Graph Theory - Lecture 3 Walks, Trials, Paths

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1 Walks

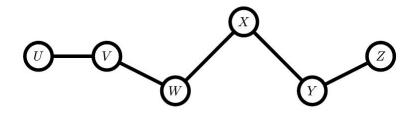


Figure 1: Simple Walk : UVWXYZ

Let us consider a graph G(V, E), where $E(G) = \{e_i \in E(G) | i = 1, 2, ..., n = |E(G)|\}$ and $V(G) = \{v_i \in V(G) | i = 1, 2, ..., k = |V(G)|\}$

pefinition 1.1 (Walk). An ordered sequence of edges such as

$$W(G) = (e_1, e_2, \dots, e_l); e_j \in E(G)$$
 (1)

where $e_i \in E(G)$ is a walk if there exists a corresponding sequence of vertices $(v_0, v_1, \dots v_l)$ such that $e_j = (v_{j-1}, v_j)$; the vertices in the walk may not be <u>distinct</u>.

Note: A walk for which $v_0 = v_l$ is a <u>closed walk</u>.

Definition 1.2. The **length** of a <u>walk</u> is the number of edges in the sequence.

E.g. in [Eq 1], the length of the walk is l.

Definition 1.3 (Distance). The **distance** d(u, v) between two vertices u and v in a graph $\mathbf{G}(\mathbf{V}, \mathbf{E})$ is the length of the shortest uv path 1 in G and for any three vertices u, v, w,

$$d(u,w) + d(w,v) \ge d(u,v). \tag{2}$$

Equation 2 is referred to as the **Tiangular Inequality**

Definition 1.4 (Eccentricity). The maximum distance from a vertex $u \in V(G)$ to any other vertex, v in G(V, E) is called the eccentricity, e(u).

If w is a vertex such that e(v) = d(v, w), then it is said that the eccentricity of the vertex v is realized by the vertex w.

Definition 1.5 (Radius). The smallest eccentricity, i.e. $\underline{Rad(G) = \min_{v \in V(G)} e(v)}$ is the radius of the $\mathbf{G}(\mathbf{V}, \mathbf{E})$.

Definition 1.6 (Diameter). The **diameter** D(G)/Diam(G) of G(V, E) is the minimum distance, d(u, v) over all pairs of $\{u, v\}$ pairs of vertices of V(G), i.e.

Moximum

OR.

The maximum eccentricity, i.e. $Diam(G) = \max_{v \in V(G)} e(v)$ is the diameter of the G(V, E)

Definition 1.7 (Central Vertex). The **Central vertex** is the vertex, for which e(v) = Rad(G).

Pheorem 1.1. For every non-trivial connected graph $\mathbf{G}(\mathbf{V},\mathbf{E})$,

$$Rad(G) \le Diam(G) \le 2 \times Rad(G)$$

Proof. The inequality $Rad(G) \leq Diam(G)$ is immediate since the smallest eccentricity cannot exceed the largest eccentricity.

Let u and v be two vertices such that d(u, v) = Diam(G) and let w be a **Central vertex** of G(V, E).

Therefore, e(w) = Rad(G). Hence, the distance between w and any other vertex of $\mathbf{G}(\mathbf{V}, \mathbf{E})$ is at most Rad(G).

By triangle inequality,

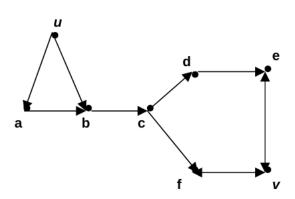


Figure 2: Diameter
$$= \infty$$
; Radius $= 3$

Froposition 1.1. If there is a walk from vertex u to v, then there is also a path from u to v of length (|V(G)|-1).

Proof. If all the vertices are distinct, then there is nothing to prove.

Otherwise, suppose some vertex v_x is repeated in the path. We can remove the segment of the walk between these two v_x 's and obtain another shorter walk from u to v. If this is not a path, we may again repeat the above procedure on this walk.

Clearly, after a finite number of steps, one can obtain a path.

