Lecture #7a

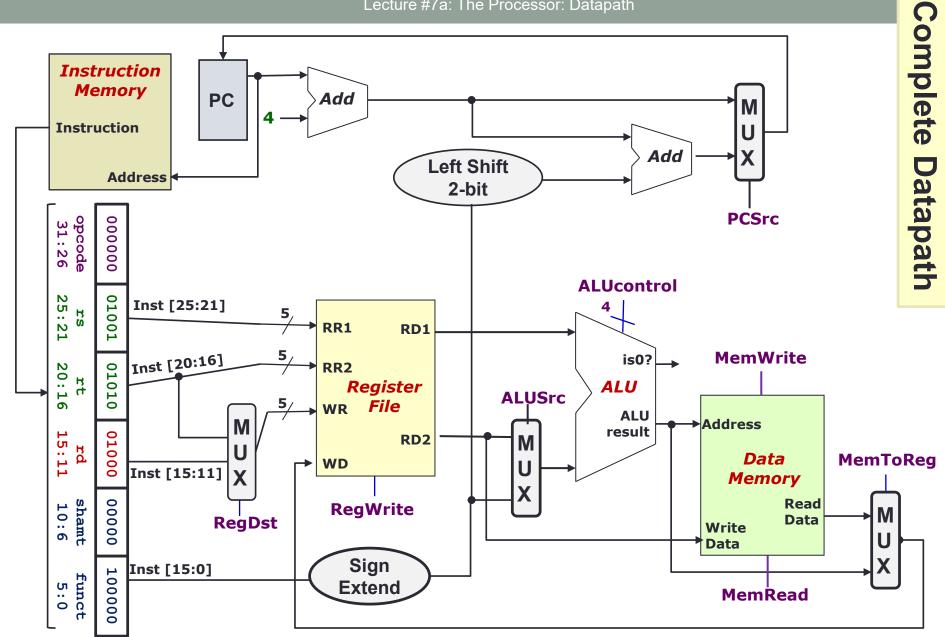
The Processor: Datapath

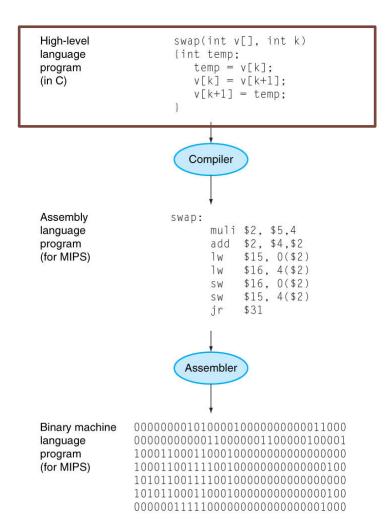


Lecture #7a: Processor: Datapath

- 1. The Complete Datapath!
- 2. Brief Recap
- 3. From C to Execution
 - 3.1 Writing C Program
 - 3.2 Compiling to MIPS
 - 3.3 Assembling to Binaries
 - 3.4 Execution (Datapath)

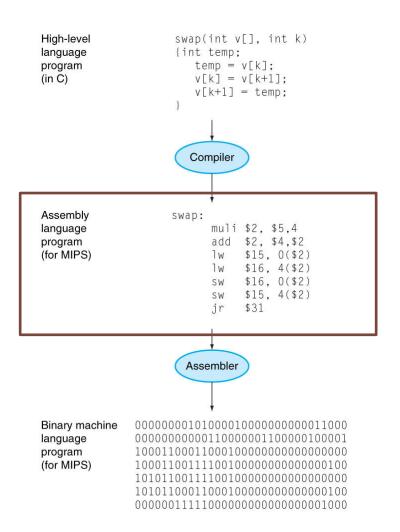
Taken from CS2100 Lecture 11a.





