CSCI317 – Final Project

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Solution 1:

(1) a short description of an improvement

ASUMPTIONS:

- I have considered using a Materialised View but since LINEITEM can be frequently updated and I_shipdate can be taken from the user input, the Materialised View would have to be frequently updated and re-compiled. Thus, I don't think using a Materialised View is a good solution in this case.
- I have also considered demoralising the database but that requires moving columns from SUPPLIER into LINEITEM. Since LINEITEM is already a large table, demoralising the database could decrease the performance of the queries rather than improving them.
- Like using a Materialised View, since LINEITEM can be frequently updated and/or inserted into, clustering the table may slow down UPDATE, INSERT and DELETE operations.
- Clustering the tables costs more than the amount of storage allowed.

However, I noticed that the VIEWS used in task1.sql are not being merged correctly (See Figure 1) so I've performed query transformation of both the query and the SELECT statement so that the Views are merged properly.

Figure 1: Original query where REVENUE view is not merged correctly

29 rows selected.

```
SOL> @showplan
SQL> SELECT * FROM TABLE(DBMS_XPLAN.DISPLAY);
PLAN TABLE OUTPUT
Plan hash value: 39222225
                          | Name | Rows | Bytes | Cost (%CPU)| Time
.....
| 0 | SELECT STATEMENT | | 1817K| 220M| 1825 (3)| 00:00:01 |
|* 1 | FILTER
 2 | SORT GROUP BY
                                                    1817K| 220M| 1825 (3)| 00:00:01
I* 3 I FILTER
i* 4 |
       HASH JOIN
          HASH JOIN | SIDX
                                                    1817K| 220M| 1779 (1)| 00:00:01
  5 I
                                                    3114 |
                                                            240K
                                                                 18 (0)| 00:00:01
 * 6
          TABLE ACCESS BY INDEX ROWID BATCHED! LINEITEM | 1816K|
                                                            83M l
                                                                 1756 (1) 00:00:01
  7
           BITMAP CONVERSION TO ROWIDS
 * 8 |
           BITMAP INDEX RANGE SCAN
  9
        SORT AGGREGATE
                                                       1 |
                                                                        (3) | 00:00:01
        SORT GROUP BY
                                                       1 |
                                                             35
                                                                 1803
  10 l
                                                                       (3) | 00:00:01
 * 11 |
         FILTER
          TABLE ACCESS BY INDEX ROWID BATCHED | LINEITEM | 1816K|
 12
                                                             60M
                                                                 1756 (1) | 00:00:01
 13 I
           BITMAP CONVERSION TO ROWIDS
|* 14 |
          BITMAP INDEX RANGE SCAN
Predicate Information (identified by operation id):
  1 - filter(SUM("L EXTENDEDPRICE"*(1-"L DISCOUNT"))= (SELECT
            MAX(SUM("L EXTENDEDPRICE"*(I-"L DISCOUNT"))) FROM "TPCHR"."LINEITEM" "LINEITEM" WHERE
            TO DATE('18-SEP-98')>=TO DATE('28-JUN-92') AND "L SHIPDATE"<='18-SEP-98' AND
            "L SHIPDATE">='28-JUN-92' GROUP BY :B1))
  3 - filter(TO DATE('18-SEP-98')>=TO DATE('28-JUN-92'))
  4 - access("L SUPPKEY"="S SUPPKEY")
  6 - filter("L SUPPKEY">=0)
  8 - access("L SHIPDATE">='28-JUN-92' AND "L SHIPDATE"<='18-SEP-98')
 11 - filter(TO_DATE('18-SEP-98')>=TO_DATE('28-JUN-92'))
 14 - access("L_SHIPDATE">='28-JUN-92" AND "L_SHIPDATE"<='18-SEP-98')
Note
  - dynamic statistics used: dynamic sampling (level=2)
39 rows selected.
```

Figured 2: solution1.sql where REVENUE is merged correctly to the SELECT statement

The transformation is performed by predicate pushing. In this case, the predicate

" l_shipdate >= '28-JUN-92' and l_shipdate <= '18-SEP-98' " is pushed down the syntax tree before the JOIN is performed. This helps reducing the size of the view REVENUE before it being joined with SUPPLIER; thus, the transformed query costs less than the original query.

