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#### 1 Basic

### 1.1 compile

```
# preset before coding
echo "cd ~/Desktop" >> ~/.bashrc
gedit -> preference -> tab width: 4
# Editor
gedit a.cpp
# Compile
g++ a.cpp -std=c++11
```

```
**All file will be compiled to a.out unless you use -o(
    not recommanded, just use a.out)**
./a.out
# Run with file input
./a.out < input.txt
# Run with file input and output
./a.out < input.txt > output.txt
# Python Run
python3 a.py < input.txt > output.txt
# Copy Paste In Ubuntu
* copy: ctrl+insert
* paste: shift+insert
# 比對文件相同
sdiff a.txt b.txt
```

#### 1.2 default code

1

2

9

9

```
#include <bits/stdc++.h>
using namespace std;
#define int long long
typedef pair<int,int> pii;
#ifdef ONLINE_JUDGE
#define cerr if(false) cerr
#endif
 int32_t main(){
#ifndef ONLINE JUDGE
   //freopen("input.txt","r",stdin);
freopen("output.txt","w",stdout);
freopen("debug.txt","w",stdcerr);
#else
   ios_base::sync_with_stdio(0);
   cin.tie(false);
#endif
| }
```

### 1.3 debug list

```
記得測試 python 的內建函數庫有哪些
bits/std++.h 跟 global variable y1 衝突,不能用
模板要記得 init
priority_queue 要清空
事先將把邊界測資加入測試
邊界條件 (過程溢位, 題目數據範圍), 會不會爆 long long
是否讀錯題目, 想不到時可以自己讀一次題目
比較容易有問題的地方換人寫
注意公式有沒有推錯或抄錯
精度誤差 sqrt(大大的東西) + EPS
喇分 random_suffle 隨機演算法
```

#### 時間複雜度 1.4

時間複雜度	可處理的最大 $N$ 數量級 (約)
O(1)	幾乎沒限制
$O(\log N)$	10 <sup>18</sup> 級別 (如快速冪)
$O(\sqrt{N})$	10 <sup>10</sup>
O(N)	10 <sup>8</sup>
$O(N \log N)$	$2 \times 10^7 \sim 5 \times 10^7$
$O(N\sqrt{N})$	$1 \times 10^5 \sim 2 \times 10^5$
$O(N^2)$	$10^4 \sim 1.5 \times 10^4$
$O(N^2 \log N)$	約 3×10 <sup>3</sup>
$O(N^3)$	$500 \sim 1000$
$O(2^N)$	$N \le 20$
O(N!)	$N \leq 10$

### 2 Dark Code

### 2.1 IO optimization

```
*if output to much, consider put all output in array
    first, then output the array.
getchar() -> getchar_unlocked()
fread() -> fread_unlocked()
inline char readchar() {
  const int S = 1<<20; // buffer size</pre>
  static char buf[S], *p = buf, *q = buf;
  if(p == q \&\& (q = (p=buf)+fread(buf,1,S,stdin)) ==
      buf) return EOF;
  return *p++;
inline int nxtint() {
 // if readchar can't use, change readchar() to
     getchar()
 int x = 0;
  int c = readchar(), neg = false;
 if (c == EOF) return -1;
 while (('0' > c | | c > '9') \&\& c != '-' \&\& c != EOF)
      c = readchar();
 if (c == '-')neg = true, c = readchar();
 while ('0' \le c \&\& c \le '9') \times = x * 10 + (c ^ '0'),
      c = readchar();
  if (neg) x = -x;
  return x;
}
```

## 3 Geometry

## 3.1 2D point

```
typedef double Double;
struct Point {
  Double x,y;
  bool operator < (const Point &b)const{</pre>
    //return tie(x,y) < tie(b.x,b.y);</pre>
    return atan2(y,x) < atan2(b.y,b.x);</pre>
  Point operator + (const Point &b)const{
    return (Point){x+b.x,y+b.y};
  Point operator - (const Point &b)const{
    return (Point){x-b.x,y-b.y};
  Point operator * (const Double &d)const{
    return Point(d*x,d*y);
  Double operator * (const Point &b)const{
    return x*b.x + y*b.y;
  Double operator % (const Point &b)const{
    return x*b.y - y*b.x;
  friend Double abs2(const Point &p){
    return p.x*p.x + p.y*p.y;
  friend Double abs(const Point &p){
    return sqrt( abs2(p) );
typedef Point Vector;
struct Line{
  Point P; Vector v;
  bool operator < (const Line &b)const{</pre>
    return atan2(v.y,v.x) < atan2(b.v.y,b.v.x);</pre>
};
```

### 3.2 兩線段交點

## 3.3 兩圓交點

### 3.4 Convex Hull

```
#include "2Dpoint.cpp"
// return H, The first will occured TWICE in vector H!
void ConvexHull(vector<Point> &P, vector<Point> &H){
    int n = P.size(), m=0;
    sort(P.begin(),P.end());
    H.clear();
    for (int i=0; i<n; i++){</pre>
        while (m>=2 && (P[i]-H[m-2]) % (H[m-1]-H[m-2])
            <0)H.pop_back(), m--;
        H.push_back(P[i]), m++;
    }
    for (int i=n-2; i>=0; i--){
        while (m>=2 && (P[i]-H[m-2]) % (H[m-1]-H[m-2])
             <0)H.pop_back(), m--;
        H.push_back(P[i]), m++;
    }
}
```

#### 4 Flow

#### 4.1 Dinic

```
4. see (b), check if it is possible.
5. answer is maxflow(ss, tt) + maxflow(s, t)
(b) Bounded Possible Flow:
1. same construction method as (a)
2. run maxflow(ss, tt)
3. for every edge connected with ss or tt:
        rule: check if their rest flow is exactly 0
4. answer is possible if every edge do satisfy the rule
5. otherwise, it is NOT possible.
(c) Bounded Minimum Flow:

    same construction method as (a)

2. answer is maxflow(ss, tt)
(d) Bounded Minimum Cost Flow:
* the concept is somewhat like bounded possible flow.

    same construction method as (a)

2. answer is maxflow(ss, tt) + (\Sigma 1 * cost for every)
   edge)
(e) Minimum Cut:

 run maxflow(s, t)

2. run cut(s)
3. ss[i] = 1: node i is at the same side with s.
const long long INF = 1LL<<60;</pre>
struct Dinic { //O(VVE), with minimum cut
    static const int MAXN = 5003;
    struct Edge{
        int u, v;
        long long cap, rest;
    int n, m, s, t, d[MAXN], cur[MAXN];
    vector<Edge> edges;
    vector<int> G[MAXN];
    void init(){
        edges.clear();
        for ( int i = 0 ; i < MAXN ; i++ ) G[i].clear()</pre>
    }
    // min cut start
    bool side[MAXN];
    void cut(int u) {
        side[u] = 1;
        for ( int i : G[u] ) {
            if ( !side[ edges[i].v ] && edges[i].rest )
                  cut(edges[i].v);
        }
    // min cut end
    void add_edge(int u, int v, long long cap){
   edges.push_back( {u, v, cap, cap} );
        edges.push_back( {v, u, 0, 0LL} );
        m = edges.size();
        G[u].push_back(m-2);
        G[v].push_back(m-1);
    }
    bool bfs(){
        memset(d, -1, sizeof(d));
        queue<int> que;
        que.push(s); d[s]=0;
        while (!que.empty()){
            int u = que.front(); que.pop();
            for (int ei : G[u]){
                 Edge &e = edges[ei];
                 if (d[e.v] < 0 && e.rest > 0){
                     d[e.v] = d[u] + 1;
                     que.push(e.v);
                 }
            }
        return d[t] >= 0;
    }
```

```
long long dfs(int u, long long a){
         if ( u == t || a == 0 ) return a;
         long long flow = 0, f;
         for ( int &i=cur[u]; i < (int)G[u].size(); i++</pre>
              ) {
             Edge &e = edges[ G[u][i] ];
            if ( d[u] + 1 != d[e.v] ) continue;
             f = dfs(e.v, min(a, e.rest));
             if (f > 0) {
                 e.rest -= f;
                 edges[ G[u][i]^1 ].rest += f;
                 flow += f;
                 a -= f;
                 if ( a == 0 )break;
            }
        return flow;
    long long maxflow(int s, int t){
         this->s = s, this->t = t;
         long long flow = 0, mf;
         while ( bfs() ){
            memset(cur, 0, sizeof(cur));
            while ( (mf = dfs(s, INF)) ) flow += mf;
         return flow;
    }
} dinic;
```

