```
print(@Readonly Object x) {
  List<@NonNull String> lst;
  ...
}
```

# Preventing errors before they happen: Lightweight verification via pluggable type-checking

Michael D. Ernst

University of Washington (Seattle, WA, USA)
University of Buenos Aires
Joint work with Werner Dietl and many others

http://CheckerFramework.org/

## **Schedule**

- Part 1 (11:00 12:30)
  - pluggable type-checking: what and why
  - demo of the Checker Framework
  - relevance to your programming problems
- Part 2 (14:00 15:30)
  - how to create your own type system
  - hands-on practice in using pluggable types

#### **Motivation**



#### Software bugs cost money

\$312 billion per year (2013)

\$440 million loss by Knight Capital Group in 30 minutes

\$6 billion: 2003 blackout in northeastern USA & Canada

#### Software bugs cost lives

2003: 11 deaths: blackout

1997: 225 deaths: jet crash caused by radar software

1991: 28 deaths: Patriot missile guidance system

1985-2000: >8 deaths: Radiation therapy



# Java's type checking is too weak

• Type checking prevents many bugs int i = "hello"; // type error

Type checking doesn't prevent enough bugs

```
System.console().readLine();

⇒ NullPointerException

Collections.emptyList().add("One");

⇒ UnsupportedOperationException
```

#### Some errors are silent

```
Date date = new Date(0);
myMap.put(date, "Java epoch");
date.setYear(70);
myMap.put(date, "Linux epoch");

⇒ Corrupted map

dbStatement.executeQuery(userInput);
⇒ SQL injection attack
```

Initialization, data formatting, equality tests, ...

Goal: Find errors at **compile time** ... before testing, customers, or hackers find them

## Solution: Pluggable type systems

- Design a type system to solve a specific problem
- Write type qualifiers in code (or, use type inference)

```
@Immutable Date date = new Date(0);
date.setTime(70); // compile-time error
```

Type checker warns about violations (bugs)

#### **Outline**

- Type qualifiers
- Pluggable type checkers
- Writing your own checker
- Verification vs. bug finding
- Conclusion

## **Type qualifiers**

• In Java 8: annotations on types

```
@Untainted String query;
List<@NonNull String> strings;
myGraph = (@Immutable Graph) tmpGraph;
@English String @ReadOnly [] words;
class UnmodifiableList<T>
  implements @Readonly List<@Readonly T> {}
```

<u>Backward-compatible</u>: with any Java compiler
 List</\*@NonNull\*/ String> strings;

## Benefits of type qualifiers

#### Find bugs in programs

Guarantee the absence of errors

#### Improve documentation

Improve code structure & maintainability

Aid compilers, optimizers, and analysis tools

Reduce number of run-time checks

#### Possible negatives:

- Must write the types (or use type inference)
- False positives are possible (can be suppressed)

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## Using a checker

- Run in IDE or on command line
- Works as a compiler plug-in (annotation processor)
- Familiar workflow and error messages

```
Console console = System.console();
console.printf("Password: ");
char[] password = console.readPassword();

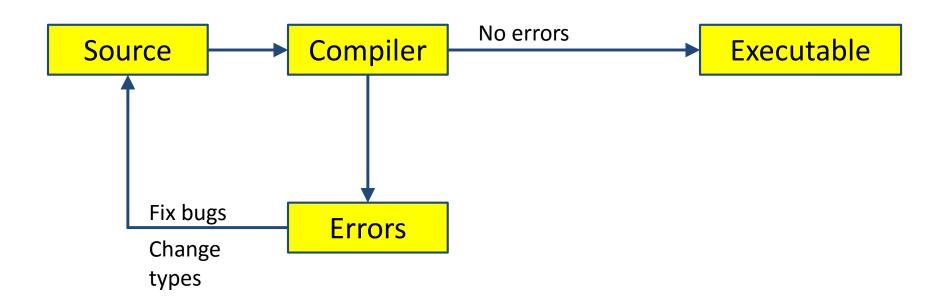
MyF

Problems 
@ Javadoc Declaration
0 errors, 1 warning, 0 others

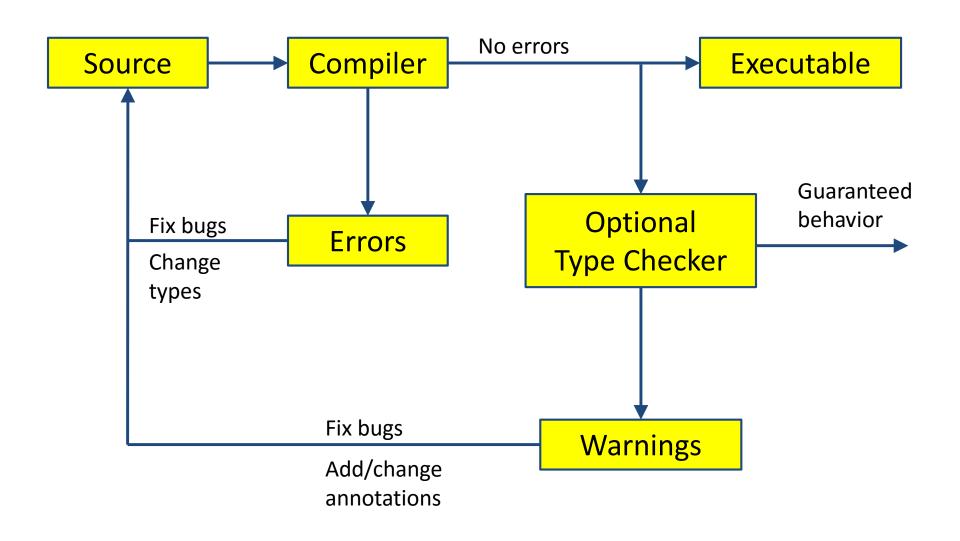
Description

Warnings (1 item)
deference of possibly-null reference console
```

