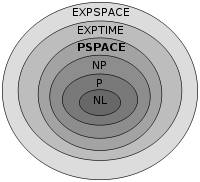
P is equal to NP

No Author Given

No Institute Given

**Abstract.** In this paper, we prove that the complexity classes P and NP are equal by showing that the 3-SAT problem can be solved in deterministic polynomial time.

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**1. Introduction**

The P versus NP problem is one of the most important problems in computer science. In 1971, Stephan Cook proved that the 3-SAT is NP-complete [1], which means that if 3-SAT can solved in polynomial time, we can state that P=NP.

**2. Preliminaries**

See Wikipedia: http://en.wikipedia.org/wiki/Boolean\_satisfiability\_problem

**3. Main theorem**

**Theorem 1**

*3-SAT can be solved in .*

**Proof.** Let us denote given 3-SAT as follows:

,

where . Since a clause is equals to , we can reform as follows:

.

By replacing clauses with variables (in case some clauses are equal, replace them with the same variables), we can obtain 2-SAT . The number of clauses of is , and the number of variables in is at most . It is known that 2-SAT can be solved in linear time [2]. ■

For instance, for 3-SAT , the corresponding 2-SAT is , where , and are corresponding to , and , respectively.

**4. Conclusion**

Give me $1,000,000!!! щ(ﾟДﾟщ)

**References**

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2. Krom, L. R.: The Decision Problem for a Class of First-Order Formulas in Which all Disjnctions are Binary. eitschrift für Mathematische Logik und Grundlagen der Mathematik 13 (1967) 15-20