

UNIVERSITÀ DI PISA

Distributed Systems and Middleware Technologies

Distributed File System with Pastry

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Contents

1 Introduction

This project consists in the implementation of a peer-to-peer distributed file storage using a distributed hash table for locating the files among the nodes. The distributed hash table is based on Pastry protocol.

2 Pastry

Pastry implements a logical ring of nodes, and each node has an ID. The files are hashed to get a key, and the key is used to identify which node is responsible for that file, checking which one has the longest common prefix. When a file is accessed or stored, a request is sent to the current node. When a node receives a request, it uses a routing table, structured as a tree, for finding the next hop for the request. Every time a row of the routing table is accessed, b bits of the prefix of the key are used to identify the correct column. If the b bits of the prefix are owned by the current node, the next row is visited, otherwise the address of the next hop is retrieved.

Whenever an address is not found or the whole table is crossed, an extra table is checked, called leaf set and containing the nodes with id close to the current node; if no node in the leaf set has a longer shared prefix with the key than the current node, then the current node is responsible for the file. This routing table, using N nodes, guarantees a maximum number of hops equal to logn.

3 Implementation

The implementation of the system is divided in two parts: one for the distributed file system, created in Erlang, and a webserver realized in Java, also thanks to the Frameworks and implementions of the Java Servlets.

3.1 Erlang

The Erlang code is divided in multiple modules, in order to isolate functions and keep them short enough.

3.1.1 Util.erl, File_handler.erl, Network.erl

This files are created for having simple utility functions. File_handler.erl contains the utility functions for dealing with files, Network.erl contains simple functions for starting the Erlang nodes and settign the cookies, in order to receive messages from different machines, and Util.erl contains generic utility functions.

3.1.2 Key_gen.erl

Key_gen.erl is the file that implements the key system for Pastry. The keys are generated hashing strings and represented as bitstrings; this type is chosen because it can be used for representing streams of any number of bits, differently from the binary type, which is exclusively for multiples of 8 bits.

Bitstrings, however, have little documentation and almost no function dedicated, so Key_gen.erl also implements the functions for comparing these numbers and handling them as hexadecimals, since Pastry requires groups of 4 bits for the routing.

3.1.3 Routing.erl

Routing.erl represents the core of Pastry protocol, implementing the Routing Table. The routing table is represented as a Tuple containing the TableKey, so the key of the node owning the table, and the table. The table is a list of 32 lists, representing the rows; the single entries of the rows are represented by a tuple containing the hexadecimal identifying the column and the node information.

During the routing or any table modification, instead of doing the whole routing, this approach requires of detecting the number of hexadecimals in common in the prefix, which represents the index of the row, and the first hexadecimal not in common, which represent the Column key.

This module contains also the functions for modifying the table and some utility functions relative to it.

3.1.4 Leaf_set.erl

 $Leaf_set.erl$ implements the Leaf Set; the Leaf Set is represented as a Tuple of two lists, Left and Right, representing the two halves of it. The two list contain at most L2 elements, represented as tuples key-node information.

The module also offerts functions for modifying the leaf set and some utility relative to it.

3.1.5 Node_actions.erl

Node_actions.erl implements some utility functions used by the upper modules to interact with files, sending messages, or broadcasting messages to lists, routing tables, or leaf sets.

The most important function is $full_route()$, which implements the full navigation of the routing table, and, in case of no match or it reaches the bottom, the leaf set is analyzed; the function returns the information of the routed node, or $route_end$ if the current node is responsible for the key in input.

3.1.6 Pastry_actions.erl

Pastry_actions.erl implements all the functions for handling the pastry requests and responses to the other nodes in the network, plus the handling of the self messages for peridoc events. The main functions are:

- join(): function used for responding to the join message with the routing table row, and for forwarding the message to the next node of the routing; function used by the node to join the network through another known node;
- join_res_handle(): function used for updating the routing table and leaf set after a join_response;
- exit(): function for handling the exit of another node from the network;
- keepalive(): sends keepalive to all known nodes:
- update_keepalive(): updates the keepalive of given nodes;
- *keepalive_res()*: responses to the keepalive to make sure that, even if a node knows another but not viceversa, the one knowing still receives keepalives from the other;
- check_expired_nodes(): checks which nodes are disconnected and updates the tables;
- *share_info()*: shares the info about the nodes with the known nodes;

3.1.7 Backup_actions.erl

 $Backup_actions.erl$ implements all the functions for sending and handling the backups. For design choice, the backups are kept only by the nodes in the leaf_set.

- join(): function used for responding to the join message with the routing table row, and for forwarding the message to the next node of the routing; function used by the node to join the network through another known node;
- join_res_handle(): function used for updating the routing table and leaf set after a join_response;
- exit(): function for handling the exit of another node from the network;
- keepalive(): sends keepalive to all known nodes;
- update_keepalive(): updates the keepalive of given nodes;
- *keepalive_res()*: responses to the keepalive to make sure that, even if a node knows another but not viceversa, the one knowing still receives keepalives from the other;

- check_expired_nodes(): checks which nodes are disconnected and updates the tables;
- share_info(): shares the info about the nodes with the known nodes;

4 Testing and Results

Once Mininet and ONOS are initialized, the network is tested to guarantee that it works as intended by the VPLS. The network is tested both with the *ping* and both with a configured broadcast ethernet frame. The broadcast is executed using the python script *frameBroadcaster.py*.

For checking that the VPLS are working as intended, we used *Wireshark*, checking which interfaces and devices receive the ethernet frames broadcasted. Different topologies were initialized as *net.json* and run. The topologies and some of the screenshots of wireshark are shown below, to prove how only the interfaces in the same VPLS are reached by the broadcast.