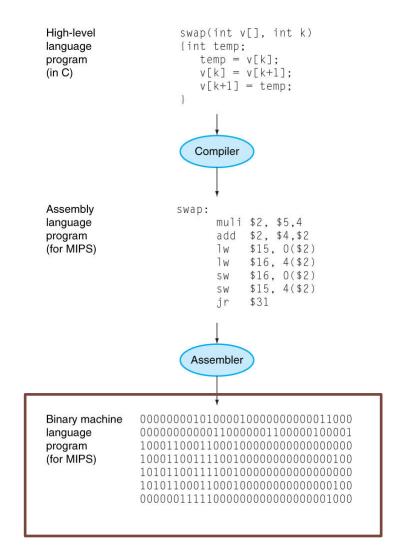
Write program in high-level language (e.g., C)

```
if(x != 0) {
  a[0] = a[1] + x;
}
```



Compiler translates to assembly language (e.g., MIPS)

```
beq $16, $0, Else
      lw $8, 4($17)
      add $8, $8, $16
      sw $8, 0($17)
Else:
```



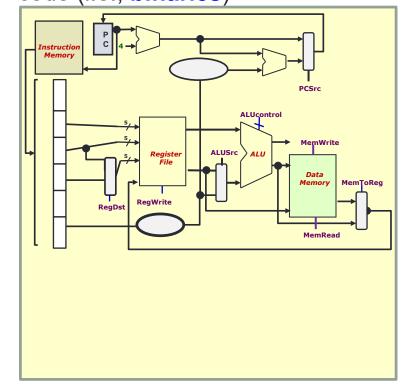
Assembler translates to machine code (i.e., binaries)

```
0001 0010 0000 0000
0000 0000 0000 0011
1000 1110 0010 1000
0000 0000 0000 0100
0000 0010 0000 1000
0100 0000 0001 0100
1010 1110 0010 1000
0000 0000 0000 0000
```

(for MIPS)

swap(int v[], int k) High-level {int temp; language program temp = v[k];v[k] = v[k+1]: (in C) v[k+1] = temp;Compiler Assembly swap: language muli \$2, \$5,4 program add \$2, \$4,\$2 (for MIPS) \$15, 0(\$2) \$16, 4(\$2) \$16, 0(\$2) SW \$15. 4(\$2) \$31 jr Assembler Binary machine 000000010100001000000000011000 language 0000000000110000001100000100001 program

Processor executes the machine code (i.e., binaries)



3. From C to Execution

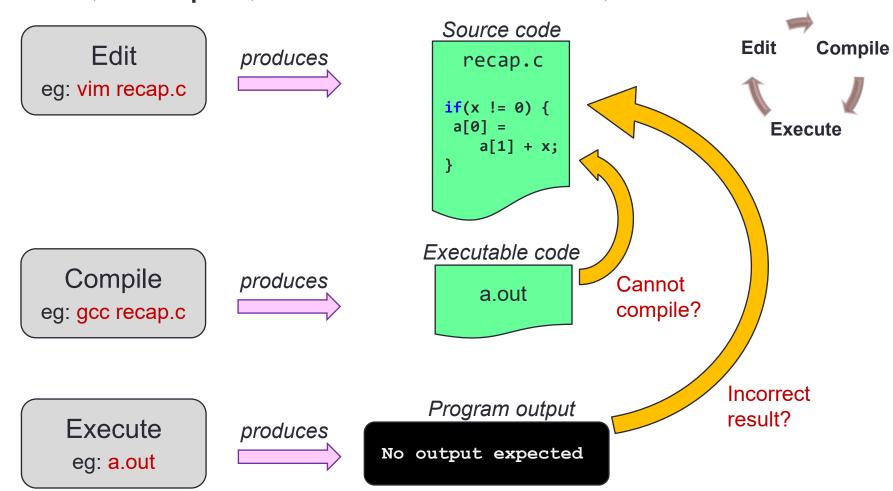
- We play the role of Programmer, Compiler, Assembler, and Processor
 - Program:

```
if(x != 0) {
  a[0] = a[1] + x;
}
```

- Programmer:
 - Show the workflow of compiling, assembling, and executing C program
- Compiler:
 - Show how the program is compiled into MIPS
- Assembler:
 - Show how the MIPS is translated into binaries
- Processor:
 - Show how the datapath is activated in the processor

3.1 Writing C Program

Edit, Compile, Execute: Lecture #2, Slide 5



Key Idea #1:

