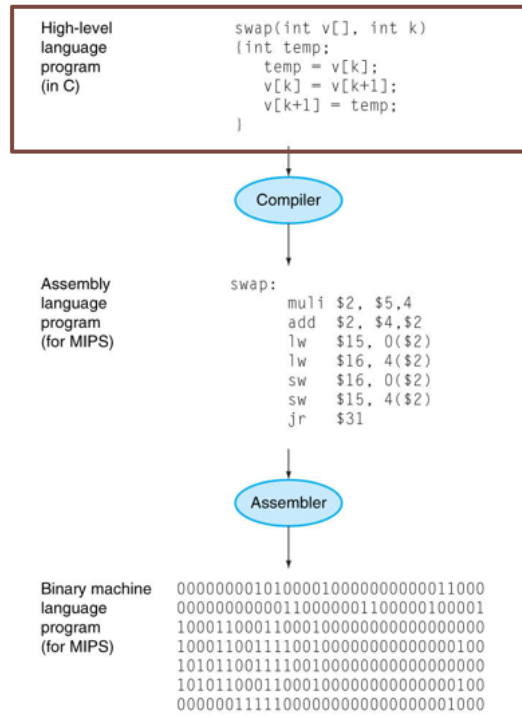


# 7 – The Processor: Datapath

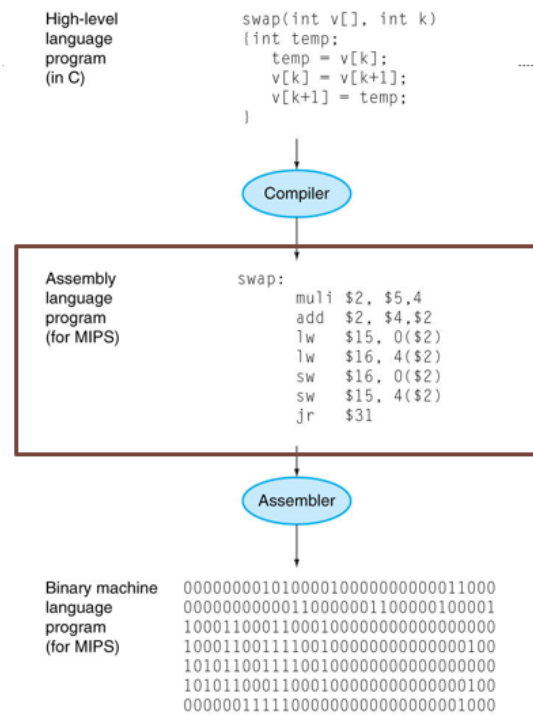
## 7.1 Brief Recap



- Write program in high-level language (e.g., **C**)

```

if(x != 0) {
    a[0] = a[1] + x;
}
  
```

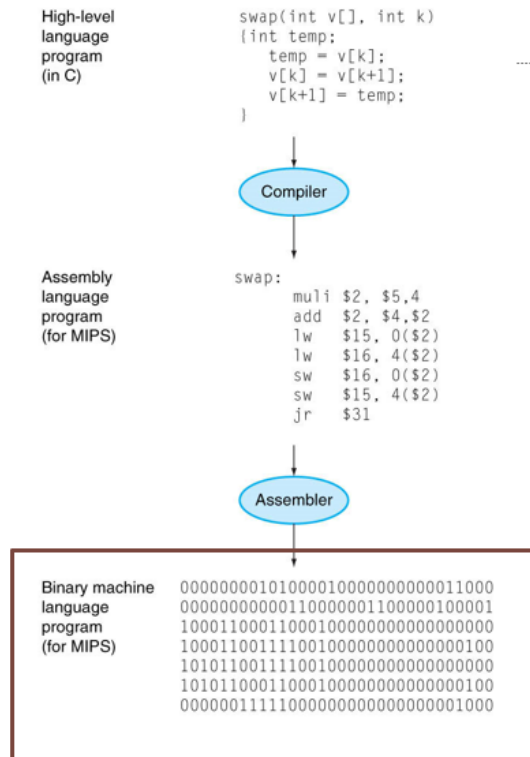


- **Compiler** translates to assembly language (e.g., **MIPS**)

```

beq $16, $0, Else
lw $8, 4($17)
add $8, $8, $16
sw $8, 0($17)

Else:
  
