### IT5100A

Industry Readiness:

Typed Functional Programming

### Monads

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```
000
       f :: [Int] -> [Int]
       f(x:xs) =
           let r = f xs
           in r ++ [x]
       main :: IO ()
       main = do
           let x = [1..10]
           print $ f x
```

#### Monads are everywhere!

Today we will learn more about monads and some typical patterns of monads, and how we can compose monads themselves!



# Monad Operations

Control.Monad module defines more operations on monads; how do we use them?

#### **Basic Monad functions** mapM :: (Traversable t, Monad m) => (a -> m b) -> t a -> m (t b) # Source Map each element of a structure to a monadic action, evaluate these actions from left to right, and collect the results. For a version that ignores the results see mapM . Examples mapM\_ :: (Foldable t, Monad m) => (a -> m b) -> t a -> m () Map each element of a structure to a monadic action, evaluate these actions from left to right, and ignore the results. For a version that doesn't ignore the results see mapM. mapM is just like traverse, but specialised to monadic actions. forM :: (Traversable t, Monad m) => t a -> (a -> m b) -> m (t b) for M is map M with its arguments flipped. For a version that ignores the results see for M . forM :: (Foldable t, Monad m) => t a -> (a -> m b) -> m () # Source for M is map M with its arguments flipped. For a version that doesn't ignore the results see for M. forM\_ is just like for\_, but specialised to monadic actions. sequence :: (Traversable t, Monad m) => t (m a) -> m (t a) # Source

# Ignoring Values

```
def my_function(x):
    print(x) # standalone statement
    return x
```

Often we want to write standalone statements to perform some monadic action; makes no sense to bind result to a variable

```
def my_function(x):
   z = print(x) # why?
   return x
```

# Ignoring Values

(>>) method is similar to >>= except we ignore the previous monad value

# Ignoring Values

```
ghci> [1, 2] >>= (\x -> [(x, 3)])
[(1, 3), (2, 3)]
ghci> [1, 2] >>= (\_ -> [3])
[3, 3]
ghci> [1, 2] >> [3]
[3, 3]
```

(>>) on lists does not seem so useful at the moment; we will see more on it shortly

## do Notation

do notation	Translation	Description
do e	е	Expression
do v <- e	e >>= (\v -> do s)	Monadic bind
do e s	e >> (do s)	Monad composition without bind
do let v = e s	let x = e in do s	Pure bind

## Monadic Functions

Standard functions have monadic equivalents with suffix M, function with ignored values have suffix \_\_

```
ghci> map (+2) [1, 2, 3]
[3, 4, 5]
ghci> map (Just . (+2)) [1, 2, 3]
[Just 3, Just 4, Just 5]
ghci> mapM (Just . (+2)) [1, 2, 3]
Just [3, 4, 5]
```

)

## Monadic Functions

```
let x' = map f x becomes do x' <- mapM f x if f is now in context!</pre>
```

```
ghci> import Text.Read
ghci> :{
ghci| toInt :: String -> Int
ghci| toInt = read
ghci| :}
ghci> toInt "123"
123
```

```
ghci> :{
ghci| toIntM :: String -> Maybe Int
ghci| toIntM = readMaybe
ghci| :}
ghci> toIntM "123"
Just 123
ghci> toIntM "hello"
Nothing
```

## Monadic Functions

let x' = map f x becomes do x' <- mapM f x if f is now in context!</pre>

```
ghci> let x = map toInt ["1","2","3"]
    in x
[1,2,3]
ghci> do x <- mapM toIntM ["1","2","3"]
        return x
Just [1,2,3]</pre>
```

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```
def f(x):
    if x > 10:
        print(x)
    return x
```

```
f x = do
   if x > 10
   then someAction x
   else return () -- basically does nothing
   return x
```

Sometimes we want to express monadic controls more elegantly, e.g. only run monadic action when condition is met

Use functions in Control. Monad module!

```
when :: Applicative f => Bool -> f () -> f ()
```

```
import Control.Monad
f x = do
    when (x > 10) (someAction x)
    return x
```

```
guard :: Alternative f :: Bool -> f ()
```

```
safeDiv1 :: Int -> Int -> Maybe Int
safeDiv1 x y = if y == 0
               then Nothing
               else Just (x `div` y)
safeDiv2 :: Int -> Int -> Maybe Int
safeDiv2 x y
    = do guard (y /= 0)
         return $ x \div y
```

```
guard :: Alternative f => Bool -> f ()
guard True = pure ()
guard False = empty
```

