

LINKED LIST:-

A linked list consists of nodes where each node contains a data element and the reference (link) to the next node in the list. First node of linked list is called "head". Last node of linked list is called "tail".

→ IMPLEMENTING STACKS THROUGH

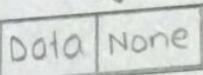
LINKED LIST

First making a class node

class Node():

```
def __init__(self, data):  
    self.data = data  
    self.next = None
```

In order to create nodes in a linked list we first create a class node with attribute data



Now making a class stack

class LinkStack():

```
def __init__(self):  
    self.head = None  
    self.tail = None
```

Now in order to implement stacks in linked list we now creating a class stack and initializing head and tail to None

Now writing a method of push

def push(self, data):

```
new-node = Node(data) → [Data | None]
```

new-node.next = self.head → [Data | None] → [Data | None]
1500 → new-node → 1600
1500 1600 → 1 head

self.head = new-node → [Data | 1600] → [Data | None]
↑ head

* head is always the first node of linked list

* In push function we create nodes at head

Now writing a function of pop

def pop(self):

temp = self.head

1500 → Data | 1600 → Data | None

↑ head ↑ temp

self.head, temp.next

1500 → Data | 1600 → Data | None

↑ 1500 ↑ temp ↑ 1600 ↑ head

temp.next = None → Data | None → Data | None

↑ temp ↑ head

⇒ In pop function we delete the first node by disconnecting its link to second node.

Now writing a function of show-all

def show-all(self):

temp = self.head

If self.head is None:

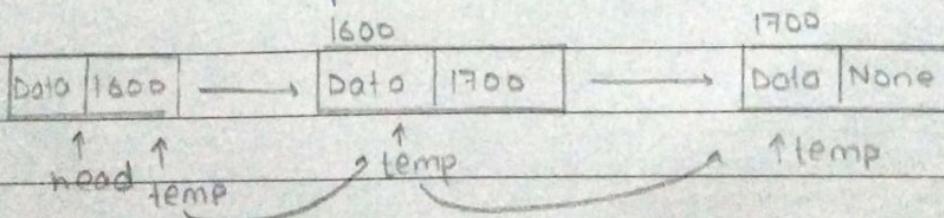
Print ("List is Empty")

return

While temp:

Print (temp.data)

temp = temp.next



In order to print our linked list we first make a temporary variable equal to head and then traverse it to last.

⇒ IMPLEMENTING QUEUES THROUGH

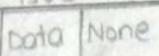
LINKED LIST :-

First making a class node

```
class node():
```

```
    def __init__(self, data):  
        self.data = data  
        self.next = None
```

→ First we make a class node in order to create nodes in a linked list i.e.; with attribute data



Now making a class queue

```
class Queue():
```

```
    def __init__(self, data):  
        self.head = None  
        self.tail = None
```

→ In order to implement queue through Linked list we now creating a class queue and initializing head & tail to None

Now writing a method of enqueue

```
def enqueue(self, data):
```

If self.head is None:

```
    new-node = node(data) → [Data | None]
```

1500

[Data | None]

↑ ↑
head tail

else:

```
    new-node = node(data)
```

1500

[Data | None]

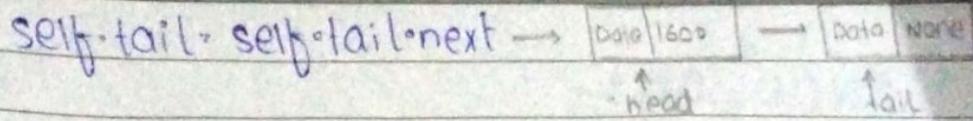
↑ ↑
head tail

1600

[Data | None]

```
    self.tail.next = new-node → [Data | 1600] → [Data | None]
```

