

ELEC 5360 Homework 4

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1.

The log-likelihood function is given by

$$\begin{aligned}\Lambda_L(\tau) &= \Re\left[\frac{1}{N_0} \int_{T_0} r(t) s_l^*(t, \tau) dt\right] \\ &= \Re\left[\frac{1}{N_0} \int_{T_0} r(t) \sum_n I_n g^*(t - nT - \tau) dt\right] \\ &= \Re\left[\sum_n I_n y_n(\tau)\right]\end{aligned}$$

where $y_n = \frac{1}{N_0} \int_{T_0} r(t) g^*(t - nT - \tau) dt$

In order to maximize the likelihood function, we set

$$\frac{d\Lambda_L(\tau)}{d\tau} = \Re\left[\sum_n I_n \frac{d}{d\tau} y_n(\tau)\right] = 0$$

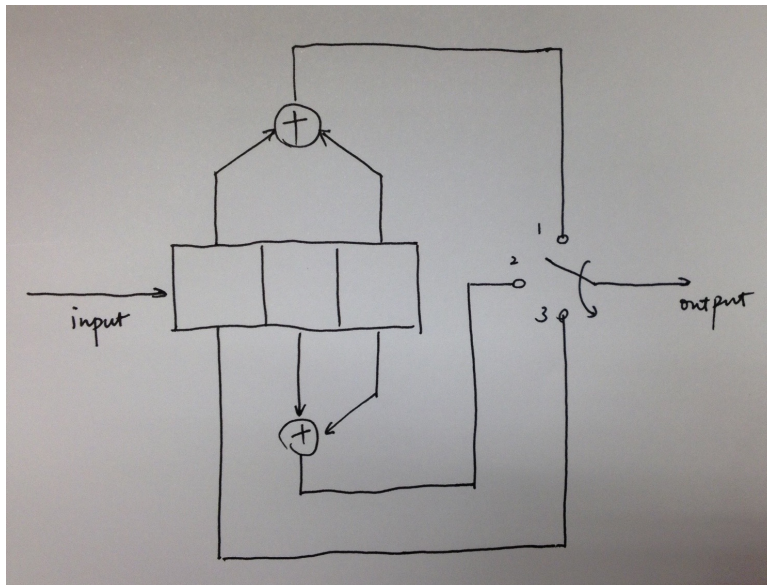
Further simplify the expression,

$$\Re\left[\sum_n I_n \frac{d}{d\tau} y_n(\tau)\right] = \sum_n \Re[I_n] \Re\left[\frac{d}{d\tau} y_n(\tau)\right] - \sum_n \Im[I_n] \Im\left[\frac{d}{d\tau} y_n(\tau)\right]$$

Therefore we obtain,

$$\sum_n \Re[I_n] \Re\left[\frac{d}{d\tau} y_n(\tau)\right] = \sum_n \Im[I_n] \Im\left[\frac{d}{d\tau} y_n(\tau)\right]$$

2. (a)

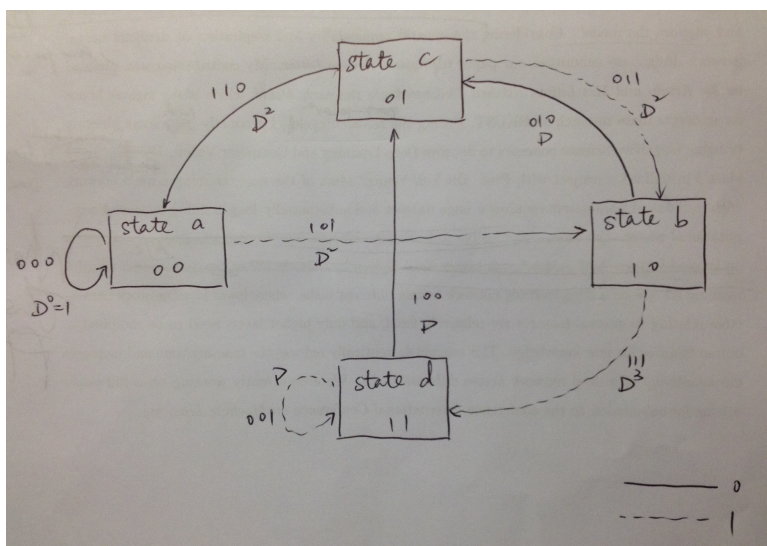


(b)

The code rate is $1/3$.

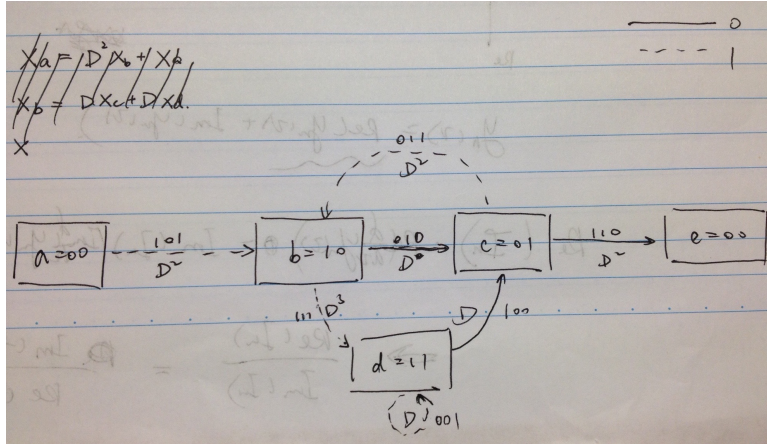
The constrain length is 3.