An Alternative is something that could be empty

```
ghci> Just () >> (return 1)
Just 1
ghci> Nothing >> (return 1)
Nothing
```

empty bound with another monad should also
give empty, return () >> m gives m

```
ghci> import Control.Monad
ghci> ls = [-2, -1, 0, 1, 2]
ghci> :{
ghci > 1s2 = do x < - 1s
ghci
    guard (x > 0)
ghci return x
ghci :}
ghci> 1s2
[1, 2]
```

Lists are also Alternatives; guard places a guard on each element in the list, basically a filter!

```
ghci> import Control.Monad
ghci> ls = [-2, -1, 0, 1, 2]
ghci> :{
ghci > 1s2 = do x < - 1s
ghci
    guard (x > 0)
ghci return $ x * 2
ghci :}
ghci> ls2
[2, 4]
```

Familiar?

List comprehension is just a specialized form of monadic binds and guards! Do notation allows us to express monadic operations in any order, giving you maximum control!

```
ghci> ls = [-2, -1, 0, 1, 2]
ghci> ls2 = [x * 2 | x <- ls, x > 0]
ghci> ls2
[2, 4]
```



# Commonly Used Monads

```
These types act as contexts or notions of computation:

Maybe a — an a or nothing

Either a b — either a or b

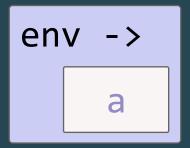
[a] — a list of possible as (nondeterminism)
```

```
Maybe | Either Int | [] | Char
```

Other common **notions of computation**:

- Reading from state
- Writing to, or editing state

## Reader



An a that depends on an environment env

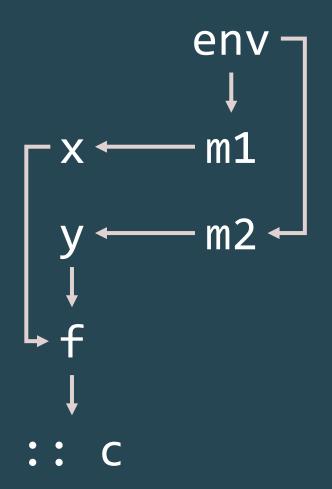
#### env -> is a monad!

```
instance Functor ((->) env) where
    fmap :: (a -> b) -> (env -> a) -> (env -> b)
    fmap f x = f \cdot x
instance Applicative ((->) env) where
    pure :: a -> (env -> a)
    pure = const
    (\langle * \rangle) :: (env -> (a -> b)) -> (env -> a) -> env -> b
    (\langle * \rangle) f g x = f x (g x)
instance Monad ((->) env) where
    return :: a -> (env -> a)
    return = pure
    (>>=) :: (env -> a) -> (a -> (env -> b)) -> env -> b
    (>>=) m f x = f (m x) x
```

## Reader: Intuition

```
m1 :: env -> a
m2 :: env -> b
f :: a -> b -> c
```

```
do x <- m1
   y <- m2
   return $ f x y</pre>
```



Example: obtaining the neighbouring nodes in a graph

```
type Reader = (->)
type Node = Int
type Graph = [(Node, [Node])]

getNeighbours :: Node -> Reader Graph [Node]
getNeighbours x = do
    neighbours <- lookup x
    return $ concat neighbours</pre>
```

```
lookup :: a -> [(a, b)] -> Maybe b
lookup :: Node -> Reader Graph (Maybe [Node])

concat :: Foldable t => t [a] -> [a]
concat :: Maybe [Node] -> [Node]
```

#### Now we can perform DFS; graph is not mentioned anywhere!

```
dfs :: Node -> Node -> Reader Graph Bool
dfs src dst = aux [] src where
  aux :: [Node] -> Node -> Reader Graph Bool
  aux visited current
      arrived = return True
      alreadyVisited = return False
      otherwise = do
      neighbours <- getNeighbours current</pre>
      ls <- mapM (aux (current : visited)) neighbours</pre>
      return $ or 1s
      where arrived = current == dst
            alreadyVisited = current `elem` visited
```

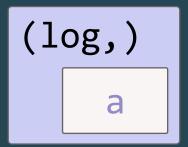
```
ghci> my_map = [(1, [2, 3])]
              , (3, [1, 4])
              (4, [3])
             , (5, [6])
          , (6, [5])
ghci> dfs 5 6 my map
True
ghci> dfs 5 2 my_map
False
ghci> dfs 1 2 [] -- empty
False
```

