1. Template

```
//I C O N I C
#include <bits/stdc++.h>
#define endl '\n'
#define ll long long int
#define ld long double
#define ff __float128
#define fore(i,a,b) for (int i = a; i < b; i++)
#define fi first
#define se second
#define pb push_back
#define all(v) v.begin(), v.end()
#define fast_io ios_base::sync_with_stdio(0);cin.tie(0);cout.tie(0);
using namespace std;
typedef pair<int,int> pii;
typedef pair<11,11> pll;
typedef vector<int> vi;
typedef vector<vi> vvi;
typedef vector<ll> vll;
typedef vector<vll> vvll;
typedef vector<pll> vpll;
const int inf = 1 << 30;
const int mod = 1e9+7;
// clear && q++ -std=c++17 -O2 -Wall template.cpp -o template &&
\rightarrow ./template
// ifstream cin("input.txt"); ofstream cout("output.txt");
```

2. Data Structure

2.1. Segment tree with lazy

```
struct Node{
  int 1, r;
  ll sum;
  int mark;
  ll lazy;
};
//Neutros: (+, 0) (*, 1) (gcd, 0) (lcm, 1) (min, inf) (max, -inf)
const int Neutro = 0;
```

```
template<typename TT> struct SegmentTree{
 int n, h;
 vector<Node> st;
 SegmentTree (int m, vector<TT> &values) : n(m){
   h = 1 << ((int)(ceil(log2(n)) + 1));
   st.resize( h );
   build(1, 1, n, values);
 TT merge(TT 1, TT r){ return 1 + r; }//for query//!CHANGE
 TT getValue(int curr){ return st[curr].sum; }//same^^
 int left(int n){ return (n << 1);}</pre>
 int right(int n){return (n << 1) | 1; }</pre>
 void initLeaf(int curr, TT value){//!CHANGE
   st[curr].mark = 0; st[curr].lazy = Neutro;
   st[curr].sum = value;//
 void updateFromChildren(int curr){//!CHANGE
   st[curr].sum = st[left(curr)].sum + st[right(curr)].sum;
 void updateNodeLazy(int curr, TT value){//updates lazy
   int 1 = st[curr].1, r = st[curr].r;
   st[curr].sum += (r - l + 1) * value; //!CHANGE
   st[curr].mark = 1; st[curr].lazy += value;//
 void propagateToChildren(int curr){//propagate lazy
   if(st[curr].mark == 0) return;
   updateNodeLazy(left(curr), st[curr].lazy );
   updateNodeLazy(right(curr), st[curr].lazy);
   st[curr].mark = 0; st[curr].lazy = Neutro;
 void build(int curr, int 1, int r, vector<TT> &values){
   st[curr].1 = 1; st[curr].r = r;
   if(1 == r) {
     initLeaf(curr, values[1]);//
   }else{
     int m = ((r - 1) >> 1) + 1;
     build(left(curr), 1, m, values);
     build(right(curr), m + 1, r, values);
     updateFromChildren(curr);
   }
 void rangeUpdate(int curr, int 1, int r, int q1, int qr, TT value){
   if( r < ql || qr < l ) return;
   else if( ql \le l \& r \le qr){
     updateNodeLazy(curr, value);
```

```
}else{
      propagateToChildren(curr);
      int m = ((r - 1) >> 1) + 1;
      rangeUpdate(left(curr), 1, m, q1, qr, value);
      rangeUpdate(right(curr), m + 1, r, ql, qr, value);
      updateFromChildren(curr);
 }
 // not lazy: //!Coment up or down
 void pointUpdate(int curr, int 1, int r, int pos, TT value){
    if(1 == r){}
      st[curr].sum += value;
    }else{
      int m = ((r - 1) >> 1) + 1;
      if(pos <= m) pointUpdate(left(curr), 1, m, pos, value);</pre>
      else pointUpdate(right(curr), m + 1, r, pos, value);
      updateFromChildren(curr);
   }
 }
 TT rangeQuery(int curr, int 1, int r, int q1, int qr){
    if( r < ql || qr < l ) return Neutro;</pre>
    else if( ql <= l \&\& r <= qr){
     return getValue(curr);
    }else{
      propagateToChildren(curr);//!CHANGE
      int m = ((r - 1) >> 1) + 1;
      return merge( rangeQuery(left(curr), 1, m, q1, qr),

→ rangeQuery(right(curr), m+1, r, ql, qr));
    }
 }
 void update( int ql, int qr, int value){
    rangeUpdate(1, 1, n, ql, qr, value);
    // pointUpdate(1, 1, n, pos, value);//!CHANGE
 }
 TT query(int ql, int qr){
    return rangeQuery(1, 1, n, ql, qr);
 }
 void printST(){
    cout << endl << "st = ";
    fore(i,0,h) cout << st[i].sum << ' '; cout << endl;</pre>
 }
};
// vector<TYPE> nums(n + 1); //1 indexed
// SegmentTree<TYPE> st(n,nums);// both must be same TYPE
// st.update(l,r,x); st.query(l, r);
```

