

Distributed Solution to the K-Queens Problem Using MPI

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Abstract

The purpose of the assignment is to write a program that solves the infamous K-queens problem. The problem involves finding the number of possible placement of K queens on a K by K chessboard. The number of possible placements exponentially increases as K increases. We are to write a parallelized program using MPI to distribute the intensive work amongst a virtual cluster.

Serial Compilation

- `gcc kqueens.c -lm -o kqueens`
- `./kqueens <k>`

MPI Compilation

- `mpicc mpi_queens.c -lm -o mpi_queens`
- `mpirun -np <nodes> --hostfile my_hosts ./mpi_queens <k>`

PARAMETERS

- k: The number of dimensions (k by k) and queens
- nodes: The number of nodes in virtual cluster

Algorithm: Serialized Version

The serialized version is an implementation of Wirth's Algorithm. We simplify our approach by using a solution vector to give hold the placements of Queens on the board. This means that only one Queen can exist in one column at any time. We can think of the starting position as being the top-left corner of the chessboard. This is a recursive solution where we start a position in column 1 and examine all possible solutions starting from that position. We use backtracking to get out of paths that lead to conflicts between Queens.

Algorithm 1 Serialized Version

```
1: procedure PLACEQUEEN(INT K , INT COLUMN, INT * SOLUTION)
2:   if  $k == col$  then
3:      $Solution++$ 
4:   else
5:     for  $j = 1..k$  do
6:       if  $checkrows() == 0$  and  $checkdiagonals() == 0$  then
7:          $solution[col] \leftarrow j$ 
8:          $PlaceQueen(k, column + 1, solution)$ 
```

While this is an elegant solution, it does not scale well. Our recursive tree grows sporadically as the number possible solutions increases with K.

Algorithm: Parallelized Version

We use the same PlaceQueens algorithm described in our serialized solution, but we partition the work among nodes on a Virtual cluster. As stated above, this is a K by K chessboard and we are examining the number of possible placements of K Queens starting from column 1. Column 1 has K possible starting solutions, so we divide those K possible starting positions among K nodes in a cluster. A master node will receive messages as to how many possible messages are received from each worker. This leads to a limitation for the program, for a K by K chessboard, we need $K + 1$ Nodes to be present in our cluster.

Algorithm 2 Parallelize Version

```
1: if myrank  $\neq$  MASTER then
2:   solution[0]  $\leftarrow$  myrank
3:   PlaceQueen(k, 1, myrank, solution)
4: else
5:   for source = 1..NumberOfNodes do
6:     recieve the responses from all worker nodes
7:     SOLUTIONS  $\leftarrow$  sum of all responses
```

We partition the work by distributing among K nodes. Each node will start their own recursive tree and find their individual number of solutions and send it to the Master node. The Master node will then go through each worker and receive their number of solutions through a response variable. The sum of all these variables will be stored as SOLUTIONS.

Accuracy of Models

In this section we analyzed how accurate were the models to one another. We can see that both approaches return the same number of solutions.

Solution matrix

Number of solutions **for** both approaches.

K = 8,10,12,14,16

| | 8 | 10 | 12 | 14 | 16 |
|------------|----|-----|-------|--------|----------|
| sequential | 92 | 724 | 14200 | 365596 | 14772512 |
| mpi | 92 | 724 | 14200 | 365596 | 14772512 |

RESULTS: SEQUENTIAL K = 8

```
[1 5 8 6 3 7 2 4 ]
[1 6 8 3 7 4 2 5 ]
[1 7 4 6 8 2 5 3 ]
[1 7 5 8 2 4 6 3 ]
[2 4 6 8 3 1 7 5 ]
[2 5 7 1 3 8 6 4 ]
[2 5 7 4 1 8 6 3 ]
[2 6 1 7 4 8 3 5 ]
[2 6 8 3 1 4 7 5 ]
[2 7 3 6 8 5 1 4 ]
[2 7 5 8 1 4 6 3 ]
[2 8 6 1 3 5 7 4 ]
[3 1 7 5 8 2 4 6 ]
[3 5 2 8 1 7 4 6 ]
[3 5 2 8 6 4 7 1 ]
[3 5 7 1 4 2 8 6 ]
[3 5 8 4 1 7 2 6 ]
[3 6 2 5 8 1 7 4 ]
[3 6 2 7 1 4 8 5 ]
[3 6 2 7 5 1 8 4 ]
[3 6 4 1 8 5 7 2 ]
```

