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Basic

1.1 compile

```
# preset before coding
echo "cd ~/Desktop" >> ~/.bashrc
gedit -> preference -> tab width: 4
# Editor
gedit a.cpp
# Compile
g++ a.cpp -std=c++11
**All file will be compiled to a.out unless you use -o(
    not recommanded, just use a.out)**
```

```
# Run
./a.out
# Run with file input
./a.out < input.txt
# Run with file input and output
./a.out < input.txt > output.txt
# Python Run
python3 a.py < input.txt > output.txt
# Copy Paste In Ubuntu
 copy: ctrl+insert
* paste: shift+insert
# 比對文件相同
sdiff a.txt b.txt
```

1.2 default code

2

8

9

```
#include <bits/stdc++.h>
using namespace std;
#define int long long
typedef pair<int,int> pii;
#ifdef ONLINE JUDGE
#define cerr if(false) cerr
#endif
int32_t main(){
#ifndef ONLINE_JUDGE
  //freopen("input.txt","r",stdin);
freopen("output.txt","w",stdout);
freopen("debug.txt","w",stdcerr);
   ios_base::sync_with_stdio(0);
   cin.tie(false);
#endif
}
```

1.3 debug list

```
記得測試 python 的內建函數庫有哪些
bits/std++.h 跟 global variable y1 衝突,不能用
模板要記得 init
priority_queue 要清空
事先將把邊界測資加入測試
邊界條件 (過程溢位, 題目數據範圍), 會不會爆 long long
是否讀錯題目, 想不到時可以自己讀一次題目
比較容易有問題的地方換人寫
注意公式有沒有推錯或抄錯
精度誤差 sqrt(大大的東西) + EPS
喇分 random_suffle 隨機演算法
```

Dark Code

2.1 IO optimization

```
*if output to much, consider put all output in array
     first, then output the array.
getchar() -> getchar_unlocked()
fread() -> fread_unlocked()
inline char readchar() {
  const int S = 1<<20; // buffer size
static char buf[S], *p = buf, *q = buf;</pre>
  if(p == q && (q = (p=buf)+fread(buf,1,S,stdin)) ==
       buf) return EOF;
  return *p++;
inline int nxtint() {
```

3 Geometry

3.1 2D point

```
typedef double Double:
struct Point {
  Double x,y;
  bool operator < (const Point &b)const{</pre>
    //return\ tie(x,y) < tie(b.x,b.y);
    return atan2(y,x) < atan2(b.y,b.x);</pre>
  Point operator + (const Point &b)const{
    return (Point){x+b.x,y+b.y};
  Point operator - (const Point &b)const{
    return (Point){x-b.x,y-b.y};
  Point operator * (const Double &d)const{
    return Point(d*x,d*y);
  Double operator * (const Point &b)const{
    return x*b.x + y*b.y;
  Double operator % (const Point &b)const{
    return x*b.y - y*b.x;
  friend Double abs2(const Point &p){
    return p.x*p.x + p.y*p.y;
  friend Double abs(const Point &p){
    return sqrt( abs2(p) );
};
typedef Point Vector;
struct Line{
  Point P; Vector v;
  bool operator < (const Line &b)const{</pre>
    return atan2(v.y,v.x) < atan2(b.v.y,b.v.x);</pre>
};
```

3.2 兩線段交點

```
const Double EPS = 1e-9;
Double cross(Point p1, Point p2, Point q1) {
    return (p2.x - p1.x) * (q1.y - p1.y) - (p2.y - p1.y)
        * (q1.x - p1.x);
}

Point interPnt(Point p1, Point p2, Point q1, Point q2,
        bool &res){
    Double f1 = cross(p2, q1, p1);
    Double f2 = -cross(p2, q2, p1);
    Double f = (f1 + f2);

if(fabs(f) < EPS) {
    res = false;
    return {};
}</pre>
```

```
res = true;
return (f2 / f) * q1 + (f1 / f) * q2;
}
```

3.3 兩圓交點

3.4 Convex Hull

```
#include "2Dpoint.cpp"
// return H, The first will occured TWICE in vector H!
void ConvexHull(vector<Point> &P, vector<Point> &H){
    int n = P.size(), m=0;
    sort(P.begin(),P.end());
    H.clear();
    for (int i=0; i<n; i++){</pre>
        while (m>=2 && (P[i]-H[m-2]) % (H[m-1]-H[m-2])
             <0)H.pop_back(), m--;
        H.push_back(P[i]), m++;
    for (int i=n-2; i>=0; i--){
        while (m>=2 && (P[i]-H[m-2]) % (H[m-1]-H[m-2])
             <0)H.pop_back(), m--;
        H.push_back(P[i]), m++;
    }
}
```

4 Flow

4.1 Dinic

```
(a) Bounded Maxflow Construction:
1. add two node ss, tt
2. add_edge(ss, tt, INF)
3. for each edge u \rightarrow v with capacity [1, r]:
        add_edge(u, tt, 1)
        add_edge(ss, v, 1)
add_edge(u, v, r-1)
4. see (b), check if it is possible.
5. answer is maxflow(ss, tt) + maxflow(s, t)
(b) Bounded Possible Flow:
1. same construction method as (a)
run maxflow(ss, tt)
3. for every edge connected with ss or tt:
        rule: check if their rest flow is exactly 0
4. answer is possible if every edge do satisfy the rule
5. otherwise, it is NOT possible.
(c) Bounded Minimum Flow:
1. same construction method as (a)
answer is maxflow(ss, tt)
(d) Bounded Minimum Cost Flow:
 the concept is somewhat like bounded possible flow.
1. same construction method as (a)
```

```
2. answer is maxflow(ss, tt) + (\sum 1 * cost for every
    edge)
                    _____
(e) Minimum Cut:

 run maxflow(s, t)

2. run cut(s)
3. ss[i] = 1: node i is at the same side with s.
const long long INF = 1LL<<60;</pre>
struct Dinic { //O(VVE), with minimum cut
    static const int MAXN = 5003;
    struct Edge{
        int u, v;
        long long cap, rest;
    int n, m, s, t, d[MAXN], cur[MAXN];
    vector<Edge> edges;
    vector<int> G[MAXN];
    void init(){
        edges.clear();
        for ( int i = 0 ; i < MAXN ; i++ ) G[i].clear()</pre>
    }
    // min cut start
    bool side[MAXN];
    void cut(int u) {
        side[u] = 1;
        for ( int i : G[u] ) {
            if ( !side[ edges[i].v ] && edges[i].rest )
                  cut(edges[i].v);
    // min cut end
    void add_edge(int u, int v, long long cap){
  edges.push_back( {u, v, cap, cap} );
        edges.push_back( {v, u, 0, 0LL} );
        m = edges.size();
        G[u].push_back(m-2);
        G[v].push_back(m-1);
    }
    bool bfs(){
        memset(d, -1, sizeof(d));
        queue<int> que;
        que.push(s); d[s]=0;
        while (!que.empty()){
            int u = que.front(); que.pop();
            for (int ei : G[u]){
                Edge &e = edges[ei];
                if (d[e.v] < 0 && e.rest > 0){
                     d[e.v] = d[u] + 1;
                     que.push(e.v);
                }
            }
        return d[t] >= 0;
    long long dfs(int u, long long a){
        if ( u == t || a == 0 ) return a;
        long long flow = 0, f;
        for ( int &i=cur[u]; i < (int)G[u].size(); i++</pre>
            Edge &e = edges[ G[u][i] ];
            if ( d[u] + 1 != d[e.v] ) continue;
            f = dfs(e.v, min(a, e.rest) );
            if (f > 0) {
                 e.rest -= f;
                edges[ G[u][i]^1 ].rest += f;
                flow += f;
                a -= f;
                if ( a == 0 )break;
            }
        return flow;
    }
```

```
long long maxflow(int s, int t){
    this->s = s, this->t = t;
    long long flow = 0, mf;
    while ( bfs() ){
        memset(cur, 0, sizeof(cur));
        while ( (mf = dfs(s, INF)) ) flow += mf;
    }
    return flow;
}
dinic;
```

