# **Implement Your Own Type System** in 45 Minutes The Checker Framework



http://CheckerFramework.org/

Twitter: @CheckerFrmwrk

Live demo: http://CheckerFramework.org/live/

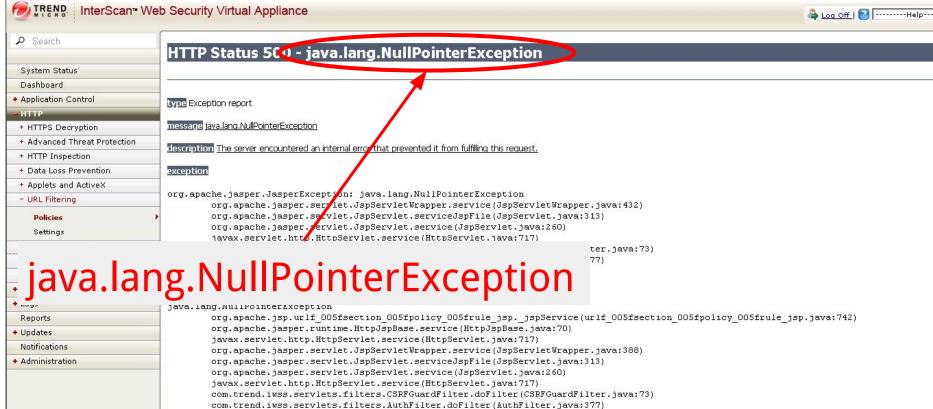
Werner Dietl, University of Waterloo Michael Ernst, University of Washington





#### **Motivation**





### Java's type system is too weak

Type checking prevents many errors
int i = "hello";

Type checking doesn't prevent enough errors

System.console().readLine();



# Java's type system is too weak

```
Type checking prevents many errors
int i = "hello";
```

```
Type checking (NullPointerException System.console().readLine();
```



### Prevent null pointer exceptions

Java 8 introduces the Optional<T> type

- Wrapper; content may be present or absent
- Constructor: of(T value)
- Methods: boolean isPresent(), T get()

Optional<String> maidenName;



### Optional reminds you to check

```
Without Optional:
                              With Optional:
         possible
                                       possible
         NullPointerException
                                       NoSuchElementException
                              Optional < St omName;
String mN/
                              omName.get().equals(...);
mName.equals(...);
                              if (omName. Present()) {
if (mName != null) {
                                              guals(...);
                                omName.ge
  mName.equals(...);
```

possible NullPointerException

Complex rules for using Optional correctly!

## **How <u>not</u> to use Optional**

#### Stuart Marks's rules:

- Never, ever, use null for an Optional variable or return value.
- 2. Never use Optional.get() unless you can prove that the Optional is present.
- 3. Prefer alternative APIs over Optional.isPresent() and Optional.get().
- 4. It's generally Let's enforce the chaining met
- 5. If an Optiona result of Optiona

nal for the specific purpose of

chain, or has an intermediate

- 6. Avoid using Optional in fields, method parameters, and collections.
- 7. Don't use an Optional to wrap any collection type (List, Set, Map). Instead, use an empty collection to represent the absence of values.





#### Which rules to enforce with a tool

#### Stuart Marks's rules:

- 1. **Never**, ever, use null for an Optional variable or return value.
- 2. **Never** use Optional.get() unless you can prove that the Optional is present.
- 3. *Prefer* alternative APIs over Optional.isPresent() and Optional.get().
- 4. It's *generally a bad idea* to create an Optional for the specific purpose of chaining methods from it to get a value.
- 5. If an Optional chain has a nested Optional chain, or has an intermediate result of Optional, it's *probably too complex*.
- 6. Avoid using Optional in fields, method parameters, and collections.
- 7. **Don't** use an Optional to wrap any collection type (List, Set, Map). Instead, use an empty collection to represent the absence of values.



