# Database CAP abilities

How to choose the right database for your service?

# Agenda

- From single to a distributed database
- Replication factor
- Consistency level
- Versions and deletes
- The CAP theorem
- The PACELC theorem



RDS



PostgreSQL



Redis



DynamoDB



MySQL



Elasticsearch



MariaDB



S3



SQLite



Timestream



Cassandra



MongoDB

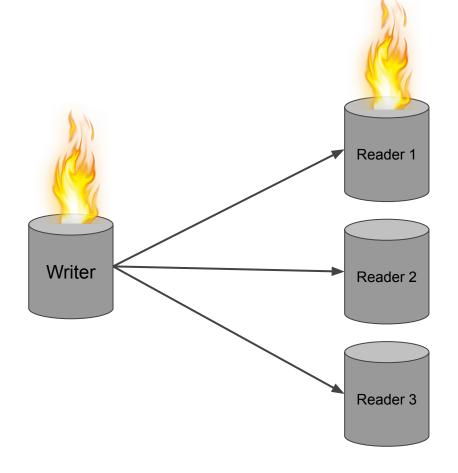
### It all started out quite simple

- Consistent data
- No overhead
- How to scale it?
- How to handle upgrades?
- What if it fails?

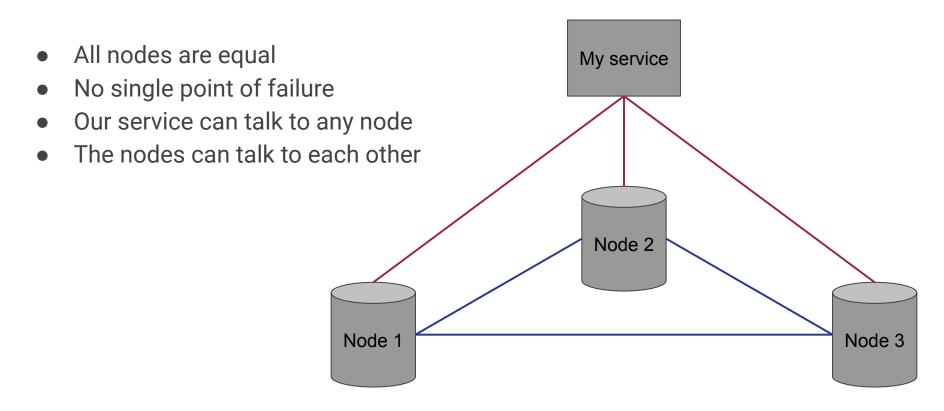


# Replication to the rescue!

- One writer, multiple readers
- More readers → more read capacity
- Can still read when writer is down
- Synchronous vs. asynchronous
- What is our strategy?

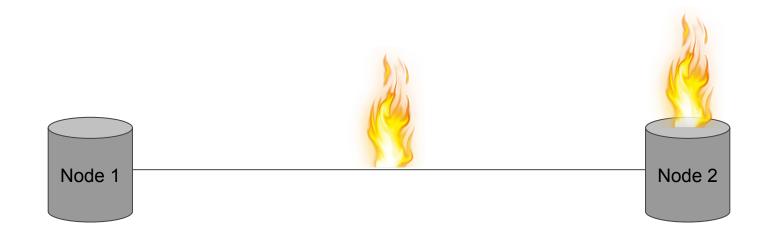


#### Let's build a distributed database!

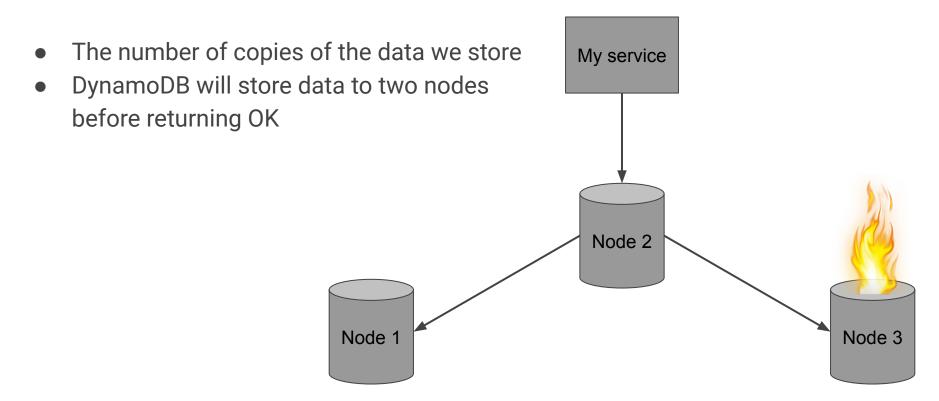


# Why not just two nodes?

- Cost, can't consume more than 50% of the resources
- Updates/failures
- To achieve a quorum, avoids split brain situations

