



Scalability in Practice: a Highly scalable web app

Yubin Xia

IPADS, Shanghai Jiao Tong University

https://www.sjtu.edu.cn

Scalable web apps are everywhere: pre-Al era









High request rate

100 billions of requests daily

Massive data

Facebook has more than 1
 billion of images uploaded
 weekly, Baidu stores tens of
 billions of web pages

Scalable web apps are everywhere: Al era

Essentially, also a web APP

- Pre-Al era: request handling = order an item
- Al era: request handling = use model to do an in

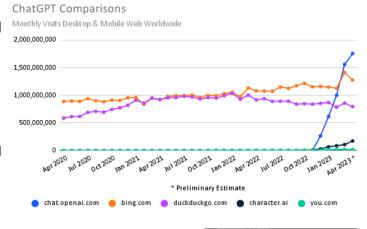
The same system requirements

- High request rate
- Massive data: both for training & for serving hist

New requirements

- Huge computation power required
- According to GPT (◄), an inference of GPT ~= sorting
 100,000,000 numbers
 Per-user sessions





Connect4 vs Tic Tac

Memory Allocation Method

☐ Teams Meeting Assistance

Sort Vector Descending

The fundamental system techniques remain similar

Al is just another (challenging) workload

Case study: a e-commerce website

Different storage systems to support different services



Computation frameworks to support different services

Hot topics



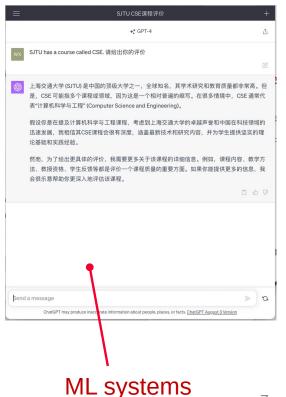
Fraud detection



Graph processing system

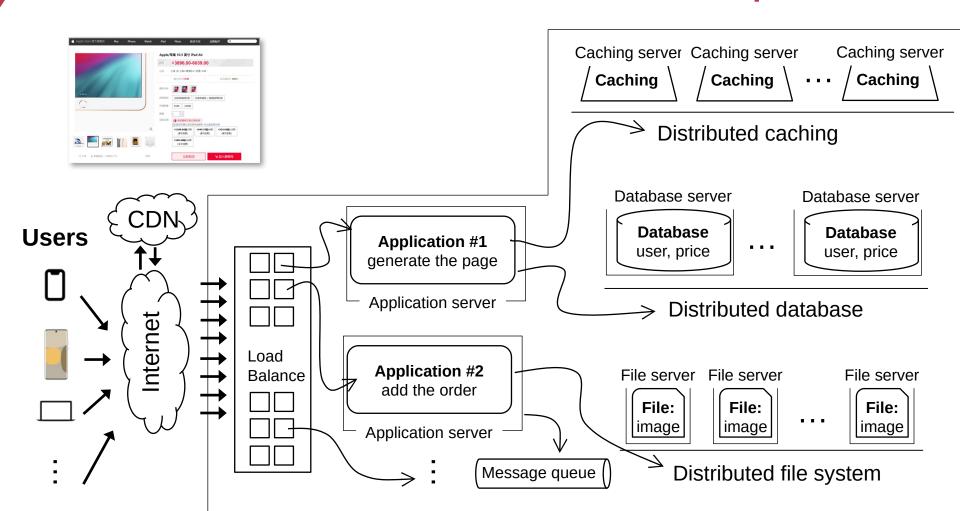
Source: https://www.graphable.ai/blog/graphdatabase-fraud-detection/

Chat Bot



Batch processing system

Each click needs thousands of servers to cooperate



How to build Taobao on a single machine (in old days)?

Operating system:

Linux (in OS class)

Serving the requests: web server

Apache, Nginx (in ICS)

Serving the data: file system & DB

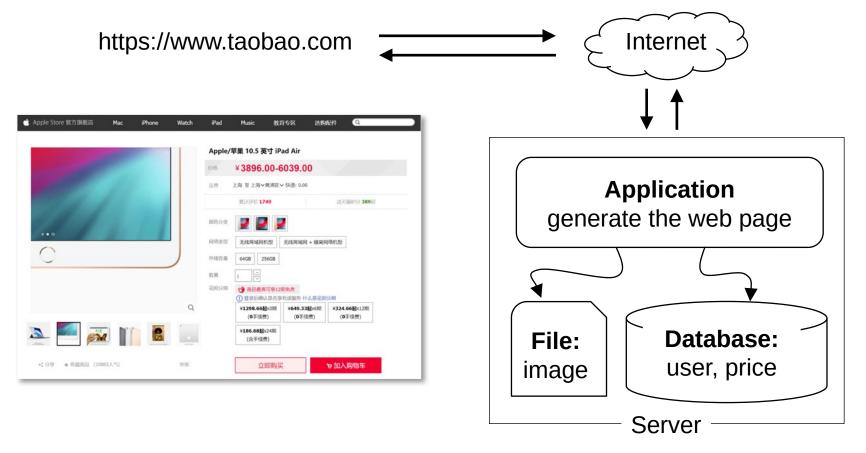
MySQL & inode file system (in CSE)

Displaying the page: HTML

PHP (in Web class)