#### Which rules to enforce with a tool

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specific purpose of

```
Addr
Obj
                                                                                                                                                                                                                                                                     = Set of Addresses \cup {null<sub>a</sub>}
                                                                                                                                                                                                                                                                     = rType, Fields
                                                                                                                                                                                                 ∈ rType
                                                                                                                                                                                                                                                                     = OwnerAddr ClassId<\(\bar{r}\)Type>
      P \in Program ::= \overline{Class}, ClassId, \underline{Expr}
                                                                                                                                                                                                 \in
                                                                                                                                                                                                                 Fields
                                                                                                                                                                                                                                                                     = FieldId \rightarrow Addr
                                   Class
                                                                                             class ClassId<TVarId
Cls ∈
                                                                                                                                                                                               \in
                                                                                                                                                                                                                OwnerAddr
                                                                                                                                                                                                                                                                     = Addr \cup \{any_a\}
                                                                                              extends ClassId< Typ
                                                                                                                                                                                                                                                                     = TVarId rType; ParId Addr
                                                                                                                                                                                                                 <sup>r</sup>Env
                                                                                              { FieldId SType; Met
                                  <sup>s</sup>Type
                                                                                              SNType | TVarId
                                                                                                                                                                                                                                                                                                      h, {}^{r}\Gamma, e_0 \rightsquigarrow h', \iota_0
                                                                                             OM ClassId < Type >
                                   <sup>s</sup>NType
                                                                                                                                                                                                                                                                                                                  \iota_0 \neq \mathtt{null}_a
                                   OM
                                                                                                                                               h, {}^{r}\Gamma, e_0 \rightsquigarrow h_0, \iota_0
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                                   Meth
                                                                                                                                                           \iota_0 \neq \mathtt{null}_a
                                   MethSig
                                                                                                                                               h_0, {}^{\mathbf{r}}\Gamma, e_2 \rightsquigarrow h_2, \iota
                                                                                           OS-Upd \frac{h' = h_2[\iota_0.f := \iota]}{h, {}^{r}\Gamma, e_0.f = e_2 \leadsto h',}
                    ∈ Purity
                  ∈ Expr
                                                                                                                                                                                                                                                                    \Gamma \vdash e_0 : N_0 \qquad N_0 = u_0 C_0 < >
                                                                                              Expr.MethId<sType>(Expr)
                                                                                                                                                                                                                                                                                       \mathtt{T}_1 = fType(\mathtt{C}_0,\mathtt{f})
                                                                                              new SType (SType) Expr
                                                                                                                                                                                                                                                                                       \Gamma \vdash e_2 : N_0 \triangleright T_1
                                                                       εГ
                                   Env
                                                                          ::= TVarId sNType; ParId sType
   h \vdash {}^{\mathbf{r}}\Gamma : {}^{\mathbf{s}}\Gamma
   h \vdash \iota_1 : dyn({}^{\mathfrak s}N, h, {}^{\mathfrak l}_{1,1})
   h \vdash \iota_2 : dyn(^{\mathfrak{s}}\mathsf{T}, \iota_1, h(\iota_1)\downarrow_1)
                                                                                                                                     \implies h \vdash \iota_2 : dun({}^{\mathtt{s}}\mathtt{N} \triangleright^{\mathtt{s}}\mathtt{T}.h.{}^{\mathtt{r}}\Gamma)
   {}^{\mathtt{s}}\mathtt{N} = \mathtt{u}_N \; \mathtt{C}_N < \!\!\!>
                                                                                                                       u_N = this_u \Rightarrow {}^{\mathbf{r}}\Gamma(this)
                                                                                                                                                           dom(C) = \overline{X}
                                                                                                                                                                                                                                                                                    free(^{s}T) \subseteq \overline{X} \circ \overline{X'}
   free(^{s}T) \subseteq dom(C_N)
                                                                                             DYN-
                                                                                                                      dyn(^{s}T, \iota, ^{r}T, (X' ^{r}T'; \bot)) = {^{s}T[\iota'/this, \iota'/peer, \iota/rep, any, \iota/rep, any,
```