The transformed query is as follows:

```
create view revenue(
        supplier_no,
        total revenue
    ) as
    select
        l suppkey,
        (l_extendedprice * (l - l_discount))
    from
        lineitem
    where
        l shipdate >= '28-JUN-92'
        and l_shipdate <= '18-SEP-98';
select
    s_suppkey,
    s name,
    s address,
    s phone,
    sum(total_revenue)
from
    revenue, supplier
where
    supplier_no = s_suppkey
group by
    s_suppkey,
    s name,
    s address,
    s phone
having sum(total revenue) = (select max(sum(total revenue)) from revenue group by s suppkey)
order by
    s_suppkey;
```

To further improve the performance of the queries, two indices were created, including:

- A Bitmap index on LINEITEM(L_SHIPDATE) and
- A B*-Tree index on SUPPLIER(S_SUPPKEY,S_NAME,S_ADDRESS,S_PHONE)

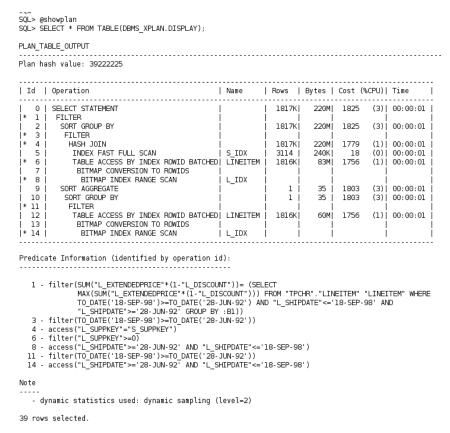
The two indices allow the query to traverse the index which only has the necessary columns, without the need to access the entire tables LINEITEM and SUPPLIER.

```
create bitmap index L_IDX on lineitem(l_shipdate);
create index S_IDX on supplier(s_suppkey,s_name,s_address,s_phone);
```

(2) information about the benefits from an improvement

- As shown in Figure 1 and Figure 2 above, the transformation is performed by predicate pushing, which helps reducing the size of the view REVENUE before it being joined with SUPPLIER; thus, the transformed query costs less than the original query.
- The two indices further improve the performance of the SELECT statement. They allow for two index scans rather than accessing to the entire tables.

The figures below demonstrate the query processing plans with and without the two indices:



The transformed query processing plan with the two indices

	ABLE_OUTPUT ash value: 3780874396							
r Laii i	iasii vatue: 3/606/4390							
Id	Operation	Name	Rows	Bytes	Cost	(%CPU)	Time	Ī
	SELECT STATEMENT		1817K	220M	12243	(1)	00:00:01	ļ
	FILTER SORT GROUP BY FILTER		 1817K 	220M	12243	(1)	00:00:01	
* 4	HASH JOIN		1817K	220M	12197	(1)	00:00:01	ĺ
5	TABLE ACCESS FULL	SUPPLIER	3114	240K	34	(0)	00:00:01	1
	TABLE ACCESS FULL	LINEITEM					00:00:01	
	SORT AGGREGATE						00:00:01	
_	SORT GROUP BY		1	35	12218	(1)	00:00:01	
	FILTER TABLE ACCESS FULL	LINEITEM	 1816K	60M	12171	(1)	00:00:01	
Predic	ate Information (identifi	ed by opera	ation id)	:				-
<pre>1 - filter(SUM("L_EXTENDEDPRICE"*(1-"L_DISCOUNT"))= (SELECT</pre>								
3 -	filter(TO DATE('18-SEP-9		_		10 02	. 50 0		-,,
	access("L_SUPPKEY"="S_SL							
6 - filter("L_SUPPKEY">=0 AND "L_SHIPDATE">='28-JUN-92' AND "L_SHIPDATE"<='18-SEP-98')								
9 -	filter(TO DATE('18-SEP-9		TE('28-JL	IN-92'))				
10 - filter("L_SHIPDATE">='28-JUN-92' AND "L_SHIPDATE"<='18-SEP-98')								