#### 4.2 min cost flow

```
// Long Long version
typedef pair<long long, long long> pll;
struct CostFlow {
    static const int MAXN = 350;
    static const long long INF = 1LL<<60;</pre>
    struct Edge {
        int to, r;
        long long rest, c;
    };
    int n, pre[MAXN], preL[MAXN]; bool inq[MAXN];
    long long dis[MAXN], fl, cost;
    vector<Edge> G[MAXN];
    void init() {
         for ( int i = 0 ; i < MAXN ; i++) G[i].clear();</pre>
    void add_edge(int u, int v, long long rest, long
         long c) {
         G[u].push\_back({v, (int)G[v].size()}, rest, c
             });
         G[v].push_back({u, (int)G[u].size()-1, 0, -c});
    pll flow(int s, int t) {
        fl = cost = 0;
         while (true) {
             fill(dis, dis+MAXN, INF);
fill(inq, inq+MAXN, 0);
             dis[s] = 0;
             queue<int> que;
             que.push(s);
             while ( !que.empty() ) {
                 int u = que.front(); que.pop();
                 inq[u] = 0;
                 for ( int i = 0 ; i < (int)G[u].size()</pre>
                      ; i++) {
                      int v = G[u][i].to;
                      long long w = G[u][i].c;
                     if ( G[u][i].rest > 0 && dis[v] >
                          dis[u] + w) {
                          pre[v] = u; preL[v] = i;
                          dis[v] = dis[u] + w;
                          if (!inq[v]) {
                              inq[v] = 1;
                              que.push(v);
                          }
                     }
                 }
             }
```

```
if (dis[t] == INF) break;
long long tf = INF;
for (int v = t, u, 1; v != s; v = u) {
        u = pre[v]; 1 = preL[v];
        tf = min(tf, G[u][1].rest);
}
for (int v = t, u, 1; v != s; v = u) {
        u = pre[v]; 1 = preL[v];
        G[u][1].rest -= tf;
        G[v][G[u][1].r].rest += tf;
}
cost += tf * dis[t];
fl += tf;
}
return {fl, cost};
}
flow;
```

### 5 Mathmatics

### 5.1 ax+by=gcd(a,b)

```
typedef pair<int, int> pii;

pii exgcd(int a, int b){
   if(b == 0) return make_pair(1, 0);
   else{
      int p = a / b;
      pii q = exgcd(b, a % b);
      int aa = q.second, bb = q.first - q.second * p;
      if(aa < 0) aa += b, bb -= a;
      return make_pair(aa, bb);
   }
}</pre>
```

### 5.2 BigInt

```
struct Bigint{
  static const int LEN = 60;
  static const int BIGMOD = 10000;
  int s;
  int v1, v[LEN];
  // vector<int> v;
  Bigint() : s(1) \{ vl = 0; \}
  Bigint(long long a) {
    s = 1; vl = 0;
    if (a < 0) \{ s = -1; a = -a; \}
    while (a) {
      push_back(a % BIGMOD);
      a /= BIGMOD;
    }
  Bigint(string str) {
    s = 1; v1 = 0;
    int stPos = 0, num = 0;
    if (!str.empty() && str[0] == '-') {
      stPos = 1;
      s = -1;
    for (int i=SZ(str)-1, q=1; i>=stPos; i--) {
      num += (str[i] - '0') * q;
      if ((q *= 10) >= BIGMOD) {
        push_back(num);
        num = 0; q = 1;
      }
    if (num) push_back(num);
  int len() const { return vl; /* return SZ(v); */ }
  bool empty() const { return len() == 0; }
void push_back(int x) { v[vl++] = x; /* v.PB(x); */ }
  void pop_back() { v1--; /* v.pop_back(); */ }
  int back() const { return v[v1-1]; /* return v.back()
  void n() { while (!empty() && !back()) pop_back(); }
  void resize(int nl) {
```

```
vl = nl; fill(v, v+vl, 0);
// v.resize(nl); // fill(ALL(v), 0);
}
void print() const {
  if (empty()) { putchar('0'); return; }
  if (s == -1) putchar('-');
  printf("%d", back());
  for (int i=len()-2; i>=0; i--) printf("%.4d",v[i]);
friend std::ostream& operator << (std::ostream& out,</pre>
    const Bigint &a) {
  if (a.empty()) { out << "0"; return out; }</pre>
  if (a.s == -1) out << "-";</pre>
  out << a.back();</pre>
  for (int i=a.len()-2; i>=0; i--) {
    char str[10];
    snprintf(str, 5, "%.4d", a.v[i]);
    out << str;
  return out;
int cp3(const Bigint &b)const {
  if (s != b.s) return s > b.s ? 1 : -1;
if (s == -1) return -(-*this).cp3(-b);
  if (len() != b.len()) return len()>b.len()?1:-1;
  for (int i=len()-1; i>=0; i--)
    if (v[i]!=b.v[i]) return v[i]>b.v[i]?1:-1;
  return 0;
bool operator < (const Bigint &b)const{return cp3(b)</pre>
bool operator <= (const Bigint &b)const{return cp3(b)</pre>
    <=0;}
bool operator >= (const Bigint &b)const{return cp3(b)
    >=0;}
bool operator == (const Bigint &b)const{return cp3(b)
bool operator != (const Bigint &b)const{return cp3(b)
    !=0;}
bool operator > (const Bigint &b)const{return cp3(b)
    ==1:}
Bigint operator - () const {
  Bigint r = (*this);
  r.s = -r.s;
  return r;
Bigint operator + (const Bigint &b) const {
  if (s == -1) return -(-(*this)+(-b));
  if (b.s == -1) return (*this)-(-b);
  Bigint r;
  int nl = max(len(), b.len());
  r.resize(nl + 1);
  for (int i=0; i<nl; i++) {</pre>
    if (i < len()) r.v[i] += v[i];</pre>
    if (i < b.len()) r.v[i] += b.v[i];</pre>
    if(r.v[i] >= BIGMOD) {
      r.v[i+1] += r.v[i] / BIGMOD;
      r.v[i] %= BIGMOD;
  }
  r.n();
  return r;
Bigint operator - (const Bigint &b) const {
  if (s == -1) return -(-(*this)-(-b));
  if (b.s == -1) return (*this)+(-b);
  if ((*this) < b) return -(b-(*this));</pre>
  Bigint r;
  r.resize(len());
  for (int i=0; i<len(); i++) {</pre>
    r.v[i] += v[i];
    if (i < b.len()) r.v[i] -= b.v[i];</pre>
    if (r.v[i] < 0) {</pre>
      r.v[i] += BIGMOD;
      r.v[i+1]--;
    }
  }
  r.n();
  return r;
Bigint operator * (const Bigint &b) {
  Bigint r;
```

```
r.resize(len() + b.len() + 1);
    r.s = s * b.s;
    for (int i=0; i<len(); i++) {</pre>
      for (int j=0; j<b.len(); j++) {</pre>
        r.v[i+j] += v[i] * b.v[j];
        if(r.v[i+j] >= BIGMOD) {
          r.v[i+j+1] += r.v[i+j] / BIGMOD;
          r.v[i+j] %= BIGMOD;
        }
      }
    }
    r.n();
    return r;
  Bigint operator / (const Bigint &b) {
    Bigint r;
    r.resize(max(1, len()-b.len()+1));
    int oriS = s;
    Bigint b2 = b; // b2 = abs(b)
    s = b2.s = r.s = 1;
    for (int i=r.len()-1; i>=0; i--) {
      int d=0, u=BIGMOD-1;
      while(d<u) {</pre>
        int m = (d+u+1)>>1;
        r.v[i] = m;
        if((r*b2) > (*this)) u = m-1;
        else d = m;
      }
      r.v[i] = d;
    }
    s = oriS;
    r.s = s * b.s;
    r.n();
    return r;
  Bigint operator % (const Bigint &b) {
    return (*this)-(*this)/b*b;
};
```

#### 5.3 GaussElimination

```
// by bcw_codebook
const int MAXN = 300:
const double EPS = 1e-8;
int n;
double A[MAXN][MAXN];
void Gauss() {
  for(int i = 0; i < n; i++) {</pre>
     bool ok = 0;
     for(int j = i; j < n; j++) {</pre>
       if(fabs(A[j][i]) > EPS) {
         swap(A[j], A[i]);
         ok = 1;
         break;
       }
    if(!ok) continue;
     double fs = A[i][i];
     for(int j = i+1; j < n; j++) {</pre>
       double r = A[j][i] / fs;
       for(int k = i; k < n; k++) {</pre>
         A[j][k] -= A[i][k] * r;
       }
    }
  }
}
template < class T>
void Gauss(vector<vector<T>> &A) {
  int n = A.size();
  for(int i = 0; i < n; i++) {</pre>
    bool ok = 0;
     for(int j = i; j < n; j++) {</pre>
       if(A[j][i] != 0) {
         swap(A[j], A[i]);
```

```
ok = 1;
break;
}

if(!ok) continue;

T fs = A[i][i];
for(int j = i+1; j < n; j++) {
    T r = A[j][i] / fs;
    for(int k = i; k < n; k++) {
        A[j][k] -= A[i][k] * r;
    }
}
}</pre>
```

#### 5.4 Inverse

```
int inverse[100000];
void invTable(int b, int p) {
   inverse[1] = 1;
   for( int i = 2; i <= b; i++ ) {
      inverse[i] = (long long)inverse[p%i] * (p-p/i) % p;
   }
}
int inv(int b, int p) {
   return b == 1 ? 1 : ((long long)inv(p % b, p) * (p-p/b) % p);
}</pre>
```

### 5.5 LinearPrime 歐拉篩

```
const int MAXP = 100; //max prime
vector<int> P; // primes
void build_prime(){
    static bitset<MAXP> ok;
    int np=0;
    for (int i=2; i<MAXP; i++){
        if (ok[i]==0)P.push_back(i), np++;
        for (int j=0; j<np && i*P[j]<MAXP; j++){
        ok[ i*P[j] ] = 1;
        if ( i%P[j]==0 )break;
    }
}</pre>
```