4.2 min cost flow

```
// long long version
typedef pair<long long, long long> pll;
struct CostFlow {
    static const int MAXN = 350;
     static const long long INF = 1LL<<60;
     struct Edge {
         int to, r;
         long long rest, c;
    int n, pre[MAXN], preL[MAXN]; bool inq[MAXN];
    long long dis[MAXN], fl, cost;
    vector<Edge> G[MAXN];
    void init() {
         for ( int i = 0 ; i < MAXN ; i++) G[i].clear();</pre>
    void add_edge(int u, int v, long long rest, long
         long c) {
         G[u].push_back({v, (int)G[v].size() , rest, c
         G[v].push_back({u, (int)G[u].size()-1, 0, -c});
    pll flow(int s, int t) {
         fl = cost = 0;
         while (true) {
             fill(dis, dis+MAXN, INF);
             fill(inq, inq+MAXN, 0);
             dis[s] = 0;
             queue<int> que;
             que.push(s);
             while ( !que.empty() ) {
                 int u = que.front(); que.pop();
                 inq[u] = 0;
                  for ( int i = 0 ; i < (int)G[u].size()</pre>
                      ; i++) {
                      int v = G[u][i].to;
                      long long w = G[u][i].c;
                      if ( G[u][i].rest > 0 && dis[v] >
                           dis[u] + w) {
                          pre[v] = u; preL[v] = i;
dis[v] = dis[u] + w;
                          if (!inq[v]) {
                               inq[v] = 1;
                               que.push(v);
                          }
                      }
                 }
             if (dis[t] == INF) break;
             long long tf = INF;
             for (int v = t, u, 1 ; v != s ; v = u ) {
    u = pre[v]; 1 = preL[v];
                 tf = min(tf, G[u][1].rest);
             for (int v = t, u, 1; v != s; v = u) {
    u = pre[v]; 1 = preL[v];
                 G[u][1].rest -= tf;
                 G[v][G[u][1].r].rest += tf;
             cost += tf * dis[t];
             fl += tf;
         return {fl, cost};
} flow;
```

5 Mathmatics

5.1 ax+by=gcd(a,b)

```
typedef pair<int, int> pii;

pii exgcd(int a, int b){
   if(b == 0) return make_pair(1, 0);
   else{
     int p = a / b;
     pii q = exgcd(b, a % b);
     int aa = q.second, bb = q.first - q.second * p;
     if(aa < 0) aa += b, bb -= a;
     return make_pair(aa, bb);
   }
}</pre>
```

5.2 BigInt

```
struct Bigint{
  static const int LEN = 60;
  static const int BIGMOD = 10000;
  int s;
 int v1, v[LEN];
  // vector<int> v;
 Bigint() : s(1) \{ vl = 0; \}
  Bigint(long long a) {
    s = 1; v1 = 0;
    if (a < 0) { s = -1; a = -a; }
    while (a) {
      push_back(a % BIGMOD);
      a /= BIGMOD;
  Bigint(string str) {
    s = 1; v1 = 0;
    int stPos = 0, num = 0;
    if (!str.empty() && str[0] == '-') {
      stPos = 1;
      s = -1;
    for (int i=SZ(str)-1, q=1; i>=stPos; i--) {
  num += (str[i] - '0') * q;
      if ((q *= 10) >= BIGMOD) {
        push_back(num);
        num = 0; q = 1;
    if (num) push_back(num);
  int len() const { return vl; /* return SZ(v); */ }
  bool empty() const { return len() == 0; }
void push_back(int x) { v[vl++] = x; /* v.PB(x); */ }
  void pop_back() { v1--; /* v.pop_back(); */ }
  int back() const { return v[vl-1]; /* return v.back()
  void n() { while (!empty() && !back()) pop_back(); }
  void resize(int nl) {
    vl = nl; fill(v, v+vl, 0);
          v.resize(nl); // fill(ALL(v), 0);
  }
  void print() const {
    if (empty()) { putchar('0'); return; }
    if (s == -1) putchar('-');
printf("%d", back());
    for (int i=len()-2; i>=0; i--) printf("%.4d",v[i]);
  friend std::ostream& operator << (std::ostream& out,</pre>
      const Bigint &a) {
    if (a.empty()) { out << "0"; return out; }</pre>
    if (a.s == -1) out << "-";</pre>
    out << a.back();
    for (int i=a.len()-2; i>=0; i--) {
      char str[10];
      snprintf(str, 5, "%.4d", a.v[i]);
      out << str;
    return out;
```

```
int cp3(const Bigint &b)const {
  if (s != b.s) return s > b.s ? 1 : -1;
  if (s == -1) return -(-*this).cp3(-b);
  if (len() != b.len()) return len()>b.len()?1:-1;
  for (int i=len()-1; i>=0; i--)
   if (v[i]!=b.v[i]) return v[i]>b.v[i]?1:-1;
  return 0;
bool operator < (const Bigint &b)const{return cp3(b)</pre>
bool operator <= (const Bigint &b)const{return cp3(b)</pre>
    <=0;}
bool operator >= (const Bigint &b)const{return cp3(b)
    >=0;}
bool operator == (const Bigint &b)const{return cp3(b)
    ==0;}
bool operator != (const Bigint &b)const{return cp3(b)
    !=0:}
bool operator > (const Bigint &b)const{return cp3(b)
    ==1;}
Bigint operator - () const {
  Bigint r = (*this);
  r.s = -r.s;
  return r:
Bigint operator + (const Bigint &b) const {
  if (s == -1) return -(-(*this)+(-b));
  if (b.s == -1) return (*this)-(-b);
  Bigint r;
  int nl = max(len(), b.len());
  r.resize(nl + 1);
  for (int i=0; i<nl; i++) {</pre>
    if (i < len()) r.v[i] += v[i];</pre>
    if (i < b.len()) r.v[i] += b.v[i];</pre>
    if(r.v[i] >= BIGMOD) {
      r.v[i+1] += r.v[i] / BIGMOD;
      r.v[i] %= BIGMOD;
    }
  }
  r.n();
  return r;
Bigint operator - (const Bigint &b) const {
  if (s == -1) return -(-(*this)-(-b));
  if (b.s == -1) return (*this)+(-b);
  if ((*this) < b) return -(b-(*this));</pre>
  Bigint r;
  r.resize(len());
  for (int i=0; i<len(); i++) {</pre>
    r.v[i] += v[i];
    if (i < b.len()) r.v[i] -= b.v[i];</pre>
    if (r.v[i] < 0) {</pre>
      r.v[i] += BIGMOD;
      r.v[i+1]--;
  r.n();
  return r;
Bigint operator * (const Bigint &b) {
  Bigint r;
  r.resize(len() + b.len() + 1);
  r.s = s * b.s;
  for (int i=0; i<len(); i++) {</pre>
    for (int j=0; j<b.len(); j++) {</pre>
      r.v[i+j] += v[i] * b.v[j];
      if(r.v[i+j] >= BIGMOD) {
        r.v[i+j+1] += r.v[i+j] / BIGMOD;
        r.v[i+j] %= BIGMOD;
      }
    }
  }
  r.n();
  return r;
Bigint operator / (const Bigint &b) {
  Bigint r;
  r.resize(max(1, len()-b.len()+1));
  int oriS = s;
  Bigint b2 = b; // b2 = abs(b)
  s = b2.s = r.s = 1;
```

```
for (int i=r.len()-1; i>=0; i--) {
    int d=0, u=BIGMOD-1;
    while(d<u) {
        int m = (d+u+1)>>1;
        r.v[i] = m;
        if((r*b2) > (*this)) u = m-1;
        else d = m;
    }
    r.v[i] = d;
}
s = oriS;
r.s = s * b.s;
r.n();
return r;
}
Bigint operator % (const Bigint &b) {
    return (*this)-(*this)/b*b;
}
};
```

5.3 GaussElimination

```
// by bcw_codebook
const int MAXN = 300;
const double EPS = 1e-8;
double A[MAXN][MAXN];
void Gauss() {
  for(int i = 0; i < n; i++) {</pre>
    bool ok = 0;
    for(int j = i; j < n; j++) {</pre>
       if(fabs(A[j][i]) > EPS) {
         swap(A[j], A[i]);
         ok = 1;
         break;
      }
    if(!ok) continue;
    double fs = A[i][i];
    for(int j = i+1; j < n; j++) {</pre>
       double r = A[j][i] / fs;
       for(int k = i; k < n; k++) {</pre>
        A[j][k] -= A[i][k] * r;
       }
    }
  }
}
template < class T>
void Gauss(vector<vector<T>> &A) {
  int n = A.size();
  for(int i = 0; i < n; i++) {</pre>
    bool ok = 0;
    for(int j = i; j < n; j++) {
  if(A[j][i] != 0) {</pre>
         swap(A[j], A[i]);
         ok = 1;
         break;
      }
    if(!ok) continue;
    T fs = A[i][i];
    for(int j = i+1; j < n; j++) {</pre>
       T r = A[j][i] / fs;
       for(int k = i; k < n; k++) {</pre>
         A[j][k] -= A[i][k] * r;
       }
    }
  }
}
```

5.4 Inverse

```
int inverse[100000];
void invTable(int b, int p) {
  inverse[1] = 1;
  for( int i = 2; i <= b; i++ ) {
    inverse[i] = (long long)inverse[p%i] * (p-p/i) % p;
  }
}
int inv(int b, int p) {
  return b == 1 ? 1 : ((long long)inv(p % b, p) * (p-p/b) % p);
}</pre>
```

5.5 LinearPrime

```
const int MAXP = 100; //max prime
vector<int> P; // primes
void build_prime(){
    static bitset<MAXP> ok;
    int np=0;
    for (int i=2; i<MAXP; i++){
        if (ok[i]==0)P.push_back(i), np++;
        for (int j=0; j<np && i*P[j]<MAXP; j++){
        ok[ i*P[j] ] = 1;
        if (i%P[j]==0 )break;
    }
}</pre>
```

