Comp 442: Compiler Design

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Winter 2024

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Assignment 1

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1 Transformed Grammar into LL(1)

Find below the complete grammar after removing all EBNF notations, left recursions and ambuigities from the provided grammar.

```
START
aParams
                             \rightarrow \text{ expr rept-aParams1} \mid \epsilon
aParamsTail
                             \rightarrow , expr
addOp
                             \rightarrow + | - | or
                             → term rightrec—arithExpr
arithExpr
arraySize
                             \rightarrow [ arraySize2
arraySize2
                             \rightarrow intNum | | ]
assignOp
                             \rightarrow arithExpr expr2
expr
expr2
                             \rightarrow relOp arithExpr | \epsilon
fParams
                             \rightarrow id : type rept-fParams3 rept-fParams4 | \epsilon
fParamsTail
                             → , id : type rept-fParamsTail4
factor
                             → id factor2 reptVariableOrFunc | intLit | floatLit
                             \rightarrow ( aParams ) | rept-idnest1
factor 2
                             \rightarrow idnest reptVariableOrFunc | \epsilon
reptVariableOrFunc
funcBody
                             \rightarrow { rept-funcBody1 }
funcDecl
                             \rightarrow funcHead;
funcDef
                             \rightarrow funcHead funcBody
funcHead
                             → func id (fParams) arrow returnType
functionCall
                             → rept-functionCallO id ( aParams )
idnest
                             \rightarrow . id idnest2
idnest2
                             \rightarrow ( aParams ) | rept-idnest1
implDef
                             \rightarrow impl id { rept-implDef3 }
indice
                             \rightarrow [ arithExpr ]
                             \rightarrow funcDecl | varDecl
memberDecl
multOp
                             \rightarrow * | / | and
                             \rightarrow inherits id rept-opt-structDecl22 | \epsilon
opt-structDecl2
                             \rightarrow \text{rept-prog}0
prog
                             → arithExpr relOp arithExpr
relExpr
                             \rightarrow eq | neq | lt | gt | leq | geq
relOp
rept-aParams1
                             \rightarrow aParamsTail rept-aParams1 | \epsilon
rept-fParams3
                             \rightarrow arraySize rept-fParams3 | \epsilon
rept-fParams4
                             \rightarrow fParamsTail rept-fParams4 | \epsilon
```