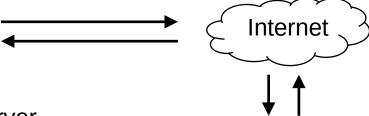


Using one server to build Taobao



The architecture of LAMP cannot scale!

https://www.taobao.com

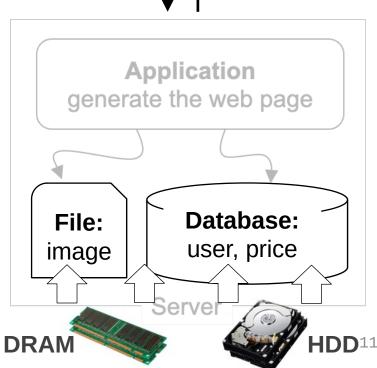


 The disk & memory of one server cannot store massive amount of data

- **DRAM**: 64 - 256 GB

- **HDD**: 2 − 40 TB

 Facebook has more than 1 billion of images uploaded weekly, Baidu stores tens of billions of web pages

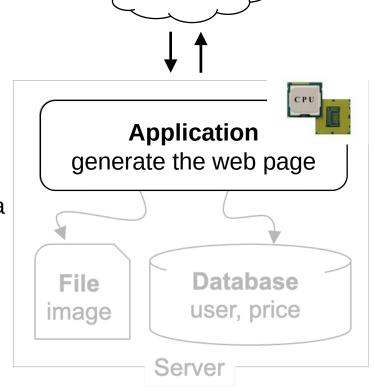


The architecture of LAMP cannot scale!

https://www.taobao.com

 Application uses the CPU for processing, while a single CPU can hardly scale due to the end of Moore's law & Dennard scaling

- Moore's law: the number of transistors in a dense integrated circuit (IC) doubles about every two years.
- Dennard scaling: as transistors get
 smaller, their power density stays constant



Internet

Step #1 for scalability: disaggregating application & data

Application: handles application logic

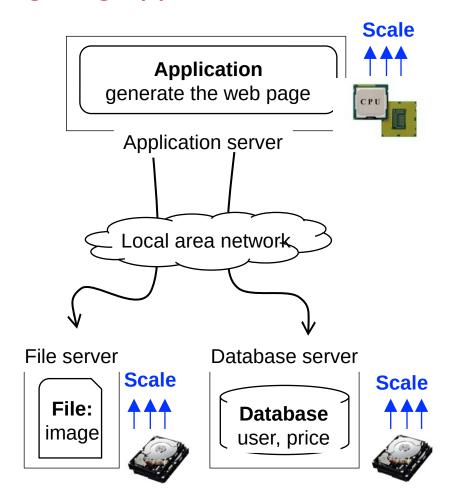
Can be scaled with more CPUs

Database: requires reading/writing disk & cache

 Can be scaled with faster disks & larger memory

File system: store large bulks of data

- E.g., images, videos
- Can be scaled with faster disks



Step #2 Avoid the slow data accesses? Caching

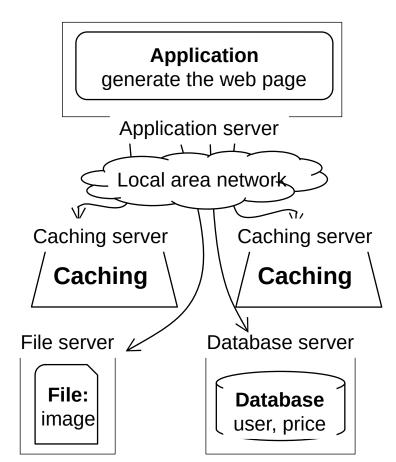
Observation: most requests access a small portion of the data (locality)

Caching system with a single node:

- E.g., page cache
- Drawbacks: limited DRAM capacity

Distributed caching server

Benefits: huge DRAM capacity,
 e.g., deploy many caching servers



Case study of distributed cache server: Memcached

Distributed caching server

Benefits: huge DRAM capacity

Memcached server

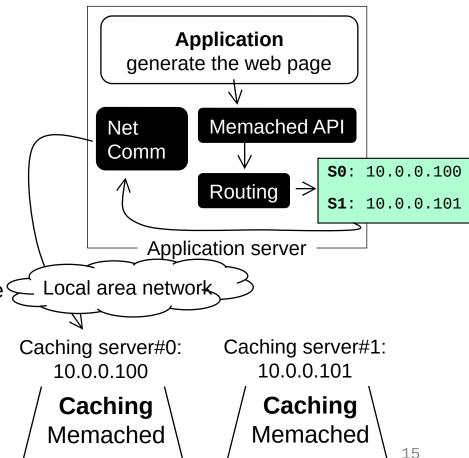
- Store all the cached data in memory
- Can scale out to use multiple servers

Memcached clients

- Check whether server has cached data
- On cache miss, fallback to database/file

Question:

How to find which server has cached the data?



Case study of distributed cache server: Memcached

Naïve method, hashing:

– address = key / #server

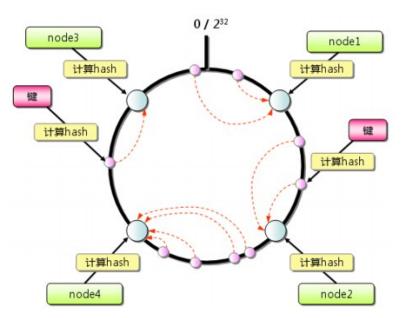
Problem:

- What if we add a new server?
- Suppose that we add server from 3 → 4
- Then there would be 75% miss!