5.6 Miller Rabin

```
typedef long long LL;
inline LL bin_mul(LL a, LL n,const LL& MOD){
  LL re=0;
  while (n>0){
    if (n&1) re += a;
    a += a; if (a>=MOD) a-=MOD;
    n>>=1:
  }
  return re%MOD;
inline LL bin_pow(LL a, LL n,const LL& MOD){
  LL re=1;
  while (n>0){
    if (n&1) re = bin_mul(re,a,MOD);
    a = bin_mul(a,a,MOD);
    n>>=1;
  }
  return re;
}
bool is_prime(LL n){
  //static LL sprp[3] = { 2LL, 7LL, 61LL};
  static LL sprp[7] = { 2LL, 325LL, 9375LL,
    28178LL, 450775LL, 9780504LL,
    1795265022LL };
  if (n==1 || (n&1)==0 ) return n==2;
  int u=n-1, t=0;
  while ( (u&1)==0 ) u>>=1, t++;
  for (int i=0; i<3; i++){</pre>
    LL x = bin_pow( sprp[i]%n, u, n);
    if (x==0 || x==1 || x==n-1)continue;
    for (int j=1; j<t; j++){</pre>
      x=x*x%n;
      if (x==1 || x==n-1)break;
    if (x==n-1)continue;
    return 0;
  }
  return 1;
| }
```

5.7 Pollard's rho

```
// from PEC
// does not work when n is prime
Int f(Int x, Int mod){
 return add(mul(x, x, mod), 1, mod);
Int pollard_rho(Int n) {
  if ( !(n & 1) ) return 2;
  while (true) {
    Int y = 2, x = rand()%(n-1) + 1, res = 1;
    for ( int sz = 2 ; res == 1 ; sz *= 2 ) {
      for ( int i = 0 ; i < sz && res <= 1 ; i++) {</pre>
        x = f(x, n);
        res = \_gcd(abs(x-y), n);
      }
      y = x;
    if ( res != 0 && res != n ) return res;
  }
}
```

5.8 數論基本工具

```
Int POW(Int a, Int n, Int mod){
    Int re=1;
    while (n>0){
        if (n&1LL) re = re*a%mod;
        a = a*a%mod;
        n>>=1;
    }
    return re;
}
Int C(Int n, Int m){
    if (m<0 || m>n)return 0;
    return J[n] * inv(J[m]*J[n-m]%MOD) %MOD;
}
```

5.9 Mobius

5.10 SG

不為 0。

```
Sprague-Grundy
1. 雙人、回合制
2. 資訊完全公開
3. 無隨機因素
4. 可在有限步內結束
5. 沒有和局
6. 雙方可採取的行動相同
SG(S) 的值為 0:後手(P)必勝
不為 0: 先手(N)必勝
int mex(set S) {
 // find the min number >= 0 that not in the S
  // e.g. S = \{0, 1, 3, 4\} mex(S) = 2
state = []
int SG(A) {
  if (A not in state) {
   S = sub_states(A)
    if( len(S) > 1 ) state[A] = reduce(operator.xor, [
       SG(B) for B in S])
    else state[A] = mex(set(SG(B) for B in next_states(
       A)))
  }
  return state[A]
```

5.11 Theorem

| }

```
Lucas's Theorem
 For non-negative integer n,m and prime P,
 C(m,n) \mod P = C(m/M,n/M) * C(m%M,n%M) \mod P
  = mult_i ( C(m_i,n_i) )
 where m_i is the i-th digit of m in base P.
Pick's Theorem
 A = i + b/2 - 1
Kirchhoff's theorem
 A_{ii} = deg(i), A_{ij} = (i,j) \in E? -1:0
 Deleting any one row, one column, and cal the det(A)
Nth Catalan recursive function:
C_0 = 1, C_{n+1} = C_n * 2(2n + 1)/(n+2)
Mobius Formula
              , if n = 1
u(n) = 1
                , 若 n 無平方數因數,且 n = p1*p2*p3
       (-1)^m
           *...*pk
                ,若 n 有大於 1 的平方數因數
- Property
1. (積性函數) u(a)u(b) = u(ab)
2. \sum \{d|n\} \ u(d) = [n == 1]
Mobius Inversion Formula
       f(n) = \sum \{d|n\} \ g(d)
if
        g(n) = \sum \{d|n\} \ u(n/d)f(d)
            = \sum \{d|n\} \ u(d)f(n/d)
- Application
the number/power of gcd(i, j) = k
- Trick
分塊, O(sqrt(n))
Chinese Remainder Theorem (m_i 兩兩互質)
 x = a_1 \pmod{m_1}
 x = a_2 \pmod{m_2}
 x = a_i \pmod{m_i}
construct a solution:
 Let M = m_1 * m_2 * m_3 * ... * m_n
```

```
Let M_-i = M \ / \ m_-i

t_-i = 1 \ / \ M_-i
t_-i * M_-i = 1 \ (mod \ m_-i)

solution x = a_-1 * t_-1 * M_-1 + a_-2 * t_-2 * M_-2 + \dots
+ a_-n * t_-n * M_-n + k * M
= k*M + \sum a_-i * t_-i * M_-i, k \text{ is positive integer.}

under mod M, there is one solution x = \sum a_-i * t_-i * M_-i

Burnside's Lemma
|G| * |X/G| = sum( |X^-g|) \text{ where } g \text{ in } G
總方法數: 每一種旋轉下不動點的個數總和 除以 旋轉的方法 數
*/
```

6 Graph

6.1 BCC

```
邊雙連通
任 意 兩 點 間 至 少 有 兩 條 不 重 疊 的 路 徑 連 接 , 找 法 :
1. 標記出所有的橋
2. 對全圖進行 DFS,不走橋,每一次 DFS 就是一個新的邊雙
    捙 捅
// from BCW
struct BccEdge {
  static const int MXN = 100005;
  struct Edge { int v,eid; };
  int n,m,step,par[MXN],dfn[MXN],low[MXN];
  vector<Edge> E[MXN];
  DisjointSet djs;
  void init(int _n) {
   n = _n; m = 0;
    for (int i=0; i<n; i++) E[i].clear();</pre>
    djs.init(n);
  void add_edge(int u, int v) {
   E[u].PB({v, m});
    E[v].PB({u, m});
   m++;
  void DFS(int u, int f, int f_eid) {
    par[u] = f;
    dfn[u] = low[u] = step++;
    for (auto it:E[u]) {
      if (it.eid == f_eid) continue;
      int v = it.v;
      if (dfn[v] == -1) {
        DFS(v, u, it.eid);
        low[u] = min(low[u], low[v]);
      } else
        low[u] = min(low[u], dfn[v]);
   }
  void solve() {
    step = 0;
    memset(dfn, -1, sizeof(int)*n);
    for (int i=0; i<n; i++) {</pre>
      if (dfn[i] == -1) DFS(i, i, -1);
    djs.init(n);
    for (int i=0; i<n; i++) {</pre>
      if (low[i] < dfn[i]) djs.uni(i, par[i]);</pre>
  }
}graph;
```

6.2 Prim

```
// edae strucute
struct edge{
 int a, b;
  double data:
 bool operator <(const edge b)const{</pre>
   return data > b.data;
// main prim algorithm
int n, m, root, aa, bb, cc;
while (cin >> n >> m){
  priority_queue<edge>yee;
  int visit[500] = {}, p[500] = {};
  double a[500][500] = {};
  //undirectional edge aa to bb is weighted cc
  for (int i = 0; i < m; i++){</pre>
   cin >> aa >> bb >> cc;
    a[aa][bb] = a[bb][aa] = cc;
 }
 cin >> root;
 yee.push({ 0, root, 0 });
  edge tmp;
  double total = 0;
 while (!yee.empty()){
    tmp = yee.top(); yee.pop();
    if (visit[tmp.b])continue;
    total += tmp.data; p[tmp.b] = tmp.a; visit[tmp.b] =
        1;
    for (int i = 1; i <= n; i++){
      if (a[tmp.b][i]!=.0&&(!visit[i])){
       yee.push({tmp.b,i,a[tmp.b][i]});
      }
   }
 }
 cout << total << endl;</pre>
```

