_i w^{(i)}) b_1 x_2 \\
&= (x_1 a_1 w_1)(a_i w_i b_i)(w^{(i)} b_1 x_2) \\
&= q_i\{x_1 := x_1 a_1 w_1, x_2 := w^{(i)} b_1 x_2\} \\
&\preceq q_i, \\
p\{x := b_i\} &= (x_1 a_1 w_1) b_i (w_1 b_1 x_2) \\
&= x_1 a_1 (w'^{(i)} a_i w_i) b_i (w_1 b_1 x_2) \\
&= (x_1 a_1 w'^{(i)}) a_i w_i b_i (w_1 b_1 x_2) \\
&= q_i\{x_1 := x_1 a_1 w'^{(i)}, x_2 := w_1 b_1 x_2\} \\
&\preceq q_i.
\end{aligned}$$

Hence, $S_1(p) \subseteq L(Q)$ holds. However, from $p \not\preceq q_i$, $L(p) \not\subseteq L(q_i)$ holds for each $i = 1, 2, \ldots, k$.

From Theorem 3, we have the following corollary.

Corollary 1 (Sato et al.(Corollary 12, [4])): Suppose $\sharp\Sigma \geq 3$. For two regular patterns p and q, the following (i), (ii) and (iii) are equivalent:

(i) $S_1(p) \subseteq L(q)$, (ii) $p \preceq q$, (iii) $L(p) \subseteq L(q)$.

2.3 Basic word equations

Proposition 1: Let w be a string in Σ^* and a, b constant symbols in Σ. If

$$wa = bw \tag{1}$$

holds, then $a = b$ holds.

Proof. Since it is trivial, we omit the proof. □

Proposition 2: Let w be a string in Σ^* and a, b, c, d constant symbols in Σ. If

$$wda = bcw \tag{2}$$

holds, then $(b, c) \in \{(a, d), (d, a)\}$ holds.

Proof. We will prove this proposition by induction on the length of w, i.e., $|w|$.

- $|w| = 0, 1, 2, 3$: it is straightforward to observe that $(b, c) \in \{(a, d), (d, a)\}$ holds.
- $|w| \geq 4$: We assume that for any string u with $0 \leq |u| < n$, if $uda = bcu$ holds, $(b, c) \in \{(a, d), (d, a)\}$ holds. Since the string w has a prefix bc and a suffix da, there exists a string u with $|u| = |w| - 4 < |w|$ such that $w = bcuda$ holds. Since $wda = bcw$, we have $bcudada = bcbcuda$, and then $uda = bcu$. Thus, from the assumption, we get $(b, c) \in \{(a, d), (d, a)\}$.

From the above, we conclude that if $wda = bcw$ holds, then $(b, c) \in \{(a, d), (d, a)\}$ holds. □

The conclusion from Proposition 2 shows that $(a, d) \in \{(b, c), (c, b)\}$. Therefore, if the equation $daw = wbc$ holds, we arrive at the same conclusion.

Proposition 3: Let w, w' be strings of constant symbols in Σ and a, b, c, d constant symbols in Σ. If

$$wdaw' = w'bcw \tag{3}$$

holds, then $(b, c) \in \{(a, d), (d, a)\}$ holds.

Proof. We will prove this proposition by an induction on $|w| + |w'|$. Without loss of generality, we assume that $|w| \geq |w'|$ because, if $|w| < |w'|$, we arrive at the same conclusion that $(a, d) \in \{(b, c), (c, b)\}$ holds.

- $|w| \geq 0$ and $|w'| = 0$: Eq. (3) reduces to $wda = bcw$. By Proposition 2, $(b, c) \in \{(a, d), (d, a)\}$ holds.

We assume that for constant strings u and u' with $|u| + |u'| < |w| + |w'|$, if $udau' = u'bcu$ holds, then $(b, c) \in \{(a, d), (d, a)\}$ holds. We divide the relations between $|w|$ and $|w'|$ into the following four cases:

- $0 < |w'| \leq |w| \leq |w'| + 1$: When either $|w| = |w'|$ or $|w| = |w'| + 1$, Eq. (3) is illustrated in Figs. 1 and 2, respectively. If $|w| = |w'|$, $(b, c) = (d, a)$ holds. If $|w| = |w'| + 1$, $d = c$ and $w = w'b = aw'$ hold. From Proposition 1, we deduce that $b = a$. Therefore, $(b, c) \in \{(a, d), (d, a)\}$ holds.
- $|w'| + 2 \leq |w| \leq 2|w'| - 1$: In Eq. 3, since $|wdaw'| = |w'bcw| = |w| + |w'| + 2$, a suffix of w overlaps with a prefix of w, as illustrated in Fig. 3. That is, there exists a constant string u of length $2|w| - (|w| + |w'| + 2) = |w| - |w'| - 2$ such that u is both a prefix and a suffix of w. Since uda has a length of $|w| - |w'|$, it is also a prefix of w. Similarly, bcu is a suffix of w. Because $|w| - (|uda| + |bcu|) = 2|w| - |w'| \geq 1$, there exists a constant string u' of length $2|w'| - |w|$ such that $w = udau'bcu$ holds. Since w' is a suffix of w and $|u'bcu| = (2|w'| - |w|) + 2 + (|w| - |w'| - 2) = |w'|$, we have $w' = u'bcu$. Similarly, $w' = udau'$. Thus, we derive the equation $u'bcu = udau'$. Since $|u| = |w| - |w'| - 2 \leq |w| - 3 < |w|$ and $|u'| = 2|w'| - |w| < |w'|$, i.e., $|u| + |u'| < |w| + |w'|$, the induction hypothesis on $|u| + |u'|$ implies that $(b, c) \in \{(a, d), (d, a)\}$ holds.
- $2|w'| \leq |w| \leq 2|w'| + 3$: When $|w| = 2|w'|$, it is straightforward to observe that $w = w'w'$. Therefore, $w'da = bcw'$ holds, as illustrated in Fig. 4. From Proposition 2, $(b, c) \in \{(a, d), (d, a)\}$ holds. When $|w| = 2|w'| + i$ $(i = 1, 2, 3)$, Eq. (3) is depicted in Figs. 5, 6, and 7, respectively. When $|w| = 2|w'| + 2$, it is clear that $(b, c) = (d, a)$. When $|w| = 2|w'| + 1$ and $|w| = 2|w'| + 3$, Proposition 1 implies that $(b, c) = (a, d)$ holds.
- $2|w'| + 4 \leq |w|$: Since the strings $w'bc$ and adw' are a prefix and a suffix of w, respectively, and $|w'bc| + |adw'| = 2|w'| + 4$, there exists a string u with $|u| \geq 0$ such that $w = w'bcudaw'$ holds. From Eq. (3), $w'bcudaw'daw' = w'bcw'bcudaw'$, i.e., $udaw' = w'bcu$ holds, as illustrated in Fig. 8. Let $u' = w'$. Since $|u| + |u'| = |w| - (2|w'| + 4) + |w'| < |w| + |w'|$, the induction hypothesis on $|u| + |u'|$ implies that $(b, c) \in \{(a, d), (d, a)\}$ holds.

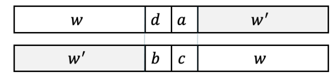


Fig. 1 Case $|w| = |w'|$ in Proposition 3

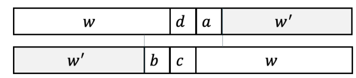


Fig. 2 Case $|w| = |w'| + 1$ in Proposition 3

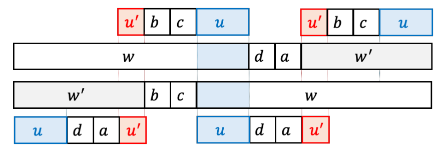


Fig. 3 Case $|w'| + 2 \leq |w| \leq 2|w'| - 1$ in Proposition 3

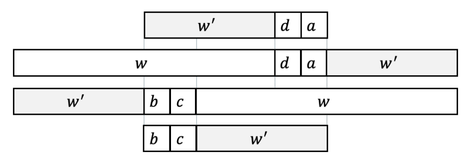


Fig. 4 Case $|w| = 2|w'|$ in Proposition 3

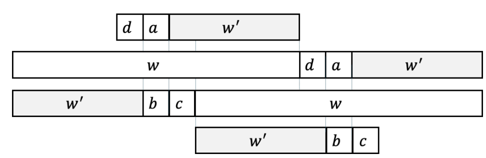


Fig. 5 Case $|w| = 2|w'| + 1$ in Proposition 3

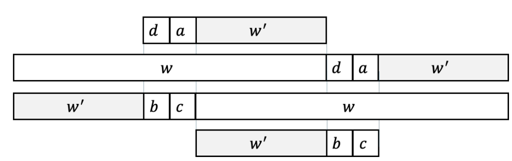


Fig. 6 Case $|w| = 2|w'| + 2$ in Proposition 3

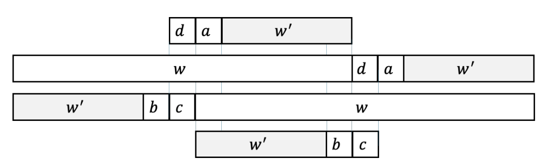


Fig. 7 Case $|w| = 2|w'| + 3$ in Proposition 3

From the above, we conclude that if $wdaw' = w'bcw$, then $(b, c) \in \{(a, d), (d, a)\}$ holds. □

3. Compactness for Sets of Regular Patterns

3.1 Compactness

In this section, we define the compactness of sets of regular patterns, formally. Then, if $\sharp\Sigma \geq 2k - 1$ holds, we show that

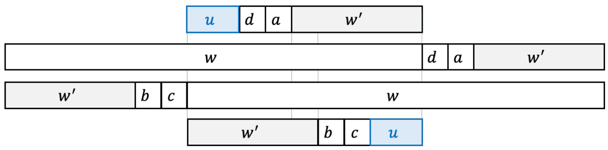


Fig. 8 Case $2|w'| + 4 \le |w|$ in Proposition 3

$\mathcal{RP}^k$ has compactness with respect to the containment.

Definition 2: Let C be a subset of $\mathcal{RP}^+$. For any regular pattern $p \in \mathcal{RP}$ and any set $Q \in C$, the set C said to have *compactness with respect to containment* if there exists a regular pattern $q \in Q$ such that $L(p) \subseteq L(q)$ holds if $L(p) \subseteq L(Q)$ holds.

Let $D \subset \mathcal{RP}$ with $\sharp D = 2$ or 3, and let p, q be regular patterns in $\mathcal{RP}$. In the following subsections (Subsecs. 3.2–3.5), we provide the conditions on D under which the implication holds: if $p\{x := r\} \preceq q$ for all $r \in D$, then $p\{x := xy\} \preceq q$. It is obvious if the variable symbol x does not appear in p. Therefore, in the following lemmas and propositions, let $p = p_1 x p_2$, where $p_i \in \mathcal{RP}$ $(i = 1, 2)$ and x is a variable symbol.

First of all, we consider the correspondence from $r \in D$ to some string in q when $p\{x := r\} \preceq q$ holds. The symbols in D correspond to either a variable or a constant symbol in q. If D has a constant string ab of length 2 for $a, b \in \Sigma$, there are three possible strings in q that correspond to ab in $p\{x := ab\}$ as follows: For $y_1 \in X$,

(a) ab, (b) ay_1, (c) $y_1 b$.

If there exists (b) ay_1 in q that corresponds to ab, i.e., there exist q_1 and $q_2 \in \mathcal{RP}$ such that

(1) $p_1 ab p_2 \preceq q_1 a y_1 q_2$,
(2) $p_1 \preceq q_1$, and
(3) either $p_2 \preceq q_2$ or $p_2 \preceq y_1' q_2$ for $y_1' \in X$.

Let $D' = (D \setminus \{ab\}) \cup \{ay\}$. It is straightforward to see that $p\{x := ay\} = p_1 a y p_2 \preceq q_1 a y_1 q_2$ holds. Thus, $p\{x := r\} \preceq q$ for all $r \in D'$ holds. Let $D'' = (D \setminus \{ab\}) \cup \{yb\}$. By a similar discussion, if there exists (c) $y_1 b$ in q that corresponds to ab, $p\{x := r\} \preceq q$ for all $r \in D''$ holds. Therefore, in this paper, we make the following definition on D:

Definition 3: Let $p, q \in \mathcal{RP}$. Let $D \subset \mathcal{RP}$ such that for all $r \in D$, $|r| = 2$ and $p\{x := r\} \preceq q$ holds. Then, if for any $ab \in D$ $(a, b \in \Sigma)$, $p\{x := ay\} \not\preceq q$ and $p\{x := yb\} \not\preceq q$ hold for any $y \in X$ that does not appear in q, the set D is said to be *maximally generalized on* (p, q).

#Σ≥2?
c≠0?

3.2 $D = \{ay, by\}$ and $D = \{ya, yb\}$

Lemma 4 (Sato et al.[4]): Let Σ be an alphabet with $\sharp\Sigma \ge 3$ and let p, q be regular patterns on $\Sigma \cup X$. Let D be the following set of regular patterns on $\Sigma \cup X$, where y is a variable symbol that does not appear in p and q:

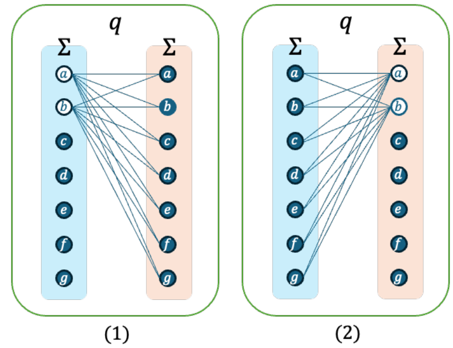


Fig. 9 Let $\Sigma = \{a, b, c, d, e, f, g\}$ and $p, q \in \mathcal{RP}$. We assume that the symbols in Σ are mutually distinct. These figures (1) and (2) express two cases $D = \{ay, by\}$ and $D = \{ya, yb\}$, respectively. In these cases, if $p\{x := r\} \preceq q$ for all $r \in D$, then $p\{x := xy\} \preceq q$ holds.

(i) $D = \{ay, by\}$ $(a \ne b)$,
(ii) $D = \{ya, yb\}$ $(a \ne b)$.

Then, if $p\{x := r\} \preceq q$ holds for all $r \in D$, then $p\{x := xy\} \preceq q$ holds.

Proof. We assume that $p\{x := xy\} \not\preceq q$ in order to derive a contradiction. In the case of (ii), by reversing the strings p and q, we can prove that the assumption $p\{x := xy\} \preceq q$ leads to a contradiction, as in the case of (i). Therefore, in the following, we consider only the case of (i): $D = \{ay, by\}$ $(a \ne b)$.

