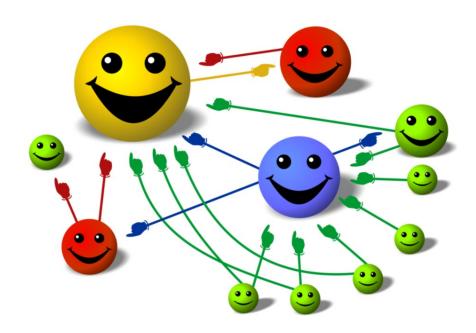
# CS5344 Link Analysis



# Web as a Graph

Nodes: Webpages

Edges: Hyperlinks

I teach a class on Database.

CS2102: Classes are in COM2 building

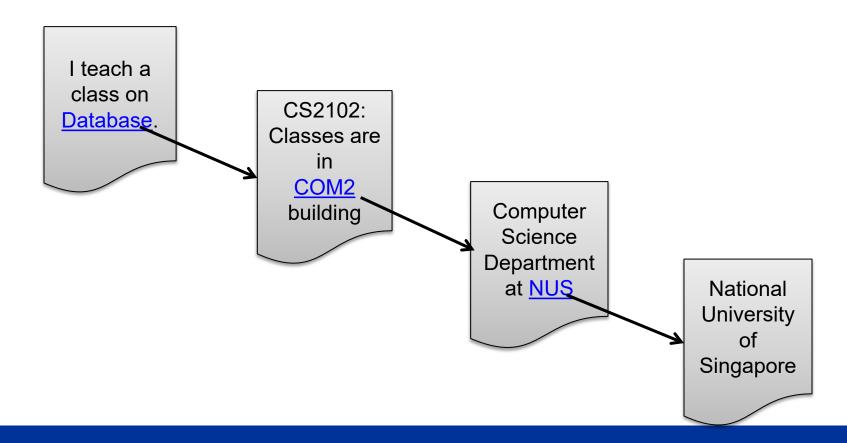
Computer Science Department at NUS

National University of Singapore

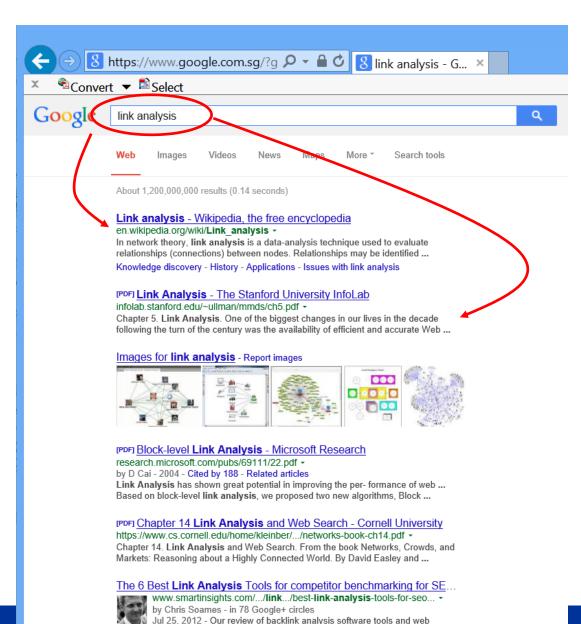
# Web as a Graph

Nodes: Webpages

Edges: Hyperlinks



#### **Web Search**



How does the search engine decide which page should be ranked higher?

# Web Search - Challenges

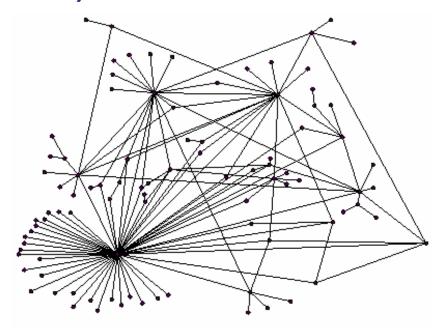
- Web contains many sources of information.
  - Who to "trust"?
- What is the "best" answer to query "newspaper"?
  - No single right answer

# **Link Analysis**

- The Web is not just a collection of documents
  - The hyperlinks are important
- A link from page A to page B may indicate
  - A is related to B, or
  - A is recommending, citing, voting for, or endorsing B
- Types of links:
  - Referential click here and get back home
  - Informational click here to get more detail
- Links influence the ranking of web pages and thus have commercial value