2.2. Segment tree Geometric sum

```
/* If you have an array [0, 0, 0, 0, 0, 0]
    update(l, r) adds x, x^2, x^3, x^4 starting from l
    Example:
                  x = 2
    update(2, 5) -> [0, 2, 4, 8, 16, 0]
    query(l, r) \rightarrow sum(arr[l], arr[l+1], ..., arr[r]) */
struct Node {
    int 1 = 0, r = 0;
    11i 1s = 0, rs = 0;
    bool flagLazy = false;
    lli lazy_ls = 0, lazy_rs = 0;
    Node() {}
    Node(int 1, int r) : 1(1), r(r) {}
    // Combine 2 nodes
    Node operator+(const Node &b) {
        Node res(1, b.r);
        res.ls = (ls + b.ls) \% MOD;
        res.rs = (rs + b.rs) \% MOD;
        return res;
    // Update range
    void updateNode(lli sum_l, lli sum_r) {
        ls = (ls + sum_l) % MOD; rs = (rs + sum_r) % MOD;
        lazy_ls = (lazy_ls + sum_l) % MOD;
        lazy_rs = (lazy_rs + sum_r) % MOD;
        flagLazy = true;
    }
    void resetLazy() {
        flagLazy = false;
        lazv_ls = 0; lazv_rs = 0;
};
struct SegmentTree {
    vector<Node> ST;
    int N;
    lli x, x_inv;
    SegmentTree(int n, lli x) : N(n), x(x) {
        x_{inv} = powerMod(x, MOD - 2);
        ST.resize(4 * N); build(1, 1, N);
    void build(int curr, int 1, int r) {
        ST[curr].1 = 1, ST[curr].r = r;
        if (1 == r) return;
        int mid = 1 + (r - 1) / 2;
        build(2 * curr, 1, mid);
        build(2 * curr + 1, mid + 1, r);
```

```
}
                                                                                     return (ans.ls - ans.rs + MOD) * powerMod(1 - x + MOD, MOD - 2) %
   void pushToChildren(int curr) {
                                                                                MOD;
       if (ST[curr].flagLazy) {
                                                                                 }
                                                                             };
           int size_child_left = (ST[2 * curr].r - ST[2 * curr].l + 1);
           int size_child_right =
                (ST[2 * curr + 1].r - ST[2 * curr + 1].l + 1);
                                                                             2.3. Implicit Treap
           lli sum_r_to_left =
               ST[curr].lazy_rs * powerMod(x_inv, size_child_right) %
\hookrightarrow MOD;
                                                                             /*----Implicit Treap 1 indexed----*/
           lli sum_l_to_right =
                                                                             struct Node {
                ST[curr].lazy_ls * powerMod(x, size_child_left) % MOD;
                                                                                 int key, priority, size, rev, value;
           ST[2 * curr].updateNode(ST[curr].lazy_ls, sum_r_to_left);
                                                                                 Node *left, *right;
           ST[2 * curr + 1].updateNode(sum_l_to_right,
                                                                                 Node(int key, int priority)
   ST[curr].lazy_rs);
                                                                                     : key(key), priority(priority), size(1), value(key), left(NULL),
           ST[curr].resetLazy();
                                                                                       right(NULL) {}
       }
                                                                             };
   }
                                                                             // pushes the lazy updates
   // UPDATE
                                                                             void push(Node *T) {
   void update(int curr, int 1, int r, int ql, int qr, int start) {
                                                                                 if (T) {
       if (1 > qr || r < ql) return;
                                                                                     if (T->rev) {
       else if (ql <= l && r <= qr) {
                                                                                          swap(T->left, T->right);
           int offset_l = l - start + 1;
                                                                                          if (T->left) T->left->rev ^= 1;
           int offset_r = offset_l + (r - l + 1) - 1;
                                                                                          if (T->right) T->right->rev ^= 1;
           ST[curr].updateNode(powerMod(x, offset_1),
                                                                                          T \rightarrow rev = 0;
                                powerMod(x, offset_r + 1));
                                                                                     }
                                                                                 }
           return;
       pushToChildren(curr);
                                                                             // modify this funciton depending on the guery that you want
       int mid = 1 + (r - 1) / 2;
                                                                             void update(Node *T) {//updateFromChildren
       update(2 * curr, 1, mid, q1, qr, start);
                                                                                 // push(T);
       update(2 * curr + 1, mid + 1, r, ql, qr, start);
                                                                                 if (T) {
       ST[curr] = ST[2 * curr] + ST[2 * curr + 1];
                                                                                     T->size = 1; T->value = T->key;
                                                                                     if (T->left) {
   void update(int ql, int qr, int start) { update(1, 1, N, ql, qr,
                                                                                          T->size += T->left->size;

    start); }

                                                                                   T->value += T->left->value;
   // QUERY
                                                                                     }
   Node query(int curr, int 1, int r, int q1, int qr) {
                                                                                     if (T->right) {
       if (l > qr || r < ql) return Node();</pre>
                                                                                          T->size += T->right->size;
       if (ql <= 1 && r <= qr) return ST[curr];</pre>
                                                                                   T->value += T->right->value;
       else {
           pushToChildren(curr);
                                                                                 }
           int mid = 1 + (r - 1) / 2;
                                                                             }
           return query(2 * curr, 1, mid, q1, qr) +
                   query(2 * curr + 1, mid + 1, r, ql, qr);
                                                                             // returns the root of the union of treaps T1 and T2
       }
                                                                             Node *merge(Node *T1, Node *T2) {
   }
                                                                                 push(T1), push(T2);
   lli query(int ql, int qr) {
                                                                                 if (T1 == NULL) return T2;
       auto ans = query(1, 1, N, ql, qr);
                                                                                 if (T2 == NULL) return T1;
```

```
if (T1->priority > T2->priority) {
                                                                                  Node *T = NULL;
        T1->right = merge(T1->right, T2);
                                                                                  fore(i, a.size()) T = insert(T, a[i], i + 1);
        update(T1); return T1;
   } else {
                                                                              }
        T2->left = merge(T1, T2->left);
                                                                              // erases the key x from the treap T
        update(T2); return T2;
                                                                              Node *erase(Node *T, int ind) {
    }
                                                                                  Node *L, *R, *AUX;
}
                                                                                  tie(L, R) = split(T, ind + 1);
// returns the roots the division of the treap T with parameter x
                                                                                  tie(L, AUX) = split(L, ind);
// T1 contains all nodes from [1,x), T2 contains all nodes from
                                                                                  return merge(L, R);
\hookrightarrow [x, T->size]
pair<Node *, Node *> split(Node *T, int x) {
                                                                              //returns the sum of all the keys in te range [l,r]
                                                                              int query(Node* T, int 1, int r){
    push(T);
    if (T == NULL) return {NULL, NULL};
                                                                                Node *L,*X,*R;
                                                                                tie(L,X) = split(T,1);
    int index;
    if (T->left) index = T->left->size + 1;
                                                                                tie(X,R) = split(X,r-1+2);
    else index = 1;
                                                                                int ans = (X) ? X \rightarrow value : 0;
    if (index < x) {
                                                                                X = merge(L,X); T = merge(X,R);
        Node *T1, *T2;
                                                                                return ans;
        tie(T1, T2) = split(T->right, x - index);
        T->right = T1;
        update(T); update(T2);
        return {T, T2};
                                                                              // returns the k-th node of the treap
    } else {
                                                                              Node *findKth(Node *T, int k) {
        Node *T1, *T2;
                                                                                  if (T == NULL) return NULL;
        tie(T1, T2) = split(T->left, x);
                                                                                  int ind = T->left ? T->left->size + 1 : 1;
        T\rightarrow left = T2;
                                                                                   if (ind > k) return findKth(T->left, k);
        update(T1); update(T);
                                                                                  else if (ind == k) return T;
        return {T1, T};
                                                                                  else return findKth(T->right, k - ind);
}
                                                                              // reverser the treap from [l,r]
                                                                              Node *reverse(Node *T, int 1, int r) {
// seed to generate random numbers
mt19937 rng(chrono::steady_clock::now().time_since_epoch().count());
                                                                                  Node *L, *R, *AUX;
// generates random numbers
                                                                                  tie(L, R) = split(T, r + 1);
uniform_int_distribution<> dis(numeric_limits<int>::min(),
                                                                                  tie(L, AUX) = split(L, 1);
                               numeric_limits<int>::max());
                                                                                  if (AUX) AUX->rev ^= 1;
// inserts a new node with key = x in the position ind of the treap T,
                                                                                  L = merge(L, AUX); T = merge(L, R);
\hookrightarrow returns
                                                                                  return T;
// the new root
Node *insert(Node *T, int x, int ind) {
                                                                              // prints the Treap inorder
    Node *TN = new Node(x, dis(rng));
                                                                              void inorder(Node *T) {
                                                                                  if (T) {
    Node *L, *R;
    tie(L, R) = split(T, ind);
                                                                                       push(T);
    TN = merge(L, TN); TN = merge(TN, R);
                                                                                       inorder(T->left);
    return TN:
                                                                                       cout << T->key << " ";
                                                                                       inorder(T->right);
// returns a treap with all the elements in a
                                                                                  }
Node *create(vi a) {
                                                                              }
```