Compilation is a structured process

```
if(x != 0) {
  a[0] = a[1] + x;
}
```

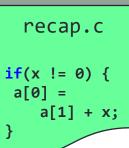
Each structure can be compiled independently

Inner Structure

```
a[0] = a[1] + x;
```

Outer Structure

```
if(x != 0) {
}
```



Key Idea #2: Variable-to-Register Mapping

```
if(x != 0) {
  a[0] = a[1] + x;
}
```

Let the mapping be:

Variable	Register Name	Register Number
X	\$s0	\$16
а	\$s1	\$17

```
recap.c

if(x != 0) {

a[0] =

a[1] + x;

}
```

Common Technique #1: Invert the condition for shorter code (Lecture #8, Slide 22)

Mapping: recap.c x: \$16 **if**(x != 0) { a: \$17 a[0] =a[1] + x;

Outer Structure

```
if(x != 0) {
```

Outer MIPS Code

```
beq $16, $0, Else
        Inner Structure
Else:
```

Common Technique #2:
 Break complex operations, use temp register (Lecture #7, Slide 29)

Inner Structure

a[0] = a[1] + x;

Simplified Inner Structure

```
$t1 = a[1];
$t1 = $t1 + x;
a[0] = $t1;
```

Mapping:

x: \$16

a: \$17

\$t1: \$8

recap.c

if(x != 0) {

a[1] + x;

a[0] =

Common Technique #3:
 Array access is 1w, array update is sw
 (Lecture #8, Slide 13)

Simplified Inner Structure

\$t1 = a[1]; \$t1 = \$t1 + x; a[0] = \$t1;

Inner MIPS Code

```
lw $8, 4($17)
add $8, $8, $16
sw $8, 0($17)
```

Mapping:

x: \$16 a: \$17 \$t1: \$8

```
recap.c

if(x != 0) {

a[0] =

a[1] + x;
```

Mapping:

x: \$16 a: \$17 \$t1: \$8

```
recap.c

if(x != 0) {

a[0] =

a[1] + x;
```

Common Error #1:

Assume that the address of the next word can be found by incrementing the address in a register by 1 instead of by the word size in bytes

Example:

```
$\fint$ $\fint$ = a[1];
is translated to
lw $\fint$8, 4($17)
instead of
lw $\fint$8, 1($17)
```

Mapping:

x: \$16 a: \$17 \$t1: \$8

```
recap.c

if(x != 0) {

a[0] =

a[1] + x;
```

Last Step:

Combine the two structures logically

Inner MIPS Code

Outer MIPS Code

```
lw $8, 4($17)
add $8, $8, $16
sw $8, 0($17)
```

```
beq $16, $0, Else

# Inner Structure

Else:
```

Combined MIPS Code

```
beq $16, $0, Else
lw $8, 4($17)
add $8, $8, $16
sw $8, 0($17)
Else:
```

recap.mips

beq \$16, \$0, Else lw \$8, 4(\$17) add \$8, \$8, \$16 sw \$8, 0(\$17)

Else:

- Instruction Types Used:
 - 1. R-Format: (Lecture #9, Slide 8)
 - opcode \$rd, \$rs, \$rt

6	5	5	5	5	6
opcode	rs	rt	rd	shamt	funct

- 2. I-Format: (Lecture #9, Slide 14)
 - opcode \$rt, \$rs, immediate

6	5	5	16
opcode	rs	rt	immediate

- 3. Branch: (Lecture #9, Slide 22)
 - Uses I-Format
 - PC = (PC + 4) + (immediate × 4)

- beq \$16, \$0, Else
 - Compute immediate value (Lecture #9, Slide 27)
 - immediate = 3
 - Fill in fields (refer to MIPS Reference Data)

```
    6
    5
    5
    16

    4
    16
    0
    3
```

Convert to binary

000100	10000	00000	00000000000011

```
beq $16, $0, Else
lw $8, 4($17)
add $8, $8, $16
sw $8, 0($17)
Else:
```

recap.mips

beq \$16, \$0, Else lw \$8, 4(\$17) add \$8, \$8, \$16 sw \$8, 0(\$17)

Else:

recap.mips

Else:

beq \$16, \$0, Else lw \$8, 4(\$17) add \$8, \$8, \$16 sw \$8, 0(\$17)

