



UNIVERSITÀ DI PISA

Distributed Systems and Middleware Technologies

## **Distributed File System with Pastry**

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# 1 Introduction

This project consists in the implementation of a peer-to-peer distributed file storage using a distributed hash table for locating the files among the nodes. The distributed hash table is based on Pastry protocol.

## 2 Pastry

Pastry implements a logical ring of nodes, and each node has an ID. The files are hashed to get a key, and the key is used to identify which node is responsible for that file, checking which one has the longest common prefix. When a file is accessed or stored, a request is sent to the current node. When a node receives a request, it uses a routing table, structured as a tree, for finding the next hop for the request. Every time a row of the routing table is accessed,  $b$  bits of the prefix of the key are used to identify the correct column. If the  $b$  bits of the prefix are owned by the current node, the next row is visited, otherwise the address of the next hop is retrieved.

Whenever an address is not found or the whole table is crossed, an extra table is checked, called leaf set and containing the nodes with id close to the current node; if no node in the leaf set has a longer shared prefix with the key than the current node, then the current node is responsible for the file. This routing table, using  $N$  nodes, guarantees a maximum number of hops equal to  $\log n$ .

## 3 Implementation

The implementation of the system is divided in two parts: one for the distributed file system, created in Erlang, and a webserver realized in Java, also thanks to the Frameworks and implementations of the Java Servlets.

### 3.1 Erlang

The Erlang code is divided in multiple modules, in order to isolate functions and keep them short enough.

#### 3.1.1 Util.erl, File\_handler.erl, Network.erl

These files are created for having simple utility functions. *File\_handler.erl* contains the utility functions for dealing with files, *Network.erl* contains simple functions for starting the Erlang nodes and setting the cookies, in order to receive messages from different machines, and *Util.erl* contains generic utility functions.

#### 3.1.2 Key\_gen.erl

*Key\_gen.erl* is the file that implements the key system for Pastry. The keys are generated hashing strings and represented as *bitstrings*; this type is chosen because it can be used for representing streams of any number of bits, differently from the *binary* type, which is exclusively for multiples of 8 bits.

*Bitstrings*, however, have little documentation and almost no function dedicated, so *Key\_gen.erl* also implements the functions for comparing these numbers and handling them as hexadecimal, since Pastry requires groups of 4 bits for the routing.

#### 3.1.3 Routing.erl

*Routing.erl* represents the core of Pastry protocol, implementing the Routing Table. The routing table is represented as a Tuple containing the TableKey, so the key of the node owning the table, and the table. The table is a list of 32 lists, representing the rows; the single entries of the rows are represented by a tuple containing the hexadecimal identifying the column and the node information.

During the routing or any table modification, instead of doing the whole routing, this approach requires of detecting the number of hexadecimal in common in the prefix, which represents the index of the row, and the first hexadecimal not in common, which represent the Column key.

This module contains also the functions for modifying the table and some utility functions relative to it.

#### 3.1.4 Leaf\_set.erl

*Leaf\_set.erl* implements the Leaf Set; the Leaf Set is represented as a Tuple of two lists, *Left* and *Right*, representing the two halves of it. The two list contain at most *L2* elements, represented as tuples key-node information.

The module also offers functions for modifying the leaf set and some utility relative to it.

#### 3.1.5 Node\_actions.erl

*Node\_actions.erl* implements some utility functions used by the upper modules to interact with files, sending messages, or broadcasting messages to lists, routing tables, or leaf sets.

The most important function is *full\_route()*, which implements the full navigation of the routing table, and, in case of no match or it reaches the bottom, the leaf set is analyzed; the function returns the information of the routed node, or *route\_end* if the current node is responsible for the key in input.

#### 3.1.6 Pastry\_actions.erl

*Pastry\_actions.erl* implements all the functions for handling the pastry requests and responses to the other nodes in the network, plus the handling of the self messages for peridoc events. The main functions are:

- *join()*: function used for responding to the *join* message with the routing table row, and for forwarding the message to the next node of the routing; function used by the node to join the network through another known node;
- *join\_res\_handle()*: function used for updating the routing table and leaf set after a *join\_response*;
- *exit()*: function for handling the exit of another node from the network;
- *keepalive()*: sends keepalive to all known nodes;
- *update\_keepalive()*: updates the keepalive of given nodes;
- *keepalive\_res()*: responses to the keepalive to make sure that, even if a node knows another but not viceversa, the one knowing still receives keepalives from the other;
- *check\_expired\_nodes()*: checks which nodes are disconnected and updates the tables;
- *share\_info()*: shares the info about the nodes with the known nodes.