Solution: consistent hashing

- Suppose that we add server from 3 → 4
- It will only incur 25% miss
- More details in later courses

How to find the data?



Consistent hashing

Step #3 for scalability: more servers

For **stateless** application servers

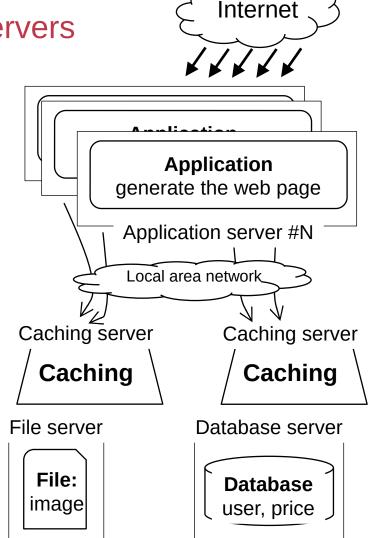
- E.g., web servers
- We can add more servers for scaling its performance

What is stateless?

 The server only executes the logic that only relies on input data but no long-term state

Benefits:

- Better fault-tolerance
- Better elasticity



Step #3 for scalability: How to do the load balance?

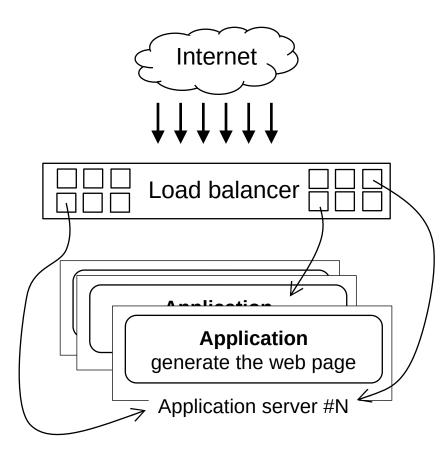
Load balance: how to route the user quests to the application servers?

Leverage the network layer

- HTTP redirection,
- reverse proxy,
- **–** ...

Load balance algorithms

- round-robin,
- random,
- hashing,
- **–** ...



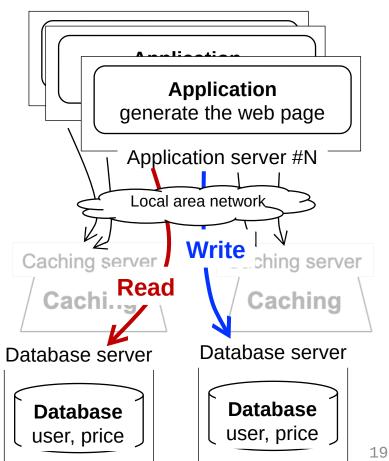
Step #4 for scalability: scaling database

#1. Separate the database for read/write

- E.g., use primary-backup replication
- The primary servers the writes and backup only serves reads

#2. Separate a table on multiple databases

- e.g., a bank has reported that a single table has **100,000,000 rows**
- However, split a table on multiple machines complex consistency management



Step #5 for scalability: using distributed file system

File:

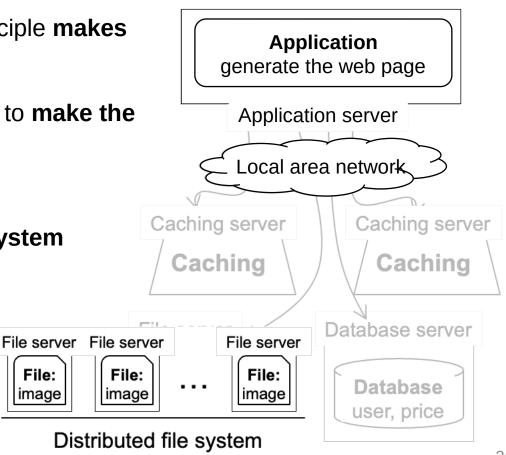
image

Using multiple database in principle makes the database distributed

We can use a similar approach to **make the** file system distributed!

Well-known distributed file system

- NFS (Network file system)
- GFS (Google file system)
- **HDFS**



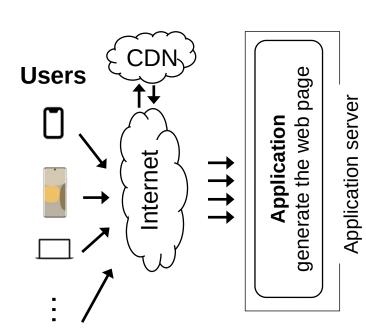
Step #6 for scalability: using CDN

Goal: return the results to user as soon as possible

Why? Amazon has reported that 100ms
 latency increase will cause 1% financial loss

Core idea: caching (again)

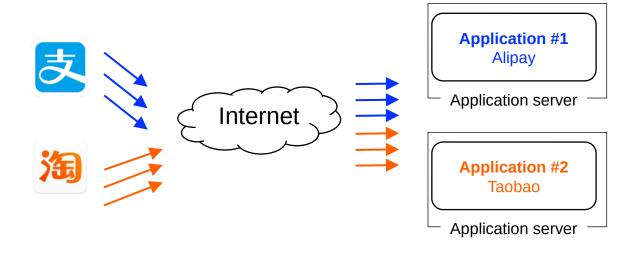
- E.g., CDN (content delivery network) caches the content at the network providers, which is closer to users
- Challenge: how to do it without users' awareness?