2



Figure 3: Shorter Path

Z Trial

Definition 2.1 (Trial). A trial is a <u>walk</u> in which all the edges e_j are distinct and a closed trial is a closed walk, that is also a trial.

A closed trial is also referred to as a Circuit.

3 Path

Definition 3.1 (Path). A path is a <u>trial</u> in which all the vertices in the sequence, Eq 1 are distinct

A path is referred to as a P_n , where $n = |P_n|$ is the number of vertices of the path OR order of the path.

Note: A closed path is a closed trial in which all the vertices are distinct except $v_0 = v_l$.

4 Cycle

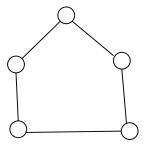


Figure 4: Cycle C_5

Definition 4.1 (Cycle). A cycle $(C_n, where n = |C_n| is the order of <math>C_n$) is a closed path that includes three or more edges.

Note:

Equivalently, a cycle is a subgraph isomorphic to one of the cycle graphs C_n , e.g. Figure 4 In set theoretic notations

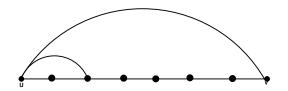
 $Walks \supset Trials \supset Paths$

Theorem 4.1. If every vertex of a graph G(V, E) has degree at least 2, then G(V, E) contains a cycle

Proof. Let P(G) be a maximal path in G(V, E), and let $u \in V(G)$ be an endpoint of P(G). Since P(G) cannot be extended, every neighbour of u must already be a vertex of P(G). Since, u has degree at least 2, it has a neighbour $v \in V(P)$ via an edge not in P(G)

The edge uv completes a cycle with the portion of P(G) from v to u

Note: If |V(G)| = Z and $E(G) = \{ij : |i - j| = 1\}$, then every vertex has degree 2, but G(V, E) has no cycle, i.e. no non-extendible path.



Eemma 4.2. If every vertex of a (finite) graph $\mathbf{D}(\mathbf{V}, \mathbf{E})$ has out-degree/ $d^+(v)$ (or in-degree/ $d^-(v)$) at least 1, then $\mathbf{D}(\mathbf{V}, \mathbf{E})$ contains a cycle.

Proof. Let P(D) be a maximal path in $\mathbf{D}(\mathbf{V}, \mathbf{E})$, and u be the last vertex on P(D). Since P(D) can not be extended, every successor of u is in V(D). There is at least one successor of u, say v. This edge uv and the path from v to u form a cycle.

Abstract

Walk: Vertices may REPEAT, edges may REPEAT

Trail: Vertices may REPEAT, edges CANNOT REPEAT

Circuit: Vertices may REPEAT (Closed Trial)

Path: Vertices CANNOT REPEAT, thus edges also CANNOT REPEAT.

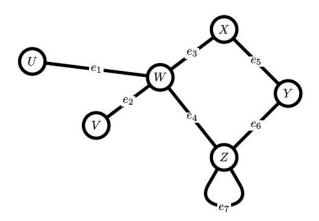


Figure 5: Graph with examples

With respect to Figure 5

- 1. Direct walk between U and Y graph is: UWXY
- 2. A more **roundabout walk** between U and Y could be: UWVWZZY
- 3. Alternative Notation : $Ue_1We_2Ve_2We_4Ze_7Ze_6Y$
- 4. The walk $Ue_1We_4Ze_7Ze_6Y$ is an open walk
- 5. The walk $Ue_1We_3Xe_5Ye_6Ze_4We_1U$ is closed walk
- 6. The walk $Ue_1We_3Xe_5Ye_6Z$ is a path
- 7. The walk $We_4Ze_7Ze_6Ye_5Xe_3W$ is a **closed trail**
- 8. The walk $We_3Xe_5Ye_6Ze_4W$ is a cycle
- 9. The walk $Ue_1We_2Ve_2We_4Ze_7Ze_6Ye_7$ is not a trail or a path
- 10. The walk $Ue_1We_4Ze_7Ze_6Y$ is a trail but not path