

# MA3227 Numerical Analysis II

Topics for Midterm Exam

Simon Etter



2019/2020

# Topics for Midterm Exam

## Finite differences

- Consistency and stability implies convergence theorem:

$$\|u - u_n\| \leq \|\Delta_n^{-1}\| \|\Delta_n u + f\|.$$

- Taylor expansion for computing the consistency error  $\|\Delta_n u + f\|$ .

## Sparse LU

- Graph of sparse matrices, structural nonzeros, fill-in.
- Path theorems for matrix powers, inverses and the LU factorisation.  
Application only, no proofs.
- Runtimes of dense and sparse LU factorisation:

$$\text{dense: } \mathcal{O}(N^3), \quad \text{sparse: } \begin{cases} \mathcal{O}(N) & \text{if } d = 1, \\ \mathcal{O}(N^{3/2}) & \text{if } d = 2, \\ \mathcal{O}(N^2) & \text{if } d = 3. \end{cases}$$

# Topics for Midterm Exam

## Orthogonal matrices

- ▶ Basic orthogonalisation procedure  $\hat{b} = b - \frac{a^T b}{a^T a} a$ .
- ▶ Use of the basic orthogonalisation procedure in the Gram-Schmidt, Arnoldi and Lanczos algorithms.
- ▶ Arnoldi relation  $AQ_k = Q_{k+1}H_k$ .
- ▶ Lanczos theorem:  $A$  symmetric  $\implies H_k$  tridiagonal.
- ▶ QR factorisation  $A = QR$ .
- ▶  $x = \arg \min \|Ax - b\|_2 \iff x = R^{-1} Q^T b$

# Topics for Midterm Exam

## Krylov methods

- ▶ Krylov algorithms:

$$x_k = p_{k-1}(A) b \quad \text{where} \quad p_{k-1} = \arg \min_{p_{k-1} \in \mathcal{P}_{k-1}} \| (Ap_{k-1}(A) - I) b \|.$$

- ▶ GMRES and MinRes minimise 2-norm, CG minimises  $A^{-1}$  norm.
- ▶ Runtime GMRES is  $N$  matvecs plus  $\mathcal{O}(Nk^2)$  other operations.  
Runtime MinRes and CG is  $N$  matvecs plus  $\mathcal{O}(Nk)$  other operations.
- ▶ GMRES is applicable to any matrix. MinRes requires symmetry.  
CG requires symmetric positive definite.

# Topics for Midterm Exam

## Krylov methods (continued)

- Convergence estimate

$$\|Ax_k - b\| \leq C \min_{q_k \in \mathcal{P}} \max_{\lambda_\ell} \frac{|q_k(\lambda_\ell)|}{|q_k(0)|}$$

and its corollaries regarding scale-invariance, outliers and finite termination.

Application only, no proofs.

You may ignore `alternating_eigenvalues()` from lab 2.

- Convergence bound for intervals: if  $\lambda_\ell \in [1, \kappa]$ , then  $\|Ax_k - b\| \leq C \rho^k$  and  $\rho \in (0, 1)$  is closer to 0 for  $\kappa$  closer to 1.

# Topics for Midterm Exam

## Jacobi and multigrid

- Jacobi and Gauss-Seidel iterations:

$$\text{Jacobi: } x_{k+1} = D^{-1} (b - (A - D) x_k),$$

$$\text{Gauss-Seidel: } x_{k+1} = (D + U)^{-1} (b - Lx_k).$$

You should know these formulae and be able to perform the iterations by hand.

- Error recursion formula

$$x_{k+1} = x_k + B(b - Ax_k) \implies x_{k+1} - x = (I - BA)(x_k - x)$$

and its consequences regarding convergence.

# Topics for Midterm Exam

## **Root-finding algorithms**

- ▶ You should be able to perform bisection, false position, Newton and secant methods by hand.
- ▶ You should be able to show the convergence result for Newton's method, but you may ignore secant and false position.
- ▶ Be informed about Broyden and gradient descent. No need to remember the formulae for Broyden's method.

# Topics for Midterm Exam

## Example question

Consider the  $2 \times 2$  matrix

$$A = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \quad \text{where } a, d \neq 0.$$

Show that Gauss-Seidel converges at least twice as fast as Jacobi.

For an actual exam question, I would clarify more precisely what is meant by twice as fast. Here, it will be clear once you see the answer on the next slide.

*Hint.* We have

$$x_k - x = R^k (x_0 - x) \quad \text{where} \quad R = \begin{cases} -D^{-1}(A - D) & \text{(Jacobi),} \\ -(U + D)^{-1}L & \text{(Gauss-Seidel).} \end{cases}$$



# Topics for Midterm Exam

*Answer.* We know that  $\|x_k - x\| \leq C |\lambda_{\max}|^k$  where  $\lambda_{\max}$  is the eigenvalue of largest magnitude of the two matrices  $R$  given above.

