

{Your Name, Your UNI}  
Homework 1 Solutions  
CSOR 4231.001 — Spring 2023

**Homework Problems**

1. (20 points) Consider the problem of computing  $N! = 1 \cdot 2 \cdot 3 \cdots N$ .
  - (a) (10 points) If  $N$  is an  $n$ -bit integer, how many bits long is  $N!$  in  $\Theta()$  notation? Give your answer in terms of both of  $N, n$ .
  - (b) (10 points) Give an algorithm to compute  $N!$  and analyze its running time. Give your answer in terms of both of  $N, n$ . Assume that the time to multiply two  $n$ -bit integers is  $O(n^2)$ .
  
2. (40 points) On sorting
  - (a) (15 points) *A lower bound for comparison-based sorting*

You may view a comparison-based sorting algorithm as a binary tree. For example, the tree in Figure 1 sorts 3 integers  $x_1, x_2, x_3$ : it first compares  $x_1$  to  $x_2$ ; if  $x_1 < x_2$ , it compares  $x_2$  to  $x_3$ , else it compares  $x_1$  to  $x_3$ , etc. Each leaf corresponds to a possible permutation of the 3 numbers. For example, if  $x_1 < x_2 < x_3$ , we get the leaf labelled  $x_1, x_2, x_3$ .

Suppose you want to sort  $n$  integers.

    - i. (2 points) Argue that every possible permutation of the  $n$  numbers must appear as a leaf in the tree corresponding to a sorting algorithm.
    - ii. (2 points) *Fill in the missing number in the following sentence:* Thus the tree has at least ... leaves.
    - iii. (3 points) The worst-case time complexity of the sorting algorithm is given by the maximum number of comparisons performed by the algorithm. What does the latter quantity correspond to in the tree?
    - iv. (3 points) Consider a binary tree of depth  $d$ . Give with a proof an upper bound on the number of leaves of the tree.
    - v. (5 points) Derive a lower bound for the worst-case time complexity of a comparison-based sorting algorithm.
  - (b) (22 points) Give an algorithm that sorts any array of  $n$  integers  $x_1, \dots, x_n$  in  $O(n + D)$  time, where  $D = \max_i x_i - \min_i x_i$ . *Hint: Think how to use extra space to achieve this running time.*
  - (c) (3 points) Suppose  $D$  is small (that is,  $O(n)$ ). Does the algorithm you proposed in part (b) contradict the lower bound you derived in part (a)?

3. (30 points) The Hadamard matrices  $H_0, H_1, H_2, \dots$  are defined as follows:

- $H_0$  is the  $1 \times 1$  matrix  $\begin{bmatrix} 1 \end{bmatrix}$
- For  $k > 0$ ,  $H_k$  is the  $2^k \times 2^k$  matrix

$$H_k = \begin{bmatrix} H_{k-1} & H_{k-1} \\ H_{k-1} & -H_{k-1} \end{bmatrix}$$

Let  $\nu$  be a column vector of length  $n = 2^k$ . Can you compute the matrix-vector product  $H_k \nu$  faster than the straightforward algorithm that requires  $O(n^2)$  time? (Assume that all the numbers involved are small enough that basic arithmetic operations like addition and multiplication take unit time.)

4. (60 points) The Fibonacci numbers are defined by the recurrence

$$F_n = \begin{cases} 0, & \text{if } n = 0 \\ 1, & \text{if } n = 1 \\ F_{n-1} + F_{n-2}, & \text{if } n \geq 2 \end{cases}$$

(a) (4 points) Show that  $F_n \geq 2^{n/2}$ ,  $n \geq 6$ .

(b) Assume that the cost of adding, subtracting, or multiplying two integers is  $O(1)$ , independent of the size of the integers.

- (6 points) Give an algorithm that computes  $F_n$  based on the recursive definition above. Develop a recurrence for the running time of your algorithm and give an asymptotic *lower* bound for it.
- (8 points) Give a non-recursive algorithm that asymptotically performs fewer additions than the recursive algorithm. Give an asymptotic upper bound for the new algorithm.
- (22 points) Give an algorithm to compute  $F_n$  in  $O(\log n)$  time using only integer additions and multiplications.

(Hint: Observe that

$$\begin{bmatrix} F_2 \\ F_1 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \cdot \begin{bmatrix} F_1 \\ F_0 \end{bmatrix}.$$

Can you build on this to compute  $F_n$ ?)

(c) (20 points) Now assume that adding two  $m$ -bit integers requires  $\Theta(m)$  time and that multiplying two  $m$ -bit integers requires  $\Theta(m^2)$  time. What is the running time of the three algorithms in part (b) under this more reasonable cost measure for the elementary arithmetic operations?