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Reader monad abstracts the environment from programmers; however, can directly get the environment using ask function

```
ask :: Reader env env
ask = id

getNeighbours :: Node -> Reader Graph [Node]
getNeighbours x = do
    my_graph <- ask -- gets the graph directly
    let neighbours = lookup x my_graph
    return $ concat neighbours</pre>
```



An a that also has some log

#### (log,) is a monad!

```
instance Functor (log,) where
    fmap :: (a -> b) -> (log, a) -> (log, b)
    fmap f (log, a) = (log, f a)
instance Monoid log => Applicative (log,) where
    pure :: a -> (log, a)
    pure = (mempty,)
    (\langle * \rangle) :: (\log, a \rightarrow b) \rightarrow (\log, a) \rightarrow (\log, b)
    (<*>) (log1, f) (log2, x) = (log1 mappend log2, f x)
instance Monoid log => Monad (log,) where
    return :: a -> (log, a)
    return = pure
    (>>=) :: (\log, a) \to (a \to (\log, b)) \to (\log, b)
    (\log, a) >>= f = let (\log 2, b) = f a
                       in (log1 mappend log2, b)
```

## Monoids

Our log must be combine-able and have an empty value  $\varepsilon$  that behaves in the most obvious way

$$E \oplus \varepsilon = \varepsilon \oplus E = E$$
  
$$E_1 \oplus (E_2 \oplus E_3) = (E_1 \oplus E_2) \oplus E_3$$

Make the log type a **Monoid, which are Semigroups with empty values!** 

```
class Semigroup a => Monoid a where
  mempty :: a
  mappend :: a -> a -> a
```

## Monoids

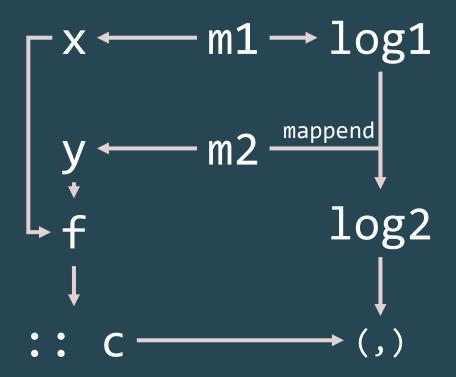
As per usual, lists are monoids with respect to the empty list and list concatenation!

```
instance Monoid [a] where
  mempty = []
  mappend = (++)
```

## Writer: Intuition

```
m1 :: (log, a)
m2 :: (log, b)
f :: a -> b -> c
```

```
do x <- m1
   y <- m2
   return $ f x y</pre>
```



ask for Readers retrieve environment, write for Writers writes to log

```
write :: w -> (w, ())
write = (,())
```

Example of logged addition:

```
type Writer = (,)
type Log = [String]

loggedAdd :: Int -> Int -> Writer Log Int
loggedAdd x y = do
    let z = x + y
    write [show x ++ " + " ++ show y ++ " = " ++ show z]
    return z
```

We can use loggedAdd to define a logged sum!

```
loggedSum :: [Int] -> Writer Log Int
loggedSum [] = return 0
loggedSum (x:xs) = do
    sum' <- loggedSum xs
    loggedAdd x sum'</pre>
```

```
ghci> y = loggedSum [1, 2, 3]
ghci> snd y
6
ghci> fst y
["3 + 0 = 3", "2 + 3 = 5", "1 + 5 = 6"]
```

## State

Stateful operation combines Reader and Writer; is a function that receives state and produces something with a new state

```
randomInt :: Seed -> (Int, Seed)
```

## State

```
newtype State s a = State { runState :: s -> (a, s) }
```

newtype declaration: same as data declaration with one constructor over one field; introduces no operational overhead

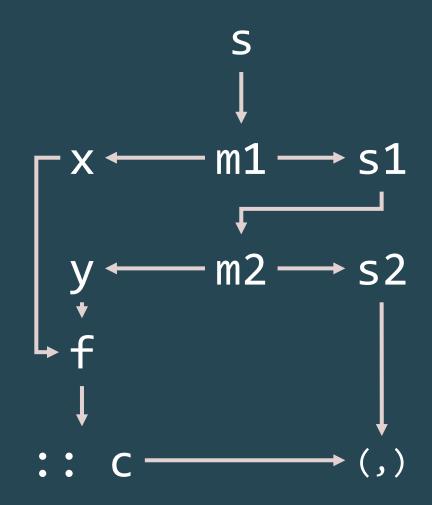
```
instance Functor (State s) where
    fmap :: (a -> b) -> State s a -> State s b
    fmap f (State f') = State $
       \slash -> let (a, s') = f' s
              in (f a, s')
instance Applicative (State s) where
    pure :: a -> State s a
    pure x = State(x,)
    (\langle * \rangle) :: State s (a -> b) -> State s a -> State s b
    (<*>) (State f) (State x) = State $
        \slashs -> let (f', s') = f s
                  (X', S'') = X S'
              in (f' x', s'')
instance Monad (State s) where
    return :: a -> State s a
    return = pure
    (>>=) :: State s a -> (a -> State s b) -> State s b
    (State f) >>= m = State $
        \slash -> let (a, s') = f s
                  State f' = m a
              in f's'
```