2.4. Policy base

```
// 20
#include <ext/pb_ds/assoc_container.hpp> // Common file
#include <ext/pb_ds/tree_policy.hpp> // Including
\rightarrow tree_order_statistics_node_update
using namespace __gnu_pbds;
template <class T>
using ordered_set =
    tree<T, null_type, less<T>, rb_tree_tag,

→ tree_order_statistics_node_update>;
ordered_set<int> s;
*s.find_by_order(k); // To get the i-th element, O-indexed
s.order_of_key(k); // The number of items strictly smaller than n
                   // For multiset, use T = pair
                   // in second parameter a global counter
                   // s.order_of_key({k, -INF}) will return
                   // how many elements < k
```

2.5. SQRT and MO's

```
//SQRT decomposition
//if RTE, change limits to min(br, n)
template <typename TT>
struct SQRT{
 int n, s;
 TT neutro = 0;
 vector<TT> A, B;
 vector<TT> lazy, marks;
 SQRT(int m, vector<TT> &arr): n(m){
    s = sqrt(n) + 1; //puede variar
    A.assign(n, neutro);
    B.assign(n / s + 1, neutro);
    lazy.assign(s, neutro); marks.assign(s, neutro);
    fore(i,0,n){ A[i] = arr[i]; B[i/s] += arr[i]; }
 void pushLazy(int block){
    if(marks[block]){
      fore(i,block,(block+1) * s && i < n) A[i] += lazy[block];</pre>
      lazy[block] = neutro; marks[block] = 0;
    }
 }
 void rangeUpdate(int 1, int r, TT value){
```

```
int bl = 1/s, br = r/s;
    if(bl == br){}
      pushLazy(bl);
      fore(i,1,r+1) A[i] += value;
      TT res = neutro;
      fore(i, bl*s, (bl+1) * s && i < n) res += A[i];
      B[bl] = res;
    }else{
      pushLazy(bl);pushLazy(br);
      fore(i,1,(bl+1) * s){ A[i] += value; B[bl] += value;}
      fore(i,bl+1, br) { B[i] += s * value; lazy[i] += value; marks[i]
      fore(i,br * s, r+1) { A[i] += value; B[br] += value;}
  void pointUpdate(int idx, TT value){//not lazy
    int block = idx / s;
    A[idx] = value;
    TT res = neutro:
    fore(i, block * s, (block + 1) * s && i < n) res += A[i];
    B[block] = res;
  TT rangeQuery(int 1, int r){
    int bl = 1/s, br = r/s;
    TT res = 0;
    if(bl == br){}
      pushLazy(bl);
      fore(i,1,r+1) res += value;
    }else{
      pushLazy(bl); pushLazy(br);
      fore(i,1,(bl+1) * s) res += A[i];
      fore(i,bl+1, br)
                           res += B[i];
      fore(i,br * s, r+1) res += A[i];
    }
 }
};
//MO's algorithm
//arr\ 1 or 0 indexed, S = sqrt(n), index = index of the query
ll answer, neutro = 0; vll arr;
struct MOquery{
  int 1, r, index, S;
  MOquery(int 1, int r, int idx, int S): 1(1), r(r), index(idx), S(S){}
  bool operator<(const MOquery & q) const{</pre>
    int bl = 1 / S, bq = q.1 / S;
    if(bl == bq && bl & 1) return r > q.r;//a little bit faster
    if(bl == bq) return r < q.r;</pre>
    return bl < bq;
```

```
};
void add(int idx){
  answer += arr[idx];
void remove(int idx){
  answer -= arr[idx]:
vector<11> MO(vector<MOquery> & queries){
 vector<ll> ans(queries.size());
 sort(queries.begin(), queries.end());
 11 current = 0;
 int prevL = 0, prevR = -1;
 int i, j;
 answer = neutro;
 for(const MOquery & q : queries){
    while (prevL > q.1) { prevL--; add(prevL); }
    while (prevR < q.r) { prevR++; add(prevR); }</pre>
    while (prevL < q.1) { remove(prevL); prevL++; }</pre>
    while (prevR > q.r) { remove(prevR); prevR--;}
    ans[q.index] = answer;
 }
 return ans;
```

2.6. Convex Hull Trick

```
const int MX = 200005;
const ll inf = 1e18;
bool Q = 0;
struct Line {
 mutable ll m, b, x;
 // Maximo: m < ot.m
 // Minimo: m > ot.m
 bool operator < (const Line ot) const {</pre>
    return Q ? x < ot.x : m < ot.m;
 }
};
ll ceil (ll a, ll b) {
 if (a < 0 != b < 0) return a / b;
 return (abs(a) + abs(b) - 1) / abs(b);
}
ll intersection (const Line &p, const Line &q) {
 return ceil(q.b - p.b, p.m - q.m);
}
```

```
struct Hull : multiset<Line> {
  bool valid (auto it) {
    if (it == begin()) {
      auto sig = it; sig++;
      if (sig != end()) sig->x = intersection(*it, *sig);
      return it->x = -inf;
    auto ant = it, sig = it;
    ant--, sig++;
    if (sig == end()) {
      it->x = intersection(*it, *ant);
      return 1;
    }
    11 x = intersection(*it, *ant);
    11 y = intersection(*it, *sig);
    if (x > y) return 0;
    it->x = x, sig->x = y;
    return 1;
  void add (ll m, ll b) {
    auto it = lower_bound({m, b, -inf});
    if (it != end() && it->m == m) {
      //Maximo: it->b > b
      //Minimo: it->b < b
      if (it->b > b) return;
      it->b = b:
    } else { it = insert({m, b, -inf}); }
    if (!valid(it)) { erase(it); return; }
    auto ant = it;
    while (ant != begin()) {
      if (valid(--ant)) break;
      erase(ant);
      if (it == begin()) { it->x = -inf; break; }
      ant = it;
    }
    auto sig = it; sig++;
    while (sig != end() && !valid(sig)) erase(sig++);
  11 query (11 x) {
    if (empty()) return 0;
    Q = 1; auto it = upper_bound(\{0, 0, x\});
    it--;
    Q = 0; return x * it->m + it->b;
};
```

2.7. BIT

```
struct Fenwick {
    int n;
    vector<long long> tree;
    Fenwick(int _n): n(_n), tree(n + 1, 0) {}
    void update(int idx, long long val) {
       for (; idx <= n; idx += idx & -idx) {
            tree[idx] += val;
       }
    }
    long long query(int idx) {
       long long ret = 0;
       for (; idx > 0; idx -= idx & -idx) {
            ret += tree[idx];
       }
       return ret;
    }
    long long query(int x, int y) { return query(y) - query(x - 1); }
};
```