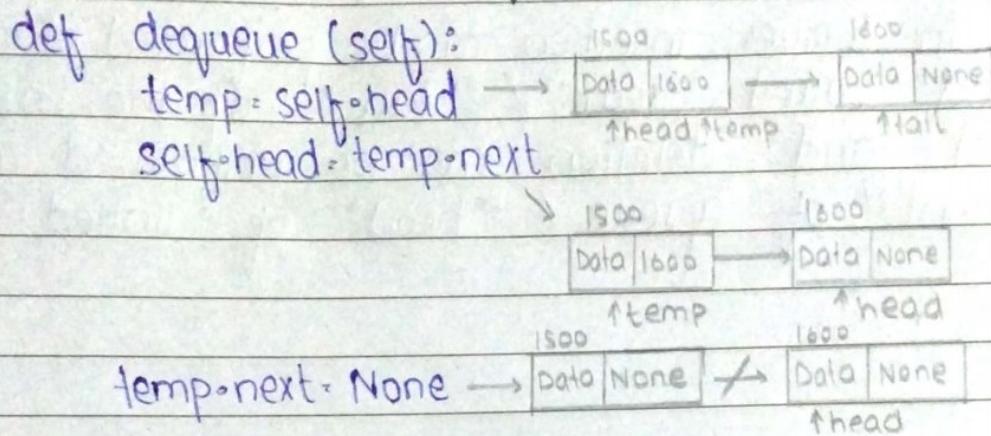
↑ ↑
head tail



* Last node of linked list is "tail"

* In enqueue function we add nodes at tail.

Now writing a method of dequeue



Now writing a method of show-all

def show-all(self):

temp = self.head

If self.head is None:

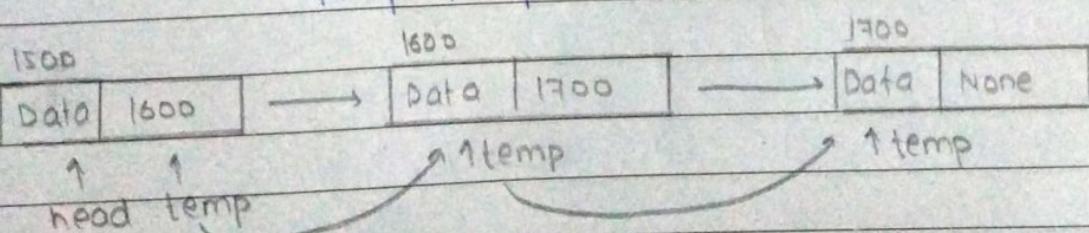
Print ("List is Empty")

Return

While temp:

Print (temp.data)

temp = temp.next



⇒ In order to print whole linked list we need to traverse whole linked list.

CONCLUSION:-

- => In stacks both insertion and deletion can be performed at one end called "top" of the stack and "head" of the linked list.
- => In Queues insertion can be performed at one end called "rear" / "tail" of the linked list and deletion can be performed at "front" of queue / "head" of linked list.

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• : ENDATP •

Using ~~node class~~ develop stack and queues.

• : T2U ANSWER

singly connected linked list with
following features

- a. add nodes
- b. Traverse all node starting from top node
- c. search any key value in all node
- d. Insert node between any two nodes.

• STACK:

A list or collection with restriction that insertion can be performed at one end (head) and deletion can be performed also at same end (head).

• STACKS :-

• STACK IMPLEMENTATION OF

LINKED LIST:-

⇒ we will select the beginning of the linked list as the top of the stack.

⇒ For push operation, a node will be inserted at the beginning / head of the linked list. bcz we consider the beginning of the linked list as the top of the stack.

⇒ For pop operation, everytime the first node will be deleted.

⇒ Time complexity of adding a node at beginning is $O(1)$

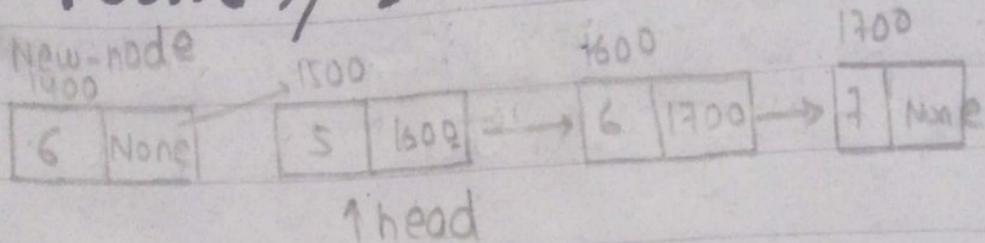
⇒ Time complexity of removing a node at beginning is also $O(1)$.