[3 6 4 2 8 5 7 1]
[3 6 8 1 4 7 5 2]
[3 6 8 1 5 7 2 4]
[3 6 8 2 4 1 7 5]
[3 7 2 8 5 1 4 6]
[3 7 2 8 6 4 1 5]
[3 8 4 7 1 6 2 5]
[4 1 5 8 2 7 3 6]
[4 1 5 8 6 3 7 2]
[4 2 5 8 6 1 3 7]
[4 2 7 3 6 8 1 5]
[4 2 7 3 6 8 5 1]
[4 2 7 5 1 8 6 3]
[4 2 8 5 7 1 3 6]
[4 2 8 6 1 3 5 7]
[4 6 1 5 2 8 3 7]
[4 6 8 2 7 1 3 5]
[4 6 8 3 1 7 5 2]
[4 7 1 8 5 2 6 3]
[4 7 3 8 2 5 1 6]
[4 7 5 2 6 1 3 8]
[4 7 5 3 1 6 8 2]
[4 8 1 3 6 2 7 5]
[4 8 1 5 7 2 6 3]
[4 8 5 3 1 7 2 6]
[5 1 4 6 8 2 7 3]
[5 1 8 4 2 7 3 6]
[5 1 8 6 3 7 2 4]
[5 2 4 6 8 3 1 7]
[5 2 4 7 3 8 6 1]
[5 2 6 1 7 4 8 3]
[5 2 8 1 4 7 3 6]
[5 3 1 6 8 2 4 7]
[5 3 1 7 2 8 6 4]
[5 3 8 4 7 1 6 2]
[5 7 1 3 8 6 4 2]
[5 7 1 4 2 8 6 3]
[5 7 2 4 8 1 3 6]
[5 7 2 6 3 1 4 8]
[5 7 2 6 3 1 8 4]

```
[5 7 4 1 3 8 6 2 ]
[5 8 4 1 3 6 2 7 ]
[5 8 4 1 7 2 6 3 ]
[6 1 5 2 8 3 7 4 ]
[6 2 7 1 3 5 8 4 ]
[6 2 7 1 4 8 5 3 ]
[6 3 1 7 5 8 2 4 ]
[6 3 1 8 4 2 7 5 ]
[6 3 1 8 5 2 4 7 ]
[6 3 5 7 1 4 2 8 ]
[6 3 5 8 1 4 2 7 ]
[6 3 7 2 4 8 1 5 ]
[6 3 7 2 8 5 1 4 ]
[6 3 7 4 1 8 2 5 ]
[6 4 1 5 8 2 7 3 ]
[6 4 2 8 5 7 1 3 ]
[6 4 7 1 3 5 2 8 ]
[6 4 7 1 8 2 5 3 ]
[6 8 2 4 1 7 5 3 ]
[7 1 3 8 6 4 2 5 ]
[7 2 4 1 8 5 3 6 ]
[7 2 6 3 1 4 8 5 ]
[7 3 1 6 8 5 2 4 ]
[7 3 8 2 5 1 6 4 ]
[7 4 2 5 8 1 3 6 ]
[7 4 2 8 6 1 3 5 ]
[7 5 3 1 6 8 2 4 ]
[8 2 4 1 7 5 3 6 ]
[8 2 5 3 1 7 4 6 ]
[8 3 1 6 2 5 7 4 ]
[8 4 1 3 6 2 7 5 ]
```