1w \$8, 4(\$17)

Fill in fields (refer to MIPS Reference Data)

6	5`	5	16
35	17	8	4

Convert to binary

100011	10001	01000	00000000000100

```
0001 0010 0000 0000 0000 0000 0000 0011

lw $8, 4($17)

add $8, $8, $16

sw $8, 0($17)

Else:
```

recap.mips

Else:

beq \$16, \$0, Else lw \$8, 4(\$17) add \$8, \$8, \$16 sw \$8, 0(\$17)

- add \$8, \$8, \$16
 - Fill in fields (refer to MIPS Reference Data)

6	5`	5	5	5_´	6
0	8	16	8	0	32

Convert to binary

000000	01000	10000	01000	00000	100000

```
0001 0010 0000 0000 0000 0000 0000 0011
1000 1110 0010 1000 0000 0000 0000 0100
add $8, $8, $16
sw $8, 0($17)
Else:
```

recap.mips

beq \$16, \$0, Else lw \$8, 4(\$17) add \$8, \$8, \$16 sw \$8, 0(\$17)

Else:

- sw \$8, 0(\$17)
 - Fill in fields (refer to MIPS Reference Data)

	6	5`	5	16			
	43	17	8	0			
•	■ Convert to hinary						

Convert to binary

101011 10001 01000 000000000000	101011		10001	01000	00000000000000
---------------------------------------	--------	--	-------	-------	----------------

```
0001 0010 0000 0000 0000 0000 0000 0011
1000 1110 0010 1000 0000 0000 0000 0100
0000 0001 0001 0000 0100 0000 0010 0000
sw $8, 0($17)
Else:
```

- Final Binary
 - Hard to read?
 - Don't worry, this is intended for machine not for human!

recap.mips

beq \$16, \$0, Else lw \$8, 4(\$17) add \$8, \$8, \$16 sw \$8, 0(\$17)

Else:

```
0001 0010 0000 0000 0000 0000 0000
                                     0011
          0010
                          0000
1000
     1110
               1000
                    0000
                                0000
                                     0100
0000
     0001 0001 0000
                     0100
                          0000
                                0010
                                     0000
1010 1110 0010 1000 0000 0000 0000
                                     0000
```

3.4 Execution (Datapath)

Given the binary

Assume two possible executions:

```
1. $16 == $0 (shorter)
2. $16 != $0 (Longer)
```

