

# Data Structures and Algorithms

Lecture 6



## Insertion Sort And Merge Sort

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# Outline

- Motivation
- Insertion Sort
- Merge Sort
- Divide and Conquer

# Motivation

How do you quickly find

- Your name in a name list?
- A book on a shelf?
- A word in a dictionary?

# Sort

them beforehand!

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# Insertion Sort

*"Let the first  $p$  items be sorted."*

# Insertion Sort

- 1) Initially  $p = 1$
- 2) Let the first  $p$  elements be sorted.
- 3) Insert the  $(p+1)$ th element properly in the list so that now  $p+1$  elements are sorted.
- 4) increment  $p$  and go to step (3)

# How is Insertion Done?

3) Insert the  $(p+1)$ th element properly in the list...

- Scan leftwards
- Move every greater element one position to the right
  - Thus making room for the new element
- Stop when
  - a smaller or equal element is found
  - the left boundary is reached
- Move the new element in
- Animation

# Pseudo Code for Insertion Sort

INSERTION-SORT(A)

1. FOR  $p = 1$  TO  $n-1$
2.      $key = A[p]$
3.      $i = p - 1$
4.     WHILE  $i \geq 0$  AND  $A[i] > key$
5.          $A[i+1] = A[i]$
6.          $i = i - 1$
7.      $A[i+1] = key$

# Discussion

- What is the **best case** for insertion sort?
  - Best case running time?
- What is the **worst case** for insertion sort?
  - Worst case running time?
- What is the **average** running time?
  - Assume that all possible inputs are of the same probability.



# Analysis of Insertion Sort

Best-case Running Time	$O(n)$
Worst-case Running Time	$O(n^2)$
Average Running Time	$O(n^2)$

- Insertion Sort's exact running time cannot be predicted in advance
  - The running time largely depends on the input
  - It is considered an  $O(n^2)$  algorithm, based on its average running time.