Found 92

RESULTS: MPI , K = 8

```
PROCESS 1 [1 5 8 6 3 7 2 4 ]
PROCESS 1 [1 6 8 3 7 4 2 5 ]
PROCESS 1 [1 7 4 6 8 2 5 3 ]
```

PROCESS 1 [1 7 5 8 2 4 6 3]
Process 1 computed 4 solutions

PROCESS 6 [6 1 5 2 8 3 7 4]
PROCESS 7 [7 1 3 8 6 4 2 5]
PROCESS 6 [6 2 7 1 3 5 8 4]
PROCESS 6 [6 2 7 1 4 8 5 3]
PROCESS 7 [7 2 4 1 8 5 3 6]
PROCESS 7 [7 2 6 3 1 4 8 5]
PROCESS 7 [7 3 1 6 8 5 2 4]
PROCESS 7 [7 3 8 2 5 1 6 4]
PROCESS 7 [7 4 2 5 8 1 3 6]
PROCESS 7 [7 4 2 8 6 1 3 5]
PROCESS 7 [7 5 3 1 6 8 2 4]
PROCESS 6 [6 3 1 7 5 8 2 4]
PROCESS 6 [6 3 1 8 4 2 7 5]
PROCESS 6 [6 3 1 8 5 2 4 7]
PROCESS 6 [6 3 5 7 1 4 2 8]
PROCESS 6 [6 3 5 8 1 4 2 7]
PROCESS 6 [6 3 7 2 4 8 1 5]
PROCESS 6 [6 3 7 2 8 5 1 4]
PROCESS 6 [6 3 7 4 1 8 2 5]
PROCESS 6 [6 4 1 5 8 2 7 3]
PROCESS 6 [6 4 2 8 5 7 1 3]
PROCESS 6 [6 4 7 1 3 5 2 8]
PROCESS 6 [6 4 7 1 8 2 5 3]
PROCESS 6 [6 8 2 4 1 7 5 3]
PROCESS 8 [8 2 4 1 7 5 3 6]
PROCESS 8 [8 2 5 3 1 7 4 6]
PROCESS 8 [8 3 1 6 2 5 7 4]
PROCESS 8 [8 4 1 3 6 2 7 5]
PROCESS 2 [2 4 6 8 3 1 7 5]
PROCESS 2 [2 5 7 1 3 8 6 4]
PROCESS 2 [2 5 7 4 1 8 6 3]
PROCESS 4 [4 1 5 8 2 7 3 6]
PROCESS 4 [4 1 5 8 6 3 7 2]
PROCESS 4 [4 2 5 8 6 1 3 7]
PROCESS 4 [4 2 7 3 6 8 1 5]
PROCESS 4 [4 2 7 3 6 8 5 1]
PROCESS 4 [4 2 7 5 1 8 6 3]

PROCESS 4 [4 2 8 5 7 1 3 6]
PROCESS 4 [4 2 8 6 1 3 5 7]
PROCESS 4 [4 6 1 5 2 8 3 7]
PROCESS 4 [4 6 8 2 7 1 3 5]
PROCESS 3 [3 1 7 5 8 2 4 6]
PROCESS 3 [3 5 2 8 1 7 4 6]
PROCESS 3 [3 5 2 8 6 4 7 1]
PROCESS 3 [3 5 7 1 4 2 8 6]
PROCESS 3 [3 5 8 4 1 7 2 6]
PROCESS 3 [3 6 2 5 8 1 7 4]
PROCESS 3 [3 6 2 7 1 4 8 5]
PROCESS 3 [3 6 2 7 5 1 8 4]
PROCESS 3 [3 6 4 1 8 5 7 2]
PROCESS 3 [3 6 4 2 8 5 7 1]
PROCESS 3 [3 6 8 1 4 7 5 2]
PROCESS 3 [3 6 8 1 5 7 2 4]
PROCESS 3 [3 6 8 2 4 1 7 5]
PROCESS 4 [4 6 8 3 1 7 5 2]
PROCESS 4 [4 7 1 8 5 2 6 3]
PROCESS 4 [4 7 3 8 2 5 1 6]
PROCESS 4 [4 7 5 2 6 1 3 8]
PROCESS 4 [4 7 5 3 1 6 8 2]
PROCESS 4 [4 8 1 3 6 2 7 5]
PROCESS 4 [4 8 1 5 7 2 6 3]
PROCESS 4 [4 8 5 3 1 7 2 6]
PROCESS 2 [2 6 1 7 4 8 3 5]
PROCESS 2 [2 6 8 3 1 4 7 5]
PROCESS 2 [2 7 3 6 8 5 1 4]
PROCESS 2 [2 7 5 8 1 4 6 3]
PROCESS 2 [2 8 6 1 3 5 7 4]
PROCESS 3 [3 7 2 8 5 1 4 6]
PROCESS 3 [3 7 2 8 6 4 1 5]
PROCESS 3 [3 8 4 7 1 6 2 5]
PROCESS 5 [5 1 4 6 8 2 7 3]
PROCESS 5 [5 1 8 4 2 7 3 6]
PROCESS 5 [5 1 8 6 3 7 2 4]
PROCESS 5 [5 2 4 6 8 3 1 7]
PROCESS 5 [5 2 4 7 3 8 6 1]
PROCESS 5 [5 2 6 1 7 4 8 3]
PROCESS 5 [5 2 8 1 4 7 3 6]