#### 3.1.7 Backup\_actions.erl

*Backup\_actions.erl* implements all the functions for sending and handling the backups. For design choice, the backups are propagated to the whole leaf set and updated according to it. The main functions are:

- *backup()*: function used for propagating a file to the leaf set for backup;
- *backup\_res()*: function used for saving a backup sent by other node;
- *backup\_remove()*: function used to communicate the leaf set to remove a file from backup;
- *new\_leaf\_backup()*: function used to check if a node is added to the new leaf set, and in that case it receives a backup of each file;
- *old\_leaf\_backup()*: function used to check if a node is removed from the new leaf set, and in that case it receives a message to delete the backup folder;

- *update\_leaf\_backup()*: function used for checking in a list of nodes which ones entered the leaf set and which ones exited, in order to propagate or remove backups;
- *backup\_update()*: function used to update the owner of the backups of a failed node. The routing is performed and, if the current node is the owner, the node moves the files to its folder, otherwise it forwards the request;
- *backup\_find()*: request to find the owner of a file. If reaches the end, a *backup\_found* response is sent;
- *backup\_found()*: updates the owner of a file on a *backup\_found*.

### 3.1.8 Web\_responses.erl

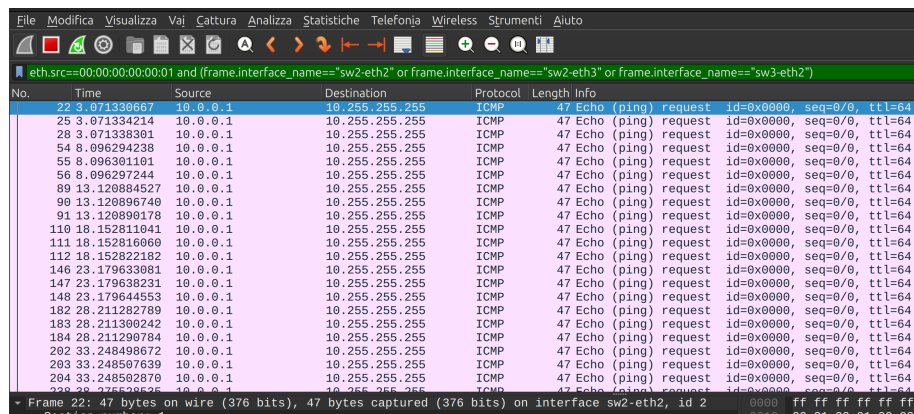
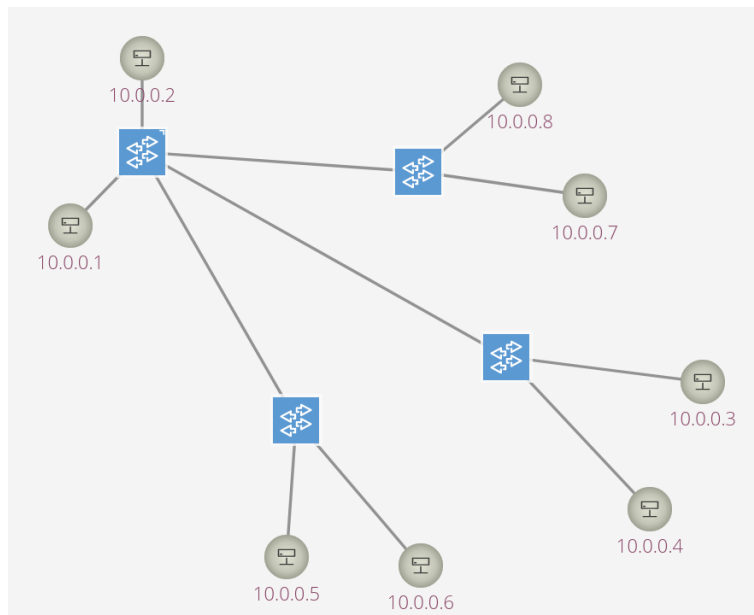
*Web\_responses.erl* implements all the functions for handling and responding to the webserver. The main functions are:

- *find\_store()*: function that, given a filename, finds a storage for it using pastry routing and responds to the web server;
- *store()*: function that, given a filename, it stores it on the node;
- *find()*: function that finds a file and sends it back to the webserver;
- *delete()*: function that, given a filename, finds for it using pastry routing, deletes it, and sends and ack to the web server;
- *get\_file\_res()*: function that starts a flooding for retrieving all the names of the files in the network; it uses a blacklist on the *Msg\_Id* for avoiding infinite flooding;
- *get\_files\_res\_handle()*: function that updates the list of files with the new nodes received during the flooding;
- *all\_files\_res()*: function that responds to the webserver with the list of files;
- *check\_expired\_blacklist()*: function called periodically for cleaning the expired entries of the blacklist.

## 4 Testing and Results

Once Mininet and ONOS are initialized, the network is tested to guarantee that it works as intended by the VPLS. The network is tested both with the *ping* and both with a configured broadcast ethernet frame. The broadcast is executed using the python script *frameBroadcaster.py*.

For checking that the VPLS are working as intended, we used *Wireshark*, checking which interfaces and devices receive the ethernet frames broadcasted. Different topologies were initialized as *net.json* and run. The topologies and some of the screenshots of wireshark are shown below, to prove how only the interfaces in the same VPLS are reached by the broadcast.