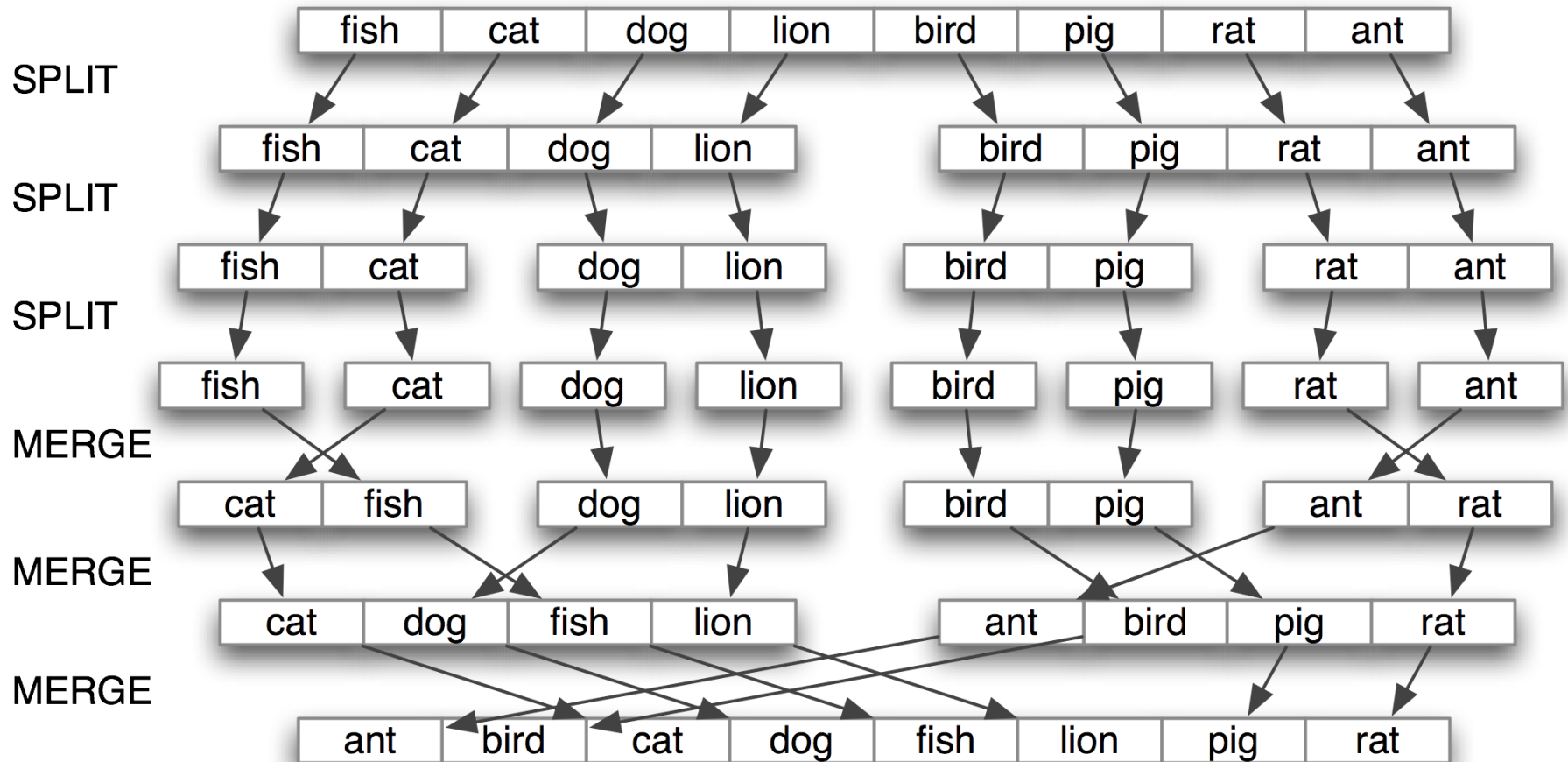
# Merge Sort

A divide-and-conquer (DC)  
algorithm

# Merge Sort

- Divide the list into two smaller lists of about equal sizes
- Sort each smaller list *recursively*
- Merge the two sorted lists to get one sorted list
- [Animation](#)

# Merge Sort Example



# Questions to Ponder

How do we divide the list?

How much time is needed?

How do we merge the two sorted lists?

How much time is needed?

# Dividing

- If the input list is a linked list, dividing takes  $\Theta(N)$  time
  - We scan the linked list, stop at the  $\lfloor N/2 \rfloor$  th entry and cut the link
- If the input list is an array  $A[0..N-1]$ : dividing takes  $O(1)$  time
  1. represent a sublist by two indexes `left` and `right`
  2. to divide `A[left..right]`, we compute `center=(left+right)/2` and obtain `A[left..center]` and `A[center+1..right]`
- `Array` is usually used as the data structure for sorting

# Mergesort

MERGESORT(A, left, right)

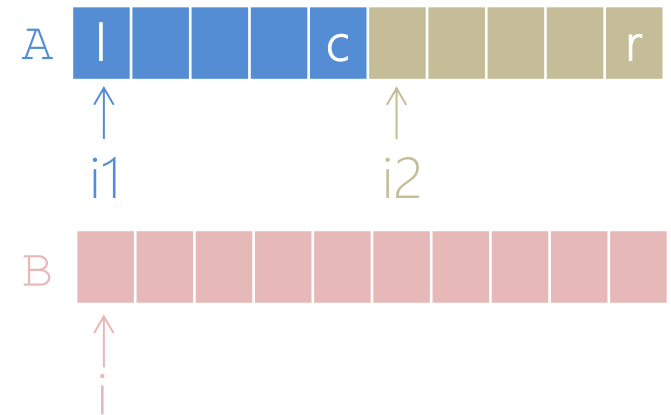
1. IF  $\text{left} \geq \text{right}$
2. RETURN
3.  $\text{center} = (\text{left} + \text{right}) / 2$
4. MERGESORT(A, left, center)
5. MERGESORT(A, center+1, right)
6. MERGE(A, left, center, right)

# Merging

MERGE(A, left, center, right)