Step #7 for scalability: separate different applications

Use dedicated servers for different applications

- E.g., micro-services

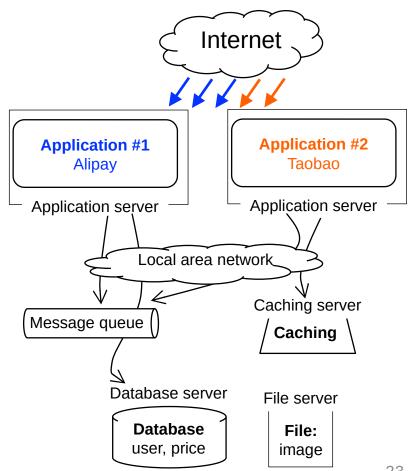


How different applications communicate? MQ or DB

Message queue (MQ): applications send the message the queue, and the queue can buffer the message (somewhere between RPC and database)

E.g., Apache Kafka

Or, applications can directly use the databases, caching (e.g. KV) or file system to communicate with shared data



How to handle complex requests?

Example: after the website becoming larger, handling requests is far more than displaying a (static) web page

Alipay (支付宝)

E.g., fraud detection

Taobao

- Hot list (热榜)
- Dynamic product ads (千人千面)
- Recommendation (you might also like)

– ...





Use distributed computing frameworks for complex queries

General batch processing systems

MapReduce, Spark, etc.

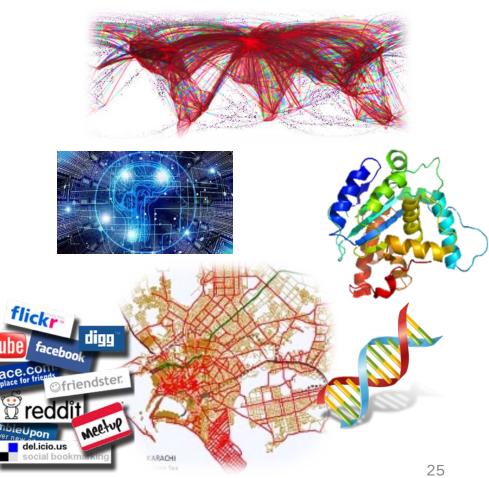
Graph processing systems

GraphLab, Pregel, etc.

Machine learning systems

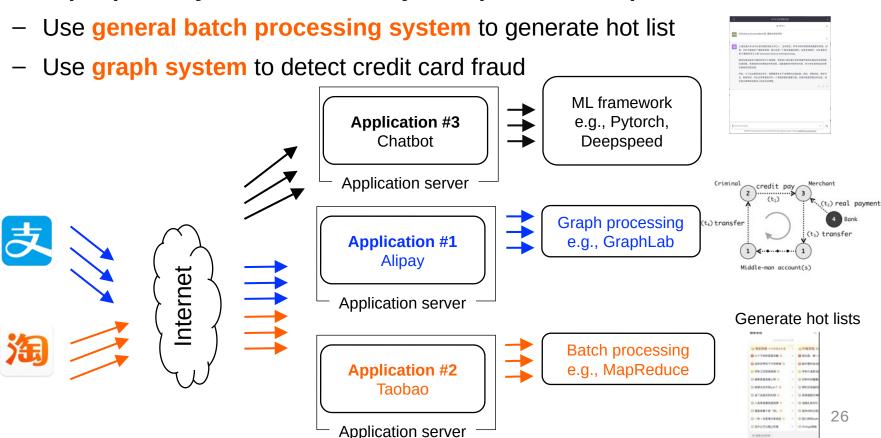
Tensorflow, pytorch, etc.



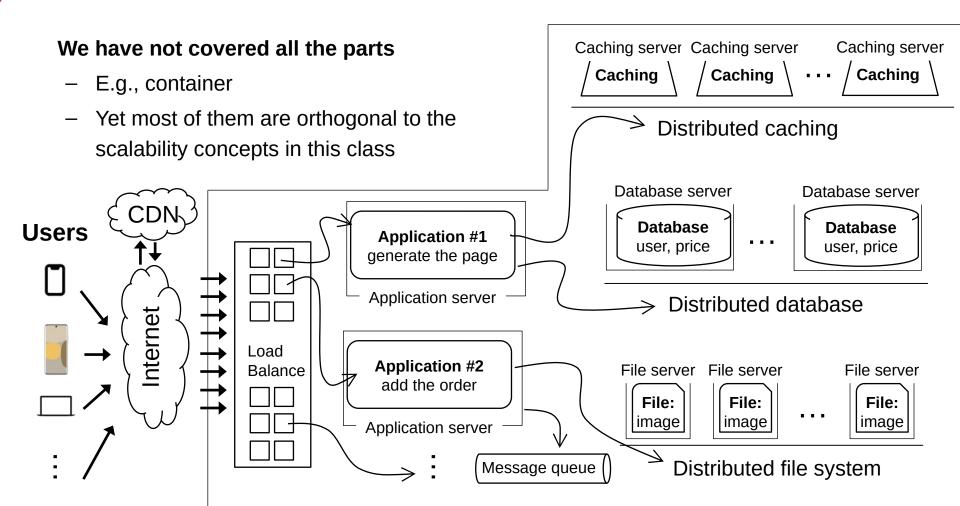


Step #7: separated applications + distributed computing

Example (Each system is backed by multiple machines)



