\overline{C}	ont ont a			4.1.14 Farey Sequence	15
C	ontents		4 2	Combinatorics	
1	Misc 1			4.2.1 Matroid Intersection	
	1.1 Contest			4.2.2 De Brujin Sequence	
	1.1.1 Makefile			4.2.3 Multinomial	
	1.2 How Did We Get Here?		4.3	Theorems	
	1.2.1 Macros			4.3.1 Kirchhoff's Theorem	
	1.2.2 Fast I/O			4.3.2 Tutte's Matrix	15
	1.2.3 Bump Allocator			4.3.3 Cayley's Formula	15
	1.3 Tools			4.3.4 Erdős-Gallai Theorem	15
	1.3.1 Floating Point Binary Search 2			4.3.5 Burnside's Lemma	15
	1.3.2 SplitMix64			4.3.6 Gram-Schmidt Process	15
	1.3.3 <random> 2</random>				
	1.3.4 x86 Stack Hack	5	Nui	meric	16
	1.3.5 ctypes		5.1	Barrett Reduction	
	1.4 Algorithms		5.2	Long Long Multiplication	
	1.4.1 Bit Hacks		5.3	Fast Fourier Transform	
	1.4.2 Aliens Trick		5.4	Fast Walsh-Hadamard Transform	
	1.4.3 Hilbert Curve		5.5	Subset Convolution	
	1.4.4 Longest Increasing Subsequence 3		5.6	Linear Recurrences	
	1.4.5 Mo's Algorithm on Tree $\dots \dots 3$			5.6.1 Berlekamp-Massey Algorithm	
				5.6.2 Linear Recurrence Calculation	
2	Data Structures 3		5.7	Matrices	
	2.1 GNU PBDS			5.7.2 Inverse	
	2.2 Segment Tree (ZKW)			5.7.3 Characteristic Polynomial	
	2.3 Line Container			5.7.4 Solve Linear Equation	
	2.4 Li-Chao Tree		F 0	-	
	2.5 adamant HLD			Polynomial Interpolation	
	2.6 van Emde Boas Tree 4 2.7 Wavelet Matrix 5		5.9	Simplex Algorithm	18
	2.8 Link-Cut Tree		~		
	2.0 Link-Out free	6		ometry	19
3	Graph 6		6.1	Point	
	3.1 Modeling			6.1.1 Quarternion	
	3.2 Matching/Flows 6			6.1.2 Spherical Coordinates	
	3.2.1 Dinic's Algorithm 6			Segments	
	3.2.2 Minimum Cost Flow		6.3	Convex Hull	
	3.2.3 Gomory-Hu Tree		0.4	6.3.1 3D Hull	
	3.2.4 Global Minimum Cut			Angular Sort	
	3.2.5 Bipartite Minimum Cover		6.5	Convex Hull Tangent	
	3.2.6 Edmonds' Algorithm		6.6	Convex Polygon Minkowski Sum	
	3.2.7 Minimum Weight Matching 8		6.7	Point In Polygon	
	3.2.8 Stable Marriage			6.7.1 Convex Version	
	3.2.9 Kuhn-Munkres algorithm			6.7.2 Offline Multiple Points Version	
			6.8	Closest Pair	22
	3.3 Shortest Path Faster Algorithm		6.9	Minimum Enclosing Circle	22
	3.4 Strongly Connected Components		6.10	Delaunay Triangulation	22
	3.4.1 2-Satisfiability			6.10.1 Quadratic Time Version	23
	3.5 Biconnected Components		6.11	Half Plane Intersection	23
	3.5.1 Articulation Points				
	3.5.2 Bridges	7	\mathbf{Stri}	ings	23
	3.6 Triconnected Components		7.1	Knuth-Morris-Pratt Algorithm	23
	3.7 Centroid Decomposition		7.2	Aho-Corasick Automaton	23
	3.8 Minimum Mean Cycle		7.3	Suffix Array	24
	3.9 Directed MST		7.4	Suffix Tree	24
	3.10 Maximum Clique		7.5	Cocke-Younger-Kasami Algorithm	24
	3.11 Dominator Tree		7.6	Z Value	25
	3.12 Manhattan Distance MST		7.7	Manacher's Algorithm	25
	3.13 Virtual Tree		7.8	Lyndon Factorization	25
1	Math 13		7.9	Palindromic Tree	25
-	4.1 Number Theory				
	4.1.1 Mod Struct	1	7	π•	
	4.1.2 Miller-Rabin	1.	. 1	Misc	
	4.1.3 Linear Sieve	-	-		
	4.1.4 Get Factors	1.	1.	Contest	
	4.1.5 Binary GCD	1	1 1	Makefile	
	4.1.6 Extended GCD	1.	1.1.	TATANGING	
	4.1.7 Chinese Remainder Theorem	1 .	PRECI	IOUS: ./p%	
	4.1.8 Baby-Step Giant-Step			•	
	4.1.9 Pohlig-Hellman Algorithm	3 %	: p%	sit a unlimited CC /d.	
	4.1.10 Pollard's Rho	5 p	ulim %: p%	nit -s unlimited && ./\$<	
	4.1.11 Tonelli-Shanks Algorithm	-		-серр -o \$බ \$< -std=c++17 -Wall -Wextra -Wshadow \	
	4.1.12 Chinese Sieve	7		g -fsanitize=address,undefined	
	4.1.13 Rational Number Binary Search				

1.2. How Did We Get Here?

1.2.1. Macros

1.2.2. Fast I/O

```
1
   struct scanner {
      static constexpr size_t LEN = 32 << 20;</pre>
      char *buf, *buf_ptr, *buf_end;
      scanner()
           : buf(new char[LEN]), buf_ptr(buf + LEN),
buf_end(buf + LEN) {}
       scanner() { delete[] buf; }
      char getc() {
        if (buf_ptr == buf_end) [[unlikely]]
           buf_end = buf + fread_unlocked(buf, 1, LEN, stdin),
           buf_ptr = buf;
11
        return *(buf_ptr++);
13
      char seek(char del) {
        char c;
while ((c = getc()) < del) {}</pre>
15
17
        return c;
      void read(int &t) {
19
        bool neg = false;

char c = seek('-');

if (c == '-') neg = true, t = 0;

else t = c ^ '0';
23
        while ((c = getc()) >= '0') t = t * 10 + (c ^ '0');
25
        if (neg) t = -t;
27
    }:
    struct printer {
29
      static constexpr size_t CPI = 21, LEN = 32 << 20;</pre>
      char *buf, *buf_ptr, *buf_end, *tbuf;
      char *int_buf, *int_buf_end;
31
      printer()
           : buf(new char[LEN]), buf_ptr(buf),
buf_end(buf + LEN), int_buf(new char[CPI + 1]()),
int_buf_end(int_buf + CPI - 1) {}
33
35
      ~printer() {
37
        flush()
        delete[] buf, delete[] int_buf;
39
      void flush() {
41
        fwrite_unlocked(buf, 1, buf_ptr - buf, stdout);
        buf_ptr = buf;
43
      void write_(const char &c) {
45
         *buf_ptr = c;
        if (++buf_ptr == buf_end) [[unlikely]]
           flush();
49
      void write_(const char *s) {
        for (; *s != '\0'; ++s) write_(*s);
51
      void write(int x) {
        if (x < 0) write_('-'), x = -x;
53
        if (x == 0) [[unlikely]]
55
           return write_('0');
        for (tbuf = int_buf_end; x != 0; --tbuf, x /= 10)
  *tbuf = '0' + char(x % 10);
57
        write_(++tbuf);
59
    };
```

Kotlin

```
import java.io.*
import java.util.*

3

@JvmField val cin = System.`in`.bufferedReader()
@JvmField val cout = PrintWriter(System.out, false)
@JvmField var tokenizer: StringTokenizer

fun nextLine() = cin.readLine()!!

fun read(): String {
```

```
while(!tokenizer.hasMoreTokens())
11
       tokenizer = StringTokenizer(nextLine())
     return tokenizer.nextToken()
13
  }
   // example
15
   fun main() {
17
     val n = read().toInt()
     val a = DoubleArray(n) { read().toDouble() }
     cout.println("omg hi")
19
     cout.flush()
21 }
```

1.2.3. Bump Allocator

```
1 // global bump allocator
   char mem[256 << 20]; // 256 MiB</pre>
   size_t rsp = sizeof mem;
3
   void *operator new(size_t s) {
     assert(s < rsp); // MLE
     return (void *)&mem[rsp -= s];
   }
7
   void operator delete(void *) {}
    // bump allocator for STL / pbds containers
   char mem[256 << 20];</pre>
   size_t rsp = sizeof mem;
   template <typename T> struct bump {
     using value_type = T;
     bump() {}
15
     template <typename U> bump(U, ...) {}
     T *allocate(size_t n) {
17
       rsp -= n * sizeof(T);
rsp &= 0 - alignof(T);
19
       return (T *)(mem + rsp);
21
     void deallocate(T *, size_t n) {}
23 };
```

1.3. Tools

1.3.1. Floating Point Binary Search

```
1 union di {
     double d:
3
     ull i:
5
   bool check(double):
   // binary search in [L, R) with relative error 2^-eps
   double binary_search(double L, double R, int eps) {
     di l = {L}, r = {R}, m;
while (r.i - l.i > 1LL << (52 - eps)) {</pre>
9
       m.i = (l.i + r.i) >> 1;
        if (check(m.d)) r = m;
11
       else l = m;
13
     return l.d;
15 }
```

1.3.2. SplitMix64

```
using ull = unsigned long long;
inline ull splitmix64(ull x) {
   // change to `static ull x = SEED;` for DRBG
   ull z = (x += 0x9E3779B97F4A7C15);
   z = (z ^ (z >> 30)) * 0xBF58476D1CE4E5B9;
   z = (z ^ (z >> 27)) * 0x94D049BB133111EB;
   return z ^ (z >> 31);
}
```

1.3.3. <random>

1.3.4. x86 Stack Hack

```
constexpr size_t size = 200 << 20; // 200MiB
int main() {
    register long rsp asm("rsp");
    char *buf = new char[size];
    asm("movq %0, %%rsp\n" ::"r"(buf + size));
    // do stuff
    asm("movq %0, %%rsp\n" ::"r"(rsp));
    delete[] buf;
}</pre>
```

1.3.5. ctypes

```
from ctypes import *

# computes 10**4300
gmp = CDLL('libgmp.so')
x = create_string_buffer(b'\x00'*16)
gmp.__gmpz_init_set_ui(byref(x), 10)
gmp.__gmpz_pow_ui(byref(x), byref(x), 4300)
gmp.__gmp_printf(b'%Zd\n', byref(x))
gmp.__gmpz_clear(byref(x))
# objdump -T `whereis libgmp.so`
```

1.4. Algorithms

1.4.1. Bit Hacks

```
// next permutation of x as a bit sequence
ull next_bits_permutation(ull x) {
  ull c = __builtin_ctzll(x), r = x + (1ULL << c);
  return (r ^ x) >> (c + 2) | r;
}
// iterate over all (proper) subsets of bitset s
void subsets(ull s) {
  for (ull x = s; x;) { --x &= s; /* do stuff */ }
}
```

1.4.2. Aliens Trick

```
// min dp[i] value and its i (smallest one)
pll get_dp(int cost);
ll aliens(int k, int l, int r) {
    while (l != r) {
        int m = (l + r) / 2;
        auto [f, s] = get_dp(m);
        if (s == k) return f - m * k;
        if (s < k) r = m;
        else l = m + 1;
}
return get_dp(l).first - l * k;
}</pre>
```

1.4.3. Hilbert Curve

```
1  ll hilbert(ll n, int x, int y) {
    ll res = 0;
3  for (ll s = n; s /= 2;) {
    int rx = !!(x & s), ry = !!(y & s);
    res += s * s * ((3 * rx) ^ ry);
    if (ry == 0) {
        if (rx == 1) x = s - 1 - x, y = s - 1 - y;
        swap(x, y);
    }
}
return res;
}
```

1.4.4. Longest Increasing Subsequence

```
template <class I> vi lis(const vector<I> &S) {
     if (S.empty()) return {};
     vi prev(sz(S));
     typedef pair<I, int> p;
     vector res;
     rep(i, \theta, sz(\hat{s})) { // change \theta -> i for longest non-decreasing subsequence
       auto it = lower_bound(all(res), p{S[i], 0});
       if (it == res.end())
          res.emplace_back(), it = res.end() - 1;
        *it = {S[i], i};
11
       prev[i] = it == res.begin() ? 0 : (it - 1)->second;
     int L = sz(res), cur = res.back().second;
     vi ans(L);
while (L--) ans[L] = cur, cur = prev[cur];
15
17
     return ans;
```

1.4.5. Mo's Algorithm on Tree

```
void MoAlgoOnTree() {
        Dfs(0, -1);
        vector<int> euler(tk);
        for (int i = 0; i < n; ++i) {
  euler[tin[i]] = i;</pre>
 5
           euler[tout[i]] = i;
        vector<int> l(q), r(q), qr(q), sp(q, -1);
for (int i = 0; i < q; ++i) {
   if (tin[u[i]] > tin[v[i]]) swap(u[i], v[i]);
 9
           int z = GetLCA(u[i], v[i]);
11
           sp[i] = z[i];
if (z == u) l[i] = tin[u[i]], r[i] = tin[v[i]];
else l[i] = tout[u[i]], r[i] = tin[v[i]];
13
15
           qr[i] = i;
        sort(qr.begin(), qr.end(), [δ](int i, int j) {
   if (l[i] / kB == l[j] / kB) return r[i] < r[j];
   return l[i] / kB < l[j] / kB;</pre>
17
19
21
        vector<bool> used(n);
         // Add(v): add/remove v to/from the path based on used[v]
        for (int i = 0, tl = 0, tr = -1; i < q; ++i) {
  while (tl < l[qr[i]]) Add(euler[tl++]);</pre>
23
           while (tl > l[qr[i]]) Add(euler[--tl]);
25
           while (tr > r[qr[i]]) Add(euler[tr--]);
while (tr < r[qr[i]]) Add(euler[++tr]);</pre>
27
            // add/remove LCA(u, v) if necessary
29
     }
```

2. Data Structures

2.1. GNU PBDS

```
#include <ext/pb_ds/assoc_container.hpp>
    #include <ext/pb_ds/priority_queue.hpp>
    #include <ext/pb_ds/tree_policy.hpp>
    using namespace __gnu_pbds;
   // most std::map + order_of_key, find_by_order, split, join
template <typename T, typename U = null_type>
    using ordered_map = tree<T, U, std::less<>, rb_tree_tag,
                                  tree_order_statistics_node_update>;
    // useful tags: rb_tree_tag, splay_tree_tag
11
    template <typename T> struct myhash {
      size_t operator()(T x) const; // splitmix, bswap(x*R), ...
    // most of std::unordered_map, but faster (needs good hash)
   template <typename T, typename U = null_type>
using hash_table = gp_hash_table<T, U, myhash<T>>;
19
   // most std::priority_queue + modify, erase, split, join
using heap = priority_queue<int, std::less<>>;
21 // useful tags: pairing_heap_tag, binary_heap_tag,
                       (rc_)?binomial_heap_tag, thin_heap_tag
    11
```

2.2. Segment Tree (ZKW)

```
1 struct segtree {
       using T = int;
T f(T a, T b) { return a + b; } // any monoid operation
static constexpr T ID = 0; // identity element
 5
        int n;
        vector<T> v;
        segtree(\textbf{int} \ n\_) \ : \ n(n\_), \ v(2 \ * \ n, \ ID) \ \{\}
        segtree(vector<T> δa) : n(a.size()), v(2 * n, ID) {
           copy_n(a.begin(), n, v.begin() + n);
for (int i = n - 1; i > 0; i--)
 9
              v[i] = f(v[i * 2], v[i * 2 + 1]);
11
13
        void update(int i, T x) {
           for (v[i += n] = x; i /= 2;)
v[i] = f(v[i * 2], v[i * 2 + 1]);
15
17
       T querv(int l. int r) {
           T tl = ID, tr = ID;
           for (l += n, r += n; l < r; l /= 2, r /= 2) {
   if (l & 1) tl = f(tl, v[l++]);
   if (r & 1) tr = f(v[--r], tr);</pre>
19
21
           return f(tl, tr);
23
       }
25 };
```

2.3. Line Container

```
mutable ll k, m, p;
        bool operator<(const Line &o) const { return k < o.k; }</pre>
        bool operator<(ll x) const { return p < x; }</pre>
 5
     // add: line y=kx+m, query: maximum y of given x
    struct LineContainer : multiset<Line, less<>> {
   // (for doubles, use inf = 1/.0, div(a,b) = a/b)
   static const ll inf = LLONG_MAX;
   ll div(ll a, ll b) { // floored division
     return a / b - ((a ^ b) < 0 %% a % b);
}</pre>
 9
11
13
       bool isect(iterator x, iterator y) {
          if (y == end()) return x -> p = inf, \theta;
           if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
else x->p = div(y->m - x->m, x->k - y->k);
15
           return x->p >= y->p;
17
19
        void add(ll k, ll m) {
          auto z = insert({k, m, 0}), y = z++, x = y;
while (isect(y, z)) z = erase(z);
21
           if (x != begin() && isect(--x, y))
          isect(x, y = erase(y));
while ((y = x) != begin() && (--x)->p >= y->p)
23
25
             isect(x, erase(y));
        ll query(ll x) {
27
           assert(!empty());
29
           auto l = *lower_bound(x);
           return l.k * x + l.m;
31
```

2.4. Li-Chao Tree

```
constexpr ll MAXN = 2e5, INF = 2e18;
   struct Line {
     ll m, b;
Line(): m(0), b(-INF) {}
     Line(ll _m, ll _b) : m(_m), b(_b) {} ll operator()(ll x) const { return m * x + b; }
7
   };
   struct Li_Chao {
9
     Line a[MAXN * 4];
      void insert(Line seg, int l, int r, int v = 1) {
        if (l == r) {
11
          if (seg(l) > a[v](l)) a[v] = seg;
          return;
15
        int mid = (l + r) >> 1;
        if (a[v].m > seg.m) swap(a[v], seg);
        if (a[v](mid) < seg(mid)) {</pre>
17
          swap(a[v], seg);
          insert(seg, l, mid, v << 1);</pre>
19
        } else insert(seg, mid + 1, r, v << 1 | 1);</pre>
21
     ll query(int x, int l, int r, int v = 1) {
23
        if (l == r) return a[v](x);
        int mid = (l + r) >> 1;
        if (x <= mid)
25
          return max(a[v](x), query(x, l, mid, v << 1));</pre>
27
          return max(a[v](x), query(x, mid + 1, r, v << 1 | 1));
29
   };
```