6.3 Bellman Ford

```
int a[100][100], d[100], p[100];
void bellman_ford(int root, int n){
  for (int i = 1; i <= n; i++)d[i] = 1e9;</pre>
  d[root] = 0, p[root] = 0;
  for (int i = 0; i<n - 1; i++){
    for (int j = 1; j <= n; j++){
  for (int k = 1; k <= n; k++){</pre>
         if (d[j] != 1e9 && a[j][k] != 1e9){
           if (d[j] + a[j][k] < d[k]){</pre>
             d[k] = d[j] + a[j][k], p[k] = j;
         }
      }
    }
  }
}
bool nega_cyc(int n){
  for (int i = 1; i <= n; i++){</pre>
    for (int j = 1; j <= n; j++){</pre>
      if (d[i] != 1e9 && a[i][j] != 1e9)
      if (d[i] + a[i][j] < d[j]){</pre>
         return 0;
      }
    }
  return 1;
int main(){
  int n, m, aa, bb, dd;
  while (cin >> n >> m){
    for (int i = 0; i \leftarrow n; i++)for (int j = 0; j \leftarrow n;
          j++){
      a[i][j] = E9;
    memset(p, 0, sizeof(p));
    for (int i = 0; i < m; i++){</pre>
      cin >> aa >> bb >> dd;
```

```
a[aa][bb] = min(a[aa][bb], dd);
}
cin >> aa;
bellman_ford(aa, n);
int t = nega_cyc(n);
if(t){
    for (int i = 1; i <= n; i++)cout << d[i] << " \n"
        [i==n];
    for (int i = 1; i <= n; i++)cout << p[i] << " \n"
        [i==n];
}
else cout << "There is a negative weight cycle in
        the graph\n";
}
</pre>
```

6.4 Kruskal

```
struct v {
  int a, b, c;
int p[200001];v a[200001];
bool sor(v a, v b) {
 return a.c < b.c;</pre>
int find(int x) {
  return(x != p[x] ? (p[x] = find(p[x])) : x);
int main() {
  int n, m, i, j, sum;
  while (cin >> n >> m) {
    for (i = 0; i < 200001; i++)p[i] = i;</pre>
    for (i = 0; i<m; i++)cin >> a[i].a >> a[i].b >> a[i
        1.c;
    sort(a, a + m, sor);
    for (i =0,j = 0;j<m; j++) {</pre>
      if(find(a[j].a) != find(a[j].b)){
        i++;
        p[find(a[j].a)] = find(a[j].b);
        sum += a[j].c;
    cout << ((i==n-1)?sum:-1) << endl;</pre>
  }
}
```

6.5 Dijkstra

```
struct node {
    int num{}, w{};
    bool operator < (const node& other)const {</pre>
        return w > other.w;
    }
};
vector<int> dijkstra(int root, const vector<vector<node</pre>
    >> &graph) {
    vector<int> d(graph.size(), INT_MAX >> 1), p(graph.
        size());
    priority_queue<node> pq;
    d[root] = p[root] = 0;
    pq.push({root, d[root]});
    while (!pq.empty()) {
        node tmp = pq.top(); pq.pop();
        for (const node &i : graph[tmp.num]) {
            if (d[i.num] > d[tmp.num] + i.w) {
                d[i.num] = d[tmp.num] + i.w;
                p[i.num] = tmp.num;
                pq.push({i.num, d[tmp.num]});
            }
        }
    return d;
```

6.6 Strongly Connected Component(SCC)

| }

```
#define MXN 100005
#define PB push_back
#define FZ(s) memset(s,0,sizeof(s))
struct Scc{
int n, nScc, vst[MXN], bln[MXN];
vector<int> E[MXN], rE[MXN], vec;
void init(int _n){
  n = _n;
  for (int i=0; i<MXN; i++){</pre>
    E[i].clear();
    rE[i].clear();
  }
}
void add_edge(int u, int v){
  E[u].PB(v);
  rE[v].PB(u);
void DFS(int u){
  vst[u]=1;
  for (auto v : E[u])
    if (!vst[v]) DFS(v);
  vec.PB(u);
void rDFS(int u){
  vst[u] = 1;
  bln[u] = nScc;
  for (auto v : rE[u])
    if (!vst[v]) rDFS(v);
void solve(){
  nScc = 0;
  vec.clear();
  FZ(vst);
  for (int i=0; i<n; i++)</pre>
    if (!vst[i]) DFS(i);
  reverse(vec.begin(), vec.end());
  FZ(vst);
  for (auto v : vec){
    if (!vst[v]){
      rDFS(v);
      nScc++;
    }
  }
}
|};
```

6.7 Hungarian

```
// Maximum Cardinality Bipartite Matching
struct Graph {
    static const int MAXN = 5005;
    vector<int> G[MAXN];
    int n:
    int match[MAXN]; // Matching Result
    int vis[MAXN];
    void init(int _n) {
         for ( int i = 0 ; i < n ; i++ ) G[i].clear();</pre>
    bool dfs(int u) {
         for ( auto v:G[u] ) {
             if (!vis[v]) {
                 vis[v] = true;
                 if (match[v] == -1 || dfs(match[v])) {
                     match[v] = u;
                     match[u] = v;
                     return true:
                 }
             }
         }
```

```
return false;
}
int solve() {
    int res = 0;
    memset(match, -1, sizeof(match));
    for (int i = 0; i < n; i++) {
        if (match[i] == -1) {
            memset(vis, 0, sizeof(vis));
            if (dfs(i)) res += 1;
        }
    }
    return res;
}
graph;</pre>
```

6.8 KM

```
Detect non-perfect-matching:

    set all edge[i][j] as INF

2. if solve() >= INF, it is not perfectmatching.
// Maximum Weight Perfect Bipartite Matching
// allow negative weight!
typedef long long Int;
struct KM {
    static const int MAXN = 1050;
static const int INF = 1LL<<60;</pre>
    int n, match[MAXN], vx[MAXN], vy[MAXN];
    Int edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[
        MAXN];
    void init(int _n){
        n = _n;
for ( int i = 0 ; i < n ; i++ )</pre>
             for ( int j = 0; j < n ; j++ )
                 edge[i][j] = 0;
    void add_edge(int x, int y, Int w){
        edge[x][y] = w;
    bool DFS(int x){
        vx[x] = 1;
        for ( int y = 0 ; y < n ; y++ ) {</pre>
             if ( vy[y] ) continue;
             if (lx[x] + ly[y] > edge[x][y]) {
                  slack[y] = min(slack[y], lx[x] + ly[y]
                        edge[x][y]);
                  vy[y] = 1;
                  if ( match[y] == -1 || DFS(match[y]) ){
                      match[y] = x;
                      return true;
                  }
             }
        return false;
    Int solve() {
         fill(match, match + n, -1);
        fill(lx, lx + n, -INF);
         fill(ly, ly + n, 0);
        for ( int i = 0; i < n; i++ )</pre>
             for ( int j = 0; j < n; j++ )
                 lx[i] = max(lx[i], edge[i][j]);
        for ( int i = 0 ; i < n; i++ ) {</pre>
             fill(slack, slack + n, INF);
             while (true){
                 fill(vx, vx + n, 0);
                  fill(vy, vy + n, 0);
                 if ( DFS(i) ) break;
                 Int d = INF;
                 for ( int j = 0 ; j < n ; j++ )
    if ( !vy[j] ) d = min(d, slack[j]);</pre>
                  for ( int j = 0 ; j < n ; j++ ) {</pre>
                      if (vx[j]) 1x[j] -= d;
                      if (vy[j]) ly[j] += d;
                      else slack[j] -= d;
                 }
             }
```

```
}
    Int res = 0;
    for ( int i = 0 ; i < n ; i++ ) {
        res += edge[ match[i] ][i];
    }
    return res;
}
} graph;</pre>
```

6.9 最小平均環

```
// from BCW
/* minimum mean cycle */
const int MAXE = 1805;
const int MAXN = 35;
const double inf = 1029384756;
const double eps = 1e-6;
struct Edge {
  int v,u;
  double c;
int n,m,prv[MAXN][MAXN], prve[MAXN][MAXN], vst[MAXN];
Edge e[MAXE];
vector<int> edgeID, cycle, rho;
double d[MAXN][MAXN];
inline void bellman_ford() {
  for(int i=0; i<n; i++) d[0][i]=0;
for(int i=0; i<n; i++) {</pre>
     fill(d[i+1], d[i+1]+n, inf);
    for(int j=0; j<m; j++) {</pre>
       int v = e[j].v, u = e[j].u;
       if(d[i][v]<inf && d[i+1][u]>d[i][v]+e[j].c) {
         d[i+1][u] = d[i][v]+e[j].c;
         prv[i+1][u] = v;
         prve[i+1][u] = j;
      }
    }
  }
double karp_mmc() {
  // returns inf if no cycle, mmc otherwise
  double mmc=inf;
   int st = -1;
  bellman_ford();
   for(int i=0; i<n; i++) {</pre>
    double avg=-inf;
    for(int k=0; k<n; k++) {</pre>
       if(d[n][i]<inf-eps) avg=max(avg,(d[n][i]-d[k][i])</pre>
           /(n-k));
       else avg=max(avg,inf);
    if (avg < mmc) tie(mmc, st) = tie(avg, i);</pre>
  for(int i=0; i<n; i++) vst[i] = 0;</pre>
  edgeID.clear(); cycle.clear(); rho.clear();
  for (int i=n; !vst[st]; st=prv[i--][st]) {
    vst[st]++:
    edgeID.PB(prve[i][st]);
    rho.PB(st);
  while (vst[st] != 2) {
    int v = rho.back(); rho.pop_back();
    cycle.PB(v);
    vst[v]++;
  reverse(ALL(edgeID));
  edgeID.resize(SZ(cycle));
  return mmc:
```

