

# Problem 4

$$\begin{aligned} a) \quad (P_1 \wedge \dots \wedge P_m) \Rightarrow Q &\equiv \neg(P_1 \wedge \dots \wedge P_m) \vee Q \quad (\text{implication elimination: } (\alpha \Rightarrow \beta) \equiv (\neg \alpha \vee \beta)) \\ &\equiv (\neg P_1 \vee \dots \vee \neg P_m \vee Q) \quad (\text{associativity of } \vee) \\ &\equiv (\neg P_1 \vee \dots \vee \neg P_m \vee Q) \end{aligned}$$

As a result, we can conclude  $(P_1 \wedge \dots \wedge P_m) \Rightarrow Q$  is equivalent to  $(\neg P_1 \vee \dots \vee \neg P_m \vee Q)$

$$\begin{aligned} b) \quad (P_1 \wedge \dots \wedge P_m) \Rightarrow (Q_1 \vee \dots \vee Q_n) &\equiv \neg(P_1 \wedge \dots \wedge P_m) \vee (Q_1 \vee \dots \vee Q_n) \\ &\equiv (\neg P_1 \vee \dots \vee \neg P_m) \vee (Q_1 \vee \dots \vee Q_n) \\ &\equiv (\neg P_1 \vee \dots \vee \neg P_m \vee Q_1 \vee \dots \vee Q_n) \end{aligned}$$

(associativity of  $\vee$ :  
 $(\alpha \vee \beta) \vee \gamma \equiv (\alpha \vee (\beta \vee \gamma))$ )

With help from the derivation of 4a and the associativity of  $\vee$ , we see that  $(P_1 \wedge \dots \wedge P_m) \Rightarrow (Q_1 \vee \dots \vee Q_n)$  is equivalent.

$$c) \text{Lother} \equiv (l_1 \vee \dots \vee l_{i-1} \vee l_{i+1} \vee \dots \vee l_K)$$

$$\text{Mother} \equiv (m_1 \vee \dots \vee m_{j-1} \vee m_{j+1} \vee \dots \vee m_n)$$

$$l_i \equiv \neg m_j$$

$$(l_1 \vee \dots \vee l_K) \equiv (l_i \vee \text{Lother}) \equiv \text{true}$$

$$(m_1 \vee \dots \vee m_n) \equiv (m_j \vee \text{Mother}) \equiv (\neg l_i \vee \text{Mother}) \equiv \text{true}$$

$$\frac{(l_i \vee \text{Lother}), (m_j \vee \text{Mother})}{\text{Lother} \vee \text{Mother}} \rightarrow (l_i \vee \text{Lother}) \wedge (\neg l_i \vee \text{Mother}) \Rightarrow (\text{Lother} \vee \text{Mother})$$

Since we know that  $(l_i \vee \text{Lother})$  is true and given, and  $(\neg l_i \vee \text{Mother})$  is also true and given. If  $l_i$  is true, then  $m_j$  is false. This would mean that ~~Lother~~ Mother needs to be true for the statement  $(l_i \vee \text{Lother}) \wedge (\neg l_i \vee \text{Mother})$  to be true.

↑	↑	↑	↑
T	unknown	false	must be true

This would mean that  $(\text{Lother} \vee \text{Mother})$  is true.

If  $m_j$  is true,  $l_i$  must be false. This would mean that Lother needs to be true for

the statement  $(l_i \vee \text{Lother}) \wedge (\neg l_i \vee \text{Mother})$  to be true. This would mean that

↑	↑	↑	↑
false	true	true	unknown

$(\text{Lother} \vee \text{Mother})$  would be true.