The query processing plan without the two indices

(3) information about the costs of an improvement, i.e. documented investments into transient and persistent storage. In terms of transient storage, no new investment into transient storage has been made compared to task1.sql. However, an investment of a total of 7.0MB into persistent storage is needed for the two indices:

- Bitmap index on LINEITEM (L SHIPDATE)
- B*-tree index on SUPPLIER(S_SUPPKEY, S_NAME, S_ADDRESS, S_PHONE) speed up the query.

	∯ Index Size (MB)
1 L_IDX	6.5
2 S_IDX	0.5

(4) a report from implementation of an improvement

```
SQL> /* IMPROVED STATEMENTS */
SQL> create bitmap index L_IDX on lineitem(l_shipdate);

Index L_IDX created.

SQL> create index S_IDX on supplier(s_suppkey,s_name,s_address,s_phone);

Index S_IDX created.
```

```
SQL> create view revenue(
    supplier_no,
    total_revenue

4    ) as
5    select
6    l_suppkey,
7     (l_extendedprice * (1 - l_discount))
8    from
9    lineitem
10    where
11         l_shipdate >= '28-JUN-92'
12    and l_shipdate <= '18-SEP-98';</pre>
```

View REVENUE created.

```
SQL> explain plan for
 2 select
 3 s_suppkey,
 4 s name,
 5 s address,
 6 s_phone,
 7 sum(total revenue)
 8 from
 9 revenue, supplier
 10 where
        supplier no = s suppkey
11
 12 group by
 s_suppkey,
s_name,
15 s_address,
16 s_phone
17 having sum(total_revenue) = (select max(sum(total_revenue)) from revenue group by s_suppkey)
 18 order by
        s suppkey;
 19
```

Explained.

SQL> @showplan

SQL> SELECT * FROM TABLE(DBMS_XPLAN.DISPLAY);

PLAN_TABLE_OUTPUT

Plan hash value: 39222225

I	d	Operation	Name	 	Rows	Bytes	Cost	(%CPU) Time	
1	0	SELECT STATEMENT		I	1817K	220M	1825	(3) 00:00:01	-
*	1	FILTER			1	1		I	I
1	2	SORT GROUP BY			1817K	220M	1825	(3) 00:00:01	-
*	3	FILTER			1	1		I	I
*	4	HASH JOIN			1817K	220M	1779	(1) 00:00:01	-
1	5	INDEX FAST FULL SCAN	S_IDX		3114	240K	18	(0) 00:00:01	-
*	6	TABLE ACCESS BY INDEX ROWID BATCHED	LINEITEM		1816K	83M	1756	(1) 00:00:01	-
1	7	BITMAP CONVERSION TO ROWIDS			1	I		1	1
*	8	BITMAP INDEX RANGE SCAN	L_IDX		1	1		1	1
1	9	SORT AGGREGATE			1	35	1803	(3) 00:00:01	.

Predicate Information (identified by operation id):

```
Note
   - dynamic statistics used: dynamic sampling (level=2)
39 rows selected.
SQL> /* Retrieve index's size*/
SQL> select segment_name,
         sum(bytes)/1024/1024 as "Index Size (MB)"
  3 from user segments
  4 where segment_name='L_IDX'
       or segment_name='S_IDX'
  6 group by segment name;
SEGMENT_NAME Index Size (MB)
L IDX
                           6.5
```

0.5

2 rows selected.

S IDX

SQL> drop view revenue;

View REVENUE dropped.

SQL> drop index L_IDX;

Index L_IDX dropped.

SQL> drop index S_IDX;

Index S_IDX dropped.