basic instances

## State: Intuition

```
m1 :: s -> (a, s)
m2 :: s -> (b, s)
f :: a -> b -> c
```

```
do x <- m1
    y <- m2
    return $ f x y</pre>
```



### State

Helper functions to retrieve and update state just like ask and write

```
put :: s -> State s ()
put s = State $ const ((), s)
get :: State s s
get = State $ \s -> (s, s)
modify :: (s \rightarrow s) \rightarrow State s ()
modify f = do s <- get
                put (f s)
```

## State

Example stateful operation: memoized fibonacci

```
type Memo = [(Integer, Integer)]
getMemoized :: Integer -> State Memo (Maybe Integer)
getMemoized n = lookup n <$> get
```

```
fib :: Integer -> Integer
fib n = fst $ runState (aux n) [] where
  aux :: Integer -> State Memo Integer
  aux 0 = return 0
  aux 1 = return 1
  aux n = do
    x <- getMemoized n
    case x of
        Just y -> return y
        Nothing -> do
            r1 <- aux (n - 1)
            r2 < -aux (n - 2)
            let r = r1 + r2
            modify ((n, r) :)
            return r
```

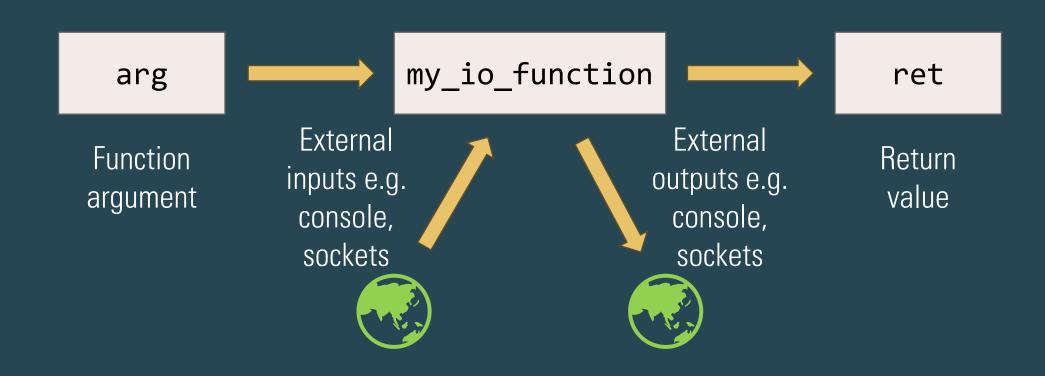
### State

Example stateful operation: memoized fibonacci

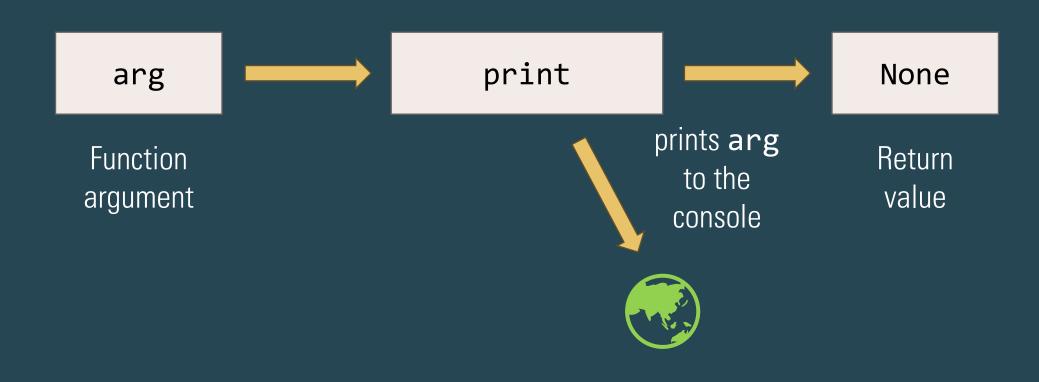
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Finally, we shall learn how to write a Hello World! Program. However, how does a purely functional program produce side effects?