6.10 偵測負環

```
#include <bits/stdc++.h>
using namespace std;

const int INF = 1000000;
const int MAXN = 200;
```

```
int n, m, q;
int d[MAXN][MAXN];
int main () {
    while ( cin >> n >> m >> q && n) {
         for ( int i = 0 ; i <= n ; i++ ) {</pre>
              for ( int j = 0 ; j <= n ; j++ ) d[i][j] =</pre>
                  (i==j?0:INF);
         }
         for ( int i = 0 ; i < m ; i++ ) {</pre>
             int a, b, c;
             cin >> a >> b >> c;
             d[a][b] = min(d[a][b], c);
         for ( int k = 0 ; k < n ; k++ ) {
              for ( int i = 0 ; i < n ; i++ ) {</pre>
                  for ( int j = 0 ; j < n ; j++ ) {</pre>
                      if ( d[i][j] > d[i][k] + d[k][j] &&
                            d[i][k] < INF \&\& d[k][j] < INF
                           //printf("%d > %d + %d\n", d[i
                                ][j], d[i][k], d[k][j]);
                           //if ( d[i][k] >= INF || d[k][j
                               ] >= INF ) cout << "NO : "
<< i << " " << j << " " << k << "--";
                           d[i][j] = min(d[i][j], d[i][k]
                                + d[k][j]);
                      }
                  }
             }
         }
         for ( int i = 0 ; i < n ; i++ ) {</pre>
             for ( int j = 0 ; j < n ; j++ ) {
                  for ( int k = 0 ; k < n && d[i][j] != -</pre>
                      INF ; k++ ) {
                      if ( d[k][k] < 0 && d[i][k] != INF</pre>
                           && d[k][j] != INF )
                           d[i][j] = -INF;
                  }
             }
         int u, v;
         for (int i=0;i<q;i++){</pre>
             scanf("%d%d",&u,&v);
             if (d[u][v] == INF) printf("Impossible\n");
             else if (d[u][v] == -INF) printf("-Infinity
                  n";
             else printf("%d\n",d[u][v]);
         puts("");
     return 0;
}
```

6.11 Tarjan

```
|割點

|點 u 為割點 if and only if 滿足 1. or 2.

|1. u 爲樹根,且 u 有多於一個子樹。

|2. u 不爲樹根,且滿足存在 (u,v) 爲樹枝邊 (或稱父子邊,即 u 爲 v 在搜索樹中的父親),使得 DFN(u) <= Low(v)。

|橋

|一條無向邊 (u,v) 是橋 if and only if (u,v) 爲樹枝邊,且滿足 DFN(u) < Low(v)。

// 0 base

struct TarjanSCC{

static const int MAXN = 1000006;

int n, dfn[MAXN], low[MAXN], scc[MAXN], scn, count;
```

```
vector<int> G[MAXN];
  stack<int> stk;
  bool ins[MAXN];
  void tarjan(int u){
    dfn[u] = low[u] = ++count;
    stk.push(u);
    ins[u] = true;
    for(auto v:G[u]){
      if(!dfn[v]){
        tarjan(v);
        low[u] = min(low[u], low[v]);
      }else if(ins[v]){
        low[u] = min(low[u], dfn[v]);
    if(dfn[u] == low[u]){
      int v;
      do {
      v = stk.top();
      stk.pop();
      scc[v] = scn;
      ins[v] = false;
      } while(v != u);
      scn++;
    }
  }
  void getSCC(){
    memset(dfn,0,sizeof(dfn));
    memset(low,0,sizeof(low));
    memset(ins,0,sizeof(ins));
    memset(scc,0,sizeof(scc));
    count = scn = 0;
    for(int i = 0 ; i < n ; i++ ){</pre>
      if(!dfn[i]) tarjan(i);
  }
\}SCC;
```

6.12 Topological Sort

```
#define N 87
bool adj[N][N];
                     // adjacency matrix
                     // record visited coordinations in
int visit[N];
    DFS
int order[N], n;
                     // save the order
bool cycle;
                    // detect the cycle
void DFS(int s)
    // back edge occured, detected the cycle
    if (visit[s] == 1) {cycle = true; return;}
    // forward edge and cross edge;C
    if (visit[s] == 2) return;
    visit[s] = 1;
    for (int t=0; t<N; ++t){</pre>
        if (adj[s][t]) DFS(t);
    visit[s] = 2;
    order[n--] = s;
                         // record the order
}
void topological_ordering()
    memset(visit, 0, sizeof(visit));
    cycle = false;
    n = N - 1:
    for (int s=0; s<9; ++s)</pre>
        if (!v[s])
            DFS(s);
    if (cycle) cout << "The graph has the cycle!";</pre>
```

7 Data Structure

7.1 2D Range Tree

}

```
// remember sort x !!!!!
typedef int T;
const int LGN = 20;
const int MAXN = 100005;
struct Point{
    friend bool operator < (Point a, Point b){</pre>
        return tie(a.x,a.y) < tie(b.x,b.y);</pre>
    }
};
struct TREE{
    Point pt;
    int toleft;
}tree[LGN][MAXN];
struct SEG{
    T mx, Mx;
    int sz;
    TREE *st;
}seg[MAXN*4];
vector<Point> P;
void build(int 1, int r, int o, int deep){
    seg[o].mx = P[1].x;
    seg[o].Mx = P[r].x;
    seg[o].sz = r-l+1;;
    if(1 == r){
        tree[deep][r].pt = P[r];
        tree[deep][r].toleft = 0;
        seg[o].st = &tree[deep][r];
        return;
    int mid = (1+r)>>1;
    build(1,mid,o+o,deep+1);
    build(mid+1,r,o+o+1,deep+1);
    TREE *ptr = &tree[deep][1];
    TREE *pl = &tree[deep+1][1], *nl = &tree[deep+1][
    TREE *pr = &tree[deep+1][mid+1], *nr = &tree[deep
        +1][r+1];
    int cnt = 0;
    while(pl != nl && pr != nr) {
        *(ptr) = pl->pt.y <= pr->pt.y ? cnt++, *(pl++):
             *(pr++);
        ptr -> toleft = cnt; ptr++;
    while(pl != nl) *(ptr) = *(pl++), ptr -> toleft =
        ++cnt, ptr++;
    while(pr != nr) *(ptr) = *(pr++), ptr -> toleft =
        cnt, ptr++;
int main(){
    int n; cin >> n;
    for(int i = 0 ;i < n; i++){</pre>
        T x,y; cin >> x >> y;
        P.push_back((Point){x,y});
    sort(P.begin(),P.end());
    build(0,n-1,1,0);
```

7.2 Sparse Table

```
const int MAXN = 200005;
 const int lgN = 20;
 struct SP{ //sparse table
  int Sp[MAXN][lgN];
  function<int(int,int)> opt;
   void build(int n, int *a){ // 0 base
     for (int i=0 ;i<n; i++) Sp[i][0]=a[i];</pre>
     for (int h=1; h<lgN; h++){</pre>
       int len = 1<<(h-1), i=0;</pre>
       for (; i+len<n; i++)</pre>
         Sp[i][h] = opt( Sp[i][h-1] , Sp[i+len][h-1] );
       for (; i<n; i++)</pre>
         Sp[i][h] = Sp[i][h-1];
     }
   int query(int 1, int r){
     int h = __lg(r-l+1);
     int len = 1<<h;</pre>
     return opt( Sp[1][h] , Sp[r-len+1][h] );
};
```

7.3 Segment Tree

```
struct Node{
    int mx; // 區間最大值
    int tag; // 子樹裡所有人都要加上 tag
};
vector<Node> seg;
// 節點 id 的整個區間要加上 tag
void addtag(int tag, int id){
    seg[id].mx += tag; // 最大值會加上 tag
    seg[id].tag += tag; // 注意可能本來就有標記了,所以
// 更新子節點資訊並把標記移到子節點身上
void push(int id){
    addtag(seg[id].tag, lc);
    addtag(seg[id].tag, rc);
    seg[id].tag = 0; // 標記被移到子節點上所以要改成 0
// 區間 [l,r] 加上 v
void modify(int 1, int r, int v, int L, int R, int id){
    if(1 <= L && R <= r){
        addtag(v, id);
        return;
    push(id);
    if(r <= M) modify(l, r, v, L, M, lc);</pre>
    else if(l > M) modify(l, r, v, M + 1, R, rc);
    else{
        modify(1, r, v, L, M, lc);
        modify(l, r, v, M + 1, R, rc);
    seg[id].mx = max(seg[lc].mx, seg[rc].mx);
int query(int 1, int r, int L, int R, int id){
    if(1 <= L && R <= r) return seg[id].mx;</pre>
    push(id);
    int M = (L + R) / 2;
    if(r <= M) return query(l, r, L, M, lc);</pre>
    else if(l > M) return query(l, r, M + 1, R, rc);
    else return max(query(1, r, L, M, lc),
                   query(l, r, M + 1, R, rc);
| }
```

