The Roots of the Automata Theory

- The present day Automata theory has its roots in the work of the logicians and mathematicians of 1930s.
 - They studied to have an answer to the question:
 - Is an algorithm possible to determine truth or falsity of a mathematical proposition?
- David Hilbert looked for an algorithm to determine if an arbitrary formula in FOPL, applied to integers, was true.
- In 1931, Kurt Godel through his Incompleteness theorem proved that no such algorithm could exist.
 - He constructed a formula in Predicate calculus, applied to integers, whose very
 definition stated that it could neither be proved nor disproved within this system.
- The quest for an answer to the question if an algorithm to any problem within a logical or mathematical framework is possible led to the development of various abstract or mathematical machines in 1930's.
 - Lambda Calculus by Alonzo Church
 - Recursive function by Kurt Godel
 - Formal systems by Stephen Kleene
 - Post's machine by Emil Post
 - Turing Machine by Alan Turing
 - Anything computable with one model is also computable with any other model.
- Turing Machine is accepted as a theoretical model of a computer as per Church's hypothesis.
 - TM is accepted as the most powerful model of computation.
- In 1940s and 1950s, some simpler machines, called Finite Automata (FA), were developed.
 - Automata, originally proposed to model the brain functions, were later found useful for variety of other purposes:
 - Lexical Analysis
 - Network Protocol Verification etc.
- In 1950s, the linguist Noam Chomsky introduced formal grammars, shown to be equivalent to automata and TMs.
 - Formal grammars have application in Compilers, and some other important software.

The Twing Machine is a simple mothermatical model of a computer.

Despite its simplicity, the Twing Machine models the computing (whose computing capability of a general purpose) computer.

The Trowing Machinel is studied both in

the class of languages it defines (called the recursively enumerable sets)

the class of integer functions it computes (called the partial recursive functions)

A parciety of other models of computation over also there, which are shown to be equivalent to the Twing Machine in computing power.

Clearly one cannot prove that the Twing machine model its equivalent to our intuitive notion of a computer, but there are compelling arguments for this equivalence, which has become known as Chwich's hypothesis. The assumption that the intuitive notion of computable function" can be identified with the class of partial recursion functions is known as Church's hypothesis or V Each Twing machine can be thought of as the Church Twing Thesis. IMO A for computed by a Tolly computing a function to integers to integers. 198 M halls on with 0) on its tape, then we say fx(i,i,i,-,i)=) at halt withing of old who there consisting of old with worther a (not halt with & every K

V Formal languages & Antomata Theory (FLAT) It deals with the study of abstract/mathematical machines, sometimes, automata, as well as the Computational problems Which can be solved on them. V Antomata - A Greek word meaning "self-acting". finite Control V Abstract Machines · various parts of such m/cs are defined logically just by their functions and not implemented on hardware Such machines may be allowed to run as long as you want (without any chance of getting heated Fig. A Finite Antomaton and may receive input as such machines compute in a stepsise fasion. That is, long as need an input symbol, perform entain punction, read entain ext input symbol and Such m/cs sometimes may have to stack of infinite memory.

I Antomata-theory is Closely related to Formal language theory.

VAn antomaton gives a fin representation of a form langnage which may be a finite or an infinite & of strungs formed with the Symbols taken from a finit Set could Alphabet), following Sp

Examples of Formal languages $\Sigma = \text{alphabet} = \{0,1\}$ $L_1 = \text{Strings with one or more leading of followed by a single 1.}$ $L_2 = \text{one or more leading of bollow by an equal number of 1's.}$

I The Roots of the Antomata Theory

V Work of Logicians in 1930's

I The Question

V David Hilbert

V Godel's Incompleteness Theorem

J Development of Abstract Machines

Trowing Machine - Atheoretical model of the Computer

V Finite Antomata: Simplen m/es V Chomprey's Grammar

· He constructed a formula in the predicate calculus, applied to intege whose very definition stated that. could neither be proved not disp within this logical system.

I The great for Lanswer to the greation framework is possible led to 8 the deve of various abstract or mathematical mo m 1930/s.

V Lambda Calulus by Alongo Churchy Development Recorsive Functions by Knort Godely Formal Systems by Stephene Klung Post's machine by Emil Post, Twing Machines by Alan Tring.

I All these models of computation are exactly equivalent, Amythin computable with one model is all computable with any other model. TM is accepted as a theoretical model of a computer as per http://www.chis hypothesis.

I That is, TM is accepted as the most powerful model of Computation

~ In 1940's and 1950s, some simpler machines, called Finite Automata (FA) were developed.

If a step by step solution to any A. Simplerthe brain functions, welater found within a logical or mathemat NCS such as Sexical Analysis, Network

Protocol Verification etc. V In 1950's, (Noam Champing introduced formal comsky's grammars, shown to be equivalent to

amounts automata and TMs v Formal grammars have application in compilers, and some other important Software.

formal = Antowala = TMb

why to Study Automata Theory

- Antomata are essential for the study of limits of computation.
- . There are two important issues
 - 1. What can a computer do at all? This study is called "decidability", and the problems that can be solved by computer are called "decidable".
 - 2. What can a computer do efficiently.

 This study is called "

 "intreactability," and the problems
 that can be solved by a compute
 using no more time than some
 slowly growing function of the
 size of the input are called

 "tractable". Often we take all

 polynomial functions to be slowly growing while functions that semests
 orow baster than holynomial week

to grow too bast.