2.5. adamant HLD

```
// subtree of v is [in[v], out[v])
   // top of heavy path of v is nxt[v]
   void dfs1(int v) {
     sz[v] = 1;
5
     for (int u : child[v]) {
       par[v] = u;
       dfs1(u);
                sz[u];
9
       if (sz[u] > sz[child[v][0]]) { swap(u, child[v][0]); }
   }
11
   void dfs2(int v) {
     in[v] = t++;
for (int u : child[v]) {
13
       nxt[u] = (u == child[v][0] ? nxt[v] : u);
15
       dfs2(u):
17
     out[v] = t;
```

```
19    }
    int lca(int a, int b) {
21        for (;; b = par[nxt[b]]) {
            if (in[b] < in[a]) swap(a, b);
            if (in[nxt[b]] <= in[a]) return a;
        }
25    }</pre>
```

```
2.6. van Emde Boas Tree
1 // stores integers in [0, 2^B)
   // find(.+) finds first >=/<= i (or -1/2^B if none)</pre>
   // space: ~2^B bits, time: 2^B init/clear, log B operation
   template <int B, typename ENABLE = void> struct VEBTree {
     const static int K = B / 2, R = (B + 1) / 2, M = (1 << B); const static int S = 1 << K, MASK = (1 << R) - 1;
     array<VEBTree<R>, S> ch;
     VEBTree<K> act;
9
     int mi, ma;
     bool empty() const { return ma < mi; }</pre>
     int findNext(int i) const {
11
        if (i <= mi) return mi;</pre>
13
        if (i > ma) return M;
        int j = i >> R, x = i & MASK;
int res = ch[j].findNext(x);
15
        if (res <= MASK) return (j << R) + res;</pre>
17
        j = act.findNext(j + 1);
        return (j >= S) ? ma : ((j << R) + ch[j].findNext(0));
19
     int findPrev(int i) const {
21
        if (i >= ma) return ma;
       if (i < mi) return -1;
int j = i >> R, x = i & MASK;
int res = ch[j].findPrev(x);
23
25
        if (res >= 0) return (j << R) + res;
        j = act.findPrev(j - 1);
27
        return (j < 0) ? mi : ((j << R) + ch[j].findPrev(MASK));</pre>
29
     void insert(int i) {
        if (i <= mi) {
          if (i == mi) return;
31
          swap(mi, i);
          if (i == M) ma = mi; // we were empty
33
          if (i >= ma) return; // we had mi == ma
        } else if (i >= ma) {
35
          if (i == ma) return;
37
          swap(ma, i);
          if (i <= mi) return; // we had mi == ma</pre>
39
        int i = i >> R:
        if (ch[j].empty()) act.insert(j);
41
        ch[j].insert(i & MASK);
43
     void erase(int i) {
45
        if (i <= mi) {
          if (i < mi) return;</pre>
47
          i = mi = findNext(mi + 1);
          if (i >= ma) {
            if (i > ma) ma = -1; // we had mi == ma
49
                                   // after erase we have mi == ma
51
        } else if (i >= ma) {
          if (i > ma) return;
53
          i = ma = findPrev(ma - 1);
55
          if (i <= mi) return; // after erase we have mi == ma</pre>
        int j = i \gg R;
57
        ch[j].erase(i & MASK);
59
        if (ch[j].empty()) act.erase(j);
61
     void clear() {
       mi = M, ma = -1;
        act.clear();
63
        for (int i = 0; i < S; ++i) ch[i].clear();</pre>
     template <class T>
     void init(const T &bts, int shift = 0, int s0 = 0,
                 int s1 = 0) {
69
        -shift + bts.findNext(shift + s0, shift + M - 1 - s1);
71
        s1 =
        (-shift + bts.findPrev(shift + M - 1 - s1, shift + s0));
73
        if (s0 + s1 >= M) clear();
75
        else {
          act.clear();
          mi = s0, ma = M - 1 - s1;
77
```

++s0:

```
++s1;
for (int j = 0; j < S; ++j) {
 79
            81
 83
            if (!ch[j].empty()) act.insert(j);
 85
 87
    template <int B> struct VEBTree<B, enable_if_t<(B <= 6)>> {
 89
      const static int M = (1 << B);</pre>
 91
      ull act;
      bool empty() const { return !act; }
void clear() { act = 0; }
 93
      int findNext(int i) const {
        return ((i < M) δδ (act >> i))
 95
               ? i + __builtin_ctzll(act >> i)
97
               : M:
99
      int findPrev(int i) const {
        return ((i != -1) && (act << (63 - i)))
                ? i - __builtin_clzll(act << (63 - i))
101
103
      void insert(int i) { act |= 1ull << i; }</pre>
      void erase(int i) { act δ= ~(1ull << i); }</pre>
      template <class T>
107
      void init(const T &bts, int shift = 0, int s0 = 0,
                int s1 = 0) {
        if (s0 + s1 >= M) act = 0;
109
        else
          act = bts.getRange(shift + s0, shift + M - 1 - s1)
111
                << s0;
113
    };
```

2.7. Wavelet Matrix

```
#pragma GCC target("popcnt,bmi2")
   #include <immintrin.h>
   // T is unsigned. You might want to compress values first
   template <typename T> struct wavelet_matrix {
     static_assert(is_unsigned_v<T>, "only unsigned T");
     struct bit vector {
        static constexpr uint W = 64;
g
        uint n, cnt0:
        vector<ull> bits;
        vector<uint> sum;
11
        bit_vector(uint n_)
        : n(n_), bits(n / W + 1), sum(n / W + 1) {}
void build() {
13
          for (uint j = 0; j != n / W; ++j)
15
            sum[j + 1] = sum[j] + _mm_popcnt_u64(bits[j]);
          cnt0 = rank0(n);
17
        void set_bit(uint i) { bits[i / W] |= 1ULL << i % W; }</pre>
19
        bool operator[](uint i) const {
          return !!(bits[i / W] & 1ULL << i % W);</pre>
23
        uint rank1(uint i) const {
          return sum[i / W]
                  mm_popcnt_u64(_bzhi_u64(bits[i / W], i % W));
       uint rank0(uint i) const { return i - rank1(i); }
27
29
     uint n, lg;
     vector<bit_vector> b;
     wavelet_matrix(const vector<T> &a) : n(a.size()) {
31
         _{
m lg(max(*max\_element(a.begin(), a.end()), T(1))) + 1;}
33
        b.assign(lg, n);
35
        vector<T> cur = a, nxt(n);
        for (int h = lg; h--;) {
  for (uint i = 0; i < n; ++i)</pre>
37
            if (cur[i] & (T(1) << h)) b[h].set_bit(i);
39
          b[h].build();
          int il = 0, ir = b[h].cnt0;
for (uint i = 0; i < n; ++i)
   nxt[(b[h][i] ? ir : il)++] = cur[i];</pre>
41
43
          swap(cur, nxt);
45
     T operator[](uint i) const {
        T res = 0;
47
        for (int h = lg; h--;)
49
          if (b[h][i])
```

```
i += b[h].cnt0 - b[h].rank0(i), res |= T(1) << h;
51
          else i = b[h].rank0(i);
       return res:
53
      // query k-th smallest (0-based) in a[l, r)
55
     T kth(uint l, uint r, uint k) const {
        T res = 0;
57
        for (int h = lg; h--;)
          uint tl = b[h].rank0(l), tr = b[h].rank0(r);
          if (k >= tr - tl) {
            k -= tr - tl;
61
            l += b[h].cnt0
            r += b[h].cnt0 - tr;
63
            res |= T(1) << h;
          } else l = tl, r = tr;
65
        }
       return res;
67
     }
      // count of i in [l, r) with a[i] < u
     uint count(uint l, uint r, T u) const {
  if (u >= T(1) << lg) return r - l;</pre>
69
        uint res = 0;
for (int h = lg; h--;) {
71
73
          uint tl = b[h].rank0(l), tr = b[h].rank0(r);
          if (u & (T(1) << h)) {
75
            l += b[h].cnt0 - tl;
            r += b[h].cnt0 - tr;
77
            res += tr - tl;
          } else l = tl, r = tr;
79
        return res;
     }
81
   };
```

2.8. Link-Cut Tree

```
const int MXN = 100005;
   const int MEM = 100005;
   struct Splay {
 5
      static Splay nil, mem[MEM], *pmem;
      Splay *ch[2], *f;
      int val, rev, size;
Splay(): val(-1), rev(0), size(0) {
  f = ch[0] = ch[1] = 8nil;
 9
     Splay(int _val) : val(_val), rev(θ), size(1) {
   f = ch[θ] = ch[1] = δnil;
11
13
      bool isr() {
        return f->ch[0] != this δδ f->ch[1] != this;
15
17
      int dir() { return f->ch[0] == this ? 0 : 1; }
      void setCh(Splay *c, int d) {
19
        ch[d] = c;
        if (c != &nil) c->f = this;
21
        pull();
      }
      void push() {
23
        if (rev)
          swap(ch[0], ch[1]);
if (ch[0] != &nil) ch[0]->rev ^= 1;
25
           if (ch[1] != &nil) ch[1]->rev ^= 1;
27
          rev = 0:
29
        }
      }
      void pull() {
31
        size = ch[0]->size + ch[1]->size + 1;
if (ch[0] != &nil) ch[0]->f = this;
33
        if (ch[1] != δnil) ch[1]->f = this;
35
    } Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem;
37
   Splay *nil = &Splay::nil;
39
   void rotate(Splay *x) {
      Splay *p = x->f;
int d = x->dir();
41
      if (!p->isr()) p->f->setCh(x, p->dir());
      else x->f = p->f
43
      p->setCh(x->ch[!d], d);
45
      x->setCh(p, !d);
      p->pull();
47
      x->pull();
49
   vector<Splay *> splayVec;
   void splay(Splay *x) {
51
      splayVec.clear();
      for (Splay *q = x;; q = q \rightarrow f) {
```

```
splayVec.push_back(q);
 55
         if (q->isr()) break;
 57
       reverse(begin(splayVec), end(splayVec));
       for (auto it : splayVec) it->push();
while (!x->isr()) {
 59
         if (x->f->isr()) rotate(x);
         else if (x->dir() == x->f->dir())
            rotate(x->f), rotate(x);
 63
         else rotate(x), rotate(x);
 65
    }
    Splay *access(Splay *x) {
 67
       Splay *q = nil;
for (; x != nil; x = x->f) {
 69
         splay(x);
 71
         x->setCh(q, 1);
         q = x;
       }
 73
       return q;
    }
     void evert(Splay *x) {
 77
       access(x);
       splay(x);
x->rev ^=
 79
       x->push();
       x->pull();
 83
    void link(Splay *x, Splay *y) {
       // evert(x);
 85
       access(x):
       splav(x):
 87
       evert(v):
       x->setCh(y, 1);
    }
 89
    void cut(Splay *x, Splay *y) {
       // evert(x);
 91
       access(y);
 93
       splay(y);
       y->push();
 95
       y->ch[0] = y->ch[0]->f = nil;
 97
    int N, Q;
    Splay *vt[MXN];
    int ask(Splay *x, Splay *y) {
       access(x);
103
       access(v);
       splay(x);
       int res = x->f->val;
105
       if (res == -1) res = x->val;
107
       return res;
100
    int main(int argc, char **argv) {
111
       scanf("%d%d", &N, &Q);
       for (int i = 1; i <= N; i++)
113
         vt[i] = new (Splay::pmem++) Splay(i);
       while (Q--) {
         char cmd[105];
115
         int u, v;
scanf("%s", cmd);
if (cmd[1] == 'i') {
    scanf("%d%d", &u, &v);
            link(vt[v], vt[u]);
         } else if (cmd[0] ==
121
           scanf("%d", &v);
cut(vt[1], vt[v]);
123
         } else {
            scanf("%d%d", &u, &v);
125
            int res = ask(vt[u], vt[v]);
            printf("%d\n", res);
127
129
       }
```

3. Graph

Modeling

- Maximum/Minimum flow with lower bound / Circulation problem 19

 - 1. Construct super source S and sink T. 2. For each edge (x, y, l, u), connect $x \to y$ with capacity u - l.
 - 3. For each vertex v, denote by in(v) the difference between the sum of incoming lower bounds and the sum of outgoing lower bounds. 23

- 4. If in(v) > 0, connect $S \to v$ with capacity in(v), otherwise, connect $v \to T$ with capacity -in(v).
 - To maximize, connect $t \to s$ with capacity ∞ (skip this in circulation problem), and let f be the maximum flow from S to T. If $f \neq \sum_{v \in V, in(v) > 0} in(v)$, there's no solution. Otherwise,

the maximum flow from s to t is the answer. — To minimize, let f be the maximum flow from S to T. Connect $t \to s$ with capacity ∞ and let the flow from S to T be f'. If $f+f'\neq \sum_{v\in V, in(v)>0} in(v),$ there's no solution. Otherwise, ' is the answer.

5. The solution of each edge e is $l_e + f_e$, where f_e corresponds to the flow of edge e on the graph.

- \bullet Construct minimum vertex cover from maximum matching M on bipartite graph (X, Y)
- 1. Redirect every edge: $y \to x$ if $(x, y) \in M$, $x \to y$ otherwise. 2. DFS from unmatched vertices in X. 3. $x \in X$ is chosen iff x is unvisited. 4. $y \in Y$ is chosen iff y is visited. Minimum cost cyclic flow

1. Consruct super source S and sink T

- 2. For each edge (x, y, c), connect $x \to y$ with (cost, cap) = (c, 1) if c>0, otherwise connect $y\to x$ with (cost,cap)=(-c,1) 3. For each edge with c<0, sum these cost as K, then increase d(y)
- by 1, decrease d(x) by 1
- 4. For each vertex v with d(v) > 0, connect $S \to v$ with (cost, cap) =(0,d(v))
- 5. For each vertex v with d(v) < 0, connect $v \to T$ with (cost, cap) =(0, -d(v))
- 6. Flow from S to T, the answer is the cost of the flow C+K
- Maximum density induced subgraph

 - 1. Binary search on answer, suppose we're checking answer T 2. Construct a max flow model, let K be the sum of all weights 3. Connect source $s \to v, \ v \in G$ with capacity K

 - 4. For each edge (u, v, w) in G, connect $u \to v$ and $v \to u$ with
- capacity w 5. For $v \in G$, connect it with sink $v \to t$ with capacity K + 2T - $\left(\sum_{e \in E(v)} w(e)\right) - 2w(v)$
- 6. T is a valid answer if the maximum flow f < K|V|
- Minimum weight edge cover
 - 1. For each $v \in V$ create a copy v', and connect $u' \to v'$ with weight w(u,v).
- 2. Connect $v \to v'$ with weight $2\mu(v)$, where $\mu(v)$ is the cost of the cheapest edge incident to v.
- 3. Find the minimum weight perfect matching on G'.
- Project selection problem
 - 1. If $p_v > 0$, create edge (s, v) with capacity p_v ; otherwise, create edge (v, t) with capacity $-p_v$.
 - 2. Create edge (u, v) with capacity w with w being the cost of choos-
- ing u without choosing v.

 3. The mincut is equivalent to the maximum profit of a subset of
- projects.
 0/1 quadratic programming

$$\sum_{x} c_{x}x + \sum_{y} c_{y}\bar{y} + \sum_{xy} c_{xy}x\bar{y} + \sum_{xyx'y'} c_{xyx'y'}(x\bar{y} + x'\bar{y'})$$

can be minimized by the mincut of the following graph:

- 1. Create edge (x,t) with capacity c_x and create edge (s,y) with capacity c_y .
- 2. Create edge (x, y) with capacity c_{xy} .
- 3. Create edge (x, y) and edge (x', y') with capacity $c_{xyx'y'}$.

3.2. Matching/Flows

3.2.1. Dinic's Algorithm

```
struct Dinic {
      struct edge {
        int to, cap, flow, rev;
 5
      static constexpr int MAXN = 1000, MAXF = 1e9;
      vector<edge>_v[MAXN];
      int top[MAXN], deep[MAXN], side[MAXN], s, t;
 7
     void make_edge(int s, int t, int cap) {
  v[s].push_back({t, cap, 0, (int)v[t].size()})
 9
        v[t].push_back({s, 0, 0, (int)v[s].size() - 1});
11
     int dfs(int a, int flow) {
  if (a == t || !flow) return flow;
13
        for (int &i = top[a]; i < v[a].size(); i++) {</pre>
          edge &e = v[a][i];
          if (deep[a] + 1 == deep[e.to] &δ e.cap - e.flow) {
             int x = dfs(e.to, min(e.cap - e.flow, flow));
17
             if (x) {
               e.flow += x, v[e.to][e.rev].flow -= <math>x;
               return x;
             }
          }
```

```
deep[a] = -1:
25
       return 0;
27
     bool bfs() {
       queue<int> q;
       fill_n(deep, MAXN, 0);
29
        q.push(s), deep[s] = 1;
31
       while (!q.empty()) {
33
          tmp = q.front(), q.pop();
          for (edge e : v[tmp])
35
            if (!deep[e.to] && e.cap != e.flow)
              deep[e.to] = deep[tmp] + 1, q.push(e.to);
37
       return deep[t];
39
     int max_flow(int _s, int _t) {
41
       s = _s, t = _t;
int flow = 0, tflow;
       while (bfs()) {
43
          fill_n(top, MAXN, 0);
          while ((tflow = dfs(s, MAXF))) flow += tflow;
45
47
       return flow;
     }
     void reset() {
49
       fill_n(side, MAXN, 0);
       for (auto &i : v) i.clear();
51
53 };
```