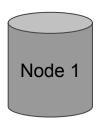


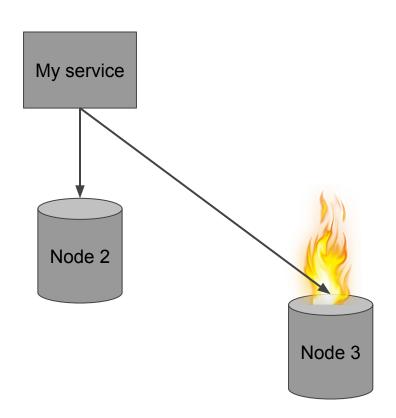
### Replication factor



# Consistency level

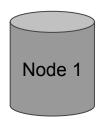
- Eventually consistency
  - Randomly pick a node and read the data from it
- Strong consistency
  - Read from at least two nodes
  - Twice as expensive as eventually consistent reads for DynamoDB

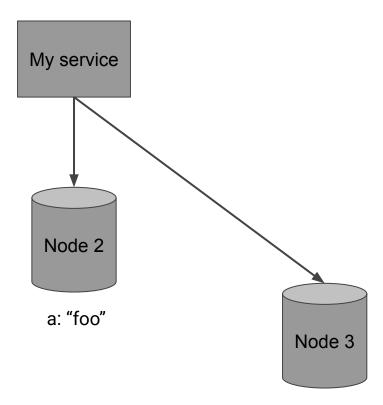




#### Which version is correct?

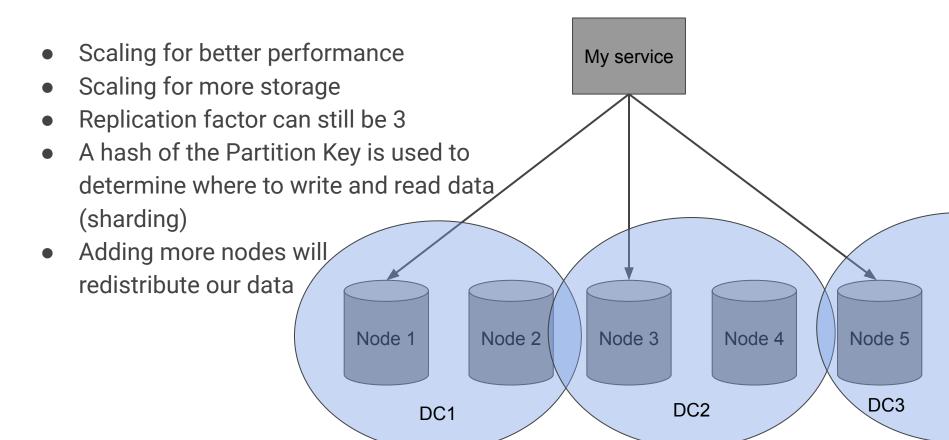
- Timestamps are stored as metadata. The most up to date version is the correct one.
- Eventually consistent reads can return stale (or no) data.
- How to handle deletes? Tombstones.





a: "bar"

# Scaling our distributed database



#### The CAP theorem

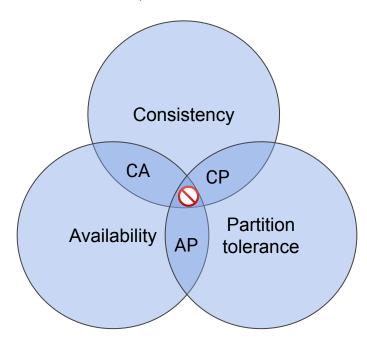
"Any distributed data store can only provide two of the three guarantees:

- Consistency
  - Every read receives the most recent write or an error.
- Availability
  - Every request receives a (non-error) response, without the guarantee that it contains the most recent write.
- Partition tolerance
  - The system continues to operate despite an arbitrary number of messages being dropped (or delayed) by the network between nodes."