2.8. merge sort tree

```
vi tree[400000];
vi vv;
void build(int a[], int v, int tl, int tr) {
    if (tl == tr) {
        tree[v] = vi(1, a[tl]);
    } else {
        int tm = (tl + tr) / 2;
        build(a, v*2, t1, tm);
        build(a, v*2+1, tm+1, tr);
        merge(tree[v*2].begin(), tree[v*2].end(), tree[v*2+1].begin(),

    tree[v*2+1].end(), back_inserter(tree[v]));
}
void query(int v, int tl, int tr, int l, int r, int x){
    if(1 > r) return;
    if(tr \ll r){
    for(auto i = tree[v].begin(); i < tree[v].end(); i++){</pre>
      if(*i <= x) vv.pb(*i);
      else break;
    }
```

```
return;
}
if(t1 > r) return;
int tm = (t1 + tr) / 2;
query(v*2, t1, tm, 1, r, x),
query(v*2+1, tm+1, tr, 1, r, x);
return;
}
```

3. Binary Search

```
//lower(x): primer elemento mayor o iqual a x
//upper(x): primer elemento mayor a x
int lowerBound(vi &nums, int a) {
    int l = 0, r = nums.size() - 1;
    while(1 \le r) {
        int m = ((r - 1) >> 1) + 1;
  // if(nums[m] == a) return m;//binary
    nums[m] < a ? 1 = m + 1 : r = m - 1; //lower & binary
  // nums[m] \le a ? l = m + 1 : r = m - 1; //upper
    return 1;//return -1; //binary
}
for(int j = 0; j < 300; j++){
  1d \ mid1 = 1 + (r - 1) / 3, \ mid2 = r - (r - 1) / 3;
  ld f1 = f(p0, Friend1[i + 1], p1, Friend2[i + 1], mid1);
  ld f2 = f(p0, Friend1[i + 1], p1, Friend2[i + 1], mid2);
  if(f1 >= f2) 1 = mid1;
  else r = mid2;
```

4. Flujos

4.0.1. Dinic

```
typedef vector<pii> vpii;
typedef tuple<int, ll, ll> edge;
class max_flow{
  private:
    int V;
    vector<edge> EL;
    vvi AL;
    vi d, last;
```

```
vpii p;
  bool BFS(int s, int t){
    d.assign(V, -1);d[s] = 0;
    queue<int> q({s});
    p.assign(V, {-1, -1});
    while( !q.empty()){
      int u = q.front(); q.pop();
      if( u== t) break;
      for(auto &idx: AL[u]){
        auto &[v, cap, flow] = EL[idx];
        if( (cap - flow > 0) && d[v] == -1){
          d[v] = d[u] + 1, q.push(v), p[v] = {u, idx};
        }
      }
    }
    return d[t] != -1;
  // ll send_one_flow(int s, int t, ll f = inf){
  // if (s == t) return f;
       auto \, \mathfrak{C}[u, idx] = p[t];
       auto \&cap = qet<1>(EL[idx]), &flow = qet<2>(EL[idx]);
       ll pushed = send_one_flow(s, u, min(f, cap-flow));
       flow += pushed;
  // return pushed;
  // }
  11 DFS(int u, int t, 11 f = inf){
    if( (u == t) || (f == 0)) return f;
    for(int &i = last[u]; i < (int) AL[u].size(); ++i){</pre>
      auto &[v, cap, flow] = EL[AL[u][i]];
      if(d[v] != d[u] + 1) continue;
      if(ll pushed = DFS(v, t, min(f, cap - flow))){
        flow += pushed;
        auto &rflow = get<2> (EL[AL[u][i] ^ 1]);
        rflow -= pushed;
        return pushed;
      }
    }
    return 0;
public:
  max_flow(int initialV) : V(initialV){
    EL.clear();
    AL.assign(V, vi());
  void add_edge(int u, int v, ll w, bool directed = true){
    if( u == v) return;
    EL.emplace_back(v, w, 0);
    AL[u].pb(EL.size() - 1);
```

```
EL.emplace_back(u, directed ? 0 : w, 0);
     AL[v].pb(EL.size() - 1);
   }
   // ll edmonds_karp(int s, int t){
   // ll mf = 0:
        while (BFS(s,t)) {
           ll f = send\_one\_flow(s, t);
           if (f == 0)break;
   //
           mf += f;
        }
   //
        return mf;
   // }
   11 dinic(int s, int t ){
     11 \text{ mf} = 0:
     while( BFS(s,t)){
       last.assign(V, 0);
       while( ll f = DFS(s,t)){
         mf += f;
       }
     }
     return mf;
   }
   vpii build(int s){
     vi visited(V, 0);
     dfs_ans(s, visited);
     vpii ans;
     //buscar aristas que van de un nodo u a uno v y cumplen alguna
\hookrightarrow condicion
     fore(u,0,V){
       fore(j,0,AL[u].size()){
         auto &[v, cap, flow] = EL[AL[u][j]];
         if(cap == flow && visited[u] && !visited[v]){
            ans.pb(\{u,v\});
       }
     }
     return ans;
   //Realizar una dfs desde s en la red residual utilizando
   //sólo aristas no saturadas (con flujo menor a la capacidad).
   void dfs_ans(int current, vi &visited){
     visited[current] = 1;
     // cout <<"entra " << current << endl;</pre>
     for(auto idx: AL[current]){
        auto &[v, cap, flow] = EL[idx];
        if(!visited[v] && flow < cap) dfs_ans(v, visited);</pre>
     }
```

```
template <typename T = int> struct MinCostFlow {
      return;
    }
                                                                                  vector<vector<Edge<T> *>> adjList;
};
                                                                                  int N;
                                                                                  MinCostFlow(int N) : N(N) { adjList.resize(N); }
4.0.2. Ford-Fulkerson
                                                                                  void addEdge(int u, int v, T capacity, T cost, int idx) {
const int sink = 37;
                                                                                      Edge<T> *uv = new Edge<T>(v, capacity, cost, idx);
int C[50][50], F[50][50], visited[50];
                                                                                      Edge<T> *vu = new Edge<T>(u, 0, -cost, idx, true);
int sendFlow(int node, int bottleneck){
                                                                                      uv->res = vu;
 if(node == sink){
                                                                                      vu->res = uv:
    return bottleneck;
                                                                                      adjList[u].push_back(uv);
 }
                                                                                      adjList[v].push_back(vu);
 visited[node] = true;
                                                                                  }
 fore(i,0,sink+1){
    int f = C[node][i] - F[node][i];
                                                                                  pair<T, T> getMinCostFlow(int s, int t) {
    if(f>0 && !visited[i]){
                                                                                      vector<T> dist(N), cap(N);
      f = sendFlow(i, min(f, bottleneck));
                                                                                      vector<bool> inQueue(N);
      if(!f) continue;
                                                                                      vector<Edge<T> *> parent(N);
      F[node][i] += f;
      F[i][node] -= f;
                                                                                      T minCost = 0, maxFlow = 0;
      return f;
   }
                                                                                      while (true) {
 }
 return 0;
                                                                                          fill(all(dist), INF_COST);
                                                                                          fill(all(cap), 0);
                                                                                          fill(all(parent), nullptr);
4.0.3. maxflow mincost
                                                                                          queue<int> q;
                                                                                          q.push(s);
template <typename T = int> struct Edge {
                                                                                          dist[s] = 0;
    int to;
                                                                                          cap[s] = INF_FLOW;
   T flow, capacity, cost;
    int idx;
                                                                                          while (!q.empty()) {
    bool rev;
                                                                                              int u = q.front();
    Edge *res;
                                                                                              q.pop();
    Edge(int to, T capacity, T cost, int idx, bool rev = false)
                                                                                              inQueue[u] = 0;
        : to(to), flow(0), capacity(capacity), cost(cost), idx(idx),
→ rev(rev) {}
                                                                                              for (auto E : adjList[u]) {
    void addFlow(T flow) {
                                                                                                  int v = E \rightarrow to;
        this->flow += flow;
                                                                                                   if (E->flow < E->capacity && dist[u] + E->cost <</pre>
        this->res->flow -= flow;
                                                                              \rightarrow dist[v]) {
};
                                                                                                       dist[v] = dist[u] + E->cost;
                                                                                                       cap[v] = min(cap[u], E->capacity - E->flow);
// MUCHO OJO CON ESTO
                                                                                                       parent[v] = E;
const lli INF_FLOW = 1e15, INF_COST = 1e9;
                                                                                                       if (!inQueue[v]) {
```

```
inQueue[v] = 1;
                               q.push(v);
                          }
                      }
                 }
             }
             if (!parent[t])
                  break;
             // if(dist[t] > 0) break;
             maxFlow += cap[t];
             minCost += cap[t] * dist[t];
             for (int u = t; u != s; u = parent[u] \rightarrow res \rightarrow to) {
                 parent[u] ->addFlow(cap[t]);
             }
        }
        return {minCost, maxFlow};
    }
};
```