Since $p\{x := xy\} \not\preceq q$, $p_1 a y p_2 \preceq q$ and $p_1 b y p_2 \preceq q$, there exist regular patterns q_1, q_2 on $\Sigma \cup X$ such that $q = q_1 a y_1 w b y_2 q_2$ or $q = q_1 b y_1 w a y_2 q_2$ for some variable symbols y_1, y_2 $(y_1 \ne y_2)$ and a constant string w $(|w| \ge 0)$ from Theorem 2. When $q = q_1 a y_1 w b y_2 q_2$ holds, the following four conditions hold: For $y_1', y_2' \in X$,

(1) $p_1 \preceq q_1$, (1') $p_2 \preceq w b y_2 q_2$ or $p_2 \preceq y_1' w b y_2 q_2$,
(2) $p_1 \preceq q_1 a y_1 w$, (2') $p_2 \preceq q_2$ or $p_2 \preceq y_2' q_2$.

From (2), there exist regular patterns p_1', p_1'' such that $p_1 = p_1' p_1''$, $p_1' \preceq q_1 a$ and $p_1'' \preceq y_1 w$ hold. Therefore, since $p = p_1 x p_2 = p_1' p_1'' x p_2$, if $p_2 \preceq w b y_2 q_2$ of (1') holds, $p \preceq q_1 a p_1'' x w b y_2 q_2 \equiv q\{y_1 := p_1'' x\}$ holds. If $p_2 \preceq y_1' w b y_2 q_2$ of (1') holds, $p \preceq q_1 a p_1'' x y_1' w b y_2 q_2 = q\{y_1 := p_1'' x y_1'\}$ holds. Thus, $p\{x := xy\} \preceq q\{y_1 := p_1'' x y y_1'\}$ holds. Hence, $p \preceq q$ holds. This contradicts the assumption. Therefore, we conclude that if $p\{x := r\} \preceq q$ for all $r \in \{ay, by\}$ $(a \ne b)$, then $p\{x := xy\} \preceq q$ holds. □

Let p, q be regular patterns in $\mathcal{RP}$. In this paper, the statement like Lemma 4 is illustrated by a bipartite graph (Σ, Σ, E) where $E = \{(a, b) \in \Sigma \times \Sigma \mid p\{x := ab\} \preceq q\}$. For example, the conditions (i) and (ii) in Lemma 4 are illustrated in (1) and (2) in Fig. 9, respectively.

3.3 $D = \{ya, bc, dy\}$

Lemma 5: Let Σ be an alphabet with $\boxed{\sharp\Sigma \ge 3}$ and p, q

Note that $p \not\preceq q$ だし $p \preceq q$ とすると $p\{x := ay\} \preceq p \preceq q$

$b \neq d$, $a \neq c$,

$b=c$ ならば, $c \neq d$, $a \neq b$, $c \neq a$, $d \neq b$,

$\boxed{b=c \text{ でもOK?}}$

必要なのは $b \neq a$ and $d \neq c$

$bc \in D$ より $p\{x := by\} \not\preceq q$, $p\{x := yc\} \not\preceq q$ → $b \neq d$, $a \neq c$

regular patterns on $\Sigma \cup X$. Let D be the following set of regular patterns on $\Sigma \cup X$, where y is a variable symbol in X that does not appear in p and q:

$$D = \{ya, bc, dy\} \ (b \notin \{a, d\} \text{ and } c \notin \{a, d\}).$$

?

これより

Then, if $p\{x := r\} \preceq q$ for all $r \in D$ and D is maximally generalized on (p, q), then $p\{x := xy\} \preceq q$.

$p = p_1 x p_2$

Proof. We assume that $p\{x := xy\} \not\preceq q$ in order to derive a contradiction. Since D is maximally generalized on (p, q), the regular pattern q can be expressed in one of the following forms: Let y_1, y_2 be distinct variable symbols in X and q_1, q_2, w, w' be either the empty string or a regular pattern on $\Sigma \cup X$.

(5-1) $q = q_1 A w B w' C q_2$, where $\{A, B, C\} = \{y_1 a, bc, dy_2\}$,
(5-2) $q = q_1 A w B q_2$, where $\{A, B\} = \{dy_1 a, bc\}$,
(5-3) $q = q_1 A w B q_2$, where $\{A, B\} = \{y_1 a y_2, bc\}$ $(a = d)$.

(5-1) Case of $q = q_1 A w B w' C q_2$, where $\{A, B, C\} = \{y_1 a, bc, dy_2\}$: At first, we prove the following three claims:

Claim 1. $B \notin \{y_1 a, dy_2\}$.

Proof of Claim 1. Suppose that $(A, B, C) = (dy_2, y_1 a, bc)$. The following conditions must be satisfied: For $y'_1, y'_2 \in X$,

(1) $p_1 \preceq q_1$, (1') $p_2 \preceq w y_1 a w' bc q_2$ or $p_2 \preceq y'_2 w y_1 a w' bc q_2$,
(2) $p_1 \preceq q_1 dy_2 w$ or $p_1 \preceq q_1 dy_2 w y'_1$, (2') $p_2 \preceq w' bc q_2$,
(3) $p_1 \preceq q_1 dy_2 w y_1 a w'$, (3') $p_2 \preceq q_2$.

$y'_2 = y_2$ であることが書いてないかも

When $p_2 \preceq w y_1 a w' bc q_2$ in (1') holds, let $q'_1 = q_1 dy_2$, $q'_2 = w y_1 a w'$, $q'_3 = bc q_2$. Since $p_1 \preceq q_1 dy_2 w y_1 a w'$ holds from (3), both $p_1 \preceq q'_1 q'_2$ and $p_2 \preceq q'_2 q'_3$ hold, and q'_2 contains a variable symbol. When $p_2 \preceq y'_2 w y_1 a w' bc q_2$ in (1') holds, let $q'_1 = q_1 d$, $q'_2 = y_2 w y_1 a w'$, $q'_3 = bc q_2$. Since $p_1 \preceq q_1 dy_2 w y_1 a w'$ holds from (3), both $p_1 \preceq q'_1 q'_2$ and $p_2 \preceq q'_2 q'_3$ hold, and q'_2 contains a variable symbol. In both cases, by Theorem 2, $p \preceq q$ holds. This contradicts the assumption that $p\{x := xy\} \not\preceq q$.

Similarly, we can show that any case where $(A, B, C) = (y_1 a, dy_2, bc)$, $(bc, y_1 a, dy_2)$, or $(bc, dy_2, y_1 a)$ also contradicts the assumption. Therefore, we have $B \notin \{y_1 a, dy_2\}$. (*End of Proof of Claim* 1)

Claim 2. $(A, B, C) = (y_1 a, bc, dy_2)$.

$(dy_2, bc, y_1 a)$ では?

Proof of Claim 2. From *Claim* 1, we have $B = bc$. Suppose that $(A, B, C) = (dy_2, bc, y_1 a)$, i.e., $q = q_1 dy_2 w bc w' y_1 a q_2$ holds. Then, the following conditions must be satisfied: For $y'_1, y'_2 \in X$,

(1) $p_1 \preceq q_1$, (1') $p_2 \preceq w bc w' y_1 a q_2$ or $p_2 \preceq y'_2 w bc w' y_1 a q_2$,
(2) $p_1 \preceq q_1 dy_2 w$, (2') $p_2 \preceq w' y_1 a q_2$,
(3) $p_1 \preceq q_1 dy_2 w bc w'$ or $p_1 \preceq q_1 dy_2 w bc w' y'_1$, (3') $p_2 \preceq q_2$.

From $p_1 \preceq q_1 dy_2 w$ in (2), p_1 is expressed as $p'_1 p''_1$ for some p'_1 and p''_1, where $p'_1 \preceq q_1 d$ and $p''_1 \preceq y_2 w$. When $p_2 \preceq w bc w' y_1 a q_2$ in (1'), we have $p = p_1 x p_2 = p'_1 p''_1 x p_2 \preceq q_1 d p''_1 x w bc w' y_1 a q_2 = q\{y_2 := p''_1 x\}$. Thus, $p\{x := xy\} \preceq q\{y_2 := p''_1 xy\}$ holds. This contradicts the assumption that $p\{x := xy\} \not\preceq q$. When $p_2 \preceq y'_2 w bc w' y_1 a q_2$ in (1'), we similarly have $p = p_1 x p_2 = p'_1 p''_1 x p_2 \preceq q_1 d p''_1 x y'_2 w bc w' y_1 a q_2 = q\{y_2 := p''_1 x y'_2\}$. Thus, $p\{x := xy\} \preceq q\{y_2 := p''_1 x y y'_2\}$ holds. This also contradicts the assumption. Therefore, we conclude that $(A, B, C) = (y_1 a, bc, dy_2)$. (*End of Proof of Claim* 2)

From *Claim* 2, The regular pattern q is expressed as $q_1 y_1 a w bc w' dy_2 q_2$, where $b \notin \{a, d\}$ and $c \notin \{a, d\}$. If $p\{x := xy\} \not\preceq q$ holds, the following conditions must be satisfied: For $y'_1, y'_2 \in X$,

(1) $p_1 \preceq q_1$ or $p_1 \preceq q_1 y'_1$, (1') $p_2 \preceq w bc w' dy_2 q_2$,
(2) $p_1 \preceq q_1 y_1 a w$, (2') $p_2 \preceq w' dy_2 q_2$,
(3) $p_1 \preceq q_1 y_1 a w bc w'$, (3') $p_2 \preceq q_2$ or $p_2 \preceq y'_2 q_2$.

Claim 3. w and w' contain no variable symbols.

Proof of Claim 3. Let $q'_1 = q_1 y_1 a$, $q'_2 = w bc w'$, and $q'_3 = dy_2 q_2$. From (1') and (3), $p_1 \preceq q'_1 q'_2$ and $p_2 \preceq q'_2 q'_3$. If q'_2 contains a variable symbol, then by Theorem 2, $p \preceq q$ holds. This contradicts the assumption. Therefore, w and w' contain no variable symbols. (*End of Proof of Claim* 3)

From *Claim* 3, w and w' are strings consisting of symbols in Σ. From (1') and (2'), $w bc w' d$ and $w' d$ are prefixes of p_2, and from (2) and (3), $a w bc w'$ and $a w$ are suffixes of p_1. It implies a contradiction in the following inductive way:

- $|w| = |w'|$: Directly, $b = d$ and $a = c$ hold.
- $|w| = |w'| + 1$: Also, $a = b$ holds.
- $|w| = |w'| + 2$: Since $a w bc w'$ and $a w$ are suffixes of p_1, and $|w| \geq 2$, a is a suffix of w. From (1') and (2'), we have $w = w' da$. Furthermore, since $a w bc w'$ and $a w$ are suffixes of p_1, it follows that $w = bc w'$. Thus, $w' da = bc w'$ holds. From Proposition 2, $(b, c) \in \{(a, d), (d, a)\}$ holds. Therefore, these cases contradict the conditions $b \notin \{a, d\}$ and $c \notin \{a, d\}$.
- $|w| \geq |w'| + 3$: From (2) and (3), there exists a string w'' of length $|w| - |w'| - 2$ such that $w = w'' bc w'$ holds. Moreover, from (2) and (3), since $|a w| < |w bc w'|$ and $a w = a w'' bc w'$, it follows that $a w''$ is a suffix of w. On the other hand, from (1') and (2'), $w' d$ is a prefix of w. Since $|w' d| + |a w''| = |w'| + |w''| + 2 = |w|$, it follows that $w = w' d a w''$ (Fig. 10). Therefore, $w' d a w'' = w'' bc w'$ holds. From Proposition 3, $(b, c) \in \{(a, d), (d, a)\}$ holds. This contradicts the conditions $b \notin \{a, d\}$ and $c \notin \{a, d\}$.

From the above, we conclude that all cases of (5-1) contradict the assertion that $p\{x := xy\} \not\preceq q$ and the conditions $b \notin \{a, d\}$ and $c \notin \{a, d\}$.

(5-2) Case of $q = q_1 A w B q_2$, where $\{A, B\} = \{dy_1 a, bc\}$:

$|w| < |w'|$ も同様にして示すことができると書く必要はありませんか？

{a, b, c, d} 3以上 b=a, b≠d a≠c
= {a, c, d} となるには b≠c

bc∈Dより p{x:=by}≺q となる
b≠d なる もし b=cであると a=c
p{x:=ycy}=p{x:=ya}≺q
maximally generalizedに反する. よって
b≠c

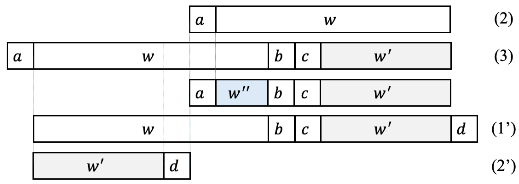


Fig. 10 Case (5-1) in Lemma 5: Relation of strings w, w', and w''

We suppose that $(A, B) = (dy_1a, bc)$, i.e., $q = q_1dy_1awbcq_2$ holds. Then, the following conditions must be satisfied for $y'_1 \in X$:

(1) $p_1 \preceq q_1$, (1') $p_2 \preceq awbcq_2$ or $p_2 \preceq y'_1awbcq_2$,
(2) $p_1 \preceq q_1d$ or $p_1 \preceq q_1dy'_1$ (2') $p_2 \preceq wbcq_2$,
(3) $p_1 \preceq q_1dy_1aw$, (3') $p_2 \preceq q_2$.

From $p_1 \preceq q_1dy_1aw$ in (3), p_1 can be expressed as $p'_1p''_1$ for some p'_1 and p''_1, where $p'_1 \preceq q_1d$ and $p''_1 \preceq y_1aw$. When $p_2 \preceq awbcq_2$ in (1'), we have

$$p = p'_1p''_1xp_2 \preceq q_1dp''_1xawbcq_2 = q\{y_1 := p''_1x\}.$$

Thus, $p\{x := xy\} \preceq q\{y_1 := p''_1xy\}$ holds. This contradicts the assumption. When $p_2 \preceq y'_1awbcq_2$ in (1'), we similarly have

$$p = p'_1p''_1xp_2 \preceq q_1dp''_1xy'_1wbcq_2 = q\{y_1 := p''_1xy'_1\}.$$

Thus, $p\{x := xy\} \preceq q\{y_1 := p''_1xyy'_1\}$ holds. This contradicts the assumption that $p\{x := xy\} \not\preceq q$. Similarly, we can show that the case $(A, B) = (bc, dy_1a)$ also contradicts the assumption.