• NOTE:

- The code of the push() function must be similar to the code of inserting the node at beginning of singly connected linked list.
- The code of the pop() operation must be similar to the removing the node at the beginning of the singly connected linked list.

CLASS STACK :-

PUSH() / INSERTION_AT-BEGINNING

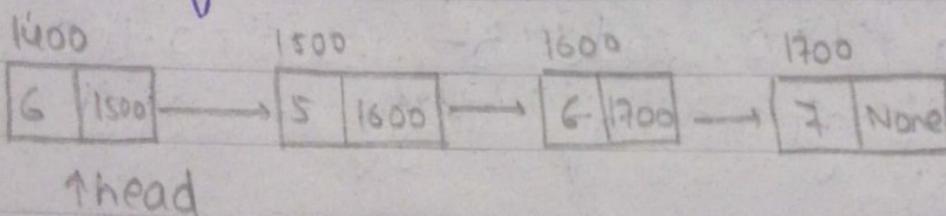


def push(self, data):

 new-node = node(data)

 new-node.next = self.head → linking

 self.head = new-node → head updated

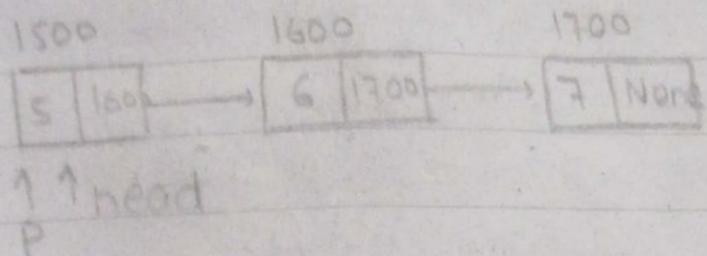


* head is always first node

* In stack.push() operation is similar to insertion at beginning.

* Time complexity of push operation is O(1)

POP / DELETION AT BEGINNING:-



def pop(self):

P = self.head

self.head = P.next

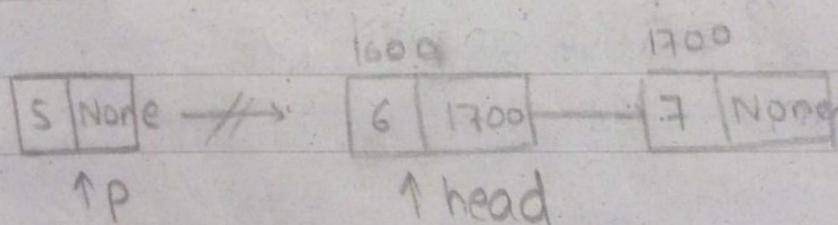
P.next = None

1600

6 | 1700

5 | None

↑ head



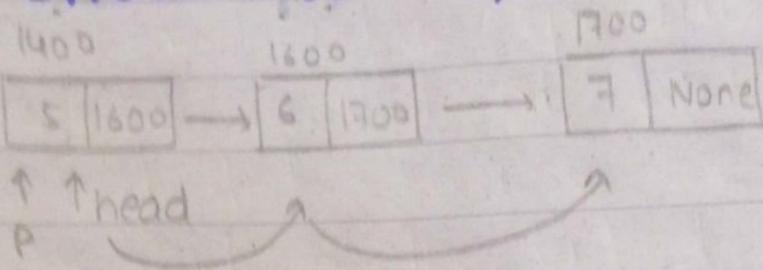
* deletion at beginning is just deleting the connection of first and second node.

* We, first update our head.

* Then disconnecting links.