## 5.6 Miller Rabin

```
typedef long long LL;
inline LL bin_mul(LL a, LL n,const LL& MOD){
 LL re=0;
 while (n>0){
   if (n&1) re += a;
    a += a; if (a>=MOD) a-=MOD;
   n>>=1:
 return re%MOD;
inline LL bin_pow(LL a, LL n,const LL& MOD){
 LL re=1;
 while (n>0){
   if (n&1) re = bin_mul(re,a,MOD);
    a = bin_mul(a,a,MOD);
   n>>=1;
 }
 return re;
}
bool is_prime(LL n){
 //static LL sprp[3] = { 2LL, 7LL, 61LL};
  static LL sprp[7] = { 2LL, 325LL, 9375LL,
    28178LL, 450775LL, 9780504LL,
    1795265022LL };
```

```
if (n=1 || (n&1)=0 ) return n==2;
int u=n-1, t=0;
while ( (u&1)=0 ) u>>=1, t++;
for (int i=0; i<3; i++){
   LL x = bin_pow( sprp[i]%n, u, n);
   if (x==0 || x==1 || x==n-1)continue;

   for (int j=1; j<t; j++){
        x=x*x%n;
        if (x==1 || x==n-1)break;
    }
   if (x==n-1)continue;
   return 0;
}
return 1;
}</pre>
```

#### 5.7 Pollard's rho

```
// from PEC
// does not work when n is prime
Int f(Int x, Int mod){
   return add(mul(x, x, mod), 1, mod);
}
Int pollard_rho(Int n) {
   if ( !(n & 1) ) return 2;
   while (true) {
      Int y = 2, x = rand()%(n-1) + 1, res = 1;
      for ( int sz = 2 ; res == 1 ; sz *= 2 ) {
        for ( int i = 0 ; i < sz && res <= 1 ; i++) {
            x = f(x, n);
            res = __gcd(abs(x-y), n);
      }
      y = x;
   }
   if ( res != 0 && res != n ) return res;
}
</pre>
```

#### 5.8 數論基本工具

```
Int POW(Int a, Int n, Int mod){
    Int re=1;
    while (n>0){
        if (n&1LL) re = re*a%mod;
        a = a*a%mod;
        n>>=1;
    }
    return re;
}

Int C(Int n, Int m){
    if (m<0 || m>n)return 0;
    return J[n] * inv(J[m]*J[n-m]%MOD) %MOD;
}
```

#### 5.9 Mobius

#### 5.10 SG

```
|Anti Nim (取走最後一個石子者敗)
先手必勝 if and only if
1. 「所有」堆的石子數都為 1 且遊戲的 SG 值為 0。
2. 「有些」堆的石子數大於 1 且遊戲的 SG 值不為 0。
Anti-SG (決策集合為空的遊戲者贏)
定義 SG 值為 0 時,遊戲結束,
則先手必勝 if and only if
1. 遊戲中沒有單一遊戲的 SG 函數大於 1 且遊戲的 SG 函數
   為 0。
2. 遊戲中某個單一遊戲的 SG 函數大於 1 且遊戲的 SG 函數
   不為 0。
______
Sprague-Grundy
1. 雙人、回合制
2. 資訊完全公開
3. 無隨機因素
4. 可在有限步內結束
5. 沒有和局
6. 雙方可採取的行動相同
SG(S) 的值為 0:後手(P)必勝
不為 0: 先手(N)必勝
int mex(set S) {
 // find the min number >= 0 that not in the S
 // e.g. S = \{0, 1, 3, 4\} mex(S) = 2
state = []
int SG(A) {
 if (A not in state) {
   S = sub_states(A)
   if( len(S) > 1 ) state[A] = reduce(operator.xor, [
      SG(B) for B in S])
   else state[A] = mex(set(SG(B) for B in next_states(
      A)))
 return state[A]
5.11
     Theorem
```

```
Lucas's Theorem
 For non-negative integer n,m and prime P,
  C(m,n) \mod P = C(m/M,n/M) * C(m%M,n%M) \mod P
  = mult_i ( C(m_i,n_i) )
  where m_i is the i-th digit of m in base P.
Pick's Theorem
 A = i + b/2 - 1
Kirchhoff's theorem
  A_{\{ii\}} = deg(i), A_{\{ij\}} = (i,j) \setminus in E ? -1 : 0
  Deleting any one row, one column, and cal the det(A)
Nth Catalan recursive function:
C_0 = 1, C_{n+1} = C_n * 2(2n + 1)/(n+2)
Mobius Formula

    , if n = 1
    , 若 n 無平方數因數,且 n = p1*p2*p3

u(n) = 1
           *...*pk
               ,若 n 有大於 1 的平方數因數
- Property
1. (積性函數) u(a)u(b) = u(ab)
2. \sum_{d|n} u(d) = [n == 1]
Mobius Inversion Formula
```

```
f(n) = \sum_{d \in \mathbb{Z}} \{d \mid n\} \ g(d)
g(n) = \sum_{d \in \mathbb{Z}} \{d \mid n\} \ u(n/d)f(d)
if
then
               = \sum \{d/n\} \ u(d)f(n/d)
- Application
the number/power of gcd(i, j) = k
- Trick
分塊, O(sqrt(n))
Chinese Remainder Theorem (m_i 兩兩互質)
  x = a_1 \pmod{m_1}
 x = a_2 \pmod{m_2}
 x = a_i \pmod{m_i}
construct a solution:
  Let M = m_1 * m_2 * m_3 * ... * m_n
 Let M_i = M / m_i
  t_i = 1 / M_i
  t_i * M_i = 1 \pmod{m_i}
 solution x = a_1 * t_1 * M_1 + a_2 * t_2 * M_2 + \dots + a_n * t_n * M_n + k * M
= k*M + \sum a_i * t_i * M_i, k is positive integer.
  under mod M, there is one solution x = \sum a_i * t_i *
      M_i
_____
Burnside's Lemma
|G| * |X/G| = sum(|X^g|) where g in G
總方法數:每一種旋轉下不動點的個數總和 除以 旋轉的方法
```

## 6 Graph

### 6.1 BCC

```
邊雙連誦
任意兩點間至少有兩條不重疊的路徑連接,找法:
1. 標記出所有的橋
2. 對全圖進行 DFS,不走橋,每一次 DFS 就是一個新的邊雙
// from BCW
struct BccEdge {
  static const int MXN = 100005;
  struct Edge { int v,eid; };
  int n,m,step,par[MXN],dfn[MXN],low[MXN];
  vector<Edge> E[MXN];
  DisjointSet djs;
  void init(int _n) {
   n = _n; m = 0;
for (int i=0; i<n; i++) E[i].clear();</pre>
    djs.init(n);
  void add_edge(int u, int v) {
    E[u].PB({v, m});
    E[v].PB({u, m});
   m++;
  void DFS(int u, int f, int f_eid) {
    par[u] = f;
    dfn[u] = low[u] = step++;
    for (auto it:E[u]) {
      if (it.eid == f_eid) continue;
     int v = it.v;
     if (dfn[v] == -1) {
       DFS(v, u, it.eid);
       low[u] = min(low[u], low[v]);
      } else {
       low[u] = min(low[u], dfn[v]);
```

```
}
}
void solve() {
    step = 0;
    memset(dfn, -1, sizeof(int)*n);
    for (int i=0; i<n; i++) {
        if (dfn[i] == -1) DFS(i, i, -1);
    }
    djs.init(n);
    for (int i=0; i<n; i++) {
        if (low[i] < dfn[i]) djs.uni(i, par[i]);
    }
}
graph;</pre>
```

#### 6.2 Prim

```
// edge strucute
struct edge{
  int a, b;
  double data;
  bool operator <(const edge b)const{</pre>
    return data > b.data;
  }
};
// main prim algorithm
int n, m, root, aa, bb, cc;
while (cin >> n >> m){
  priority_queue<edge>yee;
  int visit[500] = {}, p[500] = {};
  double a[500][500] = {};
  //undirectional edge aa to bb is weighted cc
  for (int i = 0; i < m; i++){</pre>
    cin >> aa >> bb >> cc;
    a[aa][bb] = a[bb][aa] = cc;
  cin >> root;
  yee.push({ 0, root, 0 });
  edge tmp;
  double total = 0;
  while (!yee.empty()){
    tmp = yee.top(); yee.pop();
    if (visit[tmp.b])continue;
    total += tmp.data; p[tmp.b] = tmp.a; visit[tmp.b] =
         1;
    for (int i = 1; i <= n; i++){</pre>
      if (a[tmp.b][i]!=.0&&(!visit[i])){
        yee.push({tmp.b,i,a[tmp.b][i]});
      }
    }
  cout << total << endl;</pre>
```

