



Bauman Moscow State University
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Finite State Machines and Regular Expressions



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Lecture Outline

- 1 Basic Notions
- 2 Closures and Determinisation
 - ε -Removal by Closure
 - Subset Construction and Determinisation
- 3 From NFA to Regular Expressions
 - Solving Language Equations
 - State Eliminating Method



Regular Expressions by Kleene

☹☹ Academic Definition

Given alphabet Σ , a regular expression is either a letter in Σ , ε , or a result of following operations, where r_1, r_2 are regular expressions:

- $r_1 \mid r_2$ — union (alternation). $\mathcal{L}(r_1 \mid r_2) = \mathcal{L}(r_1) \cup \mathcal{L}(r_2)$;
- $r_1 r_2$ — concatenation (sequencing).
 $\mathcal{L}(r_1 r_2) = \{\omega_1 \omega_2 \mid \omega_1 \in \mathcal{L}(r_1) \ \& \ \omega_2 \in \mathcal{L}(r_2)\}$;
- $(r_1)^*$ — iteration (0 or more concatenations of r_1 with itself);

$$\mathcal{L}((r_1)^*) = \{\varepsilon\} \bigcup_{i=1}^{\infty} \mathcal{L}(r_1).$$

Syntactic Sugar

- r^+ — positive iteration (shortcut for $r r^*$);
- $r?$ — option (shortcut for $(r \mid \varepsilon)$).



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Priorities: star > concatenation > union.

$$ab^* \mid c^*d \Leftrightarrow \left(a(b^*) \right) \mid \left((c^*)d \right).$$



Terminological Clash

Academic regexes

- |, ., * (sometimes +, ?) operations;
- define regular languages;
- studied in university courses (compilers & formal languages)

REGEX (extended regexes)

- lookaheads, backreferences, etc;
- define non-context-free languages;
- used in practice (PCRE2 standart).

- Almost identical names are used for completely different (although related) notions.



Occam Razor: Non-Deterministic Finite Automata

Only excitatory signals are left on there, and all inner neurons fire whenever there is at least one input signal.

☹☹ Definition

A non-deterministic finite automaton (NFA) is a tuple

$\mathcal{A} = \langle Q, \Sigma, q_0, F, \delta \rangle$, where:

- *Q — state set;*
- *Σ — terminal alphabet;*
- *$\delta : Q \times (\Sigma \cup \{\varepsilon\}) \rightarrow 2^Q$ — transition rules;*
- *$q_0 \in Q$ — starting state;*
- *$F \subseteq Q$ — final states.*

Sometimes we use notation:

$\langle q_1, a, q_2 \rangle \in \delta \Leftrightarrow \langle q_1, a, M \rangle \in \delta \ \& \ q_2 \in M.$

Or, usually, simply: $q_1 \xrightarrow{a} q_2.$



Asymmetry of NFA Definition

- Classical works (Kleene, Brzozowski): multiple NFA starting states are allowed.
- Modern formal language theory: the unique starting state in NFA is assumed.
- Equivalent (we can add an unique starting state with ε -transitions to the multiple states), but confusing (e.g. in Brzozowski minimisation).

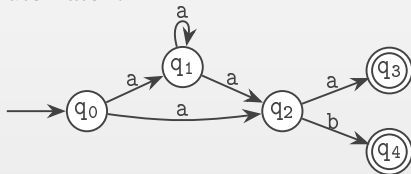


Encoding into Grammars

Observation

- Transition $q_1 \xrightarrow{a} q_2$ can be seen as a rewriting rule $[q_1] \rightarrow a[q_2]$, assuming that $[q_i]$ are some intermediate constructors, while $a \in \Sigma$ is a terminal constructor.
- In order to model computation termination, for every final state q_F , we can add the rewriting rule $[q_F] \rightarrow \varepsilon$.