(5-3) Case of $q = q_1AwBq_2$, where $\{A, B\} = \{y_1ay_2, bc\}$ $(a = d)$: Suppose that $(A, B) = (y_1ay_2, bc)$, i.e., $q = q_1y_1ay_2wbcq_2$ holds. Then, the following conditions must be satisfied: For $y'_1, y'_2 \in X$,

(1) $p_1 \preceq q_1$ or $p_1 \preceq q_1y'_1$ (1') $p_2 \preceq y_2wbcq_2$,
(2) $p_1 \preceq q_1y_1ay_2$, (2') $p_2 \preceq wbcq_2$ or $p_2 \preceq y'_2wbcq_2$,
(3) $p_1 \preceq q_1y_1ay_2w$, (3') $p_2 \preceq q_2$.

Let $q'_1 = q_1y_1a$, $q'_2 = y_2w$, $q'_3 = bcq_2$. From (3) and (1'), we have $p_1 \preceq q'_1q'_2$ and $p_2 \preceq q'_2q'_3$, respectively. Since q'_2 contains a variable symbol, Theorem 2 implies that $p \preceq q$ holds. This contradicts the assumption. Similarly, we can show that the case $(A, B) = (bc, y_1ay_2)$ also contradicts the assumption.

#Σ=2でもよいか?

From the above, we conclude that if $p\{x := r\} \preceq q$ for all $r = \{ya, bc, dy\}$ $(b \notin \{a, d\}$ and $c \notin \{a, d\})$, then $p\{x := xy\} \preceq q$ holds. □

The condition in Lemma 5 is illustrated in four cases (3)–(6) in Fig. 11.

Lemma 6: Let Σ be an alphabet with $\boxed{\sharp\Sigma \geq 3}$ and let p, q be regular patterns on $\Sigma \cup X$. Let D be one of the following sets of regular patterns on $\Sigma \cup X$, where y is a variable symbol in X that does not appear in p and q.

不要では?

$$D = \{ya, bc, dy\} \ (b = a, \boxed{b \neq d}, \text{ and } c \notin \{a, d\}),$$

Then, if $p\{x := r\} \preceq q$ for all $r \in D$ and D is maximally generalized on (p, q), then $p\{x := xy\} \preceq q$.

We note that if $b = d$, then, because $p\{x := dy\} \preceq q$, $p\{x := bc\} \preceq q$ is always satisfied. In this sense, D essentially consists of only two elements. To avoid this, we assume $b \neq d$

Dがmaximally generalizedになるため

ピリオド

Proof. We assume that $p\{x := xy\} \not\preceq q$ in order to derive a contradiction. The proof is almost the same as the proof of Lemma 5. Since $p\{x := r\} \preceq q$ for all $r \in D$ and D is maximally generalized on (p, q), there are three strings of length 2 corresponding to ya, bc, dy in q. The symbols appearing in D correspond to either a variable or a constant symbol in q. Let y_1 and y_2 be variable symbols appearing in q. The strings ya and dy must correspond to the strings y_1a and dy_2 in q, respectively. For the same reasons stated at the beginning of Lemma 5, the string bc corresponds to the string bc in q as well. Let A, B, C be regular patterns on $\Sigma \cup X$, where $\{A, B, C\} = \{y_1a, ac, dy_3\}$. Since $p\{x := xy\} \not\preceq q$, q can be expressed in one of the following four forms: Let y_1, y_2 be distinct variable symbols in X, and q_1, q_2, w, w' either the empty string or a regular pattern on $\Sigma \cup X$. From the conditions $b = a$ and $b \neq d$, it follows that $a \neq d$.

(6-1) $q = q_1AwBw'Cq_2$, where $\{A, B, C\} = \{y_1a, ac, dy_2\}$,
(6-2) $q = q_1AwBq_2$, where $\{A, B\} = \{y_1ac, dy_2\}$,
(6-3) $q = q_1Aq_2$, where $A = dy_1ac$.

集合でよいの? どうか?

In cases (6-1) and (6-2), similar to Lemma 5, it is shown that $q = q_1y_1awacw'dy_2q_2$ and $q = q_1y_1acwdy_2q_2$, respectively, where w and w' contain no variable symbols.

(6-1) Case of $q = q_1AwBw'Cq_2$, where $\{A, B, C\} = \{y_1a, ac, dy_2\}$: The following conditions must be satisfied:

(1) $p_1 \preceq q_1$, (1') $p_2 \preceq wacw'dy_2q_2$,
(2) $p_1 \preceq q_1y_1aw$, (2') $p_2 \preceq w'dy_2q_2$,
(3) $p_1 \preceq q_1y_1awacw'$, (3') $p_2 \preceq q_2$.

From (1') and (2'), $wacw'd$ and $w'd$ are prefixes of p_2, and from (2) and (3), $awacw'$ and aw are suffixes of p_1. It implies a contradiction in the following inductive way:

- $|w| = |w'|$: $c = a$ holds.
- $|w| = |w'| + 1$: $w = w'd = cw'$ holds. Thus, from Proposition 1, $c = d$ holds.
- $|w| = |w'| + 2$: $w = w'da = acw'$ holds. From Proposition 2, $c \in \{a, d\}$ holds.
- $|w| \geq |w'| + 3$: From (2) and (3), there exists a string w'' of length $|w| - |w'| - 2$ such that $w = w''acw'$ holds. Moreover, from (2) and (3), since $|aw| < |wacw'|$ and $aw = aw''acw'$, it follows that aw'' is a suffix of w. On the other hand, from (1') and (2'), $w'd$ is a prefix of w.

D = dya, ay, bby
π1 = xaby
π2 = xaybz
L(π2)⊂L(π1)

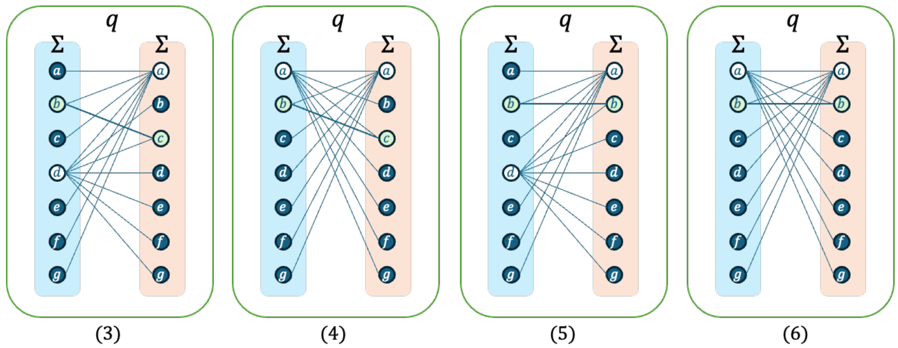


Fig. 11 Let $\Sigma = \{a, b, c, d, e, f, g\}$ and $p, q \in \mathcal{RP}$. We assume that the symbols in Σ are mutually distinct. The figure (3) expresses case $D = \{ya, bc, dy\}$ in Lemma 5. The figures (4), (5), and (6) express three cases $D = \{ya, bc, ay\}$, $D = \{ya, bb, dy\}$, and $D = \{ya, bb, ay\}$, respectively. In these cases, if $p\{x := r\} \preceq q$ for all $r \in D$ and D is maximally generalized on (p, q), then $p\{x := xy\} \preceq q$ holds.

ya

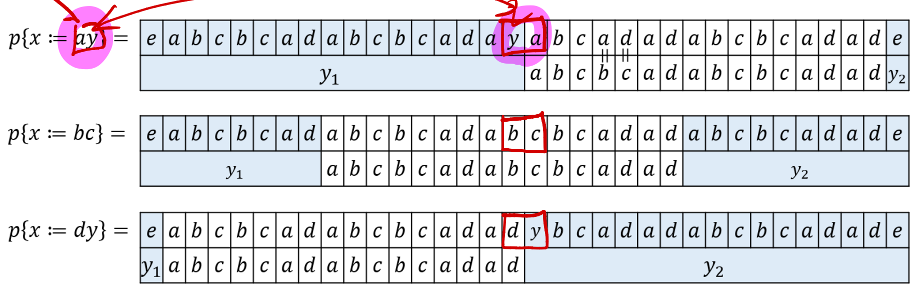


Fig. 12 Substitutions for p and each correspondence to q.

P.9で参照
順序がおかしいかも

Since $|w'd| + |aw''| = |w'| + |w''| + 2 = |w|$, we have $w = w'daw''$. Therefore, $w'daw'' = w''acw'$ holds (Fig. 13). From Proposition 3, we have $c \in \{a, d\}$.

- $|w'| = |w| + 1$: From (1') and (2'), $c = d$ holds.
- $|w'| = |w| + 2$: From (1') and (2'), d is a prefix of w'. Thus, from (2) and (3), $w' = wac = daw$ holds. From Proposition 2, $c \in \{a, d\}$ holds.
- $|w'| \geq |w| + 3$: From (1') and (2'), there exists a string w'' of length $|w| - |w'| - 2$ such that $w' = wacw''$ holds. Moreover, from (1') and (2'), since $|w'd| < |wacw'|$ and $w'd = wacw''d$, $w'd$ is a prefix of w'. On the other hand, from (1') and (2'), awd is a suffix of w'. Since $|w''d| + |aw| = |w''| + |w| + 2 = |w'|$, we have $w' = w''daw$. Therefore, $w''daw = wacw''$ holds. From Proposition 3, we have $c \in \{a, d\}$.

w''

w'

2

3

All the cases contradict the condition $c \notin \{a, d\}$. Therefore, if $b = a$, $b \neq d$, and $c \notin \{a, d\}$ are satisfied, case (6-1) is impossible.

他の場合も証明を書くか省略することを書くほうがいいかと

(6-2) Case of $q = q_1 AwBq_2$, where $\{A, B\} = \{y_1 ac, dy_2\}$: For $q = q_1 y_1 acwdy_2 q_2$, the following conditions must be satisfied:

不要では?

(1) $p_1 \preceq q_1$, (1') $p_2 \preceq cwdy_3 q_2$,

(2) $p_1 \preceq q_1 y_1$, (2') $p_2 \preceq wdy_3 q_2$,

(3) $p_1 \preceq q_1 y_1 acwdy_3$, (3') $p_2 \preceq q_2$.

- If $|w| = 0$, from (1') and (2'), the prefix of p_2 is cd and d. Thus, we have $c = d$.

集合でよいのでしょうか?

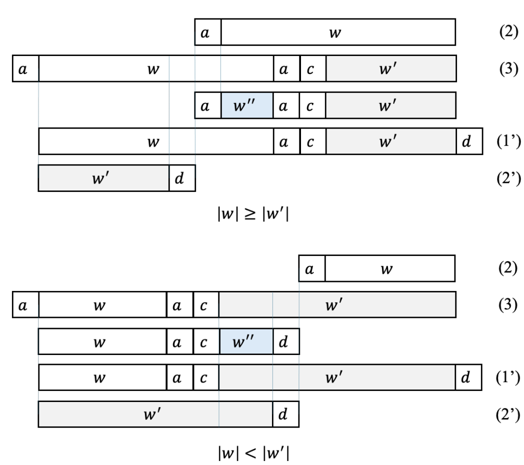
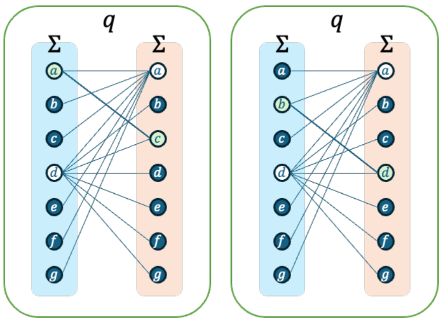


Fig. 13 Case (6-1) in Lemma 6: Relation of strings w, w', and w''

- If $|w| = 1$, from (1') and (2'), the prefix of p_2 is cwd and wd. Thus, we have $w = c = d$.
- If $|w| \geq 2$, then from (1') and (2'), cwd and wd are prefixes of p_2. Thus, we have $cw = wd$. From Proposition 2, $c = d$ holds.

All of these cases do not meet $b = a$, $b \neq d$, and $c \notin \{a, d\}$. Therefore, if $b = a$, $b \neq d$, and $c \notin \{a, d\}$ are satisfied, case



(7) (8)

Fig. 14 Let $\Sigma = \{a, b, c, d, e, f, g\}$ and $p, q \in \mathcal{RP}$. We assume that the symbols in Σ are mutually distinct. The figures (7) and (8) express two cases $D = \{ya, ac, dy\}$ and $D = \{ya, bd, dy\}$ in Lemmas 6 and 7, respectively. In these cases, if $p\{x := r\} \preceq q$ for all $r \in D$ and D is maximally generalized on (p, q), then $p\{x := xy\} \preceq q$ holds.

(6-2) is also impossible.

(6-3) Case of $q = q_1 A q_2$, where $A = dy_1ac$: For $q = q_1dy_1acq_2$, the following conditions must be satisfied for $y'_1, y''_1 \in X$:

順序が他のものと異なる

ダミー/2　y_1　削除

(1) $p_1 \preceq$ ~~q_1d or~~ $p_1 \preceq q_1dy'_1$	(1') $p_2 \preceq \cancel{c}q_2$,
(2) $p_1 \preceq q_1dy_1$,	(2') $p_2 \preceq cq_2$,
(3) $p_1 \preceq q_1$,	(3') $p_2 \preceq acq_2$ or
	$p_2 \preceq y''_1acq_2$.

cq_2

For $p_1 \preceq q_1d$ in (1) and $p_2 \preceq acq_2$ in (3'), $p = p_1xp_2 \preceq q_1dxacq_2 \preceq q\{y_1 := x\}$ holds. From this, we have $p\{x := xy\} \preceq q\{y_1 := x\}$. This contradicts the assumption that $p\{x := xy\} \npreceq q$. Similarly, we can show that the other cases of (1) and (3') also contradict the assumption.

From the above, we conclude that if $p\{x := r\} \preceq q$ for all $r \in \{ya, bc, dy\}$ ($b = a$, $b \neq d$, and $c \notin \{a, d\}$), then $p\{x := xy\} \preceq q$ holds. □

すべて

The conditions in Lemmas 6 and 7 are illustrated in (7) and (8) in Fig. 14, respectively.

修正が必要かも

Lemma 7: Let Σ be an alphabet with $\sharp\Sigma \geq 3$ and let p, q be regular patterns on $\Sigma \cup X$. Let D be one of the following sets of regular patterns on $\Sigma \cup X$, where y is a variable symbol in X that does not appear in p and q.

不要かなと？　より

$$D = \{ya, bc, dy\}\ (b \notin \{a, d\}, \boxed{c \neq a}, \text{ and } c = d).$$

Then, if $p\{x := r\} \preceq q$ for all $r \in D$ and D is maximally generalized on (p, q), then $p\{x := xy\} \preceq q$.