Heap

= Addr  $\rightarrow$  Obj



- 1. **Type hierarchy** (subtyping)
- 2. **Type rules** (what operations are illegal)
- 3. **Type introduction** (what types for literals, ...)
- 4. **Dataflow** (run-time tests)

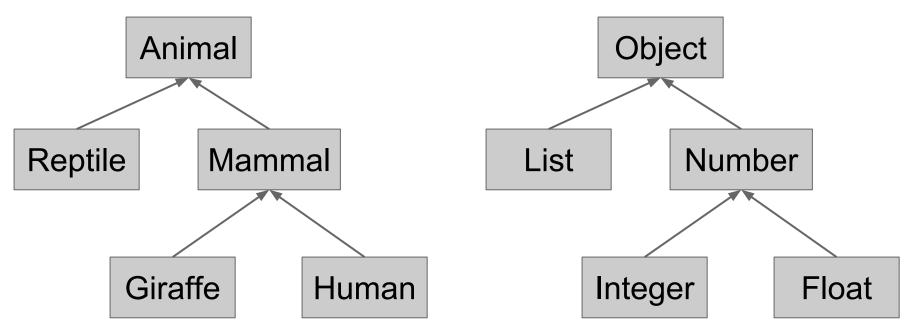
We will define two type systems: Nullness and Optional



- 1. Type hierarchy (subtyping)
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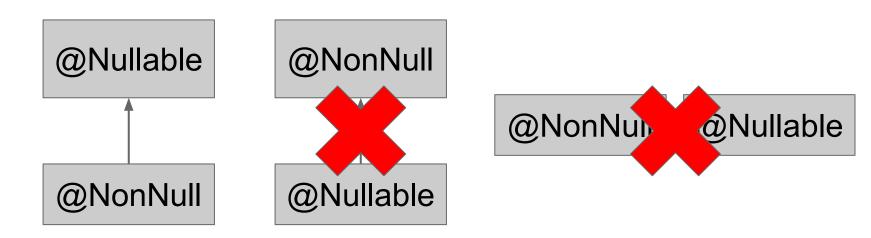
# 1. Type hierarchy



- the types
- their relationships (lower = fewer values, more properties)



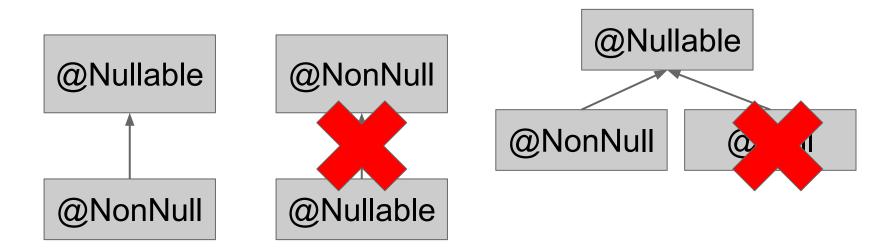
## Type hierarchy for nullness



- the types
- their relationships



## Type hierarchy for nullness

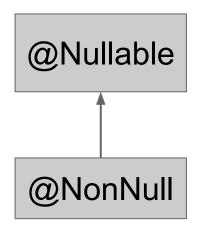


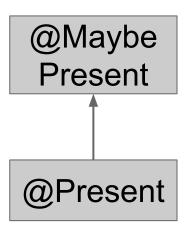
- the types
- their relationships



# **Type hierarchy for Optional**

"Never use Optional.get() unless you can prove that the Optional is present."