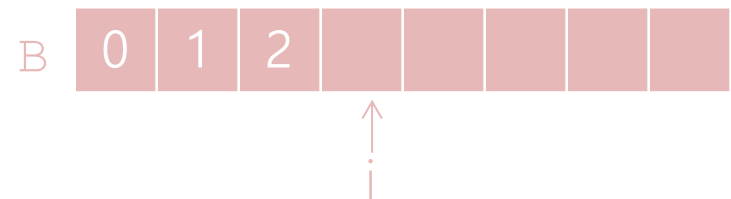
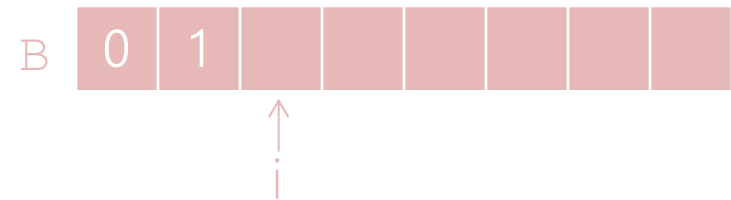
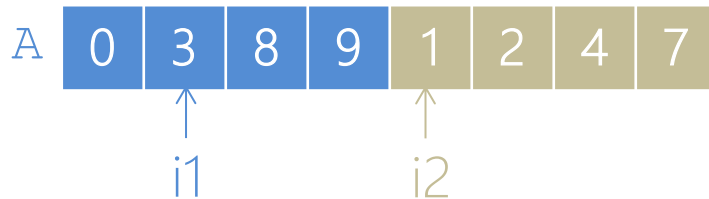
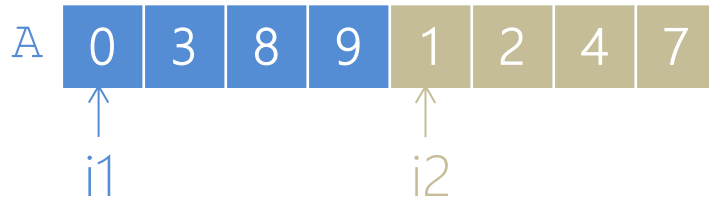
1.  $i1 = \text{left}, i2 = \text{center} + 1, i = 0$
2. WHILE  $i1 \leq \text{center}$  AND  $i2 \leq \text{right}$
3.     IF  $A[i1] < A[i2]$
4.          $B[i++] = A[i1++]$
5.     ELSE
6.          $B[i++] = A[i2++]$
7. WHILE  $i1 \leq \text{center}$
8.      $B[i++] = A[i1++]$
9. WHILE  $i2 \leq \text{right}$
10.      $B[i++] = A[i2++]$
11. Copy B to  $A[\text{left}..\text{right}]$