7.4 ZKW 線段樹

```
const int M=1e5+111;
int n,m,q;
int sum[M<<2],mn[M<<2],mx[M<<2],add[M<<2];</pre>
int read() {
    int x;
    cin >> x;
    return x;
void build(){
    for(m=1;m<=n;m<<=1);</pre>
    for(int i=m+1;i<=m+n;++i)</pre>
        sum[i]=mn[i]=mx[i]=read();
    for(int i=m-1;i;--i){
        sum[i]=sum[i<<1]+sum[i<<1|1];
        mn[i]=min(mn[i<<1],mn[i<<1|1]),</pre>
        mn[i << 1]-=mn[i], mn[i << 1|1]-=mn[i];
        mx[i]=max(mx[i<<1],mx[i<<1|1]),
        mx[i << 1] -= mx[i], mx[i << 1|1] -= mx[i];
    }
void update_node(int x,int v,int A=0){
    x+=m, mx[x]+=v, mn[x]+=v, sum[x]+=v;
    for(;x>1;x>>=1){
        sum[x]+=v:
        A=min(mn[x],mn[x^1]);
        mn[x]-=A, mn[x^1]-=A, mn[x>>1]+=A;
        A=\max(\max[x],\max[x^1]),
        mx[x]-=A, mx[x^1]-=A, mx[x>>1]+=A;
    }
void update_part(int s,int t,int v){
    int A=0,lc=0,rc=0,len=1;
    for(s+=m-1,t+=m+1;s^t^1;s>>=1,t>>=1,len<<=1){</pre>
        if(s&1^1) add[s^1]+=v,lc+=len, mn[s^1]+=v,mx[s
             ^1]+=v;
        if(t&1)
                   add[t^1]+=v,rc+=len, mn[t^1]+=v,mx[t]
            ^1]+=v;
        sum[s>>1]+=v*lc, sum[t>>1]+=v*rc;
        A=min(mn[s],mn[s^1]),mn[s]-=A,mn[s^1]-=A,mn[s
            >>1]+=A.
        A=min(mn[t],mn[t^1]),mn[t]-=A,mn[t^1]-=A,mn[t]
             >>1]+=A;
        A=max(mx[s],mx[s^1]),mx[s]-=A,mx[s^1]-=A,mx[s
            >>1]+=A,
        A=\max(\max[t],\max[t^1]),\max[t]-=A,\max[t^1]-=A,\max[t]
            >>1]+=A;
    for(lc+=rc;s;s>>=1){
        sum[s>>1]+=v*lc;
        A=min(mn[s],mn[s^1]),mn[s]-=A,mn[s^1]-=A,mn[s
            >>1]+=A,
        A=\max(\max[s],\max[s^1]),\max[s]-=A,\max[s^1]-=A,\max[s
            >>1]+=A;
    }
int query node(int x,int ans=0){
    for(x+=m;x;x>>=1) ans+=mn[x]; return ans;
int query_sum(int s,int t){
    int lc=0,rc=0,len=1,ans=0;
    for(s+=m-1,t+=m+1;s^t^1;s>>=1,t>>=1,len<<=1){</pre>
        if(s&1^1) ans+=sum[s^1]+len*add[s^1],lc+=len;
        if(t&1) ans+=sum[t^1]+len*add[t^1],rc+=len;
        if(add[s>>1]) ans+=add[s>>1]*lc;
        if(add[t>>1]) ans+=add[t>>1]*rc;
    for(lc+=rc,s>>=1;s;s>>=1) if(add[s]) ans+=add[s]*lc
    return ans;
int query_min(int s,int t,int L=0,int R=0,int ans=0){
    if(s==t) return query_node(s);
    for(s+=m,t+=m;s^t^1;s>>=1,t>>=1){
        L+=mn[s],R+=mn[t];
        if(s&1^1) L=min(L,mn[s^1]);
        if(t&1) R=min(R,mn[t^1]);
    for(ans=min(L,R),s>>=1;s;s>>=1) ans+=mn[s];
    return ans;
int query_max(int s,int t,int L=0,int R=0,int ans=0){
```

```
if(s==t) return query_node(s);
for(s+=m,t+=m;s^t^1;s>>=1,t>>=1){
    L+=mx[s],R+=mx[t];
    if(s&1^1) L=max(L,mx[s^1]);
    if(t&1) R=max(R,mx[t^1]);
}
for(ans=max(L,R),s>>=1;s;s>>=1) ans+=mx[s];
return ans;
}
```

7.5 Lazy Tag

```
| void modify(type value, int 1, int r, int L, int R, vertex v){
    if(1 == L && r == R){
        //打懶標在v上;
        return;
    }
    int M = (L + R) / 2;
    if(r <= M) modify(value, 1, r, L, M, //v的左子節點);
    else if(1 > M) modify(value, 1, r, M + 1, R, //v的 右子節點);
    else{
        modify(value, 1, M, L, M, v的左子節點);
        modify(value, M + 1, r, M + 1, R, //v的右子節點);
        modify(value, M + 1, r, M + 1, R, //v的右子節點);
    }
    //用兩個子節點的答案更新v的答案;
}
```

7.6 BIT 樹狀樹組

```
class Bitree {
public:
    vector<int> data;
     Bitree(const vector<int> &nums) {
         data.resize(nums.size() + 1, 0);
         for(int i = 0; i < nums.size(); i++ ) {</pre>
             update(i, nums[i]);
    void update(int x, int val) {
         for(; x < data.size(); x += lowbit(x)) {</pre>
             data[x] += val;
    int query(int x) {
         x++;
         int result = 0;
         for(; x > 0; x \rightarrow 1 lowbit(x)) {
             result += data[x];
         return result;
    int lowbit(int x) {
         return x & (-x);
};
```

8 String

8.1 KMP

```
template < typename T>
void build_KMP(int n, T *s, int *f){ // 1 base
  f[0]=-1, f[1]=0;
  for (int i=2; i<=n; i++){
    int w = f[i-1];
    while (w>=0 && s[w+1]!=s[i])w = f[w];
    f[i]=w+1;
}
```

```
template < typename T>
int KMP(int n, T *a, int m, T *b) {
  build_KMP(m,b,f);
  int ans=0;

for (int i=1, w=0; i<=n; i++) {
    while ( w>=0 && b[w+1]!=a[i] )w = f[w];
    w++;
    if (w==m) {
        ans++;
        w=f[w];
    }
  }
  return ans;
}
```

8.2 smallest rotation

```
| string mcp(string s){
    int n = s.length();
    s += s;
    int i=0, j=1;
    while (i<n && j<n){
        int k = 0;
        while (k < n && s[i+k] == s[j+k]) k++;
        if (s[i+k] <= s[j+k]) j += k+1;
        else i += k+1;
        if (i == j) j++;
    }
    int ans = i < n ? i : j;
    return s.substr(ans, n);
}
/*

Booth 演算法

用於尋找一個字串的字典序最小的循環旋轉
*/
Contact GitHub API Training Shop Blog About
```

8.3 Suffix Array

```
/*he[i]保存了在後綴數組中相鄰兩個後綴的最長公共前綴長度
*sa[i]表示的是字典序排名為i的後綴是誰(字典序越小的排
     名越靠前)
*rk[i]表示的是後綴我所對應的排名是多少 */
const int MAX = 1020304;
int ct[MAX], he[MAX], rk[MAX];
int sa[MAX], tsa[MAX], tp[MAX][2];
void suffix_array(char *ip){
 int len = strlen(ip);
  int alp = 256;
  memset(ct, 0, sizeof(ct));
  for(int i=0;i<len;i++) ct[ip[i]+1]++;</pre>
  for(int i=1;i<alp;i++) ct[i]+=ct[i-1];</pre>
  for(int i=0;i<len;i++) rk[i]=ct[ip[i]];</pre>
  for(int i=1;i<len;i*=2){</pre>
    for(int j=0;j<len;j++){</pre>
      if(j+i>=len) tp[j][1]=0;
      else tp[j][1]=rk[j+i]+1;
      tp[j][0]=rk[j];
    memset(ct, 0, sizeof(ct));
    for(int j=0;j<len;j++) ct[tp[j][1]+1]++;</pre>
    for(int j=1;j<len+2;j++) ct[j]+=ct[j-1];</pre>
    for(int j=0;j<len;j++) tsa[ct[tp[j][1]]++]=j;</pre>
    memset(ct, 0, sizeof(ct));
    for(int j=0;j<len;j++) ct[tp[j][0]+1]++;</pre>
    for(int j=1;j<len+1;j++) ct[j]+=ct[j-1];</pre>
    for(int j=0;j<len;j++)</pre>
      sa[ct[tp[tsa[j]][0]]++]=tsa[j];
    rk[sa[0]]=0;
    for(int j=1;j<len;j++){</pre>
      if( tp[sa[j]][0] == tp[sa[j-1]][0] &&
        tp[sa[j]][1] == tp[sa[j-1]][1] )
```

```
rk[sa[j]] = rk[sa[j-1]];
else
    rk[sa[j]] = j;
}

for(int i=0,h=0;i<len;i++){
    if(rk[i]==0) h=0;
    else{
        int j=sa[rk[i]-1];
        h=max(0,h-1);
        for(;ip[i+h]==ip[j+h];h++);
    }
    he[rk[i]]=h;
}</pre>
```