3.2.2. Minimum Cost Flow

```
1
   struct MCF {
      struct edge {
        ll to, from, cap, flow, cost, rev;
      } *fromE[MAXN];
      vector<edge> v[MAXN];
      ll n, s, t, flows[MAXN], dis[MAXN], pi[MAXN], flowlim;
      void make_edge(int s, int t, ll cap, ll cost) {
        if (!cap) return;
9
        v[s].pb(edge{t, s, cap, 0LL, cost, v[t].size()});
        v[t].pb(edge\{s, t, 0LL, 0LL, -cost, v[s].size() - 1\});
11
     bitset<MAXN> vis
      void dijkstra() {
13
        vis.reset();
         _gnu_pbds::priority_queue<pair<ll, <mark>int</mark>>> q;
15
        vector<decltype(q)::point_iterator> its(n);
17
        q.push({0LL, s})
        while (!q.empty()) {
19
          int now = q.top().second;
          q.pop();
          if (vis[now]) continue;
          vis[now] = 1;
          ll ndis = dis[now] + pi[now];
23
          for (edge &e : v[now]) {
            if (e.flow == e.cap || vis[e.to]) continue;
            if (dis[e.to] > ndis + e.cost - pi[e.to]) {
  dis[e.to] = ndis + e.cost - pi[e.to];
               flows[e.to] = min(flows[now], e.cap - e.flow);
29
               fromE[e.to] = &e;
              if (its[e.to] == q.end())
                 its[e.to] = q.push({-dis[e.to], e.to});
              else q.modify(its[e.to], {-dis[e.to], e.to});
            }
33
          }
35
        }
     bool AP(ll &flow) {
37
        fill_n(dis, n, INF);
39
        fromE[s] = 0;
        dis[s] = 0;
        flows[s] = flowlim - flow;
41
        dijkstra();
        if (dis[t] == INF) return false;
flow += flows[t];
43
        for (edge *e = fromE[t]; e; e = fromE[e->from]) {
          e->flow += flows[t];
47
          v[e->to][e->rev].flow -= flows[t];
        for (int i = 0; i < n; i++)
          pi[i] = min(pi[i] + dis[i], INF);
51
        return true:
     pll solve(int _s, int _t, ll _flowlim = INF) {
   s = _s, t = _t, flowlim = _flowlim;
53
        pll re:
55
        while (re.F != flowlim && AP(re.F))
```

```
57
        for (int i = 0; i < n; i++)
59
          for (edge δe : v[i])
            if (e.flow != 0) re.S += e.flow * e.cost;
        re.S /= 2;
61
       return re;
63
     void init(int _n) {
65
       n = n;
        fill_n(pi, n, 0);
        for (int i = 0; i < n; i++) v[i].clear();</pre>
67
69
     void setpi(int s) {
        fill_n(pi, n, INF);
       pi[s] = 0;
71
        for (ll it = 0, flag = 1, tdis; flag \delta\delta it < n; it++) {
73
          flag = 0;
          for (int i = 0; i < n; i++)
if (pi[i] != INF)</pre>
75
              for (edge &e : v[i])
77
                 if (e.cap && (tdis = pi[i] + e.cost) < pi[e.to])
                   pi[e.to] = tdis, flag = 1;
79
     }
81 };
```

3.2.3. Gomory-Hu Tree

Requires: Dinic's Algorithm

```
1 int e[MAXN][MAXN];
   int p[MAXN];
Dinic D; // original graph
 3
    void gomory_hu() {
      fill(p, p + n, 0);
fill(e[0], e[n], INF);
      for (int s = 1; s < n; s++) {
         int t = p[s];
Dinic F = D;
 9
         int tmp = F.max_flow(s, t);
for (int i = 1; i < s; i++)</pre>
11
            e[s][i] = e[i][s] = min(tmp, e[t][i]);
         for (int i = s + 1; i <= n; i++)
13
            if (p[i] == t && F.side[i]) p[i] = s;
15
      }
```

3.2.4. Global Minimum Cut

```
// weights is an adjacency matrix, undirected
   pair<int, vi> getMinCut(vector<vi> &weights) {
     int N = sz(weights);
     vi used(N), cut, best_cut;
5
     int best_weight = -1;
     for (int phase = N - 1; phase >= 0; phase--) {
  vi w = weights[0], added = used;
7
9
        int prev, k = 0;
        rep(i, 0, phase) {
11
          prev = k;
          k = -1;
13
          rep(j, 1, N) if (!added[j] &&
                             (k == -1 \mid \mid w[j] > w[k])) k = j;
          if (i == phase - 1) {
15
            rep(j, 0, N) weights[prev][j] += weights[k][j];
            rep(j, 0, N) weights[j][prev] = weights[prev][j];
17
            used[k] = true;
19
            cut.push back(k);
            if (best_weight == -1 || w[k] < best_weight) {
  best_cut = cut;</pre>
21
              best_weight = w[k];
23
          } else {
            rep(j, 0, N) w[j] += weights[k][j];
25
            added[k] = true;
27
       }
29
     return {best_weight, best_cut};
31 }
```

3.2.5. Bipartite Minimum Cover

Requires: Dinic's Algorithm

```
// maximum independent set = all vertices not covered
// x : [0, n), y : [0, m]
struct Bipartite_vertex_cover {
   Dinic D;
```

```
5
      int n, m, s, t, x[maxn], y[maxn];
void make_edge(int x, int y) { D.make_edge(x, y + n, 1); }
       int matching() {
          int re = D.max_flow(s, t);
 q
          for (int i = 0; i < n; i++)
            for (Dinic::edge &e : D.v[i])
               if (e.to != s && e.flow == 1) {
11
                  x[i] = e.to - n, y[e.to - n] = i;
                  break;
13
15
          return re;
       // init() and matching() before use
void solve(vector<int> &vx, vector<int> &vy) {
17
19
          bitset<maxn * 2 + 10> vis;
          queue<int> q;
         for (int i = 0; i < n; i++)
   if (x[i] == -1) q.push(i), vis[i] = 1;
while (!q.empty()) {
   int i = 0; i < n; i++)
   if (x[i] == -1) q.push(i), vis[i] = 1;</pre>
21
23
            int now = q.front();
25
            q.pop();
            if (now < n) {
27
               for (Dinic::edge &e : D.v[now])
                  if (e.to != s && e.to - n != x[now] && !vis[e.to])
29
                    vis[e.to] = 1, q.push(e.to);
            } else {
               if (!vis[y[now - n]])
                  vis[y[now - n]] = 1, q.push(y[now - n]);
33
            }
          for (int i = 0; i < n; i++)
35
         if (!vis[i]) vx.pb(i);
for (int i = 0; i < m; i++)
  if (vis[i + n]) vy.pb(i);</pre>
37
39
       void init(int _n, int _m) {
          n = _n, m = _m, s = n + m, t
for (int i = 0; i < n; i++)
41
                                                 = s + 1;
43
            x[i] = -1, D.make_edge(s, i, 1);
          for (int i = 0; i < m; i++)
45
            y[i] = -1, D.make_edge(i + n, t, 1);
47 };
```

3.2.6. Edmonds' Algorithm

```
struct Edmonds {
 1
      int n, T;
vector<vector<int>> g;
       vector<<mark>int</mark>> pa, p, used, base;
       Edmonds(int n)
            : n(n), T(\theta), g(n), pa(n, -1), p(n), used(n),
              base(n) {}
       void add(int a, int b) {
 9
         g[a].push_back(b);
         g[b].push_back(a);
11
       int getBase(int i) {
         while (i != base[i])
13
            base[i] = base[base[i]], i = base[i];
15
       vector<int> toJoin;
       void mark_path(int v, int x, int b, vector<int> &path) {
  for (; getBase(v) != b; v = p[x]) {
    p[v] = x, x = pa[v];
19
            toJoin.push_back(v);
toJoin.push_back(x);
21
            if (!used[x]) used[x] = ++T, path.push_back(x);
23
         }
       bool go(int v) {
         for (int x : g[v]) {
27
            int b, bv = getBase(v), bx = getBase(x);
            if (bv == bx) {
29
               continue;
            } else if (used[x]) {
31
               vector<int> path;
33
               toJoin.clear();
               if (used[bx] < used[bv])</pre>
               mark_path(v, x, b = bx, path);
else mark_path(x, v, b = bv, path);
for (int z : toJoin) base[getBase(z)] = b;
35
37
               for (int z : path)
  if (go(z)) return 1;
39
            } else if (p[x] == -1) {
              p[x] = v;
if (pa[x] == -1) {
41
```

```
43
              for (int y; x != -1; x = v)
                y = p[x], v = pa[y], pa[x] = y, pa[y] = x;
              return 1;
45
            if (!used[pa[x]]) {
47
              used[pa[x]] = ++T;
49
              if (go(pa[x])) return 1;
51
         }
53
       return 0;
55
     void init_dfs() {
        for (int i = 0; i < n; i++)
          used[i] = 0, p[i] = -1, base[i] = i;
57
59
     bool dfs(int root) {
       used[root] = ++T;
61
       return go(root):
63
     void match() {
       int ans = 0;
for (int v = 0; v < n; v++)
65
          for (int x : g[v])
            if (pa[v] == -1 \&\& pa[x] == -1) {
67
              pa[v] = x, pa[x] = v, ans++;
              break;
69
            }
71
        init_dfs();
        for (int i = 0; i < n; i++)
        if (pa[i] == -1 88 dfs(i)) ans++, init_dfs(); cout << ans * 2 << "\n";
73
75
        for (int i = 0; i < n; i++)
          if (pa[i] > i)
            cout << i + 1 << " " << pa[i] + 1 << "\n";
77
79 };
```

3.2.7. Minimum Weight Matching

```
1 struct Graph {
      static const int MAXN = 105;
      int n, e[MAXN][MAXN];
      int match[MAXN], d[MAXN], onstk[MAXN];
      vector<int> stk;
 5
      void init(int _n) {
         n = _n;
for (int i = 0; i < n; i++)</pre>
           for (int j = 0; j < n; j++)
 9
              // change to appropriate infinity
// if not complete graph
11
              e[i][j] = 0;
13
      void add_edge(int u, int v, int w) {
        e[u][v] = e[v][u] = w;
15
17
      bool SPFA(int u) {
         if (onstk[u]) return true;
         stk.push_back(u);
onstk[u] = 1;
19
         for (int v = 0; v < n; v++) {
21
           if (u != v && match[u] != v && !onstk[v]) {
              int m = match[v];
if (d[m] > d[u] - e[v][m] + e[u][v]) {
    d[m] = d[u] - e[v][m] + e[u][v];
    d[m] = d[u] - e[v][m] + e[u][v];
23
25
                 onstk[v] = 1:
27
                 stk.push_back(v);
                 if (SPFA(m)) return true;
29
                stk.pop_back();
onstk[v] = 0;
31
           }
33
         onstk[u] = 0;
35
         stk.pop_back();
         return false;
37
      int solve() {
  for (int i = 0; i < n; i += 2) {</pre>
39
           match[i] = i + 1;
41
           match[i + 1] = i;
43
         while (true) {
           int found = 0;
           for (int i = 0; i < n; i++) onstk[i] = d[i] = 0; for (int i = 0; i < n; i++) {
45
47
              stk.clear();
              if (!onstk[i] && SPFA(i)) {
49
                 found = 1;
```

```
while (stk.size() >= 2) {
51
                int u = stk.back();
                stk.pop_back();
53
                int v = stk.back();
                stk.pop_back();
                match[u] = v;
55
                match[v] = u;
57
            }
59
          if (!found) break;
61
63
        for (int i = 0; i < n; i++) ret += e[i][match[i]];</pre>
        ret /= 2:
65
        return ret;
67
   } graph;
```

3.2.8. Stable Marriage

```
// normal stable marriage problem
    /* input:
 3
   3
    Albert Laura Nancy Marcy
   Brad Marcy Nancy Laura
    Chuck Laura Marcy Nancy
    Laura Chuck Albert Brad
    Marcy Albert Chuck Brad
   Nancy Brad Albert Chuck
11
13
   using namespace std;
    const int MAXN = 505;
15
   int n;
int favor[MAXN][MAXN]; // favor[boy_id][rank] = girl_id;
int order[MAXN][MAXN]; // order[girl_id][boy_id] = rank;
int current[MAXN]; // current[boy_id] = rank;
17
19
    // boy_id will pursue current[boy_id] girl.
21
   int girl_current[MAXN]; // girl[girl_id] = boy_id;
    void initialize() {
23
      for (int i = 0; i < n; i++) {
25
        current[i] = 0;
        girl_current[i] = n;
        order[i][n] = n;
29
    map<string, int> male, female;
31
    string bname[MAXN], gname[MAXN];
33
    int fit = 0;
35
    void stable_marriage() {
      queue<int> que;
for (int i = \theta; i < n; i++) que.push(i);
37
      while (!que.empty()) {
39
        int boy_id = que.front();
41
        que.pop();
43
        int girl_id = favor[boy_id][current[boy_id]];
        current[boy_id]++;
45
        if (order[girl_id][boy_id] <</pre>
           order[girl_id][girl_current[girl_id]]) {
if (girl_current[girl_id] < n)</pre>
47
49
             que.push(girl_current[girl_id]);
           girl_current[girl_id] = boy_id;
51
        } else {
           que.push(boy_id);
53
      }
55 }
57
   int main() {
      cin >> n:
59
      for (int i = 0; i < n; i++) {
61
        string p, t;
        cin >> p;
male[p] = i;
63
        bname[i] = p;
        for (int j = 0; j < n; j++) {
65
           cin >> t;
           if (!female.count(t)) {
67
             gname[fit] = t;
```

```
69
           female[t] = fit++;
71
         favor[i][j] = female[t];
       }
73
75
     for (int i = 0; i < n; i++) {
       string p, t;
77
       cin >> p;
       for (int j = 0; j < n; j++) {
79
         order[female[p]][male[t]] = j;
81
83
     initialize();
     stable_marriage();
     for (int i = 0; i < n; i++) {
87
       cout << bname[i] <<
89
            << gname[favor[i][current[i] - 1]] << endl;</pre>
91 }
```

3.2.9. Kuhn-Munkres algorithm

```
1 // Maximum Weight Perfect Bipartite Matching
    // Detect non-perfect-matching:
   // 1. set all edge[i][j] as INF
    // 2. if solve() >= INF, it is not perfect matching.
 5
   typedef long long ll;
   struct KM {
      static const int MAXN = 1050;
 9
      static const ll INF = 1LL << 60;</pre>
      int n, match[MAXN], vx[MAXN], vy[MAXN];
ll edge[MAXN][MAXN], lx[MAXN], ly[MAXN], slack[MAXN];
11
      void init(int _n) {
        n = _n;
for (int i = 0; i < n; i++)
13
           for (int j = 0; j < n; j++) edge[i][j] = 0;</pre>
15
      void add_edge(int x, int y, ll w) { edge[x][y] = w; }
17
      bool DFS(int x) {
19
        vx[x] = 1:
        for (int y = 0; y < n; y++) {
           if (vy[y]) continue;
21
           if (lx[x] + ly[y] > edge[x][y]) {
23
             slack[v]
             min(slack[y], lx[x] + ly[y] - edge[x][y]);
           } else {
25
27
             if (match[y] == -1 || DFS(match[y])) {
               match[y] = x;
29
               return true;
31
          }
33
        return false;
35
      ll solve() {
        fill(match, match + n, -1);
        fill(lx, lx + n, -INF);
fill(ly, ly + n, 0);
37
        for (int i = 0; i < n; i++)
for (int j = 0; j < n; j++)
39
        lx[i] = max(lx[i], edge[i][j]);
for (int i = 0; i < n; i++) {</pre>
41
           fill(slack, slack + n, INF);
43
           while (true) {
             fill(vx, vx + n, 0);
fill(vy, vy + n, 0);
if (DFS(i)) break;
45
47
             ll d = INF;
49
             for (int j = 0; j < n; j++)
               if (!vy[j]) d = min(d, slack[j]);
             for (int j = 0; j < n; j++) {
                if (vx[j]) lx[j] -= d;
53
                if (vy[j]) ly[j] += d;
               else slack[j] -= d;
55
             }
          }
57
        il res = 0;
for (int i = 0; i < n; i++) {</pre>
59
          res += edge[match[i]][i];
61
        return res;
```

```
63 | } graph;
```

3.4.1. 2-Satisfiability

Requires: Strongly Connected Components

3.3. Shortest Path Faster Algorithm

```
struct SPFA {
      static const int maxn = 1010, INF = 1e9;
      int dis[maxn];
      bitset<maxn> inq, inneg;
      queue<<mark>int</mark>> q, tq;
      vector<pii> v[maxn];
      void make_edge(int s, int t, int w) {
        v[s].emplace_back(t, w);
 9
      void dfs(int a) {
        inneg[a] = 1;
for (pii i : v[a])
11
13
          if (!inneg[i.F]) dfs(i.F);
      bool solve(int n, int s) { // true if have neg-cycle
for (int i = 0; i <= n; i++) dis[i] = INF;</pre>
15
        dis[s] = 0, q.push(s);
17
        for (int i = 0; i < n; i++) {
19
          inq.reset();
          int now;
          while (!q.empty()) {
            now = q.front(), q.pop();
23
             for (pii &i : v[now]) {
               if (dis[i.F] > dis[now] + i.S) {
                 dis[i.F] = dis[now] + i.S;
                 if (!inq[i.F]) tq.push(i.F), inq[i.F] = 1;
27
            }
          q.swap(tq);
31
        bool re = !q.empty();
33
        inneg.reset();
        while (!q.empty()) {
          if (!inneg[q.front()]) dfs(q.front());
35
          q.pop();
37
        return re;
39
      void reset(int n) {
41
        for (int i = 0; i <= n; i++) v[i].clear();</pre>
43 };
```

3.4. Strongly Connected Components

```
1 struct TarjanScc {
      int n, step;
vector<int> time, low, instk, stk;
       vector<vector<int>> e, scc;
      TarjanScc(int n )
      : n(n_), step(0), time(n), low(n), instk(n), e(n) {}
void add_edge(int u, int v) { e[u].push_back(v); }
      void dfs(int x) {
         time[x] = low[x] = ++step;
 9
         stk.push_back(x);
         instk[x] = 1;
for (int y : e[x])
11
13
           if (!time[y]) {
              dfs(y);
15
              low[x] = min(low[x], low[y]);
              else if (instk[y]) {
              low[x] = min(low[x], time[y]);
17
19
         if (time[x] == low[x]) {
           scc.emplace_back();
for (int y = -1; y != x;) {
21
              v = stk.back();
              stk.pop_back();
instk[y] = 0;
23
              scc.back().push_back(y);
           }
         }
27
29
       void solve() {
         for (int i = 0; i < n; i++)
  if (!time[i]) dfs(i);</pre>
31
         reverse(scc.begin(), scc.end());
33
         // scc in topological order
35
   };
```

```
1 // 1 based, vertex in SCC = MAXN * 2
  // (not i) is i + n
  struct two_SAT {
    int n, ans[MAXN];
5
    SCC S;
    void imply(int a, int b) { S.make_edge(a, b); }
    bool solve(int _n) {
      n = n;
      9
11
13
      return true;
15
    void init(int _n) {
17
      fill_n(ans, n + 1, 0);
      S.init(n * 2);
19
21 } SAT;
```

3.5. Biconnected Components

3.5.1. Articulation Points

```
1 void dfs(int x, int p) {
     tin[x] = low[x] = ++t;
3
     int ch = 0;
     for (auto u : g[x])
5
        if (u.first != p) {
          if (!ins[u.second])
            st.push(u.second), ins[u.second] = true;
          if (tin[u.first]) {
9
            low[x] = min(low[x], tin[u.first]);
            continue:
11
          ++ch;
         dfs(u.first, x);
low[x] = min(low[x], low[u.first]);
13
15
          if (low[u.first] >= tin[x]) {
            cut[x] = true;
            ++SZ;
17
            while (true) {
19
              int e = st.top();
              st.pop();
              bcc[e] = sz;
21
              if (e == u.second) break;
23
25
     if (ch == 1 && p == -1) cut[x] = false;
```