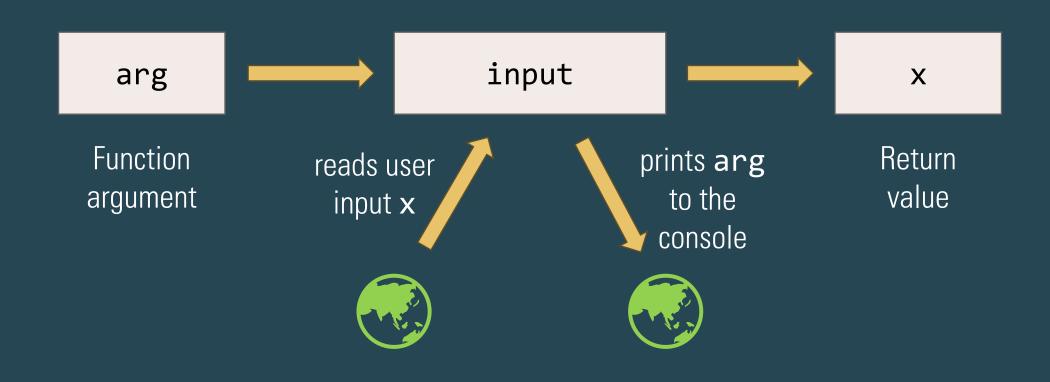
# I/O in Other Languages



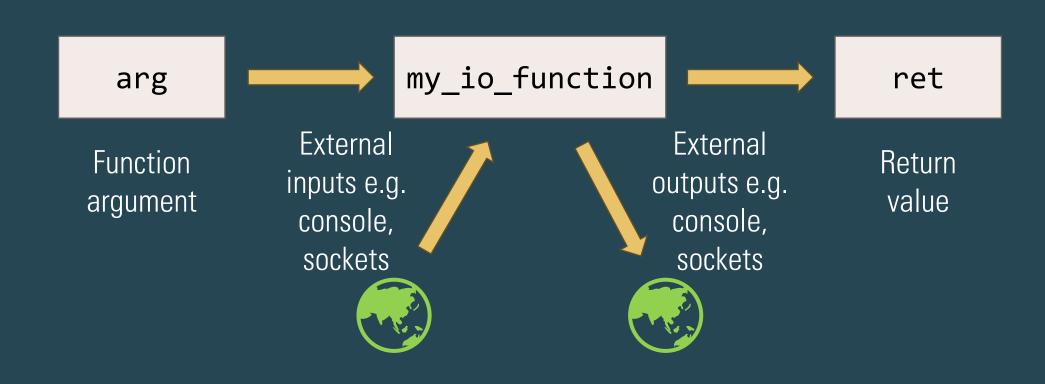
# I/O in Other Languages



# I/O in Other Languages



#### What if the world (external environment) was a term? :0



#### Group inputs as tuple, outputs as tuple

arg



my\_io\_function



ret

Function argument

External inputs e.g. console, sockets

External outputs e.g. console, sockets

Return value

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#### Currying

arg



my\_io\_function



ret

Function argument

External inputs e.g. console, sockets

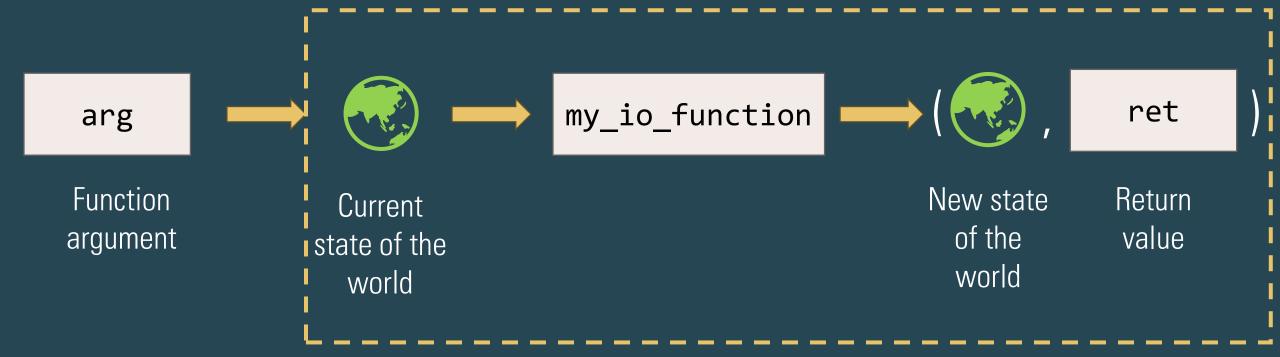
External outputs e.g. console, sockets

Return value

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#### State monad?