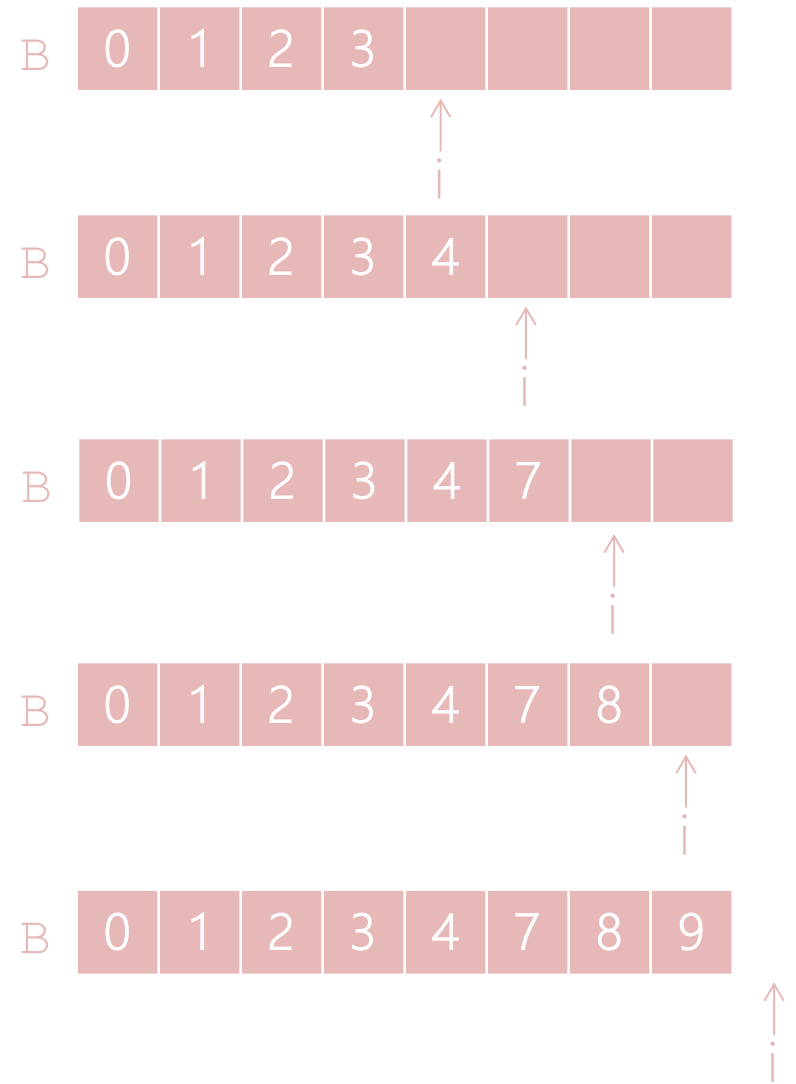
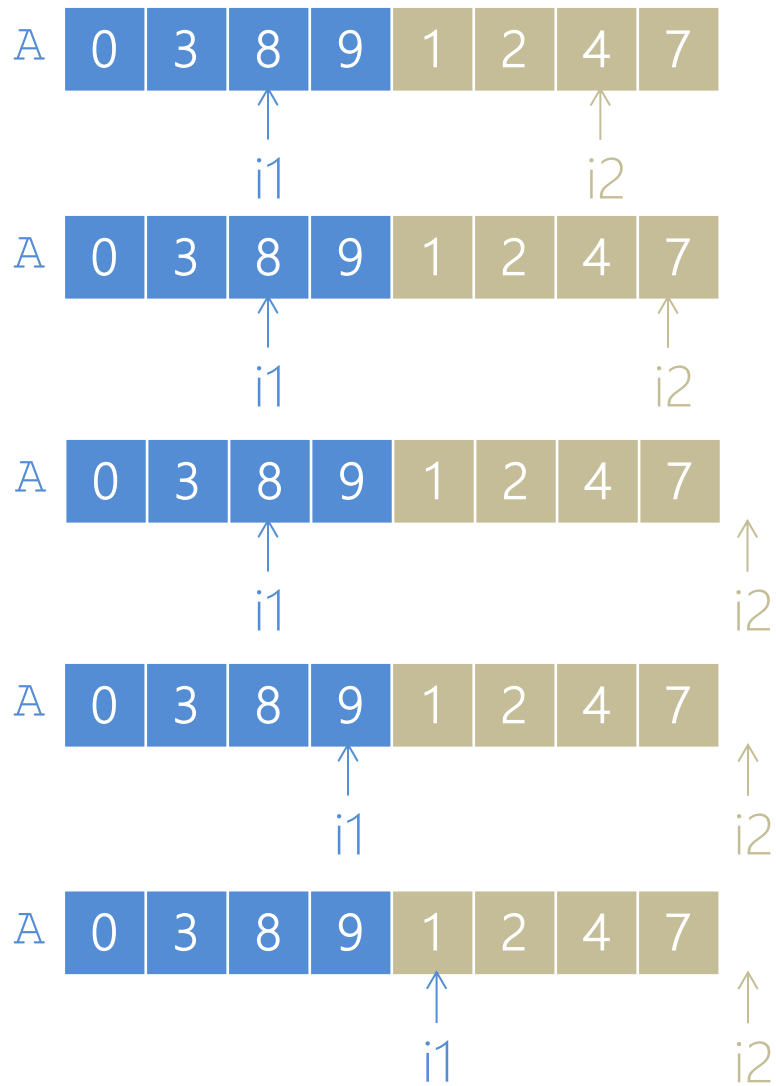
- Merge two sorted sub-arrays  
 $A[\text{left}..\text{center}]$   
and  
 $A[\text{center} + 1, \text{right}]$   
into  $A[\text{left}..\text{right}]$
- Use an extra array, B.





# Merge Example





# Discussion on Merge

- Suppose that `A[left..right]` contains `n` elements
  - What is the `worst-case` running time?
  - What is the `best-case` running time?
  - What is the `extra storage cost`?

# Analysis of Merge Sort

- Let  $T(n)$  denote the worst-case running time of MergeSort where  $n$  is the number of items to be sorted
- Assume that  $n$  is a power of 2.