```



➤ **Assembler** translates to machine code (i.e., **binaries**)

```

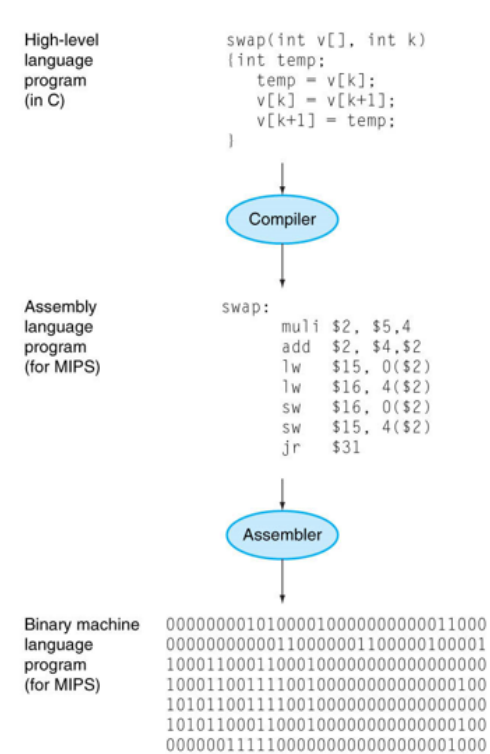
0001 0010 0000 0000
0000 0000 0000 0011

1000 1110 0010 1000
0000 0000 0000 0100

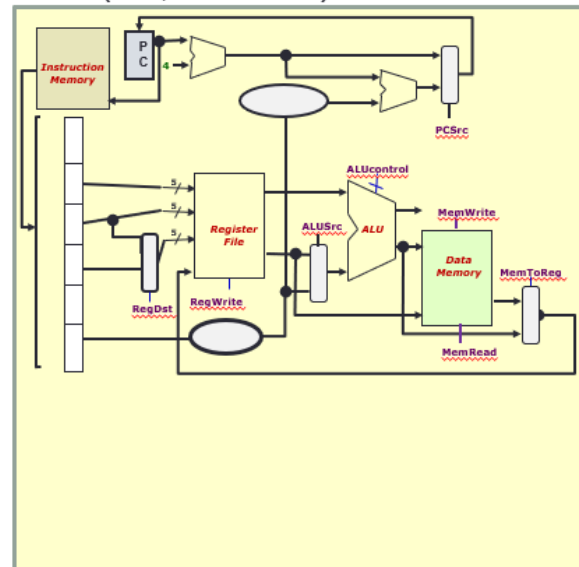
0000 0010 0000 1000
0100 0000 0001 0100

1010 1110 0010 1000
0000 0000 0000 0000

```



➤ **Processor** executes the machine code (i.e., **binaries**)