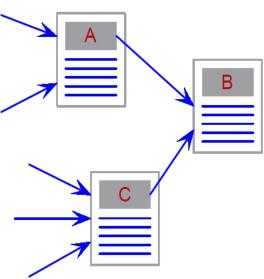
# Importance of Web Pages

- Not all web pages are equally important
- A page is important if it is pointed to by other important pages (recursion)

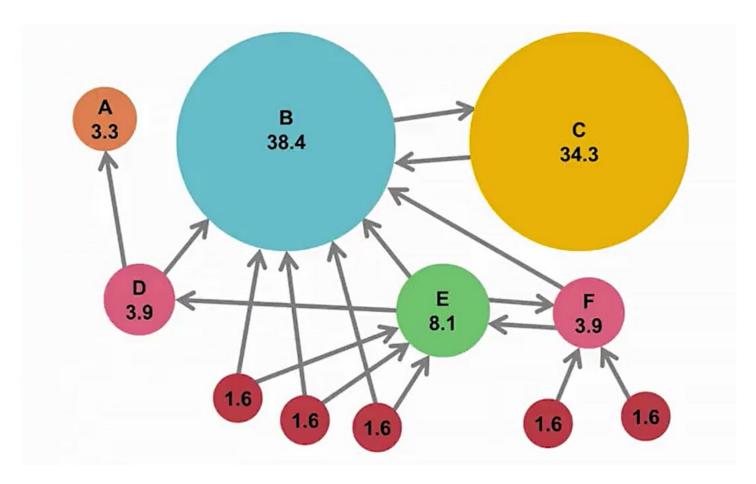


# **PageRank**

- Idea: Links as votes
  - A page is more important if it has more links
- Incoming links to a page is a measure of importance and authority of the page
  - www.stanford.edu has 23,400 in-links
  - www.joe-schmoe.com has 1 in-link
- Are all incoming links equal?
  - Links from important pages count more



#### **Example PageRank Scores**



- A "vote" from an important page is worth more
- A page is important if it is pointed to by other important pages

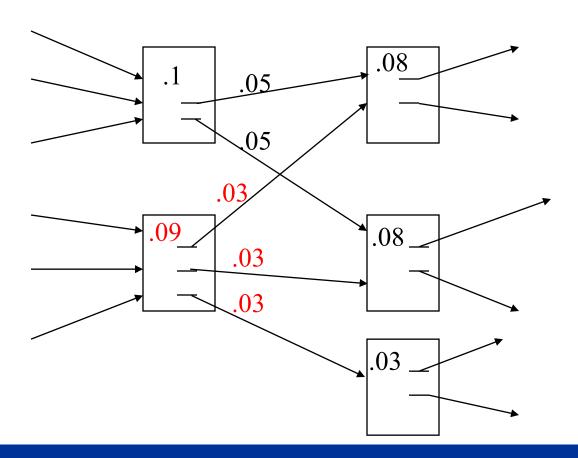
#### **Recursive Formulation**

- A link's vote is proportional to the importance of its source page
- If page j with importance  $r_j$  has n out-links, each link gets  $r_i/n$  votes
- Page j's own importance is the sum of the votes on its in-links

$$r_j = r_i/3 + r_k/4$$

#### Flow Model

 Can view it as a process of PageRank "flowing" from pages to the pages they point to



#### Flow Model

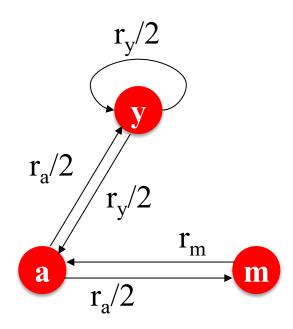
Define a rank r<sub>i</sub> for page j

$$r_j = \sum_{i \to j} \frac{r_i}{d_i}$$

 $d_i$  is the out-degree of node i

Flow Equations

$$r_y = r_y/2 + r_a/2$$
  
 $r_a = r_y/2 + r_m$   
 $r_m = r_a/2$ 



#### **Matrix Formulation**

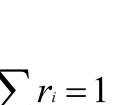
- Stochastic adjacency matrix M
  - Let page i has d<sub>i</sub> outlinks
  - If  $i \rightarrow j$ , then  $M_{ii} = 1/d_i$  else  $M_{ii} = 0$