5. Strings

5.1. aho-corasik

```
//#define feach(f, g) for(auto &f: g)
const int N=1e5+10, MOD=1e9+7, SIG=26;
int id=1, dp[N];
string s;
vector<int> adj[2*N];
struct node{
 int fail,ch[SIG]={};
 vector<int> lens;
}t[2*N];
void insert(string s){
 int u=1:
 for(auto &c: s){
    c-= 'a';
    if(!t[u].ch[c]) t[u].ch[c]=++id;
    u=t[u].ch[c];
 }
 t[u].lens.pb(s.size());
```

```
void dfs(int u){
  t[u].lens.insert(t[u].lens.end(), t[t[u].fail].lens.begin(),

    t[t[u].fail].lens.end());
  for(auto &v: adj) dfs(v);
void build(){
  queue<int> q;
  int u=1;
  t[1].fail=1;
  fore(i,0,SIG) {
    if(t[u].ch[i]) t[t[u].ch[i]].fail=u, q.push(t[u].ch[i]);
    else t[u].ch[i]=1;
  while(!q.empty()){
    u=q.front(); q.pop();
    fore(i,0,SIG){
      if(t[u].ch[i]) t[t[u].ch[i]].fail=t[t[u].fail].ch[i],

    q.push(t[u].ch[i]);
      else t[u].ch[i]=t[t[u].fail].ch[i];
  fore(i,2,id+1) adj[t[i].fail].pb(i);
  dfs(1);
}
```

5.2. Rolling Hash

```
// Code from: https://github.com/JorgeIba/Reference/
// requires lli = __int128, otherwise change the values of MOD, P ans PI
// (inverse of P)
struct Hash {
   lli MOD = 212345678987654321LL, P = 1777771, PI =
→ 106955741089659571LL;
   // lli \ MOD = 1e9 + 7, P = 31, PI = 1290322591l; // For lowercase
\hookrightarrow english
   // letters lli MOD = 1e9 + 7, P = 53, PI = 56603774ll; // For
   // lower/uppercase english letters lli MOD = 1824261409, P = 53, PI =
   // 1411221090;
   vector<lli> h, pi;
   Hash(string &s) {
        assert((P * PI) % MOD == 1);
       h.resize(s.size() + 1);
        pi.resize(s.size() + 1);
       h[0] = 0;
       pi[0] = 1;
       lli p = 1;
        forr(i, 1, s.size()) {
```

```
h[i] = (h[i-1] + p * s[i-1] % MOD) % MOD;
                                                                                 int n = 1;
            pi[i] = pi[i - 1] * PI % MOD;
                                                                                 while (n < a.size() + b.size() - 1) n <<= 1;
            p = p * P \% MOD;
                                                                               fa.resize(n); fb.resize(n);
       }
    }
                                                                               fft(fa, 1); fft(fb, 1);
    // returns the hash of substring [s,e)
                                                                               fore(i,0,n) fa[i] *= fb[i];
    lli get(int s, int e) { return ((h[e] - h[s] + MOD) % MOD) * pi[s] %
                                                                                 fft(fa, -1);
   MOD; }
};
                                                                                 vector<ll> result(n);
                                                                                 for (int i = 0; i < n; i++) result[i] = round(fa[i].real());</pre>
                                                                                 return result;
     Math
                                                                             /* if it's numbers and not polynomials, we have to normalise */
                                                                             void normalise(vll &ans) {
6.1. FFT
                                                                               int carry = 0;
                                                                               for (ll i = 0; i < ans.size(); ++i) {
                                                                                 ans[i] += carry;
using cd = complex<double>;
const double PI = acos(-1);
                                                                                 carry = ans[i] / 10;
                                                                                 ans[i] %= 10;
void fft(vector<cd> & a, int inv) {
                                                                               if(carry > 0) ans.pb(carry);
    int n = a.size();
                                                                               int t = ans.size();
                                                                               while(t > 0 && ans[t-1] == 0){
  for(int i = 1, j = 0; i < n - 1; ++i){
   for(int k = n >> 1; (j ^= k) < k; k >>= 1);
                                                                                 ans.pop_back();
   if(i < j) swap(a[i], a[j]);
                                                                                 t--;
 }
                                                                               if(ans.size() == 0) ans.push_back(0);
 vector<cd> w(n >> 1);
    for (int k = 2; k \le n; k \le 1) {
                                                                             6.2. NTT
       // cd w1 = polar(1.0, 2 * PI / k * inv);
       w[0] = 1;
       for (int j = 1; j < k >> 1; j++) // best precision but slower
                                                                             const int mod = 998244353, g = 3;
           w[j] = polar(1.0, 2 * j * PI / k * inv);
         // w[j] = w[j-1]*w1;
                                                                             lli powMod(lli a, lli b, lli mod){
       for (int i = 0; i < n; i += k) {
                                                                               lli ans = 1;
            for (int j = 0; j < k >> 1; j++) {
                                                                               b \%= mod-1;
                cd u = a[i + j], v = a[i + j + (k >> 1)] * w[j];
                                                                               if(b < 0) b+= mod-1;
                a[i + j] = u + v;
                                                                               while(b){
                a[i + j + (k >> 1)] = u - v;
                                                                                 if (b\&1) ans = ans * a % mod;
           }
                                                                                 a = a*a \mbox{\mbox{$mod$}};
       }
                                                                                 b>>=1;
    if (inv == -1) for (cd & x : a) x \neq n;
                                                                               return ans;
}
                                                                             lli inverse(lli a, lli mod){return powMod(a,mod-2, mod);}
```