* Time complexity of pop operation is O(1)

SHOW-ALL() / TRAVERSAL:-



def show-all(self):

P = self.head

While self.head is None:

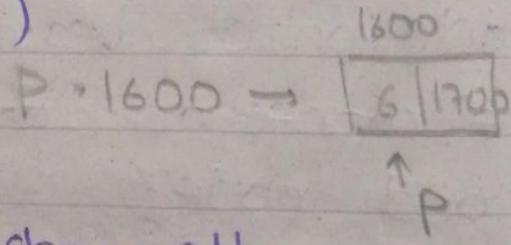
Print ("List is Empty")

return

While P:

Print (P.data)

P = P.next



* Time-complexity of show-all function is O(n)

* We must have to traverse all nodes of linked list from head to last node.

• QUEUES:-

QUEUE IMPLEMENTATION OF

LINKED LIST:-

• Queue:

A list or collection with restriction that insertion can be performed at one end (tail) and deletion can be performed at other hand (head)

Operations:

1. enqueue () $\rightarrow O(1)$
2. dequeue () $\rightarrow O(n)$
3. show_all () $\rightarrow O(n)$

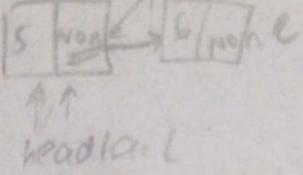
=> we will select the tail of the linked list as the top of the queue.

mtld queue mai node hamesha
tail pr ke insert hogi

- ⇒ For enqueue operation, a node will be inserted at tail of the linked list, because we consider the tail of the linked list as the top of the queue.
- ⇒ For pop operation, first node will be deleted.
- ⇒ Time complexity of enqueue is $O(1)$.
- ⇒ Time complexity of dequeue is $O(1)$.

NOTE:-

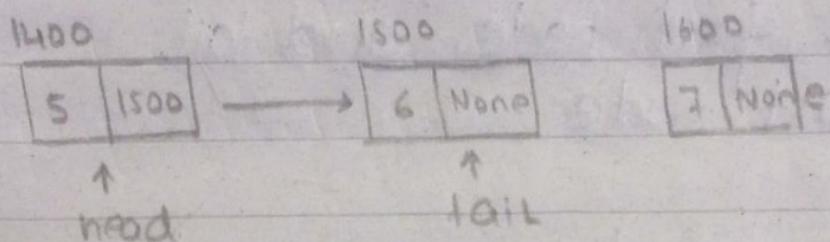
- The code of enqueue operation must be similar to the insertion of node at tail.



- * The code of deque operation must be similar to the deletion of node at beginning.

CLASS QUEUE:-

ENQUEUE / INSERTION AT TAIL :-



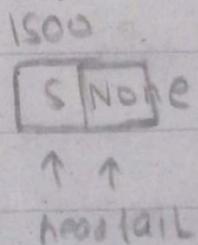
```
def enqueue(self, data):
```

If self.tail is None:

```
    new_node = node(data)
```

```
    self.head = new_node
```

```
    self.tail = self.head
```

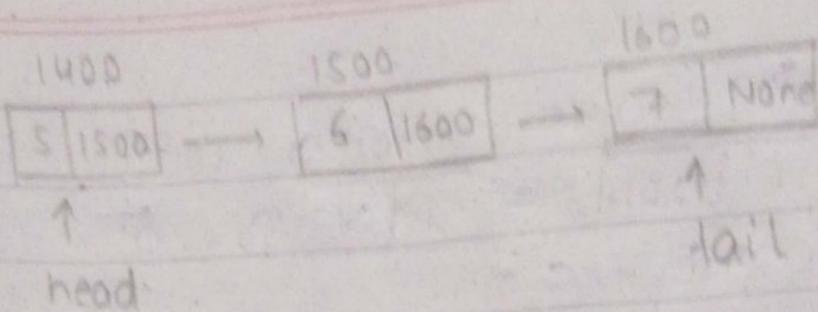


else:

```
    new_node = node(data)
```

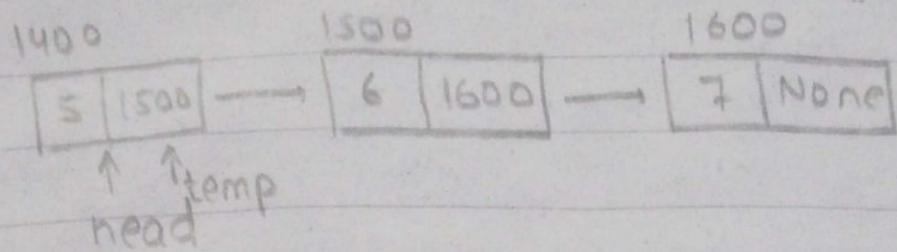
```
    self.tail.next = new_node
```

```
    self.tail = self.tail.next
```



- Time complexity of enqueue operation is $O(1)$
- head is always first node and tail is always last node
- In enqueue operation, we insert nodes at tail.