## 7.2 From C to Execution

- We play the role of Programmer, Compiler, Assembler, and Processor
  - Program:

```

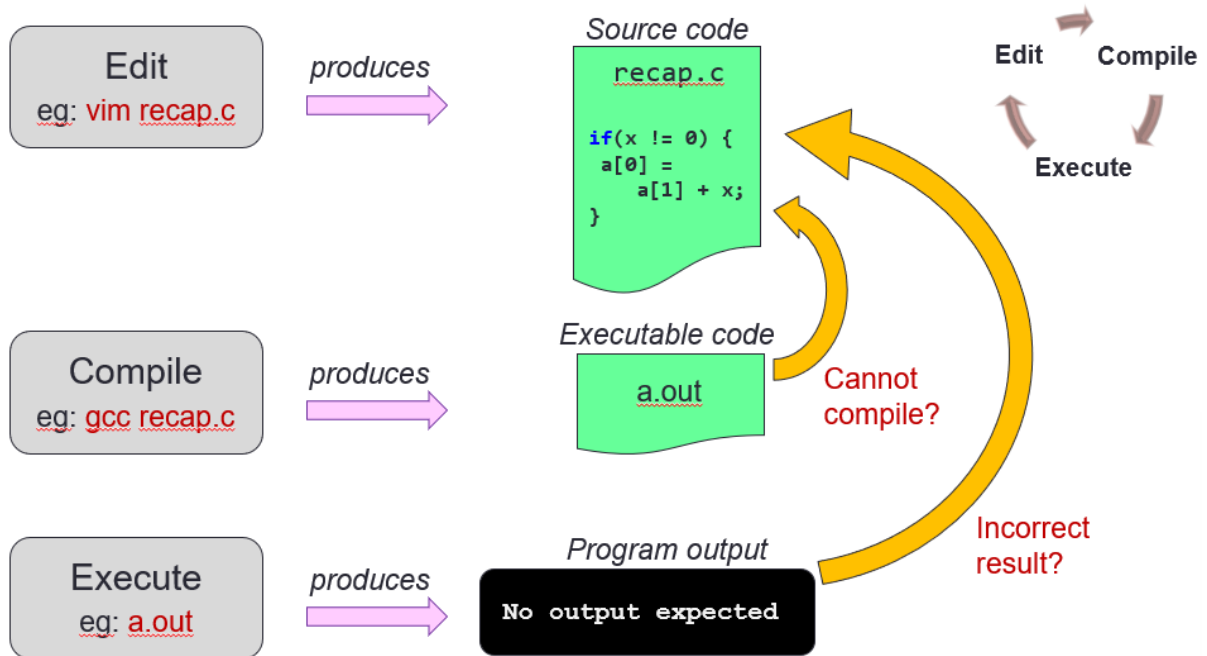
1  if (x != 0) {
2      a[0] = a[1] + x;
3  }

```

- Programmer: Show the workflow of compiling, assembling, and executing C program
- Compiler: Show how the program is compiled into MIPS
- Assembler: Show how the MIPS is translated into binaries
- Processor: Show how the datapath is activated in the processor

### 7.2.1 Writing C Program

#### ■ Edit, Compile, Execute: Lecture #2, Slide 5



### 7.2.2 Compiling to MIPS

#### Key Idea

- Key Idea #1:  
Compilation is a structured process

```
1  if (x != 0) {
2      a[0] = a[1] + x;
3  }
```

Each structure can be compiled independently

#### Inner Structure

```
a[0] = a[1] + x;
```

#### Outer Structure

```
if (x != 0) {

}
}
```

- Key Idea #2:

Variable-to-Register Mapping

Let the mapping be:

Variable	Register Name	Register Number
x	\$s0	\$16
a	\$s1	\$17

### Common Technique

- Common Technique #1:

Invert the condition for shorter code

#### Outer Structure

```
if (x != 0) {  
  
  
  
}
```

#### Outer MIPS Code

```
beq $16, $0, Else  
  
# Inner Structure  
  
Else:
```

- Common Technique #2:

Break complex operations, use temp register

#### Inner Structure

```
a[0] = a[1] + x;
```

#### Simplified Inner Structure

```
$t1 = a[1];  
$t1 = $t1 + x;  
a[0] = $t1;
```

- Common Technique #3:

Array access is `lw`, array update is `sw`

#### Simplified Inner Structure

```
$t1 = a[1];  
$t1 = $t1 + x;  
a[0] = $t1;
```

#### Inner MIPS Code

```
lw $8, 4($17)  
add $8, $8, $16  
sw $8, 0($17)
```

## Common Error

- Common Error #1:

Assume that the address of the next word can be found by incrementing the address in a register by 1 instead of by the word size in bytes

Example:

```
$t1 = a[1]
```

is translated to:

```
lw $8, 4($17)
```

instead of

```
lw $s8, 1($17)
```

## Finalize

- Last Step:

Combine the two structures logically

### Inner MIPS Code

```
lw  $8, 4($17)
add $8, $8, $16
sw  $8, 0($17)
```

### Outer MIPS Code

```
beq $16, $0, Else
```

```
# Inner Structure
```

```
Else:
```

### Combined MIPS Code

```
beq $16, $0, Else
lw  $8, 4($17)
add $8, $8, $16
sw  $8, 0($17)

Else:
```

## 7.2.3 Assembling to Binaries

- Instruction Types Used:

1. R-Format: `opcode $rd, $rs, $rt`



2. I-Format: `opcode $rt, $rs, immediate`



3. Branch:

- Use I-format

- $PC = (PC+4) + (\text{immediate} \times 4)$

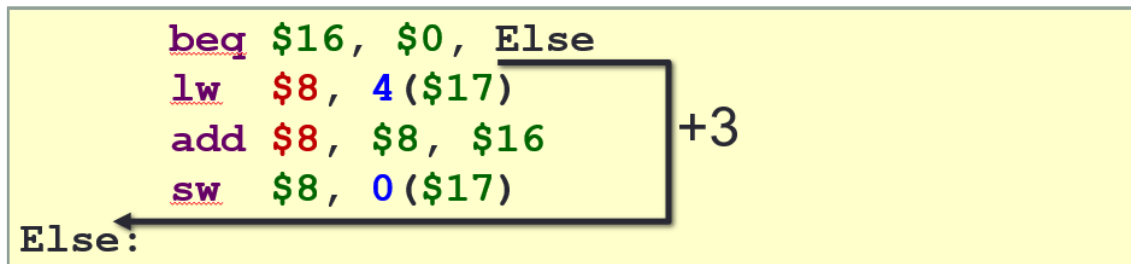
4. `beq $16, $0, Else`

- Compute immediate value
  - `immediate` = 3
- Fill in fields

6	5	5	16
4	16	0	3

- Convert to binary

6	5	5	16
4	16	0	3

5. `lw $8, 4($17)`

- Filled in fields (Refer to MIPS Reference data)

6	5	5	16
35	17	8	4

- Convert to binary

100011	10001	01000	0000000000000100
--------	-------	-------	------------------

0001 0010 0000 0000 0000 0000 0000 0011  
`lw $8, 4($17)`  
`add $8, $8, $16`  
`sw $8, 0($17)`  
 Else:

6. `add $8, $8, $16`

- Filled in fields

6	5	5	5	5	6
0	8	16	8	0	32

- Convert to binary

000000	01000	10000	01000	00000	100000
--------	-------	-------	-------	-------	--------

```

0001 0010 0000 0000 0000 0000 0000 0011
1000 1110 0010 1000 0000 0000 0000 0100
    add $8, $8, $16
    sw  $8, 0($17)
Else:

```

7. `sw $8, 0($17)`

- Filled in fields

6	5	5	16
43	17	8	0

- Convert to binary

101011	10001	01000	0000000000000000
--------	-------	-------	------------------

```

0001 0010 0000 0000 0000 0000 0000 0011
1000 1110 0010 1000 0000 0000 0000 0100
0000 0001 0001 0000 0100 0000 0010 0000
    sw  $8, 0($17)
Else:

```

## 7.2.4 Execution (Datapath)

- Given the binary
  - Assume two possible executions
    1. `$16 == $0` (shorter)
    2. `$16 != $0` (larger)
  - Convention:

Fetch: 

Decode: 

