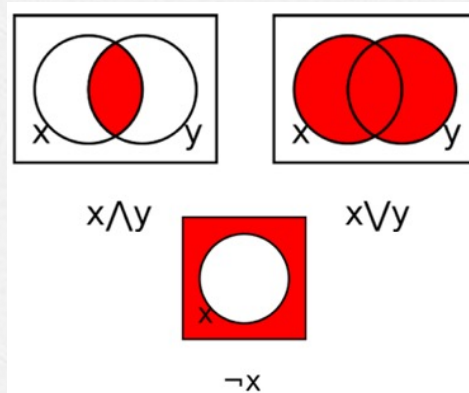


COMP1003 Computer Organization

Lecture 3 Boolean Algebra: From Bits to Logic



United International College

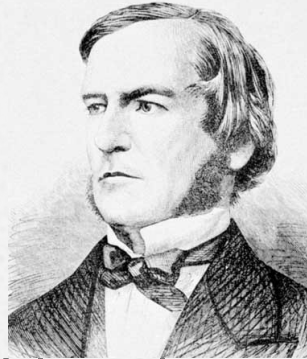
From Bits to Boolean Logic

- Computers represent information by bit (Binary digit)
- A bit has two possible values, namely 0 and 1.
- A bit can be used to represent a truth value, true and false.
- Therefore bits operations correspond to the logical operations in Boolean Algebra.

Algebra

- Algebra is the study of **mathematical symbols** and the **rules** for manipulating these symbols.
- The word comes from the Arabic language from the title of the book *Ilm al-jabr wa'l-muqābala* by *al-Khwarizmi*.
- The original meaning of **algebra** (**al-jabr**) is “reunion of broken parts”

George Boole (1815-1864)



- o Logic was not considered as part of mathematics
- o In 1854 Boole published **An investigation into the Laws of Thought**, on Which are founded the Mathematical Theories of Logic and Probabilities.
- o He pointed out the analogy between algebraic symbols and those that represent logical forms.
- o It began the algebra of logic called **Boolean algebra**

Boolean Algebra

- o **Algebra** is a branch of mathematics that substitutes letters for numbers.
 - o $x+3=y$
- o **Boolean algebra**, sometimes referred to as the algebra of logic, is **a two-valued system of algebra** that represents logical relationships and operations.

Boolean Variables & Operators

- Boolean variables are variables that can take only two values: true/false; or 1/0

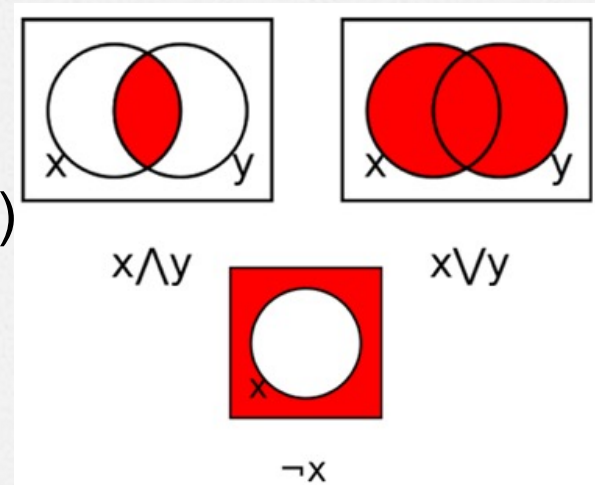
- $A, B, C = \{0, 1\}$

- Boolean Operators

- AND ($A \text{ AND } B$, AB , $A \wedge B$)

- OR ($A \text{ OR } B$, $A+B$, $A \vee B$)

- NOT NOT A , A' , $\neg A$



Venn diagrams

Truth Table of Boolean Operators

AND		
A	B	$A \cdot B$
0	0	0
0	1	0
1	0	0
1	1	1

OR		
A	B	$A + B$
0	0	0
0	1	1
1	0	1
1	1	1

NOT	
A	\bar{A}
0	1
1	0

Boolean Functions

- o A **function** is a relation that uniquely associates members of one set with members of another set
- o A **Boolean function** has
 - o At least one Boolean variable
 - o At least one Boolean operator
 - o At least one input from the set $\{0,1\}$
- o It produces an output that is also a member of the set $\{0,1\}$
- o **$F = X + YZ'$**

Precedence of Boolean Operators

- There might be many Boolean operators in one Boolean function.
- Which operator to apply first?
- The rules of precedence
 - NOT top priority,
 - followed by AND
 - then OR

$$F = X + YZ'$$

Truth Table of a Boolean Function

$$F = X + YZ'$$

X	Y	Z	F
0	0	0	0
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	1
1	1	0	1
1	1	1	1

Exercise

- Write the truth table for

$$F = A' + B'$$

- Write the truth table for

$$Q = A.B.\overline{C} + A.\overline{B}.C + \overline{A}.B.C$$

Simplify a Boolean Function

- o Digital computers contain circuits that implement Boolean functions
- o The simpler that we can make a Boolean function, the smaller the circuit that will result
- o Simpler circuits are cheaper to build, consume less power, and run faster than complex circuits.
- o With this in mind, we always want to reduce our Boolean functions to their simplest form.
- o There are a number of Boolean identities that help us to do this.