PROCESS 5 [5 3 1 6 8 2 4 7]
PROCESS 5 [5 3 1 7 2 8 6 4]
PROCESS 5 [5 3 8 4 7 1 6 2]
PROCESS 5 [5 7 1 3 8 6 4 2]
PROCESS 5 [5 7 1 4 2 8 6 3]
PROCESS 5 [5 7 2 4 8 1 3 6]
PROCESS 5 [5 7 2 6 3 1 4 8]
PROCESS 5 [5 7 2 6 3 1 8 4]
PROCESS 5 [5 7 4 1 3 8 6 2]
PROCESS 5 [5 8 4 1 3 6 2 7]
PROCESS 5 [5 8 4 1 7 2 6 3]
Process 2 computed 16 solutions

Process 3 computed 16 solutions

Process 4 computed 8 solutions

Process 5 computed 4 solutions

Process 6 computed 8 solutions

Process 7 computed 18 solutions

Process 8 computed 18 solutions

Total Number of Solutions 92

Visualization

Figure 1:

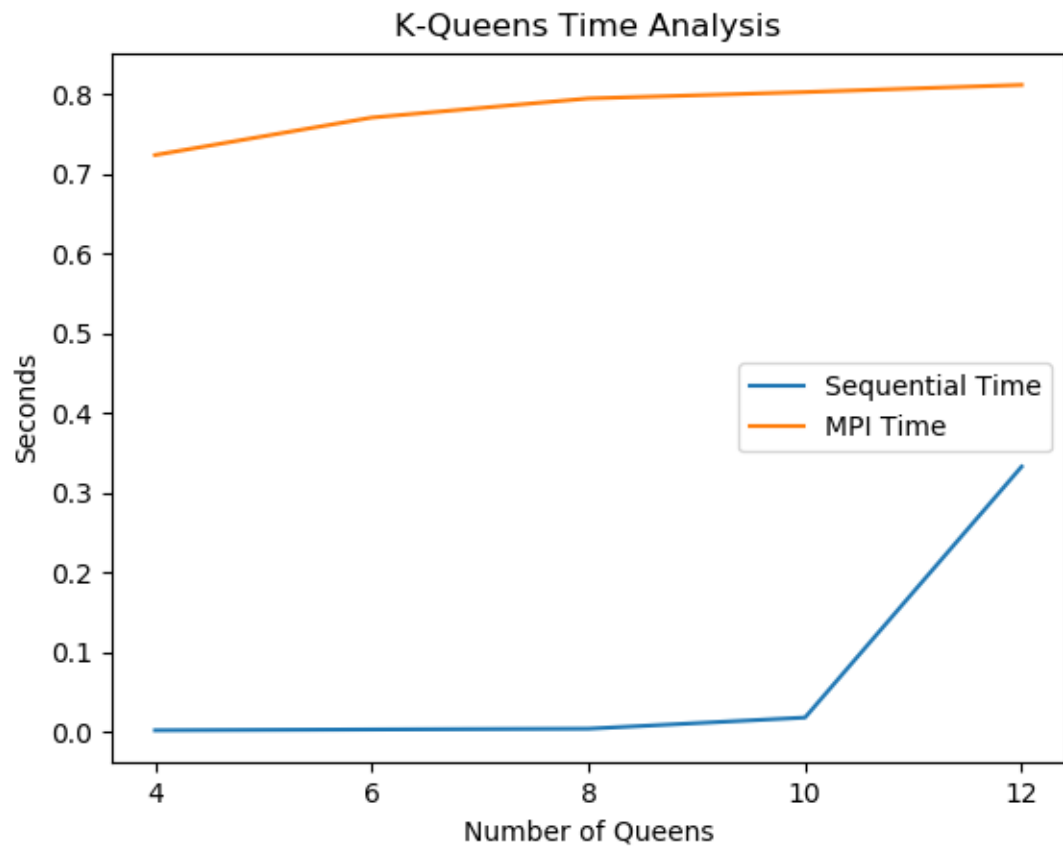
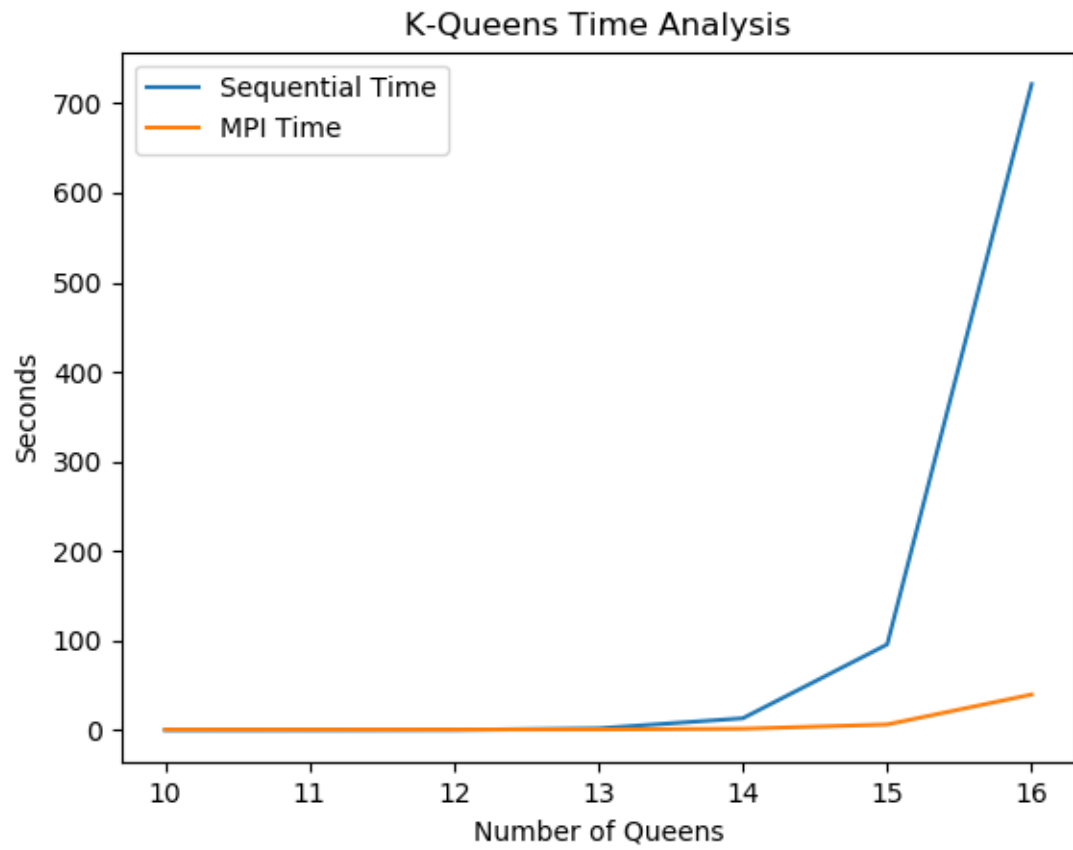