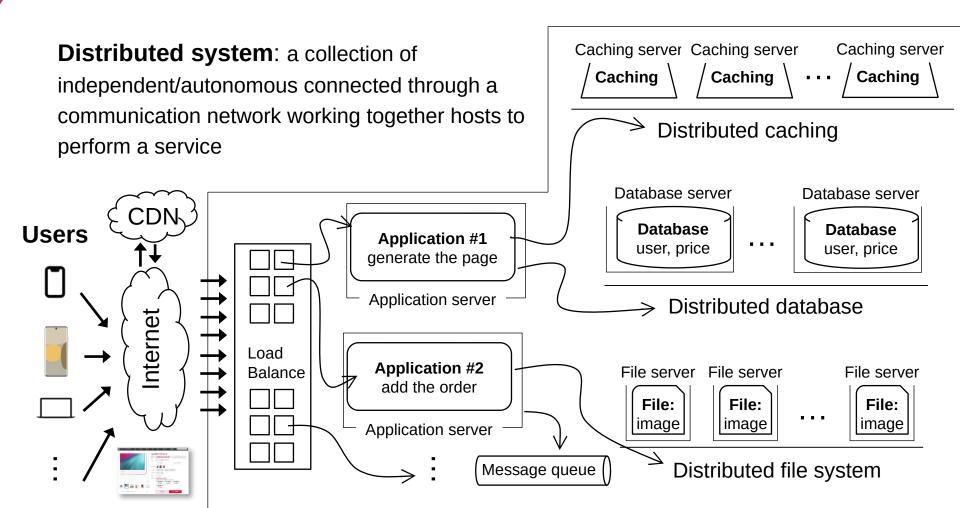
A scalable website: overall picture



Recall: we have talked about distributed system

But what is a distributed system?

A scalable website: powered by distributed systems!



Modern scalable system systems are

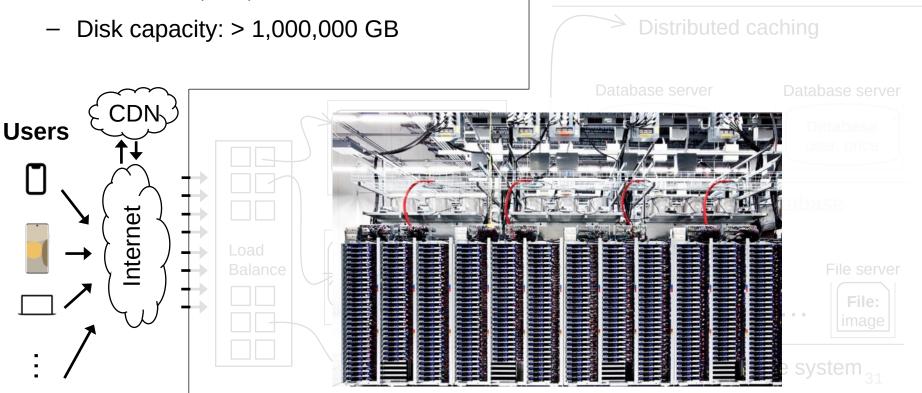
large & distributed

Powered by large-scale datacenter

Scalable websites are powered by modern datacenters

Large server and storage farms

CPUs: > 20,000,000 cores



Datacenters that power the scalable website

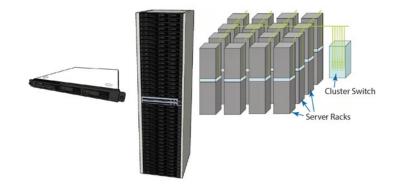
Large-scale distributed systems: 10K – 100K servers

Each rack: 40 servers

Network: 10Gbps – 100Gbps in rack

10-100 MW of power





Source: "The Datacenter as a Computer ---An Introduction to the Design of Warehouse-Scale Machines"

Geo-replicated datacenters

The **locations** of Alibaba datacenters



Regions outside Mainland China

Alibaba Cloud offers an expanding network of CDN nodes and deployment regions, including the first public cloud data center regions in the Middle East (Dubai) and Indonesia, a string of strategic data centers in Asia, and a strong presence in North America, Europe, and Australia.

O Regions in Mainland China

Data center regions in China offer BGP backbone network lines providing high-quality coverage countrywide to ensure stable and fast access inside the Mainland. In general, we recommend customers to select the data center closest to their end-users to further speed up online access.



Fault is common, how to handle it?

Fault is common: fault, error, failure

Fault can be latent or active

if active, get wrong data or control signals

Error is the results of active fault

- e.g. violation of assertion or invariant of spec
- discovery of errors is ad hoc (formal specification?)

Failure happens if an error is not detected and masked

not producing the intended result at an interface

Fault is common, especially in large distributed systems

Why faults are common especially in distributed systems? Scale!

 "Suppose a cluster has ultra-reliable server nodes with a stellar mean time between failures (MTBF) of 30 years (10,000 days)—a cluster of 10,000 servers will see an average of one server failure per day."

Fault is common

especially in large distributed systems

What are the causes?