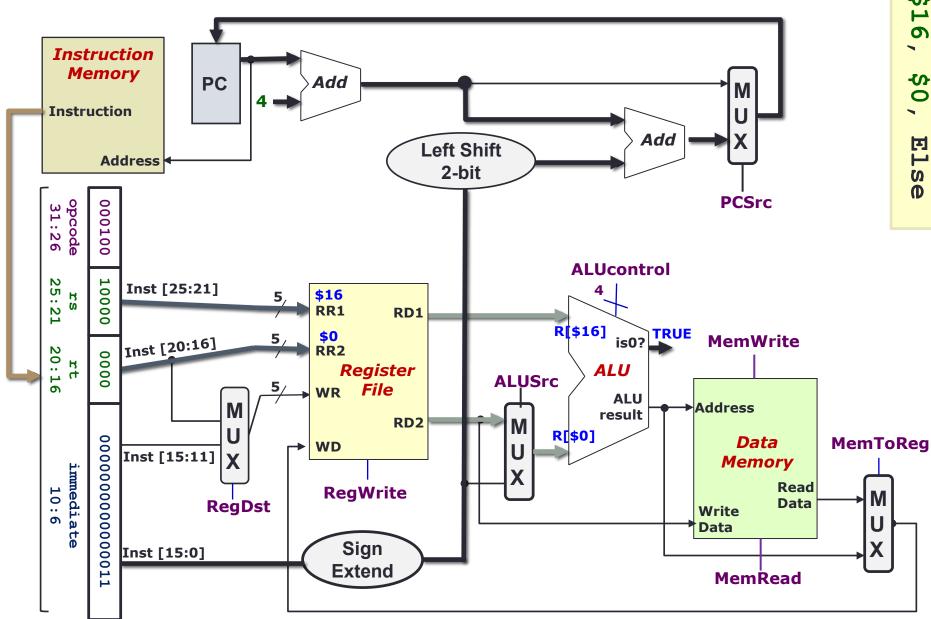
Convention:

```
Fetch: Memory:

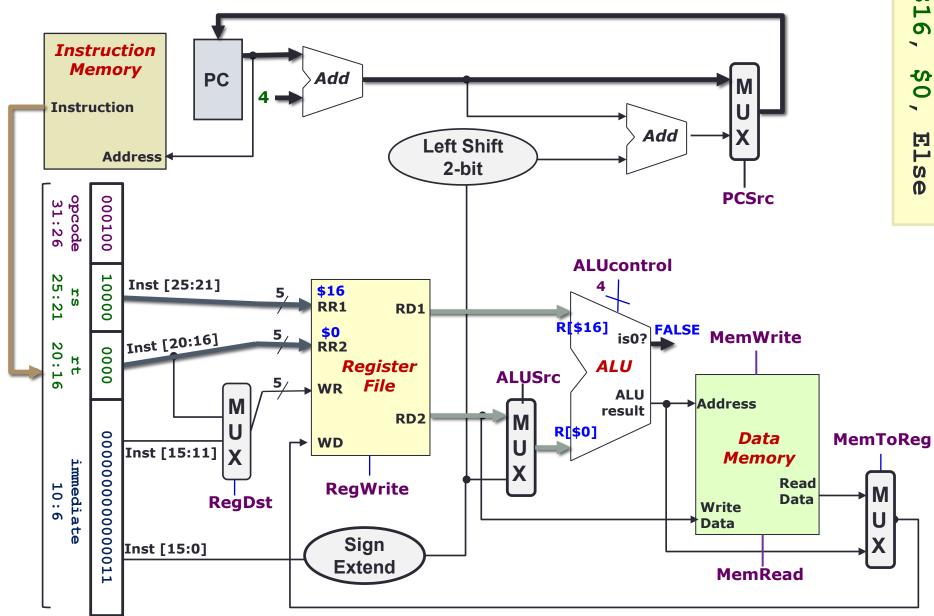
Decode: Reg Write: Other:
```

```
0000
                0000
     0010
                     0000
                           0000
                                 0000
                                      0011
           0010
                1000
                     0000
                           0000
1000
     1110
                                 0000
                                      0100
0000
     0001
           0001
                0000
                      0100
                           0000
                                 0010
                                      0000
     1110 0010 1000 0000 0000 0000
                                      0000
```