For Jacobi, we compute

$$-D^{-1}(A - D) = -\begin{pmatrix} 0 & \frac{b}{a} \\ \frac{c}{d} & 0 \end{pmatrix}$$

and thus  $|\lambda_{\max}| = \sqrt{\frac{bc}{ad}}$ .

For Gauss-Seidel, we compute

$$-(U + D)^{-1}L = -\begin{pmatrix} 0 & \frac{b}{a} \\ 0 & \frac{bc}{ad} \end{pmatrix}$$

and thus  $|\lambda_{\max}| = \frac{bc}{ad}$ .

We conclude  $|\lambda_{\max}^{(Jac)}|^2 = |\lambda_{\max}^{(GS)}|$ , i.e. Jacobi requires two steps to achieve the same error reduction as Gauss-Seidel achieves in a single step.

# Topics for Midterm Exam

## Example question

Let  $A \in \mathbb{R}^{N \times N}$  be a dense and symmetric matrix with eigenvalues in the interval  $[1, \kappa]$ . Determine  $\alpha \in \mathbb{R}$  such that if  $\kappa = \mathcal{O}(N^{\tilde{\alpha}})$  with  $\tilde{\alpha} < \alpha$ , then MinRes is asymptotically faster than LU factorisation, and if  $\tilde{\alpha} > \alpha$  then MinRes is asymptotically slower than LU factorisation.

*Hint.* Under the assumptions above, the MinRes iterates  $x_k$  satisfy

$$\|Ax_k - b\| \leq C \rho^k \text{ with } \rho = 1 - \mathcal{O}(\kappa^{-1/2}),$$

and

$$k = \mathcal{O}(\log(\varepsilon) (1 - \rho)^{-1}) = \mathcal{O}(\log(\varepsilon) \kappa^{1/2})$$

steps are required to achieve  $\|Ax_k - b\| \leq \varepsilon$ .

# Topics for Midterm Exam

*Answer.* Since  $A$  is dense, the runtime of LU is  $\mathcal{O}(N^3)$ .

For MinRes, the runtime is  $k$  matrix-vector products and  $\mathcal{O}(Nk)$  other operations. Since  $A$  is dense, the matrix-vector products dominate and the runtime is  $\mathcal{O}(N^2k)$ .

We know that MinRes requires  $k = \mathcal{O}(\kappa^{1/2})$  to achieve a fixed error. The critical  $\alpha$  is therefore  $\alpha = 2$ .

# Topics for Midterm Exam

## Example question

Consider the inner product

$$\langle p, q \rangle = \int_0^1 p(x) q(x) dx$$

defined for all polynomials  $p, q \in \mathcal{P}_k = \{p(x) \mid p(x) = \sum_{\ell=0}^k c_\ell x^\ell\}$ .

Determine three polynomials  $p_0, p_1, p_2$  such that  $\langle p_k, p_\ell \rangle = 0$  if  $k \neq \ell$  and

$$\text{span}\{p_0\} = \mathcal{P}_0, \quad \text{span}\{p_0, p_1\} = \mathcal{P}_1, \quad \text{span}\{p_0, p_1, p_2\} = \mathcal{P}_2.$$

# Topics for Midterm Exam

*Answer.* The polynomials can be found by applying the Gram-Schmidt algorithm to  $1, x, x^2$ . We compute:

$$p_0(x) = 1, \quad p_1(x) = x - \frac{\langle 1, x \rangle}{\langle 1, 1 \rangle} 1 = x - \frac{1}{2}$$

$$\begin{aligned} p_2(x) &= x^2 - \frac{\langle x - \frac{1}{2}, x^2 \rangle}{\langle x - \frac{1}{2}, x - \frac{1}{2} \rangle} \left(x - \frac{1}{2}\right) - \frac{\langle 1, x^2 \rangle}{\langle 1, 1 \rangle} 1 \\ &= x^2 - \left(x - \frac{1}{2}\right) - \frac{1}{3} \\ &= x^2 - x + \frac{1}{6}. \end{aligned}$$

# Topics for Midterm Exam

## Example question

Perform a single step of Newton's method for

$$f(x, y) = \begin{pmatrix} x^2 - xy + 1 \\ 2xy - 2x - y + 5 \end{pmatrix}$$

and initial guess  $(x_0, y_0) = (0, 1)$ .

# Topics for Midterm Exam

*Answer.* We compute

$$\nabla f(x, y) = \begin{pmatrix} 2x - y & -x \\ 2y - 2 & 2x - 1 \end{pmatrix} \implies \nabla f(0, 1) = \begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix}$$

and

$$f(0, 1) = \begin{pmatrix} 1 \\ 4 \end{pmatrix}.$$

Thus,

$$\begin{pmatrix} x_1 \\ y_1 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} - \begin{pmatrix} -1 & 0 \\ 0 & -1 \end{pmatrix}^{-1} \begin{pmatrix} 1 \\ 4 \end{pmatrix} = \begin{pmatrix} 1 \\ 5 \end{pmatrix}$$