#### 6.3 Bellman Ford

```
int a[100][100], d[100], p[100];
void bellman_ford(int root, int n){
  for (int i = 1; i <= n; i++)d[i] = 1e9;</pre>
  d[root] = 0, p[root] = 0;
  for (int i = 0; i<n - 1; i++){</pre>
    for (int j = 1; j <= n; j++){</pre>
      for (int k = 1; k <= n; k++){
         if (d[j] != 1e9 && a[j][k] != 1e9){
           if (d[j] + a[j][k] < d[k]){</pre>
             d[k] = d[j] + a[j][k], p[k] = j;
           }
        }
      }
    }
  }
}
bool nega_cyc(int n){
  for (int i = 1; i <= n; i++){</pre>
    for (int j = 1; j <= n; j++){</pre>
```

```
if (d[i] != 1e9 && a[i][j] != 1e9)
      if (d[i] + a[i][j] < d[j]){</pre>
        return 0;
    }
  }
  return 1;
int main(){
  int n, m, aa, bb, dd;
  while (cin >> n >> m){
    for (int i = 0; i <= n; i++)for (int j = 0; j <= n;
          j++){
      a[i][j] = E9;
    memset(p, 0, sizeof(p));
    for (int i = 0; i < m; i++){</pre>
      cin >> aa >> bb >> dd;
      a[aa][bb] = min(a[aa][bb], dd);
    cin >> aa;
    bellman_ford(aa, n);
    int t = nega_cyc(n);
    if(t){
      for (int i = 1; i <= n; i++)cout << d[i] << " \n"</pre>
           [i==n];
      for (int i = 1; i <= n; i++)cout << p[i] << " \n"</pre>
           [i==n];
    else cout << "There is a negative weight cycle in</pre>
         the graph\n";
  }
}
```

#### 6.4 Kruskal

```
struct v {
  int a, b, c;
int p[200001];v a[200001];
bool sor(v a, v b) {
 return a.c < b.c;</pre>
int find(int x) {
  return(x != p[x] ? (p[x] = find(p[x])) : x);
}
int main() {
  int n, m, i, j, sum;
  while (cin >> n >> m) {
    for (i = 0; i < 200001; i++)p[i] = i;</pre>
    for (i = 0; i<m; i++)cin >> a[i].a >> a[i].b >> a[i
         ].c;
    sort(a, a + m, sor);
    for (i = 0, j = 0; j < m; j++) {
      if(find(a[j].a) != find(a[j].b)){
        p[find(a[j].a)] = find(a[j].b);
        sum += a[j].c;
    cout << ((i==n-1)?sum:-1) << endl;</pre>
  }
}
```

### 6.5 Dijkstra

```
struct node {
   int num{}, w{};
   bool operator < (const node& other)const {
      return w > other.w;
   }
};
```

```
vector<int> dijkstra(int root, const vector<vector<node
    >> &graph) {
    vector<int> d(graph.size(), INT_MAX >> 1), p(graph.
        size());
    priority_queue<node> pq;
    d[root] = p[root] = 0;
    pq.push({root, d[root]});
    while (!pq.empty()) {
        node tmp = pq.top(); pq.pop();
        for (const node &i : graph[tmp.num]) {
            if (d[i.num] > d[tmp.num] + i.w) {
                 d[i.num] = d[tmp.num] + i.w;
                p[i.num] = tmp.num;
                pq.push({i.num, d[tmp.num]});
            }
        }
    return d;
}
```

### 6.6 Strongly Connected Component(SCC)

```
#define MXN 100005
 #define PB push_back
#define FZ(s) memset(s,0,sizeof(s))
 struct Scc{
int n, nScc, vst[MXN], bln[MXN];
 vector<int> E[MXN], rE[MXN], vec;
 void init(int _n){
  n = _n;
   for (int i=0; i<MXN; i++){</pre>
    E[i].clear();
     rE[i].clear();
  }
 void add_edge(int u, int v){
  E[u].PB(v);
  rE[v].PB(u);
void DFS(int u){
  vst[u]=1;
   for (auto v : E[u])
    if (!vst[v]) DFS(v);
  vec.PB(u);
void rDFS(int u){
  vst[u] = 1;
  bln[u] = nScc;
   for (auto v : rE[u])
     if (!vst[v]) rDFS(v);
 void solve(){
  nScc = 0;
  vec.clear();
   FZ(vst);
   for (int i=0; i<n; i++)</pre>
     if (!vst[i]) DFS(i);
   reverse(vec.begin(),vec.end());
  FZ(vst);
   for (auto v : vec){
    if (!vst[v]){
      rDFS(v);
       nScc++;
     }
  }
}
|};
```

### 6.7 Hungarian

```
// Maximum Cardinality Bipartite Matching
struct Graph {
    static const int MAXN = 5005;
    vector<int> G[MAXN];
    int n;
```

```
int match[MAXN]; // Matching Result
    int vis[MAXN];
    void init(int _n) {
         for ( int i = 0 ; i < n ; i++ ) G[i].clear();</pre>
    bool dfs(int u) {
         for ( auto v:G[u] ) {
              if (!vis[v]) {
                  vis[v] = true;
                  if (match[v] == -1 || dfs(match[v])) {
                       match[v] = u;
                       match[u] = v;
                       return true;
                  }
              }
         return false;
    int solve() {
         int res = 0;
         memset(match, -1, sizeof(match));
for (int i = 0; i < n; i++) {</pre>
              if (match[i] == -1) {
                  memset(vis, 0, sizeof(vis));
                  if (dfs(i)) res += 1;
         return res;
} graph;
```

#### 6.8 KM

```
Detect non-perfect-matching:

    set all edge[i][j] as INF

2. if solve() >= INF, it is not perfectmatching.
// Maximum Weight Perfect Bipartite Matching
// allow negative weight!
typedef long long Int;
struct KM {
    static const int MAXN = 1050;
    static const int INF = 1LL<<60;</pre>
    int n, match[MAXN], vx[MAXN], vy[MAXN];
    Int edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[
        MAXN];
    void init(int _n){
        n = _n;
        for ( int i = 0 ; i < n ; i++ )</pre>
            for ( int j = 0; j < n ; j++ )
                edge[i][j] = 0;
    void add_edge(int x, int y, Int w){
        edge[x][y] = w;
    bool DFS(int x){
        vx[x] = 1;
        for ( int y = 0; y < n; y++) {
            if ( vy[y] ) continue;
            if (lx[x] + ly[y] > edge[x][y]) {
                slack[y] = min(slack[y], lx[x] + ly[y]
                      edge[x][y]);
                vy[y] = 1;
                if ( match[y] == -1 || DFS(match[y]) ){
                    match[y] = x;
                    return true;
                }
            }
        return false;
    Int solve() {
        fill(match, match + n, -1);
        fill(lx, lx + n, -INF);
```

```
fill(ly, ly + n, 0);
         for ( int i = 0; i < n; i++ )</pre>
              for ( int j = 0; j < n; j++ )</pre>
                  lx[i] = max(lx[i], edge[i][j]);
         for ( int i = 0 ; i < n; i++ ) {</pre>
              fill(slack, slack + n, INF);
              while (true){
                  fill(vx, vx + n, 0);
                  fill(vy, vy + n, 0);
                  if ( DFS(i) ) break;
                  Int d = INF;
                  for ( int j = 0 ; j < n ; j++ )</pre>
                      if ( !vy[j] ) d = min(d, slack[j]);
                  for ( int j = 0 ; j < n ; j++ ) {</pre>
                      if (vx[j]) lx[j] -= d;
                       if (vy[j]) ly[j] += d;
                      else slack[j] -= d;
                  }
             }
         Int res = 0;
         for ( int i = 0 ; i < n ; i++ ) {</pre>
             res += edge[ match[i] ][i];
         return res:
     }
|} graph;
```

### 6.9 最小平均環

```
// from BCW
/* minimum mean cycle */
const int MAXE = 1805;
 const int MAXN = 35;
 const double inf = 1029384756;
 const double eps = 1e-6;
 struct Edge {
  int v,u;
  double c;
int n,m,prv[MAXN][MAXN], prve[MAXN][MAXN], vst[MAXN];
 Edge e[MAXE];
 vector<int> edgeID, cycle, rho;
 double d[MAXN][MAXN];
 inline void bellman_ford() {
  for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {</pre>
     fill(d[i+1], d[i+1]+n, inf);
     for(int j=0; j<m; j++) {</pre>
       int v = e[j].v, u = e[j].u;
       if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
         d[i+1][u] = d[i][v]+e[j].c;
         prv[i+1][u] = v;
         prve[i+1][u] = j;
       }
    }
  }
 double karp_mmc() {
  // returns inf if no cycle, mmc otherwise
   double mmc=inf;
   int st = -1;
   bellman_ford();
   for(int i=0; i<n; i++) {</pre>
     double avg=-inf;
     for(int k=0; k<n; k++) {</pre>
       if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i])</pre>
           /(n-k):
       else avg=max(avg,inf);
     if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
   for(int i=0; i<n; i++) vst[i] = 0;</pre>
   edgeID.clear(); cycle.clear(); rho.clear();
   for (int i=n; !vst[st]; st=prv[i--][st]) {
    vst[st]++:
     edgeID.PB(prve[i][st]);
     rho.PB(st);
```

```
while (vst[st] != 2) {
   int v = rho.back(); rho.pop_back();
   cycle.PB(v);
   vst[v]++;
}
reverse(ALL(edgeID));
edgeID.resize(SZ(cycle));
return mmc;
}
```