template<int mod, int g>

vector<ll> multiply(vector<ll> const& a, vector<ll> const& b) {
 vector<cd> fa(a.begin(), a.end()), fb(b.begin(), b.end());

```
void ntt(vector<int> & X, int inv){
                                                                              EX (11 re = 0, 11 im = 0) : re(re), im(im){}
 int n = X.size();
 for(int i = 1, j = 0; i < n - 1; ++i){
                                                                              EX& operator = (EX oth) {
    for(int k = n >> 1; (j ^= k) < k; k >>= 1);
                                                                                 return re = oth.re, im = oth.im, *this;
    if(i < j) swap(X[i], X[j]);</pre>
 }
                                                                              int norm () const {
 vector<lli> wp(n>>1, 1);
                                                                                 return trim((111 * re * re - 111 * sq * im % mod * im) % mod);
 for(int k = 1; k < n; k <<= 1){
    lli wk = powMod(g, inv * (mod - 1) / (k<<1), mod);
                                                                              EX conj () const {
    for(int j = 1; j < k; ++j)
                                                                                 return {re, trim(-im)};
      wp[j] = wp[j - 1] * wk % mod;
    for(int i = 0; i < n; i += k << 1){
                                                                              EX operator * (EX ot) const {
      for(int j = 0; j < k; ++j){
                                                                                 return {
       int u = X[i + j], v = X[i + j + k] * wp[j] % mod;
                                                                                   int((111 * re * ot.re + 111 * sq * im % mod * ot.im) % mod),
       X[i + j] = u + v < mod ? u + v : u + v - mod;
                                                                                   int((111 * re * ot.im + 111 * im * ot.re) % mod)
       X[i + j + k] = u - v < 0 ? u - v + mod : u - v;
                                                                                };
     }
                                                                              };
   }
                                                                              EX& operator *= (const EX& ot) {
 }
                                                                                 *this = *this * ot; return *this;
 if(inv == -1){
   lli nrev = inverse(n, mod);
                                                                              EX operator * (ll k) const {
   for(int i = 0; i < n; ++i)
                                                                                 k = ((k \% mod) + mod) \% mod;
     X[i] = X[i] * nrev % mod;
                                                                                 return { (re * k) % mod, (im * k) % mod };
 }
}
                                                                              EX operator / (ll n) const {
                                                                                 return { re * inv(n) % mod, im * inv(n) % mod };
template<int mod, int g>
vector<int> multiply(vector<int> const& a, vector<int> const& b) {
                                                                              EX operator / (EX ot) const {
    vector<int> fa(a.begin(), a.end()), fb(b.begin(), b.end());
                                                                                 return *this * ot.conj() / ot.norm();
    int n = 1;
    while (n < a.size() + b.size() - 1) n <<= 1;
                                                                              EX& operator /= (const EX& ot) {
 fa.resize(n); fb.resize(n);
                                                                                 *this = *this / ot; return *this;
 ntt<mod, g>(fa, 1); ntt<mod, g>(fb, 1);
                                                                              EX operator + (EX ot) const {
 fore(i,0,n) fa[i] = (1LL * fa[i] * fb[i]) % mod;
                                                                                 return { trim(re + ot.re), trim(im + ot.im) };
    ntt < mod, g > (fa, -1);
                                                                              EX& operator += (const EX& ot) {
    return fa;
                                                                                 *this = *this + ot; return *this;
}
                                                                              EX operator - (EX ot) const {
                                                                                 return { trim(re - ot.re), trim(im - ot.im) };
6.3. Field extension
                                                                              EX& operator -= (const EX& ot) {
                                                                                 *this = *this - ot; return *this;
int sq = 5;
// const lli sqrt5 = 383008016;//mod1e9+9
const 11 mod = 1000000007;//is important to be CONST
                                                                              EX pow (11 p) const {
struct EX {
                                                                                 EX res(1), b = *this;
 ll re, im;
                                                                                 while (p) {
```

```
if (p & 1) res *= b; b *= b; p /= 2;
                                                                                  for (int i = 2; i * i <= n; ++i)
   }
                                                                                      if (n \% i == 0)
                                                                                          fact.push_back(i); while (n % i == 0) n /= i;
    return res;
 }
                                                                                  if (n > 1) fact.push_back(n);
                                                                                  for (int res = 2; res <= p; ++res) {
 bool operator == (EX ot) const {
                                                                                      bool ok = true:
    return re == ot.re && im == ot.im;
                                                                                      for (int factor : fact)
 }
                                                                                          if (powmod(res, phi / factor, p) == 1)
 bool operator != (EX ot) const {
                                                                                              ok = false; break;
   return !(*this == ot);
                                                                                      if (ok) return res;
                                                                                 }
 }
                                                                                  return -1;
  static ll trim(ll a) {
                                                                             }
    if (a \ge mod) a -= mod:
                                                                              vi rootK() \{// \text{ finds all numbers } x \text{ such that } x^k = a \pmod{n}
   if (a < 0) a += mod;
                                                                                  int n, k, a;
                                                                                  scanf("%d %d %d", &n, &k, &a);
    return a;
                                                                                  if (a == 0) return vi(1,1);
  static ll inv (ll b) {
                                                                                  int g = generator(n);
    11 \text{ res} = 1, p = mod - 2;
                                                                                  // Baby-step giant-step discrete logarithm algorithm
    while (p) {
                                                                                  int sq = (int) sqrt (n + .0) + 1;
      if (p & 1) res = 1ll * res * b % mod;
                                                                                  vector<pair<int, int>> dec(sq);
      b = 111 * b * b % mod;
                                                                                  for (int i = 1; i \le sq; ++i)
      p /= 2;
                                                                                      dec[i-1] = \{powmod(g, i * sq * k % (n - 1), n), i\};
                                                                                  sort(dec.begin(), dec.end());
                                                                                  int any_ans = -1;
    return res;
 };
                                                                                  for (int i = 0; i < sq; ++i) {
};
                                                                                      int my = powmod(g, i * k % (n - 1), n) * a % n;
                                                                                      auto it = lower_bound(dec.begin(), dec.end(), make_pair(my, 0));
                                                                                      if (it != dec.end() && it->first == my) {
6.4. nCr
                                                                                          any_ans = it->second * sq - i;break;
                                                                                      }
                                                                                 }
vll fact(200007,0), inv(200007,0), invfact(200007,0);
void factorial(){
                                                                                  if (any_ans == -1) return vi(1,-1);
 fact[0] = 1; inv[0] = inv[1] = 1; invfact[0] = 1;
                                                                                  // Print all possible answers
 fore(i,1,200005) fact[i] = (fact[i-1] * i) \% MOD;
                                                                                  int delta = (n-1) / gcd(k, n-1);
 fore(i,2,200005) inv[i] = inv[MOD % i] * (MOD - MOD / i) % MOD;
                                                                                  vi ans;
                                                                                  for (int cur = any_ans % delta; cur < n-1; cur += delta)
 fore(i,1,200005) invfact[i] = invfact[i-1] * inv[i] % MOD;
                                                                                      ans.push_back(powmod(g, cur, n));
ll ncr(ll n, ll r){
                                                                                  sort(ans.begin(), ans.end());
 return fact[n] * invfact[n - r] % MOD * invfact[r] % MOD ;
                                                                                  return ans;
                                                                             }
}
                                                                              6.6. Miller Rabin
6.5. Discrete Root
int generator(int p) {
                                                                              typedef __int128 i128;
    vector<int> fact;
                                                                              //powerModCode
```