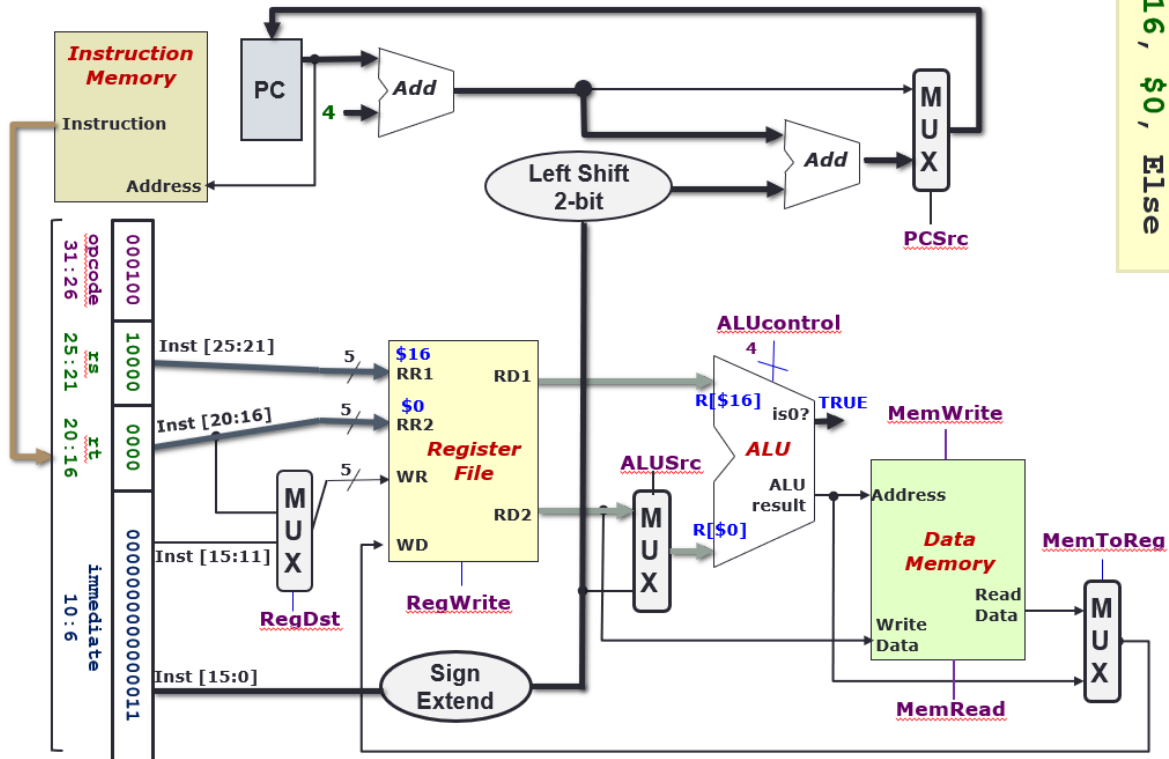
ALU: 

Memory: 

Reg Write: 

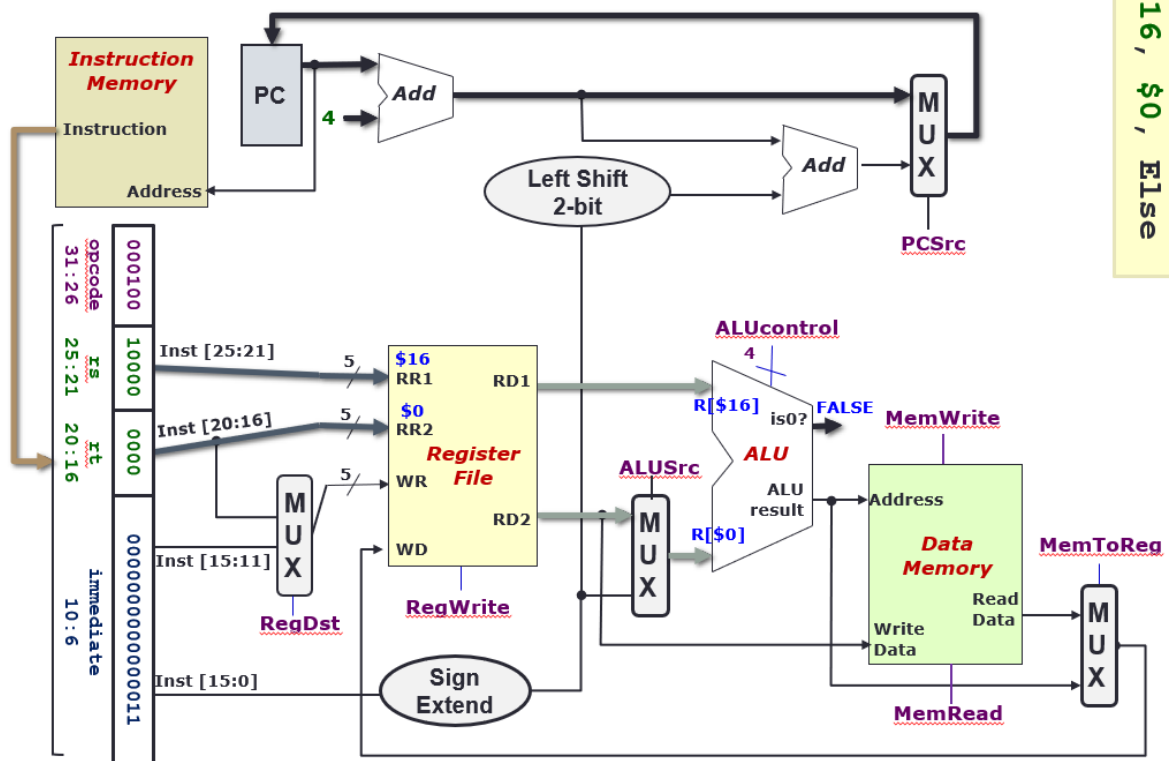
Other: 