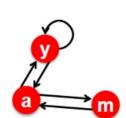


- Vector with one entry per page
- $r_i$  is the importance score of page i



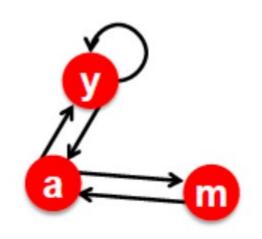
• Flow equations can be written in matrix form 
$$r = M \cdot r$$

$$r_j = \sum_{i \to j} \frac{r_i}{d_i}$$



	y	a	1	m
y	1/2	1/2	2	0
a	1/2	0		1
m	0	1/2	2	0
,			T	

### Example Flow Equations and M



	y	a	m
y	1/2	1/2	0
a	1/2	0	1
m	0	1/2	0

$$\begin{vmatrix}
y \\
a \\
m
\end{vmatrix} = \begin{vmatrix}
\frac{1}{2} & \frac{1}{2} & 0 \\
\frac{1}{2} & 0 & 1 \\
0 & \frac{1}{2} & 0
\end{vmatrix}$$

#### **Power Iteration Method**

- Given a web graph with n nodes, where the nodes are pages and edges are hyperlinks
- Power iteration simple iterative scheme
  - Suppose there are N web pages
  - Initialize  $\mathbf{r}^{(0)} = [1/N, ...., 1/N]^T$
  - Iterate:  $r^{(t+1)} = M \cdot r^{(t)}$

- $r_j^{(t+1)} = \sum_{i \to j} \frac{r_i^{(t)}}{d_i}$
- $d_i$  .... out-degree of node i

- Stop when  $|\mathbf{r}^{(t+1)} \mathbf{r}^{(t)}|_1 < \varepsilon$ 
  - $|\mathbf{x}|_1 = \sum_{i \in [1,N]} |x_i|$  is the L<sub>1</sub> norm
  - Can use any other vector norm e.g., Euclidean

#### **Example**

#### Power Iteration:

Set 
$$r_i = 1/N$$

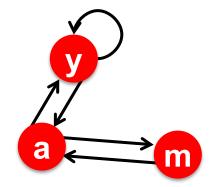
1: 
$$r'_j = \sum_{i \to j} \frac{r_i}{d_i}$$

**2**: 
$$r = r'$$

Goto 1

#### Example:

$$\begin{pmatrix} \mathbf{r}_{\mathbf{y}} \\ \mathbf{r}_{\mathbf{a}} \\ \mathbf{r}_{\mathbf{m}} \end{pmatrix} = \frac{1/3}{1/3}$$



	у	а	m
у	1/2	1/2	0
а	1/2	0	1
m	0	1/2	0

$$r_y = r_y/2 + r_a/2$$

$$r_a = r_y/2 + r_m$$

$$r_m = r_a/2$$

Iteration 0, 1, 2, ...

#### **PageRank**

$$r_j^{(t+1)} = \sum_{i \to j} \frac{r_i^{(t)}}{\mathbf{d_i}}$$
 or equivalently  $r = Mr$ 

#### • Questions:

- 1. Does this converge?
- 2. Does it converge to what we want?

# Does this converge?

$$r_j^{(t+1)} = \sum_{i \to j} \frac{r_i^{(t)}}{d_i}$$

Iteration 0, 1, 2, ...

#### Does it converge to what we want?

$$r_j^{(t+1)} = \sum_{i \to j} \frac{r_i^{(t)}}{d_i}$$

$$\frac{\mathbf{r_a}}{\mathbf{r_b}} = \frac{1}{0} \quad \frac{0}{1} \quad \frac{0}{0}$$

Iteration 0, 1, 2, ...

#### **Problems on Real Web**

#### Imagine a random web surfer

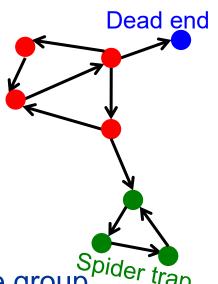
- At any time t, surfer is on some page i
- At time t+1, surfer follows an out-link from i uniformly at random
- Ends up on some page j linked from i

#### Dead ends

- A page has no out-links
- Random walk has "nowhere" to go to
- Such pages cause importance to "leak out"

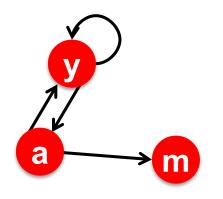
#### Spider traps

- A group of pages have no out-links out of the group place training
- Random walk gets "stuck" in a trap
- Eventually spider traps absorb all importance