### 6.10 偵測負環

```
#include <bits/stdc++.h>
using namespace std;
const int INF = 1000000;
const int MAXN = 200;
int n, m, q;
int d[MAXN][MAXN];
int main () {
    while ( cin >> n >> m >> q && n) {
        for ( int i = 0 ; i <= n ; i++ ) {
            for ( int j = 0 ; j <= n ; j++ ) d[i][j] =</pre>
                 (i==j?0:INF);
        for ( int i = 0 ; i < m ; i++ ) {</pre>
            int a, b, c;
            cin >> a >> b >> c;
            d[a][b] = min(d[a][b], c);
        for ( int k = 0 ; k < n ; k++ ) {
            for ( int i = 0 ; i < n ; i++ ) {</pre>
                for ( int j = 0 ; j < n ; j++ ) {</pre>
                     if ( d[i][j] > d[i][k] + d[k][j] &&
                          d[i][k] < INF && d[k][j] < INF
                          ) {
                         //printf("%d > %d + %d\n", d[i
                             ][j], d[i][k], d[k][j]);
                         k << "--";
                         d[i][j] = min(d[i][j], d[i][k]
                             + d[k][j]);
                    }
                }
            }
        }
        for ( int i = 0 ; i < n ; i++ ) {</pre>
            for ( int j = 0 ; j < n ; j++ ) {</pre>
                 for ( int k = 0 ; k < n && d[i][j] != -
                     INF ; k++ ) {
                     if ( d[k][k] < 0 && d[i][k] != INF</pre>
                         && d[k][j] != INF )
                         d[i][j] = -INF;
                }
            }
        int u, v;
        for (int i=0;i<q;i++){</pre>
            scanf("%d%d",&u,&v);
            if (d[u][v] == INF) printf("Impossible\n");
            else if (d[u][v] == -INF) printf("-Infinity
                n'';
            else printf("%d \ n", d[u][v]);
        puts("");
    return 0;
}
```

### 6.11 Tarjan

```
點 u 為割點 if and only if 滿足 1. or 2.
1. u 爲樹根,且 u 有多於一個子樹。
2. u 不爲樹根,且滿足存在 (u,v) 爲樹枝邊 (或稱父子邊,
    即 u 爲 v 在搜索樹中的父親),使得 DFN(u) <= Low(v)
一條無向邊 (u,v) 是橋 if and only if (u,v) 爲樹枝邊,且
    滿足 DFN(u) < Low(v)。
// 0 base
struct TarjanSCC{
 static const int MAXN = 1000006;
  int n, dfn[MAXN], low[MAXN], scc[MAXN], scn, count;
  vector<int> G[MAXN];
  stack<int> stk;
  bool ins[MAXN];
  void tarjan(int u){
    dfn[u] = low[u] = ++count;
    stk.push(u);
    ins[u] = true;
   for(auto v:G[u]){
     if(!dfn[v]){
       tarjan(v);
       low[u] = min(low[u], low[v]);
      }else if(ins[v]){
       low[u] = min(low[u], dfn[v]);
   if(dfn[u] == low[u]){
     int v;
     do {
      v = stk.top();
      stk.pop();
      scc[v] = scn;
      ins[v] = false;
      } while(v != u);
      scn++;
 }
  void getSCC(){
    memset(dfn,0,sizeof(dfn));
    memset(low,0,sizeof(low));
   memset(ins,0,sizeof(ins));
    memset(scc,0,sizeof(scc));
    count = scn = 0;
    for(int i = 0 ; i < n ; i++ ){</pre>
     if(!dfn[i]) tarjan(i);
 }
}SCC;
```

#### 6.12 Topological Sort

```
visit[s] = 1;
    for (int t=0; t<N; ++t){</pre>
         if (adj[s][t]) DFS(t);
    visit[s] = 2;
    order[n--] = s;
                           // record the order
void topological_ordering()
    memset(visit, 0, sizeof(visit));
    cycle = false;
    n = N - 1;
    for (int s=0; s<9; ++s)</pre>
         if (!v[s])
             DFS(s);
    if (cycle) cout << "The graph has the cycle!";</pre>
         for (int i=0; i<N; ++i)</pre>
             cout << order[i];</pre>
  }
}
```

## 7 Data Structure

### 7.1 2D Range Tree

```
// remember sort x !!!!!
typedef int T;
const int LGN = 20;
const int MAXN = 100005;
struct Point{
   Тх, у;
    friend bool operator < (Point a, Point b){</pre>
        return tie(a.x,a.y) < tie(b.x,b.y);</pre>
struct TREE{
    Point pt;
    int toleft;
}tree[LGN][MAXN];
struct SEG{
    T mx, Mx;
    int sz;
    TREE *st;
}seg[MAXN*4];
vector<Point> P;
void build(int 1, int r, int o, int deep){
    seg[o].mx = P[1].x;
    seg[o].Mx = P[r].x;
    seg[o].sz = r-l+1;;
    if(1 == r){
        tree[deep][r].pt = P[r];
        tree[deep][r].toleft = 0;
        seg[o].st = &tree[deep][r];
        return;
    int mid = (l+r)>>1;
    build(1,mid,o+o,deep+1);
    build(mid+1,r,o+o+1,deep+1);
    TREE *ptr = &tree[deep][1];
    TREE *pl = &tree[deep+1][l], *nl = &tree[deep+1][
    TREE *pr = &tree[deep+1][mid+1], *nr = &tree[deep
        +1][r+1];
    int cnt = 0:
    while(pl != nl && pr != nr) {
        *(ptr) = pl->pt.y <= pr->pt.y ? cnt++, *(pl++):
              *(pr++);
```

### 7.2 Sparse Table

```
const int MAXN = 200005;
const int lgN = 20;
/* Sp[i][j] 為 區間 [i, i + 2^j - 1] 的值 */
/* 從 i 開始 長度為 2 ^ j */
/*解決可重複貢獻問題 */
struct SP{ //sparse table
  int Sp[MAXN][lgN];
  function<int(int,int)> opt;
  void build(int n, int *a){ // 0 base
    for (int i=0 ;i<n; i++) Sp[i][0]=a[i];</pre>
    for (int h=1; h<lgN; h++){</pre>
      int len = 1<<(h-1), i=0;</pre>
      for (; i+len<n; i++)</pre>
        Sp[i][h] = opt(Sp[i][h-1], Sp[i+len][h-1]);
      for (; i<n; i++)</pre>
        Sp[i][h] = Sp[i][h-1];
    }
  int query(int 1, int r){
    int h = __lg(r-l+1);
    int len = 1<<h;</pre>
    return opt( Sp[1][h] , Sp[r-len+1][h] );
};
```

### 7.3 Segment Tree

```
struct Node{
   int mx; // 區間最大值
   int tag; // 子樹裡所有人的'值'都要加上 tag
};
vector<Node> seg;
// 節點 id 的整個區間要加上 tag
void addtag(int tag, int id){
   seg[id].mx += tag; // 最大值會加上 tag
   seg[id].tag += tag; // 注意可能本來就有標記了,所以
       是 +=
}
// 更新子節點資訊並把標記移到子節點身上
void push(int id){
   addtag(seg[id].tag, lc);
   addtag(seg[id].tag, rc);
   seg[id].tag = 0; // 標記被移到子節點上所以要改成 0
// 區間 [l,r] 加上 v
void modify(int 1, int r, int v, int L, int R, int id){
   if(1 <= L && R <= r){
      addtag(v, id);
       return;
   push(id);
```

```
if(r <= M) modify(1, r, v, L, M, lc);</pre>
    else if(1 > M) modify(1, r, v, M + 1, R, rc);
    else{
        modify(1, r, v, L, M, lc);
        modify(l, r, v, M + 1, R, rc);
    seg[id].mx = max(seg[lc].mx, seg[rc].mx);
}
int query(int 1, int r, int L, int R, int id){
    if(1 <= L && R <= r) return seg[id].mx;</pre>
    push(id);
    int M = (L + R) / 2;
    if(r <= M) return query(l, r, L, M, lc);</pre>
    else if(1 > M) return query(1, r, M + 1, R, rc);
    else return max(query(1, r, L, M, lc),
                     query(1, r, M + 1, R, rc));
}
```