bool singleTest(i128 a, i128 n) $\frac{1}{2k} = 1 \pmod{n}$

int phi = p-1, n = phi;

```
i128 \exp = n - 1;
 while (~exp & 1) exp >>= 1;// while exp is even
 if( powerMod(a, exp, n) == 1 ) return true; // a^{exp} = a^{(n-1)/2k}
\rightarrow ) = 1 (mod n)
 while (\exp < n - 1) {
    if( powerMod(a, exp, n) == n - 1 ) return true;
    exp <<= 1;
 }
 return false;
bool MillerRabin(lli n){
    if(n < 2) return false:
    if(n <= 3) return true:
    if( ~n & 1) return false;
    for(lli a: {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
        if(n == a) return true;
    if(!singleTest(a, n)) return false;
    return true; //Probability = 1 - (1/4) ^size_of(vector_a)
}
bool MillerRabinOther(lli n){
    if(n < 2) return false;
    if(n <= 3) return true;
    if( ~n & 1) return false;
    lli d = n-1, s = 0; //n-1 = 2^s *k
    for(;(^d\&1); d >>= 1, s++); //d = k
    for(lli a: {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
        if(n == a) return true;
        i128 residuo = powerMod(a, d, n);
        if(residuo == 1 or residuo == n-1) continue;
        lli x = s;
        while(--x){}
            residuo = (residuo * residuo) % n;
            if(residuo == n-1) break;
        }
        if(x==0) return false;
    return true; //Probability = 1 - (1/4) ^size_of(vector_a)
}
```

7. Graphs

7.1. bfs-dfs

```
void bfs(vector<vector<int>> &graph, int start){
 int n = graph.size();
 queue<int> q; q.push(start);
 vector<int> visited(n,0);
 while(!q.empty()){
   int current = q.front(); q.pop();
   for(auto next: graph[current])
     if(!visited[next])
       q.push(next);
 }
 return;
void dfs(int current, vector<vector<int>> &graph, vector<int> &visited){
 visited[current] = 1:
 for(auto next: graph[current])
   if(!visited[next]) dfs(next, graph, visited);
 return;
```

7.2. cicles

```
int hasCycleDirected(int current, vector<vector<int>> &graph, vector<int>
if(color[current] == 1) return 1;
 if(color[current] == 2) return 0;
 color[current] = 1;
 int ans = 0;
 for(int child: graph[current]){
   ans = ans | hasCycleDirected(child, graph, color);
 color[current] = 2;
 return ans;
int hasCycleUndirected(int current, int father, vector<vector<int>>
if(color[current] == 1) return 1;
 color[current] = 1;
 int ans = 0;
 for(int child: graph[current]) if(child != father)
   ans = ans | hasCycleUndirected(child, current, graph, color);
 return ans:
```

```
int isBipartiteDFS(int current, vector<vector<int>> &graph, vector<int>
}
 bool is = 1:
 for(int child : graph[current]){
                                                                             return pesos;
    if(color[child] == 0){//no esta coloreado
      color[child] = color[current] == 1 ? 2 : 1;// coloreo con el color
→ opuesto a mi nodo current
                                                                           7.4. DSU
      is = is & isBipartiteDFS(child, graph, color);
    else if(color[child] == color[current]) return 0;//si son iquales no
                                                                            struct DSU {
   es bipartito
                                                                             vector<int> parent;
                                                                             vector<int> rank;
 return is;
                                                                             DSU(int n) : parent(n), rank(n) {
                                                                               fore(i,0,n) parent[i] = i;
bool isBipartiteBFS(int start, vector<vector<int>> &graph, vector<int>
int Find(int u) {
 int n = graph.size();
                                                                               return parent[u] = (u == parent[u]? u : Find(parent[u]));
  queue<pair<int, int>> q; //node, color
  q.push({start, 1});
                                                                             void Union(int u, int v) {
  color[start] = 1;
                                                                               int pv = Find(v);
  while(!q.empty()){
                                                                               int pu = Find(u);
    auto current = q.front(); q.pop();
                                                                               if(pv != pu) {
    for (int child :graph[current.first]){
                                                                                 if(rank[pu] < rank[pv]) swap(pu, pv);</pre>
      if(color[child] == current.second) return false;
                                                                                 if(rank[pu] == rank[pv]) rank[pu]++;
      if(color[child] == 0){
                                                                                 parent[pv] = pu;
        color[child] = current.second == 1 ? 2 : 1;
        q.push({child, color[child]});
                                                                             bool is_same(int u, int v) {
                                                                               return Find(u) == Find(v);
 return true;
                                                                           };
                                                                           7.5. FloydWarshall
7.3. Dijkstra
                                                                           void floydWarshall(vvi &mtx){
vi dijkstra(int start, int n, vvpii graph){
                                                                             int n = mtx.size();
 priority_queue<pii, vpii, greater<pii>>> pq;
                                                                             fore(k,0,n) fore(i,0,n) fore(j,0,n)
  vi pesos(n, INT_MAX);
                                                                               mtx[i][j] = min(mtx[i][j], mtx[i][k] + mtx[k][j]);
                                                                           }
 pq.push({0, start});
 pesos[start] = 0;
  while(!pq.empty()){
                                                                           7.6. Kruskal
    int v = pq.top().second;pq.pop();
    for (auto next : graph[v]){
      int u = next.first, w = next.second;
                                                                           vector<int> parent, rank;
      if(pesos[u] > pesos[v] + w){
                                                                           void make_set(int v) {
       pesos[u] = pesos[v] + w;
                                                                               parent[v] = v; rank[v] = 0;
```

}

pq.push({pesos[u], u });

```
int find set(int v) {
    if (v == parent[v]) return v;
    return parent[v] = find_set(parent[v]);
}
void union_sets(int a, int b) {
    a = find set(a):
    b = find_set(b);
    if (a != b) {
        if (rank[a] < rank[b]) swap(a, b);</pre>
        parent[b] = a;
        if (rank[a] == rank[b]) rank[a]++;
    }
}
struct Edge {
    int u, v, weight;
    bool operator<(Edge const& other) {</pre>
        return weight < other.weight;
    }
};
void Kruskal(){
 vector<Edge> edges, result;
 int n, cost = 0;
 parent.resize(n); rank.resize(n);
 for (int i = 0; i < n; i++) make_set(i);
  sort(edges.begin(), edges.end());
 for (Edge e : edges) {
    if (find_set(e.u) != find_set(e.v)) {
      cost += e.weight;
      result.push_back(e);
      union_sets(e.u, e.v);
    }
 }
}
       topoSort
vi toposortBFS(vvi &graph){
    int n = graph.size();
    vi incoming_edges(n);
 fore(i,0,n) for(auto v: graph[i]) incoming_edges[v]++;
    queue<int> q;
    fore(i,0,n) if(incoming_edges[i] == 0) q.push(i);
    int cnt = 0;
    vi ans;
    while(!q.empty()){
```