DEQUEUE / DELETION AT BEGINNING:



def dequeue(self):

temp = self.head \leftarrow first make temporary variable
self.head = temp.next \rightarrow update head
temp.next = None \rightarrow disconnect 1st node and second node

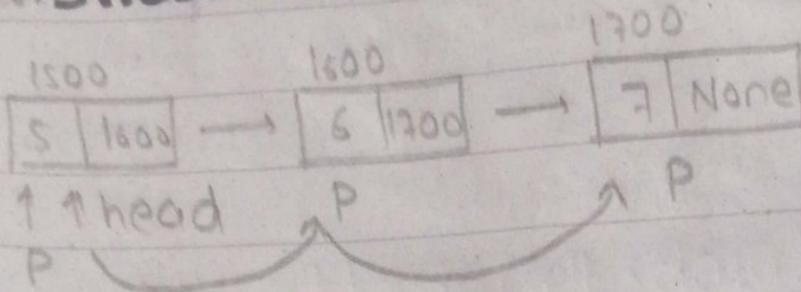
\Rightarrow Time complexity of dequeue operation is $O(1)$

\Rightarrow We first make temp variable -

\Rightarrow Then update our head

\Rightarrow Then disconnecting first and second node links.

SHOW-ALL / TRAVERSAL:



```
def show_all(self):  
    P = self.head  
    if self.head is None:  
        print("List is Empty")
```

```
    while P:  
        print(P.data)  
        P = P.next
```

- * Time complexity of show-all function is $O(n)$
- * We must have to traverse all nodes from head to tail.

LINKED LIST WITH FEATURES:-

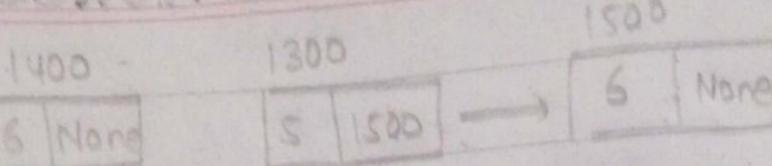
```
class node():
    def __init__(self, data):
        self.data = data
        self.next = None
```

```
class linked_list():
    def __init__(self):
        self.head = None
```

(a) add node

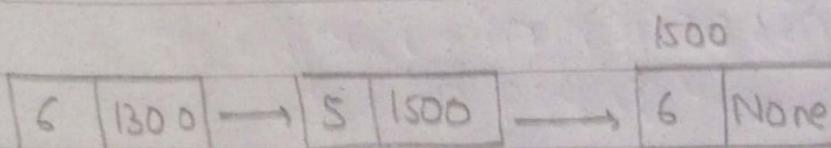
=> hum add node head pai bhi
Kiskay and tail pai bhi.

=> function of add_node is same
as insertion at beginning or
insertion at end.