Automaton:



Grammar:

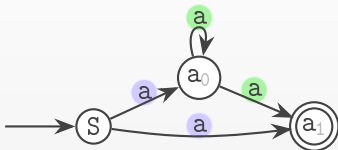
$S \rightarrow a[q_1]$	$[q_2] \rightarrow a[q_3]$
$S \rightarrow a[q_2]$	$[q_2] \rightarrow b[q_4]$
$[q_1] \rightarrow a[q_1]$	$[q_3] \rightarrow \varepsilon$
$[q_1] \rightarrow a[q_2]$	$[q_4] \rightarrow \varepsilon$

We rename the starting nonterminal $[q_0]$ to S , for uniformity.

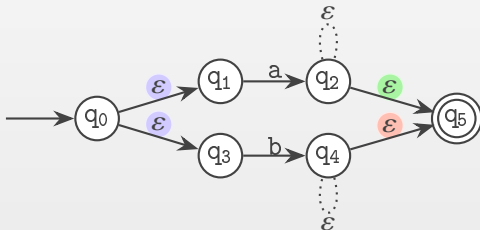


Sources of Non-Determinism in an NFA

- Transition sets wrt (with respect to) a letter $\gamma \in \Sigma$ that are not singletons.



- ε -transitions (so-called silent actions).

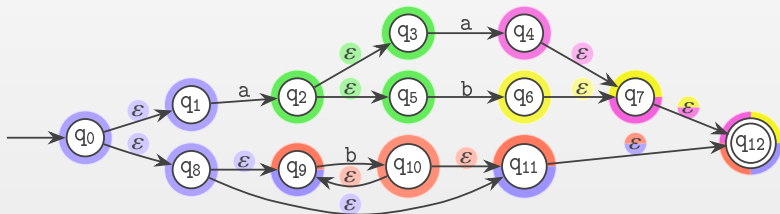


Closures

Given $\omega \in \Sigma^*$, a ω -closure of a state q in NFA \mathcal{A} is a set of states reachable from q by the action ω .

We say that ω is in the language of the NFA \mathcal{A} ($\omega \in \mathcal{L}(\mathcal{A})$)
 $\Leftrightarrow \omega$ -closure of the starting state of \mathcal{A} contains a final state.

Special case: ε -closures: sets of states reachable via doing nothing.

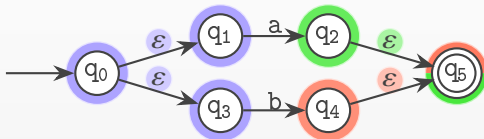


Given such closures, they can be considered as new «states».

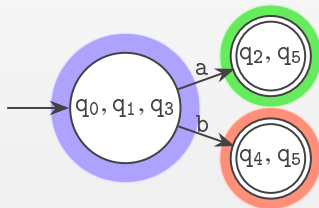


Simple Example of ε -Removal

An NFA \mathcal{A} with the ε -closures of its states being highlighted:



The closures are then merged into single states, and given a transition from $q_i \xrightarrow{\gamma} q_j$, where q_i belongs to closure $M(q_i)$, and q_j to $M(q_j)$, transition $M(q_i) \xrightarrow{\gamma} M(q_j)$ is added.



A closure is marked as a final \Leftrightarrow it contains at least one final state.



ε -Closures and Chain Rules

- Any transition $q_i \xrightarrow{\varepsilon} q_j$ corresponds to **a chain rule** $[q_i] \rightarrow [q_j]$ in the corresponding grammar G .
- state ε -closure is a closure set of the corresponding non-terminal N :
$$C(N) = \left\{ N_i \mid \exists N'_1, \dots, N'_k (N \rightarrow N'_1 \ \& \ \dots \ \& \ N'_k \rightarrow N_i) \right\}$$

I.e. $\langle N, N_i \rangle$ are pairs in **a transitive closure** \rightarrow_c^+ of the relation \rightarrow_c :
 $A_i \rightarrow_c A_j \Leftrightarrow (A_i \rightarrow A_j \in G)$.
- Before removing all chain rules, for every $N' \in C(N)$ and a non-chain rule $N' \rightarrow \Phi$, we add the transition $N \rightarrow \Phi$ to the set of grammar rules. Exactly as in the ε -closure algorithm for NFA.