Why faults are common especially in distributed systems? Scale!

 "Suppose a cluster has ultra-reliable server nodes with a stellar mean time between failures (MTBF) of 30 years (10,000 days)—a cluster of 10,000 servers will see an average of one server failure per day."

Causes:

- Operation error (human, configuration, etc.)
- Software error (e.g., bug)
- Hardware error
- Power outage
- Natural disaster

首页 > 新闻 > 要闻

腾讯称微信故障因市政施工挖断光缆

一财网 • 2013-07-22 17:22



▲ 关注帖子

字节一个实习生,把公司所有lite的模型都删了醫醫醫

Fault is common

especially in large distributed systems

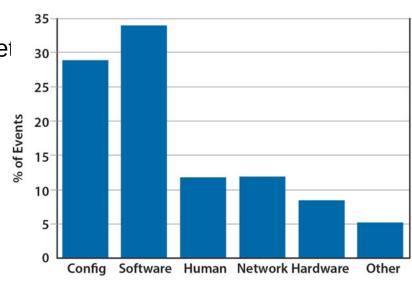
What are the causes?

Why faults are common especially in distributed systems? Scale!

 "Suppose a cluster has ultra-reliable server nodes with a stellar mean time between failures (MTBF) of 30 years (10,000 days)—a cluster of 10,000 servers will see an average of one server failure per day."

Causes:

- Operation error (human, configuration, et
- Software error (e.g., bug)
- Hardware error
- Power outage
- Natural disaster



Source: The Datacenter as a Computer

Faults and partial failures

On a single computer, it either works or it doesn't

But in a distributed system, some parts of the system can be broken in some unpredictable way

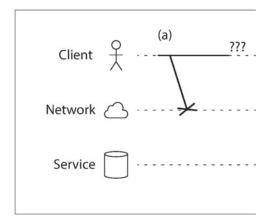
Such failure is partial (aka., grey failure)

Since **most** other parts of the system are **OK**, we want the system **still working**!

Faults and partial failures

Example: unreliable network

- A client (e.g., Smartphone),
 sends requests to the service (e.g., Taobao),
 through network (5G, Wifi)
- The client gets: "网络竟然奔溃了", how can a network break?

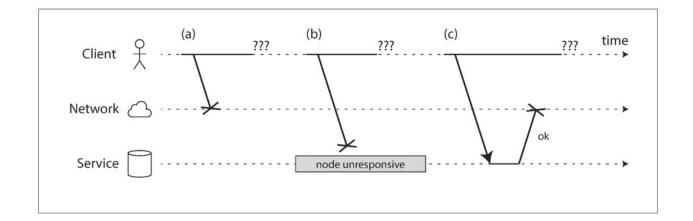




Example: unreliable network

A user sends a request but the server does not reply, possible reasons:

- 1. The request may have been **lost** (e.g., someone unplugged a network cable).
- 2. The request may be waiting in a queue and will be delivered **later** (e.g., the network or the recipient is overloaded).
- 3. The remote node may have **failed** (e.g., it crashed or it was powered down).



Example: unreliable network

A user sends a request but the server does not reply, possible reasons:

- 4. The remote node may **have temporarily stopped** responding (e.g., it is experiencing a long **garbage collection pause**)
- 5. The remote node may have processed your request, but the **response** has been **lost** on the network (e.g., a network switch has been misconfigured).
- 6. The remote node may have processed your request, but the response has been **delayed** and will be delivered later (e.g., the network or your own machine is

overloaded).

A **network partition** refers to network decomposition into relatively independent <u>subnets</u>

- Can happen when a switch is being upgraded in a datacenter
- Can even happen when being attacked by sharks

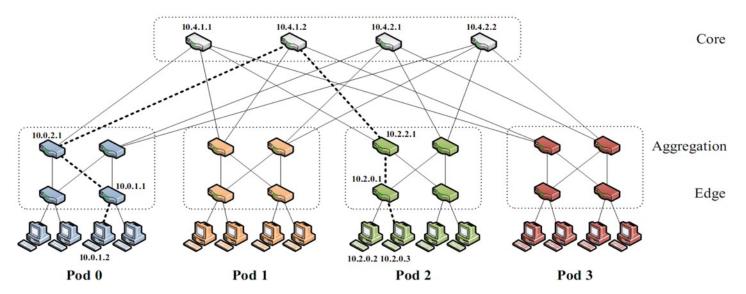
Network partition is usually a reality

You can never count on the network connectivity

A **network partition** refers to network decomposition into relatively independent <u>subnets</u>

Can happen when a switch is being upgraded in a datacenter

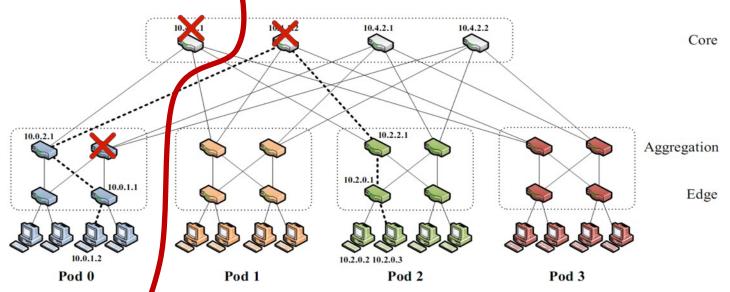
For example, suppose the datacenter adopts a **fat-tree** topology



A **network partition** refers to network decomposition into relatively independent <u>subnets</u>

Can happen when a switch is being upgraded in a datacenter

For example, suppose the datacenter adopts a **fat-tree** topology



A **network partition** refers to network decomposition into relatively independent <u>subnets</u>

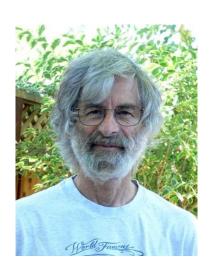
- Can happen when a switch is being upgraded in a datacenter
- Can even happen when being attacked by sharks



You know you have a distributed system when the crash of a computer you've never heard of stops you from getting any work done.