Proof. The proof follows by reversing p and q and subsequently applying Lemma 6. □

When the conditions of Lemmas 5, 6, and 7 are not satisfied, counterexamples can be constructed as follows:

Proposition 4: Let Σ be an alphabet with $\sharp\Sigma \geq 3$. For a variable symbol y, let $D = \{ya, bc, dy\}$ ($b = a$ and $c = d$). There exist regular patterns p and q on $\Sigma \cup X$ such that $p\{x := r\} \preceq q$ for any $r \in D$, but $p\{x := xy\} \npreceq q$.

$c=d$, $a \neq c$, $b \neq d$, $b \neq c$
$a \neq d$, ということは $|\{a, b, c\}| \geq 3$ 以上が必要.

Proof. We provide an example to demonstrate this proposition. Let a, b, c, d, e be constant symbols in Σ, and let x, y, y_1, y_2 be variable symbols in X. Define the regular patterns p and q as follows:

$$p = eabcbcadabcbcada\boxed{x}bcadadabcbcadade,$$
$$q = y_1abcbcadabcbcadady_2 \quad (b = a \text{ and } c = d).$$

Obviously $p\{x := xy\} \npreceq q$ holds. For these p and q, the condition for Proposition 4 holds as follows (see also Fig. 12):

順序が分かりにくいかも

$$
\begin{aligned}
& p\ \{x := \boxed{ya}\} \\
&= (eabcbcadabcbcaday)abcadadabcbcadade \\
&= q\{y_1 := eabcbcadabcbcaday,\ y_2 := e\} \\
&\preceq q, \\
& p\ \{x := bc\} \\
&= (eabcbcad)abcbcadabcbcadad(abcbcadade) \\
&= q\{y_1 := eabcbcad,\ y_2 := abcbcadade\} \\
&\preceq q, \\
& p\ \{x := dy\} \\
&= eabcbcadabcbcadad(ybcadadabcbcadade) \\
&= q\{y_1 := e,\ y_2 := ybcadadabcbcadade\} \\
&\preceq q.
\end{aligned}
$$

□

3.4 $D = \{a_1b_1, a_2b_2, a_3y\}$ and $D = \{a_1b_1, a_2b_2, yb_3\}$

Lemma 8: Let Σ be an alphabet with $\sharp\Sigma \geq 3$ and p, q regular patterns on $\Sigma \cup X$. Let D be the following set of regular patterns on $\Sigma \cup X$, where y is a variable symbol in X that does not appear in p and q:

可能性として $\sharp\Sigma \geq 3$ となる
$a_1 \neq a_2$, $a_2 \neq a_3$, $a_3 \neq a_1$

$$D = \{a_1b_1, a_2b_2, a_3y\},$$
$$\text{where } a_i \neq a_j \text{ and } b_i \neq b_j\ (i \neq j, 1 \leq i, j \leq 3).$$

この条件より
$a_1 \neq a_3$
$a_2 \neq a_3$

Then, if $p\{x := r\} \preceq q$ holds for all $r \in D$ and D is maximally generalized on (p, q), then $p\{x := xy\} \preceq q$ holds.

Proof. We assume that $p\{x := xy\} \npreceq q$ holds. Since D is maximally generalized on (p, q), from the same argument as in the proof of Lemma 6, it is sufficient to consider the following five cases (8-1)–(8-5) of q: For $y_1 \in X$,

(8-1) $q = q_1a_1b_1wa_2b_2w'a_3y_1q_2$,
(8-2) $q = q_1a_1b_1b_2y_1q_2$ ($a_2 = b_1$ and $a_3 = b_2$),
(8-3) $q = q_1a_1b_1b_2wa_3y_1q_2$ ($b_1 = a_2$),
(8-4) $q = q_1a_3y_1wa_1b_1b_2q_2$ ($b_1 = a_2$),
(8-5) $q = q_1a_1b_1y_1wa_2b_2q_2$ ($b_1 = a_3$),

where no variable symbol appears in both w and w'.

(8-1) Case of $q = q_1a_1b_1wa_2b_2w'a_3y_1q_2$: The following conditions must be satisfied: For $y'_1 \in X$,

(1) $p_1 \preceq q_1$,	(1') $p_2 \preceq wa_2b_2w'a_3y_1q_2$,
(2) $p_1 \preceq q_1a_1b_1w$,	(2') $p_2 \preceq w'a_3y_1q_2$,
(3) $p_1 \preceq q_1a_1b_1wa_2b_2w'$,	(3') $p_2 \preceq q_2$ or $p_2 \preceq y'_1q_2$.

|w|=|w'|の場合が欠けているような気がします.

いきなり場合分けなので < 1行説明を入れたほうがいいかもしれません

$a_2b_2w' = a_2b_2w_1w$
$a_1b_1w =$

- $|w| + 1 = |w'|$: From (2) and (3), $a_1b_1wa_2b_2w'$ and a_1b_1w are suffixes of p_1. Since there exists a constant symbol w_1 such that $w' = w_1w$ and $b_2w_1w = a_1b_1w$ hold, then $b_2 = a_1$. Moreover, $wa_2b_2w'a_3$ and $w'a_3$ are prefixes of p_2 from (1') and (2'). Since there exists a constant symbol w_2 such that $w' = ww_2$ and $wa_2b_2 = ww_2a_3$ hold, then $b_2 = a_3$. Thus, $a_1 = a_3$ holds. This contradicts the assumption of $a_1 \neq a_3$.
- これより $|w| + 1 < |w'|$: From (2) and (3), $a_1b_1wa_2b_2w'$ and a_1b_1w are suffixes of p_1. Hence, a_1b_1 is suffixes of w. 本当? ? Moreover, $wa_2b_2w'a_3$ and $w'a_3$ are prefixes of p_2 from ~~(1') and (2'). Hence, there exist constant symbols~~ strings w_1 and w_2 such that $w' = w_1w$, $w' = ww_2$ and $|a_2b_2w_1| = |w_2a_3|+1$ hold. Thus, since the second-to-last symbol of w_1 is a_3, $a_1 = a_3$ holds. This contradicts the assumption of $a_1 \neq a_3$.

$wa_2b_2w_1w$
$\neq w_2a_3$

- $|w| = |w'| + 1$: From (1') and (2'), $wa_2b_2w'a_3$ and $w'a_3$ are prefixes of p_2. Since there exists a constant symbol w_1 such that $w = w'w_1$ and $w'w_1 = w'a_3$ hold, then $w_1 = a_3$ holds. Moreover, since $a_1b_1wa_2b_2w'$ and a_1b_1w are suffixes of p_1 from (2) and (3), there exists a constant symbol w_2 such that $w = w_2w'$ and $|w_1a_2b_2w'| = |a_1b_1w_2w'|$ hold. Hence, $w_1 = a_1$ holds. Thus, $a_1 = a_3$ holds. This contradicts the assumption of $a_1 \neq a_3$.
- $|w| > |w'| + 1$: Since $wa_2b_2w'a_3$ ($w'w_1$) and $w'a_3$ are prefixes of p_2 from (1') and (2'), there exists a constant string w_1 such that $w = w'w_1$ and the first symbol of w_1 is a_3. Moreover, since there exists a constant string w_2 such that $w = w_2w'$ and $|w_1a_2b_2| = |a_1b_1w_2|$ hold, a_1b_1 is a prefix of w_1. Thus, $a_3 = a_1$ holds. This contradicts the assumption of $a_1 \neq a_3$.

説明が少し足らない気がします.

$a_1b_1w = a_1b_1w_2w$
$w_1a_2b_2w$
$|w_1| = |w_2|$

(8-2) Case of $q = q_1a_1b_1b_2y_1q_2$ ($a_2 = b_1$ and $a_3 = b_2$): The following conditions must be satisfied: For $y'_1 \in X$,

$$
\begin{array}{ll}
(1)\ p_1 \preceq q_1, & (1')\ p_2 \preceq b_2y_1q_2, \\
(2)\ p_1 \preceq q_1a_1, & (2')\ p_2 \preceq y_1q_2, \\
(3)\ p_1 \preceq q_1a_1b_1, & (3')\ p_2 \preceq q_2 \text{ or } p_2 \preceq y'_1q_2.
\end{array}
$$

From (2) and (3), a_1b_1 and a_1 are suffixes of p_1. Hence, $b_1 = a_1$ holds. Thus, from the assumption of $b_1 = a_2$, $a_1 = a_2$ holds. This contradicts the assumption of $a_1 \neq a_2$.

(8-3) Case of $q = q_1a_1b_1b_2wa_3y_1q_2$ ($b_1 = a_2$): The following conditions must be satisfied: For $y'_1 \in X$,

$$
\begin{array}{ll}
(1)\ p_1 \preceq q_1, & (1')\ p_2 \preceq b_2wa_3y_1q_2, \\
(2)\ p_1 \preceq q_1a_1, & (2')\ p_2 \preceq wa_3y_1q_2, \\
(3)\ p_1 \preceq q_1a_1b_1b_2w, & (3')\ p_2 \preceq q_2 \text{ or } p_2 \preceq y'_1q_2.
\end{array}
$$

- $|w| = 0$: From (2) and (3), a_1 and $a_1b_1b_2$ are suffixes of p_1. Hence, $a_1 = b_2$ holds. Moreover, since b_2a_3 and a_3 is prefixes of p_2, $b_2 = a_3$ holds. Thus, $a_1 = a_3$ holds. This contradicts the assumption of $a_1 \neq a_3$.
- $|w| \geq 1$: Since a_1 and $a_1b_1b_2w$ are suffixes of p_1 from (2) and (3), the last symbol of w is a_1. Moreover, since b_2wa_3 and wa_3 are prefixes of p_2 from (1') and (2'), the last symbol of w is a_3. Thus, $a_1 = a_3$ holds. This contradicts the assumption of $a_1 \neq a_3$.

(8-4) Case of $q = q_1a_3y_1wa_1b_1b_2q_2$ ($b_1 = a_2$): The following conditions must be satisfied: For $y'_1 \in X$,

$$
\begin{array}{ll}
(1)\ p_1 \preceq q_1, & (1')\ p_2 \preceq wa_1b_1b_2q_2 \text{ or} \\
 & \quad\ p_2 \preceq y'_1wa_1b_1b_2q_2, \\
(2)\ p_1 \preceq q_1a_3y_1w, & (2')\ p_2 \preceq b_2q_2, \\
(3)\ p_1 \preceq q_1a_3y_1wa_1, & (3')\ p_2 \preceq q_2.
\end{array}
$$

From (3), there exist regular patterns p'_1 and p''_1 such that $p_1 = p'_1p''_1$, $p'_1 \preceq q_1a_3$, and $p''_1 \preceq y_1wa_1$ hold. Hence, if $p_2 \preceq wa_1b_1b_2q_2$ of (1') holds, since $p = p_1xp_2 = p'_1p''_1xp_2 \preceq q_1a_3p''_1xwa_1b_1b_2q_2 = q\{y_1 := p''_1x\}$, then $p \preceq q$ holds. Thus, this contradicts the assumption. Similarly, $p_2 \preceq y'_1wa_1b_1b_2q_2$ of (1') leads to a contradiction.

(8-5) Case of $q = q_1a_1b_1y_1wa_2b_2q_2$ ($b_1 = a_3$): The following conditions must be satisfied: For $y'_1 \in X$,

$$
\begin{array}{ll}
(1)\ p_1 \preceq q_1, & (1')\ p_2 \preceq y_1wa_2b_2q_2, \\
(2)\ p_1 \preceq q_1a_1, & (2')\ p_2 \preceq wa_2b_2q_2 \text{ or} \\
 & \quad\ p_2 \preceq y'_1wa_2b_2q_2, \\
(3)\ p_1 \preceq q_1a_1b_1y_1w, & (3')\ p_2 \preceq q_2.
\end{array}
$$

Let $q'_1 = q_1a_1b_1$, $q'_2 = y_1w$, $q'_3 = a_2b_2q_2$. From (3), $p_1 \preceq q'_1q'_2$ holds, and from (1'), $p_2 \preceq q'_2q'_3$ holds. Since q'_2 contains a variable symbol y_1, $p \preceq q$ holds from Theorem 2. This contradicts the assumption. □

Lemma 9: Let Σ be an alphabet with $\sharp\Sigma \geq 3$ and p, q regular patterns on $\Sigma \cup X$. Let D be the following set of regular patterns on $\Sigma \cup X$, where y is a variable symbol in X that does not appear in p and q:

$$D = \{a_1b_1, a_2b_2, yb_3\},$$
$$\text{where } a_i \neq a_j \text{ and } b_i \neq b_j\ (i \neq j, 1 \leq i, j \leq 3).$$

これより
$p\{x := a_1y\} \not\preceq q$
$p\{x := yb_1\} \not\preceq q$
$p\{x := a_2x\} \not\preceq q$
$p\{x := yb_2\} \not\preceq q$
⇓
$b_3 \neq b_1$, $b_2 \neq b_1$

Then, if $p\{x := r\} \preceq q$ for all $r \in D$ and D is maximally generalized on (p, q), then $p\{x := xy\} \preceq q$.

Proof. The proof follows by reversing p and q and subsequently applying Lemma 8. □

3.5 $D = \{a_1b_1, a_2b_2, a_3b_3\}$

In Lemma 14 (ii) of [4], they stated that, when $\sharp\Sigma \geq 3$, for regular patterns p, q, if $p\{x := r\} \preceq q$ for any $r \in D$, then $p\{x := xy\} \preceq q$ holds, where $D = \{a_1b_1, a_2b_2, a_3b_3\}$ ($a_i \neq a_j$ and $b_i \neq b_j$ for each i, j ($i \neq j, 1 \leq i, j \leq 3$)). Unfortunately, there exist the following counterexamples of Lemma 14 (ii) of [4].

CKに関する記述なし

Example 2: Assume that $a_1 = b_2$ and $a_3 = b_1$ hold.

(1) Let $p = ca_1x'a_3c$ and $q = xa_1a_3y$. It is clear that $p\{x := xy\} \not\preceq q$ holds. However, we can see that $p\{x' := a_1b_1\} \preceq q$, $p\{x' := a_2b_2\} \preceq q$ and $p\{x' := a_3b_3\} \preceq q$

Pでしょうか?

定数が定数に対応

$a_1=b_2,\ a_3=b_1$ $q = xa_1a_3y = xa_1b_1y$

$ca_1ybb_1c = ca_1ya_1b_1c$

hold, since $p\{x' := a_1b_1\} = ca_1a_1b_1a_3c = q\{x := ca_1, y := a_3c\}$, $p\{x' := a_2b_2\} = ca_1a_2b_2a_3c = q\{x := ca_1a_2, y := c\}$ and $p\{x' := a_3b_3\} = ca_1a_3b_3a_3c = q\{x := c, y := b_3a_3c\}$ hold.