- Assume \$16 == \$0



`beq $16, $0, Else`

- Assume \$16 != \$0

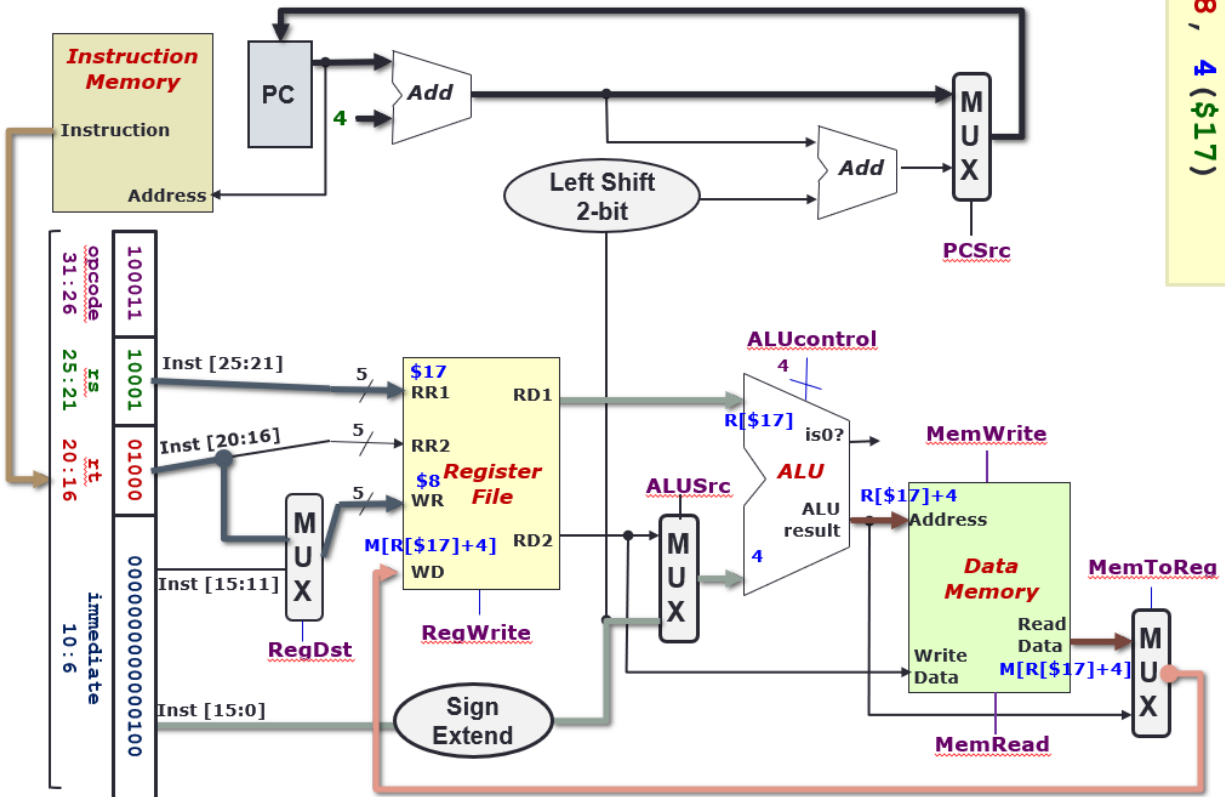


`beq $16, $0, Else`