Boolean Identities

1) $X \cdot 0 = 0$	10A) $X \cdot Y = Y \cdot X$	} Commutative Law
2) $X \cdot 1 = X$	10B) $X + Y = Y + X$	
3) $X \cdot X = X$	11A) $X(YZ) = (XY)Z$	} Associative Law
	11B) $X + (Y + Z) = (X + Y) + Z$	
4) $X \cdot \bar{X} = 0$	12A) $X(Y + Z) = XY + XZ$	} Distributive Law
	12B) $(X + Y)(W + Z) = XW + XZ + YW + YZ$	
5) $X + 0 = X$	13A) $X + \bar{X}Y = X + Y$	} Consensus Theorem
6) $X + 1 = 1$	13B) $\bar{X} + XY = \bar{X} + Y$	
7) $X + X = X$	13C) $X + \bar{X}\bar{Y} = X + \bar{Y}$	
8) $X + \bar{X} = 1$	13D) $\bar{X} + X\bar{Y} = \bar{X} + \bar{Y}$	
9) $\overline{\bar{X}} = X$	14A) $\overline{X\bar{Y}} = \bar{X} + Y$	} DeMorgan's
	14B) $\overline{X + Y} = \bar{X} \bar{Y}$	

How to Prove These Boolean Identities

- o All of the above identities can be proved using truth tables
- o To do this, you use truth tables to show all of the possible values of both sides of the equation
- o If they are identical, then the identity is true

Proof of $(X + Y)' = X'Y'$

X	Y	X + Y	$(X + Y)'$	X'	Y'	$X'Y'$
0	0	0	1	1	1	1
0	1	1	0	1	0	0
1	0	1	0	0	1	0
1	1	1	0	0	0	0

- the values in the columns under $(X + Y)'$ and $X'Y'$ are identical

Simplifying Boolean Expression

- Which identity is used at each step?

$$\begin{aligned}X'Y' + XYZ + X'Y &= X'Y' + X'Y + XYZ \\&= X'(Y' + Y) + XYZ \\&= X'(1) + XYZ \\&= X' + XYZ \\&= X' + YZ\end{aligned}$$

$$(A+\overline{B}+\overline{C})(A+\overline{B}+C)(A+B+\overline{C})$$

$$A(A+\overline{B}+C)(A+B+\overline{C}) + \overline{B}(A+\overline{B}+C)(A+B+\overline{C}) + \overline{C}(A+\overline{B}+C)(A+B+\overline{C})$$

$$(AA+A\overline{B}+AC)(AA+AB+A\overline{C}) + (\overline{A}\overline{B}+\overline{B}\overline{B}+\overline{B}\overline{C})(\overline{A}\overline{B}+\overline{B}\overline{B}+\overline{B}\overline{C}) + (\overline{A}\overline{C}+\overline{B}\overline{C}+\overline{C}\overline{C})(\overline{A}\overline{C}+\overline{B}\overline{C}+\overline{C}\overline{C})$$

$$AA = A \quad (\text{Idempotent Law})$$

$$\overline{B}\overline{B} = \overline{C}\overline{C} = 0 \quad (\text{Complement Law})$$

$$(A+A\overline{B}+AC)(A+AB+A\overline{C}) + (\overline{A}\overline{B}+\overline{B}\overline{B}+\overline{B}\overline{C})(\overline{A}\overline{B}+0+\overline{B}\overline{C}) + (\overline{A}\overline{C}+\overline{B}\overline{C}+0)(\overline{A}\overline{C}+\overline{B}\overline{C}+\overline{C})$$

$$A + AB = A \quad (\text{Absorption Law})$$

$$A + AC = A \quad (\text{Absorption Law})$$

$$A + 0 = A \quad (\text{Identity Law})$$

$$(A+AC)(A+AB+A\overline{C}) + (\overline{A}\overline{B}+\overline{B}\overline{B}+\overline{B}\overline{C})(\overline{A}\overline{B}+\overline{B}\overline{C}) + (\overline{A}\overline{C}+\overline{B}\overline{C})(\overline{A}\overline{C}+\overline{B}\overline{C}+\overline{C})$$

$$A(A+AB+A\overline{C}) + (\overline{B}+\overline{B}\overline{C})(\overline{A}\overline{B}+\overline{B}\overline{C}) + (\overline{A}\overline{C}+\overline{B}\overline{C})(\overline{B}\overline{C}+\overline{C})$$

$$A(A+A\overline{C}) + \overline{B}(\overline{A}\overline{B}+\overline{B}\overline{C}) + \overline{C}(\overline{A}\overline{C}+\overline{B}\overline{C})$$

$$AA + (\overline{A}\overline{B}+\overline{B}\overline{B}\overline{C}) + (\overline{A}\overline{C}+\overline{B}\overline{C}\overline{C})$$

$$A + (\overline{A}\overline{B}+\overline{B}\overline{C}) + (\overline{A}\overline{C}+\overline{B}\overline{C})$$

$$A + (\overline{A}\overline{B}) + (\overline{B}\overline{C}) + (\overline{A}\overline{C})$$

$$A + (\overline{B}\overline{C}) + (\overline{A}\overline{C})$$

$$A + (\overline{B}\overline{C})$$

Complements

- Sometimes it is more economical to build a circuit using the complement of a function
- DeMorgan's law provides an easy way of finding the complement of a Boolean function

$$\overline{(XY)} = \overline{X} + \overline{Y}, \quad \overline{(X + Y)} = \overline{X} \cdot \overline{Y}$$