/が無いと異なる

(2) Let $p = cb_2a_1b_1b_2x'a_1b_1b_2a_3c$ and $q = xb_2a_1b_1b_2a_3y$. It is clear that $p\{x' := xy\} \not\preceq q$ holds. However, we have $p\{x' := a_1b_1\} \preceq q$, $p\{x' := a_2b_2\} \preceq q$, and $p\{x' := a_3b_3\} \preceq q$, since $p\{x' := a_1b_1\} = cb_2a_1b_1b_2a_1b_1a_1b_1b_2a_3c = q\{x := cb_2a_1b_1, y := b_2a_3c\}$, $p\{x' := a_2b_2\} = cb_2a_1b_1b_2a_2b_2a_1b_1b_2a_3c = q\{x := cb_2a_1b_1b_2a_2, y := c\}$, and $p\{x' := a_3b_3\} = cb_2a_1b_1b_2a_3b_3a_1b_1b_2a_3c = q\{x := c, y := b_3a_1b_1b_2a_3c\}$ hold.

The conditions in Lemmas 8, 9, and 10 are illustrated in the cases (9), (10), and (11) in Fig. 15.

Lemma 10: Let Σ be an alphabet with $\sharp\Sigma \geq 3$ and p, q regular patterns on $\Sigma \cup X$. Let D be the following set of regular patterns on $\Sigma \cup X$, where y is a variable symbol in X that does not appear in p and q:

maximally generalized から明らかでは?

Example 9 と異なる点

$$D = \{a_1b_1, a_2b_2, a_3b_3\},$$

where $a_i \neq a_j$ and $b_i \neq b_j$ $(i \neq j, 1 \leq i, j \leq 3)$.

Then, if $p\{x := r\} \preceq q$ for all $r \in D$ and D is maximally generalized on (p, q), then $p\{x := xy\} \preceq q$.

Proof. We assume that $p\{x := xy\} \not\preceq q$ holds. Since D is maximally generalized on (p, q), it is sufficient to consider the following four cases (10-1)-(10-4) of q for some regular patterns q_1, q_2 and some constant strings w, w' ($|w| \geq 0$ and $|w'| \geq 0$):

(10-1) $q = q_1a_1b_1wa_2b_2w'a_3b_3q_2$,
(10-2) $q = q_1a_1b_1a_3b_3q_2$ ($b_1 = a_2$ and $a_3 = b_2$),
(10-3) $q = q_1a_1b_1b_2wa_3b_3q_2$ ($b_1 = a_2$),
(10-4) $q = q_1a_1b_1wa_2b_2b_3q_2$ ($b_2 = a_3$).

(10-1) Case of $q = q_1a_1b_1wa_2b_2w'a_3b_3q_2$: The following conditions must be satisfied:

(1) $p_1 \preceq q_1$,	(1') $p_2 \preceq wa_2b_2w'a_3b_3q_2$,
(2) $p_1 \preceq q_1a_1b_1w$,	(2') $p_2 \preceq w'a_3b_3q_2$,
(3) $p_1 \preceq q_1a_1b_1wa_2b_2w'$,	(3') $p_2 \preceq q_2$.

From

- $|w| = |w'|$: From (2) and (3), $a_1b_1wa_2b_2w'$ and a_1b_1w are suffixes of p_1. Then, $a_1b_1w = a_2b_2w'$. Hence, $a_1b_1 = a_2b_2$. This contracts the assumption of $a_1 \neq a_2$ and $b_1 \neq b_2$.
- 前に $|w| + 1 = |w'|$: The strings $wa_2b_2w'a_3b_3$ and $w'a_3b_3$ are prefixes of p_2. If there exists a constant symbol w_1 such that $w'a_3b_3 = ww_1a_3b_3$, then b_2 and a_3 are the same symbol from $wa_2b_2 = ww_1a_3$. from (2) and (3), $a_1b_1wa_2b_2w'$ and a_1b_1w are suffixes of p_1. Then, there exists a constant symbol w_2 such that $w' = w_2w$, then b_2 and a_1 are the same symbol from $b_2w_2w = a_1b_1w$. Hence, from $b_2 = a_3$, a_3 and a_1 are same symbol. This contradicts the assumption of $a_3 \neq a_1$.
- $|w| + 1 < |w'|$: From (2) and (3), $a_1b_1wa_2b_2w'$ and

$wa_1b_2w_1w$ a_3b_3

$ww_2a_3b_3$

a_1b_1w are suffixes of p_1. If there exists a constant string w_1 ($|w_1| \geq 2$) such that $w' = w_1w$, then a_1b_1 is a suffix of w_1. From conditions (1′) and (2′), $wa_2b_2w'a_3b_3$ and $w'a_3b_3$ are prefixes of p_2. If there exist constant strings w_1 and w_2 such that $w' = w_1w = ww_2$ holds, then [redacted] a_3b_3 are a suffixes of w_1 from $|w_1| = |w_2|$ and $ww_2a_3b_3 = wa_2b_2w_1$. Hence, $a_1b_1 = a_3b_3$. This contradicts the assumption of $a_1 \neq a_3$ and $b_1 \neq b_3$.
- $|w| > |w'|$: We can prove the contradiction in a similar way as $|w| \leq |w'|$.

(10-2) Case of $q = q_1a_1b_1a_3b_3q_2$ ($b_1 = a_2$ and $a_3 = b_2$): The following conditions must be satisfied:

(1) $p_1 \preceq q_1$,	(1') $p_2 \preceq a_3b_3q_2$,
(2) $p_1 \preceq q_1a_1$,	(2') $p_2 \preceq b_3q_2$,
(3) $p_1 \preceq q_1a_1b_1$,	(3') $p_2 \preceq q_2$.

From (2) and (3), since a_1b_1 and a_1 are suffixes of p_1, $b_1 = a_1$ holds. From the assumption of $b_1 = a_2$, the equation $a_1 = a_2$ holds. This contradicts the assumption of $a_1 \neq a_2$.

(10-3) Case of $q = q_1a_1b_1b_2wa_3b_3q_2$ ($b_1 = a_2$): The following conditions must be satisfied:

(1) $p_1 \preceq q_1$,	(1') $p_2 \preceq b_2wa_3b_3q_2$,
(2) $p_1 \preceq q_1a_1$,	(2') $p_2 \preceq wa_3b_3q_2$,
(3) $p_1 \preceq q_1a_1b_1b_2w$,	(3') $p_2 \preceq q_2$.

- $|w| = 0$: From (2) and (3), a_1 and $a_1b_1b_2$ are suffixes of p_1. Moreover, from (1') and (2'), $b_2a_3b_3$ and a_3b_3 are prefixes of p_2. Since $b_2 = a_1$ and $b_2a_3 = a_3b_3$, $a_1 = a_3$ holds. This contradicts the assumption of $a_1 \neq a_3$.
- $|w| \geq 1$: From (2) and (3), a_1 and $a_1b_1b_2w$ are suffixes of p_1. Hence, the last symbol of w is a_1. Moreover, $b_2wa_3b_3$ and wa_3b_3 are prefixes of p_2 from (1') and (2'). Hence, the last symbol of w is a_3. Therefore, $a_1 = a_3$ holds. This contradicts the assumption of $a_1 \neq a_3$.

(10-4) Case of $q = q_1a_1b_1wa_2b_2b_3q_2$ ($b_2 = a_3$): The following conditions must be satisfied:

(1) $p_1 \preceq q_1$,	(1') $p_2 \preceq wa_2b_2b_3q_2$,
(2) $p_1 \preceq q_1a_1b_1w$,	(2') $p_2 \preceq b_3q_2$,
(3) $p_1 \preceq q_1a_1b_1wa_2$,	(3') $p_2 \preceq q_2$.

- $|w| = 0$: From (2) and (3), a_1b_1 and $a_1b_1a_2$ are suffixes of p_1. And from (1') and (2'), $a_2b_2b_3$ and b_3 are prefixes of p_2. Since $b_1 = a_2$ and $a_2 = b_3$, then $b_1 = b_3$ holds. This contradicts the assumption of $b_1 \neq b_3$.
- $|w| \geq 1$: Since a_1b_1w and $a_1b_1wa_2$ are suffixes of p_1 from (2) and (3), the first symbol of w is b_1. Moreover, since $wa_2b_2b_3$ and b_3 are prefixes of p_2 from (1') and (2'), the first symbol of w is b_3. Therefore, $b_1 = b_3$ holds. This contradicts the assumption of $b_1 \neq b_3$. □

$a_1 = a_3$ などとかけば ページの省略になるかも

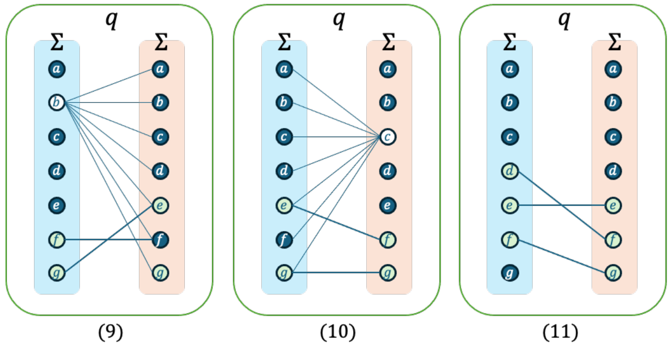


Fig. 15 Let $\Sigma = \{a,b,c,d,e,f,g\}$ and $p, q \in \mathcal{RP}$. We assume that the symbols in Σ are mutually distinct. The figures (9), (10,) and (11) express cases $D = \{a_1b_1, a_2b_2, a_3y\}$, $D = \{a_1b_1, a_2b_2, yb_3\}$, and $D = \{a_1b_1, a_2b_2, a_3b_3\}$ in Lemmas 8, 9, and 10, respectively, where $a_i \neq a_j$ and $b_i \neq b_j$ for each i, j $(i \neq j, 1 \leq i, j \leq 3)$. In these cases, if $p\{x := r\} \preceq q$ for all $r \in D$ and D is maximally generalized on (p, q), then $p\{x := xy\} \preceq q$ holds.

$a_1=f$, $b_1=f_1$
$a_2=g$, $b_2=e$
$a_3=c$ でしょうか?
$\Sigma=\{a,b,\cdots,g\}$との関係性が少し読みづらいかも..

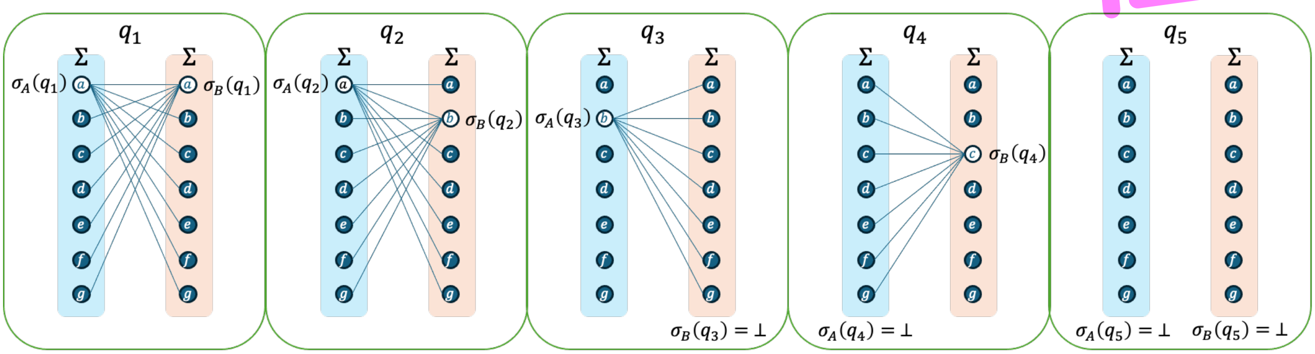


Fig. 16 Let $\Sigma = \{a,b,c,d,e,f,g\}$, $Q = \{q_1, q_2, q_3, q_4, q_5\}$. We set $A(q_1) = \{a\}$ and $B(q_1) = \{a\}$, and then $\sigma_A(q_1) = a$ and $\sigma_B(q_1) = a$, and so on. For each regular pattern q_i $(i = 1, \ldots, 5)$, we represent a string $w \in \Sigma \cdot \Sigma$ satisfying that $p\{x := w\} \preceq q_i$ by the edge between the left (first) and right (second) symbols of w. For example, the leftmost figure shows that $p\{x := ay\} \preceq q_1$ and $p\{x := ya\} \preceq q_1$ for a variable symbol y. We note that these figures may contain more edges than those illustrated. From these figures, we get $\ell_A = 1$, $\ell_B = 0$, and $Q^{(\perp,\perp)} = \{q_5\}$, $Q^{(\perp,\cdot)} = \{q_4\}$, $Q^{(\cdot,\perp)} = \{q_3\}$, $Q^{(\cdot,\cdot)} = \{q_1, q_2\}$.

S_2よりも条件が弱い.

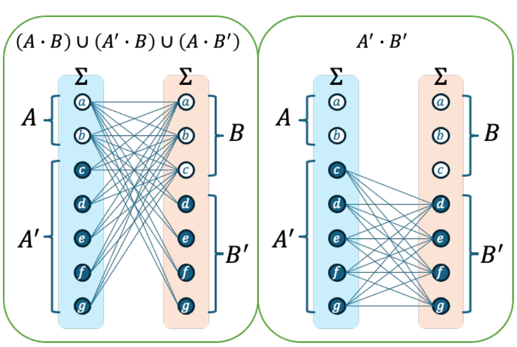


Fig. 17 In the left figure, we aggregate all of the edges appearing in Fig. 16. For all $w = a'b' \in A' \cdot B'$, there must be a regular pattern q_i $(1 \leq i \leq 5)$ that satisfies that $p\{x := w\} \preceq q_i$.

3.6 Characteristic sets for finite union of regular patterns

Lemma 11: Let k be an integer with $k \geq 1$. Let Σ be an alphabet with $\sharp\Sigma = k + 2$. Let $p \in \mathcal{RP}$ in which a variable symbol x appears, and let $Q \in \mathcal{RP}^k$. If for any string $w \in \Sigma^*$ with $|w| = 2$, there exists a regular pattern $q_w \in Q$ such that $p\{x := w\} \preceq q_w$ holds, then there exists a regular pattern $q \in Q$ such that $p\{x := xy\} \preceq q$ holds, where y is a variable symbol that does not appear in q.