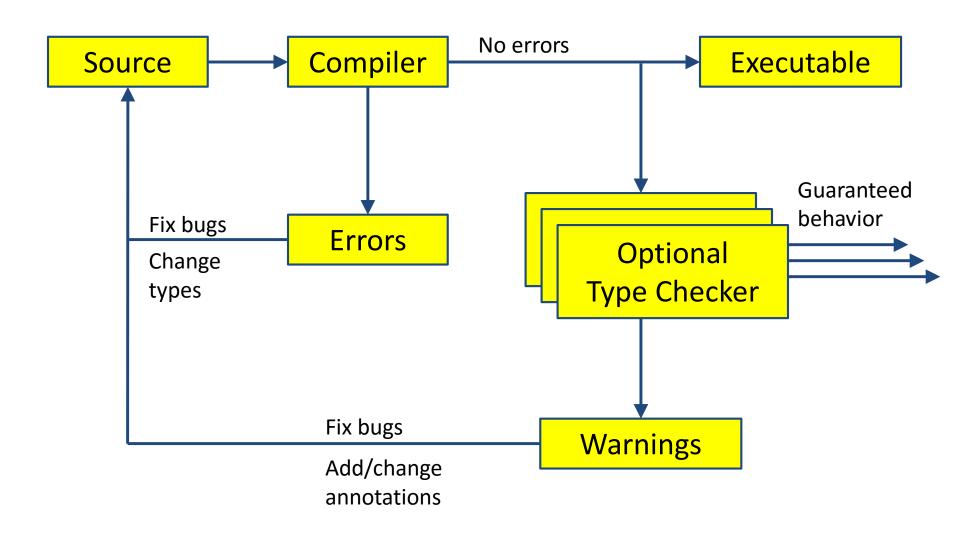
# **Type Checking**



# **Optional Type Checking**



# **Optional Type Checking**



## **Nullness and mutation demo**

- Detect errors
- Guarantee the absence of errors
- Verify the correctness of optimizations

## **Checkers are effective**

**Practical**: in daily use at Google, on Wall Street, etc.

Scalable: > 6 MLOC checked at UW

#### Selected case study results:

- Signature strings: 28 errors in OpenJDK, ASM, AFU
- Nullness: >200 errors in Google Collections, javac, Daikon
- Interning: >200 problems in Xerces, Lucene
- Format strings: 104 errors, only 107 annotations required
- Regular expressions: 56 errors in Apache, etc.; 200 annos
- Fake enumerations: problems in Swing, JabRef
- Compiler messages: 8 wrong keys in Checker Framework

## **Comparison: other nullness tools**

	Null pointer errors		False	Annotations
	Found	Missed	warnings	written
Checker				
Framework	8	0	4	35
FindBugs	0	8	1	0
Jlint	0	8	8	0
PMD	0	8	0	0

- Checking the Lookup program for file system searching (4KLOC)
- False warnings are suppressed via an annotation or assertion