3.5.2. Bridges

```
1 // if there are multi-edges, then they are not bridges
   void dfs(int x, int p) {
     tin[x] = low[x] = ++t;
3
     st.push(x);
     for (auto u : g[x])
5
       if (u.first != p) {
7
         if (tin[u.first]) {
           low[x] = min(low[x], tin[u.first]);
9
           continue:
11
         dfs(u.first, x);
         low[x] = min(low[x], low[u.first]);
13
         if (low[u.first] == tin[u.first]) br[u.second] = true;
15
     if (tin[x] == low[x]) {
       ++SZ;
       while (st.size()) {
17
         int u = st.top();
19
         st.pop();
         bcc[u] = sz;
         if (u == x) break;
21
23
     }
```

3.6. Triconnected Components

```
1
   // requires a union-find data structure
   struct ThreeEdgeCC {
     int V, ind;
      vector<int> id, pre, post, low, deg, path;
      vector<vector<<mark>int</mark>>> components;
     UnionFind uf;
      template <class Graph>
      void dfs(const Graph &G, int v, int prev) {
9
        pre[v] = ++ind;
        for (int w : G[v])
          if (w != v) {
11
            if (w == prev) {
               prev = -1;
13
               continue;
15
             if (pre[w] != -1) {
               if (pre[w] < pre[v]) {
17
                 deg[v]++;
                 low[v] = min(low[v], pre[w]);
19
               } else ·
                 deg[v]--;
21
                 int &u = path[v];
23
                 for (; u != -1 && pre[u] <= pre[w] &&
                         pre[w] <= post[u];) {</pre>
                   uf.join(v, u);
deg[v] += deg[u];
25
                   u = path[u];
27
29
               continue:
31
             dfs(G, w, v);
            if (path[w] == -1 88 deg[w] <= 1) {
  deg[v] += deg[w];
  low[v] = min(low[v], low[w]);</pre>
33
35
               continue;
37
            if (deg[w] == 0) w = path[w];
            if (low[v] > low[w]) {
39
               low[v] = min(low[v], low[w]);
41
               swap(w, path[v]);
43
             for (; w != -1; w = path[w]) {
              uf.join(v, w);
               deg[v] += deg[w];
            }
47
        post[v] = ind;
49
      template <class Graph>
      ThreeEdgeCC(const Graph &G)
          : V(G.size()), ind(-1), id(V, -1), pre(V, -1),
            post(V), low(V, INT_MAX), deg(V, \theta), path(V, -1),
53
            uf(V) {
        for (int v = 0; v < V; v++)
if (pre[v] == -1) dfs(G, v, -1);
55
        components.reserve(uf.cnt);
57
        for (int v = 0; v < V; v++)
          if (uf.find(v) == v) {
59
            id[v] = components.size();
61
             components.emplace_back(1, v);
            components.back().reserve(uf.getSize(v));
63
        for (int v = 0; v < V; v++)
          if (id[v] == -1)
65
            components[id[v] = id[uf.find(v)]].push_back(v);
67
   };
```

3.7. Centroid Decomposition

```
void get_center(int now) {
                                                                          17
     v[now] = true
      vtx.push_back(now);
                                                                          19
      sz[now] = 1;
     mx[now] = 0;
                                                                          21
      for (int u
                   G[now])
       if (!v[u]) {
                                                                          23
          get center(u);
          mx[now] = max(mx[now], sz[u]);
                                                                          25
          sz[now] += sz[u];
11
                                                                          27
   void get_dis(int now, int d, int len) {
  dis[d][now] = cnt;
                                                                          29
13
15
                                                                          31
     v[now] = true:
     for (auto u : G[now])
```

```
17
        if (!v[u.first]) { get_dis(u, d, len + u.second); }
19 void dfs(int now, int fa, int d) {
     get_center(now);
21
     int c = -1;
     for (int i : vtx) {
        if (max(mx[i], (int)vtx.size() - sz[i]) <=</pre>
23
            (int)vtx.size() / 2)
25
        v[i] = false;
27
     get_dis(c, d, 0);
for (int i : vtx) v[i] = false;
29
     v[c] = true;
31
     vtx.clear();
     dep[c] = d;
p[c] = fa;
33
     for (auto u : G[c])
       if (u.first != fa δδ !v[u.first]) {
35
          dfs(u.first, c, d + 1);
37
   }
```

3.8. Minimum Mean Cycle

```
// d[i][j] == 0 if {i,j} !in E
               long long d[1003][1003], dp[1003][1003];
               pair<long long, long long> MMWC() {
                       memset(dp, 0x3f, sizeof(dp));
for (int i = 1; i <= n; ++i) dp[0][i] = 0;
    5
                       for (int i = 1; i <= n; ++i) {
  for (int i = 1; i <= n; ++i) {
    for (int j = 1; j <= n; ++j) {
      for (int k = 1; k <= n; ++k) {
          dp[i][k] = min(dp[i - 1][j] + d[j][k], dp[i][k]);
          dp[i][k] = min(dp[i - 1][i] + d[i][k], dp[i][k]);
          dp[i][k] = min(dp[i - 1][i] + d[i][k], dp[i][k]);
          dp[i][k] = min(dp[i - 1][i][k], dp[i][k], dp[i][k]);
          dp[i][k] = min(dp[i - 1][i][k], dp[i][k], 
   9
11
                                 }
13
                         long long au = 1ll << 31, ad = 1;</pre>
15
                         for (int i = 1; i <= n; ++i) {
                                   if (dp[n][i] == 0x3f3f3f3f3f3f3f3f3f) continue;
                                   long long u = 0, d = 1;
17
                                   for (int j = n - 1; j >= 0; --j) {
  if ((dp[n][i] - dp[j][i]) * d > u * (n - j)) {
19
                                                    u = dp[n][i] - dp[j][i];
                                                    d = n - j;
21
                                          }
23
                                  if (u * ad < au * d) au = u, ad = d;
25
                                                                                                 _gcd(au, ad);
                         long long g =
                        return make_pair(au / g, ad / g);
27
```

3.9. Directed MST

5

9

11

13

15

```
template <typename T> struct DMST {
  T g[maxn][maxn], fw[maxn];
  int n, fr[maxn];
  bool vis[maxn], inc[maxn];
  void clear() {
     for (int i = 0; i < maxn; ++i) {
  for (int j = 0; j < maxn; ++j) g[i][j] = inf;
  vis[i] = inc[i] = false;</pre>
     }
  }
  void addedge(int u, int v, T w) {
  g[u][v] = min(g[u][v], w);
  T operator()(int root, int _n) {
     if (dfs(root) != n) return -1;
     T ans = 0:
     while (true) {
       for (int i = 1; i <= n; ++i) fw[i] = inf, fr[i] = i;
       for (int i = 1; i <= n; ++i)
          if (!inc[i]) {
             for (int j
                             1; j <= n; ++j) {
               if (!inc[j] && i != j && g[j][i] < fw[i]) {</pre>
                  fw[i] = g[j][i];
fr[i] = j;
               }
            }
          }
       int x = -1;
for (int i = 1; i <= n; ++i)
if (i != root && !inc[i]) {
             int j = i, c = 0;
```

```
33
               while (i != root && fr[i] != i && c <= n)
                 ++c, j = fr[j];
               if (j == root || c > n) continue;
35
               else {
37
                 x = i:
                 break;
39
               }
          if (!~x) {
41
             for (int i = 1; i <= n; ++i)
               if (i != root δδ !inc[i]) ans += fw[i];
43
             return ans;
45
          for (int i = 1; i <= n; ++i) vis[i] = false;</pre>
          do {
49
            ans += fw[y];
            y = fr[y];
vis[y] = inc[y] = true;
51
          } while (y != x);
          inc[x] = false;
53
          for (int k = 1; k <= n; ++k)
if (vis[k]) {
55
               for (int j = 1; j <= n; ++j)
                 if (!vis[j]) {
57
                   if (g[x][j] > g[k][j]) g[x][j] = g[k][j];
                   if (g[j][k] < inf &&
    g[j][k] - fw[k] < g[j][x])</pre>
59
61
                      g[j][x] = g[j][k] - fw[k];
                 }
63
65
        return ans:
67
      int dfs(int now) {
69
        vis[now] = true;
        for (int i = 1; i \le n; ++i)
          if (g[now][i] < inf && !vis[i]) r += dfs(i);
71
        return r;
73
   }:
```

3.10. Maximum Clique

```
1 // source: KACTL
   typedef vector<bitset<200>> vb;
   struct Maxclique {
      double limit = 0.025, pk = 0;
      struct Vertex {
        int i, d = 0;
 9
      typedef vector<Vertex> vv;
      vb e;
11
      vv V:
      vector<vi> C;
13
      vi qmax, q, S, old;
      void init(vv &r) {
        for (auto \delta v : r) v.d = 0;
        for (auto &v : r)
17
          for (auto j :
                          r) v.d += e[v.i][j.i];
        sort(all(r), [](auto a, auto b) { return a.d > b.d; });
        int mxD = r[0].d;
19
        rep(i, 0, sz(r)) r[i].d = min(i, mxD) + 1;
21
      void expand(vv &R, int lev = 1) {
   S[lev] += S[lev - 1] - old[lev];
   old[lev] = S[lev - 1];
23
        while (sz(R)) {
25
          if (sz(q) + R.back().d <= sz(qmax)) return;</pre>
27
          q.push_back(R.back().i);
29
          for (auto v : R)
            if (e[R.back().i][v.i]) T.push_back({v.i});
31
          if (sz(T)) {
             if (S[lev]++ / ++pk < limit) init(T);</pre>
             int j = 0, mxk = 1,
33
                 mnk = max(sz(qmax) - sz(q) + 1, 1);
             C[1].clear(), C[2].clear();
             for (auto v : T) {
               int k = 1;
               auto f = [δ](int i) { return e[v.i][i]; };
39
               while (any_of(all(C[k]), f)) k++;
               if (k > mxk) mxk = k, C[mxk + 1].clear();
if (k < mnk) T[j++].i = v.i;</pre>
41
               C[k].push_back(v.i);
43
             if (j > 0) T[j - 1].d = 0;
```

```
45
           rep(k, mnk, mxk + 1) for (int i : C[k]) T[j].i = i,
                                                     T[j++].d =
47
           expand(T, lev + 1);
         } else if (sz(q) > sz(qmax)) qmax = q;
49
         q.pop_back(), R.pop_back();
       }
51
     vi maxClique() {
53
       init(V), expand(V);
55
       return qmax;
57
     Maxclique(vb conn)
          : e(conn), C(sz(e) + 1), S(sz(C)), old(S) {
59
       rep(i, 0, sz(e)) V.push_back({i});
     }
61 };
```

```
3.11. Dominator Tree
1 // idom[n] is the unique node that strictly dominates n but
   // does not strictly_dominate any other node that strictly
   // dominates n. idom[n] = 0 if n is entry or the entry
   // cannot reach n.
   struct DominatorTree {
     static const int MAXN = 200010;
     int n, s;
     vector<int> g[MAXN], pred[MAXN];
vector<int> cov[MAXN];
     int dfn[MAXN], nfd[MAXN], ts;
     int par[MAXN]
11
     int sdom[MAXN], idom[MAXN];
13
     int mom[MAXN], mn[MAXN];
15
     inline bool cmp(int u, int v) { return dfn[u] < dfn[v]; }</pre>
17
     int eval(int u) {
       if (mom[u] == u) return u;
        int res = eval(mom[u])
19
        if (cmp(sdom[mn[mom[u]]), sdom[mn[u]]))
21
          mn[u] = mn[mom[u]];
       return mom[u] = res;
23
     void init(int _n, int _s) {
25
       n = _n;
s = _s;
27
       REP1(i, 1, n) {
          g[i].clear();
29
          pred[i].clear();
31
          idom[i] = 0;
       }
33
     }
     void add_edge(int u, int v) {
35
       g[u].push_back(v);
       pred[v].push_back(u);
37
     void DFS(int u) {
39
        dfn[\dot{u}] = ts;
        nfd[ts] = u;
        for (int v
                    : g[u])
          if (dfn[v] == 0) {
43
            par[v] = u;
45
            DFS(v);
          }
47
     void build() {
49
       ts = 0;
       REP1(i, 1, n) {
   dfn[i] = nfd[i] = 0;
51
          cov[i].clear();
53
          mom[i] = mn[i] = sdom[i] = i;
       DFS(s);
55
        for (int i = ts; i >= 2; i--) {
57
          int u = nfd[i];
          if (u == 0) continue;
59
          for (int v : pred[u])
            if (dfn[v]) {
61
              eval(v):
              if (cmp(sdom[mn[v]], sdom[u]))
63
                sdom[u] = sdom[mn[v]]:
          cov[sdom[u]].push_back(u);
65
          mom[u] = par[u];
for (int w : cov[par[u]]) {
67
            eval(w):
69
            if (cmp(sdom[mn[w]], par[u])) idom[w] = mn[w];
```

```
else idom[w] = par[u];

cov[par[u]].clear();

REP1(i, 2, ts) {
   int u = nfd[i];
   if (u == 0) continue;
   if (idom[u] != sdom[u]) idom[u] = idom[idom[u]];

}

}
}
dom;
```

3.12. Manhattan Distance MST

```
// returns [(dist, from, to), ...]
   // then do normal mst afterwards
   typedef Point<int> P;
3
   vector<array<int, 3>> manhattanMST(vector<P> ps) {
     vi id(sz(ps));
     iota(all(id), 0);
     vector<array<int, 3>> edges;
     rep(k, 0, 4) {
       sort(all(id), [8](int i, int j) {
         return (ps[i] - ps[j]).x < (ps[j] - ps[i]).y;
11
       map<int, int> sweep;
for (int i : id) {
13
          for (auto it = sweep.lower_bound(-ps[i].y);
               it != sweep.end(); sweep.erase(it++)) {
15
            int j = it->second;
           P d = ps[i] - ps[j];
17
           if (d.y > d.x) break;
19
           edges.push_back({d.y + d.x, i, j});
21
          sweep[-ps[i].y] = i;
       for (P &p : ps)
         if (k \& 1) p.x = -p.x;
25
         else swap(p.x, p.y);
     return edges;
```

3.13. Virtual Tree

Requires: adamant HLD

```
// id[u] is the index of u in pre-order traversal
   vector<pii> build(vector<int> h) {
      sort(h.begin(), h.end(),
      [8](int u, int v) { return id[u] < id[v]; });
int root = h[0], top = 0;
for (int i : h) root = lca(i, root);
      vector<int> stk(h.size(), root);
      vector<pii> e;
9
      for (int u : h) {
        if (u == root) continue;
11
        int l = lca(u, stk[top]);
        if (l != stk[top]) {
           while (id[l] < id[stk[top - 1]])</pre>
13
             e.emplace_back(stk[top - 1], stk[top]), top--;
          e.emplace_back(stk[top], l), top--;
if (l != stk[top]) stk[++top] = l;
15
17
        stk[++top] = u;
19
      while (top) e.emplace_back(stk[top - 1], stk[top]), top--;
      return e:
   }
```

4. Math

4.1. Number Theory

4.1.1. Mod Struct

 $\begin{array}{l} A\ list\ of\ safe\ primes:\ 26003,27767,28319,28979,29243,29759,30467\\ 910927547,919012223,947326223,990669467,1007939579,1019126699\\ 929760389146037459,975500632317046523,989312547895528379 \end{array}$

```
 \begin{array}{|c|c|c|c|c|c|} \hline NTT \ prime \ p & p-1 & primitive \ root \\ \hline 65537 & 1 \ll 16 & 3 \\ 469762049 & 7 \ll 26 & 3 \\ 998244353 & 119 \ll 23 & 3 \\ 2748779069441 & 5 \ll 39 & 3 \\ 1945555039024054273 & 27 \ll 56 & 5 \\ \hline \end{array}
```

Requires: Extended GCD

```
1 template <typename T> struct M {
      static T MOD; // change to constexpr if already known
 3
      T v;
      M(T x = 0) \{
 5
        v = (-MOD \le x \&\& x \le MOD) ? x : x % MOD;
         if (v < 0) v += MOD;
 7
      explicit operator T() const { return v; }
 9
      bool operator==(const M &b) const { return v == b.v; }
      bool operator!=(const M &b) const { return v != b.v; }
      M operator-() { return M(-v); }
11
      M operator-() { return m( v /, )
M operator-(M b) { return M(v + b.v); }
M operator-(M b) { return M(v - b.v); }
M operator*(M b) { return M((_int128)v * b.v % MOD); }
13
      M operator/(M b) { return *this * (b ^{\land} (MOD - 2)); }
15
       // change above implementation to this if MOD is not prime
17
      M inv() {
        auto [p, _, g] = extgcd(v, MOD);
19
         return assert(g == 1), p;
21
      friend M operator^(M a, ll b) {
         M ans(1);
         for (; b; b >>= 1, a *= a)
23
           if (b & 1) ans *= a;
25
         return ans;
27
      friend M &operator+=(M &a, M b) { return a = a + b; }
      friend M & operator = (M & a, M b) { return a = a - b; } friend M & operator *= (M & a, M b) { return a = a * b; }
29
      friend M \deltaoperator/=(M \deltaa, M b) { return a = a / b; }
31 }:
   using Mod = M<int>;
template <> int Mod::MOD = 1'000'000'007;
   int &MOD = Mod::MOD;
```

4.1.2. Miller-Rabin

Requires: Mod Struct

```
1  // checks if Mod::MOD is prime
bool is_prime() {
3    if (MOD < 2 || MOD % 2 == 0) return MOD == 2;
    Mod A[] = {2, 7, 61}; // for int values (< 2^31)
5    // ll: 2, 325, 9375, 28178, 450775, 9780504, 1795265022
    int s = __builtin_ctzll(MOD - 1), i;
6    for (Mod a : A) {
        Mod x = a ^ (MOD >> s);
6    for (i = 0; i < s && (x + 1).v > 2; i++) x *= x;
        if (i && x != -1) return 0;
11    }
12    return 1;
13 }
```

4.1.3. Linear Sieve

```
constexpr ll MAXN = 1000000;
    bitset<MAXN> is_prime;
    vector<ll> primes;
    ll mpf[MAXN], phi[MAXN], mu[MAXN];
    void sieve() {
      is_prime.set();
      is_prime[1] = 0;
      mu[1] = phi[1] = 1;
for (ll i = 2; i < MAXN; i++) {
 9
11
         if (is_prime[i]) {
            mpf[\overline{i}] = i;
13
            primes.push_back(i);
            phi[i] = i - 1;
15
            mu[i] = -1;
17
         for (ll p : primes) {
           if (p > mpf[i] || i * p >= MAXN) break;
is_prime[i * p] = 0;
mpf[i * p] = p;
mpf[i * p] = p;
19
            mu[i * p] = -mu[i];
if (i % p == 0)
   phi[i * p] = phi[i] * p, mu[i * p] = 0;
21
23
            else phi[i * p] = phi[i] * (p - 1);
25
      }
27 }
```