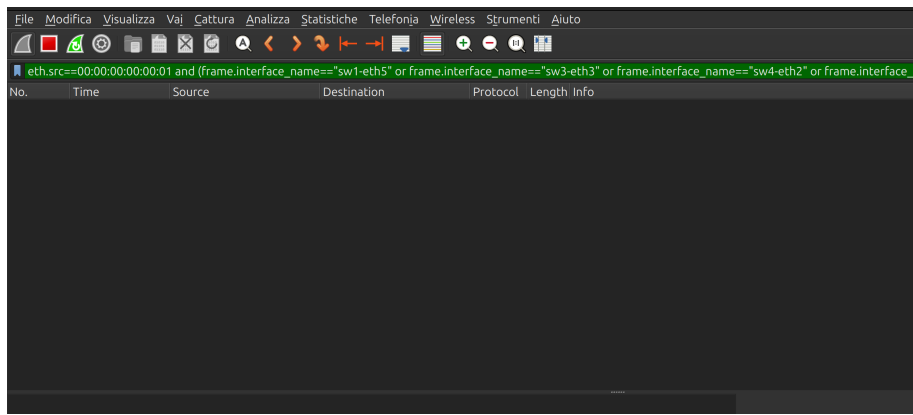


Figure 3: Broadcast example 1: all other interfaces not reached

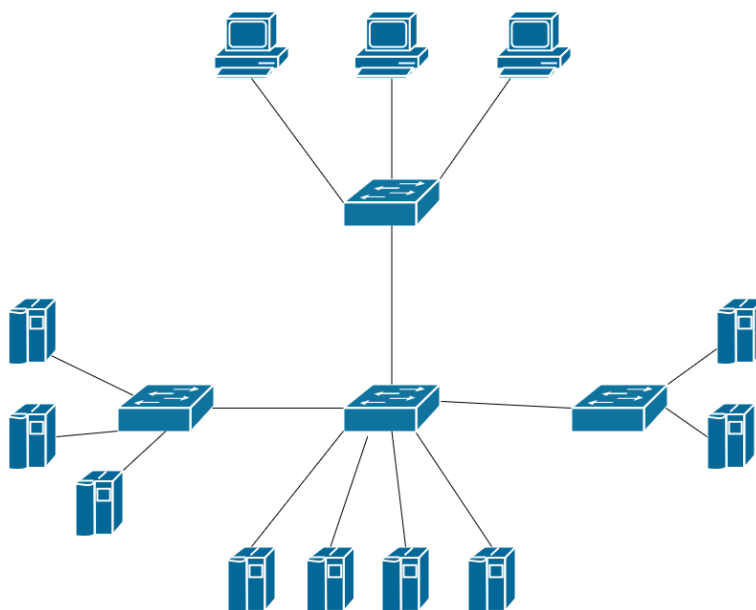


Figure 4: Topology 2

*12 interfaces									
File Modifica Visualizza Vaj Cattura Analizza Statistiche Telefonia Wireless Strumenti Aiuto									
eth.src==00:00:00:00:00:02 and (frame.interface_name=="sw2-eth4" or frame.interface_name=="sw4-eth2" or frame.interface_name=="sw4-eth3")									
No.	Time	Source	Destination	Protocol	Length	Info			
1	0.000000000	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
2	-0.000003096	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
4	-0.000006963	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
54	5.025707506	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
55	5.025715000	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
56	5.025711654	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
81	10.049673299	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
82	10.049678138	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
84	10.049667017	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
134	15.091721768	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
135	15.091731136	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
136	15.091727529	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
188	20.125684666	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
189	20.125690912	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
190	20.125694989	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
222	25.150915487	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
223	25.150909446	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
225	25.150908916	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
284	30.190250855	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
287	30.190269430	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
288	30.190280521	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		
319	35.226042214	10.0.0.2	10.255.255.255	ICMP	51	Echo (ping) request	id=0x0000, seq=0/0, ttl=64		

Figure 5: Broadcast example 2: servers reached

\*12 Interfaces

FileModificaVisualizzaVajCatturaAnalizzaStatisticheTelefoniaWirelessStrumentiAiuto

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Figure 6: Broadcast example 2: all other interfaces not reached

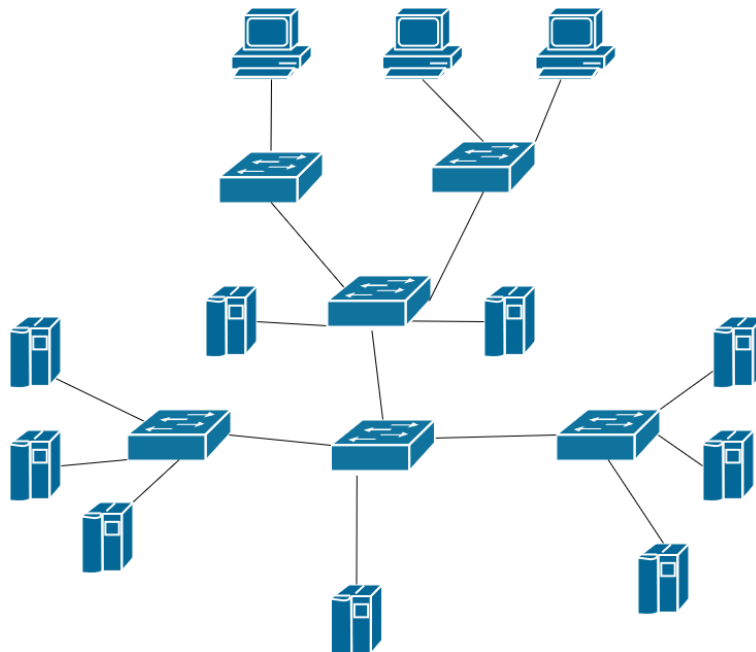


Figure 7: Topology 3



No.	Time	Source	Destination	Protocol	Length	Info
64	4.522976433	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
67	4.522974751	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
70	4.522980352	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
103	9.557274631	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
104	9.557278019	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
105	9.557282678	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
156	14.579213250	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
157	14.579219762	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
158	14.579227537	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
184	19.606206576	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
185	19.606211977	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
186	19.606218429	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
236	24.639032979	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
237	24.639038459	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
238	24.639045061	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64
288	29.676294408	10.0.0.3	10.255.255.255	ICMP	51	Echo (ping) request id=0x0000, seq=0/0, ttl=64

Figure 8: Broadcast example 3: servers reached

No.	Time	Source	Destination	Protocol	Length	Info
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Figure 9: Broadcast example 3: all other interfaces not reached