## Checkers are featureful

- Full type systems: inheritance, overriding, generics (type polymorphism), etc.
- Type qualifier polymorphism
- Flow-sensitive type qualifier inference
  - no need to write annotations within method bodies
- Qualifier defaults
- Pre-/post-conditions, side effect annotations
- Warning suppression

#### Checkers are usable

- Integrated with toolchain
  - javac, Eclipse, Ant, Maven
- Annotations are not too verbose
  - @NonNull: 1 per 75 lines
    - with program-wide defaults, 1 per 2000 lines
  - @Interned: 124 annotations in 220 KLOC revealed 11 bugs
  - @Format: 107 annotations in 2.8 MLOC revealed 104 bugs
  - Possible to annotate part of program
  - Fewer annotations in new code
- Inference tools add annotations to your program
- Few false positives
- First-year CS majors preferred using checkers to not

## What a checker guarantees

- The program satisfies the type property. There are:
  - no bugs (of particular varieties)
  - no wrong annotations
- Caveat 1: only for code that is checked
  - Native methods
  - Reflection
  - Code compiled without the pluggable type checker
  - Suppressed warnings
    - Indicates what code a human should analyze
  - Checking part of a program is still useful
- Caveat 2: The checker itself might contain an error

## Formalizations

```
Set of Addresses \cup \{\text{null}_a\}
                                                                                                                           Addr
                                                                                                                           Obj
                                                                                                                                                                   Type, Fields
                                                                                                                           <sup>r</sup>Type
                                                                                                                                                                    OwnerAddr ClassId<Type>
                                         ::= Class, ClassId, Expr
            \in Program
                                                                                                                  \in
                                                                                                                           Fields
                                                                                                                                                                   FieldId \rightarrow Addr
                                                       class ClassId<TVarId
Cls
                     Class
                                                                                                                                                                   Addr \cup \{any_a\}
                                                                                                                  \in
                                                                                                                           OwnerAddr
                                                       extends ClassId<sTyp
                                                                                                                  \in
                                                                                                                                                                   TVarId Type; ParId Addr
                                                                                                                           <sup>r</sup>Env
                                                       { FieldId SType; Met
                                                       *NType | TVarId
             ∈ <sup>s</sup>Type
                                                                                                                                                                             h, {}^{r}\Gamma, e_0 \rightsquigarrow h', \iota_0
            \in {}^{\mathtt{s}}\mathtt{NType}
                                                       OM ClassId<SType>
                                                                                                                                                                                  \iota_0 \neq \mathtt{null}_a
                     OM
                                                                                  h, {}^{r}\Gamma, e_0 \rightsquigarrow h_0, \iota_0
                                           ::=
                                                                                                                                               OS-Read \frac{\iota = h'(\iota_0) \downarrow_2 (f)}{h, {}^{\mathbf{r}}\Gamma, e_0.f \leadsto h', \iota}
  mt
                     Meth
                                           ::=
                                                                                    \iota_0 \neq \mathtt{null}_a
                     MethSig
                                                                                    h_0, {}^{\mathbf{r}}\Gamma, e_2 \rightsquigarrow h_2, \iota
            \in Purity
             \in Expr
                                                                                                                                                        \Gamma \vdash e_0 : N_0 \qquad N_0 = u_0 C_0 < >
                                                       Expr.MethId<sType>(Expr) |
                                                                                                                                                                    T_1 = fType(C_0, f)
                                                       new SType | (SType) Expr
                                                                                                                                                                      \Gamma \vdash e_2 : \mathbb{N}_0 \triangleright \mathbb{T}_1
                                           <sup>s</sup>Env
                                           ::= TVarId sNType; ParId sType
 h \vdash {}^{\mathbf{r}}\Gamma : {}^{\mathbf{s}}\Gamma
 h \vdash \iota_1 : dyn({}^{\mathfrak s} \mathbb{N}, h, {}^{\mathfrak l}_{\mathfrak s})
 h \vdash \iota_2 : dyn({}^{\mathtt{s}}\mathtt{T}, \iota_1, \mathbf{h}(\iota_1) \downarrow_1)
                                                                              \implies h \vdash \iota_2 : dyn({}^{\mathtt{s}}\mathsf{N} \triangleright {}^{\mathtt{s}}\mathsf{T}, \mathsf{h}, {}^{\mathtt{r}}\Gamma)
  {}^{\mathtt{s}}\mathtt{N} = \mathtt{u}_N \; \mathtt{C}_N < >
                                                                      {}^{\mathbf{r}}\mathbf{T} = \iota' _<> \iota \vdash {}^{\mathbf{r}}\mathbf{T} \stackrel{\mathbf{r}}{<}: \iota' \stackrel{\mathbf{r}}{<}\overline{\mathbf{T}} > \iota \vdash {}^{\mathbf{r}}\mathbf{T} \stackrel{\mathbf{r}}{<}: \iota' \stackrel{\mathbf{r}}{<}\overline{\mathbf{T}} > \Rightarrow \iota \vdash \overline{{}^{\mathbf{r}}}\mathbf{T} \stackrel{\mathbf{r}}{<}: \overline{{}^{\mathbf{r}}}\mathbf{T}_a
 \mathtt{u}_N = \mathtt{this}_u \Rightarrow {}^{\mathbf{r}}\Gamma(\mathtt{this})
                                                                                                                                                \mathrm{free}({}^{\mathtt{s}}\mathtt{T})\subseteq \overline{\mathtt{X}}\circ \overline{\mathtt{X}'}
                                                                                           dom(C) = \overline{X}
 free({}^{\mathtt{s}}\mathtt{T}) \subseteq dom(\mathtt{C}_N)
                                                     DYN-
                                                                     dyn({}^{\mathtt{s}}\mathtt{T},\iota,{}^{\mathtt{r}}\mathtt{T},(\overline{\mathtt{X}'\ {}^{\mathtt{r}}\mathtt{T}'};\underline{\ })) = {}^{\mathtt{s}}\mathtt{T}[\iota'/\mathtt{this},\iota'/\mathtt{peer},\iota/\mathtt{rep},\mathtt{any}_a/\mathtt{any}_u,\overline{{}^{\mathtt{r}}\mathtt{T}/\mathtt{X}},\overline{{}^{\mathtt{r}}\mathtt{T}'/\mathtt{X}'}]
```