- $F = X' + YZ'$, what is F' ?

$$\overline{\overline{X} + Y\overline{Z}} \rightarrow X(\overline{Y} + Z)$$

Canonical Forms

- There are many ways of representing the same Boolean expression
- Logically equivalent expressions** have identical truth tables. For example, $(X+Y)' = X'Y'$
- In order to eliminate confusion, designers express Boolean functions in standardized or canonical form
- There are two **canonical forms**
 - Sum-of-products**
 - Product-of-sums**

Sum-of-Products

- different “product” terms from inputs are “summed” together
- Also called **Disjunctive Normal Form** (DNF)

Inputs			Output	Product
C	B	A	Q	
0	0	0	0	
0	0	1	0	
0	1	0	0	
0	1	1	1	$A.B.\bar{C}$
1	0	0	0	
1	0	1	1	$A.\bar{B}.C$
1	1	0	1	$\bar{A}.B.C$
1	1	1	0	

$$Q = A.B.\bar{C} + A.\bar{B}.C + \bar{A}.B.C$$

Example

○ Write the sum-of-products expression

Inputs			Output
C	B	A	Q
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

Product-of-Sums

- different “sum” terms from inputs are “product-ed” together
- Also called **Conjunctive Normal Form** (CNF)

Inputs			Output	Product
C	B	A	Q	
0	0	0	0	$A + B + C$
0	0	1	1	
0	1	0	0	$A + \bar{B} + C$
0	1	1	1	
1	0	0	1	
1	0	1	1	
1	1	0	0	$A + \bar{B} + \bar{C}$
1	1	1	1	

$$Q = (A + B + C)(A + \bar{B} + C)(A + \bar{B} + \bar{C})$$

Example

Write the product-of-sums expression

Inputs			Output
C	B	A	Q
0	0	0	0
0	0	1	0
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	1
1	1	0	1
1	1	1	0

Binary Addition

- o Binary arithmetic operations are different from Boolean operations
- o How to express binary addition in Boolean algebra?
 1. Draw the truth table
 2. Write out the expression in standard form
 3. Simplify it

o $S = ?$

o $C = ?$

A	B	Sum	Carry
0	0	0	0
0	1	1	0
1	0	1	0
1	1	0	1

Why must the Boolean expressions be simplified?

- o Both SOP and POS may not be the simplest form; only guarantee a standard form
- o The expression must eventually be converted to its simplest form
 - o A one-to-one correspondence exists between a Boolean expression and its implementation using electrical circuits
 - o Unnecessary product terms in the expression lead to unnecessary components in the physical circuit, which in turn yield a more expensive circuit.

Karnaugh map

- The Karnaugh map (KM or K-map 卡诺图) is a method of simplifying Boolean algebra expressions introduced by Maurice Karnaugh (1924~) in 1953.



		A, B			
		0, 0	0, 1	1, 1	1, 0
C, D	1, 0	0	0	1	1
	1, 1	0	0	1	1
	0, 1	0	0	0	1
	0, 0	0	1	1	1

$AC + AB + B\bar{C}\bar{D}$

Be careful of the order!

Example

$$\text{Out} = \bar{A}\bar{B}CD + \bar{A}BCD + ABCD + A\bar{B}CD + AB\bar{C}\bar{D} + AB\bar{C}D + ABC\bar{D}$$

		CD			
		00	01	11	10
A B	00			1	
	01			1	
	11	1	1	1	1
	10			1	

$$\text{Out} = AB + CD$$

Example

$$\begin{aligned} \text{Out} = & \overline{A}\overline{B}\overline{C}\overline{D} + \overline{A}\overline{B}C\overline{D} + \overline{A}B\overline{C}\overline{D} + \overline{A}BC\overline{D} + AB\overline{C}\overline{D} \\ & + AB\overline{C}D + ABCD + A\overline{B}\overline{C}\overline{D} + A\overline{B}C\overline{D} \end{aligned}$$

		CD			
		00	01	11	10
A \ B	00	1		1	
	01	1		1	
	11	1	1	1	
	10	1		1	

$$\text{Out} = \overline{C}\overline{D} + CD + AB\overline{C}$$

		CD			
		00	01	11	10
A \ B	00	1		1	
	01	1		1	
	11	1	1	1	
	10	1		1	

$$\text{Out} = \overline{C}\overline{D} + CD + ABD$$

What's Next

- How to **physically build the logic device** to compute Boolean functions?