## The CAP theorem (by Eric Brewer in 1998)

"When a network partition failure happens, it must be decided whether to do one of the following:

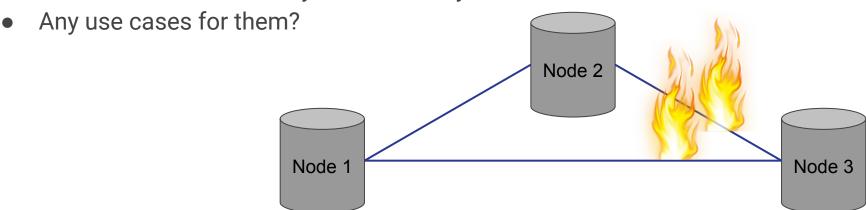
- cancel the operation and thus decrease the availability but ensure consistency
- proceed with the operation and thus provide availability but risk inconsistency."



https://en.wikipedia.org/wiki/CAP\_theorem

### Network partition example

- CA (consistency and availability)
- CP (consistency and partition tolerance)
- AP (availability and partition tolerance)
- The choice is consistency vs. availability



### CAP theorem for popular databases

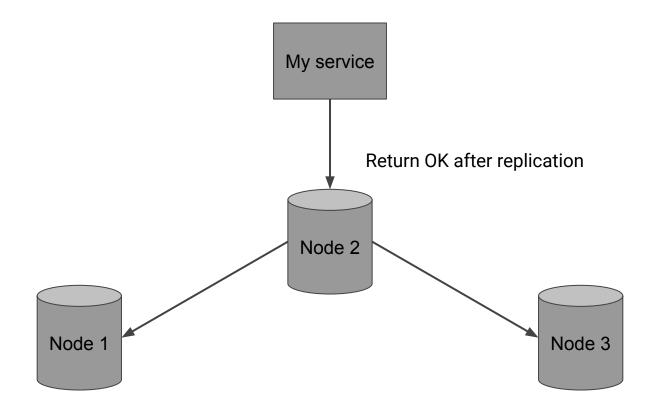
- Availability, Consistency
  - MySQL, PostgreSQL, SQLite
- Consistency, Partition tolerance
  - o MongoDB, Redis
- Availability, Partition tolerance
  - DynamoDB, Cassandra

#### **PACELC**

- There is another trade-off for distributed data stores.
- "In case of network partitioning (P) in a distributed computer system, one has
  to choose between availability (A) and consistency (C) (as per the CAP
  theorem), but else (E), even when the system is running normally in the
  absence of partitions, one has to choose between latency (L) and consistency
  (C)."

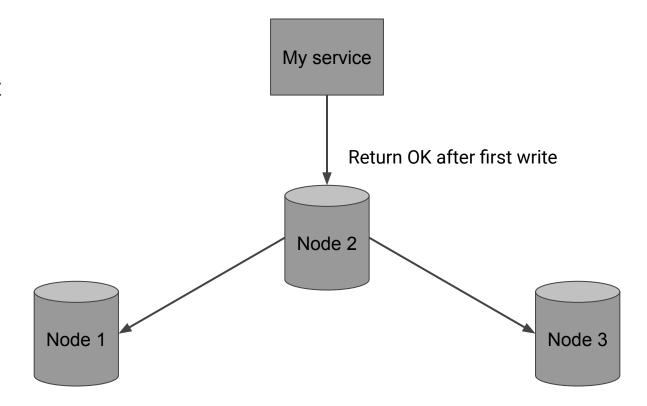
### Consistency

- Strong consistency
- Longer latency



### Latency

- Low latency
- Eventually consistent

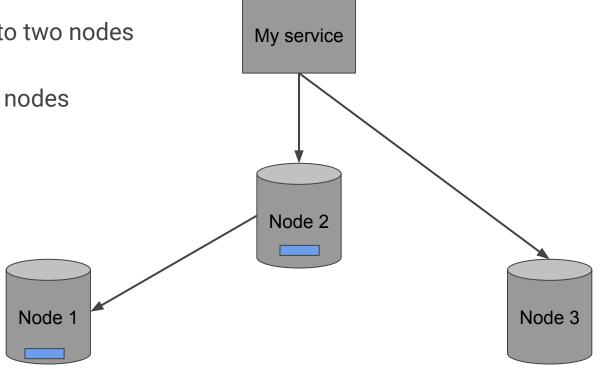


#### Best of both worlds?

 Return OK after writing to two nodes (quorum)
 Read the data from two nodes

 Read the data from two nodes (quorum)

- Medium latency
- Strong consistency
- DynamoDB example



### Summary

- From single to a distributed database
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- The CAP theorem
- The PACELC theorem