Heap

 $Addr \rightarrow Obj$ 

## **Annotating libraries**

- Each checker comes with JDK annotations
  - For signatures, not bodies
  - Finds errors in clients, but not in the library itself
- Inference tools for annotating new libraries

## What bugs can you detect & prevent?

The property you care about:

Null dereferences

Mutation and side-effects

Concurrency: locking

Security: encryption,

tainting

Aliasing

**Equality tests** 

Strings: localization, regular expression syntax, signature representation, format string syntax

**Enumeractions** 

Typestate (e.g., open/closed files)

Users can write their own checkers!

The annotation you write:

@NonNull

@Immutable

@GuardedBy

@Encrypted

@OsTrusted, @Untaint...

@Linear

@Interned

@Localized

@Regex

@FullyQualified

@Format

@Fenum

@State

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## **Example: Regular expressions**

```
// Prints the first matching group.
// For example:
     java RegexExample ([0-9]*):([0-9]*) 23:59
// prints "Group 1 = 23"
public static void main(String[] args) {
  String regex = args[0];
                             PatternSyntaxException
  String content = args[1];
  Pattern pat = Pattern.compile(regex);
 Matcher mat = pat.matcher(content);
  if (mat.matches()) {
    System.out.println("Group 1 = "
                        + mat.group(1));
                              IndexOutOfBoundsException
```

## Regular expression type system

- What runtime errors to prevent?
   PatternSyntaxException and IndexOutOfBoundsException.
- What operations are legal?
   Pattern.compile only on valid regex.
   Matcher.group(i) only if >i groups.
- What properties of data should hold?
   Strings: valid regex vs. invalid.
   Number of groups in a regex.

## **Example: Encrypted communication**

```
void send(@Encrypted String msg) {...}
@Encrypted String msg1 = ...;
send(msg1); // OK
String msg2 = ...;
send(msg2); // Warning!
```

## **Encryption type system**

- What runtime exceptions to prevent?
   Invalid information flow.
- What operations are legal?
   send() only on encrypted data.
- What properties of data should hold?
   Separate encrypted from plaintext strings.

## **Brainstorming new type checkers**

- What runtime exceptions to prevent?
- What operations are legal and illegal?
- What properties of data should hold?

- Type-system checkable properties:
  - Dependency on values
  - Not on program structure, timing, ...

## **Brainstorming**

- What runtime exceptions to prevent?
- What operations are legal and illegal?
- What properties of data should hold?