Divide:  $O(1)$  time

Conquer:  $2T(n/2)$  time

Combine:  $O(n)$  time

- Recurrence equation:

$$T(n) = \begin{cases} 2T(n/2) + O(n), & n > 1 \\ O(1), & n = 1 \end{cases}$$

# Analysis of Merge Sort

Solve the recurrence relation  $T(n) = O(n \log n)$

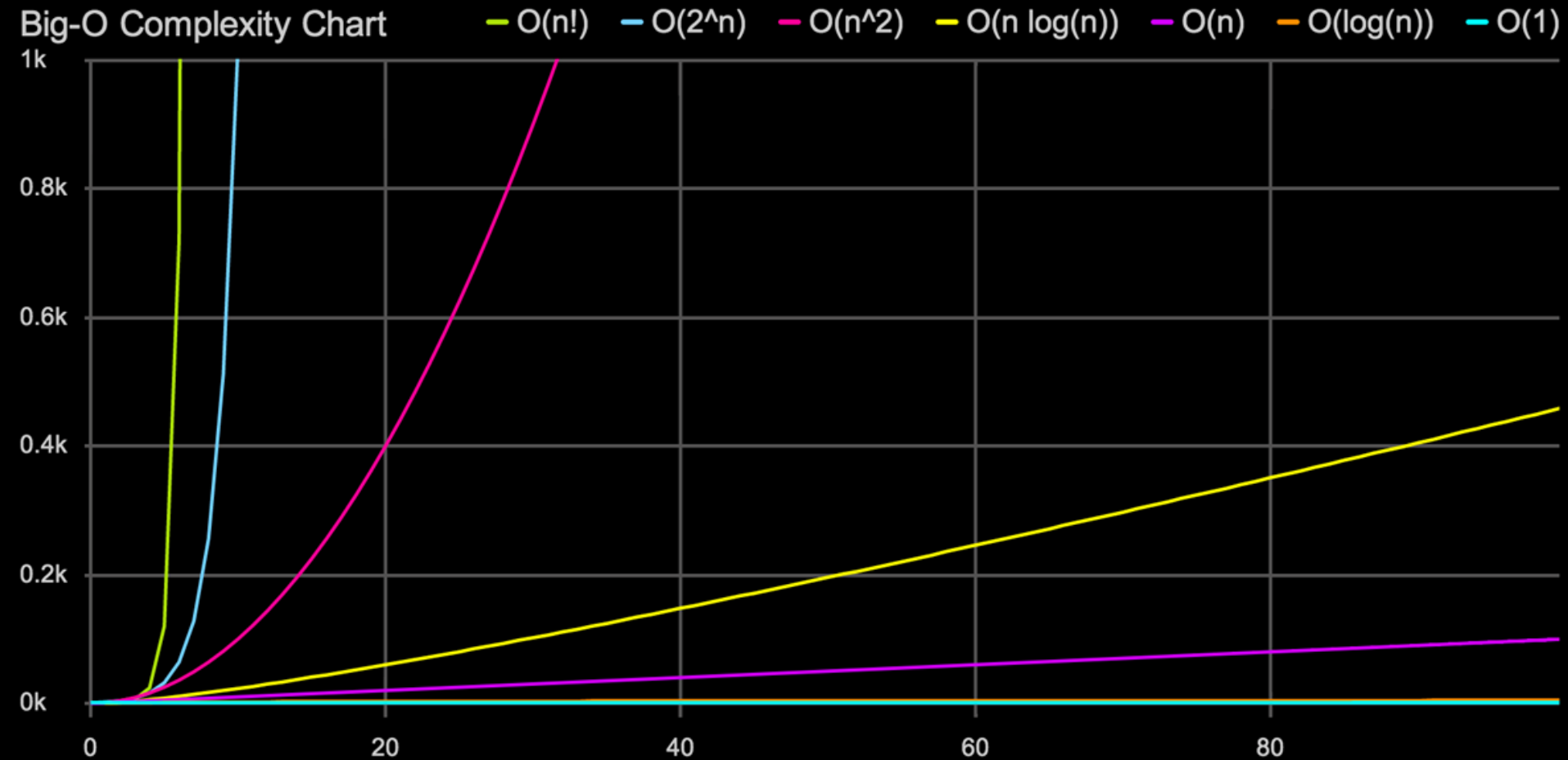
$$\begin{aligned}T(n) &= 2T(n/2) + n \\&= 2[2T(n/2^2) + n/2] + n \\&= 2^2T(n/2^2) + 2n \\&= 2^3T(n/2^3) + 3n \\&= 2^iT(n/2^i) + i*n\end{aligned}$$