State ret

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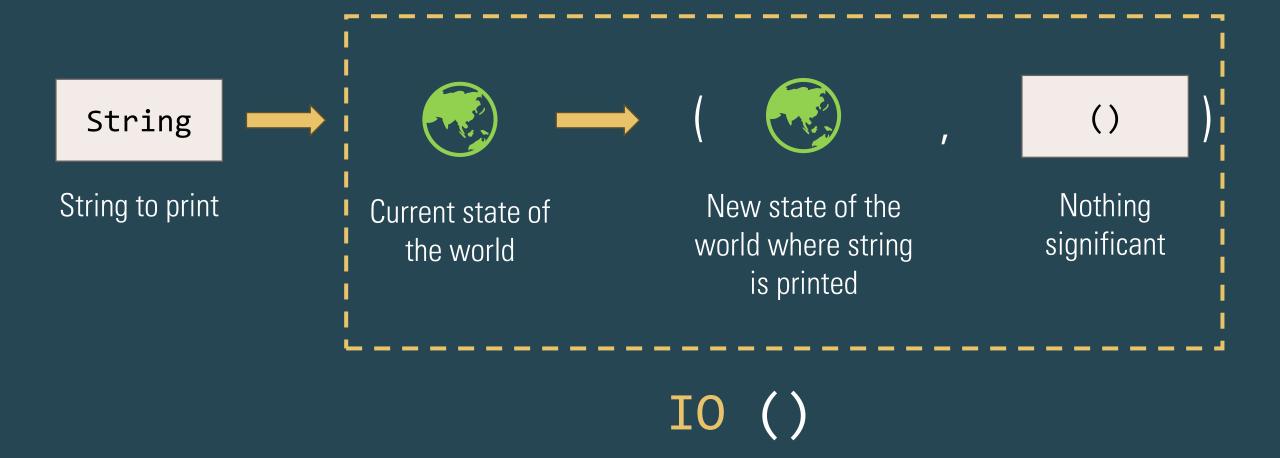
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# 1/0

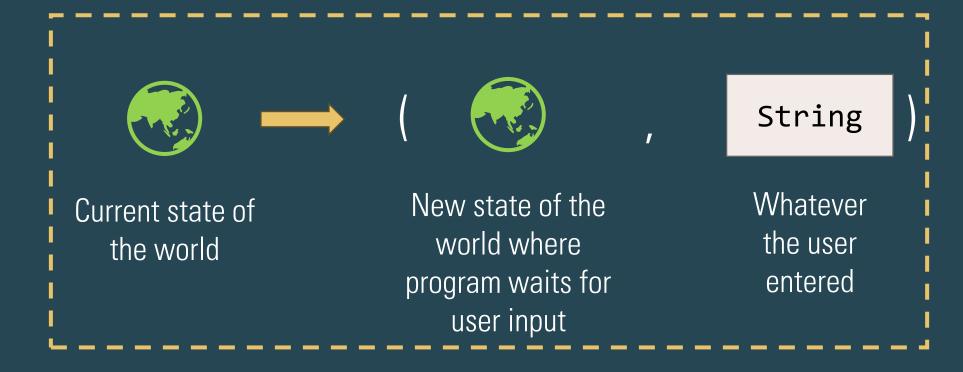
IO monad is the State monad whose state is the real world

```
newtype IO a = IO (State# RealWorld -> (# State# RealWorld, a #))
```