### 7.4 ZKW 線段樹

```
const int M=1e5+111;
int n,m,q;
int sum[M<<2],mn[M<<2],mx[M<<2],add[M<<2];</pre>
int read() {
    int x;
    cin >> x;
    return x;
void build(){
    for ( m=1; m <= n; m < <= 1);</pre>
    for(int i=m+1;i<=m+n;++i)</pre>
        sum[i]=mn[i]=mx[i]=read();
    for(int i=m-1;i;--i){
        sum[i]=sum[i<<1]+sum[i<<1|1];</pre>
        mn[i]=min(mn[i<<1],mn[i<<1|1]),</pre>
        mn[i << 1] -= mn[i], mn[i << 1|1] -= mn[i];
        mx[i]=max(mx[i<<1],mx[i<<1|1]),
        mx[i << 1] -= mx[i], mx[i << 1|1] -= mx[i];
    }
void update_node(int x,int v,int A=0){
    x+=m, mx[x]+=v, mn[x]+=v, sum[x]+=v;
    for(;x>1;x>>=1){
        sum[x]+=v;
        A=min(mn[x],mn[x^1]);
        mn[x]-=A, mn[x^1]-=A, mn[x>>1]+=A;
        A=max(mx[x],mx[x^1]),
        mx[x]-=A, mx[x^1]-=A, mx[x>>1]+=A;
void update_part(int s,int t,int v){
    int A=0,lc=0,rc=0,len=1;
    for(s+=m-1,t+=m+1;s^t^1;s>>=1,t>>=1,len<<=1){</pre>
        if(s&1^1) add[s^1]+=v,lc+=len, mn[s^1]+=v,mx[s
             ^1]+=v;
        if(t&1)
                   add[t^1]+=v,rc+=len, mn[t^1]+=v,mx[t]
             ^1]+=v;
        sum[s>>1]+=v*lc, sum[t>>1]+=v*rc;
        A=min(mn[s],mn[s^1]),mn[s]-=A,mn[s^1]-=A,mn[s
             >>1]+=A.
        A=min(mn[t],mn[t^1]),mn[t]-=A,mn[t^1]-=A,mn[t
            >>1]+=A;
        A=max(mx[s],mx[s^1]),mx[s]-=A,mx[s^1]-=A,mx[s
             >>1]+=A,
        A=\max(\max[t],\max[t^1]),\max[t]-=A,\max[t^1]-=A,\max[t
             >>1]+=A;
    for(lc+=rc;s;s>>=1){
        sum[s>>1]+=v*lc;
        A=min(mn[s],mn[s^1]),mn[s]-=A,mn[s^1]-=A,mn[s
             >>1]+=A,
        A=max(mx[s],mx[s^1]),mx[s]-=A,mx[s^1]-=A,mx[s
             >>1]+=A;
    }
int query_node(int x,int ans=0){
    for(x+=m;x;x>>=1) ans+=mn[x]; return ans;
```

```
int query_sum(int s,int t){
    int lc=0,rc=0,len=1,ans=0;
    for(s+=m-1,t+=m+1;s^t^1;s>>=1,t>>=1,len<<=1){</pre>
        if(s&1^1) ans+=sum[s^1]+len*add[s^1],lc+=len;
        if(t&1) ans+=sum[t^1]+len*add[t^1],rc+=len;
        if(add[s>>1]) ans+=add[s>>1]*lc;
        if(add[t>>1]) ans+=add[t>>1]*rc;
    for(lc+=rc,s>>=1;s;s>>=1) if(add[s]) ans+=add[s]*lc
    return ans;
int query_min(int s,int t,int L=0,int R=0,int ans=0){
    if(s==t) return query_node(s);
    for(s+=m,t+=m;s^t^1;s>>=1,t>>=1){
        L+=mn[s],R+=mn[t];
        if(s&1^1) L=min(L,mn[s^1]);
        if(t&1) R=min(R,mn[t^1]);
    for(ans=min(L,R),s>>=1;s;s>>=1) ans+=mn[s];
    return ans;
int query_max(int s,int t,int L=0,int R=0,int ans=0){
    if(s==t) return query_node(s);
    for(s+=m,t+=m;s^t^1;s>>=1,t>>=1){
        L+=mx[s],R+=mx[t];
        if(s&1^1) L=max(L,mx[s^1]);
        if(t&1) R=max(R,mx[t^1]);
    for(ans=max(L,R),s>>=1;s;s>>=1) ans+=mx[s];
    return ans;
|}
```

### 7.5 Lazy Tag

## 7.6 BIT 樹狀樹組

```
class Bitree {
public:
   /* bit 一定是 1 indexed */
    vector<int> data;
    Bitree(const vector<int> &nums) {
        data.resize(nums.size() + 1, 0);
        for(int i = 0; i < nums.size(); i++ ) {</pre>
            update(i, nums[i]);
    void update(int x, int val) {
        x++; /*變成 1 indexed*/
        for(; x < data.size(); x += lowbit(x)) {</pre>
            data[x] += val;
    int query(int x) {
        x++; /*變成 1 indexed*/
        int result = 0;
        for(; x > 0; x -= lowbit(x)) {
```

```
result += data[x];
}
return result;
}
static int lowbit(int x) {
    return x & (-x);
}
};
```

### 7.7 並查集 union find

```
struct DisjointSet {
    vector<int> parent, sz; // parent[i] = 父節點, sz[
        i] = 集合大小
    void init(int n) {
        parent.resize(n + 1);
        sz.assign(n + 1, 1);
for (int i = 0; i <= n; i++) {</pre>
            parent[i] = i;
        }
    int find(int x) {
        if (parent[x] != x) {
            parent[x] = find(parent[x]); // 路徑壓縮
        return parent[x];
    bool unite(int x, int y) {
        x = find(x);
        y = find(y);
        if (x == y) return false; // 已在同一集合
        // 啟發式合併:小的掛到大的下面
        if (sz[x] < sz[y]) swap(x, y);
        parent[y] = x;
        sz[x] += sz[y];
        return true;
    }
    bool same(int x, int y) {
        return find(x) == find(y);
    }
};
```

# 8 String

### 8.1 KMP

```
template<typename T>
void build_KMP(int n, T *s, int *f){ // 1 base
  f[0]=-1, f[1]=0;
  for (int i=2; i<=n; i++){</pre>
    int w = f[i-1];
    while (w \ge 0 \& s[w+1]! = s[i])w = f[w];
    f[i]=w+1;
  }
}
template < typename T>
int KMP(int n, T *a, int m, T *b){
  build_KMP(m,b,f);
  int ans=0;
  for (int i=1, w=0; i<=n; i++){</pre>
    while ( w>=0 && b[w+1]!=a[i] )w = f[w];
    w++;
    if (w==m){
      ans++;
      w=f[w];
    }
  return ans;
```

### 8.2 smallest rotation

```
| string mcp(string s){
    int n = s.length();
    s += s;
    int i=0, j=1;
    while (i<n && j<n){
        int k = 0;
        while (k < n && s[i+k] == s[j+k]) k++;
        if (s[i+k] <= s[j+k]) j += k+1;
        else i += k+1;
        if (i == j) j++;
    }
    int ans = i < n ? i : j;
    return s.substr(ans, n);
}
/*

Booth 演算法

用於尋找一個字串的字典序最小的循環旋轉
*/
Contact GitHub API Training Shop Blog About
```

### 8.3 Suffix Array

```
/*he[i]保存了在後綴數組中相鄰兩個後綴的最長公共前綴長度
 *sa[i]表示的是字典序排名為i的後綴是誰(字典序越小的排
      名越靠前)
 *rk[i]表示的是後綴我所對應的排名是多少 */
const int MAX = 1020304;
int ct[MAX], he[MAX], rk[MAX];
int sa[MAX], tsa[MAX], tp[MAX][2];
void suffix_array(char *ip){
  int len = strlen(ip);
  int alp = 256;
  memset(ct, 0, sizeof(ct));
  for(int i=0;i<len;i++) ct[ip[i]+1]++;</pre>
  for(int i=1;i<alp;i++) ct[i]+=ct[i-1];</pre>
  for(int i=0;i<len;i++) rk[i]=ct[ip[i]];</pre>
  for(int i=1;i<len;i*=2){</pre>
    for(int j=0;j<len;j++){</pre>
       if(j+i>=len) tp[j][1]=0;
       else tp[j][1]=rk[j+i]+1;
       tp[j][0]=rk[j];
    memset(ct, 0, sizeof(ct));
    for(int j=0;j<len;j++) ct[tp[j][1]+1]++;</pre>
    for(int j=1;j<len+2;j++) ct[j]+=ct[j-1];</pre>
    for(int j=0;j<len;j++) tsa[ct[tp[j][1]]++]=j;</pre>
    memset(ct, 0, sizeof(ct));
for(int j=0;j<len;j++) ct[tp[j][0]+1]++;</pre>
    for(int j=1;j<len+1;j++) ct[j]+=ct[j-1];</pre>
    for(int j=0;j<len;j++)</pre>
       sa[ct[tp[tsa[j]][0]]++]=tsa[j];
    rk[sa[0]]=0;
    for(int j=1;j<len;j++){</pre>
       if( tp[sa[j]][0] == tp[sa[j-1]][0] &&
         tp[sa[j]][1] == tp[sa[j-1]][1] )
         rk[sa[j]] = rk[sa[j-1]];
       else
         rk[sa[j]] = j;
    }
  for(int i=0,h=0;i<len;i++){</pre>
    if(rk[i]==0) h=0;
    else{
       int j=sa[rk[i]-1];
       h=max(0,h-1);
       for(;ip[i+h]==ip[j+h];h++);
    he[rk[i]]=h;
  }
| }
```

### 8.4 Z-value

```
z[0] = 0;
for ( int bst = 0, i = 1; i < len ; i++ ) {</pre>
 if ( z[bst] + bst <= i ) z[i] = 0;</pre>
  else z[i] = min(z[i - bst], z[bst] + bst - i);
  while ( str[i + z[i]] == str[z[i]] ) z[i]++;
  if ( i + z[i] > bst + z[bst] ) bst = i;
// 回文版
void Zpal(const char *s, int len, int *z) {
    // Only odd palindrome len is considered
    // z[i] means that the longest odd palindrom
        centered at
    // i is [i-z[i] .. i+z[i]]
    z[0] = 0;
    for (int b=0, i=1; i<len; i++) {</pre>
        if (z[b] + b >= i) z[i] = min(z[2*b-i], b+z[b]-
            i):
        else z[i] = 0;
        while (i+z[i]+1 < len and i-z[i]-1 >= 0 and
                s[i+z[i]+1] == s[i-z[i]-1]) z[i] ++;
        if(z[i] + i > z[b] + b) b = i;
    }
}
```