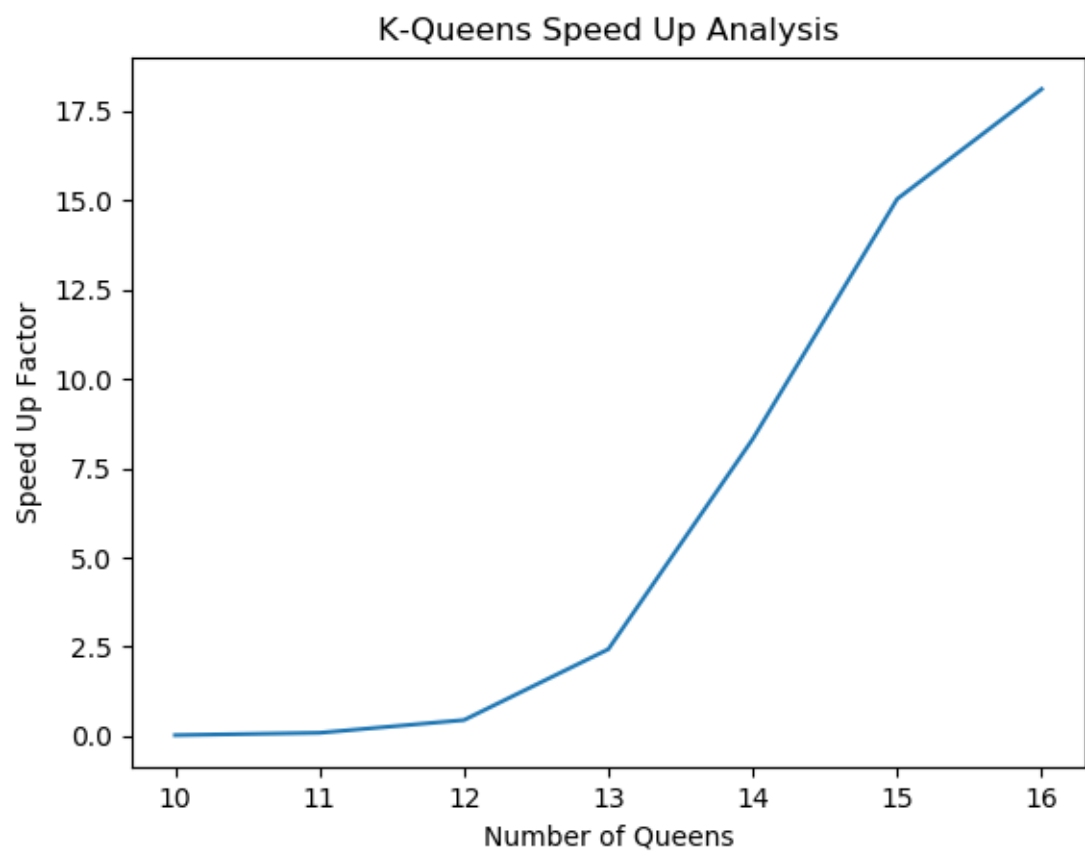
Figure 2:



Time Complexity Analysis

Our parallel program suffers from the overhead communication for small k values. For large k -values we see that our parallelized version does not suffer from such a drastic exponential growth. However, it does still grow exponentially. This means that the payoff for partitioning the work will still eventually suffer from a bottleneck. Our parallel version is ideal for k -values 14-18, but our speedup after that begins to wane off. The work of each individual worker is still growing exponentially. In a sense, our parallel version is essentially prolonging the inevitable growth curve. We also limit our parallelism to K , which is not ideal if we wanted to mitigate that exponential growth by adding more and more nodes to high values of k .

Figure 3:



CODE: Serialized Version

```
#include <stdio.h>
#include <math.h>
#include <stdlib.h>

void place_queen(int,int, int []);
int check_rows(int, int, int[]);
int check_diagonals(int, int, int []);
void initialize_array(int, int []);
void printArray(int ,int []);

int Solutions;

int main (int argc, char *argv[]) {

    int k = 0;
    if (argc != 2) {
        fprintf (stderr, "need a k for K Queens Problem!!!\n");
        exit(-1);
    }

    k = atoi (argv[1]);

    int solution[k];
    //int start = 0;

    initialize_array(k, solution);

    place_queen(k, 0, solution);
    printf("Found %d\n", Solutions);
}
```

```

void place_queen(int k, int col, int * solution)
{

    if (k == col)
    {
        printArray(k,solution);
        Solutions++;
    }

    for(int row=1; row <= k; row++)
    {
        // check if queen can be placed safely
        if(check_rows(row, col,solution) == 0 && check_diagonals(row,
            col, solution) == 0)
        {
            //place queen on this row
            solution[col] = row;
            // try rest of problem
            place_queen (k,col+1, solution);

        }//end inner if

    }//end loop
}

//will take a row and column and see if a queen is in danger
int check_rows(int row, int col, int * solution){
    //queens cant exist on the same row
    for(int i = 0; i < col; i ++){
        if(solution[i] == row ){
            return -1;
        }
    }
    return 0;
}

int check_diagonals(int row, int col, int * solution){
    //values we will be examining