new-node 1
head

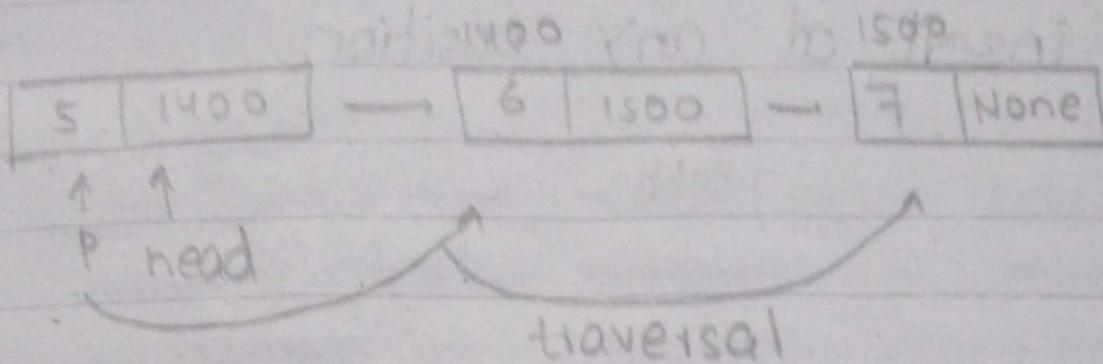
```
def addnode(self, data):
    new-node = node(data)
    new-node.next = self.head
    self.head = new-node
```



↑
head

- head is always first node
- Time complexity of new-node is $O(1)$

(b) Traversal



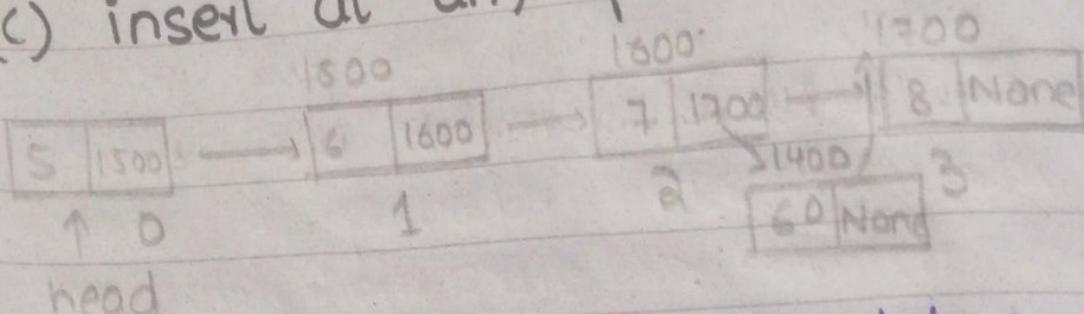
```
def traversal(self):  
    P = self.head
```

If self.head is None:
 print ("List is Empty")

While P:
 Print (P.data)
 P = P.next

=> Time complexity of traversal is
 $O(n)$.

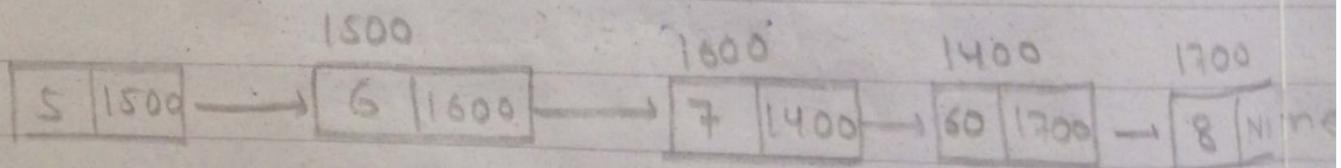
(c) insert at any position



```
def insert_at_index(self, data, index):
    new_node = node(data)
    p = self.head
    for i in range(index - 1):
        p = p.next
```

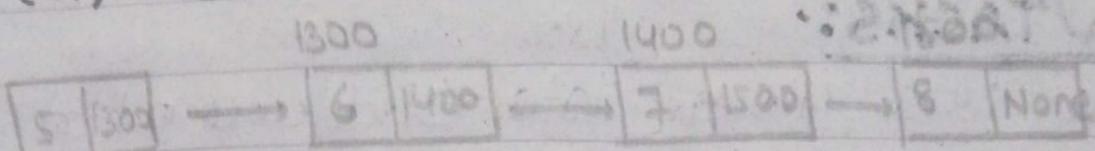
new_node.next = p.next

p.next = new_node



PC : SAL

(a) Search any key value in alt-node



↑ ↑ P
head.

def search(self, key):
 P = self.head
 i = 0

While P:

If key == P.data:

Print(f" {key} is present at
index {i} ")

i = i + 1

return

else:

P = P.next

Print("Key is absent")

↳ If key not found -

DAB: 09

STACKS:

Stacks is a linear data structure which follows a particular order in which the operations are performed. The order may be LIFO (last in first out).