## 8.5 旋轉哈希

```
typedef unsigned __int128 ulll;
ull1 power(ull1 a, ull1 n, ull1 m) {
    ulll re = 1:
    while (n > 0) {
        if (n & 1) re = re * a % m;
        a = a * a % m;
        n >>= 1;
    return re;
}
ulll inv(ulll a, ulll m) {
    return power(a, m - 2, m);
struct Rh {
    const ulll p, mod;
    vector<ulll> ps{1};
    Rh(ulll p, ulll mod) : p(p), mod(mod) {}
    vector<ull>> build(const string &s) {
        vector<ulll> h(s.size() + 1);
        h[0] = 0;
        ps.resize(s.size() + 1);
        for (int i = 0; i < s.size(); i++) {</pre>
            ps[i + 1] = ps[i] * p % mod;
            h[i + 1] = (h[i] + s[i] * ps[i + 1] % mod)
                % mod;
        return h;
    ulll subhash(const vector<ulll> &h, int 1, int r) {
        // [l, r] 指原字串
        return ((h[r + 1] - h[1]) * inv(ps[1], mod)) %
            mod:
};
constexpr uint64_t mod = (1ull<<61) - 1;</pre>
uint64_t modmul(uint64_t a, uint64_t b){
  uint64_t 11 = (uint32_t)a, h1 = a>>32, 12 = (uint32_t)a
      )b, h2 = b >> 32;
  uint64_t l = 11*12, m = 11*h2 + 12*h1, h = h1*h2;
  uint64_t ret = (1&mod) + (1>>61) + (h << 3) + (m >>
      29) + (m << 35 >> 3) + 1;
  ret = (ret & mod) + (ret>>61);
  ret = (ret & mod) + (ret>>61);
  return ret-1;
}
```

### 8.6 後綴自動機

```
struct state {
    int len{}, link{};
     array<int, 26> next{};
};
 struct SAM {
     int sz{}, last{};
     vector<state> st;
     SAM(int maxlen) : st(maxlen * 2) {
         st[0].len = 0;
         st[0].link = -1;
         sz++;
         last = 0;
     void insert(char c) {
         insert_impl(c - 'a');
     void insert_impl(char c) {
         int cur = sz++;
         st[cur].len = st[last].len + 1;
         int p = last;
         while(p != -1 && !st[p].next[c]) {
             st[p].next[c] = cur;
             p = st[p].link;
         if(p == -1) {
             st[cur].link = 0;
         else {
             int q = st[p].next[c];
             if(st[p].len + 1 == st[q].len) {
                 st[cur].link = q;
             else {
                 int clone = sz++;
                 st[clone].len = st[p].len + 1;
                 st[clone].next = st[q].next;
                 st[clone].link = st[q].link;
                 while(p != -1 && st[p].next[c] == q) {
                     st[p].next[c] = clone;
                     p = st[p].link;
                 st[q].link = st[cur].link = clone;
             }
         last = cur;
     }
|};
```

## 9 Others

### 9.1 矩陣樹定理

```
新的方法介紹
```

下面我們介紹一個新的方法一Matrix-Tree定理(Kirchhoff矩陣 - 樹定理)。

Matrix-Tree定理是解決生成樹數問題最有力的武器之一。它首 先於1847年被Kirchhoff證明。在介紹定理之前,我們先先 明確幾個概念:

- 1.G的度數矩陣D[G]是一個n\*n的矩陣,並且滿足:當i≠j時,dij =0;當i=j時,dij等於vi的度數。
- 2.G的鄰接矩陣A[G]也是一個n\*n的矩陣, 且滿足:若vi、vj之間有邊直接相連,則aij=1,否則為0。
- 我們定義G的Kirchhoff矩陣(也稱為拉普拉斯算子)C[G]為C[G]=D[G]-A[G],
- 則Matrix-Tree定理可以描述為:G的所有不同的生成樹的個數等於其Kirchhoff矩陣C[G]任何一個n-1階主子式的行列式的絕對值。
- 所謂n-1階主子式,就是對於r(1≤r≤n),將C[G]的第r行、第r列 同時去掉後所得到的新矩陣,以Cr[G]表示。

```
生成樹計數
演算法步驟:
1、 建構拉普拉斯矩陣
Matrix[i][j] =
degree(i) , i==j
-1, i-j有邊
0,其他情況
2、 去掉第r行,第r列(r任意)
3、 計算矩陣的行列式
#include <stdio.h>
#include <string.h>
#include <algorithm>
#include <iostream>
#include <math.h>
using namespace std;
const double eps = 1e-8;
const int MAXN = 110;
int sgn(double x)
    if(fabs(x) < eps)return 0;</pre>
    if(x < 0) return -1;
    else return 1;
double b[MAXN][MAXN];
double det(double a[][MAXN],int n)
    int i, j, k, sign = 0;
    double ret = 1;
    for(i = 0;i < n;i++)</pre>
    for(j = 0;j < n;j++) b[i][j] = a[i][j];</pre>
    for(i = 0;i < n;i++)</pre>
        if(sgn(b[i][i]) == 0)
            for(j = i + 1; j < n; j++)
            if(sgn(b[j][i]) != 0) break;
            if(j == n)return 0;
            for (k = i; k < n; k++) swap (b[i][k], b[j][k]);
            sign++;
        ret *= b[i][i];
        for(k = i + 1; k < n; k++) b[i][k]/=b[i][i];
        for(j = i+1; j < n; j++)
        for(k = i+1; k < n; k++) b[j][k] -= b[j][i]*b[i][
            k];
    if(sign & 1)ret = -ret;
    return ret;
double a[MAXN][MAXN];
int g[MAXN][MAXN];
int main()
    int T;
    int n,m;
    int u,v;
    scanf("%d",&T);
    while(T--)
        scanf("%d%d",&n,&m);
        memset(g,0,sizeof(g));
        while(m--)
            scanf("%d%d",&u,&v);
            g[u][v] = g[v][u] = 1;
        memset(a,0,sizeof(a));
        for(int i = 0;i < n;i++)</pre>
        for(int j = 0; j < n; j++)</pre>
        if(i != j && g[i][j])
        {
            a[i][i]++;
            a[i][j] = -1;
        double ans = det(a,n-1);
        printf("%.0lf\n",ans);
```

return 0;

## 9.2 1D/1D dp 優化

| }

```
#include < bits / stdc++.h>
int t, n, L;
int p;
char s[MAXN][35];
11 sum[MAXN] = \{0\};
long double dp[MAXN] = {0};
int prevd[MAXN] = {0};
long double pw(long double a, int n) {
    if ( n == 1 ) return a;
    long double b = pw(a, n/2);
    if ( n & 1 ) return b*b*a;
    else return b*b;
long double f(int i, int j) {
     cout << (sum[i] - sum[j]+i-j-1-L) << endl;</pre>
    return pw(abs(sum[i] - sum[j]+i-j-1-L), p) + dp[j];
struct INV {
    int L, R, pos;
INV stk[MAXN*10];
int top = 1, bot = 1;
void update(int i) {
    while ( top > bot && i < stk[top].L && f(stk[top].L</pre>
        , i) < f(stk[top].L, stk[top].pos) ) {</pre>
        stk[top - 1].R = stk[top].R;
        top--;
    int lo = stk[top].L, hi = stk[top].R, mid, pos =
        stk[top].pos;
    //if (i >= lo) lo = i + 1;
    while ( lo != hi ) {
        mid = lo + (hi - lo) / 2;
        if ( f(mid, i) < f(mid, pos) ) hi = mid;</pre>
        else lo = mid + 1;
    if ( hi < stk[top].R ) {</pre>
        stk[top + 1] = (INV) { hi, stk[top].R, i };
        stk[top++].R = hi;
    }
}
int main() {
    cin >> t:
    while ( t-- ) {
        cin >> n >> L >> p;
        dp[0] = sum[0] = 0;
        for ( int i = 1 ; i <= n ; i++ ) {</pre>
            cin >> s[i];
            sum[i] = sum[i-1] + strlen(s[i]);
            dp[i] = numeric_limits<long double>::max();
        stk[top] = (INV) \{1, n + 1, 0\};
        for ( int i = 1 ; i <= n ; i++ ) {</pre>
            if ( i >= stk[bot].R ) bot++;
            dp[i] = f(i, stk[bot].pos);
            update(i);
//
              cout << (ll) f(i, stk[bot].pos) << endl;</pre>
        if ( dp[n] > 1e18 ) {
            cout << "Too hard to arrange" << endl;
        } else {
            vector<PI> as;
            cout << (11)dp[n] << endl;</pre>
    return 0;
```