Proof. Without loss of generality, we suppose that $\sharp Q = k$ holds. Otherwise, for some regular pattern q already in Q, we can add a new regular pattern q' equivalent to q, i.e., $q' \equiv q$, to Q repeatedly until $\sharp Q = k$ is satisfied. For any $q \in Q$, we define the sets $A(q), B(q) \subseteq \Sigma$ as follows:

$$A(q) = \{a \in \Sigma \mid p\{x := ay\} \preceq q,\ y \in X\},$$
$$B(q) = \{b \in \Sigma \mid p\{x := yb\} \preceq q,\ y \in X\}.$$

If there exists $q \in Q$ such that $|A(q)| \geq 2$ or $|B(q)| \geq 2$, from Lemma 4, $p\{x := xy\} \preceq q$ holds. Below, we suppose that $|A(q)| \leq 1$ and $|B(q)| \leq 1$. Let $\perp$ be a constant symbol that is not a member in Σ. We define the functions $\sigma_A : Q \to \Sigma \cup \{\perp\}$ and $\sigma_B : Q \to \Sigma \cup \{\perp\}$ as follows:

ℓ_A, ℓ_Bの意味..

線が1本不足しているような

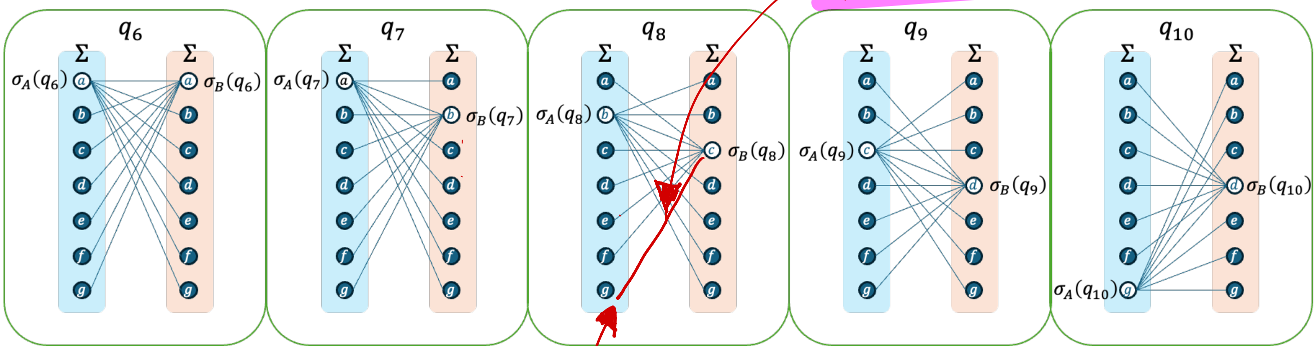


Fig. 18 Let $\Sigma = \{a, b, c, d, e, f, g\}$, $Q = \{q_6, q_7, q_8, q_9, q_{10}\}$. From these figures, we get $\ell_A = 1$, $\ell_B = 1$, $Q^{(\perp,\perp)} = Q^{(\perp,\cdot)} = Q^{(\cdot,\perp)} = \emptyset$, and $Q^{(\cdot,\cdot)} = Q$.

$\sigma_A^{-1}(a) \subset Q$
$\sigma_B^{-1}(a) \subset Q$

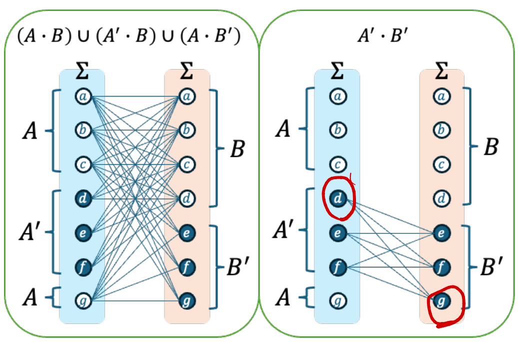


Fig. 19 In the left figure, we aggregate all of the edges appearing in Fig. 18. From Fig. 18 and this right figure, we get $Q_1^{(\cdot,\cdot)} = \{q_6, q_7, q_8, q_9\}$ and $Q_2^{(\cdot,\cdot)} = \{q_{10}\}$. From Proposition 4, even if the string $dg \in A' \cdot B'$ satisfies $p\{x := gd\} \preceq q_{10}$, it does not imply that $p\{x := xy\} \preceq q_{10}$.

$$\sigma_A(q) = \begin{cases} a & \text{if } A(q) = \{a\}, \\ \perp & \text{if } A(q) = \emptyset. \end{cases}$$

$$\sigma_B(q) = \begin{cases} b & \text{if } B(q) = \{b\}, \\ \perp & \text{if } B(q) = \emptyset. \end{cases}$$

The inverse functions of σ_A and σ_B are denoted by σ_A^{-1} and σ_B^{-1}, respectively. That is, for $a, b \in \Sigma \cup \{\perp\}$, let $\sigma_A^{-1}(a) = \{q \in Q \mid \sigma_A(q) = a\}$ and $\sigma_B^{-1}(b) = \{q \in Q \mid \sigma_B(q) = b\}$. We give an example in Fig. 16.

A and B denotes the following subsets of Σ:

$A \subset \Sigma$
$B \subset \Sigma$

$$A = \bigcup_{q \in Q \setminus \sigma_A^{-1}(\perp)} A(q), \quad B = \bigcup_{q \in Q \setminus \sigma_B^{-1}(\perp)} B(q).$$

Then, let $A' = \Sigma \setminus A$ and $B' = \Sigma \setminus B$. For any $a, b \in \Sigma$, we use the following notations:

$$\ell_A = \sum_{a \in A} (\sharp\sigma_A^{-1}(a) - 1), \quad \ell_B = \sum_{b \in B} (\sharp\sigma_B^{-1}(b) - 1).$$

These ℓ_A and ℓ_B represent the numbers of excess duplicate symbols in A and B. We easily see the following claim:

Claim 1.

(i) $\sharp A + \sharp A' = \sharp B + \sharp B' = k + 2$,

(ii) $\sharp A + \ell_A + \sharp\sigma_A^{-1}(\perp) = \sharp B + \ell_B + \sharp\sigma_B^{-1}(\perp) = k.$

定数文字
Lemma 21は

Since $\sharp\Sigma = k + 2$ and $\sharp Q = k$, $\sharp A' \geq 2$ and $\sharp B' \geq 2$ hold. We partition Q into the following subsets:

$q \in Q^{(\cdot,\cdot)}$

$$Q^{(\perp,\perp)} = \sigma_A^{-1}(\perp) \cap \sigma_B^{-1}(\perp),$$
$$Q^{(\perp,\cdot)} = \sigma_A^{-1}(\perp) \cap (Q \setminus \sigma_B^{-1}(\perp)),$$
$$Q^{(\cdot,\perp)} = (Q \setminus \sigma_A^{-1}(\perp)) \cap \sigma_B^{-1}(\perp),$$
$$Q^{(\cdot,\cdot)} = (Q \setminus \sigma_A^{-1}(\perp)) \cap (Q \setminus \sigma_B^{-1}(\perp)).$$

$Q^{(\cdot,\cdot)} \subset Q$ subset

From the condition of this lemma, for any string $w \in \Sigma^*$ with $|w| = 2$, there exists a regular pattern $q_w \in Q$ such that $p\{x := w\} \preceq q_w$ holds. In particular, for $w = a'b' \in A' \cdot B'$, we must have $q_w \in Q$ that satisfies that $p\{x := w\} \preceq q_w$ (Fig. 17). It is easy to see that if $w \in (A \cdot B) \cup (A' \cdot B) \cup (A \cdot B')$, there exists a regular pattern $q_w \in Q^{(\perp,\cdot)} \cup Q^{(\cdot,\perp)} \cup Q^{(\cdot,\cdot)}$ such that $p\{x := w\} \preceq q_w$ holds. We have the following two claims:

$A' \cdot B'$より maximally generalized pが5

Claim 2. If there exist $q \in Q^{(\perp,\perp)}$ and distinct 5 strings $w_i \in A' \cdot B'$ $(1 \leq i \leq 5)$ such that $p\{x := w_i\} \preceq \boxed{q}$ holds $(1 \leq i \leq 5)$, then $p\{x := xy\} \preceq q$ holds.

共通

Proof of Claim 2. Let $W = \{a_1b_1, \ldots, a_5b_5\} \subset A' \cdot B'$. Because, for any i $(1 \leq i \leq 5)$, $|W \cap \{a_ic \mid c \in \Sigma\}| \leq 2$ and $|W \cap \{cb_i \mid c \in \Sigma\}| \leq 2$, it can be proven that there are 3 strings $a_{i_1}b_{i_1}, a_{i_2}b_{i_2}, a_{i_3}b_{i_3} \in W$ such that $a_{i_j} \neq a_{i_{j'}}$ and $b_{i_j} \neq b_{i_{j'}}$ for any $i_j, i_{j'}$ $(i_j \neq i_{j'}, 1 \leq j, j' \leq 3)$. Therefore, from Lemma 10, this claim holds. (*End of Proof of Claim*)

2

Claim 3. If there exist $q \in Q^{(\perp,\cdot)} \cup Q^{(\cdot,\perp)}$ and distinct 3 strings $w_i \in A' \cdot B'$ $(1 \leq i \leq 3)$ such that $p\{x := w_i\} \preceq \boxed{q}$ holds $(1 \leq i \leq 3)$, then $p\{x := xy\} \preceq q$ holds.

重要

Proof of Claim 3. Let $W = \{a_1b_1, a_2b_2, a_3b_3\} \subset A' \cdot B'$. Because, for any i $(1 \leq i \leq 3)$, $|W \cap \{a_ic \mid c \in \Sigma\}| \leq 2$ and $|W \cap \{cb_i \mid c \in \Sigma\}| \leq 2$, it can be proven that there are 2 strings $a_{i_1}b_{i_1}, a_{i_2}b_{i_2} \in W$ such that $a_{i_1} \neq a_{i_2}$ and $b_{i_1} \neq b_{i_2}$. Therefore, from Lemmas 8 and 9, this claim holds. (*End of Proof of Claim*)

3

If there exist a regular pattern $q \in Q^{(\perp,\perp)} \cup Q^{(\perp,\cdot)} \cup Q^{(\cdot,\perp)}$ and enough strings $w \in A' \cdot B'$ such that either of the conditions of *Claims* 2 and 3 is satisfied, this lemma holds. Then, we assume that it is not the case.

ちょっと難しい

Assumption 1. There is no regular pattern $q \in Q^{(\perp,\perp)}$ and 5 strings $w \in A' \cdot B'$ such that the condition of *Claim* 2 is satisfied and there is no regular pattern $q \in Q^{(\perp,\cdot)} \cup Q^{(\cdot,\perp)}$ and 3 strings $w \in A' \cdot B'$ such that the condition of *Claim* 3 is satisfied.

Let $\mathcal{L}_1 = \sharp\{w \in A' \cdot B' \mid \exists q \in Q^{(\perp,\perp)} \cup Q^{(\perp,\cdot)} \cup Q^{(\cdot,\perp)}$ s.t. $p\{x := w\} \preceq q\}$. Under *Assumption* 1, each $q \in Q^{(\perp,\perp)}$ has at most 4 strings $w \in A' \cdot B'$ such that the condition of *Claim* 2 is satisfied, and each $q \in Q^{(\perp,\cdot)} \cup Q^{(\cdot,\perp)}$ has at most 2 strings $w \in A' \cdot B'$ such that the condition of *Claim* 3 is satisfied. Then, by *Claim* 1,

$$\begin{aligned}
\mathcal{L}_1 &\le 4\sharp Q^{(\perp,\perp)} + 2\sharp Q^{(\perp,\cdot)} + 2\sharp Q^{(\cdot,\perp)} \\
&= 2(\sharp Q^{(\perp,\perp)} + \sharp Q^{(\perp,\cdot)}) + 2(\sharp Q^{(\perp,\perp)} + \sharp Q^{(\cdot,\perp)}) \\
&= 2\sharp\sigma_A^{-1}(\perp) + 2\sharp\sigma_B^{-1}(\perp) \\
&= 2(k - \sharp A - \ell_A) + 2(k - \sharp B - \ell_B) \\
&= 2(\sharp A' - \ell_A - 2) + 2(\sharp B' - \ell_B - 2) \\
&= 2(\sharp A' + \sharp B') - 2(\ell_A + \ell_B) - 8.
\end{aligned}$$

あとで A'が?

$A \subset \Sigma$
$B \subset \Sigma$

$\sigma_B(q) \in A'$
$\sigma_B(q) \notin A$

Next, we partition $Q^{(\cdot,\cdot)}$ into the following two subsets:

$$Q_1^{(\cdot,\cdot)} = \{q \in Q^{(\cdot,\cdot)} \mid \sigma_A(q) \in B \text{ or } \sigma_B(q) \in A\},$$
$$Q_2^{(\cdot,\cdot)} = \{q \in Q^{(\cdot,\cdot)} \mid \sigma_A(q) \in B' \text{ and } \sigma_B(q) \in A'\}.$$

Qのsubset $\sigma_A(q) \in B$

We show the next two claims on $Q_1^{(\cdot,\cdot)}$ and $Q_2^{(\cdot,\cdot)}$:

Claim 4. If there exist $q \in Q_1^{(\cdot,\cdot)}$ and a string $a'b' \in A' \cdot B'$ such that $p\{x := a'b'\} \preceq q$ holds, then $p\{x := xy\} \preceq q$ holds.

$a' \in A'$ maximally generalized

Proof of Claim 4. Suppose that both $\sigma_A(q) \in B$ and $\sigma_B(q) \in A$ hold. Then, since $a' \notin \{\sigma_A(q), \sigma_B(q)\} \subseteq A \cap B$ and $b' \notin \{\sigma_A(q), \sigma_B(q)\} \subseteq A \cap B$, from Lemma 5, $p\{x := xy\} \preceq q$ holds. Suppose that $\sigma_A(q) \in B$ and $\sigma_B(q) \in A'$. If $a' = \sigma_B(q)$, since $a' \in B$, $a' \neq b'$ holds. Since $\sigma_A(q) \in B$, $b' \neq \sigma_A(q)$ holds. That is, $a' = \sigma_B(q)$, $a' \neq \sigma_A(q)$, and $b' \notin \{\sigma_A(q), \sigma_B(q)\}$ hold. Therefore, from Lemmas 6 and 7, $p\{x := xy\} \preceq q$ holds. If $a' \neq \sigma_B(q)$, since $b' \neq \sigma_A(q)$, from Lemma 5, $p\{x := xy\} \preceq q$ holds. Similarly, the case that $\sigma_A(q) \in B'$ and $\sigma_B(q) \in A$ is proven. (*End of Proof of Claim* 4)

or b $b' = \sigma_A(q)$ $a' \notin \{\sigma_A(q), \sigma_B(q)\}$

Claim 5. If there exist $q \in Q_2^{(\cdot,\cdot)}$ and a string $a'b' \in A' \cdot B'$ such that $(a' \neq \sigma_B(q)$ or $b' \neq \sigma_A(q))$ and $p\{x := a'b'\} \preceq q$ hold, then $p\{x := xy\} \preceq q$ holds.