```

```

int danger1;
int danger2;

for(int i = 0; i < col; i++){
    danger1 = solution[i]+i; // diagonals from left side
    danger2 = i - solution[i]; //diagonals from rights side

    if(row + col == danger1){ //left hand diagonals
        return -1;
    }
    if( col - row == danger2){ //right hand diagonals
        return -1;
    }
}
return 0;
}

//fill array with zeros
void initialize_array(int k, int * array){
    for (int i = 0; i < k; i++){
        array[i] = 0;
    }
}

//print array
void printArray(int k, int * array){
    printf("[");
    for (int i = 0; i < k; i++){
        printf("%d ", array[i]);
    }
    printf("]\n");
}

```

CODE: Parallelized Version

```

#include <stdio.h>
#include <math.h>
#include <stdlib.h>
#include <mpi.h>
#include <time.h>
#include <unistd.h>

void place_queen(int,int,int , int []);
int check_rows(int, int, int[]);
int check_diagonals(int, int, int []);
void initialize_array(int, int []);
void printArray(int,int ,int []);

int Solutions;

#define MASTER 0
#define TAG    0

int main (int argc, char *argv[]) {
    int k = 0;
    if (argc != 2) {
        fprintf (stderr, "need a k for K Queens Problem!!!!\n");
        exit(-1);
    }
    int my_rank, num_nodes, source;

    MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
    MPI_Comm_size(MPI_COMM_WORLD, &num_nodes);
    k = atoi (argv[1]);

    int solution[k];
    //int start = 0;

    if (my_rank != MASTER) {

```

```

        solution[0] = my_rank;
        place_queen(k, 1, my_rank, solution );

        MPI_Send(&Solutions, 1, MPI_INT, MASTER, 0, MPI_COMM_WORLD);
    }
    else {

        //master calculates a reduction
        for (source = 1; source < num_nodes; source++) {
            int response;
            MPI_Recv(&response, 1, MPI_INT, MPI_ANY_SOURCE, 0,
                    MPI_COMM_WORLD, MPI_STATUS_IGNORE);
            printf("Process %d computed %d solutions\n", source,
                    response);
            Solutions += response;
            printf("-----\n");

        }
        printf("Total Number of Solutions %d\n", Solutions);
    }

    MPI_Finalize();
    return 0;
}

void place_queen(int k, int col, int process, int * solution)
{
    // srand(time(NULL));

    if (k == col)
    {
        Solutions++;
    }
    else{

```

```

solution[0] = process;
for(int row = 1; row <= k; row++)
{
    // check if queen can be placed safely
    if(check_rows(row, col,solution) == 0 &&
        check_diagonals(row, col, solution) == 0)
    {
        //place queen on this row
        solution[col] = row;

        place_queen (k,col+1, process,solution);

    }//end inner if

} //end loop
}
} //end method

//will take a row and column and see if a queen is in danger
int check_rows(int row, int col, int * solution){
    //queens cant exist on the same row
    for(int i = 0; i < col; i ++){
        if(solution[i] == row ){
            return -1;
        }
    }
    return 0;
}

int check_diagonals(int row, int col, int * solution){
    //values we will be examining
    int danger1;
    int danger2;

    for(int i = 0; i < col; i ++){
        danger1 = solution[i]+i; // diagonals from left side
        danger2 = i - solution[i]; //diagonals from rights side
    }
}

```

```

        if(row + col == danger1){ //left hand diagonals
            return -1;
        }
        if( col - row == danger2){ //right hand diagonals
            return -1;
        }
    }
    return 0;
}

//fill array with zeros
void initialize_array(int k, int * array){
    for (int i = 0; i < k; i++){
        array[i] = 0;
    }
}

//print array
void printArray(int k,int process, int * array){
    printf("PROCESS %d [", process);
    for (int i = 0; i < k; i++){
        printf("%d ", array[i]);
    }
    printf("]\n");
}

```