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```
rept-fParamsTail4
                            \rightarrow arraySize rept-fParamsTail4 | \epsilon
rept-funcBody1
                            \rightarrow varDeclOrStat rept-funcBody1 | \epsilon
rept-functionCall0
                            \rightarrow idnest rept-functionCall0 | \epsilon
rept-idnest1
                            \rightarrow indice rept-idnest1 | \epsilon
                            \rightarrow funcDef rept-implDef3 | \epsilon
rept-implDef3
                           \rightarrow , id rept-opt-structDecl22 | \epsilon
rept-opt-structDecl22
rept-prog0
                            \rightarrow structOrImplOrFunc rept-prog0 | \epsilon
rept-statBlock1
                            \rightarrow statement rept-statBlock1 | \epsilon
rept-structDecl4
                            \rightarrow visibility memberDecl rept-structDecl4 | \epsilon
rept-varDecl4
                            \rightarrow arraySize rept-varDecl4 | \epsilon
rept-variable0
                            \rightarrow idnest rept-variable 0 | \epsilon
                            \rightarrow indice rept-variable 2 | \epsilon
rept-variable2
returnType
                            \rightarrow type | void
rightrec-arithExpr
                                    {\it addOp\ term\ rightrec-arithExpr}
rightrec-term
                                    multOp factor rightrec-term
sign
                            \rightarrow + \mid -
                            \rightarrow { rept-statBlock1 } | statement | \epsilon
statBlock
                            \rightarrow if (relExpr) then statBlock else statBlock; | v
statement
                            \rightarrow rept-idnest1 statement3 | (aParams) statement4
statement2
                            \rightarrow . id statement2 | assignOp expr ;
statement3
                            \rightarrow . id statement2 | ;
statement4
structDecl
                            → struct id opt-structDecl2 { rept-structDecl4 } ;
                            → structDecl | implDef | funcDef
structOrImplOrFunc
                            → factor rightrec-term
term
                            \rightarrow integer | float | id
type
                            → let id : type rept-varDecl4 ;
varDecl
varDeclOrStat
                            → varDecl | statement
variable
                            \rightarrow id variable2
variable2
                            → rept-idnest1 reptvariable | (aParams) varIdnest
                            \rightarrow varIdnest reptvariable | \epsilon
reptvariable
varIdnest
                            \rightarrow . id varIdnest2
varIdnest2
                            \rightarrow ( aParams ) varIdnest | rept-idnest1
visibility
                           → public | private
```

2 First and Follow Sets

Table 1: First Set

nonterminal	first set
ADDOP	plus minus or
ARRAYSIZE2	intlit rsqbr
EXPR2	eq neq lt gt leq geq
FACTOR2	lpar lsqbr
FUNCBODY	lcurbr
FUNCHEAD	func
FPARAMS	id

FUNCTIONCALL	id dot
IDNEST2	lpar lsqbr
FUNCDECL	func
ARITHEXPR	intlit floatlit lpar not plus minus
RELOP	
	eq neq lt gt leq geq
APARAMSTAIL DEPTA DA DA MG1	comma
REPTAPARAMS1	comma
REPTFPARAMS3	lsqbr
FPARAMSTAIL	comma
REPTFPARAMS4	comma
REPTFPARAMSTAIL4	lsqbr
REPTFUNCBODY1	let if while read write return
REPTFUNCTIONCALL0	dot
REPTIMPLDEF3	func
REPTOPTSTRUCTDECL22	comma
REPTPROG0	struct impl func
MEMBERDECL	let func
ARRAYSIZE	lsqbr
REPTVARIABLE0	dot
INDICE	lsqbr
REPTVARIABLE2	lsqbr
IDNEST	dot
REPTVARIABLEORFUNC	dot
RETURNTYPE	void integer float id
RIGHTRECARITHEXPR	plus minus or
MULTOP	mult div and
SIGN	plus minus
START	struct impl func
PROG	struct impl func
REPTSTATBLOCK1	if while read write return
RELEXPR	intlit floatlit lpar not plus minus
STATBLOCK	lcurbr if while read write return
STATEMENT3	dot equal
ASSIGNOP	equal
EXPR	intlit floatlit lpar not plus minus
STATEMENT4	dot semi
ID	ϵ
STATEMENT2	lpar dot lsqbr equal
OPTSTRUCTDECL2	inherits
REPTSTRUCTDECL4	public private
STRUCTORIMPLORFUNC	struct impl func
STRUCTDECL STRUCTDECL	struct
IMPLDEF	impl
FUNCDEF	func
TERM	intlit floatlit lpar not plus minus
FACTOR	intlit floatlit lpar not plus minus

RIGHTRECTERM	mult div and
TYPE	integer float id
REPTVARDECL4	lsqbr
VARDECLORSTAT	let if while read write return
VARDECL	let
STATEMENT	if while read write return
VARIABLE	id
VARIABLE2	lpar lsqbr dot
REPTVARIABLE	dot
VARIDNEST2	lpar lsqbr
APARAMS	intlit floatlit lpar not plus minus
VARIDNEST	dot
REPTIDNEST1	lsqbr
VISIBILITY	public private

Table 2: Follow Set

Non-Terminal	Follow Set
ADDOP	intlit floatlit lpar not plus minus
ARRAYSIZE2	semi lsqbr rpar comma
EXPR2	comma rpar semi
FACTOR2	mult div and dot rsqbr
	eq neq lt gt leq geq plus
	minus or comma rpar semi
FUNCBODY	struct impl func reurbr
FUNCHEAD	semi lcurbr
FPARAMS	rpar
FUNCTIONCALL	ϵ
IDNEST2	mult div and dot id rsqbr
	eq neq lt gt leq geq plus
	minus or comma rpar semi
FUNCDECL	rcurbr public private
ARITHEXPR	rsqbr eq neq lt gt leq geq comma rpar semi
RELOP	intlit floatlit lpar not plus minus
APARAMSTAIL	comma rpar
REPTAPARAMS1	rpar
REPTFPARAMS3	rpar comma
FPARAMSTAIL	comma rpar
REPTFPARAMS4	rpar
REPTFPARAMSTAIL4	comma rpar
REPTFUNCBODY1	rcurbr
REPTFUNCTIONCALL0	id
REPTIMPLDEF3	rcurbr
REPTOPTSTRUCTDECL22	lcurbr
REPTPROG0	ϵ

MEMBERDECL	rcurbr public private
ARRAYSIZE	semi lsqbr rpar comma
REPTVARIABLE0	ϵ
INDICE	equal mult div and lsqbr