Proof of Claim 5. When $a' = b'$, since $\sigma_A(q) \neq \sigma_B(q)$, from Lemma 5, this claim holds. Similarly, when $a' \neq b'$, from Lemmas 5, 6, and 7, this holds. (*End of Proof of Claim* 5)

前の方がいいかも

We give an example in Fig. 18 and Fig. 19.

If there exist a regular pattern $q \in Q_2^{(\cdot,\cdot)}$ and a string $w \in A' \cdot B'$ such that the condition of *Claim* 5 is satisfied, this lemma holds. Then, we also assume that it is not the case.

この方が理解しやすいかも

Assumption 2. There is no $q \in Q_2^{(\cdot,\cdot)}$ and a string $a'b' \in A' \cdot B'$ such that the condition of *Claim* 5 is satisfied.

Let $\mathcal{L}_2 = \sharp\{a'b' \in A' \cdot B' \mid \exists q \in Q_2^{(\cdot,\cdot)}$ s.t. $p\{x := a'b'\} \preceq q\}$. For any $a'b' \in A' \cdot B'$ and $q \in Q_2^{(\cdot,\cdot)}$, if $a' = \sigma_B(q)$ and

$a' \neq b'$ $a' \in A'$, $b' \in B'$ $a' \notin A$, $b' \notin B$ $a' \in B$

$a' = \sigma_B(q)$
$b' = \sigma_A(q)$

何故 Proposition 4 が出てくる??

$b' = \sigma_A(q)$ hold (it is the condition of Proposition 4), by considering the duplicate numbers ℓ_A and ℓ_B, we have the following inequality:

$$\mathcal{L}_2 \le \min\{\sharp A' + \ell_B, \sharp B' + \ell_A\}.$$

不明

We show the last claim:

Claim 6. $\sharp A' \times \sharp B' - \mathcal{L}_1 - \mathcal{L}_2 \ge 2$.

Proof of Claim 6. First we prove the inequality when $\sharp A \le k-1$ and $\sharp B \le k-1$, i.e., $\sharp A' \ge 3$ and $\sharp B' \ge 3$ hold. Since $\mathcal{L}_2 \le \frac{1}{2}(\sharp A' + \sharp B' + \ell_A + \ell_B)$,

$$\begin{aligned}
&\sharp A' \times \sharp B' - \mathcal{L}_1 - \mathcal{L}_2 \\
&\ge \sharp A' \times \sharp B' - (2(\sharp A' + \sharp B') - 2(\ell_A + \ell_B) - 8) \\
&\quad - \frac{1}{2}(\sharp A' + \sharp B' + \ell_A + \ell_B) \\
&= \sharp A' \times \sharp B' - \frac{5}{2}(\sharp A' + \sharp B') + \frac{3}{2}(\ell_A + \ell_B) + 8 \\
&= (\sharp A' - \frac{5}{2})(\sharp B' - \frac{5}{2}) + \frac{3}{2}(\ell_A + \ell_B) + \frac{7}{4} \ge 2.
\end{aligned}$$

$\frac{25}{4} + \frac{25}{4} = \frac{25}{2}$

When $\sharp A = k$ and $\sharp B \le k$, i.e., $\sharp A' = 2$ and $\sharp B' \ge 2$ hold, since $\ell_A = 0$, $\mathcal{L}_1 \le 2\sharp B' - 2\ell_B - 4$ holds. Moreover, $\mathcal{L}_2 \le \min\{\sharp B', \ell_B + 2\}$ holds. From *Claim* 1, $\ell_B + 2 = k - \sharp\sigma_B^{-1}(\perp) - \sharp B = \sharp B' - \sharp\sigma_B^{-1}(\perp)$ holds. Therefore, $\mathcal{L}_2 \le \ell_B + 2$ holds. Thus,

+2が不足

$$\begin{aligned}
&\sharp A' \times \sharp B' - \mathcal{L}_1 - \mathcal{L}_2 \\
&\ge 2\sharp B' - (2\sharp B' - 2\ell_B - 4) - (\ell_B + 2) \\
&= \ell_B + 2 \ge 2.
\end{aligned}$$

Similarly, the case when $\sharp A \le k$ and $\sharp B = k$ is proven. (*End of Proof of Claim* 6)

Under *Assumptions* 1 and 2, from *Claim* 6, there exist at least two $w \in A' \cdot B'$ and a regular pattern $q \in Q_1^{(\cdot,\cdot)}$ such that the condition of *Claim* 4 is satisfied. Therefore, for such a regular pattern q, $p\{x := xy\} \preceq q$ holds. □

Lemma 12 (Sato et al.[4]): Let Σ be a finite alphabet with $\sharp\Sigma \ge 3$ and p, q regular patterns. If there exists a constant symbol $a \in \Sigma$ such that $p\{x := a\} \preceq q$ and $p\{x := xy\} \preceq q$, then $p \preceq q$ holds, where y is a variable symbol that does not appear in q.

From Lemma 11 and Lemma 12, we have the following theorem.

Theorem 4: Let $k \ge 3$, $\sharp\Sigma \ge 2k-1$, $P \in \mathcal{RP}^+$ and $Q \in \mathcal{RP}^k$. Then, the following (i),(ii) and (iii) are equivalent:

$$\text{(i) } S_2(P) \subseteq L(Q), \text{ (ii) } P \sqsubseteq Q, \text{ (iii) } L(P) \subseteq L(Q).$$

Proof. It is clear that (ii) implies (iii) and (iii) implies (i). From Theorem 3, if $\sharp\Sigma \ge 2k+1$, then (i) implies (ii). Let $\sharp Q = k$, $p \in P$, $\sharp\Sigma = 2k-1$ or $2k$. Then, we show that (i) implies (ii). It suffices to show that $S_2(p) \subseteq L(Q)$ implies $P \sqsubseteq Q$ for any regular pattern $p \in \mathcal{RP}$. The proof is done by mathematical induction on n, where n is the number of

Since $a' \notin \{\sigma_A(q), \sigma_B(q)\}$ and $b' \notin \{\sigma_A(q), \sigma_B(q)\}$ がいいかも

$L_2 = \#\{a'b' \in A' \cdot B' \mid \exists q_2 \in Q_2^{(\cdot,\cdot)} \text{ s.t. } P\{x := a'b'\} \leq q\}$

基本的には

$L_2 \leq \#A' \cdot \#B'$ は明らか.

$Q_2^{(\cdot,\cdot)} = \{q \in Q^{(\cdot,\cdot)} \mid \sigma_A(q) \in B' \text{ and } \sigma_B(q) \in A'\}$

$\sigma_A(q) \in B'$ より

$\alpha = \sigma_A(q)$ とすると $P\{x := \alpha y\} \leq q$ and $\alpha \in A$

$L_2 \leq \min\{\#A' + l_B, \#B' + l_A\}$

$\#A \leq k$ は何故

claim 1 の (ii) より明らか.

$\#B$ も同様

variable symbols appears in p.

In case $n = 0$, $S_2(p) = \{p\}$ holds. By (i), we have $\{p\} = L(Q)$. Thus, $p \preceq q$ for some $q \in Q$.

For $n \geq 0$, we assume that it is valid for any regular pattern p with n variable symbols. Let p be a regular pattern such that $n + 1$ variable symbols appear in p and $S_2(p) \subseteq L(Q)$. We assume that $p \not\sqsubseteq Q$, that is, $p \not\preceq q_i$ for any $i \in \{1, \ldots, k\}$. Let $Q = \{q_1, \ldots, q_k\}$ and p_1, p_2 regular patterns, x a variable symbol with $p = p_1 x p_2$. For $a, b \in \Sigma$, let $p_a = p\{x := a\}$ and $p_{ab} = p\{x := ab\}$. Both p_a and p_{ab} have n variable symbols, respectively. Thus, $S_2(p_a) \subseteq L(Q)$ and $S_2(p_{ab}) \subseteq L(P)$ hold. By the induction hypothesis, there exist $i, i' \in \{1, \ldots, k\}$ such that $p_a \preceq q_i$ and $p_{ab} \preceq q_{i'}$. Let $D_i = \{a \in \Sigma \mid p\{x := a\} \preceq q_i\}$ $(i = 1, \ldots, k)$. We assume that $\sharp D_i \geq 3$ for some $i \in \{1, \ldots, k\}$. By Lemma 2, we have $p \preceq q_i$. This contradicts the assumption. Thus, we have $\sharp D_i \leq 2$ for any $i \in \{1, \ldots, k\}$. If $\sharp\Sigma = 2k - 1$, then $\sharp D_i = 2$ or $\sharp D_i = 1$ for any $i \in \{1, \ldots, k\}$. Moreover, If $\sharp\Sigma = 2k$, then $\sharp D_i = 2$ for any $i \in \{1, \ldots, k\}$. Since $k \geq 3$, $2k + 1 \geq k + 2$ holds. By Lemma 11, there exists $i \in \{1, \ldots, k\}$ such that $p\{x := xy\} \preceq q_i$. Therefore, by Lemma 12, we have $p \preceq q_i$. This contradicts the assumption. Thus, (i) implies (ii). □

From Theorem 4, the following Corollary 2 holds.

Corollary 2: Let $k \geq 3$, $\sharp\Sigma \geq 2k - 1$ and $P \in \mathcal{RP}^+$. Then, $S_2(P)$ is a characteristic set for $L(P)$ within $\mathcal{RPL}^k$.

Lemma 13 (Sato et al.[4]): Let $k \geq 3$ and $\sharp\Sigma \leq 2k - 2$. Then, $\mathcal{RP}^k$ does not have compactness with respect to containment.

Proof. Let $\Sigma = \{a_1, \ldots, a_{k-1}, b_1, \ldots, b_{k-1}\}$ and p, q_i regular patterns, $w_i \in \Sigma^*$ $(i = 1, \ldots, k - 1)$ defined in a similar way to Example 1. Let $q_k = x_1 a_1 w_1 x y w_1 b_1 x_2$. Since $p\{x := a_i\} = x_1 a_1 w_1 a_i w_1 b_1 x_2 \preceq q_i$ and $p\{x := b_i\} = x_1 a_1 w_1 b_i w_1 b_1 x_2 \preceq q_i$ for any $i \in \{1, \ldots, k - 1\}$, we have $S_1(p) \subseteq \bigcup_{i=1}^{k-1} L(q_i)$. For any $w \in \{s \in \Sigma^+ \mid |s| \geq 2\}$, $p\{x := w\} = x_1 a_1 w_1 w w_1 b_1 x_2 \preceq q_k$. Thus, we have $L(p) \subseteq L(Q)$. By Theorem 1, since $p \not\preceq q_i$, $L(p) \not\subseteq L(q_i)$ for any $i \in \{1, \ldots, k\}$. Therefore, $\mathcal{RP}^k$ does not have compactness with respect to containment. □

From Theorem 4 and Lemma 13, we have the following Theorem 5.

Theorem 5: Let $k \geq 3$ and $\sharp\Sigma \geq 2k - 1$. Then, $\mathcal{RP}^k$ has compactness with respect to containment.

In case $k = 2$, we have the following theorem.

Theorem 6: Let $\sharp\Sigma \geq 4$, $P \in \mathcal{RP}^+$ and $Q \in \mathcal{RP}^2$. The following (i), (ii) and (iii) are equivalent:

(i) $S_2(P) \subseteq L(Q)$, (ii) $P \sqsubseteq Q$, (iii) $L(P) \subseteq L(Q)$.

Proof. It is clear that (ii) implies (iii), and (iii) implies (i). Thus, we show that (i) implies (ii). It suffices to show that $S_2(p) \subseteq L(Q)$ implies $P \sqsubseteq Q$ for any regular pattern $p \in \mathcal{RP}$. Let $Q = \{q_1, q_2\}$. The proof is done by mathematical induction on n, where n is the number of variable symbols appearing in p. In case $n = 0$, $p \in \Sigma^+$. Since $S_2(p) = \{p\} \subseteq L(Q)$, we have $p \preceq q$ for some $q \in Q$. For $n \geq 0$, we assume that it is valid for any regular pattern p with n variable symbols. Let p be a regular pattern such that $n + 1$ variable symbols appear in p, and $S_2(p) \subseteq L(Q)$ hold. We assume that $p \not\preceq q_i$ $(i = 1, 2)$. Let p_1, p_2 be regular patterns and x a variable symbol with $p = p_1 x p_2$. For $a, b \in \Sigma$, let $p_a = p\{x := a\}$ and $p_{ab} = p\{x := ab\}$. Note that p_a and p_{ab} have n variable symbols. Thus, by the assumption, $S_2(p_a) \subseteq L(Q)$ and $S_2(p_{ab}) \subseteq L(Q)$ imply $p_a \preceq q_i$ and $p_{ab} \preceq q_{i'}$ for some $i, i' \in \{1, 2\}$. Let $D_i = \{a \in \Sigma \mid p\{x := a\} \preceq q_i\}$ $(i = 1, 2)$. By Lemma 2, if $\sharp D_i \geq 3$ for some $i \in \{1, 2\}$, then $p \preceq q_i$. This contradicts that $p \not\preceq q_i$ $(i = 1, 2)$. Thus, we have $\sharp D_i \leq 2$ for any $i \in \{1, 2\}$. Since $\sharp\Sigma \geq 4$, we consider that $\sharp D_1 = 2$ and $\sharp D_2 = 2$. From Lemma 11, $p\{x := xy\} \preceq q_i$ for some $i \in \{1, 2\}$. From Lemma 12, we have $p \preceq q_i$ for some $i \in \{1, 2\}$. This contradicts that $p \not\preceq q_i$ $(i = 1, 2)$. Hence, (i) implies (ii). □

The next example gives a set of regular patterns $P \in \mathcal{RP}^+$ and a set of regular patterns $Q \in \mathcal{RP}^2$ that, in case $\sharp\Sigma = 3$, the three conditions (i),(ii) and (iii) in Theorem 6 are not equivalent

Example 3: Let $\Sigma = \{a, b, c\}$, p, q_1, q_2 regular patterns and x, x', x'' variable symbols such that $p = x' a x b x''$, $q_1 = x' a b x''$ and $q_2 = x' c x''$. Let $w \in \Sigma^+$. If w contains c, then $p\{x := w\} \preceq q_2$. On the other hand, if w does not contain c, then $p\{x := w\} \preceq q_1$. Thus, $L(p) \subseteq L(q_1) \cup L(q_2)$. However, $p \not\preceq q_1$ and $p \not\preceq q_2$.