- Leslie

Lamport



Lamport received the 2013 <u>Turing Award</u> for "fundamental contributions to the theory and practice of distributed and concurrent systems, notably the invention of concepts such as causality and logical clocks, safety and liveness, replicated state machines, and sequential consistency" in 2014.





Yam Peleg 🤣 @Yampeleg · 1h **Training Cost**

GPT-4'

It is ove

Everyth

7:22 AM

OpenAl's training FLOPS for GPT-4 is ~2.15e25, on ~25,000 A100s for 90 to 100 days at about 32% to 36% MFU.

Part of this extremely low utilization is due to an absurd number of failures requiring checkpoints that needed to be restarted from.

Availability and reliability

Availability: A measure of the time that a system was usable, as a fraction of the time that it was intended to be usable (x nines), corresponding **downtime**:

- e.g. 3-nines -> 8 hour/year
- e.g. 5-nines -> 5 min/year
- e.g. 7-nines -> 3 sec/year

Metrics to measure reliability

- MTTF: mean time to failure
- MTTR: mean time to repair
- MTBF: mean time between failure
- MTBF = MTTF + MTTR

$$MTTF = \frac{1}{N} \sum_{i=1}^{N} TTF_i$$

$$MTTR = \frac{1}{N} \sum_{i=1}^{N} TTR_i$$

Large-scale systems can be highly available

Example#1: Internet, the BGP (Border Gateway Protocol) is highly available

Example#2: WeChat, Taobao & Baidu rarely (but not never) become outage

Example #3: OpenAI successfully trained GPT despite many failures ◀





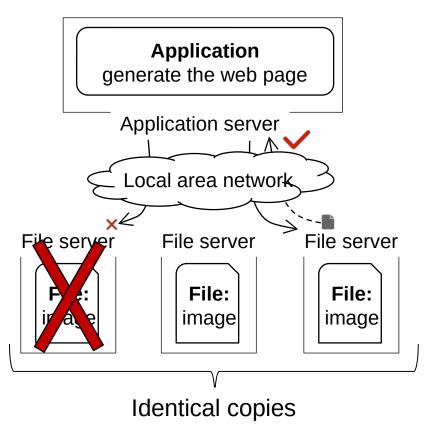
Achieving high availability: handling failures w/ replications

Replication

replicas: identical multiple copies

Example: replicated file servers

If one copy survives, the application is available



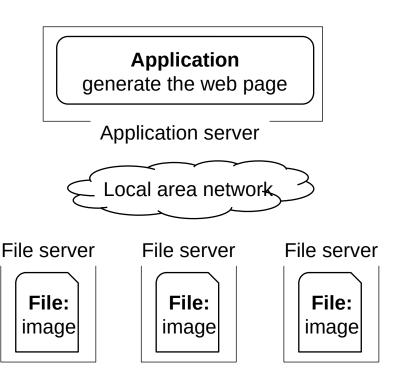
Replication

replicas: identical multiple copies

Example: replicated file servers

If one copy survives, the application is available

Challenge: consistency



Replication

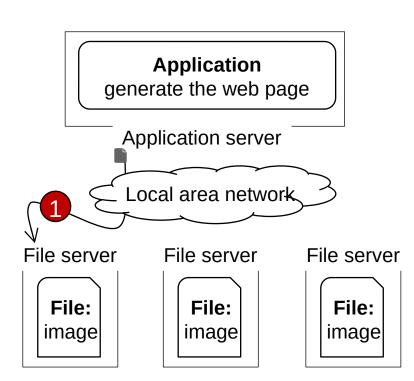
replicas: identical multiple copies

Example: replicated file servers

If one copy survives, the application is available

Challenge: consistency

1. Application put file A to one server



Replication

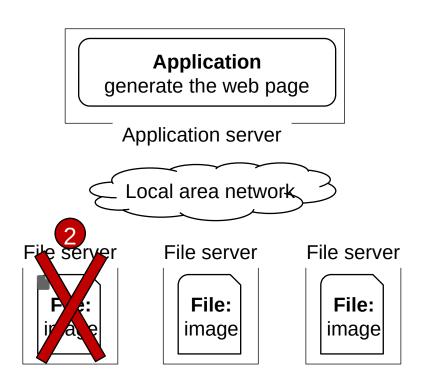
replicas: identical multiple copies

Example: replicated file servers

If one copy survives, the application is available

Challenge: consistency

- 1. Application put file A to one server
- 2. The server crashed

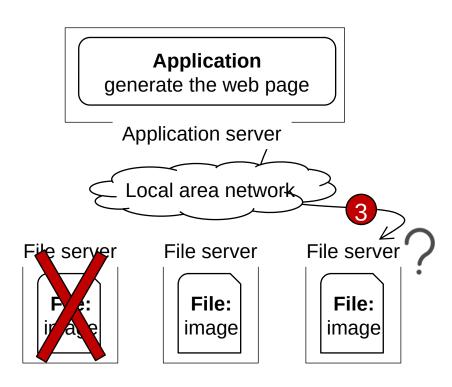


Replication

replicas: identical multiple copies

Example: replicated file servers

- If one copy survives, the application is available
- 1. Application put file A to one server
- 2. The server crashed
- 3. What happens when read A again?