Initial grammar:

$S \rightarrow Q_1$ $S \rightarrow Q_3$ $Q_1 \rightarrow aQ_2$
 $Q_3 \rightarrow bQ_4$ $Q_2 \rightarrow Q_5$ $Q_4 \rightarrow Q_5$
 $Q_5 \rightarrow \varepsilon$

After removing chain rules:

$S \rightarrow aQ_2$ $S \rightarrow bQ_4$
 $Q_2 \rightarrow \varepsilon$ $Q_4 \rightarrow \varepsilon$

Note: unreachable non-terminals Q_1, Q_3, Q_5 are deleted from the resulting grammar.



ω -Closures and Subset Construction

The closure sets wrt transitions by non- ε actions can be also merged in similar sense.

Subset Automaton Construction

Let an ε -free NFA \mathcal{A} be given. Its **subset automaton** $D(\mathcal{A})$ can be constructed as follows.

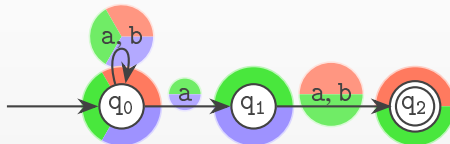
- q_0 becomes the starting state $\{q_0\}$ of $D(\mathcal{A})$.
- Given a state M in $D(\mathcal{A})$ and $\gamma \in \Sigma$, construct a closure set $M_\gamma = \{q_i \mid \exists q_j \in M (q_j \xrightarrow{\gamma} q_i)\}$. If M_γ is non-empty and does not yet introduced as a state of $D(\mathcal{A})$, add it to set of states of $D(\mathcal{A})$.
- The final states of $D(\mathcal{A})$ are labelled with the sets containing at least one final state of \mathcal{A} .

In fact, the states of $D(\mathcal{A})$ are ω -closures of \mathcal{A} -states, where $\omega \in \Sigma^*$.



Subset Automaton: a Simple Example

Let us consider the following NFA with γ -closures of its states:

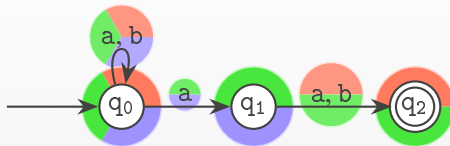


- a-closure of starting state $\{q_0\}$ is $\{q_0, q_1\}$.
- b-closure of starting state $\{q_0\}$ is the state $\{q_0\}$ itself.
- a-closure of the state $\{q_0, q_1\}$ is $\{q_0, q_1, q_2\}$.
- b-closure of the state $\{q_0, q_1\}$ is $\{q_0, q_2\}$.
- a-closure of the state $\{q_0, q_1, q_2\}$ is $\{q_0, q_1, q_2\}$ itself, while b-closure is the state $\{q_0, q_2\}$.
- a-closure of the state $\{q_0, q_2\}$ is $\{q_0, q_1\}$, while b-closure is $\{q_0\}$.

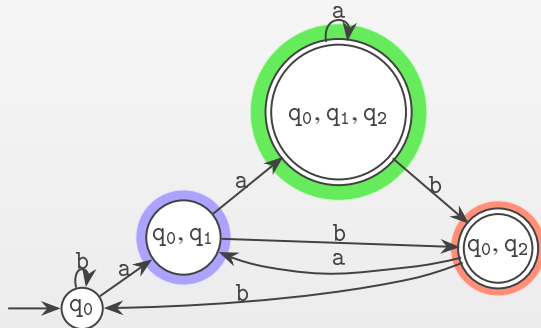


Subset Automaton: a Simple Example

Let us consider the following NFA with γ -closures of its states:



Hence, its subset automaton is:



Deterministic Finite Automata

👁👁 Definition

A deterministic finite automaton (DFA) is a tuple

$\mathcal{A} = \langle Q, \Sigma, q_0, F, \delta \rangle$, where:

- Q is a state set, Σ is a terminal alphabet;
- δ is a transition set $\langle q_i, \gamma, q_j \rangle$, where $q_i, q_j \in Q$, $\gamma \in \Sigma$, and for any q_i , γ there is at most one q_j such that $q_i \xrightarrow{\gamma} q_j \in \delta$;
- $q_0 \in Q$ is a starting state, $F \subseteq Q$ is a set of final states.

Language $\mathcal{L}(\mathcal{A})$ of DFA \mathcal{A} is a set $\{\omega \mid \exists q \in F (q_0 \xrightarrow{\omega} q)\}$, i.e. there exists a final state that is ω -closure of q_0 .

By construction, the subset automaton has no non-determinism in the transition set:

- ε -transitions are eliminated in the preliminary ε -free NFA;
- the non-singleton transition sets are processed in the subset construction.



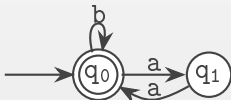
Traps and Trims

- A **trap state** is a state s.t. any its ω -closure is non-final.

Trim DFA

- For any q_i, γ there is **at most one** q_j s.t. $q_i \xrightarrow{\gamma} q_j \in \delta$;
- is naturally constructed via subset technique;
- default in RoFL course, useful for most operations.

Trim DFA example

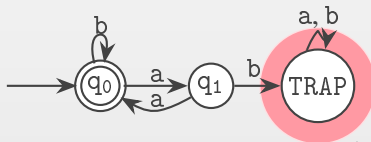


Complete DFA

- For any q_i, γ there is **exactly one** q_j s.t. $q_i \xrightarrow{\gamma} q_j \in \delta$;
- usually requires introducing **trap (sink) states**;
- useful for constructing complementation.

DFA with the trap state

for $\Sigma = \{a, b\}$



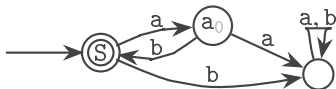
Complementation and Traps

- By switching finality of all states in DFA \mathcal{A} , we can construct a DFA \mathcal{A}' accepting exactly the set of words that are rejected by the initial DFA, i.e. $\mathcal{L}(\mathcal{A}') = \Sigma^* \setminus \mathcal{L}(\mathcal{A})$.
- The language complementation requires complete DFA.

Initial trim DFA for $(ab)^*$



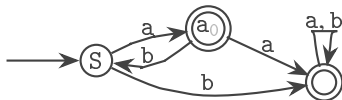
Complete DFA for $(ab)^*$



Invalid complement DFA
recognising $a(ba)^*$



Valid complement DFA



- Without a trap state, complementation operation loses words starting with b , or containing either aa or bb .



One More Encoding: Equations

Sometimes it is convenient to gather all the right-hand sides of the rules with a same left-hand side together. Then, if we replace \rightarrow by $=$ sign, we get an equation system determining non-terminal languages:

$$\begin{array}{llll} S \rightarrow a[q_1] & [q_2] \rightarrow a[q_3] & & S = a[q_1] \mid a[q_2] \\ S \rightarrow a[q_2] & [q_2] \rightarrow b[q_4] & \rightarrow & [q_1] = a[q_1] \mid a[q_2] \\ [q_1] \rightarrow a[q_1] & [q_3] \rightarrow \varepsilon & & [q_2] = a[q_3] \mid b[q_4] \\ [q_1] \rightarrow a[q_2] & [q_4] \rightarrow \varepsilon & & [q_3] = \varepsilon \\ & & & [q_4] = \varepsilon \end{array}$$

If there is no rule part $[q_1] = a[q_1]$, these languages could be found by exhaustive substitutions of the right-hand sides.