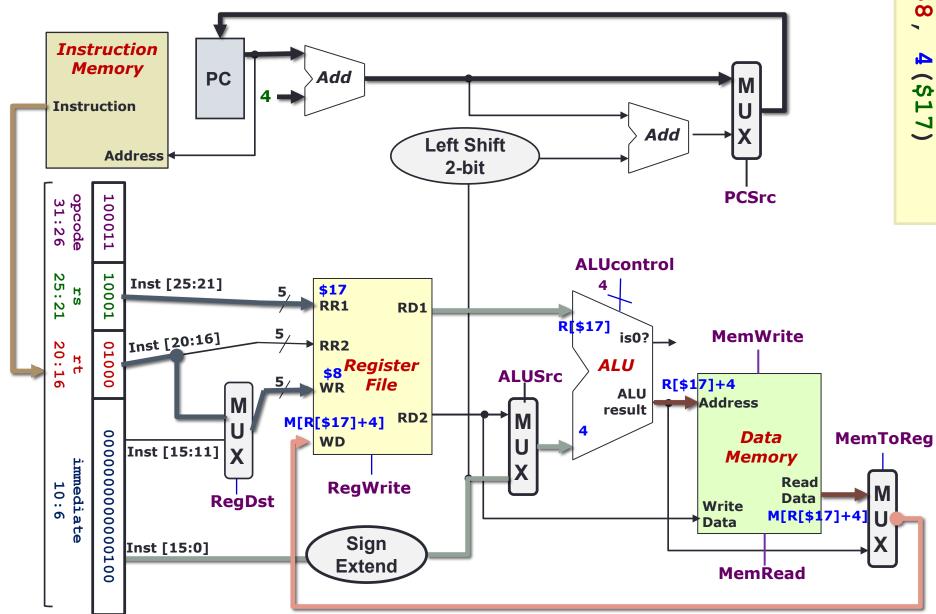


Assume \$16 != \$0

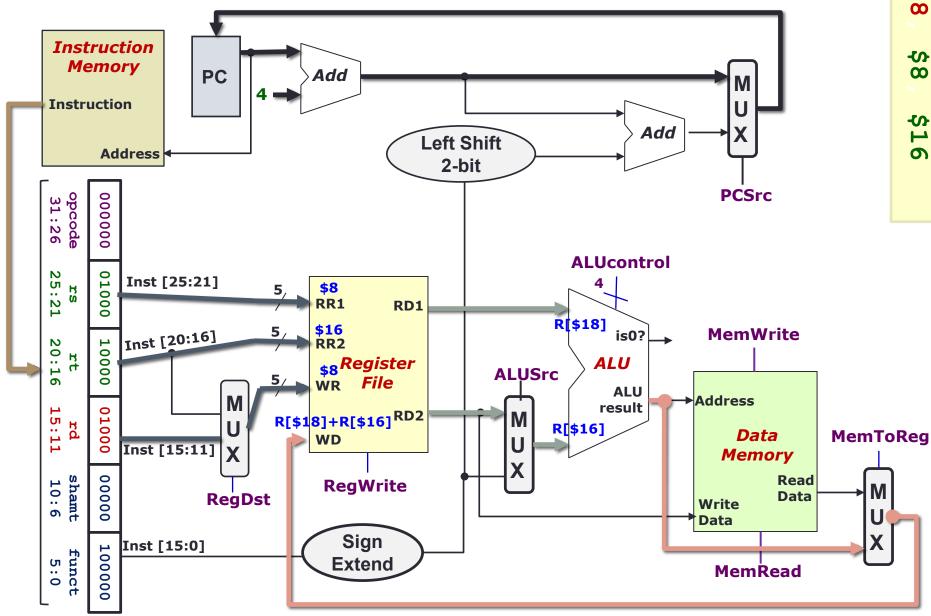


1w

Assume \$16 != \$0

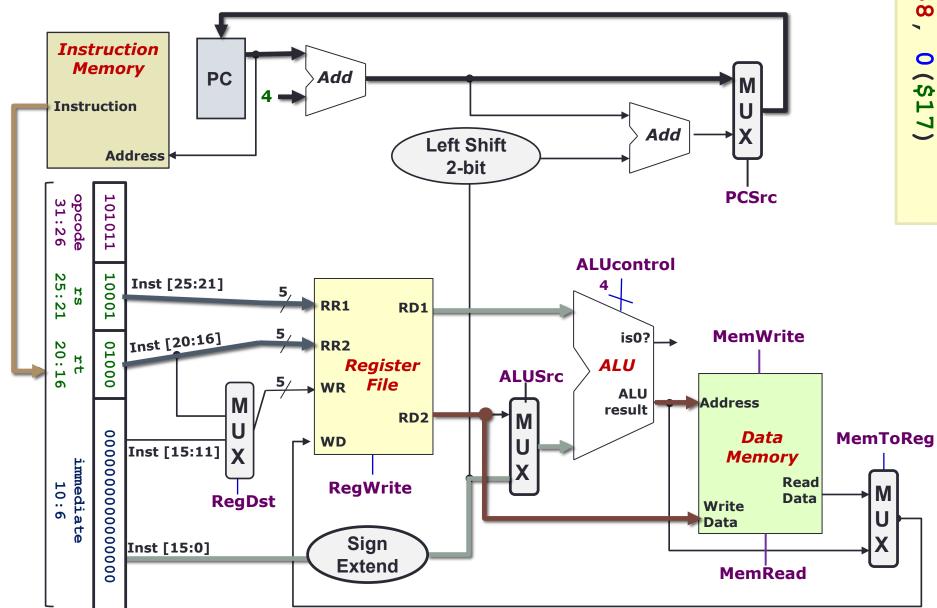






WS

Assume \$16 != \$0



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