Achieving high availability: handling failures

Replication

replicas: identical multiple copies

Techniques to cope with consistency:

- Primary-backup replication
- Replicated state machine
- **–** ...



Achieving high availability: handling failures via retry

Restart or reconstruct

- Monitoring and catching errors
- Restart or reconstruct the system (sub-system)
- E.g., restart the stateless application server is ok

What about consistency? Must made trade-off

- Stateless applications does not have consistency issue
- Some applications, like Google search, can even tolerate occasional inconsistency
 - Can you notice the inconsistency of search results?

The CAP theorem

Consistency, Availability & Partition tolerance

CAP: an example

Amazon has two zones: US and Euro

- All US users connect to the US zone
- All Euro users connect to the Euro zone

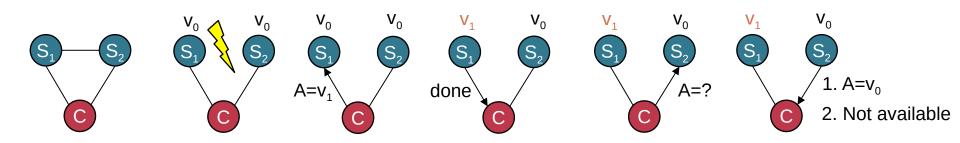
Define Availability and Consistency

- C: all users see the actual number of items
- Not-C: users see 1 item left, but actually is 0
- A: all users can buy items at any time (if there are some)
- Not-A: users might get "cannot buy now, please wait and retry"

The CAP theorem: 2 out of 3

It is impossible for a distributed computer system to simultaneously provide all three of the following guarantees

- Consistency (all nodes see the same data at the same time)
- Availability (a guarantee that every request receives a response about whether it succeeded or failed)
- Partition tolerance (the system continues to operate despite arbitrary message loss or failure of part of the system)



Partition Tolerance

"P" is usually a reality

You can never count on the network connectivity

AP: sacrifice C

- If you have one book but sell it to two customers
- Maybe just an apology and a small gift coupon

CP: sacrifice A

- If you are selling train ticket but cannot deliver
- Customer may sue you
- But the user can fail to buy the ticket, e.g., 12306 during spring festival in the last few years





CAP: not a Binary Decision

Example: a CP system does not mean no "A" at all during network partition

- If the Euro zone and US zone are disconnected, the US zone can still keep available
- Only the Euro zone is not available
- When the network connects again, apply US zone data to Euro zone for consistency

Summary of the (ideal) properties of distributed systems

Can our system handle a larger workload?

Scalability
How easy the developer
can use our system?

Ease of programming

How fast is our system?

Performance

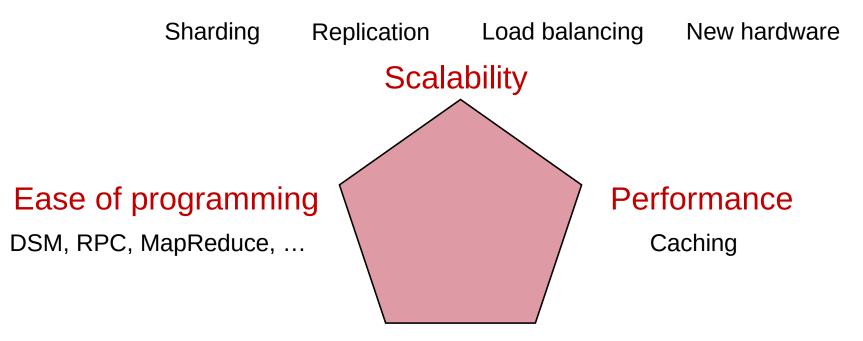
Consistency

Is our data/computation correct?

Fault tolerance

What if some component crash?

Summary of the (ideal) properties of distributed systems



Consistency

Sequential, eventual, consensus, isolation, etc.

Fault tolerance

High availability: replication & checkpointing

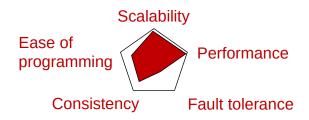
Durability: atomicity, logging, etc.

Summary of the (ideal) properties of distributed systems

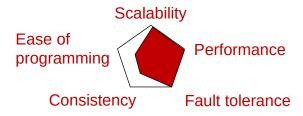
These properties typically cannot achieve at the same time

The adults want them all, but the reality forces them to make trade-offs

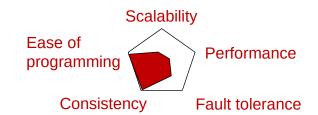
Remote Procedure Call (RPC)



NoSQL databases



Distributed Shared Memory (DSM)



NewSQL databases (e.g., Spanner)

