

Algorithm 2022 spring
PA3 Report
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1. Method

I used a $n \times n$ metric `weights[n][n]` to restore the weights within every two edges, which n is the number of vertices. And I also used another metric `PI[n][2]` to restore the states including a vertex's predecessor and whether it's selected. I used the same data structure for both undirected and directed graphs. Besides, I used some constants to define each situation. `LIMIT` is the minimum value of weight `-100`, `BOUNDARY` is a value less than `LIMIT` which means no edge between the two edges, `OUB`(out of boundary) is a value less than `BOUNDARY`, and `INFINITY` is an extremely large number. Then I run Prim's algorithm to reweight the graphs for n times and find out the maximum spanning tree. For directed graphs, I used BFS everytime to check whether there was a cycle when I added an edge to the tree. If it formed a cycle, the edge was what I was going to remove. Finally, I summed up the weights to get the total removed weight and returned.

2. Conclusion

- a. Using arrays as the main data structure could be a little faster than using vectors.
- b. For undirected graphs, we just find the maximum spanning tree then we can get the removed weights.
- c. For directed graphs, they are NP-complete problems.
- d. The Algorithm is really too hard to fully understand in one semester. I still need more practice.