4.1.4. Get Factors

Requires: Linear Sieve

```
vector<ll> all_factors(ll n) {
1
     vector<ll> fac = {1};
3
     while (n > 1) {
        const ll p = mpf[n];
        vector<ll> cur = {1};
        while (n % p == 0) {
         n /= p;
          cur.push_back(cur.back() * p);
9
        vector<ll> tmp;
       for (auto x : fac)
  for (auto y : cur) tmp.push_back(x * y);
11
13
        tmp.swap(fac);
15
     return fac:
```

4.1.5. Binary GCD

```
1  // returns the gcd of non-negative a, b
    ull bin gcd(ull a, ull b) {
        if (!a || !b) return a + b;
        int s = _builtin_ctzll(a | b);
        a >>= _builtin_ctzll(a);
        while (b) {
            if ((b >>= _builtin_ctzll(b)) < a) swap(a, b);
            b -= a;
        }
        return a << s;
11 }</pre>
```

4.1.6. Extended GCD

```
1
// returns (p, q, g): p * a + q * b == g == gcd(a, b)
// g is not guaranteed to be positive when a < 0 or b < 0
tuple<ll, ll, ll> extgcd(ll a, ll b) {
    ll s = 1, t = 0, u = 0, v = 1;
    while (b) {
        ll q = a / b;
        swap(a -= q * b, b);
        swap(s -= q * t, t);
        swap(u -= q * v, v);
    }
11
    return {s, u, a};
}
```

4.1.7. Chinese Remainder Theorem

Requires: Extended GCD

```
1  // for 0 <= a < m, 0 <= b < n, returns the smallest x >= 0
  // such that x % m == a and x % n == b
3  ll crt(ll a, ll m, ll b, ll n) {
    if (n > m) swap(a, b), swap(m, n);
    auto [x, y, g] = extgcd(m, n);
    assert((a - b) % g == 0); // no solution
7  x = ((b - a) / g * x) % (n / g) * m + a;
    return x < 0 ? x + m / g * n : x;
9 }</pre>
```

4.1.8. Baby-Step Giant-Step

Requires: Mod Struct

```
1  // returns x such that a ^ x = b where x \in [l, r)
  ll bsgs(Mod a, Mod b, ll l = 0, ll r = MOD - 1) {
    int m = sqrt(r - l) + 1, i;
    unordered_map<ll, ll> tb;
    Mod d = (a ^ l) / b;
    for (i = 0, d = (a ^ l) / b; i < m; i++, d *= a)
    if (d == 1) return l + i;
    else tb[(ll)d] = l + i;
    Mod c = Mod(1) / (a ^ m);
    for (i = 0, d = 1; i < m; i++, d *= c)
    if (auto j = tb.find((ll)d); j != tb.end())
        return j->second + i * m;
    return assert(0), -1; // no solution
}
```

4.1.9. Pohlig-Hellman Algorithm

```
Goal: Find an integer x such that g^x=h in an order p^e group.

1. Let x=0 and \gamma=g^{p^{e-1}}.

2. For k=0,1,\ldots,e-1:
Let c=(g^{-x}h)^{p^{e-1-k}}, and compute d such that \gamma^d=c.
Set x=x+p^kd.
```

4.1.10. Pollard's Rho

```
1 | ll f(ll x, ll mod) { return (x * x + 1) % mod; }
   // n should be composite
3
  ll pollard_rho(ll n) {
     if (!(n & 1)) return 2;
     while (1) {
5
       ll y = 2, x = RNG() % (n - 1) + 1, res = 1;
       for (int sz = 2; res == 1; sz *= 2) {
         for (int i = 0; i < sz && res <= 1; i++) {
9
           x = f(x, n);
           res = \_gcd(abs(x - y), n);
         ļ
11
         y = x;
13
       if (res != 0 && res != n) return res;
15
     }
```

4.1.11. Tonelli-Shanks Algorithm

Requires: Mod Struct

```
1 int legendre(Mod a) {
       if (a == 0) return 0;
return (a ^ ((MOD - 1) / 2)) == 1 ? 1 : -1;
 3
 5 Mod sqrt(Mod a) {
       assert(legendre(a) != -1); // no solution
       ll p = MOD, s = p - 1;
if (a == 0) return 0;
       if (p == 2) return 1;
 9
       if (p % 4 == 3) return a ^ ((p + 1) / 4);
       int r, m;
for (r = 0; !(s & 1); r++) s >>= 1;
Mod n = 2;
11
13
       while (legendre(n) != -1) n += 1;
Mod x = a ^ ((s + 1) / 2), b = a ^ s, g = n ^ s;
while (b != 1) {
15
17
          Mod t = b;
          for (m = 0; t != 1; m++) t *= t;
Mod gs = g ^ (1LL << (r - m - 1));
19
          g = gs * gs, x *= gs, b *= g, r = m;
21
       }
       return x:
23
    // to get sqrt(X) modulo p^k, where p is an odd prime: 
// c = x^2 (mod p), c = X^2 (mod p^k), q = p^(k-1)
// X = x^q * c^((p^k-2q+1)/2) (mod p^k)
```

4.1.12. Chinese Sieve

```
1 const ll N = 1000000;
    // f, g, h multiplicative, h = f (dirichlet convolution) g
   ll pre_g(ll n);
ll pre_h(ll n);
 5
    // preprocessed prefix sum of f
   ll pre_f[N];
   // prefix sum of multiplicative function f
ll solve_f(ll n) {
      static unordered_map<ll, ll> m;
      if (n < N) return pre_f[n];</pre>
11
      if (m.count(n)) return m[n];
      ll ans = pre_h(n);
     for (ll l = 2, r; l <= n; l = r + 1) {
    r = n / (n / l);
13
        ans -= (pre_g(r) - pre_g(l - 1)) * djs_f(n / l);
15
17
      return m[n] = ans;
```

4.1.13. Rational Number Binary Search

```
mid.p > N || mid.q > N || dir ^ pred(mid))
    t++;
else len += step;
swap(lo, hi = hi.go(lo, len));
(dir ? L : H) = !!len;
}
return dir ? hi : lo;
}
```

4.1.14. Farey Sequence

```
// returns (e/f), where (a/b, c/d, e/f) are
// three consecutive terms in the order n farey sequence
// to start, call next_farey(n, 0, 1, 1, n)
pll next_farey(ll n, ll a, ll b, ll c, ll d) {
    ll p = (n + b) / d;
    return pll(p * c - a, p * d - b);
}
```

4.2. Combinatorics

4.2.1. Matroid Intersection

This template assumes 2 weighted matroids of the same type, and that removing an element is much more expensive than checking if one can be added. Remember to change the implementation details.

The ground set is $0, 1, \ldots, n-1$, where element i has weight w[i]. For the unweighted version, remove weights and change BF/SPFA to BFS.

```
constexpr int N = 100;
   constexpr int INF = 1e9;
3
                             // represents an independent set
   struct Matroid {
     Matroid(bitset<N>); // initialize from an independent set
                             // if adding will break independence
      bool can_add(int);
      Matroid remove(int); // removing from the set
   auto matroid_intersection(int n, const vector<int> &w) {
     bitset<N> S;
for (int sz = 1; sz <= n; sz++) {
11
        Matroid M1(S), M2(S);
13
        vector<vector<pii>>> e(n + 2);
15
        for (int j = 0; j < n; j++)
17
          if (!S[j]) {
            if (M1.can_add(j)) e[n].emplace_back(j, -w[j]);
19
            if (M2.can_add(j)) e[j].emplace_back(n + 1, 0);
        for (int i = 0; i < n; i++)
21
          if (S[i]) {
            Matroid T1 = M1.remove(i), T2 = M2.remove(i);
            for (int j = 0; j < n; j++)</pre>
              if (!S[j]) {
25
                if (T1.can_add(j)) e[i].emplace_back(j, -w[j]);
                if (T2.can_add(j)) e[j].emplace_back(i, w[i]);
27
          }
29
        vector<pii> dis(n + 2, {INF, 0});
31
        vector<int> prev(n + 2, -1);
33
       // change to SPFA for more speed, if necessary bool upd = 1;
        dis[n] = \{0, 0\};
35
        while (upd) {
37
          upd = 0;
          for (int u = 0; u < n + 2; u++)
39
            for (auto [v, c] : e[u]) {
              pii x(dis[u].first + c, dis[u].second + 1);
if (x < dis[v]) dis[v] = x, prev[v] = u, upd = 1;</pre>
45
        if (dis[n + 1].first < INF)</pre>
          for (int x = prev[n + 1]; x != n; x = prev[x])
            S.flip(x):
47
        else break:
49
        // S is the max-weighted independent set with size sz
51
      return S;
53 }
```

4.2.2. De Brujin Sequence

```
int res[kN], aux[kN], a[kN], sz;
void Rec(int t, int p, int n, int k) {
  if (t > n) {
```

```
if (n \% p == 0)
5
          for (int i = 1; i <= p; ++i) res[sz++] = aux[i];</pre>
     } else {
7
        aux[t] = aux[t - p];
       Rec(t + 1, p, n, k);
for (aux[t] = aux[t - p] + 1; aux[t] < k; ++aux[t])
9
          Rec(t + 1, t, n, k);
11
13
   int DeBruijn(int k, int n) {
     // return cyclic string of length k^n such that every
15
     // string of length n using k character appears as a
     // substring.
     if (k == 1) return res[0] = 0, 1;
     fill(aux, aux + k * n, 0);
     return sz = 0, Rec(1, 1, n, k), sz;
19
```

4.2.3. Multinomial

```
// ways to permute v[i]
ll multinomial(vi &v) {
    ll c = 1, m = v.empty() ? 1 : v[0];
    for (int i = 1; i < v.size(); i++)
        for (int j = 0; i < v[i]; j++) c = c * ++m / (j + 1);
    return c;
}</pre>
```

4.3. Theorems

4.3.1. Kirchhoff's Theorem

Denote L be a $n \times n$ matrix as the Laplacian matrix of graph G, where $L_{ii} = d(i)$, $L_{ij} = -c$ where c is the number of edge (i,j) in G.

- The number of undirected spanning in G is $|\det(\tilde{L}_{11})|$.
- The number of directed spanning tree rooted at r in G is $|\det(\tilde{L}_{rr})|$.

4.3.2. Tutte's Matrix

Let D be a $n \times n$ matrix, where $d_{ij} = x_{ij}$ (x_{ij} is chosen uniformly at random) if i < j and $(i, j) \in E$, otherwise $d_{ij} = -d_{ji}$. $\frac{rank(D)}{2}$ is the maximum matching on G.

4.3.3. Cayley's Formula

• Given a degree sequence d_1, d_2, \ldots, d_n for each labeled vertices, there are

$$\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$$

spanning trees.

• Let $T_{n,k}$ be the number of labeled forests on n vertices with k components, such that vertex $1,2,\ldots,k$ belong to different components. Then $T_{n,k}=kn^{n-k-1}$.

4.3.4. Erdős-Gallai Theorem

A sequence of non-negative integers $d_1 \geq d_2 \geq \ldots \geq d_n$ can be represented as the degree sequence of a finite simple graph on n vertices if and only if $d_1 + d_2 + \ldots + d_n$ is even and

$$\sum_{i=1}^{k} d_i \le k(k-1) + \sum_{i=k+1}^{n} \min(d_i, k)$$

holds for all $1 \le k \le n$.

4.3.5. Burnside's Lemma

Let X be a set and G be a group that acts on X. For $g \in G$, denote by X^g the elements fixed by g:

$$X^g = \{ x \in X \mid gx \in X \}$$

Then

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|.$$

4.3.6. Gram-Schmidt Process

Let $\mathbf{v}_1,\mathbf{v}_2,\dots$ be linearly independent vectors, then the orthogonalized vectors are

$$\mathbf{u}_i = \mathbf{v}_i - \sum_{j=1}^{i-1} \frac{\langle \mathbf{u}_j, \mathbf{v}_k \rangle}{\langle \mathbf{u}_j, \mathbf{u}_j \rangle} \mathbf{u}_j$$

Requires: Mod Struct

Numeric 5.

5.1. Barrett Reduction

```
1 using ull = unsigned long long;
   using uL = __uint128_t;
   // very fast calculation of a % m
3
   struct reduction {
     const ull m, d;
     explicit reduction(ull m) : m(m), d(((uL)1 << 64) / m) {}
     inline ull operator()(ull a) const {
       ull q = (ull)(((uL)d * a) >> 64);
       return (a -= q * m) >= m ? a - m : a;
11 };
```

5.2. Long Long Multiplication

```
1 using ull = unsigned long long;
  using ll = long long;
  using ld = long double;
   // returns a * b % M where a, b < M < 2**63
  ull mult(ull a, ull b, ull M) {
    ll ret = a * b - M * ull(ld(a) * ld(b) / ld(M));
    return ret + M * (ret < 0) - M * (ret >= (ll)M);
```

5.3. Fast Fourier Transform

```
template <tvpename T>
    void fft_(int n, vector<T> &a, vector<T> &rt, bool inv) {
        vector<int> br(n);
        for (int i = 1; i < n; i++) {
  br[i] = (i & 1) ? br[i - 1] + n / 2 : br[i / 2] / 2;
  if (br[i] > i) swap(a[i], a[br[i]]);
        for (int len = 2; len <= n; len *= 2)</pre>
          for (int i = 0; i < n; i += len)

for (int j = 0; j < len / 2; j++) {
    int pos = n / len * (inv ? len - j : j);
    T u = a[i + j], v = a[i + j + len / 2] * rt[pos];
 9
11
                 a[i + j] = u + v, a[i + j + len / 2] = u - v;
13
        if (T minv = T(1) / T(n); inv)
15
           for (T &x : a) x *= minv;
17 }
```

11 Requires: Mod Struct

5

7

9

3

5

11 13

```
void ntt(vector<Mod> &a, bool inv, Mod primitive_root) {
     int n = a.size();
     Mod root = primitive_root ^ (MOD - 1) / n;
     vector<Mod> rt(n + 1, 1);
     for (int i = 0; i < n; i++) rt[i + 1] = rt[i] * root; fft_(n, a, rt, inv);
7
   }
   void fft(vector<complex<double>> &a, bool inv) {
9
     int n = a.size();
     vector<complex<double>> rt(n + 1);
11
     double arg = acos(-1) * 2 / n;
     for (int i = 0; i <= n; i++)
13
       rt[i] = {cos(arg * i), sin(arg * i)};
     fft_(n, a, rt, inv);
15 }
```

5.4. Fast Walsh-Hadamard Transform

Requires: Mod Struct ¹⁵

```
17
    void fwht(vector<Mod> &a, bool inv) {
1
                                                                                             19
       int n = a.size();
       for (int d = 1; d < n; d <<= 1)
for (int m = 0; m < n; m++)
                                                                                             21
            if (!(m & d)) {
               inv ? a[m] -= a[m | d] : a[m] += a[m | d]; // AND
inv ? a[m | d] -= a[m] : a[m | d] += a[m]; // OR
Mod x = a[m], y = a[m | d]; // XOR
                                                                                             23
                                                                                             25
9
               a[m] = x + y, a[m | d] = x - y;
                                                                             // XOR
                                                                                             27
11
       if (Mod iv = Mod(1) / n; inv) // XOR
                                                                                             29
          for (Mod &i : a) i *= iv; // XOR
13 }
```

5.5. Subset Convolution

```
#pragma GCC target("popcnt")
   #include <immintrin.h;</pre>
 3
   void fwht(int n, vector<vector<Mod>> &a, bool inv) {
     for (int h = 0; h < n; h++)
for (int i = 0; i < (1 << n); i++)
           if (!(i & (1 << h)))
             for (int k = 0; k <= n; k++)
inv ? a[i | (1 << h)][k] -= a[i][k]
                    : a[i | (1 << h)][k] += a[i][k];
11
    // c[k] = sum(popcnt(i & j) == sz && i | j == k) a[i] * b[j]
13
   vector<Mod> subset_convolution(int n, int sz
                                         const vector<Mod> &a_,
15
                                         const vector<Mod> &b_) {
      int len = n + sz + 1, N = 1 << n;</pre>
      vector<vector<Mod>> a(1 << n, vector<Mod>(len, 0)), b = a;
17
      for (int i = 0; i < N; i++)
19
        a[i][_mm_popcnt_u64(i)] = a_[i]
        b[i][_mm_popcnt_u64(i)] = b_[i];
      fwht(n, a, 0), fwht(n, b, 0);
for (int i = 0; i < N; i++) {</pre>
21
23
        vector<Mod> tmp(len);
        for (int j = 0; j < len; j++)
for (int k = 0; k <= j; k++)
25
             tmp[j] += a[i][k] * b[i][j - k];
27
        a[i] = tmp;
29
      fwht(n, a, 1);
     vector<Mod> c(N);
for (int i = 0; i < N; i++)</pre>
31
        c[i] = a[i][_mm_popcnt_u64(i) + sz];
33
```

5.6. Linear Recurrences

5.6.1. Berlekamp-Massey Algorithm

```
1 template <typename T>
   vector<T> berlekamp_massey(const vector<T> δs) {
  int n = s.size(), l = 0, m = 1;
      vector<T> r(n), p(n);
      r[0] = p[0] = 1;
      T b = 1, d = 0;
      for (int i = 0; i < n; i++, m++, d = 0) {
         for (int j = 0; j <= l; j++) d += r[j] * s[i - j];
         if ((d /= b) == 0) continue; // change if T is float
        for (int j = m; j < n; j++) r[j] -= d * p[j - m]; if (l * 2 <= i) l = i + 1 - l, b *= d, m = 0, p = t;
13
      return r.resize(l + 1), reverse(r.begin(), r.end()), r;
15 }
```

5.6.2. Linear Recurrence Calculation

```
1 template <typename T> struct lin_rec {
       using poly = vector<T>
       poly mul(poly a, poly b, poly m) {
          int n = m.size();
          poly r(n);
          for (int i = n - 1; i >= 0; i--) {
            r.insert(r.begin(), 0), r.pop_back();
T c = r[n - 1] + a[n - 1] * b[i];
// c /= m[n - 1]; if m is not monic
for (int j = 0; j < n; j++)
r[j] += a[j] * b[i] - c * m[j];</pre>
 9
         return r;
      poly pow(poly p, ll k, poly m) {
  poly r(m.size());
         r[0] = 1;
          for (; k; k >>= 1, p = mul(p, p, m))
           if (k & 1) r = mul(r, p, m);
          return r;
       T calc(poly t, poly r, ll k) {
          int n = r.size();
         poly p(n);
p[1] = 1;
         poly q = pow(p, k, r);
          for (int i = 0; i < n; i++) ans += t[i] * q[i];
          return ans;
31 };
```