- The operations of stacks are:
- push operation
- pop operation.

Class stack():

```
def __init__(self):  
    self.container = list()
```

```
def push(self, data):  
    self.container.append(data)
```

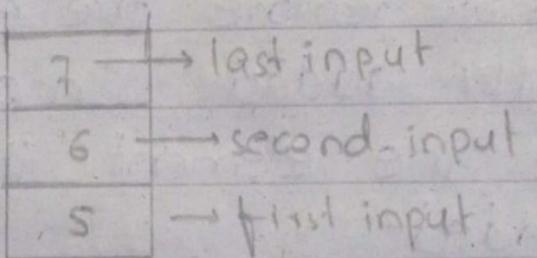
```
def pop(self, da):  
    self.container.pop(0)
```

```
def show_stack(self): : 303030 .  
    return self.container
```

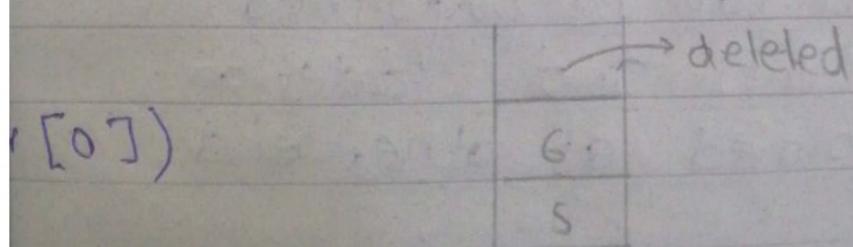
* Stack is a linear data structure in which the insertion and deletion are allowed only at end, called the top of the stack.

OPERATIONS OF STACK:-

• **PUSH (DATA)**: Insert data onto stack.



• **POP()**:- Deletes the last inserted element from the stack.



• QUEUE:

Queue is a linear data structure which follows a particular order in which the operations are performed.

The order may be FIFO.

(First in first out). A good example of queue is any queue of consumers for a resource where the consumer that came first is served first. The difference between stacks and queues is in removing.

In stacks, we remove the element the most recently added. In queues, we remove the element the least recently added.

Or

A list or collection with restriction that insertion can be performed at one end (tail) and deletion can be performed at other end (head).

$$\begin{array}{r}
 0 \\
 n=4 \quad 4 \ 3 \ 2 \ 1 \ \underline{\text{len}-1} \\
 n=1 \quad 0 \ 1 \ 2 \ 3
 \end{array}$$

operations of queue:

- * enqueue

→ ~~first in last out~~

- * dequeue

- * Insertion at one end and deletion at other end.

class Queue():

```
def __init__(self):
    self.container = list()
```

```
def enqueue(self, data):
    self.container.append(data)
```

```
def dequeue(self):
```

```
    self.container.pop(0)
    n = len(self.container)
```

```
    self.container.pop(n-1)
```

```
def show-queue(lj):
    self
    return self.container
```

DYNAMIC SET:-

A set that supports insertion or deletion of elements. This data structure is frequently used in database access.

```
class dynamic-set():
    def __init__(self):
        self.set = list()

    def insert(self, data):
        self.set.append(data)

    def search(self, key):
        n = len(self.set)
        for i in range(n):
            if self.set[i] == key:
                print("KEY IS PRESENT")
                return
```

Print ("KEY IS NOT PRESENT")

def delete(self):

 self.set.remove(self.set[0])

def maximum(self):

 n = len(self.set)

 for i in range(n):

 for j in range(0, n-i-1):

 if self.set[j] > self.set[j+1]:
 self.set[j], self.set[j+1] =
 self.set[j+1], self.set[j]

 return self.set[n-1]

def minimum(self):

 n = len(self.set)

 for i in range(n):

 for j in range(n-i-1):

 if self.set[j] > self.set[j+1]:
 self.set[j], self.set[j+1] =
 self.set[j+1], self.set[j]

return self.set[0]