8.4 Z-value

```
z[0] = 0:
for ( int bst = 0, i = 1; i < len ; i++ ) {</pre>
  if (z[bst] + bst <= i) z[i] = 0;
  else z[i] = min(z[i - bst], z[bst] + bst - i);
  while ( str[i + z[i]] == str[z[i]] ) z[i]++;
  if ( i + z[i] > bst + z[bst] ) bst = i;
// 回文版
void Zpal(const char *s, int len, int *z) {
    // Only odd palindrome len is considered
    // z[i] means that the longest odd palindrom
         centered at
    // i is [i-z[i] .. i+z[i]]
    z[0] = 0;
    for (int b=0, i=1; i<len; i++) {</pre>
        if (z[b] + b >= i) z[i] = min(z[2*b-i], b+z[b]-
            i);
        else z[i] = 0;
        while (i+z[i]+1 < len and i-z[i]-1 >= 0 and
               S[i+z[i]+1] == S[i-z[i]-1]) z[i] ++;
        if(z[i] + i > z[b] + b) b = i;
    }
}
```

8.5 旋轉哈希

```
typedef unsigned __int128 ull1;
ulll power(ulll a, ulll n, ulll m) {
    ulll re = 1;
    while (n > 0) {
        if (n & 1) re = re * a % m;
        a = a * a % m;
        n >>= 1;
    return re;
}
ulll inv(ulll a, ulll m) {
    return power(a, m - 2, m);
}
struct Rh {
    const ulll p, mod;
    vector<ulll> ps{1};
    Rh(ulll p, ulll mod) : p(p), mod(mod) {}
    vector<ull>> build(const string &s) {
        vector<ulll> h(s.size() + 1);
        h[0] = 0;
        ps.resize(s.size() + 1);
        for (int i = 0; i < s.size(); i++) {</pre>
            ps[i + 1] = ps[i] * p % mod;
            h[i + 1] = (h[i] + s[i] * ps[i + 1] % mod)
                % mod:
        return h;
    }
```

```
ulll subhash(const vector<ulll> &h, int 1, int r) {
        // [l, r] 指原字串
        return ((h[r + 1] - h[l]) * inv(ps[l], mod)) %
            mod:
    }
};
constexpr uint64_t mod = (1ull<<61) - 1;</pre>
uint64_t modmul(uint64_t a, uint64_t b){
 uint64_t l1 = (uint32_t)a, h1 = a>>32, l2 = (uint32_t
      )b, h2 = b >> 32;
 uint64_t l = 11*12, m = 11*h2 + 12*h1, h = h1*h2;
 uint64_t ret = (1\&mod) + (1>>61) + (h << 3) + (m >>
      29) + (m << 35 >> 3) + 1;
 ret = (ret & mod) + (ret>>61);
 ret = (ret & mod) + (ret>>61);
 return ret-1;
```

9 Others

新的方法介绍

9.1 矩陣樹定理

```
下面我们介绍一种新的方法——Matrix-Tree定理(Kirchhoff矩
Matrix-Tree定理是解决生成树计数问题最有力的武器之一。它
   首先于1847年被Kirchhoff证明。在介绍定理之前,我们首
   先 明 确 几 个 概 念 :
1、G的度数矩阵D[G]是一个n*n的矩阵,并且满足:当i≠j时,
   dij=0;当i=j时,dij等于vi的度数。
2、G的邻接矩阵A[G]也是一个n*n的矩阵, 并且满足:如果vi
    、vj之间有边直接相连,则aij=1,否则为0。
我们定义G的Kirchhoff矩阵(也称为拉普拉斯算子)C[G]为C[G]=
   D[G]-A[G] ,
则Matrix-Tree定理可以描述为:G的所有不同的生成树的个数
   等于其Kirchhoff矩阵C[G]任何一个n-1阶主子式的行列式
   的绝对值。
所谓n-1阶主子式,就是对于r(1≤r≤n),将C[G]的第r行、第r列
   同时去掉后得到的新矩阵,用Cr[G]表示。
生成树计数
算法步骤:
1、 构建拉普拉斯矩阵
    Matrix[i][i] =
degree(i) , i==j
        -1,i-j有边
        0,其他情况
2、 去掉第r行,第r列(r任意)
3、 计算矩阵的行列式
#include <stdio.h>
#include <string.h>
#include <algorithm>
#include <iostream>
#include <math.h>
using namespace std;
const double eps = 1e-8;
const int MAXN = 110;
int sgn(double x)
   if(fabs(x) < eps)return 0;</pre>
   if(x < 0) return -1;
   else return 1;
double b[MAXN][MAXN];
double det(double a[][MAXN],int n)
   int i, j, k, sign = 0;
   double ret = 1;
   for(i = 0;i < n;i++)</pre>
   for(j = 0;j < n;j++) b[i][j] = a[i][j];</pre>
   for(i = 0;i < n;i++)</pre>
```

```
if(sgn(b[i][i]) == 0)
              for(j = i + 1; j < n;j++)</pre>
              if(sgn(b[j][i]) != 0) break;
              if(j == n)return 0;
              for(k = i;k < n;k++) swap(b[i][k],b[j][k]);</pre>
              sign++;
         ret *= b[i][i];
         for(k = i + 1;k < n;k++) b[i][k]/=b[i][i];</pre>
         for(j = i+1; j < n; j++)</pre>
         for(k = i+1; k < n; k++) b[j][k] -= b[j][i]*b[i][
     if(sign & 1)ret = -ret;
     return ret;
double a[MAXN][MAXN];
int g[MAXN][MAXN];
 int main()
     int T;
     int n,m;
     int u,v;
     scanf("%d",&T);
     while(T--)
         scanf("%d%d",&n,&m);
         memset(g,0,sizeof(g));
         while(m--)
         {
              scanf("%d%d",&u,&v);
              u--:v--:
              g[u][v] = g[v][u] = 1;
         memset(a,0,sizeof(a));
         for(int i = 0;i < n;i++)</pre>
         for(int j = 0; j < n; j++)</pre>
         if(i != j && g[i][j])
              a[i][i]++;
              a[i][j] = -1;
         double ans = det(a,n-1);
         printf("%.0lf \setminus n",ans);
     return 0:
|}
```

9.2 1D/1D dp 優化

```
#include < bits / stdc++.h>
int t, n, L;
int p;
char s[MAXN][35];
11 \text{ sum}[MAXN] = \{0\};
long double dp[MAXN] = {0};
int prevd[MAXN] = {0};
long double pw(long double a, int n) {
    if ( n == 1 ) return a;
    long double b = pw(a, n/2);
    if ( n & 1 ) return b*b*a;
    else return b*b;
long double f(int i, int j) {
     cout << (sum[i] - sum[j]+i-j-1-L) << endl;</pre>
    return pw(abs(sum[i] - sum[j]+i-j-1-L), p) + dp[j];
struct INV {
    int L, R, pos;
INV stk[MAXN*10];
int top = 1, bot = 1;
void update(int i) {
    while ( top > bot && i < stk[top].L && f(stk[top].L</pre>
         , i) < f(stk[top].L, stk[top].pos) ) {</pre>
         stk[top - 1].R = stk[top].R;
        top--;
```

```
int lo = stk[top].L, hi = stk[top].R, mid, pos =
         stk[top].pos;
    //if ( i >= lo ) lo = i + 1;
    while ( lo != hi ) {
        mid = lo + (hi - lo) / 2;
        if ( f(mid, i) < f(mid, pos) ) hi = mid;</pre>
        else lo = mid + 1;
    if ( hi < stk[top].R ) {</pre>
        stk[top + 1] = (INV) { hi, stk[top].R, i };
        stk[top++].R = hi;
    }
}
int main() {
    cin >> t:
    while ( t-- ) {
        cin >> n >> L >> p;
        dp[0] = sum[0] = 0;
        for ( int i = 1 ; i <= n ; i++ ) {</pre>
             cin >> s[i];
             sum[i] = sum[i-1] + strlen(s[i]);
             dp[i] = numeric_limits<long double>::max();
        stk[top] = (INV) \{1, n + 1, 0\};
        for ( int i = 1 ; i <= n ; i++ ) {</pre>
             if ( i >= stk[bot].R ) bot++;
             dp[i] = f(i, stk[bot].pos);
             update(i);
//
               cout << (ll) f(i, stk[bot].pos) << endl;</pre>
        if ( dp[n] > 1e18 ) {
             cout << "Too hard to arrange" << endl;
        } else {
             vector<PI> as;
             cout << (11)dp[n] << endl;</pre>
    return 0;
}
```