### 9.3 Theorm - DP optimization

```
Monotonicity & 1D/1D DP & 2D/1D DP
Definition xD/vD
1D/1D \ DP[j] = min(0 \le i < j) \ \{ \ DP[i] + w(i, j) \ \}; \ DP[0] = k
2D/1D DP[i][j] = min(i < k \le j) \{ DP[i][k - 1] + DP[k][j] \}
    + w(i, j); DP[i][i] = 0
Monotonicity
a | w(a, c) w(a, d)
b \mid w(b, c) w(b, d)
Monge Condition
Concave(凹四邊形不等式): w(a, c) + w(b, d) >= w(a, d) +
     w(b, c)
Convex (凸四邊形不等式): w(a, c) + w(b, d) <= w(a, d) +
     w(b, c)
Totally Monotone
Concave(凹單調): w(a, c) <= w(b, d) ----> w(a, d) <= w
    (b, c)
Convex (凸單調): w(a, c) >= w(b, d) ----> w(a, d) >= w
   (b, c)
1D/1D DP O(n^2) -> O(nlgn)
**CONSIDER THE TRANSITION POINT**
Solve 1D/1D Concave by Stack
Solve 1D/1D Convex by Deque
2D/1D Convex DP (Totally Monotone) O(n^3) \rightarrow O(n^2)
h(i, j - 1) \le h(i, j) \le h(i + 1, j)
```

### 9.4 Stable Marriage

```
// normal stable marriage problem
// input:
//3
//Albert Laura Nancy Marcy
//Brad Marcy Nancy Laura
//Chuck Laura Marcy Nancy
//Laura Chuck Albert Brad
//Marcy Albert Chuck Brad
//Nancy Brad Albert Chuck
#include < bits / stdc++.h>
using namespace std;
const int MAXN = 505;
int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id
int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank
int current[MAXN]; // current[boy_id] = rank; boy_id
    will pursue current[boy_id] girl.
int girl_current[MAXN]; // girl[girl_id] = boy_id;
void initialize() {
  for ( int i = 0 ; i < n ; i++ ) {</pre>
    current[i] = 0;
    girl_current[i] = n;
    order[i][n] = n;
map<string, int> male, female;
string bname[MAXN], gname[MAXN];
int fit = 0;
void stable_marriage() {
  queue<int> que;
  for ( int i = 0 ; i < n ; i++ ) que.push(i);</pre>
  while ( !que.empty() ) {
    int boy_id = que.front();
    que.pop();
    int girl_id = favor[boy_id][current[boy_id]];
```

```
current[boy_id] ++;
    if ( order[girl_id][boy_id] < order[girl_id][</pre>
         girl_current[girl_id]] ) {
      if ( girl_current[girl_id] < n ) que.push(</pre>
           girl_current[girl_id]); // if not the first
      girl_current[girl_id] = boy_id;
    } else {
      que.push(boy_id);
 }
int main() {
 cin >> n;
  for ( int i = 0 ; i < n; i++ ) {</pre>
    string p, t;
    cin >> p;
    male[p] = i;
    bname[i] = p;
    for ( int j = 0 ; j < n ; j++ ) {
      cin >> t;
      if ( !female.count(t) ) {
        gname[fit] = t;
        female[t] = fit++;
      favor[i][j] = female[t];
 }
  for ( int i = 0 ; i < n ; i++ ) {</pre>
    string p, t;
    cin >> p;
    for ( int j = 0 ; j < n ; j++ ) {</pre>
      cin >> t;
      order[female[p]][male[t]] = j;
  initialize();
 stable_marriage();
 for ( int i = 0 ; i < n ; i++ ) {
  cout << bname[i] << " " << gname[favor[i][current[i]</pre>
         ] - 1]] << endl;
```

#### 9.5 莫隊

```
/* nums 長度 N ;; query 長度為 M */
/* O(N * sqrt(M))*/
struct Query {
   int l, r, id;
void add(int pos) {
   /*更新狀態*/
   /*將pos所在的移入集合*/
void del(int pos) {
   /*更新狀態*/
   /*將pos所在的移出集合*/
sort(query.begin(), query.end(), [bsz](const Query &a,
    const Query &b){
   if(a.1 / bsz != b.1 / bsz) {
       return a.l < b.l;</pre>
    return (a.1 / bsz) & 1 ? a.r < b.r : a.r > b.r;
|});
```

```
int l = 1;
int r = 0;

vector<pair<int, int>> res(m);

for(int i = 0; i < query.size(); i++ ) {
    auto &q = query[i];
    /*順序不能換*/
    while (l > q.l) add(--l);
    while (r < q.r) add(++r);
    while (l < q.l) del(l++);
    while (r > q.r) del(r--);
    res[q.id] = /* 根據當前狀態求解 */
}
```

### 9.6 矩陣乘法

## 9.7 c++ 小抄

```
//pbds tree
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
tree<int, null_type, less<int>, rb_tree_tag,
   tree_order_statistics_node_update> tr;
tree::find_by_order(k) // O(LogN) 取得第k大的元素
tree::order_of_key(ele) // O(logN) 得到ele是tree中第幾
    大(有幾個元素小於ele)
//pbds pair priority_queue
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using namespace __gnu_pbds;
priority_queue<int, less<int>, pairing_heap_tag> pq;
auto it = pq.push(x);
// type of it = priority_queue<int, less<int>,
   pairing_heap_tag>::point_iterator
pq.pop();
pq.top();
pq.join(b);
pq.empty();
pq.size();
pq.modify(it,6);
                   // O(LogN)
pq.erase(it);
//builtin functions
__builtin_popcount(x); // 1的個數
__builtin_popcountll(x); // for long long
                       // 前導0的個數
__builtin_clz(x);
                        // 後導0的個數
__builtin_ctz(x);
                        // 奇偶性
__builtin_parity(x);
//溢位檢查
ret = __builtin_add_overflow(a, b, &res) // if ret = 1
    a+b 溢位
ret = __builtin_sub_overflow(a, b, &res) // if ret = 1
   a-b 溢位
```

```
_builtin_mul_overflow(a, b, &res) // if ret = 1
   a*b 溢位
      ret =
    溢位 第三個參數是判斷的類型
//vector SIMD
typedef int v4si __attribute__ ((vector_size(4 * sizeof
   (int))));
//大質數表
{1000000007, 1000000009, 1000000021, 1000000033,
    1000000087, 1000000093, 1000000097, 1000000123,
    1000000321};
//mt19937
#include <random>
#include <chrono>
int getRendom(int 1, int r) {
   static auto seed = std::chrono::system_clock::now()
       .time_since_epoch().count();
   static std::mt19937 gen(seed);
   std::uniform_int_distribution<int> dis(l, r);
   return dis(gen);
//sorted vector 去重
vec.erase(unique(vec.begin(), vec.end()), vec.end());
//std::valarray
valarray<int> a(初始值,數量);
valarray<int> a(10);
valarray<int> b(10);
valarray<int> c = a + b;
valarray<int> d = a * b;
valarray<int> e = a + 10;
valarray<int> f = a * 10;
valarray<int> g = a.cshift(1); //循環左移
valarray<bool> equal = a == b;
int sum = a.sum();
int max = a.max();
int min = a.min();
std::valarray<int> g = a.apply([](int x) { return x * x
   ; });
//regex ***very slow***
#include <regex>
using namespace std;
bool res = regex_match("abc", regex("a.c"));
bool res = regex_match("abc", regex("A.c", regex::icase
   )); //忽略大小寫
```

# 9.8 python 小抄

```
#!/usr/bin/env python3
# 帕斯卡三角形
n = 10
dp = [ [1 for j in range(n)] for i in range(n) ]
for i in range(1,n):
    for j in range(1,n):
        dp[i][j] = dp[i][j-1] + dp[i-1][j]
for i in range(n):
    print(
              '.join( '{:5d}'.format(x) for x in dp[i] )
# EOF1
while True:
    try:
        n, m = map(int, input().split())
    except:
        break
# EOF2
import sys
for s in sys.stdin:
    print(eval(s.replace("/", "//")))
```

```
# input a sequence of number
a = [ int(x) for x in input().split() ]
a.sort()
        ,
''.join( str(x)+' ' for x in a ) )
print(
ncase = int( input() )
for _ in range(ncase):
    n, m = [int(x) for x in input().split()]
a, b = "$"+input(), "$"+input()
    dp = [ [int(0) for j in range(m+1)] for i in range(
        n+1) ]
    for i in range(1,n+1):
        for j in range(1,m+1):
             dp[i][j] = max(dp[i-1][j],dp[i][j-1])
             if a[i]==b[j]:
                 dp[i][j] = max(dp[i][j],dp[i-1][j-1]+1)
    for i in range(1,n+1):
        print(dp[i][1:])
    print('a=\{:s\}, b=\{:s\}, |LCS(a,b)|=\{:d\}'.format(a,b)\}
        [1:],b[1:],dp[n][m]))
# list, dict, string
a = [1, 3, 4, 65, 65]
b = list.copy() # b = [1, 3, 4, 65], list a 跟 llst b
    互相獨立
cnt = list.count(65) # cnt == 2
loc = list.index(65) # loc == 3, find the leftmost
    element, {\tt if} \ {\tt not} \ {\tt found} \ {\tt then} \ {\tt return} \ {\tt ERROR}
list.sort(reverse = True|False, key = none|lambda x:x
    [1]) # list.sort has side effect but no reture
    value
# stack
                 # C++
stack = [3,4,5]
stack.append(6) # push()
stack.pop()
               # pop()
stack[-1]
                 # top()
len(stack)
                # size() 0(1)
                 # C++
from collections import deque
queue = deque([3,4,5])
queue.append(6) # push()
queue.popleft() # pop()
queue[0]
                 # front()
                 # size() 0(1)
len(queue)
```

### 10 Persistence