- the types
- their relationships



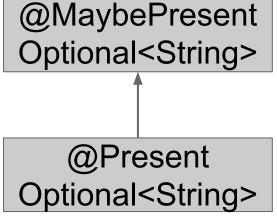
# Type = type qualifier + Java basetype



Default qualifier = @MaybePresent

Type

- @MaybePresent Optional<String>
- Optional<String>



equivalent



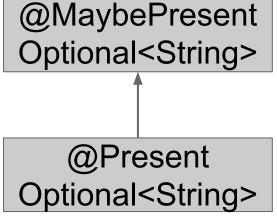
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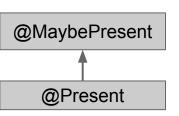


### 2. Type rules

To prevent <u>null pointer exceptions</u>:

- expr.field
   expr.getValue()
   receiver must be non-null
- synchronized (expr) { ... }
   monitor must be non-null
- ...





"Never use Optional.get() unless you can prove that the Optional is present."

Only call Optional.get() on a receiver of type

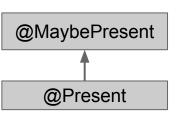
@Present Optional.

```
class Optional<T> {
   T get() { ... }
}
```

```
example call:
myOptional.get()
```

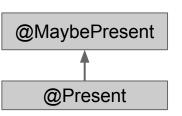
```
example call:
a.equals(b)
```





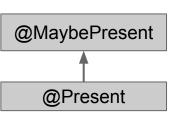
```
Only call Optional.get() on a receiver of type
@Present Optional.
                           example call:
                           myOptional.get()
class Optional<T> {
  T get(Optional<T> this) { ... }
```





```
Only call Optional.get() on a receiver of type
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                          example call:
                          myOptional.get()
class Optional<T> {
  T get(@Present Optional<T> this) {...}
```





```
Only call Optional.get() on a receiver of type
@Present Optional.
                           example call:
                           myOptional.get()
class Optional<T> {
  T get(@Present Optional<T> this) {...}
  T orElseThrow(@Present O... this, ...) {...}
```

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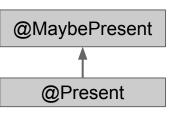
### Type introduction rules

For Nullness type system:

- null : @Nullable
- "Hello World" : @NonNull



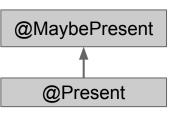
# **Type introduction for Optional**



```
Optional<T> of(T value) {...}
Optional<T> ofNullable(T value){...}
```



# **Type introduction for Optional**



```
@Present Optional<T> of(T value) {...}
Optional<T> ofNullable(@Nullable T value){...}
```



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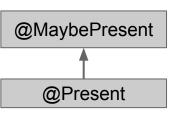


## Flow-sensitive type refinement

After an operation, give an expression a more specific type

```
@Nullable Object x;
if (x != null) {
    ...    x is @NonNull here
}
    y = new SomeType();
    ...    y is @NonNull here
y = unknownValue;
    ...    y is @Nullable again
...    y is @Nullable again
```

# **Type refinement for Optional**



"Never use Optional.get() unless you can prove that the Optional is present."

After receiver.isPresent() returns true, the receiver's type is @Present

```
@MaybePresent Optional<String> x;
if (x.isPresent()) {
    ...     x is @Present here
}
...     x is @MaybePresent again
```



### Now, let's implement it

Follow the instructions in the Checker Framework Manual

https://checkerframework.org/manual/#creating-a-checker



### You can use the Optional Checker

Distributed with the Checker Framework Checks 6 of the 7 rules for using Optional



#### More at CodeOne 2018

Preventing Errors Before They Happen: The Checker Framework TUT4991, Tue Oct 23, 8:45 – 10:45 Moscone West - Room 2009

Using Type Annotations to Improve Your Code BOF4992, Tue Oct 23, 19:30 – 20:15 Moscone West - Room 2009



# Pluggable type-checking improves code

Checker Framework for creating type checkers

• Featureful, effective, easy to use, scalable

Prevent bugs at compile time

Create custom type-checkers

Improve your code!

http://CheckerFramework.org/