9.3 Theorm - DP optimization

```
Monotonicity & 1D/1D DP & 2D/1D DP
Definition xD/yD
1D/1D DP[j] = min(0 \le i < j) \{ DP[i] + w(i, j) \}; DP[0] = k
2D/1D DP[i][j] = min(i < k \le j) \{ DP[i][k - 1] + DP[k][j] \}
    + w(i, j); DP[i][i] = 0
Monotonicity
a \mid w(a, c) w(a, d)
b | w(b, c) w(b, d)
Monge Condition
Concave(凹四邊形不等式): w(a, c) + w(b, d) >= w(a, d) +
     w(b, c)
Convex (凸四邊形不等式): w(a, c) + w(b, d) <= w(a, d) +
     w(b, c)
Totally Monotone
Concave(凹單調): w(a, c) <= w(b, d) ----> w(a, d) <= w
    (b, c)
Convex (凸單調): w(a, c) >= w(b, d) ----> w(a, d) >= w
    (b, c)
1D/1D DP O(n^2) \rightarrow O(nlgn)
**CONSIDER THE TRANSITION POINT**
Solve 1D/1D Concave by Stack
Solve 1D/1D Convex by Deque
2D/1D Convex DP (Totally Monotone) O(n^3) \rightarrow O(n^2)
h(i, j - 1) \le h(i, j) \le h(i + 1, j)
```

9.4 Stable Marriage

```
// normal stable marriage problem
// input:
//3
//Albert Laura Nancy Marcy
//Brad Marcy Nancy Laura
//Chuck Laura Marcy Nancy
//Laura Chuck Albert Brad
//Marcy Albert Chuck Brad
//Nancy Brad Albert Chuck
#include < bits / stdc++.h>
using namespace std;
const int MAXN = 505;
int n;
int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id
int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank
int current[MAXN]; // current[boy_id] = rank; boy_id
    will pursue current[boy_id] girl.
int girl_current[MAXN]; // girl[girl_id] = boy_id;
void initialize() {
  for ( int i = 0 ; i < n ; i++ ) {</pre>
    current[i] = 0;
    girl_current[i] = n;
    order[i][n] = n;
  }
}
map<string, int> male, female;
string bname[MAXN], gname[MAXN];
int fit = 0;
void stable_marriage() {
  queue<int> que;
  for ( int i = 0 ; i < n ; i++ ) que.push(i);</pre>
  while ( !que.empty() ) {
    int boy_id = que.front();
    que.pop();
    int girl_id = favor[boy_id][current[boy_id]];
    current[boy_id] ++;
    if ( order[girl_id][boy_id] < order[girl_id][</pre>
         girl_current[girl_id]] ) {
      if ( girl_current[girl_id] < n ) que.push(</pre>
           girl_current[girl_id]); // if not the first
      girl_current[girl_id] = boy_id;
    } else {
      que.push(boy_id);
  }
int main() {
  cin >> n:
  for ( int i = 0 ; i < n; i++ ) {</pre>
    string p, t;
    cin >> p;
    male[p] = i;
    bname[i] = p;
    for ( int j = 0 ; j < n ; j++ ) {
      cin >> t;
      if ( !female.count(t) ) {
        gname[fit] = t;
         female[t] = fit++;
      favor[i][j] = female[t];
  }
  for ( int i = 0 ; i < n ; i++ ) {</pre>
    string p, t;
```

```
cin >> p;
    for ( int j = 0 ; j < n ; j++ ) {
        cin >> t;
        order[female[p]][male[t]] = j;
    }
}
initialize();
stable_marriage();

for ( int i = 0 ; i < n ; i++ ) {
    cout << bname[i] << " " << gname[favor[i][current[i ] - 1]] << endl;
}
</pre>
```

9.5 矩陣乘法

9.6 c++ 小抄

```
//pbds tree
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
tree<int, null_type, less<int>, rb_tree_tag,
   tree_order_statistics_node_update> tr;
tree::find_by_order(k) // O(logN) 取得第k大的元素
tree::order_of_key(ele) // O(LogN) 得到ele是tree中第幾
    大(有幾個元素小於ele)
//pbds pair priority_queue
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/priority_queue.hpp>
using namespace __gnu_pbds;
priority_queue<int, less<int>, pairing_heap_tag> pq;
pq.push(x);
pq.pop();
pq.top();
pq.join(b);
pq.empty();
pa.size():
pq.modify(it,6);
pq.erase(it);
//builtin functions
__builtin_popcount(x); // 1的個數
__builtin_popcountll(x); // for long long
                        // 前導0的個數
__builtin_clz(x);
                        // 後導0的個數
__builtin_ctz(x);
                        // 奇偶性
__builtin_parity(x);
//溢位檢查
ret = __builtin_add_overflow(a, b, &res) // if ret = 1
   a+b 溢位
ret = __builtin_sub_overflow(a, b, &res) // if ret = 1
   a-b 溢位
ret = __builtin_mul_overflow(a, b, &res) // if ret = 1
   a*b 溢位
```

```
_builtin_add_overflow_p(a, b, 0LL) // if ret = 1
     溢位 第三個參數是判斷的類型
//vector SIMD
typedef int v4si __attribute__ ((vector_size(4 * sizeof
    (int))));
//大質數表
{1000000007, 1000000009, 1000000021, 1000000033
    1000000087, 1000000093, 1000000097, 1000000123,
    1000000321};
//mt19937
#include <random>
#include <chrono>
int getRendom(int 1, int r) {
    static auto seed = std::chrono::system_clock::now()
         .time_since_epoch().count();
    static std::mt19937 gen(seed);
    std::uniform_int_distribution<int> dis(l, r);
    return dis(gen);
//sorted vector 去重
vec.erase(unique(vec.begin(), vec.end()), vec.end());
//std::valarray
valarray<int> a(初始值,數量);
valarray<<mark>int</mark>> a(10);
valarray<<mark>int</mark>> b(10);
valarray<int> c = a + b;
valarray<int> d = a * b;
valarray<int> e = a + 10;
valarray<<mark>int</mark>> f = a * 10;
valarray<int> g = a.cshift(1); //循環左移
valarray<bool> equal = a == b;
int sum = a.sum();
int max = a.max();
int min = a.min();
std::valarray<int> g = a.apply([](int x) { return x * x
//regex
#include <regex>
using namespace std;
bool res = regex_match("abc", regex("a.c"));
bool res = regex_match("abc", regex("A.c", regex::icase
    )); //忽略大小寫
```

9.7 python 小抄

```
#!/usr/bin/env python3
#帕斯卡三角形
n = 10
dp = [ [1 for j in range(n)] for i in range(n) ]
for i in range(1,n):
    for j in range(1,n):
        dp[i][j] = dp[i][j-1] + dp[i-1][j]
for i in range(n):
    print( '
             '.join( '{:5d}'.format(x) for x in dp[i] )
# EOF1
while True:
    try:
        n, m = map(int, input().split())
    except:
        break
# EOF2
import sys
for s in sys.stdin:
    print(eval(s.replace("/", "//")))
# input a sequence of number
a = [ int(x) for x in input().split() ]
```

```
ncase = int( input() )
for _ in range(ncase):
   n, m = [int(x) for x in input().split()]
a, b = "$"+input(), "$"+input()
    dp = [ [int(0) for j in range(m+1)] for i in range(
       n+1) ]
    for i in range(1,n+1):
        for j in range(1,m+1):
            dp[i][j] = max(dp[i-1][j],dp[i][j-1])
            if a[i]==b[j]:
                dp[i][j] = max(dp[i][j],dp[i-1][j-1]+1)
   for i in range(1,n+1):
        print(dp[i][1:])
    print('a=\{:s\}, b=\{:s\}, |LCS(a,b)|=\{:d\}'.format(a
        [1:],b[1:],dp[n][m]))
# list, dict, string
a = [1, 3, 4, 65, 65]
b = list.copy() # b = [1, 3, 4, 65], list a 跟 llst b
    互相獨立
cnt = list.count(65) # cnt == 2
loc = list.index(65) # loc == 3, find the leftmost
    element, if not found then return ERROR
list.sort(reverse = True|False, key = none|lambda x:x
    [1]) # list.sort has side effect but no reture
    value
                # C++
# stack
stack = [3,4,5]
stack.append(6) # push()
stack.pop()
              # pop()
stack[-1]
                # top()
len(stack)
               # size() 0(1)
# queue
                # C++
from collections import deque
queue = deque([3,4,5])
queue.append(6) # push()
queue.popleft() # pop()
queue[0]
              # front()
len(queue)
                # size() 0(1)
```

10 Persistence