39

41

65

67

69

71

73

75

77

79

81

83

85

87

5.7. Matrices

11

13

17

21

5.7.1. Determinant

```
Requires: Mod Struct
                                                                                                      43
    Mod det(vector<vector<Mod>> a) {
        int n = a.size();
 3
        Mod ans = 1;
        for (int i = 0; i < n; i++) {
                                                                                                      47
           int b = i;
           for (int j = i + 1; j < n; j++)
  if (a[j][i] != 0) {</pre>
                                                                                                      51
                 break;
                                                                                                      53
           if (i != b) swap(a[i], a[b]), ans = -ans;
          if (1 != u) Swap(all), ----,

ans *= a[i][i];

if (ans == 0) return 0;

for (int j = i + 1; j < n; j++) {

    Mod v = a[j][i] / a[i][i];
                                                                                                      55
                                                                                                      57
15
                                                                                                      59
              if (v != 0)
                 for (int k = i + 1; k < n; k++)
a[j][k] -= v * a[i][k];</pre>
                                                                                                      61
19
           }
        }
                                                                                                      63
        return ans;
```

```
double det(vector<vector<double>> a) {
      int n = a.size();
3
      double ans = 1:
      for (int i = 0; i < n; i++) {</pre>
        int b = i;
        for (int j = i + 1; j < n; j++)
  if (fabs(a[j][i]) > fabs(a[b][i])) b = j;
        if (i != b) swap(a[i], a[b]), ans = -ans;
9
        ans *= a[i][i];
        if (ans == 0) return 0;
        for (int j = i + 1; j < n; j++) {
  double v = a[j][i] / a[i][i];</pre>
11
13
           if (v != 0)
             for (int k = i + 1; k < n; k++)
                a[j][k] -= v * a[i][k];
15
        }
17
      return ans;
19 }
```

5.7.2. Inverse

```
// Returns rank.
    // Result is stored in A unless singular (rank < n).</pre>
   // For prime powers, repeatedly set
// A^{-1} = A^{-1} (2I - A*A^{-1}) (mod p^k)
   // where A^{-1} starts as the inverse of A mod p,
   // and k is doubled in each step.
   int matInv(vector<vector<double>> &A) {
      int n = sz(A);
 9
      vi col(n);
11
      vector<vector<double>> tmp(n, vector<double>(n));
      rep(i, \theta, n) tmp[i][i] = 1, col[i] = i;
13
      rep(i, 0, n) {
    int r = i, c = i;
    rep(j, i, n)
15
        rep(k, i, n) if (fabs(A[j][k]) > fabs(A[r][c])) r = j,
17
19
        if (fabs(A[r][c]) < 1e-12) return i;</pre>
        A[i].swap(A[r])
        tmp[i].swap(tmp[r])
        rep(j, \theta, n) swap(A[j][i], A[j][c]),
        swap(tmp[j][i], tmp[j][c]);
23
        swap(col[i], col[c]);
25
        double v = A[i][i];
        rep(j, i + 1, n)
           double f = A[j][i] / v;
27
          A[j][i] = 0;
29
          rep(k, i + 1, n) A[j][k] -= f * A[i][k];
          rep(k, \theta, n) tmp[j][k] -= f * tmp[i][k];
31
        rep(j, i + 1, n) A[i][j] /= v;
rep(j, 0, n) tmp[i][j] /= v;
33
        A[i][i] = 1;
35
37
      for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
```

```
double v = A[j][i];
          rep(k, 0, n) tmp[j][k] -= v * tmp[i][k];
     rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] = tmp[i][j];
     return n;
45 }
   int matInv_mod(vector<vector<ll>>> &A) {
     int n = sz(A);
     vi col(n);
     vector<vector<ll>> tmp(n, vector<ll>(n));
     rep(i, 0, n) tmp[i][i] = 1, col[i] = i;
     rep(i, 0, n) {
        int r = i, c = i;
        rep(j, i, n) rep(k, i, n) if (A[j][k]) {
          r = j;
c = k;
          goto found;
        return i:
      found:
        A[i].swap(A[r]);
        tmp[i].swap(tmp[r]);
        rep(j, 0, n) swap(A[j][i], A[j][c]),
swap(tmp[j][i], tmp[j][c]);
        swap(col[i], col[c])
        ll v = modpow(A[i][i], mod - 2);
        rep(j, i + 1, n) {
    ll f = A[j][i] * v % mod;
          A[j][i] = 0;
          rep(k,
                 i + 1, n) A[j][k]
          (A[j][k] - f * A[i][k]) \% mod;
          rep(k, \theta, n) tmp[j][k]
          (tmp[j][k] - f * tmp[i][k]) % mod;
       rep(j, i + 1, n) A[i][j] = A[i][j] * v \% mod;
       rep(j, 0, n) tmp[i][j] = tmp[i][j] * v % mod;
        A[i][i] = 1;
     for (int i = n - 1; i > 0; --i) rep(j, 0, i) {
          ll v = A[j][i];
          rep(k, 0, n) tmp[j][k] =
          (tmp[j][k] - v * tmp[i][k]) % mod;
     rep(i, 0, n) rep(j, 0, n) A[col[i]][col[j]] =
tmp[i][j] % mod + (tmp[i][j] < 0 ? mod : 0);</pre>
     return n;
```

5.7.3. Characteristic Polynomial

```
// calculate det(a - xI)
    \textbf{template} \hspace{0.1cm} < \hspace{-0.1cm} \textbf{typename} \hspace{0.1cm} \hspace{0.1cm} \hspace{0.1cm} \hspace{0.1cm} \textbf{T} \hspace{-0.1cm} > \hspace{-0.1cm}
 3
    vector<T> CharacteristicPolynomial(vector<vector<T>> a) {
       int N = a.size();
       for (int j = 0; j < N - 2; j++) {
  for (int i = j + 1; i < N; i++) {
    if (a[i][j] != 0) {</pre>
 7
               swap(a[j + 1], a[i]);
for (int k = 0; k < N; k++)
 9
                  swap(a[k][j + 1], a[k][i]);
11
               break;
             }
13
         15
17
19
                T coe = inv * a[i][j];
                for (int l = j; l < N; l++)
21
                  a[i][l] -= coe * a[j + 1][l];
                for (int k = 0; k < N; k++)
                  a[k][j + 1] += coe * a[k][i];
23
            }
25
          }
27
       vector<vector<T>> p(N + 1);
       p[0] = \{T(1)\};
29
       for (int i = 1; i <= N; i++) {
          p[i].resize(i + 1);
31
          for (int j = 0; j < i; j++) {
  p[i][j + 1] -= p[i - 1][j];</pre>
33
             p[i][j] += p[i - 1][j] * a[i - 1][i - 1];
```

```
35
           T x = 1;
           for (int m = 1; m < i; m++) {
   x *= -a[i - m][i - m - 1];</pre>
37
              T coe = x * a[i - m - 1][i - 1];
for (int j = 0; j < i - m; j++)
39
41
                 p[i][j] += coe * p[i - m - 1][j];
           }
43
        return p[N];
45 }
```

5.7.4. Solve Linear Equation

```
27
   typedef vector<double> vd;
   const double eps = 1e-12;
 3
      solves for x: A * x = b
                                                                             31
   int solveLinear(vector<vd> &A, vd &b, vd &x) {
      int n = sz(A), m = sz(x), rank = 0, br, bc;
                                                                             33
      if (n) assert(sz(A[0]) == m);
      vi col(m);
                                                                             35
 9
      iota(all(col), 0);
                                                                             37
11
      rep(i, 0, n) {
        double v, bv = 0;
                                                                             39
13
        rep(r, i, n) rep(c, i, m) if ((v = fabs(A[r][c])) > bv)
                                                                             41
15
        bc = c, bv = v;
        if (bv <= eps) {
                                                                             43
          rep(j, i, n) if (fabs(b[j]) > eps) return -1;
17
                                                                             45
19
        swap(A[i], A[br]);
                                                                             47
        swap(b[i], b[br]);
21
        swap(col[i], col[bc]);
rep(j, 0, n) swap(A[j][i], A[j][bc]);
bv = 1 / A[i][i];
                                                                             49
23
                                                                             51
        rep(j, i + 1, n) {
    double fac = A[j][i] * bv;
    b[j] -= fac * b[i];
25
                                                                             53
          rep(k, i + 1, m) A[j][k] -= fac * A[i][k];
29
        rank++;
                                                                             57
      }
31
                                                                             59
33
      x.assign(m, 0);
      for (int i = rank; i--;) {
                                                                             61
        b[i] /= A[i][i]
35
        x[col[i]] = b[i];
                                                                             63
37
        rep(j, 0, i) b[j] -= A[j][i] * b[i];
                                                                             65
      return rank; // (multiple solutions if rank < m)</pre>
39
   }
                                                                             67
```

5.8. Polynomial Interpolation

```
// returns a, such that a[0]x^0 + a[1]x^1 + a[2]x^2 + ...
                                                                              71
   // passes through the given points
   typedef vector<double> vd;
3
                                                                              73
   vd interpolate(vd x, vd y, int n) {
      vd res(n), temp(n);
                                                                              75
      rep(k, 0, n - 1) rep(i, k + 1, n) y[i] = (y[i] - y[k]) / (x[i] - x[k]);
                                                                              77
      double last = 0;
9
      temp[0] = 1;
                                                                              79
      rep(k, 0, n) rep(i, 0, n) {
        res[i] += y[k] * temp[i];
11
                                                                              81
        swap(last, temp[i]);
temp[i] -= last * x[k];
13
                                                                              83
      return res;
                                                                              85
                                                                              87
```

69

5.9. Simplex Algorithm

```
89
1 // Two-phase simplex algorithm for solving linear programs
                                                                         91
     of the form
3 //
                                                                         93
   //
           maximize
                         c^T x
5 //
                         Ax <= b
           subject to
   11
                                                                         95
                         x >= 0
// INPUT: A -- an m x n matrix 9 // h -- an m x
7 //
                                                                         97
              b -- an m-dimensional vector
              c -- an n-dimensional vector
                                                                         99
11 //
              \boldsymbol{x} -- a vector where the optimal solution will be
              stored
```

```
13 | //
    // OUTPUT: value of the optimal solution (infinity if
15 // unbounded
                 above, nan if infeasible)
17 //
    // To use this code, create an LPSolver object with A, b,
19 // and c as arguments. Then, call Solve(x).
   typedef long double ld;
21
    typedef vector<ld> vd;
23
   typedef vector<vd> vvd;
    typedef vector<int> vi;
    const ld EPS = 1e-9;
    struct LPSolver {
      int m, n;
      vi B. N:
      vvd D;
      LPSolver(const vvd &A, const vd &b, const vd &c)
           : m(b.size()), n(c.size()), N(n + 1), B(m),
         D(m + 2, vd(n + 2)) {
for (int i = 0; i < m; i++)
         for (int j = 0; j < n; j++) D[i][j] = A[i][j];
for (int i = 0; i < m; i++) {
           B[i] = n + i;
D[i][n] = -1;
           D[i][n + 1] = b[i];
         for (int j = 0; j < n; j++) {
           N[j] = j;
           D[m][j] = -c[j];
        N[n] = -1;
        D[m + 1][n] = 1;
      void Pivot(int r, int s) {
  double inv = 1.0 / D[r][s];
         for (int i = 0; i < m + 2; i++)
           if (i != r)
              for (int j = 0; j < n + 2; j++)
        if (j != s) D[i][j] -= D[r][j] * D[i][s] * inv;
for (int j = 0; j < n + 2; j++)
    if (j != s) D[r][j] *= inv;
for (int i = 0; j < n + 2; j++)</pre>
         for (int i = 0; i < m + 2; i++)
           if (i != r) D[i][s] *= -inv;
         D[r][s] = inv;
         swap(B[r], N[s]);
      bool Simplex(int phase) {
         int x = phase == 1 ? m + 1 : m;
         while (true) {
           int s = -1;
           for (int j = 0; j <= n; j++) {
  if (phase == 2 && N[j] == -1) continue;
  if (s == -1 || D[x][j] < D[x][s] ||</pre>
                   D[x][j] == D[x][s] && N[j] < N[s])
           if (D[x][s] > -EPS) return true;
           int r = -1;
for (int i = 0; i < m; i++) {</pre>
              if (D[i][s] < EPS) continue;</pre>
              if (r == -1 ||
                   D[i][n + 1] / D[i][s] < D[r][n + 1] / D[r][s] |
                   (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r][s]) &&
                   B[i] < B[r]
                r = i;
           if (r == -1) return false;
           Pivot(r, s);
      ld Solve(vd &x) {
         int r = 0;
         for (int i = 1; i < m; i++)
  if (D[i][n + 1] < D[r][n + 1]) r = i;</pre>
         if (D[r][n + 1] < -EPS) {
           Pivot(r, n);
           if (!Simplex(1) || D[m + 1][n + 1] < -EPS)</pre>
           return -numeric_limits<ld>::infinity();
for (int i = 0; i < m; i++)</pre>
              if (B[i] == -1) {
```

```
int s = -1;
101
                   for (int j = 0; j <= n; j++)
if (s == -1 || D[i][j] < D[i][s] ||
103
                           D[i][j] == D[i][s] \delta\delta N[j] < N[s])
105
                   Pivot(i, s);
107
           if (!Simplex(2)) return numeric_limits<ld>::infinity();
109
           for (int i = 0; i < m; i++)
             if (B[i] < n) \times [B[i]] = D[i][n + 1];
           return D[m][n + 1];
113
115 };
117 int main() {
119
        const int m = 4:
        const int n = 3;
        ld _A[m][n] = {
121
        ld _-lim][1] - [-1, -5, 0], {1, 5, 1}, {-1, -5, -1}};
ld _b[m] = {10, -4, 5, -5};
ld _c[n] = {1, -1, 0};
123
125
        vvd A(m);
        vd b(_b, _b + m);
127
        vd c(_c, _c + n);
for (int i = θ; i < m; i++) A[i] = vd(_A[i], _A[i] + n);
129
        LPSolver solver(A, b, c);
        vd x;
133
        ld value = solver.Solve(x);
        cerr << "VALUE: " << value << endl; // VALUE: 1.29032
cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
for (size_t i = 0; i < x.size(); i++) cerr << " " << x[i];</pre>
135
137
        cerr << endl;
139
        return 0;
```

6. Geometry

6.1. Point

```
template <typename T> struct P {
      T x, y;
P(T x = 0, T y = 0) : x(x), y(y) {}
 3
       bool operator<(const P &p) const {
         return tie(x, y) < tie(p.x, p.y);</pre>
      bool operator==(const P &p) const {
         return tie(x, y) == tie(p.x, p.y);
9
       P operator-() const { return {-x, -y}; }
      P operator (P p) const { return {x + p.x, y + p.y}; } P operator (P p) const { return {x - p.x, y - p.y}; } P operator (T d) const { return {x + d, y + d}; } P operator (T d) const { return {x + d, y + d}; } P operator (T d) const { return {x / d, y / d}; }
11
13
15
      T dist2() const { return x * x + y * y;
       double len() const { return sqrt(dist2()); }
      P unit() const { return *this / len(); }
17
       friend T dot(P a, P b) { return a.x * b.x + a.y * b.y; }
       friend T cross(P a, P b) { return a.x * b.y - a.y * b.x; }
19
       friend T cross(P a, P b, P o) {
         return cross(a - o, b - o);
21
23 };
    using pt = P<ll>;
```

6.1.1. Quarternion

```
constexpr double PI = 3.141592653589793;
   constexpr double EPS = 1e-7;
   struct Q {
     using T = double;
     T x, y, z, r;
     Q(T r = 0): x(0), y(0), z(0), r(r) {}
Q(T x, T y, T z, T r = 0): x(x), y(y), z(z), r(r) {}
      friend bool operator==(const Q &a, const Q &b) {
       return (a - b).abs2() <= EPS;</pre>
9
     friend bool operator!=(const Q &a, const Q &b) {
11
       return !(a == b):
13
     Q operator-() { return Q(-x, -y, -z, -r); }
15
     Q operator+(const Q &b) const {
```

```
return Q(x + b.x, y + b.y, z + b.z, r + b.r);
17
     Q operator-(const Q &b) const {
19
       return Q(x - b.x, y - b.y, z - b.z, r - b.r);
21
     Q operator*(const T &t) const {
       return Q(x * t, y * t, z * t, r * t);
23
     Q operator*(const Q &b) const {
       return Q(r * b.x + x * b.r + y * b.z - z * b.y,
r * b.y - x * b.z + y * b.r + z * b.x,
25
27
                  r * b.z + x * b.y - y * b.x + z * b.r,
                  r * b.r - x * b.x - y * b.y - z * b.z);
29
     Q operator/(const Q &b) const { return *this * b.inv(); }
     T abs2() const { return r * r + x * x + y * y + z * z; }
31
     T len() const { return sqrt(abs2()); }
     Q conj() const { return Q(-x, -y, -z, r); }
Q unit() const { return *this * (1.0 / len()); }
33
     Q inv() const { return conj() * (1.0 / abs2()); } friend T dot(Q a, Q b) {
35
        return a.x * b.x + a.y * b.y + a.z * b.z;
37
39
     friend Q cross(Q a, Q b) {
       return Q(a.y * b.z - a.z * b.y, a.z * b.x - a.x * b.z,
a.x * b.y - a.y * b.x);
41
43
     friend Q rotation_around(Q axis, T angle) {
       return axis.unit() * sin(angle / 2) + cos(angle / 2);
45
     Q rotated_around(Q axis, T angle) {
       Q u = rotation_around(axis, angle);
47
       return u * *this / u:
49
     friend Q rotation_between(Q a, Q b) {
51
        a = a.unit(), b = b.unit();
        if (a == -b) {
53
          // degenerate case
          Q ortho = abs(a.y) > EPS ? cross(a, Q(1, 0, 0))
55
                                       : cross(a, Q(0, 1, 0));
          return rotation_around(ortho, PI);
57
        return (a * (a + b)).conj();
59
     }
   };
```

6.1.2. Spherical Coordinates

```
1 struct car p {
     double x, y, z;
 3
   };
   struct sph_p {
 5
     double r, theta, phi;
 7
   sph_p conv(car_p p) {
9
     double r = sqrt(p.x * p.x + p.y * p.y + p.z * p.z);
      double theta = asin(p.y / r);
11
      double phi = atan2(p.y, p.x);
     return {r, theta, phi}
13 }
   car_p conv(sph_p p) {
     double x = p.r * cos(p.theta) * sin(p.phi);
double y = p.r * cos(p.theta) * cos(p.phi);
15
      double z = p.r * sin(p.theta);
      return {x, y, z};
19 }
```

6.2. Segments

```
// for non-collinear ABCD, if segments AB and CD intersect
    bool intersects(pt a, pt b, pt c, pt d) {
  if (cross(b, c, a) * cross(b, d, a) > 0) return false;
  if (cross(d, a, c) * cross(d, b, c) > 0) return false;
 3
 5
      return true:
 7
    // the intersection point of lines AB and CD
    pt intersect(pt a, pt b, pt c, pt d) {
9
      auto x = cross(b, c, a), y = cross(b, d, a);
       if (x == y) {
         // if(abs(x, y) < 1e-8) {
// is parallel
11
13
      } else {
         return d * (x / (x - y)) - c * (y / (x - y));
15
      }
```