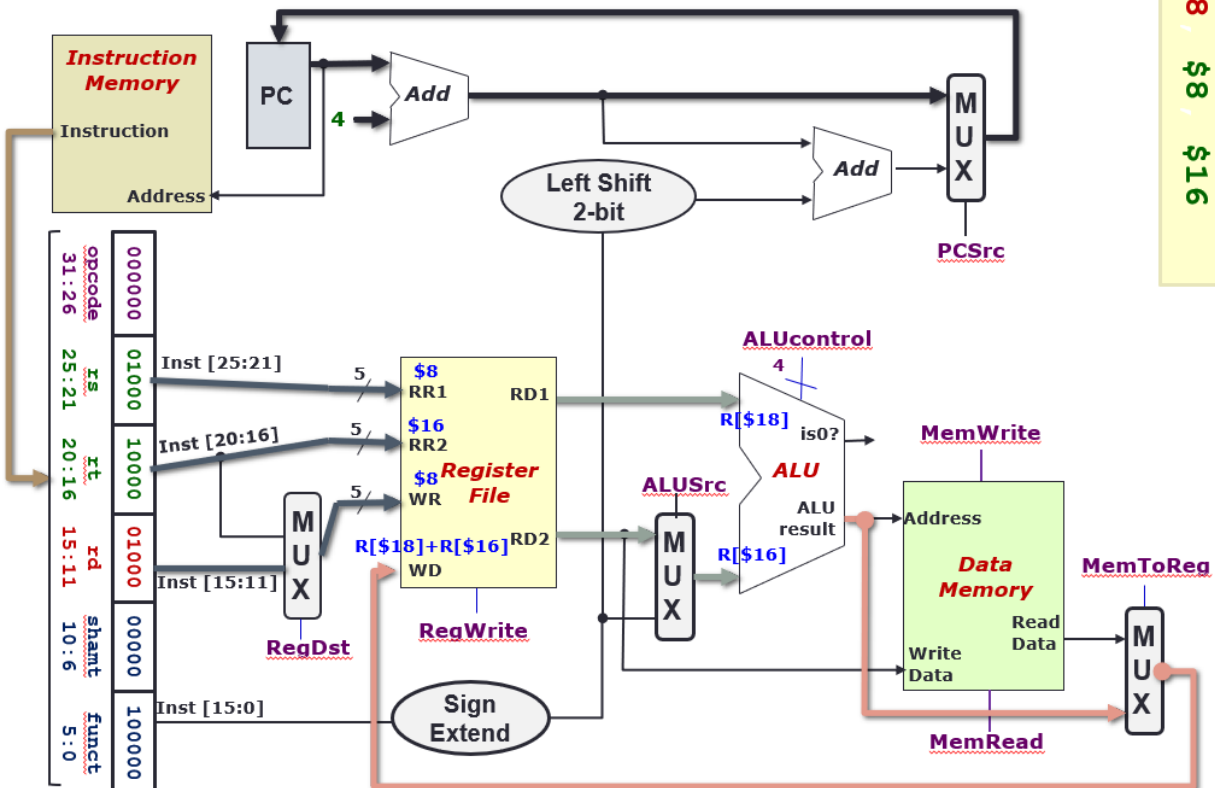
- Assume \$16 != \$0

lw \$8, 4(\$17)



- Assume \$16 != \$0

add \$8, \$8, \$16



- Assume \$16 != \$0

SW \$8, 0(\$17)

