

Recap:

→ We recall the Yoneda lemma. It states that for any  $F: C \rightarrow \text{Set}$  and  $a \in C$  we have

$$[C, \text{Set}](C(a, -), F) \cong F a$$

Moreover, each natural transformation

$$\eta: C(a, -) \Rightarrow F$$

has the form

$$\eta_x(f: a \rightarrow x) = (F f) (\eta_a \text{id}_a)$$

→ Let's see what happens if  $F = C(b, -)$ :

$$[C, \text{Set}](C(a, -), C(b, -)) \cong C(b, a)$$

I.e.

$$\forall x. (a \rightarrow x) \rightarrow b \rightarrow x \cong b \rightarrow a$$

and

$$\begin{aligned} \eta_x(f: a \rightarrow x) &= C(b, f) (\eta_a \text{id}_a) \\ &= f \circ (\eta_a \text{id}_a) \end{aligned}$$

→ This works because  $\eta_a \text{id}_a: b \rightarrow a$  so  $\eta_x f: b \rightarrow x$  as desired

$$C(a, x) \xrightarrow{\eta_x} C(b, x)$$

With this we can define a contravariant functor

$$F: C \rightarrow [C, \text{Set}]$$

$$F a = C(a, -)$$

$$F f = \lambda h. h \circ f$$

$$\begin{array}{ccc} a & \xrightarrow{F} & C(a, -) \\ \uparrow f & & \downarrow F f \\ b & \xrightarrow{F} & C(b, -) \end{array}$$

→ This is called the Yoneda embedding of  $C$ .

① Express the Co-Yoneda embedding in Haskell.

Solution:

The Co-Yoneda embedding states that

$$[C, \text{Set}](C(-, a), C(-, b)) \cong C(a, b)$$

In Haskell this is

$$\forall x. (x \rightarrow a) \rightarrow x \rightarrow b \cong a \rightarrow b$$

② Let

$$\mathcal{C} :: \forall x. ((a \rightarrow x) \rightarrow (b \rightarrow x)) \rightarrow (b \rightarrow a)$$

$$\mathcal{C} f = \lambda u. f \text{ id}_a u = f \text{ id}_a$$

$$\Psi :: (b \rightarrow a) \rightarrow (\forall x. (a \rightarrow x) \rightarrow (b \rightarrow x))$$

$$\Psi f = \lambda g_x. g_x \circ f$$

Show that  $\mathcal{C}$  and  $\Psi$  are inverses from one another.

Solution:

$$\begin{aligned} \mathcal{C} \circ \Psi f &= \mathcal{C} \lambda g_x. g_x \circ f \\ &= (\lambda g_x. g_x \circ f) \text{ id}_a \\ &= \text{id}_a \circ f \\ &= f \quad \checkmark \end{aligned}$$

$$\begin{aligned} \Psi \circ \mathcal{C} f &= \Psi (f \text{ id}_a) \\ &= \lambda g_x. g_x \circ (f \text{ id}_a) \\ &= \lambda g_x. c(b, g_x) (f \text{ id}_a) \\ &\stackrel{\text{Yoneda embedding}}{=} \lambda g_x. f g_x \\ &= f \quad \checkmark \end{aligned}$$

③

Work out the Yoneda embedding for a Monoid.

Solution:

Let  $M$  be a monoid category with a single object  $a$ .

The embedding says that

$$[M, \text{Set}] (M(a, -), M(a, -)) \cong M(a, a)$$

i.e. the number of natural transformations from the morphisms of  $M$  to itself matches the number of morphisms in  $M$ .

This is because there are  $\#M(a, a)$  choices to map the identity, and all other maps follow from the naturality condition.

(4)

Apply the covariant Yoneda embedding to a preorder category.

Solution:

Take  $b \leq a$ . By the embedding we get

$$\begin{array}{ccc} a & \xrightarrow{F} & C(a, -) \\ \uparrow \eta & & \downarrow \eta \\ b & \xrightarrow{F} & C(b, -) \end{array}$$

$$\text{So } b \leq a \Leftrightarrow \exists \eta: C(a, -) \Rightarrow C(b, -)$$

$$\Updownarrow$$

$$\exists \eta:: \forall x. (a \rightarrow x) \rightarrow (b \rightarrow x)$$

$$\Updownarrow$$

$$\forall x. a \leq x \Rightarrow b \leq x$$

⑤ Embed the functor category  $[C, D]$  into  $\text{Set}$ .

Solution:

Define the contravariant functor  $e$

$$e: [C, D] \rightarrow [[C, D], \text{Set}]$$

$$e F = [C, D](F, -)$$

$$e(\varphi: G \Rightarrow F) = \lambda \beta: [C, D](F, -) . \beta \circ \varphi$$

I.e. we get