Let  $i = \log_2(n)$ :

$$\begin{aligned}T(n) &= nT(n/n) + n*\log(n) \\&= O(n*\log(n))\end{aligned}$$

Note: this is also the best running time so in fact  $T(n) = \Theta(n \log n)$

# $n \cdot \log(n)$ is much faster than $n^2$ !



# Divide and Conquer

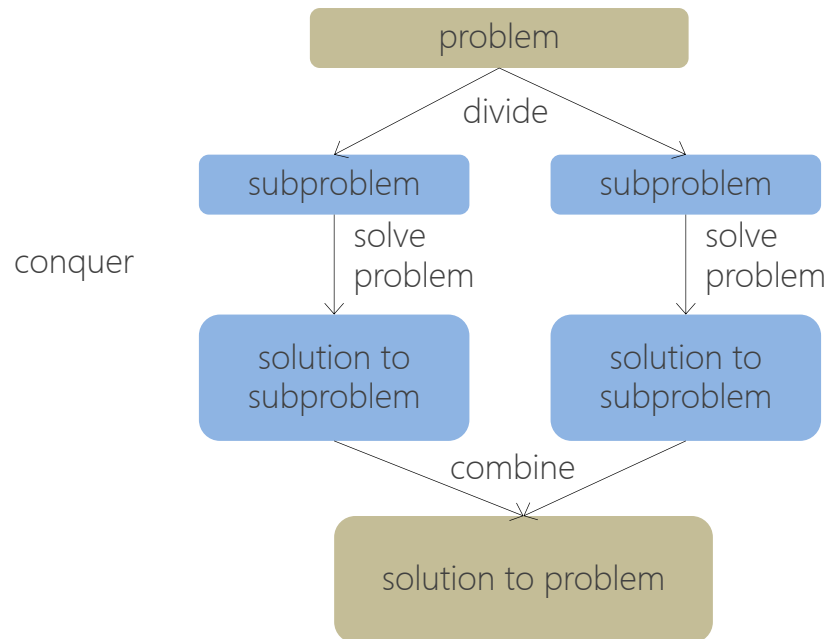
If the problem is large, **break** it into sub-problems that are **smaller in size** but are **similar in structure** to the original problem, **recursively** solve the sub-problems, and finally **combine** the sub-solutions into a final solution that solves the original problem.