6.3. Convex Hull

```
// returns a convex hull in counterclockwise order
// for a non-strict one, change cross >= to >
vector<pt> convex_hull(vector<pt> p) {
    sort(ALL(p));
    if (p[0] == p.back()) return {p[0]};
    int n = p.size(), t = 0;
    vector<pt> h(n + 1);
    for (int _ = 2, s = 0; _--; s = --t, reverse(ALL(p)))
    for (pt i : p) {
        while (t > s + 1 && cross(i, h[t - 1], h[t - 2]) >= 0)
        t--;
        h[t++] = i;
    }
    return h.resize(t), h;
}
```

```
6.3.1. 3D Hull
   typedef Point3D<double> P3;
       void ins(int x) { (a == -1 ? a : b) = x; }
       void rem(int x) { (a == x ? a : b) = -1; }
       int cnt() { return (a != -1) + (b != -1); }
      int a. b:
    };
 9
    struct F {
11
      int a, b, c;
    };
13
    // collinear points will kill it, please remove before use // skip between -snip- comments if no 4 coplanar points
15
    vector<F> hull3d(vector<P3> A) {
      int n = A.size(), t2 = 2, t3 = 3;
17
      vector<vector<PR>>> E(n, vector<PR>(n, {-1, -1}));
19
      vector<F> FS;
21
       for (int i = 2; i < n; i++) // -snip-
         for (int j = i + 1; j < n; j++) {
    ll v = cross(A[0], A[1], A[i]).dot(A[j] - A[0]);</pre>
23
            if (v != 0) {
              if (v < 0) swap(i, j);
swap(A[2], A[t2 = i]), swap(A[3], A[t3 = j]);
25
27
              goto ok;
29
         }
      assert(!"all coplanar");
31
    ok:; // -snip-
    #define E(x, y) E[min(f.x, f.y)][max(f.x, f.y)]
#define C(a, b)
33
35
      if (E(a, b).cnt() != 2) mf(f.a, f.b, i);
      auto mf = [8](int i, int j, int k) {
  F f = {i, j, k};
37
39
         E(a, b).ins(k);
         E(a, c).ins(j);
         E(b, c).ins(i)
41
         FS.push_back(f);
43
      auto in = [8](int i, int j, int k, int l) {
  P3 a = cross(A[i], A[j], A[l]),
  b = cross(A[j], A[k], A[l]),
  c = cross(A[k], A[i], A[i]);
}
45
47
         c = cross(A[k], A[i], A[l]);
return a.dot(b) > 0 && b.dot(c) > 0;
49
      mf(0, 2, 1), mf(0, 1, 3), mf(1, 2, 3), mf(0, 3, 2);
51
       for (int i = 4; i < n; i++) {
  for (int j = 0; j < FS.size(); j++) {
    F f = FS[j];
}</pre>
53
55
            ll d =
            cross(A[f.a], A[f.b], A[f.c]).dot(A[i] - A[f.a]);
57
            if (d > 0 || (d == 0 && in(f.a, f.b, f.c, i))) {
              E(a, b).rem(f.c);
59
              E(a, c).rem(f.b);
61
              E(b, c).rem(f.a);
              swap(FS[j--], FS.back());
63
              FS.pop_back();
65
         for (int j = 0, s = FS.size(); j < s; j++) {</pre>
            F f = FS[j];
67
            C(c, b);
69
            C(b, a);
```

```
C(a, c):
         }
71
      }
73
      vector<int> idx(n), ri(n); // -snip-
iota(idx.begin(), idx.end(), 0);
75
      swap(idx[t3], idx[3]), swap(idx[t2], idx[2]);
      for (int i = 0; i < n; i++) ri[idx[i]] = i;
for (auto 8[a, b, c] : FS)
77
79
        a = ri[a], b = ri[b], c = ri[c]; // -snip-
      return FS;
81
    #undef E
83 #undef C
```

6.4. Angular Sort

```
auto angle_cmp = [](const pt &a, const pt &b) {
    auto btm = [](const pt &a) {
        return a.y < 0 || (a.y == 0 && a.x < 0);
        };
    return make_tuple(btm(a), a.y * b.x, abs2(a)) <
            make_tuple(btm(b), a.x * b.y, abs2(b));
};

void angular_sort(vector<pt> &p) {
    sort(p.begin(), p.end(), angle_cmp);
}
```

6.5. Convex Hull Tangent

```
1 // before calling, do
   // int top = max_element(c.begin(), c.end()) -
   // c.begin();
   // c.push_back(c[0]), c.push_back(c[1]);
   pt left_tangent(const vector<pt> &c, int top, pt p) {
     int n = c.size() - 2;
     int ans = -1;
     do {
9
        if (cross(p, c[n], c[n + 1]) >= 0 &&
            (cross(p, c[top + 1], c[n]) > 0 | |
             cross(p, c[top], c[top + 1]) < 0))
11
          break;
        int l = top + 1, r = n + 1;
13
        while (l < r - 1) {
          int m = (l + r) / 2;
15
          if (cross(p, c[m - 1], c[m]) > 0 \delta\delta
17
              cross(p, c[top + 1], c[m]) > 0)
            l = m;
19
          else r = m;
21
        ans = l;
     } while (false);
23
        if (cross(p, c[top], c[top + 1]) >= 0 &&
            (cross(p, c[1], c[top]) > 0 ||
cross(p, c[0], c[1]) < 0))
25
        break;
int l = 1, r = top + 1;
27
        while (l < r - 1) {
29
          int m = (l + r) / 2;
if (cross(p, c[m - 1], c[m]) > 0 &&
31
              cross(p, c[1], c[m]) > 0)
            l = m;
33
          else r = m;
35
        ans = l;
     } while (false);
37
     return c[ans] - p;
39 }
```

6.6. Convex Polygon Minkowski Sum

```
1 // O(n) convex polygon minkowski sum
    // must be sorted and counterclockwise
   vector<pt> minkowski_sum(vector<pt> p, vector<pt> q) {
      auto diff = [](vector<pt> &c) {
        auto rcmp = [](pt a, pt b) {
  return pt{a.y, a.x} < pt{b.y, b.x};</pre>
 5
        rotate(c.begin(), min_element(ALL(c), rcmp), c.end());
 9
        c.push_back(c[0]);
        vector<pt> ret;
for (int i = 1; i < c.size(); i++)</pre>
11
          ret.push_back(c[i] - c[i - 1]);
13
        return ret;
     }:
     auto dp = diff(p), dq = diff(q);
15
     pt cur = p[0] + q[0];
```

```
17
     vector<pt> d(dp.size() + dq.size()), ret = {cur};
     // include angle_cmp from angular-sort.cpp
     merge(ALL(dp), ALL(dq), d.begin(), angle_cmp);
19
     // optional: make ret strictly convex (UB if degenerate)
     int now = 0;
21
     for (int i = 1; i < d.size(); i++) {</pre>
       if (cross(d[i], d[now]) == 0) d[now] = d[now] + d[i];
23
       else d[++now] = d[i];
25
     d.resize(now + 1);
     // end optional part
27
     for (pt v : d) ret.push_back(cur = cur + v);
29
     return ret.pop_back(), ret;
```

17

41

6.7. Point In Polygon

```
bool on_segment(pt a, pt b, pt p) {
  return cross(a, b, p) == θ δδ dot((p - a), (p - b)) <= θ;</pre>
3
     // p can be any polygon, but this is O(n)
    bool inside(const vector<pt> &p, pt a) {
      int cnt = 0, n = p.size();
for (int i = 0; i < n; i++) {</pre>
         pt l = p[i], r = p[(i + 1) \% n];
          // change to return 0; for strict version
g
         if (on_segment(l, r, a)) return 1; cnt \hat{} (a.y < l.y) - (a.y < r.y)) * cross(l, r, a) > 0;
11
13
      return cnt;
```

6.7.1. Convex Version

```
// no preprocessing version
    // p must be a strict convex hull, counterclockwise
// if point is inside or on border
    bool is_inside(const vector<pt> δc, pt p) {
      int n = c.size(), l = 1, r = n - 1;
if (cross(c[0], c[1], p) < 0) return false;
if (cross(c[n - 1], c[0], p) < 0) return false;</pre>
                                                                                  47
                                                                                  49
      while (l < r - 1) {
         int m = (l + r) / 2
 q
                                                                                  51
         T a = cross(c[\theta], c[m], p);
         if (a > 0) l = m;
                                                                                  53
         else if (a < 0) r = m;
13
         else return dot(c[0] - p, c[m] - p) <= 0;
                                                                                  55
      if (l == r) return dot(c[0] - p, c[l] - p) <= 0;</pre>
15
                                                                                  57
      else return cross(c[l], c[r], p) \Rightarrow 0;
17
                                                                                  59
    // with preprocessing version
19
                                                                                  61
    vector<pt> vecs;
21
   pt center;
                                                                                  63
       p must be a strict convex hull, counterclockwise
    // BEWARE OF OVERFLOWS!!
                                                                                  65
    void preprocess(vector<pt> p) {
      for (auto \delta v : p) v = v * 3;
center = p[0] + p[1] + p[2];
25
                                                                                  67
      center.x /= 3, center.y /= 3;
for (auto &v : p) v = v - center;
                                                                                  69
29
      vecs = (angular_sort(p), p);
                                                                                  71
   bool intersect_strict(pt a, pt b, pt c, pt d) {
31
                                                                                  73
      if (cross(b, c, a) * cross(b, d, a) > 0) return false;
      if (cross(d, a, c) * cross(d, b, c) >= 0) return false;
33
                                                                                  75
      return true;
35
                                                                                  77
     // if point is inside or on border
37
   bool query(pt p) {
                                                                                  79
      p = p * 3 - center;
      auto pr = upper_bound(ALL(vecs), p, angle_cmp);
39
                                                                                  81
      if (pr == vecs.end()) pr = vecs.begin();
auto pl = (pr == vecs.begin()) ? vecs.back() : *(pr - 1);
41
                                                                                  83
      return !intersect_strict({0, 0}, p, pl, *pr);
43 }
```

6.7.2. Offline Multiple Points Version

Requires: Point, GNU PBDS

87

```
89
  using Double =
                       float128:
  using Point = pt<Double, Double>;
                                                                               91
                                                                               93
  vector<Point> poly;
vector<Point> query;
                                                                               95
7 vector<int> ans;
```

```
9 struct Segment {
     Point a, b;
11
     int id;
   vector<Segment> segs;
15
   Double Xnow;
   inline Double get_y(const Segment &u, Double xnow = Xnow) {
     const Point &a = u.a;
     const Point &b = u.b;
19
     return (a.y * (b.x - xnow) + b.y * (xnow - a.x)) /
             (b.x - a.x);
21
   bool operator<(Segment u, Segment v) {</pre>
23
     Double yu = get_y(u);
     Double yv = get_y(v);
if (yu != yv) return yu < yv;
     return u.id < v.id;
27
   ordered_map<Segment> st;
29
   struct Event {
     int type; // +1 insert seg, -1 remove seg, 0 query
31
     Double x, y;
33
     int id;
35
   bool operator<(Event a, Event b) {</pre>
     if (a.x != b.x) return a.x < b.x;
     if (a.type != b.type) return a.type < b.type;</pre>
     return a.y < b.y;</pre>
30
   }
   vector<Event> events;
   void solve() {
     set<Double> xs;
     set<Point> ps;
     for (int i = 0; i < n; i++) {
       xs.insert(poly[i].x);
       ps.insert(poly[i]);
     for (int i = 0; i < n; i++) {
       Segment s\{poly[i], poly[(i + 1) \% n], i\};
        if (s.a.x > s.b.x ||
            (s.a.x == s.b.x \&\& s.a.y > s.b.y)) {
          swap(s.a, s.b);
        segs.push_back(s);
        if (s.a.x != s.b.x) {
          events.push_back(\{+1, s.a.x + 0.2, s.a.y, i\});
          events.push_back(\{-1, s.b.x - 0.2, s.b.y, i\});
       }
     for (int i = 0; i < m; i++) {
       events.push_back({0, query[i].x, query[i].y, i});
     sort(events.begin(), events.end());
     int cnt = 0;
     for (Event e : events) {
        int i = e.id;
        Xnow = e.x;
        if (e.type == 0) {
          Double x = e.x;
          Double y = e.y';
Segment tmp = \{\{x - 1, y\}, \{x + 1, y\}, -1\};
auto it = st.lower_bound(tmp);
          if (ps.count(query[i]) > 0) {
            ans[i] = 0;
          } else if (xs.count(x) > 0) {
            ans[i] =
          } else if (it != st.end() 88
                      get_y(*it) == get_y(tmp)) {
            ans[i] = \overline{0};
          } else if (it != st.begin() &&
                      get_y(*prev(it)) == get_y(tmp)) {
            ans[i] = 0:
          } else {
            int rk = st.order_of_key(tmp);
            if (rk % 2 == 1) {
              ans[i] = 1;
            } else {
              ans[i] = -1;
            }
       } else if (e.type == 1) {
          st.insert(segs[i]);
```

25

27

29

31

51

53

57

87

```
assert((int)st.size() == ++cnt);
97
        } else if (e.type == -1) {
          st.erase(segs[i]);
99
          assert((int)st.size() == --cnt);
101
```

6.8. Closest Pair

```
vector<pll> p; // sort by x first!
   bool cmpy(const pll &a, const pll &b) const {
                                                                                 37
      return a.y < b.y;</pre>
                                                                                 39
   ll sq(ll x) { return x * x; }
// returns (minimum dist)^2 in [l, r)
                                                                                 41
   ll solve(int l, int r) {
      if (r - l <= 1) return 1e18;</pre>
      int m = (l + r) / 2;
ll mid = p[m].x, d = min(solve(l, m), solve(m, r));
                                                                                 43
                                                                                 45
      auto pb = p.begin();
11
      inplace_merge(pb + l, pb + m, pb + r, cmpy);
                                                                                 47
      vector<pll> s;
for (int i = l; i < r; i++)</pre>
13
                                                                                 49
        if (sq(p[i].x - mid) < d) s.push_back(p[i]);</pre>
15
      for (int i = 0; i < s.size(); i++)</pre>
        for (int j = i + 1;
j < s.size() && sq(s[j].y - s[i].y) < d; j++)
17
           d = min(d, dis(s[i], s[j]));
19
      return d;
                                                                                 55
21 }
```

6.9. Minimum Enclosing Circle

```
59
   typedef Point<double> P;
   double ccRadius(const P &A, const P &B, const P &C) {
                                                                         61
      return (B - A).dist() * (C - B).dist() * (A - C).dist() /
             abs((B - A).cross(C - A)) / 2;
                                                                         63
5
     ccCenter(const P &A, const P &B, const P &C) {
                                                                         65
     P b = C - A, c = B - A;
     return A + (b * c.dist2() - c * b.dist2()).perp() /
                                                                         67
9
                  b.cross(c) / 2;
                                                                         69
   pair<P, double> mec(vector<P> ps) {
11
      shuffle(all(ps), mt19937(time(0)));
                                                                         71
13
     P o = ps[0];
      double r = 0, EPS = 1 + 1e-8;
                                                                         73
      rep(i, \theta, sz(ps)) if ((o - ps[i]).dist() > r * EPS) {
15
        o = ps[i], r = 0;
rep(j, 0, i) if ((o - ps[j]).dist() > r * EPS) {
                                                                         75
17
          o = (ps[i] + ps[j]) / 2;
                                                                         77
          r = (o - ps[i]).dist();
19
          rep(k, 0, j) if ((o - ps[k]).dist() > r * EPS) {
  o = ccCenter(ps[i], ps[j], ps[k]);
                                                                         79
            r = (o - ps[i]).dist();
                                                                         81
23
       }
                                                                         83
25
     return {o, r};
                                                                         85
27 }
```

6.10. Delaunay Triangulation

```
89
   // O(n * log(n)), T_large must be able to hold O(T^4) (can // be long long if coord <= 2e4)
1
                                                                                91
   struct quad_edge {
  int o = -1; // origin of the arc
                                                                                93
      quad_edge *onext, *rot;
bool mark = false;
                                                                                95
      quad_edge() {}
      quad_edge(int o) : o(o) {}
                                                                                97
      int d() { return sym()->o; } // destination of the arc
quad_edge *sym() { return rot->rot; }
9
                                                                                99
      quad_edge *oprev() { return rot->onext->rot; }
quad_edge *lnext() { return sym()->oprev(); }
11
                                                                               101
      static quad_edge *make_sphere(int a, int b) {
        103
15
         new quad_edge{}}};
17
        for (auto i = 0; i < 4; ++i)
           q[i]->onext = q[-i \ \delta \ 3], \ q[i]->rot = q[i + 1 \ \delta \ 3];
                                                                               107
19
        return q[0];
                                                                               109
21
      static void splice(quad_edge *a, quad_edge *b) {
        swap(a->onext->rot->onext, b->onext->rot->onext);
                                                                               111
23
        swap(a->onext, b->onext);
```

```
static quad_edge *connect(quad_edge *a, quad_edge *b) {
  quad_edge *q = make_sphere(a->d(), b->o);
    splice(q, a->lnext()), splice(q->sym(), b);
    return q;
  }
template <class T, class T_large, class F1, class F2>
bool delaunay_triangulation(const vector<point<T>>> &a,
                                F1 process_outer_face,
                                F2 process_triangles) {
  vector<int> ind(a.size());
  iota(ind.begin(), ind.end(), 0);
sort(ind.begin(), ind.end(),
        [8](int i, int j) { return a[i] < a[j]; });
  ind.erase(
  unique(ind.begin(), ind.end(),
    [8](int i, int j) { return a[i] == a[j]; }),
  ind.end());
  int n = (int)ind.size():
  if (n < 2) return {};</pre>
  auto circular = [8](point<T> p, point<T> a, point<T> b,
                        point<T> c) {
    a -= p, b -= p, c -= p;
    return ((T_large)a.squared_norm() * (b ^ c) + (T_large)b.squared_norm() * (c ^ a) + (T_large)c.squared_norm() * (a ^ b)) *
            (doubled\_signed\_area(a, b, c) > 0 ? 1 : -1) >
  auto recurse = [8](auto self, int l,
                        int r) -> array<quad_edge *, 2> {
    if (r - l <= 3) {
      quad edge *p =
      quad_edge::make_sphere(ind[l], ind[l + 1]);
      if (r - l == 2) return {p, p->sym()};
      quad_edge *q =
      quad_edge::make_sphere(ind[l + 1], ind[l + 2]);
      quad_edge::splice(p->sym(), q);
      auto side = doubled_signed_area(
      a[ind[l]], a[ind[l + 1]], a[ind[l + 2]]);
      quad_edge *c = side ? quad_edge::connect(q, p) : NULL;
      return {side < 0 ? c->sym() :
               side < 0 ? c : q->sym()};
    int m = l + (r - l >> 1);
    auto [ra, A] = self(self, l, m);
auto [B, rb] = self(self, m, r);
    while (
    doubled_signed_area(a[B->o], a[A->d()], a[A->o]) < 0 &&
    (A = A->lnext()) \mid \mid
    doubled_signed_area(a[A->o], a[B->d()], a[B->o]) > 0 \delta\delta
    (B = B->sym()->onext))
    quad_edge *base = quad_edge::connect(B->sym(), A);
    if (A->o == ra->o) ra = base->sym();
    if (B->o == rb->o) rb = base;
#define valid(e)
  (doubled_signed_area(a[e->d()], a[base->d()],
                          a[base->o]) > 0)
#define DEL(e, init, dir)
  quad_edge *e = init->dir;
  if (valid(e))
    while (circular(a[e->dir->d()], a[base->d()],
                      a[base->o], a[e->d()])) {
      quad_edge *t = e->dir;
      quad_edge::splice(e, e->oprev());
      quad_edge::splice(e->sym(), e->sym()->oprev());
      delete e->rot->rot->rot;
      delete e->rot->rot;
      delete e->rot;
      delete e;
      e = t;
    while (true) {
      DEL(LC, base->sym(), onext);
      DEL(RC, base, oprev());
      if (!valid(LC) && !valid(RC)) break;
      if (!valid(LC) ||
           valid(RC) ชิชิ circular(a[RC->d()], a[RC->o],
a[LC->d()], a[LC->o]))
         base = quad_edge::connect(RC, base->sym());
      else
         base = quad_edge::connect(base->sym(), LC->sym());
    ļ
    return {ra, rb};
  };
  auto e = recurse(recurse, 0, n)[0];
  vector<quad_edge *> q = {e}, rem;
```

```
113
      while (doubled signed area(a[e->onext->d()], a[e->d()].
                                    a[e->o]) < 0)
         e = e->onext;
115
       vector<int> face;
117
       face.reserve(n);
      bool colinear = false;
119
    #define ADD
121
         quad_edge *c = e;
         face.clear();
123
         do {
           c->mark = true;
           face.push_back(c->o);
125
           q.push_back(c->sym());
           rem.push_back(c);
           c = c->lnext();
129
         } while (c != e);
131
      ADD:
      process_outer_face(face);
for (auto qi = 0; qi < (int)q.size(); ++qi) {</pre>
133
         if (!(e = q[qi])->mark) {
135
           ADD:
           colinear = false;
           process_triangles(face[0], face[1], face[2]);
137
139
      for (auto e : rem) delete e->rot, delete e;
141
       return !colinear;
```