	id dot rsqbr eq neq lt gt leq geq
	plus minus or comma rpar semi
REPTVARIABLE2	ϵ
IDNEST	mult div and dot id
	rsqbr eq neq lt gt leq
	geq plus minus or comma rpar semi
REPTVARIABLEORFUNC	mult div and rsqbr
	eq neq lt gt leq geq
	plus minus or comma rpar semi
RETURNTYPE	semi lcurbr
RIGHTRECARITHEXPR	rsqbr eq neq lt gt leq geq comma rpar semi
MULTOP	intlit floatlit lpar not plus minus
SIGN	intlit floatlit lpar not plus minus
START	ϵ
PROG	ϵ
REPTSTATBLOCK1	rcurbr
RELEXPR	rpar
STATBLOCK	else semi
STATEMENT3	else semi let if while
Lagranop	read write return reurbr
ASSIGNOP	intlit floatlit lpar not plus minus
EXPR	comma rpar semi
STATEMENT4	else semi let if while
	read write return reurbr
ID	mult div and id lpar dot lsqbr equal
	rsqbr eq neq lt gt leq geq plus
CEATED ADMES	minus or comma rpar semi
STATEMENT2	else semi let if while
	read write return rcurbr
OPTSTRUCTDECL2 REPTSTRUCTDECL4	leurbr
STRUCTORIMPLORFUNC	reurbr
STRUCTORIMPLORFUNC	struct impl func
IMPLDEF	struct impl func
FUNCDEF	struct impl func
TERM	struct impl func rcurbr rsqbr eq neq lt gt leq
1 171 (171	geq plus minus or
FACTOR	comma rpar semi mult div and rsqbr eq neq
11101011	lt gt leq geq plus
	minus or comma rpar semi
RIGHTRECTERM	rsqbr eq neq lt gt
MIGHTREOTERM	rsqur eq neq n gt

	leq geq plus minus
	or comma rpar semi
TYPE	rpar lcurbr comma lsqbr semi
REPTVARDECL4	semi
VARDECLORSTAT	let if while read write return reurbr
VARDECL	public private let if while read write return rcurbr
STATEMENT	else semi let if while read write return rcurbr
VARIABLE	rpar
VARIABLE2	rpar
REPTVARIABLE	rpar
VARIDNEST2	rpar dot
APARAMS	rpar
VARIDNEST	rpar dot
REPTIDNEST1	equal mult div and
	id dot rsqbr eq neq
	lt gt leq geq plus minus
	or comma rpar semi
VISIBILITY	let func

3 Design of recursive parser

In this section I will detail the overall structure of my implementation and explain some important helper methods. First and foremost, a lot of time was spent in the preparation phase before touching a single line of code. This consisted of familiarising myself with the provided grammar tools and applying techniques learned in class to derive an LL(1) grammar. The parser's goal is to validate the syntax of the code, it will get tokens from the lexer and make sure they form a correct sequence. To help in the implementation, it is crucial that every derivation must be deterministic, hence we need to remove ambiguities and left recursions.

Once I had obtained an LL(1) grammar, it was time to start the implementation. I once again chose to stray away from the table driven method because I was unfamiliar with the style and it seemed less intuitive to me. I began by creating a parser class that was initialized with a corresponding lexer. Then I proceeded by defining all the non-terminals from my grammar as methods of the parser. After that, the implementation simply consisted of calling the appropriate functions of the RHS of these non-terminals. I also made use of helper functions such as match and nextToken to iterate through the list of tokens provided by the lexer.

After making sure that my parser passed through the right sequence of function calls, I moved on to the derivation. This part was quite simple and consisted of keeping track of previous derivations and using the built-in Python string method replace() with the corresponding RHS of the non-terminals. In essence, the derivation serves as a visual representation of the steps that the predictive parser is taking.

Finally, I spent some time implementing the FIRST and FOLLOW sets of every non-terminal in my grammar. I made a crucial realisation that it would allow me to almost entirely remove the necessity to return a boolean for the non-terminal methods. In particular, I made heavy use of the FIRST set to see which sequence of tokens to consume

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given a non-terminal that had multiple options. I also use these sets to implement the error recovery method proposed in the lecture. This step required some fine-tuning of my implementation to allow for efficient skipping of syntax errors. However, it remains that I need to do a lot more testing to ensure robustness of my predictive recursive parser.

4 Tools

In this section I will be describing the set of tools and libraries I used in the analysis and implementation of my predictive parser. As mentioned above, I made use of the Java grammar converter tool along with the sample partial solution provided by Professor Paquet. This allowed me to get a head start towards deriving the LL(1) grammar. Then I made use of both recommended tools, namely FLACI and uCalgary CFG checker. The former served as my main tool for grammar verification along with testing valid parsing sequences provided by my classmates. The latter served its purpose by providing suggestions to make the grammar LL(1), which hinted at a general strategy to resolve ambiguities, along with a very valuable automatic generation of the FIRST and FOLLOW sets of every non-terminal. I also once again drew inspiration for the initial structure of my parser from the TeenyTinyCompiler tutorial. And finally, I have to declare that I also made heavy use of chatGPT along with a HTML table to LaTeX converter to turn the FIRST and FOLLOW sets from the uCalgary tool into it's equivalent Python code, using the appropriate TokenType value from the lexer.