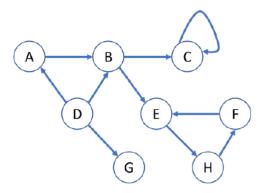
CS4225/CS5425 BIG DATA SYSTEMS FOR DATA SCIENCE

Tutorial 5: Graph and Test Practice

- 1. Given the following graph,
- 1) how many dead ends are there in the graph? For each dead end (if any), please indicate the set of vertices forming the dead end.
- 2) how many spider traps are there in the graph? For each spider trap (if any), please indicate the set of vertices forming the spider trap.

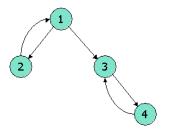


Answer:

- 1) Dead ends: $\{G\}$,
- 2) Spider traps: {C}, {E, F, H}, {B, C, E, F, H}, {A, B, C, E, F, H}
- **2**. True/False: In Topic-specific PageRank, random walker will teleport to any page with equal probability.

Answer: False. Random walker will only teleport to a topic-specific set of "relevant" pages.

3. Consider the following link topology.



Write down the Topic-Specific PageRank equations for the following link topology. Assume that pages selected for the teleport set are nodes 1 and 2 (where teleports go to either node with equal probability). Assume further that the teleport probability, $(1 - \beta)$, is 0.3.

Recall the topic sensitive pagerank (TSPR) equations:

```
A_{ij} = \begin{cases} \beta M_{ij} + (1 - \beta)/|S| & \text{if } i \in S \\ \beta M_{ij} + 0 & \text{otherwise} \end{cases}
```

Basically, the only difference here (compared to the non-topic sensitive case in the last 2 questions), is that the teleport terms are only distributed among the nodes in the teleport S, instead of being distributed over all nodes. Let r_1 , r_2 , r_3 , r_3 denote the importance of the 4 nodes.

```
r_1 = 0.7 r_2 + 0.15

r_2 = 0.35 r_1 + 0.15

r_3 = 0.35 r_1 + 0.7 r_4

r_4 = 0.7 r_3
```

4. Show pseudocode for the compute() function for the PageRank with teleport (β = 0.85) over vertices algorithm in Pregel / Giraph. Set the initial PageRank value as 1/N (N is the number of vertices), Run 30 iterations and then stop. You can (if you choose) use the functions: getValue(), setValue(), getNumVertices(), getSuperStep(), getOutEdgeIterator().

Answer:

```
Compute(v, messages):
  if getSuperStep() == 0:
    v.setValue(1 / getNumVertices())
  if getSuperStep() >= 1:
    sum = 0
    for m in messages:
        sum += m
    v.setValue(0.15 / getNumVertices() + 0.85 * sum)
  if getSuperStep() < 30:
        sendMsgToAllEdges(v.getValue() / len(getOutEdgeIterator()))
  else:
    voteToHalt()</pre>
```

Writing pseudo code is fine; no worries about the syntax. E.g. len(x) or x.size() or "size of x" are all fine.

5. On a certain large Spark cluster, Rose creates a data frame named **traceA**, and writes the shown program to process the trace. The **traceA** data frame keeps the logs, and each log entry represents the log information of one web page access, including various fields: **ip**, the IP address of the log entry and **time**, the amount of time of that access.

```
Line 0: maxSql = spark.sql("""
Line 1: SELECT ip, sum(time) as access_total
Line 2: FROM traceA
Line 3: WHERE time>0.1
Line 4: GROUP BY ip
Line 5: ORDER BY sum(time) DESC
Line 6: """)
Line 7: maxSql.collect()
```

For the above codes, what are the potential performance bottlenecks? Please identify which lines cause the bottleneck and justify your answer.

Answer:

Depending on whether the dataframe has been in the RAM, if not reading traceA (Line 1-3) may incur potential I/O cost

Depending on the size of data generated from Line 3, Line 4 & 5 can also be the bottleneck as both operations are wide transformation and thus require data shuffling through the network I/O.

If the grouped data after Line 4 is highly skewed (e.g. a super big size of data for certain ip address), Line 5 will have task straggler issue, i.e. certain tasks takes much longer time than the other tasks to complete.

Line 7 may return too many results. We can add LIMIT to the SQL query to limit the number of results.