#### **Problem: Dead Ends**

- A page with no out-links
- Random walk has "nowhere" to go to
- All importance "leaks out of" the Web!
- Matrix is not stochastic so initial assumptions are not met



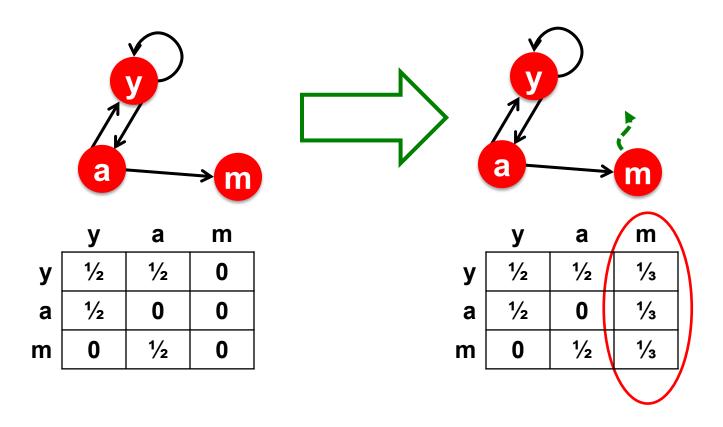
	у	а	m
У	1/2	1/2	0
a	1/2	0	0
n	0	1/2	0 /

$$\begin{bmatrix} r_y \\ r_a \\ r_m \end{bmatrix} = \begin{bmatrix} 1/3 \\ 1/3 \\ 1/6 \\ 1/6 \end{bmatrix} \begin{bmatrix} 1/4 \\ 1/6 \\ 1/12 \end{bmatrix} \begin{bmatrix} 5/24 \\ 1/8 \\ 1/12 \end{bmatrix} \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$

Iteration 0, 1, 2, ...

# **Solution: Teleport**

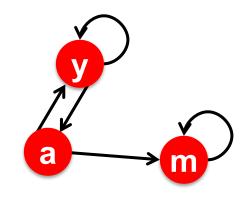
 Adjust the matrix to allow a surfer to jump to some random page from dead ends



### **Problem: Spider Traps**

- A group of pages with no links out of the group
- Random walk gets "stuck" in a trap
- Accumulate all the importance of the Web



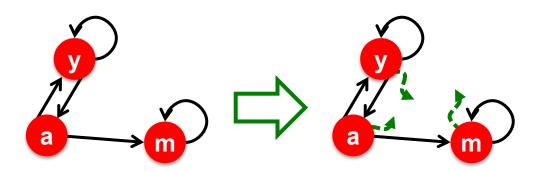


	У	а	m
y	1/2	1/2	0
a	1/2	0	0
n	0	1/2	1

Г -	] [ ]		Г 7		Г
ry	1/3	1/3	1/4	5/24	0
<b>r</b> a	= 1/3	1/6	1/6	1/8	 0
rm	$\begin{vmatrix} = & 1/3 \\ 1/3 & \end{vmatrix}$	1/2	7/12	2/3	(1)
L _	]	ᆫ	ᆸ	_	

### **Solution: Teleport**

- At each time step, a random surfer has two options
  - With probability  $\beta$ , follow a link at random
  - With probability 1-β, jump to some random page
  - Common values for  $\beta$  are in the range 0.8 to 0.9
- Surfer will teleport out of spider trap within a few time steps