# Three Phases of DC

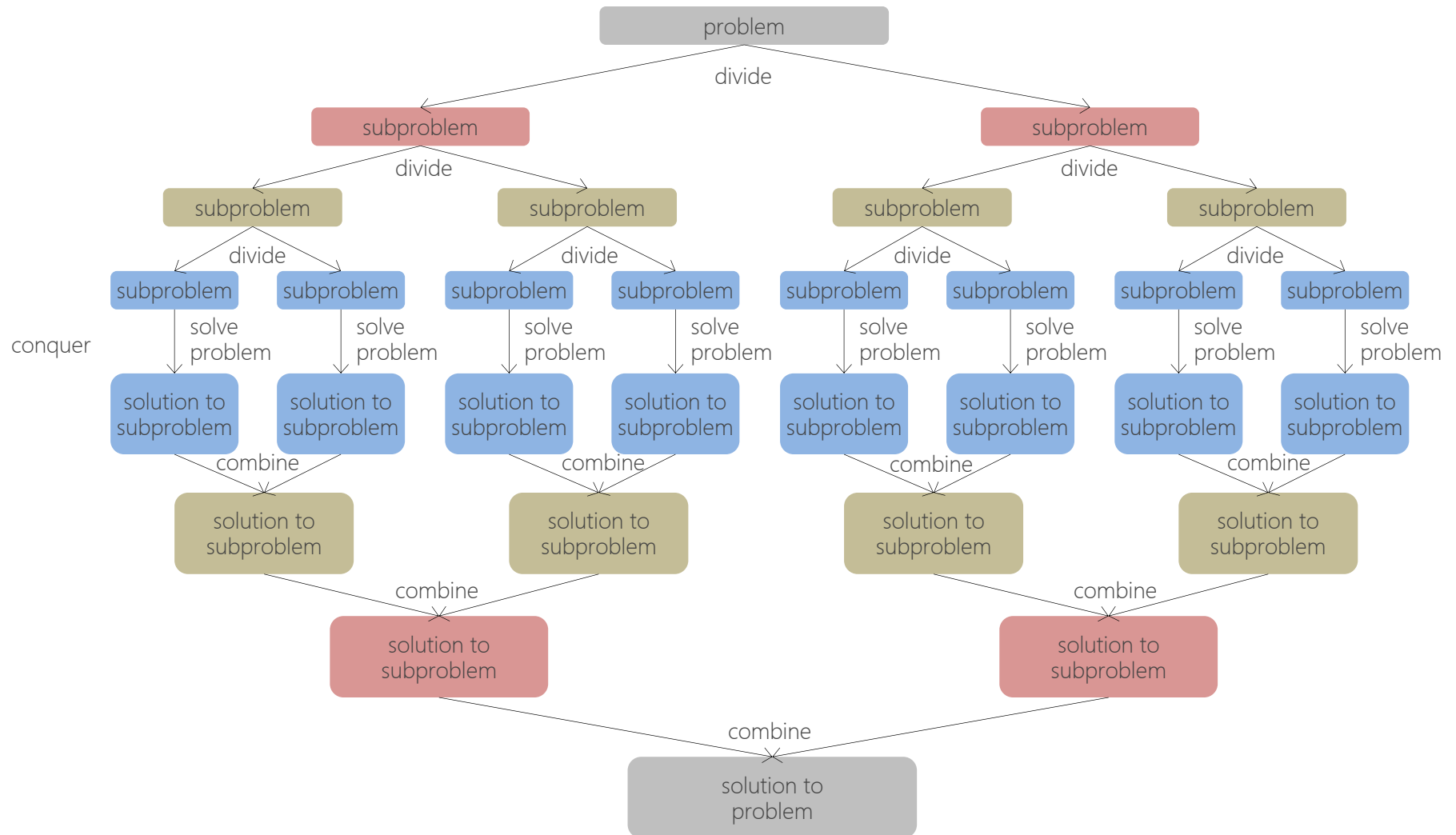
- Divide: **top** → **bottom**
  - Divide a problem into sub-problems
- Conquer: **bottom level**
  - Solve the sub-problems **recursively**
  - If the sub-problems are small enough, solve them as **base cases**
- Combine: **bottom** → **top**
  - Combine the solutions to the sub-problems into that of the original problem
  - Usually the key!



# Divide-Conquer-Combine



# Bigger Divide-Conquer-Combine



# Task

- Create a class, `Sorting2`, which includes at least four static methods
  - public static void `insertionSort(int[] A)`
  - Two overloading `mergeSort` methods
    - public static void `mergeSort(int[] A)`
    - private static void `mergeSort(int[] A, int left, int right)`
    - The first one is public and is for the user to call
    - The second one is private and recursive
    - The second one is called by the first one
  - public static void `main(String[] args)`
- Auxiliary methods may be defined
- Submit `Sorting2.java` to iSpace

# Task

- `public static void insertionSort(int[] A)`
  - $A$  is an array of integers
  - Sort  $A$  using insertion sort
- `public static void mergeSort(int[] A)`
  - It calls the recursive `mergeSort` method to sort  $A$
  - *`mergeSort(A, 0, A.length-1)`*
- `private static void mergeSort(int[] A, int left, int right)`
  - Sort sub-array  $A[\textit{left}..\textit{right}]$  using merge sort
- `public static void main(String[] args)`
  - Generate an array  $A1$  consisting of  $10^5$  random integers which are in range  $[0, 999]$
  - Generate another array  $A2$  which is identical to  $A1$
  - Sort  $A1$  using `insertionSort` and  $A2$  using `mergeSort`
  - Print the elapsed time in `milliseconds` during which both search functions run, respectively