(c)



(d)

Split the state a to *start* state a and *end* state e as shown below.



Then we obtain

$$X_b = D^2 X_a + D^2 X_c$$

$$X_c = D X_b + D X_d$$

$$X_d = D^3 X_b + D X_d$$

$$X_e = D^2 X_c$$

Simplify the set of equations, we will get

$$X_e = D^2 (D X_b + D X_d)$$

$$X_d = D^3 / (1 - D) X_b$$

$$D^2 X_a = X_b - D^2 (D X_b + D X_d)$$

Then

$$\begin{aligned}\frac{X_e}{X_a} &= \frac{(D^3 + D^6/(1-D))D^2}{1 - D^3 + D^6/(1-D)} \\ &= \frac{D^5(1-D) + D^8}{(1-D^3)(1-D) + D^6} \\ &= \frac{D^5 - D^6 + D^8}{1 - D - D^3 + D^4 + D^6}\end{aligned}$$

(e)

From last equation, we obtain

$$\frac{X_e}{X_a} = (D^5 - D^6 + D^8)[1 + (D + D^3 - D^4 - D^6) + (D + D^3 - D^4 - D^6)^2 + \dots]$$

Therefore the minimum free distance d_{free} of the code is 5.

3.

Notice that $f_k - f_j = n/T$ for $n = 1, 2, \dots, N-1$ (as $f_k \neq f_j$), also under the assumption that $f_k + f_j \gg 1$, Then

$$\begin{aligned}& \int_0^T \cos(2\pi f_k t + \phi_k) \cos(2\pi f_j t + \phi_j) dt \\ &= \frac{1}{2} \int_0^T \cos(2\pi f_k t + \phi_k + 2\pi f_j t + \phi_j) dt + \frac{1}{2} \int_0^T \cos(2\pi f_k t + \phi_k - 2\pi f_j t - \phi_j) dt \\ &= \frac{1}{2} \int_0^T \cos(2\pi(f_k + f_j)t + \phi_k + \phi_j) dt + \frac{1}{2} \int_0^T \cos(2\pi n t/T + \phi_k - \phi_j) dt \\ &= \frac{1}{2\pi \frac{n}{T}} \frac{1}{2} \sin(2\pi \frac{nt}{T} + \phi_k - \phi_j) \Big|_0^T + \frac{1}{2\pi(f_k + f_j)} \frac{1}{2} \sin(2\pi(f_k + f_j)t + \phi_k + \phi_j) \Big|_0^T\end{aligned}$$

Notice that

$$\sin(2\pi \frac{nt}{T} + \phi_k - \phi_j) \Big|_0^T = \sin(2\pi n + \phi_k - \phi_j) - \sin(\phi_k - \phi_j) = 0$$

and

$$\frac{1}{2\pi(f_k + f_j)} \simeq 0 \text{ for large } f_k + f_j$$

Hence we get $\int_0^T \cos(2\pi f_k t + \phi_k) \cos(2\pi f_j t + \phi_j) dt = 0$ and therefore subcarriers of the corresponding sub-channels are mutually orthogonal.

4. (a)

If we don't account the cyclic prefix, the maximum symbol rate for each subcarrier relates to the Δf . So the ideal maximum symbol rate will be 15k symbol/sec. However, in this specific case, taking the CP in to consideration, using the least amount of CP, we get the maximum symbol rate $(\text{size}(FFT) - \text{size}(CP)) / (\text{size}(FFT) \times T_{\text{sample}}) = (2048 - 144) / (2048 / (\Delta f \times 2048)) = 13.95\text{k symbol/sec}$

(b)

The symbol duration is given by $T = NT_{tb}$, where $T_{tb} = 1/f_{tb}$ and $N = 20 \text{ MHz} / 15 \text{ kHz}$. Plugging the numbers in, we get $20 \text{ MHz} / 15 \text{ kHz} / 20 \text{ MHz} = 1/15 \approx 0.0667 \text{ sec}$.

(c)

One view of the need for CP in OFDM is, in the guarding interval preventing the inter-OFDM system-interference, we may as well arising something to further prevent information loss. More formally, in the channel matrix point of view, cyclic prefix will lead to a circular matrix which can further leads to diagonal matrix decomposition, this help simplify the matrix a lot.

(d)

The power loss can be defined as the ratio between redundant information propagated and symbol transmitted.

Therefore, for normal CP, we get 144/2048, for extended CP, 512/2048

(e)

We need to require the guarding interval spread wider than the delay, namely $CP / (\text{size}(FFT) \Delta f) \geq 7\mu s$, where $\Delta f = 15 \text{ kHz}$. Then we get $CP \geq 215.04$. Therefore the extended 512 CP is needed.