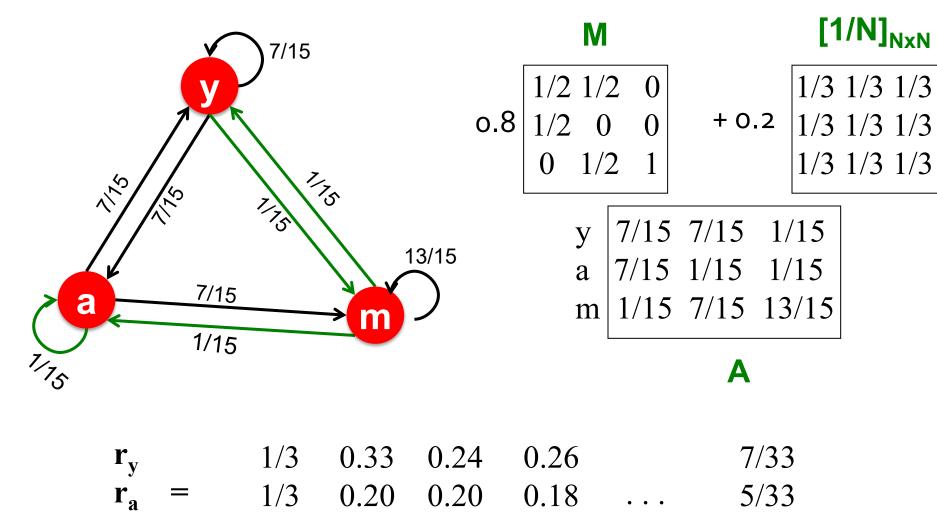


# Random Teleports ( $\beta$ = 0.8)

1/3

 $\mathbf{r}_{\mathbf{m}}$ 

0.46



0.52

0.56

21/33

# Limitations of PageRank

- Measures generic popularity of a page
  - Ignore or miss topic-specific authorities
  - Solution: Topic-specific PageRank
- Susceptible to link spam
  - Artificial link topologies created in order to boost page rank
  - Solution: TrustRank
- Uses a single measure of importance
  - Other models of importance
  - Solution: Hubs-and-Authorities

# **Topic-Specific PageRank**

- Instead of generic popularity, can we measure popularity within a topic?
- Goal: Evaluate Web pages not just according to their popularity, but by how close they are to a particular topic, e.g. "sports" or "history"
- Allows search queries to be answered based on interests of the user
  - Example: Is "Jaguar" an animal, the automobile, or a version of MAC OS?

# **Topic-Specific PageRank**

- Recall random walker has a small probability of teleporting at any step
  - Standard PageRank: Any page with equal probability
  - Topic Specific PageRank: Teleport set is restricted to a topic-specific set of "relevant" pages
- Idea: Bias the random walk
  - When random walker teleports, pick a page from a set S of web pages
  - S contains only pages that are relevant to the topic
  - Get a different rank vector r<sub>s</sub> for each teleport set S

# **Topic-Specific PageRank**

- Decide on topics to create PageRank vectors
  - Open Directory (DMOZ) (www.dmoz.org)
  - The 16 DMOZ top-level categories: arts, business, sports, ...
- Pick a teleport set for each of these topics, and compute the topic-sensitive PageRank vector for that topic
- Determine the topic that is most relevant for a query
  - User picks from a menu
  - Query context e.g., query from a web page on a known topic
  - User context e.g., user's bookmarks
- Use the PageRank vectors for that topic to order results to the search query

# TrustRank – Combating Web Spam

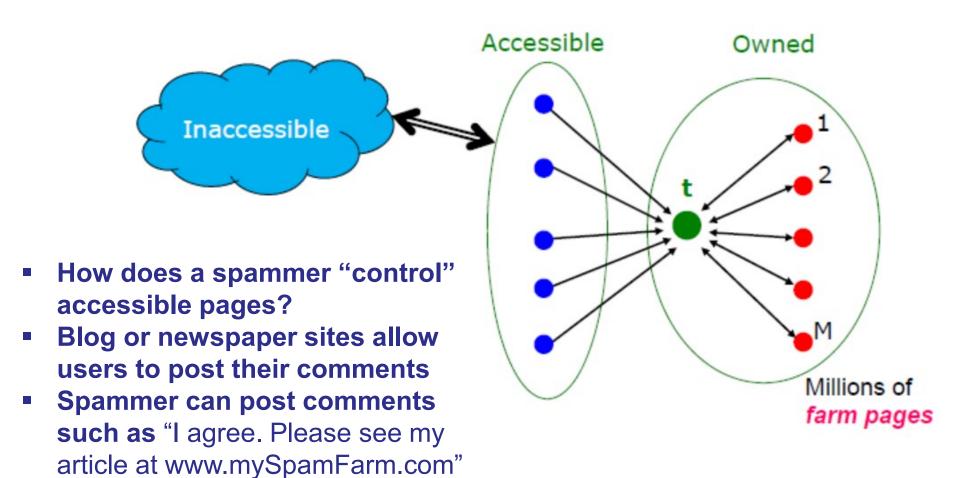
#### Spamming

- Any deliberate action to boost a web page's position in search engine results, incommensurate with page's real value
- Approximately 10-15% of web pages are spam