int u = q.front(); q.pop();

ans.push_back(u);

```
cnt++;
        for(auto v: graph[u]){
            incoming_edges[v]--;
            if(incoming_edges[v] == 0) q.push(v);
        }
    }
    if(cnt != n) return {-1};
    return ans;
void dfs(int curr, vector<bool> &visited, vector<vector<int>> &graph,

    vector<int> &ans){
    if(visited[curr]) return;
    visited[curr] = true;
    for(auto nextNode: graph[curr]) dfs(nextNode, visited, graph, ans);
    ans.push_back(curr);
    return;
vector<int> toposort_dfs(vector<vector<int>> &graph){
    int n = graph.size();
    vector<bool> visited(n);
    vector<int> ans;
    fore(i,0,n) if(!visited[i]) dfs(i, visited, graph, ans);
    reverse(ans.begin(), ans.end());//dfs => reverse
    return ans;
}
```

8. Tree

```
struct Tree { // 1 indexed
  int n, root;
  vvi ancestros; // Binary lifting
  vi level, sz; // Binary lifting
  vi tin, tout, aplanado; // Euler tour
  vi distancias; //Diameter or distance from a node

Tree(int n, int root, vvi& adj): root(root), n(n),
  ancestros(21, vi(n+1, 0)), sz(n+1, 1), level(n+1),
  tin(n+1, 0), tout(n+1, 0), aplanado(n+1,0) {
    int idx = 1;
    precalc(root, root, adj, idx); // precalculos
    fore(k, 1, 21) fore(x, 1, n + 1) //binary lifting
        ancestros[k][x] = ancestros[k - 1][ancestros[k - 1][x]];
    tout[0] = n + 1; // just in case
}

void precalc(int curr, int father, vvi &adj, int &idx) {//DFS O(n)}
```

```
tin[curr] = idx;
   aplanado[idx] = curr;
   for (auto& next : adj[curr]) {
     if (next == father) continue;
      ancestros[0][next] = curr:
     level[next] = level[curr] + 1;
     idx++; precalc(next, curr, adj, idx);
     sz[curr] += sz[next];
   tout[curr] = idx;
 }
 int kAncestorOfx(int k, int x) \{//O(log(n))\}
   fore(j,0,21) if (k & (1 << j)) x = ancestros[j][x];
   return x;
 }
 int LCA(int u, int v) \{//O(\log(n))\}
   if(level[u] < level[v]) swap(u,v);</pre>
   int k = level[u] - level[v];
   u = kAncestorOfx(k, u);
   if(u == v) return u;
   for(int i = 20; i >= 0; i--){
     if(ancestros[i][u] != ancestros[i][v]){
       u = ancestros[i][u]:
       v = ancestros[i][v];
     }
   return ancestros[0][u];
 }
 int isAncestor(int u, int v) { // O(1)
   return tin[u] <= tin[v] && tout[v] <= tout[u];</pre>
 int LCA2(int u, int v) \{//O(\log(n))\}
   if (isAncestor(u, v)) return u;
   if (isAncestor(v, u)) return v;
   for(int k = 20; k >= 0; k--){
     if(!isAncestor(ancestros[k][u], v))
       u = ancestros[k][u];
   return ancestros[0][u];
 }
 void dfs_distances(int curr, int father, vvi &adj, vi &distancias, int
\rightarrow dist){//0(n)}
```

```
distancias[curr] = dist;
  for(auto &next : adj[curr]){
    if(next == father) continue;
    dfs_distances(next, curr, adj, distancias, dist + 1);
 }
int getDiameter(vvi &adj){// O(2n)
  distancias.assign(n+1,0);
  dfs_distances(root, root, adj, distancias, 0);
  int farthest = root;
  int max_dist = 0;
  fore(i,1,n+1){
    if(distancias[i] > max_dist){
      farthest = i:
      max_dist = distancias[i];
   }
 }
  distancias.assign(n+1,0);
  max_dist = 0;
  dfs_distances(farthest, farthest, adj, distancias, 0);
  fore(i,1,n+1){
    max_dist = max(max_dist, distancias[i]);
 return max_dist;
int getCentroid(vvi& adj){ //
  int vf = 1:
  while(vf){
    vf = 0:
    for(auto next: adj[root]){
      if(sz[next] > n/2){
        int aux = sz[root];
        sz[root] -= sz[next];
        sz[next] = aux;
        root = next;
        vf = 1;
      }
    }
  return root;
void printData(){
  // fore(i,0,n+1) cout << tin[i] << ' ' ; cout << endl;
  // fore(i,0,n+1) cout << tout[i] << ' '; cout << endl;
```

```
// fore(i,0,4) {fore(j,0,n+1) cout << ancestros[i][j] <<' '; cout <<
   endl;}
 }
};
8.1. HLD
struct HeavyLightDecomp {
  int n, root, idx;
 vi level, size, head, pos, newVals;
 vvi ancestros;
 HeavyLightDecomp(int n, int root, vvi & adj, vi& vals):n(n),
\rightarrow root(root), ancestros(21, vi(n+1,0)),level(n+1), size(n+1, 1),
\rightarrow head(n+1), pos(n+1), newVals(1){
    dfs(root, root, adj);
    idx = 1; hld(root, root, adj, vals);
 }
 void dfs(int curr, int father, vvi & adj) {
    for (auto& next : adj[curr]) {
      if (next == father) continue;
      ancestros[0][next] = curr;
      level[next] = level[curr] + 1;
      dfs(next, curr, adj);
      size[curr] += size[next]:
 }
  void hld(int curr, int nodeHead, vvi &adj, vi &vals){
    head[curr] = nodeHead;
    pos[curr] = idx++;
    newVals.pb(vals[curr]);
    int sz_heavy = 0, heavy = -1;
    for(auto next: adj[curr]){
      if(next == ancestros[0][curr]) continue;
      if(size[next] > sz_heavy)
        sz_heavy = size[next], heavy = next;
    if(heavy != -1) hld(heavy, nodeHead, adj, vals);
    for(auto next: adj[curr])
      if(next != heavy && next != ancestros[0][curr])
        hld(next, next, adj, vals);
 }
  template <class BinaryOperation>
  void traversePath(int u, int v, BinaryOperation op){
    for(; head[u] != head[v]; u = ancestros[0][head[u]]){
      if(level[head[u]] < level[head[v]]) swap(u,v);</pre>
```

```
op(pos[head[u]], pos[u]);
}
if(pos[u] > pos[v]) swap(u,v);
op(pos[u], pos[v]);
}
template <class DSType> ll query(int u, int v, DSType *st){
  int ans = -1;
    traversePath(u, v, [this, &ans, st](int l, int r){ans = max(ans, st->query(l,r));});
  return ans;
}
template <class DSType> void update(int u, ll val, DSType *st){
    traversePath(u, u, [this, &val, st](int l, int r){st->update(l, val);});
}
```

9. DP

9.1. Knapsack