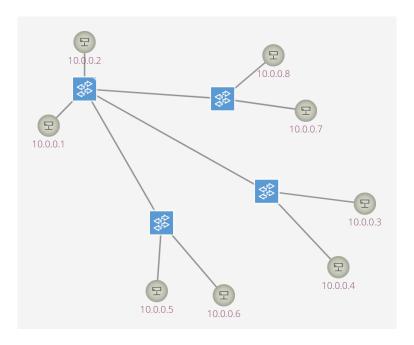


Figure 1: Topology 1

File	Modifica Visualizza	Va <u>i C</u> attura <u>A</u> naliza	za <u>S</u> tatistiche Telefon <u>i</u> a <u>W</u> irel	ess S <u>t</u> rumenti <u>A</u> iuto	0		
			> 🍫 🗠 → 📜 🔳 (Ð 👄 📵 🎹			
■ eth.src==00:00:00:00:00:00:1 and (frame.interface_name=="sw2-eth2" or frame.interface_name=="sw2-eth3" or frame.interface_name=="sw3-eth2")							
No.		Source	Destination	Protocol Length I			
	22 3.071330667	10.0.0.1	10.255.255.255	ICMP 47 E	Echo (ping)		seq=0/0, ttl=64
	25 3.071334214	10.0.0.1	10.255.255.255		Echo (ping)		seq=0/0, ttl=64
	28 3.071338301	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	54 8.096294238	10.0.0.1	10.255.255.255		Echo (ping)		seq=0/0, ttl=64
	55 8.096301101	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	56 8.096297244	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	89 13.120884527	10.0.0.1	10.255.255.255		Echo (ping)		seq=0/0, ttl=64
	90 13.120896740	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	91 13.120890178	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	110 18.152811041	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	111 18.152816060	10.0.0.1	10.255.255.255		Echo (ping)		seq=0/0, ttl=64
	112 18.152822182	10.0.0.1	10.255.255.255		Echo (ping)		seq=0/0, ttl=64
	146 23.179633081	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	147 23.179638231	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	148 23.179644553	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	182 28.211282789	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	183 28.211300242	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	184 28.211290784	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	202 33.248498672	10.0.0.1	10.255.255.255				seq=0/0, ttl=64
	203 33.248507639	10.0.0.1	10.255.255.255		Echo (ping)		seq=0/0, ttl=64
	204 33.248502870	10.0.0.1	10.255.255.255		Echo (ping)		seq=0/0, ttl=64
y Fi			s), 47 bytes captured (376				ff ff ff ff ff
	Section number: 1	NI WITC (OTO DICE	on the syces capeared (ore	bics, on interi	ucc 3M2 CCH2		00 21 00 01 00 00

Figure 2: Broadcast example 1: servers reached

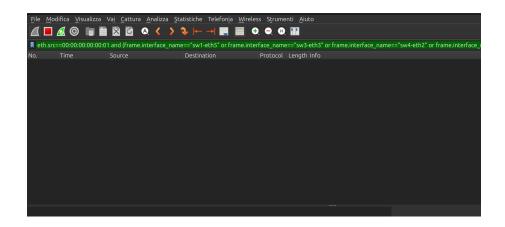


Figure 3: Broadcast example 1: all other interfaces not reached

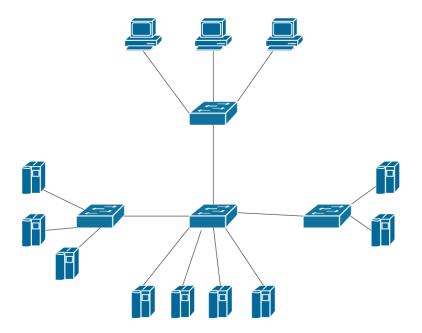


Figure 4: Topology 2

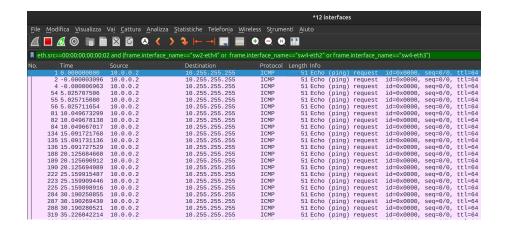


Figure 5: Broadcast example 2: servers reached

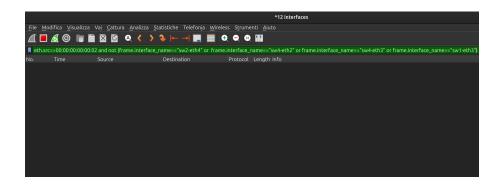


Figure 6: Broadcast example 2: all other interfaces not reached

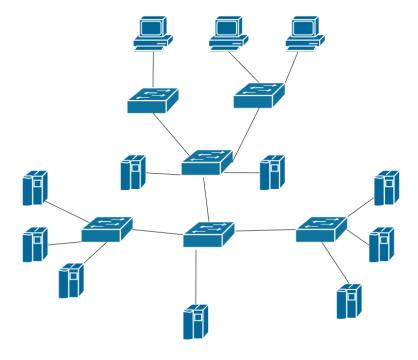


Figure 7: Topology 3

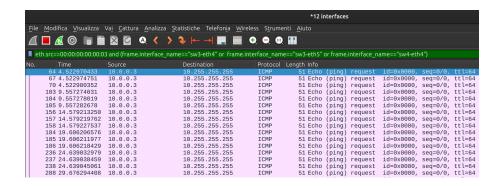


Figure 8: Broadcast example 3: servers reached



Figure 9: Broadcast example 3: all other interfaces not reached