Practical assignment 2 Algorithms & Datastructures

Carlo Jessurun s1013793 Tony Lopar s1013792

Nijmegen, January 11, 2018

Frits Vaandrager Joshua Moerman 2017-2018 Radboud University Nijmegen

Contents

1	$\mathbf{E}\mathbf{x}\mathbf{p}$	lanation	3
	1.1	Reading the input	:
	1.2	Algorithm	3
	1.3	Computing output	
2	Ana	lysis	4
	2.1	Correctness	4
		2.1.1 Reading the input	4
		2.1.2 Algorithm	4
		2.1.3 Computing output	4
	2.2	Complexity	4
		2.2.1 Reading the input	4
		2.2.2 Algorithm	4
		2.2.3 Writing output	4
3	Ref	erence	5

1 Explanation

In this chapter we will explain how our algorithm works. The explentation is split into three parts.

1.1 Reading the input

For the reading of the input we made a custom reader which replaces the scanner in Java. The custom reader works faster than the standard scanner in Java. Uitleg waarom sneller?

1.2 Algorithm

The algorithm first tries to make profitable groups of the input. Products are added to the previous group when after addition the group still gives profit.

After this the algorithm optimizes the list with profits by putting groups with a higher profit in front of the list.

Finally the first D profits are substracted from the total sum of products where D is the number of dividers.

1.3 Computing output

2 Analysis

2.1 Correctness

In this section we will discuss the Correctness of the processes.

2.1.1 Reading the input

2.1.2 Algorithm

2.1.3 Computing output

2.2 Complexity

In this chapter we will describe the complexity of our algorithm. We will first describe the complexity in detail for the smaller parts and after this compute the complexity for the whole algorithm.

2.2.1 Reading the input

The reading of the output has to process all items in the output for the algorithm. Before we read the prices of the products, we have to read the number of products and the number of dividers. This means that in total n + 2 items should be read from the input. The complexity of this reading is $\mathcal{O}(n)$.

2.2.2 Algorithm

2.2.3 Writing output

The output gets the total sum of products with the possible profit already substracted. The output should only be rounded, since the substraction may cause values that does not end with 0. The complexity of this process is $\mathcal{O}(1)$.

3 Reference