

• PSEUDO CODE:-

Selection sort (A)

1. for $i=1$ to $A.length$
2. $min_pos = i$
3. for $J=i+1$ to length of list
4. if $list[min_pos] > list[J]$
5. $min_pos = J$
6. $temp = list[i]$
7. $list[i] = list[min_pos]$
8. $list[min_pos] = temp$

• DRY RUNNING

$A = [4, 3, 2, 1] \rightarrow$ worst case

• 1ST ITERATION OF FOR:-

1. $i=1$
2. $min_pos = 1$
3. $J=2 \rightarrow$ 1st iteration of for
4. if $4 > 3$
5. $min_pos = 2$

2nd iteration of for 3. $J=3$

4. if $3 > 2$
5. $min_pos = 3$

3rd iteration of for

3. $J=4$
4. if $2 > 1$

5.

min-pos = 4

4th iteration of for

3. J = 5

4. _____

5. _____

6. temp = 4

7. list [1] = 1

8. list [4] = 4 \rightarrow A: [1, 3, 2, 4]

• 2ND ITERATION OF FOR:-

1. i = 2

2. min-pos = 2

1st iteration of for:-

3. J = 3

4. If $3 > 2$

5. min-pos = 3

2nd iteration of for:-

3. J = 4

4. If $2 > 4$

5. _____

3rd iteration of for:-

3. J = 5

4. _____

5. _____

6. temp = 3

7. list [2] = 3

8. list [3] = 3 \rightarrow A: [1, 2, 3, 4]



Teacher's Signature

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• 3RD ITERATION OF FOR:-

1. $i = 3$

2. $\text{min-pos} = 3$

3. 1st iteration of for:- $J = 4$

4. $\text{if } 3 > 4$

5. _____

2nd iteration of for

3. $J = 5$

4. _____

5. _____

6. $\text{temp} = 3$

7. $\text{list}[3] = 3 \rightarrow A = [1, 2, 3, 4]$

8. $\text{list}[3] = 3$

• 4TH ITERATION OF FOR

1. $i = 4$

2. $\text{min-pos} = 4$

1st iteration of for:-

3. $J = 5$

4. _____

5. _____

6. $\text{temp} = 4$

7. $\text{list}[4] = 4 \rightarrow A = [1, 2, 3, 4]$

8. $\text{list}[4] = 4$

• 5TH ITERATION OF FOR:-

1. $i = 5 \rightarrow$ loop terminates

Line no.	Time/ instruction	frequency (worst case)	(best case)
1	C	$n+1$	$n+1$
2	C	n	n
3	C	$\sum_{j=1}^n j \cdot \frac{n(n+1)}{2}$	$\sum_{j=1}^n j \cdot \frac{n(n+1)}{2}$
4	C	$\sum_{j=1}^n j-1 \cdot \frac{n(n-1)}{2}$	$\sum_{j=1}^n j-1 \cdot \frac{n(n-1)}{2}$
5	C	$\sum_{j=1}^n j-1 \cdot \frac{n(n-1)}{2}$	0
6	C	n	n
7	C	n	n
8	C	n	n

Worst case:

$$T(n) = c \left[n+1 + n + \frac{n(n+1)}{2} + \frac{n(n-1)}{2} + \frac{n(n-1)}{2} + n + n + n \right]$$

$$T(n) = c \left[5n+1 + \frac{n^2+n}{2} + \frac{n^2-n}{2} + \frac{n^2-n}{2} \right]$$

$$T(n) = c \left[10n+2 + \frac{n^2+n}{2} + \frac{n^2-n}{2} + \frac{n^2-n}{2} \right]$$

$$T(n) = c \left[\frac{3n^2 + 9n + 2}{2} \right]$$

$$T(n) = K_1 n^2 + K_2 n + K_3$$

• DISCUSSION:-

- selection sort in worst case grows quadratically.
- Worst case time complexity of selection sort is $O(n^2)$
- Means that, maximum time which a code can take.

(Best case):-

$$T(n) = \left[n+1 + n + \frac{n(n+1)}{2} + \frac{n(n-1)}{2} + 0 + n + n + n \right] c$$

$$T(n) = \left[5n + 1 + \frac{n^2 + n}{2} + \frac{n^2 - n}{2} \right] c$$

$$T(n) = \left[\frac{10n + 2 + 2n^2}{2} \right] c$$

$$T(n) = [K_1 n^2 + K_2(n) + K_3] c$$

• DISCUSSION:-

- selection sort in best case grows quadratically.
- best case time complexity of selection sort is $\Omega(n^2)$
- Means that, minimum time which a code can take