From Theorem 6, the following two corollaries holds.

Corollary 3: Let $\sharp\Sigma \geq 4$ and $P \in \mathcal{RP}^+$. Then, $S_2(P)$ is a characteristic set for $L(P)$ within $\mathcal{RPL}^2$.

Corollary 4: Let $\sharp\Sigma \geq 4$. Then, $\mathcal{RP}^2$ has compactness with respect to containment.

4. Regular Pattern without Adjacent Variable Symbols

A regular pattern p is said to be a *non-adjacent variable regular pattern* (*NAV* regular pattern) if p does not contain consecutive variable symbols. For example, the regular pattern $p = axybc$ is not an *NAV* regular pattern because xy is appeared in p. Let $\mathcal{RP}_{NAV}$ be the set of all *NAV* regular patterns. Let $\mathcal{RP}^+_{NAV}$ be the set of all finite subsets S of $\mathcal{RP}_{NAV}$ such that S is not the empty set, i.e., $\mathcal{RP}^+_{NAV} = \{S \subseteq \mathcal{RP}_{NAV} \mid \sharp S \leq 1\}$, and $\mathcal{RP}^k_{NAV}$ the set of all subsets P of $\mathcal{RP}^+_{NAV}$ such that P consists of at most k $(k \geq 1)$ *NAV* regular patterns, i.e., $\mathcal{RP}^k_{NAV} = \{P \in \mathcal{RP}^+_{NAV} \mid \sharp P \leq k\}$. We define the compactness with respect to containment for $\mathcal{RP}^k_{NAV}$ in a similar way as Def.2. For any *NAV* regular pattern $p \in \mathcal{RP}_{NAV}$ and any set $Q \in \mathcal{RP}^k_{NAV}$ with k $(k \geq 1)$, the set $\mathcal{RP}^k_{NAV}$ said to have *compactness with respect to containment* if there exists an *NAV* regular pattern $q \in Q$ such that $L(p) \subseteq L(q)$ holds if $L(p) \subseteq L(Q)$ holds. Then, the following Theorem 7 holds.

Theorem 7: For an integer k $(k \geq 2)$, let $\sharp\Sigma \geq k + 2$, $P \in$

$p = x'cadadaadacbadadaadaxadadaadacbadadaadabx''$,
$q_1 = x'cadadaadacbadadaadacx''$,
$q_2 = x'badadaadacx''$,
$q_3 = x'aadadx''$.

Fig. 20 *NAV* regular patterns p, q_1, q_2, and q_3

$\mathcal{RP}^{+}_{NAV}$, $Q \in \mathcal{RP}^{k}_{NAV}$. Then, the following (i), (ii) and (iii) are equivalent:

$$\text{(i) } S_2(P) \subseteq L(Q), \text{(ii) } P \sqsubseteq Q, \text{(iii) } L(P) \subseteq L(Q).$$

Proof. From the definitions of $\mathcal{RP}^{+}_{NAV}$ and $\mathcal{RP}^{k}_{NAV}$, it is clear that (ii) implies (iii) and (iii) implies (i). Hence, we will show that (i) implies (ii) by mathematical induction on the number n of variable symbols that appear in an *NAV* regular pattern $p \in P$ as follows: If $n = 0$, then we have $S_2(\{p\}) = \{p\}$. Hence, $p \in L(Q)$. Therefore, there exists $q \in Q$ such that $p \preceq q$ holds.

If $n \geq 0$, we assume that the proposition holds for any regular *NAV* regular pattern containing $n \geq 0$ variable symbols. Let p be an *NAV* regular pattern containing $n + 1$ variable symbols such that $S_2(\{p\}) \subseteq L(Q)$ holds and p contains a variable symbol x. There exist two *NAV* regular patterns p_1, p_2 such that $p = p_1 x p_2$ holds. By the induction hypothesis, for any constant string $w \in \Sigma^*$ with $|w| = 2$, $\{p\{x := w\}\} \preceq Q$ holds because $p\{x := w\}$ contains n variable symbols. Hence, there exists an *NAV* regular pattern $q_w \in Q$ such that $p\{x := w\} \preceq q_w$ holds. From Lemma 11, there exists a regular pattern $q \in Q$ such that $p\{x := xy\} \preceq q$ holds, where y is a variable symbol that does not appear in q. This contradicts the condition $Q \in \mathcal{RP}^{k}_{NAV}$. Thus, we have that (i) implies (ii). □

Corollary 5: Let $k \geq 2$, $\sharp\Sigma \geq k+2$ and $P \in \mathcal{RP}^{+}_{NAV}$. Then, $S_2(P)$ is a characteristic set of $\mathcal{RPL}^{k}_{NAV}$.

Lemma 14: Let $k \geq 2$ and $\sharp\Sigma \leq k + 1$. Then, $\mathcal{RP}^{k}_{NAV}$ does not have compactness with respect to containment.

Proof. Let $\Sigma = \{a_1, \ldots, a_{k+1}\}$. We assume that for $i = 1, 2, \ldots, k$, $p\{x := a_i y\} \preceq q_i$ and $p\{x := y a_{i+1}\} \preceq q_i$ holds. If $p\{x := a_{k+1} a_1\} \preceq q_1$ holds, $S_2(p) \backslash S_1(p) \subseteq \bigcup_{i=1}^{k} L(q_i)$ holds. This show that $L(p) \subseteq L(Q)$ holds. However, for $i = 1, 2, \ldots, k$, since $p \not\preceq q_i$ holds, we have that $L(p) \not\subseteq L(q_i)$ holds. Hence, $\mathcal{RP}^{k}_{NAV}$ does not have compactness with respect to containment. □

Next, in Example 4, we give an example for Lemma 14.

Example 4: Let Σ be the set of four constant symbols a, b, c, d, i.e., $\Sigma = \{a, b, c, d\}$ and x, x', x'' three distinct variable symbols. Let p, q_1, q_2, q_3 be the *NAV* regular patterns given in Fig. 20. Then, we have $L(p) \subseteq L(q_1) \cup L(q_2) \cup L(q_3)$. This show that for $P = \{p\}$, $Q = \{q_1, q_2, q_3\}$, (iii) of Theorem 7 holds. However, since $p \not\preceq q_1$, $p \not\preceq q_2$ and $p \not\preceq q_3$ hold, we have $P \not\sqsubseteq Q$, that is, (ii) of Theorem 7 does not hold.

From Theorem 7 and Lemma 14, we have the following theorem.

Table 2 The conditions on the number $\sharp\Sigma$ of constant symbols in Σ required for compactness with respect to containment.

Class	$k = 2$	$k \geq 3$
$\mathcal{RP}^k$	$\sharp\Sigma \geq 4$	$\sharp\Sigma \geq 2k - 1$
$\mathcal{RP}^{k}_{NAV}$	$\sharp\Sigma \geq k + 2$	

Theorem 8: Let $k \geq 2$ and $\sharp\Sigma \geq k + 2$. Then, the set $\mathcal{RPL}^{k}_{NAV}$ has compactness with respect to containment.

5. Conclusion

In this paper, for an integer k ($k \geq 2$), we have shown the conditions on the number of constant symbols in Σ, summarized in Table 2, required for the classes $\mathcal{RP}^k$ of all the set of k regular pattern languages and $\mathcal{RP}^{k}_{NAV}$ of all the set of k non-adjacent variable regular patterns in $\mathcal{RP}_{NAV}$ to have compactness with respect to containment. This result leads to design an efficient learning algorithm for finite unions of languages of non-adjacent variable regular patterns in $\mathcal{RP}_{NAV}$, based on the learning algorithm for $\mathcal{RP}^k$ proposed by Arimura et al. [8].

Extending the notion of strong compactness, as introduced by Arimura et al. [9], to finite unions of regular pattern languages with non-adjacent variables remains as a topic for future research. Furthermore, based on the characteristic set for $\mathcal{RP}^{k}_{NAV}$, we plan to propose a polynomial-time inductive inference algorithm that identifies finite unions of regular pattern languages with non-adjacent variables in the limit from positive examples. Ishinada et al. [17] investigated a query learning model that employs high-precision Graph Convolution Networks (GCNs) as oracles for tree patterns. Applying the findings of the present study to tree pattern languages, with the aim of enabling the extension of their work to finite unions of tree pattern languages, remains an important direction for future research.

Acknowledgements

This work was partially supported by JSPS KAKENHI Grant Numbers JP20K04973, JP21K12021, JP22K12172, JP24K15074, and JP24K15090. We thank Mr. K. Horii, a master's student at the Graduate School of Information Sciences, Hiroshima City University, for fruitful discussions and constructive comments.

References

[1] D. Angluin, "Finding Patterns Common to a Set of Strings," Journal of Computer and System Sciences, 21(1):46–62, 1980, DOI:10.1016/0022-0000(80)90041-0.
[2] D. Angluin, "Inductive Inference of Formal Languages from Positive Data," Information and Control, 45(2):117–135, 1980, DOI:10.1016/S0019-9958(80)90285-5.
[3] Y. Mukouchi, "Characterization of Pattern Languages," in Proc. ALT '91, Ohmusha, pp.93-104, 1991.
[4] D. Sato, Y. Mukouchi and D. Zheng, "Characteristic Sets for Unions of Regular Pattern Languages and Compactness," in Proc. ALT '98,

Springer LNAI 1501, pp.220-233, 1998.
[5] Y. Mukouchi, “Containment Problems for Pattern Languages,” IEICE Transactions on Information and Systems, E75-D(4):420-425, 1992.
[6] K. Wright, “Identification of Unions of Languages Drawn from an Identifiable Class,” in Proc. COLT 1989, pp.328-333, 1989.
[7] T. Shinohara and H. Arimura, “Inductive inference of unbounded unions of pattern languages from positive data,” Theoretical Computer Science, 241(1-2): 135–161, 2000, DOI:10.1016/S0304-3975(99)00270-4.
[8] H. Arimura, T. Shinohara and S. Otsuki, “Finding Minimal Generalizations for Unions of Pattern Languages and Its Application to Inductive Inference from Positive Data,” in Proc. STACS ’94, Springer LNCS 775, pp.649-660, 1994.
[9] H. Arimura and T. Shinohara, “Strong Compactness of Containment for Unions of Regular Pattern Languages (in Japanese),” RIMS Kôkyûroku of Kyoto Univ., Vol.950, pp.246-249, 1996.
[10] J.D. Day, D. Reidenbach and M.L. Schmid, “Closure Properties of Pattern Languages,” Journal of Computer and System Sciences 84:11-31, 2017, DOI:10.1016/j.jcss.2016.07.003.
[11] S. Matsumoto, T. Uchida, T. Shoudai, Y. Suzuki, and T. Miyahara, “An Efficient Learning Algorithm for Regular Pattern Languages Using One Positive Example and a Linear Number of Membership Queries,” IEICE Trans. Inf. & Syst., vol.E103-D, No.3, pp.526-539, 2020, DOI:10.1587/transinf.2019FCP0009.
[12] N. Taketa, T. Uchida, T. Shoudai, S. Matsumoto, Y. Suzuki, and T. Miyahara, “Visualizing the Prediction Basis of Deep Learning Models using a Query Learning Algorithm for Linear Pattern Languages (in Japanese)”, JSAI2022(The 36th Annual Conference of the Japanese Society for Artificial Intelligence), 2G4-GS-2-03, 2022.
[13] S. Arikawa, T. Shinohara, and A. Yamamoto, “Learning elementary formal system,” Theoretical Computer Science, vol.95, no.1, pp.97–113, 1992, DOI:10.1016/0304-3975(92)90068-Q
[14] H. Arimura, H. Ishizaka and T. Shinohara,“Learning Unions of Tree Patterns Using Queries,” Theor. Comput. Sci., vol.185, No.1, pp.47-62, 1997, DOI:10.1016/S0304-3975(97)00015-7.
[15] Y. Suzuki, T. Shoudai, T. Uchida, and T. Miyahara, “Ordered Term Tree Languages Which are Polynomial Time Inductively Inferable from Positive Data,” Theoretical Computer Science, vol.350, No.1, pp.63-90, 2006, DOI:10.1016/j.tcs.2005.10.022.
[16] T. Uchida, T. Shoudai, and S. Miyano, “Parallel Algorithms for Refutation Tree Problem on Formal Graph Systems,” IEICE Trans. Inf. & Syst., vol.E78-D, No.2 pp.99-112, 1995.
[17] K. Ishinada, T. Shoudai, T. Uchida, and S. Matsumoto, “Analysis of Query Learning Models with High-Accuracy GCN Oracles for Unordered Tree Patterns (in Japanese),” IPSJ SIG Technical Report on Mathematical Modeling and Problem Solving (MPS), 15, 2023.

Naoto Taketa received the B.S. and M.S. degrees in Information Sciences from Hiroshima City University, in 2022 and 2024, respectively. He is currently with Rakuten Card Co., Ltd., System Strategy Department, System Division.

Tomoyuki Uchida received the B.S. degree in Mathematics, the M.S. and Dr. Sci. degrees in Information Systems all from Kyushu University, in 1989, 1991 and 1994, respectively. Currently, he is a professor of Graduate School of Information Sciences, Hiroshima City University. His research interests include data mining from semistructured data, algorithmic graph theory and algorithmic learning theory.

Takayoshi Shoudai received the B.S. in 1986, the M.S. degree in 1988 in Mathematics and the Dr. Sci. in 1993 in Information Science all from Kyushu University. Currently, he is a professor of Department of Computer Science and Engineering, Fukuoka Institute of Technology. His research interests include graph algorithms, computational learning theory and machine learning.

Satoshi Matsumoto is a professor of Department of Mathematical Sciences, Tokai University, Kanagawa, Japan. He received the B.S. degree in Mathematics, the M.S. and Dr. Sci. degrees in Information Systems all from Kyushu University, Fukuoka, Japan in 1993, 1995 and 1998, respectively. His research interests include algorithmic learning theory.

Yusuke Suzuki received the B.S. degree in Physics, the M.S. and Dr. Sci. degrees in Informatics all from Kyushu University, in 2000, 2002 and 2007, respectively. He is currently a lecturer of Graduate School of Information Sciences, Hiroshima City University, Hiroshima, Japan. His research interests include machine learning and data mining.

Tetsuhiro Miyahara is an associate professor of Graduate School of Information Sciences, Hiroshima City University, Hiroshima, Japan. He received the B.S. degree in Mathematics, the M.S. and Dr. Sci. degrees in Information Systems all from Kyushu University, Fukuoka, Japan in 1984, 1986 and 1996, respectively. His research interests include algorithmic learning theory, knowledge discovery and machine learning.