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## **SQL** injection attack

 Server code bug: SQL query constructed using unfiltered user input

```
query = "SELECT * FROM users "
+ "WHERE name='" + userInput + "';";
```

- User inputs: a' or '1'='1
- Result:

```
query \Rightarrow SELECT * FROM users
WHERE name=\a' or \1'=\1';
```

Query returns information about all users

#### **Taint checker**

```
@TypeQualifier
@SubtypeOf(Unqualified.class)
@ImplicitFor(trees = {STRING_LITERAL})
public @interface Untainted { }
```

#### To use it:

- 1. Write @Untainted in your program
  List getPosts(@Untainted String category) {...}
- 2. Compile your program

```
javac <u>-processor BasicChecker -Aquals=Untainted</u>
MyProgram.java
```

#### Taint checker demo

- Detect SQL injection vulnerability
- Guarantee absence of such vulnerabilities

# Defining a type system

```
@TypeQualifier
public @interface NonNull { }
```

## Defining a type system

- 1. Qualifier hierarchy
- rules for assignment
- 2. Type introduction
- types for expressions

3. Type rules

checker-specific errors

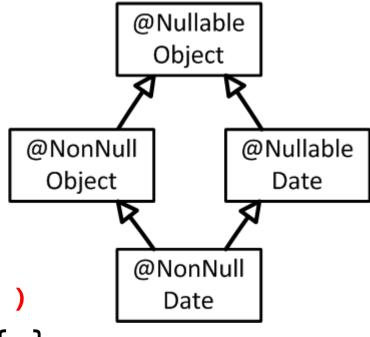
```
@TypeQualifier
public @interface NonNull { }
```

## Defining a type system

- 1. Qualifier hierarchy
- 2. Type introduction
- 3. Type rules

```
@TypeQualifier
@SubtypeOf( Nullable.class )
public @interface NonNull { }
```

What assignments are legal:



## Defining a type system

- 1. Qualifier hierarchy
- 2. Type introduction
- 3. Type rules

Gives the type of expressions:

```
new Date()
"hello " + getName()
Boolean.TRUE
```

## Defining a type system

- 1. Qualifier hierarchy
- 2. Type introduction
- 3. Type rules

#### Errors for unsafe code:

```
synchronized (expr) {
    ...
}
Warn if expr may be null
```

```
void visitSynchronized(SynchronizedTree node) {
   ExpressionTree expr = node.getExpression();
   AnnotatedTypeMirror type = getAnnotatedType(expr);
   if (! type.hasAnnotation(NONNULL))
      checker.report(Result.failure(...), expr);
}
```

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### Verification

- Goal: prove that no bug exists
- Specifications: user provides
- False negatives: none
- False positives: user suppresses warnings
- Downside: user burden

## **Bug-finding**

- Goal: find some bugs at low cost
- **Specifications**: infer likely specs
- False negatives: acceptable
- False positives: heuristics focus on most important bugs
- Downside: missed bugs

Neither is "better"; each is appropriate in certain circumstances. The approaches are converging.

## Other design considerations

- Visibility of specifications and warning suppressions
  - In the source code
    - documentation aids programer understanding
  - In the tool
    - reduces code clutter
- Analysis comprehensibility
  - A transparent tool gives understandable outcomes
    - requires more upfront effort; more false positives
  - An opaque tool can use more powerful analyses
    - requires more effort to understand warnings

#### **Outline**

- Type qualifiers
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- Hands-on practice
- Conclusion

## How to get started

#### 1. Write the specification

Search the Javadoc for occurrences of "null" Replace the wordy English text by @Nullable Can also search code, but no annos in methods

# **2. Run Nullness Checker**: verify/improve spec For each warning:

- Reason about whether the code is safe
- Express that reasoning as annotations
- Consider improving the code's design

## **Tips**

#### What to type-check:

- Only type-check properties that matter to you
  - Use subclasses (not type qualifiers) if possible
- Choose part of your code to type-check first
  - Eliminate raw types such as List; use List<String>

#### While you are doing type-checking:

- Write the spec first (and think of it as a spec)
- Avoid warning suppressions when possible

## Your turn to improve your code!

- 1. Choose a project you care about
- 2. Improve it
  - Apply an existing checker to your code, or
  - Create a new domain-specific type checker

#### Or, try the tutorial:

http://types.cs.washington.edu/checker-framework/tutorial

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## Pluggable type-checking

- Java 8 syntax for type annotations
- Checker Framework for creating type checkers
  - Featureful, effective, easy to use, scalable
- Prevent bugs at compile time
- Create custom type-checkers
- Learn more, or download the Checker Framework: <a href="http://CheckerFramework.org/">http://CheckerFramework.org/</a>
   <a href="http://checkerFramework">(or, web search for "Checker Framework")</a>