6.10.1. Quadratic Time Version

Requires: 3D Hull

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```
1
   template <class P, class F>
   void delaunay(vector<P> &ps, F trifun) {
     if (sz(ps) == 3) {
3
       int d = (ps[0].cross(ps[1], ps[2]) < 0);</pre>
       trifun(0, 1 + d, 2 - d);
     vector<P3> p3;
     for (P p : ps) p3.emplace_back(p.x, p.y, p.dist2());
9
     if (sz(ps) > 3)
       for (auto t : hull3d(p3))
11
         if ((p3[t.b] - p3[t.a])
             .cross(p3[t.c] - p3[t.a])
13
              .dot(P3(0, 0, 1)) < 0)
           trifun(t.a, t.c, t.b);
15 }
```

6.11. Half Plane Intersection

```
struct Line {
     Point P;
     Vector v;
     bool operator<(const Line 8b) const {
       return atan2(v.y, v.x) < atan2(b.v.y, b.v.x);</pre>
   bool OnLeft(const Line &L, const Point &p) {
9
     return Cross(L.v, p - L.P) > 0;
   Point GetIntersection(Line a, Line b) {
11
     Vector u = a.P - b.P;
     Double t = Cross(b.v, u) / Cross(a.v, b.v);
13
     return a.P + a.v * t;
15
   int HalfplaneIntersection(Line *L, int n, Point *poly) {
17
     sort(L, L + n);
19
     int first, last;
     Point *p = new Point[n];
21
     Line *q = new Line[n];
     q[first = last = 0] = L[0];
23
     for (int i = 1; i < n; i++) {
       while (first < last δδ !OnLeft(L[i], p[last - 1]))</pre>
25
       while (first < last && !OnLeft(L[i], p[first])) first++;
27
       q[++last] = L[i];
       if (fabs(Cross(q[last].v, q[last - 1].v)) < EPS) {</pre>
29
          last--
         if (OnLeft(q[last], L[i].P)) q[last] = L[i];
31
       if (first < last)</pre>
         p[last - 1] = GetIntersection(q[last - 1], q[last]);
35
     while (first < last && !OnLeft(q[first], p[last - 1]))</pre>
       last--;
```

```
37
     if (last - first <= 1) return 0:
     p[last] = GetIntersection(q[last], q[first]);
39
     for (int i = first; i <= last; i++) poly[m++] = p[i];</pre>
41
     return m;
43 }
```

Strings

7.1. Knuth-Morris-Pratt Algorithm

```
1 vector<int> pi(const string &s) {
       vector<int> p(s.size());
       for (int i = 1; i < s.size(); i++) {</pre>
          int g = p[i - 1];
          while (g \delta \delta s[i] != s[g]) g = p[g - 1];
         p[i] = g + (s[i] == s[g]);
 7
       }
       return p;
9
    }
    vector<int> match(const string &s, const string &pat) {
  vector<int> p = pi(pat + '\0' + s), res;
  for (int i = p.size() - s.size(); i < p.size(); i++)</pre>
11
          if (p[i] == pat.size())
13
            res.push_back(i - 2 * pat.size());
15
       return res;
```

7.2. Aho-Corasick Automaton

```
1 struct Aho_Corasick {
     static const int maxc = 26, maxn = 4e5;
     struct NODES {
       int Next[maxc], fail, ans;
    NODES T[maxn];
     int top, qtop, q[maxn];
     int get node(const int &fail) {
      fill_n(T[top].Next, maxc, 0);
T[top].fail = fail;
       T[top].ans = 0;
       return top++;
    int insert(const string &s) {
       int ptr = 1;
       for (char c : s) \{ // change char id \}
         if (!T[ptr].Next[c]) T[ptr].Next[c] = get_node(ptr);
         ptr = T[ptr].Next[c];
       return ptr;
     } // return ans_last_place
     void build_fail(int ptr) {
       int tmp;
       for (int i = 0; i < maxc; i++)</pre>
         if (T[ptr].Next[i]) {
           tmp = T[ptr].fail;
           while (tmp != 1 && !T[tmp].Next[i])
           tmp = T[tmp].fail;
if (T[tmp].Next[i] != T[ptr].Next[i])
           if (T[tmp].Next[i]) tmp = T[tmp].Next[i];
T[T[ptr].Next[i]].fail = tmp;
           q[qtop++] = T[ptr].Next[i];
     void AC_auto(const string &s) {
       int ptr = 1;
       for (char c : s) {
         while (ptr != 1 && !T[ptr].Next[c]) ptr = T[ptr].fail;
         if (T[ptr].Next[c])
           ptr = T[ptr].Next[c];
           T[ptr].ans++;
         }
      }
     void Solve(string &s) {
       for (char &c : s) // change char id
       for (int i = 0; i < qtop; i++) build_fail(q[i]);</pre>
       AC_auto(s);
for (int i = qtop - 1; i > -1; i--)
         T[T[q[i]].fail].ans += T[q[i]].ans;
     void reset() {
       qtop = top = q[0] = 1;
       get_node(1);
```

```
} AC;
   // usage example
59
   string s, S;
   int n, t, ans_place[50000];
61
   int main() {
63
     Tie cin >> t;
      while (t--)
65
        AC.reset();
        cin >> S >> n;
        for (int i = 0; i < n; i++) {
67
          cin >> s;
69
          ans_place[i] = AC.insert(s);
        AC.Solve(S);
for (int i = 0; i < n; i++)
71
73
          cout << AC.T[ans_place[i]].ans << '\n';</pre>
75 }
```

7.3. Suffix Array

```
1 // sa[i]: starting index of suffix at rank i
                   0-indexed, sa[0] = n (empty string)
    // lcp[i]: lcp of sa[i] and sa[i - 1], lcp[\overline{0}] = 0
    struct SuffixArray {
       vector<int> sa, lcp;
       SuffixArray(string &s,
                         int lim = 256) { // or basic_string<int>
          int n = sz(s) + 1, k = 0, a, b;
vector<int> x(all(s) + 1), y(n), ws(max(n, lim)),
 9
          rank(n);
          sa = lcp = y, iota(all(sa), 0);
for (int j = 0, p = 0; p < n;
    j = max(1, j * 2), lim = p) {
    p = j, iota(all(y), n - j);
}</pre>
11
13
             for (int i = 0; i < n; i++)
if (sa[i] >= j) y[p++] = sa[i] - j;
15
17
             fill(all(ws), 0);
             for (int i = 0; i < n; i++) ws[x[i]]++;
             for (int i = 1; i < lim; i++) ws[i] += ws[i - 1];
19
             for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
             swap(x, y), p = 1, x[sa[0]] = 0;
for (int i = 1; i < n; i++)
  a = sa[i - 1], b = sa[i],</pre>
23
25
                x[b] = (y[a] == y[b] \delta \delta y[a + j] == y[b + j])
                          ? p - 1 : p++;
27
29
          for (int i = 1; i < n; i++) rank[sa[i]] = i;
          for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)
for (k &&-, j = sa[rank[i] - 1];
   s[i + k] == s[j + k]; k++)</pre>
31
33
35 };
```

7.4. Suffix Tree

```
struct SAM {
                                     // char range
     static const int maxc = 26;
     static const int maxn = 10010; // string len
     struct Node {
       Node *green, *edge[maxc];
       int max_len, in, times;
     } *root, *last, reg[maxn * 2];
     int top;
9
     Node *get_node(int _max) {
       Node *re = &reg[top++];
11
       re->in = 0, re->times = 1;
       re->max_len = _max, re->green = 0;
13
       for (int i = 0; i < maxc; i++) re->edge[i] = 0;
15
     void insert(const char c) { // c in range [0, maxc)
       Node *p = last;
       last = get_node(p->max_len + 1);
19
       while (p && !p->edge[c])
         p->edge[c] = last, p = p->green;
       if (!p) last->green = root;
21
       else {
23
         Node *pot_green = p->edge[c];
         if ((pot_green->max_len) == (p->max_len + 1))
25
           last->green = pot_green;
         else {
           Node *wish = get_node(p->max_len + 1);
27
           wish->times = 0:
           while (p && p->edge[c] == pot_green)
29
```

```
p->edge[c] = wish, p = p->green;
             for (int i = 0; i < maxc; i++)
31
               wish->edge[i] = pot_green->edge[i];
33
             wish->green = pot_green->green;
             pot_green->green = wish;
35
             last->green = wish;
37
        }
39
      Node *q[maxn * 2];
     int ql, qr;
41
      void get_times(Node *p) {
        ql = 0, qr = -1, reg[0].in = 1;
for (int i = 1; i < top; i++) reg[i].green->in++;
for (int i = 0; i < top; i++)</pre>
43
45
          if (!reg[i].in) q[++qr] = &reg[i];
        while (ql \ll qr) {
          q[ql]->green->times += q[ql]->times;
          if (!(--q[ql]->green->in)) q[++qr] = q[ql]->green;
49
        }
51
      void build(const string &s) {
53
        top = 0:
        root = last = get_node(0);
for (char c : s) insert(c - 'a'); // change char id
55
        get_times(root);
57
      // call build before solve
59
     int solve(const string &s) {
        Node *p = root;
        for (char c : s)
61
          if (!(p = p->edge[c - 'a'])) // change char id
63
             return 0;
        return p->times;
   };
```

7.5. Cocke-Younger-Kasami Algorithm

```
1 struct rule {
     // s -> xy // if y == -1, then s -> x (unit rule)
     int s, x, y, cost;
  };
5
   int state:
  // state (id) for each letter (variable)
// lowercase letters are terminal symbols
  map<char, int> rules;
   vector<rule> cnf;
11
   void init() {
     state = 0;
13
     rules.clear();
     cnf.clear();
15 }
   // convert a cfg rule to cnf (but with unit rules) and add
17
   void add_to_cnf(char s, const string &p, int cost) {
     if (!rules.count(s)) rules[s] = state++;
19
     for (char c : p)
     if (!rules.count(c)) rules[c] = state++;
if (p.size() == 1) {
21
23
       cnf.push_back({rules[s], rules[p[0]], -1, cost});
     } else {
25
       // length >= 3 -> split
       int left = rules[s];
27
       int sz = p.size();
       for (int i = 0; i < sz - 2; i++)
29
         cnf.push_back({left, rules[p[i]], state, 0});
         left = state++;
31
       cnf.push_back(
       {left, rules[p[sz - 2]], rules[p[sz - 1]], cost});
33
35 }
   constexpr int MAXN = 55;
   vector<long long> dp[MAXN][MAXN];
   // unit rules with negative costs can cause negative cycles
39
   vector<bool> neg_INF[MAXN][MAXN];
   43
45
       if (neg_c || neg_INF[l][r][c.x]) {
         dp[l][r][c.s] = 0;
neg_INF[l][r][c.s] = true;
47
49
       } else {
```

```
dp[l][r][c.s] = cost;
51
      }
53
   }
    void bellman(int l, int r, int n) {
      for (int k = 1; k <= state; k++)</pre>
55
        for (rule c : cnf)
57
           if (c.y == -1)
             relax(l, r, c, dp[l][r][c.x] + c.cost, k == n);
59
    void cyk(const string &s) {
61
      vector<int> tok;
      for (char c : s) tok.push_back(rules[c]);
      for (int i = 0; i < tok.size(); i++) {
  for (int j = 0; j < tok.size(); j++) {</pre>
63
           dp[i][j] = vector<long long>(state + 1, INT_MAX);
65
           neg_INF[i][j] = vector<bool>(state + 1, false);
67
        dp[i][i][tok[i]] = 0;
69
        bellman(i, i, tok.size());
71
      for (int r = 1; r < tok.size(); r++) {</pre>
        for (int l = r - 1; l >= 0; l--) {
           for (int k = 1; k < r; k++)
73
             for (rule c : cnf)
75
               if (c.y != -1)
                  relax(l, r,
                         dp[l][k][c.x] + dp[k + 1][r][c.y] +
77
                         c.cost):
79
           bellman(l, r, tok.size());
81
      }
    }
83
    // usage example
   int main() {
85
      init();
      add_to_cnf('S', "aSc", 1);
add_to_cnf('S', "BBB", 1);
add_to_cnf('S', "SB", 1);
add_to_cnf('B', "b", 1);
87
89
91
      cyk("abbbbc");
      // dp[0][s.size() - 1][rules[start]] = min cost to
      // generate
      cout << dp[0][5][rules['S']] << '\n'; // 7</pre>
95
      cyk("acbc")
      cout << dp[0][3][rules['S']] << '\n'; // INT_MAX</pre>
      add_to_cnf('S', "S", -1);
cyk("abbbbc");
99
      cout << neg_INF[0][5][rules['S']] << '\n'; // 1</pre>
```

7.6. Z Value

```
int z[n];
void zval(string s) {
    // z[i] => longest common prefix of s and s[i:], i > 0
    int n = s.size();
    z[0] = 0;
    for (int b = 0, i = 1; i < n; i++) {
        if (z[b] + b <= i) z[i] = 0;
        else z[i] = min(z[i - b], z[b] + b - i);
        while (s[i + z[i]] == s[z[i]]) z[i]++;
        if (i + z[i] > b + z[b]) b = i;
}
```

7.7. Manacher's Algorithm

```
int z[n];
1
   void manacher(string s) {
3
      // z[i] => longest odd palindrome centered at i is
                  s[i - z[i] ... i + z[i]]
      //
// to get all palindromes (including even length),
// insert a '#' between each s[i] and s[i + 1]
      int n = s.size();
      z[0] = 0;
      for (int b = 0, i = 1; i < n; i++) {
9
        if (z[b] + b >= i)
11
          z[i] = min(z[2 * b - i], b + z[b] - i);
        else z[i] = 0;
        while (i + z[i] + 1 < n \ \&\& \ i - z[i] - 1 >= 0 \ \&\&
13
                s[i + z[i] + 1] == s[i - z[i] - 1])
15
        if (z[i] + i > z[b] + b) b = i;
17
```

7.8. Lyndon Factorization

```
vector<string> duval(string s) {
     // s += s for min rotation
     int n = s.size(), i = 0, ans;
     vector<string> res;
     while (i < n) { // change to i < n / 2 for min rotation
5
       ans = i;
int j = i + 1, k = i;
       for (; j < n \delta \delta s[k] <= s[j]; j++)
         k = s[k] < s[j] ? i : k + 1;
9
       while (i <= k)
11
         res.push_back(s.substr(i, j - k));
         i += j - k;
       }
13
     // min rotation is s.substr(ans, n / 2)
15
     return res;
17 }
```

```
7.9. Palindromic Tree
   struct palindromic_tree {
      struct node {
 3
        int next[26], fail, len;
 5
        num; // cnt: appear times, num: number of pal. suf.
        \label{eq:node_int} \begin{array}{ll} \mbox{node(int l = 0) : fail(0), len(l), cnt(0), num(0) \{} \\ \mbox{for (int i = 0; i < 26; ++i) next[i] = 0;} \end{array}
 7
 9
      };
      vector<node> St;
11
      vector<char> s;
      int last, n;
13
     palindromic_tree() : St(2), last(1), n(0) {
        St[0].fail = 1, St[1].len = -1, s.pb(-1);
15
      inline void clear() {
        St.clear(), s.clear(), last = 1, n = 0;
17
        St.pb(0), St.pb(-1);
        St[0].fail = 1, s.pb(-1);
19
21
     inline int get_fail(int x) {
        while (s[n - St[x].len - 1] != s[n]) x = St[x].fail;
23
        return x;
25
      inline void add(int c) {
        s.push_back(c -= 'a'), ++n;
        int cur = get_fail(last);
27
        if (!St[cur].next[c]) {
29
          int now = SZ(St);
          St.pb(St[cur].len + 2);
          St[now].fail = St[get_fail(St[cur].fail)].next[c];
31
          St[cur].next[c] = now;
          St[now].num = St[St[now].fail].num + 1;
33
35
        last = St[cur].next[c], ++St[last].cnt;
      inline void count() { // counting cnt
37
        auto i = St.rbegin();
for (; i != St.rend(); ++i) {
39
          St[i->fail].cnt += i->cnt;
41
43
     inline int size() { // The number of diff. pal.
        return SZ(St) - 2;
45
```