E.g. $\mathcal{L}([q_3]) = \mathcal{L}([q_4]) = \{\varepsilon\}$, while
 $\mathcal{L}([q_2]) = \{a\mathcal{L}([q_3])\} \cup \{b\mathcal{L}([q_4])\} = \{a, b\}$.

How to deal with self-referring rules as $[q_1] = a[q_1]$?



Arden's Lemma

👁👁 Theorem

If a language \mathcal{L} satisfies the equation $\mathcal{L} = \mathcal{L}_1\mathcal{L} \cup \mathcal{L}_2$, where $\varepsilon \notin \mathcal{L}_1$, then $\mathcal{L} = \mathcal{L}_1^*\mathcal{L}_2$.

Proof: Let us consider arbitrary $\omega \in \mathcal{L}$.

- If $\omega \in \mathcal{L}_2$, then the statement trivially holds.
- Otherwise, $\exists \omega_1 \in \mathcal{L}_1, \omega' \in \mathcal{L} (\omega = \omega_1\omega')$. The suffix ω' also belongs to $\mathcal{L}_1\mathcal{L} \cup \mathcal{L}_2$, and $|\omega'| < |\omega|$, since $\omega_1 \neq \varepsilon$. Now we can repeat the same reasoning for ω' , and due to finiteness of $|\omega|$ and well-foundedness of $(\mathbb{N}, <)$ we will eventually get $\omega' \in \mathcal{L}_2$. \square

Arden's lemma allows one to solve the equation systems in Gaussian style, via non-terminal elimination + substitution, assuming there are no chain rules in the grammar.



Equation Solving Example

Let us construct the language of the grammar:

$$S \rightarrow aT \quad S \rightarrow aS$$

$$T \rightarrow aT \quad T \rightarrow bT \quad T \rightarrow bF \quad F \rightarrow \varepsilon$$

First, construct the system and substitute F :

$$S = (aS) \mid (aT)$$

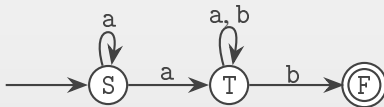
$$T = ((a \mid b)T) \mid b(\varepsilon)$$

Solve the second equation: $T = (a \mid b)^*b$

Then substitute the solution: $S = (aS) \mid (a(a \mid b)^*b)$.

The resulting language is: $S = a^*a(a \mid b)^*b$

The NFA that corresponds to the grammar is given below:



Equation Solving Example

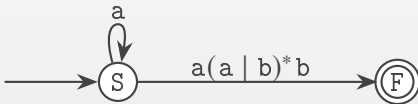
Let us construct the language of the grammar:

$$\begin{aligned} S &\rightarrow aT & S &\rightarrow aS \\ T &\rightarrow aT & T &\rightarrow bT & T &\rightarrow bF & F &\rightarrow \varepsilon \end{aligned}$$

The resulting language is: $S = a^*a(a \mid b)^*b$

The NFA that corresponds to the grammar is given below .

After solving T -based equation and substituting F value, in fact we again constructed an NFA, whose transitions are marked with regexes.



If we assume that S is preceded by the “very starting state” S' , then $\mathcal{L}(S)$ can be also considered as a transition in the NFA containing only S' and F states.



Finding NFA Language

The extended NFAs allow one to use transitions marked with regexes.

State Exclusion Method

- For the sake of uniformity, we introduce “the very starting state” S , having ε -transition to q_0 , and “the very final state” T , having ingoing ε -transitions from $q \in F$. All the states except S and T are now ordinary.
- In order to exclude the state q s.t. $q \xrightarrow{\tau} q$, for all pairs q_A, q_B , where $q_A \xrightarrow{\Phi} q$, $q \xrightarrow{\Psi} q_B$, add the transition $q_A \xrightarrow{\Phi(\tau)^*\Psi} q_B$, then we can delete q .
- When only S and T are left, where $S \xrightarrow{\rho} T$, the expression ρ is the regex equivalent to the NFA.