#### putStrLn :: String -> IO ()



#### getLine :: IO String



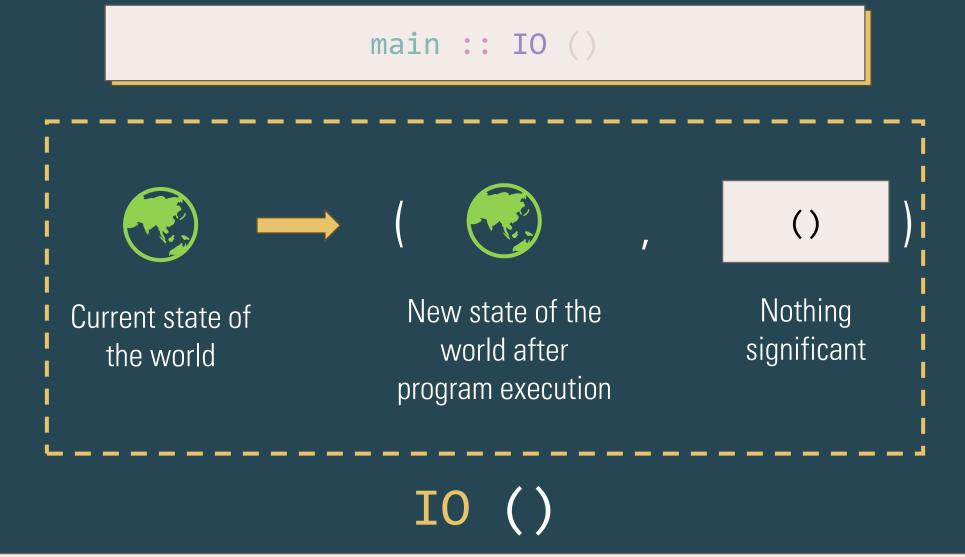
### IO String

Composing I/O actions is as simple as using the State monad!

RealWorld term never needs to be accessed or used

```
-- Main.hs
main :: IO ()
main = do
    name <- getLine
    putStrLn $ "Hello " ++ name ++ "!"</pre>
```

IO monad abstracts RealWorld state and prevents impurity in pure functions. All functions performing I/O actions must be done within the IO monad... enforced by type system!



When a program is run, the Haskell runtime "passes in the current state of the world" into main, and "returns the new state of the world after program execution"



Length of path between two connected edges in graph, where each node only has one neighbouring node

```
type Node = Int
type Graph = [(Node, Node)]
dfs :: Node -> Node -> Graph -> Maybe Int
dfs src dst gph = aux src [] gph where
    aux :: Node -> [Node] -> Graph -> Maybe Int
    aux current visited gph'
        arrived = return ∅
        alreadyVisited = Nothing
        otherwise = do
          n <- lookup current gph
          (+1) <$> aux n (current : visited) gph'
      where arrived = current == dst
            alreadyVisited = current `elem` visited
```

Currently using the Maybe monad, can we also add Reader?

```
type Node = Int
type Graph = [(Node, Node)]
dfs :: Node -> Node -> Reader Graph (Maybe Int)
dfs src dst = aux src [] where
    aux :: Node -> [Node] -> Reader Graph (Maybe Int)
    aux current visited
       arrived = return 0
       alreadyVisited = Nothing
       otherwise = do
          n <- lookup current</pre>
          (+1) <$> aux n (current : visited)
      where arrived = current == dst
            alreadyVisited = current `elem` visited
```

Naively changing the type to use **Reader** and **Maybe** and removing occurrences of graph will not type check as **do** block works in **Reader** context, not **Maybe**!

## **Enriching Readers with Maybe**

```
newtype ReaderMaybe env a = ReaderMaybe { runReaderMaybe :: Reader env (Maybe a)
instance Monad (ReaderMaybe env) where
  return :: a -> ReaderMaybe env a
  return = pure
  (>>=) :: ReaderMaybe env a -> (a -> ReaderMaybe env b) -> ReaderMaybe env b
  (ReaderMaybe 1s) >>= f = ReaderMaybe $ do
   m < -1s
    case m of
      Just x -> runReaderMaybe $ f x
     Nothing -> return Nothing
```

ReaderMaybe monad is a combination of Reader and Maybe!

## **Enriching Readers with Maybe**

```
dfs :: Node -> Node -> Graph -> Maybe Int
dfs src dst = runReaderMaybe (aux src []) where
    aux :: Node -> [Node] -> ReaderMaybe Graph Int
    aux current visited
        arrived = return 0
        alreadyVisited = ReaderMaybe $ return Nothing
        otherwise = do
          n <- ReaderMaybe $ lookup current</pre>
          (+1) <$> aux n (current : visited)
      where arrived = current == dst
            alreadyVisited = current `elem` visited
```

Now we can use the **ReaderMaybe** monad in our **dfs** function; graph not mentioned anywhere and we still have **Maybe** behaviour!

## **Enriching IO with Maybe**

```
newtype IOMaybe a = IOMaybe { runIOMaybe :: IO (Maybe a) }
instance Monad IOMaybe where
  return :: a -> IOMaybe a
  return = pure
  (>>=) :: IOMaybe a -> (a -> IOMaybe b) -> IOMaybe b
  (IOMaybe m) >>= f = IOMaybe $ do
   maybe m <- m
    case maybe m of
      Just x -> runIOMaybe $ f x
     Nothing -> return Nothing
```

Combining IO and Maybe requires us to go through the same process!