#### Link Spam

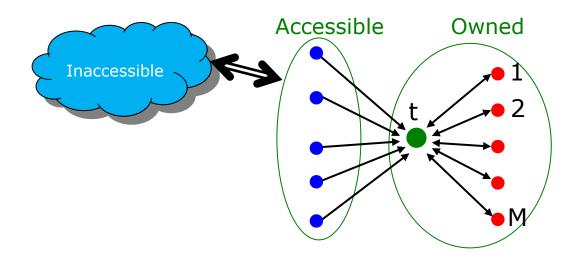
Create link structures that boost the PageRank of a particular page

#### Spammer's View of the Web

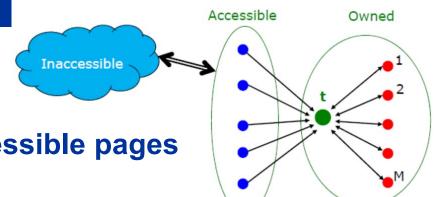


#### **Link Farms**

- Spammer's goal is to maximize the PageRank of target page t
- Get as many links from accessible pages as possible to target page t
- Construct "link farm" to get PageRank multiplier effect



#### **Analysis**



x: PageRank contributed by accessible pages

y: PageRank of target page t

N: Total number of web pages

Rank of each "farm" page 
$$=\frac{\beta y}{M} + \frac{1-\beta}{N}$$

$$y = x + \beta M \left[ \frac{\beta y}{M} + \frac{1 - \beta}{N} \right] + \frac{1 - \beta}{N}$$

$$y = \frac{x}{1-\beta^2} + c\frac{M}{N}$$
 where  $c = \frac{\beta}{1+\beta}$ 

Let  $\beta$  = 0.85. Then  $1/(1 - \beta^2)$  = 3.6, and c = 0.46

- External PageRank (x) increased by 360%!
- Obtain additional amount of PageRank that is 46% of the fraction of the Web, M/N, that is in the spam farm
- By making M large, we can make y as large as we want

### **How to Combat Link Spam?**

- Detect and blacklist structures that look like spam farms
  - One page links to a very large number of pages, each of which links back to it
  - Leads to more sophisticated way of hiding spam farms, and detecting them...
- TrustRank: Topic-specific PageRank with a teleport set of trusted pages
  - e.g, .edu domains, .gov domains, etc
  - Lower the score of spam pages

#### **TrustRank**

- Basic principle: Approximate isolation
  - It is rare for a "good" (trustworthy) page to point to a "bad" (spam) page
- Sample a set of seed pages from the web
- An oracle (human) identifies the good pages and the spam pages in the seed set
  - Expensive task, so keep seed set small
  - Subset of pages in the seed set that are identified as good are called the trusted pages

#### **Trust Propagation**

- Perform a topic-sensitive PageRank with the trusted pages as the teleport set
- Propagate trust through links
  - Each page gets a trust value between 0 and 1
- Use a threshold value and mark all pages below the trust threshold as spam

### **Trust Propagation (Simple Model)**

- Set trust of each trusted page to 1
- Suppose trust of page p is t<sub>p</sub>
  - p has a set of out-links  $o_p$
- For each  $q \in o_p$ , p confers the trust to q
  - $\beta t_p / |o_p|$  for  $0 < \beta < 1$
- Trust is additive
  - Trust of p is the sum of the trust conferred on p by all its inlinked pages
- Trust attenuation
  - Degree of trust conferred by a trusted page decreases with the distance in the graph
- Trust splitting
  - The larger the number of out-links, the less scrutiny the page author gives each out-link; trust is split across out-links

# Picking the Seed Set

#### Two conflicting considerations:

- Human has to inspect each seed page → seed set must be small
- Must ensure every good page gets adequate trust rank → need to make all good pages reachable from seed set by short paths

#### 1. Use PageRank to pick the top-k pages

Theory is a bad page cannot have very high rank

#### 2. Use trusted domains with controlled membership

E.g. university pages (.edu)or government pages (.gov)

### **Summary**

- Link analysis in social network graphs to find communities
- Girvan-Newman algorithm use edge betweenness measure to separate nodes into communities
- Content of web pages and hyperlinks are important in web search
- Page Rank algorithm determine importance of web pages
- Trust Rank algorithm to overcome link spams