## **Enriching IO with Maybe**

```
instance Monad (ReaderMaybe env) where
return = pure

(ReaderMaybe ls) >>= f = ReaderMaybe $ do
    m <- ls
    case m of
    Just x -> runReaderMaybe $ f x
    Nothing -> return Nothing
```

```
instance Monad IOMaybe where
  return = pure

(IOMaybe m) >>= f = IOMaybe $ do
    maybe_m <- m
    case maybe_m of
    Just x -> runIOMaybe $ f x
    Nothing -> return Nothing
```

Notice any similarities?

### **Monad Transformers**

```
newtype MaybeT m a = MaybeT { runMaybeT :: m (Maybe a) }
instance (Monad m) => Monad (MaybeT m) where
  return = MaybeT . return . Just
  x >>= f = MaybeT $ do
    v <- runMaybeT x
    case v of
        Nothing -> return Nothing
        Just y -> runMaybeT (f y)
```

MonadT enriches m with Monad

### **Monad Transformers**

```
dfs :: Node -> Node -> Graph -> Maybe Int
dfs src dst = runMaybeT (aux src []) where
    aux :: Node -> [Node] -> MaybeT (Reader Graph) Int
    aux current visited
        arrived = return 0
        alreadyVisited = MaybeT $ return Nothing
        otherwise = do
          n <- MaybeT $ lookup current</pre>
          (+1) <$> aux n (current : visited)
      where arrived = current == dst
            alreadyVisited = current `elem` visited
```

Directly use the MaybeT monad transformer in our dfs implementation!

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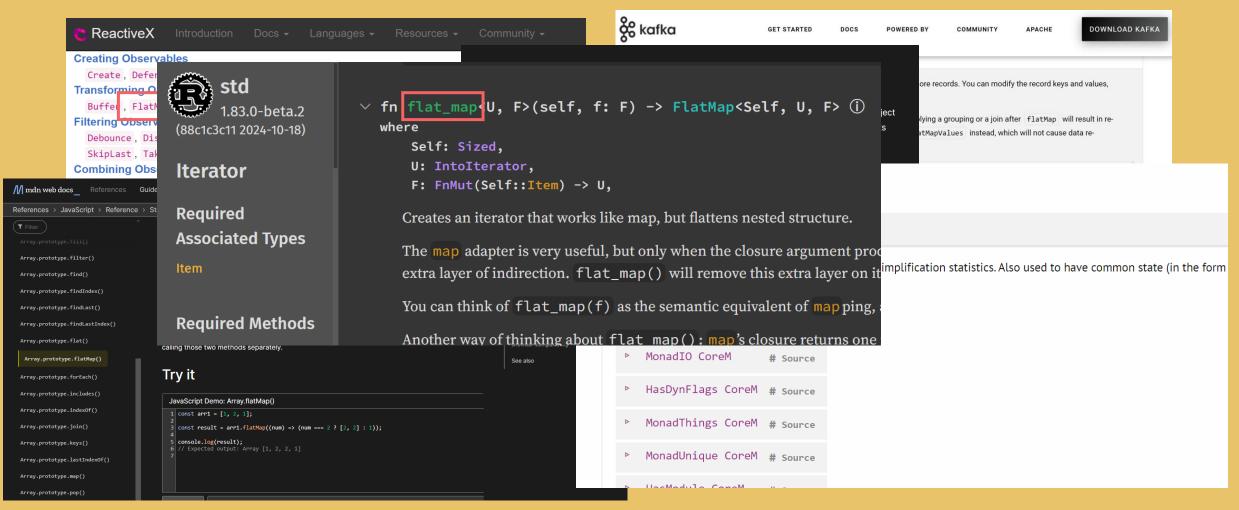
### **Monad Transformer Libraries**

Monads are so common—transformers and mtl libraries contain common monad transformers. Download dependencies and use them!

To work with projects with dependencies easily, use a package manager/build tool like cabal or stack!



# Monads Are Everywhere



# Creating Monads?

- Does your library involve nondeterminism or streams/lists of data?
- Does your library perform I/O?
- Does your library produce potentially empty computation?
- Does your library potentially fail?
- Does your library read from an environment?
- Does your library write to state?
- Does your library process state?

If yes, create a monad!

# Thank you

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