HONG KONG UNIVERSITY OF SCIENCE & TECHNOLOGY

COMP3031 (Principles of Programming Languages)

Fall 2015

FINAL EXAMINATION

12:30PM - 3:30PM Dec 10, 2015 Thursday LG5 Multi-function Room

Name	
Student ID	ITSC Account

- 1. About the exam:
 - a. This is a closed-book, closed-note examination.
 - b. You CANNOT use any electronic devices including calculators during the examination. Please TURN OFF all of your electronic devices (e.g., mobile phone) and put them into your bag.
 - c. You CANNOT leave during the last 15 minutes of the examination.
- 2. About this paper:
 - a. This paper contains 14 pages, including this title page.
 - b. The total number of points is 100, distributed to seven problems.
- 3. About your answers:
 - a. Write your answers in the designated space following each question.
 - b. Make sure your final answers are clearly recognizable.
 - c. Rough work can be done in the provided "additional blank paper for draft work". Do not, however, write answers there as they will NOT be graded.
 - d. Attempt all questions.

Problem	1	2	3	4	5	6	7	Total
Marks								

Problem 1. SML Programming (20 points)

Given the following SML datatype:

a) Write a function path to return a list of labels of the tree nodes that form the path from the root to the specified label, if there is such a path in the tree; otherwise, return nil. The list is generated through pre-order traversal, and only the first path found is returned.

```
val path = fn : ''a -> ''a tree -> ''a list
   Examples:
-val x = node(0, node(1, leaf(2), leaf(3)), node(2, leaf(3), empty_tree));
val x = node (0, node (1, leaf #, leaf #), node (2, leaf #, empty_tree)) :
int tree
- path 0 x;
val it = [0] : int list
- path 1 x;
val it = [0,1]: int list
- path 2 x;
val it = [0,1,2] : int list
- path 3 x;
val it = [0,1,3] : int list
- path 4 x;
val it = [] : int list
- path true (leaf false);
val it = [] : bool list
```

```
fun search x _ empty_tree = []
  | search x L (leaf(y)) = if x = y then L@[x] else []
  | search x L (node(y, Left, Right)) =
  if x = y then L@[x]
  else
  let
  val L1 = L@[y]
  val el = search x L1 Left
  in
  if el = []
  then search x L1 Right
  else el
  end;
  fun path x y = search x [] y;
```

```
fum path x empty_tree = []
  | path x (leaf(y)) = if x = y then [x] else []
  | path x (node(y, Left, Right)) =
  if x = y then [x]
  else
  let
  val el = path x Left
  in
  if el = []
  then
  let val er = path x Right
  in if er = [] then [] else y::er end
  else y::el
  end;
```

Grading criteria:

Base case 1: 1 pt

Base case 2: 2 pts

val it = 1 : int
- height "e" x;
val it = ~1 : int

val it = ~1 : int
- height ~1 (leaf ~1);

val it = 1 : int

- height true empty_tree;

Base case 3: equal case 1 pts; left subtree traversal 2 pts; right subtree traversal 2 pts; "list append" or "list construct" 2 pts.

Syntax error: -1 pt for one category of similar errors

val height = fn : ''a -> ''a tree -> int

b) Write a function height to return the maximum height of a specified label in the tree. If the label is not in the tree, return ~1.

```
Sol:
fun measure empty_tree = 0
| measure (leaf(_)) = 1
measure (node(_, L, R)) =
let
val 1 = measure L
val r = measure R
if 1 < r then 1+r else 1+l
end;
fun height x empty_tree = ~1
| height x (leaf(y)) = if x = y then 1 else \sim1
height x (node(y, L, R)) =
if x = y then measure (node(y, L, R))
else let
val l = height x L
val r = height x R
if 1 < r then r else 1
end;
```

Grading criteria:

Measure the height of a subtree: 4 pts (1, 1, 2 for each base case)

Function height: 6 pts (1, 2, 3 for each base case) Syntax error: -1 pt for one category of similar errors

Problem 2. Prolog Programming (20 points)

Given a knowledge base of facts in the form of node(X, L, R), where X is the label of a tree node, and L and R labels of the left and right children of the node. The set of facts represents a binary tree where all labels of the tree nodes are unique:

```
node(a,b,c).
node(b,d,e).
node(e,f,g).
```

a) Define a predicate path(X, L) in which X is the label of a tree node, and L is a list of labels that form the path from the root node of the tree to X, if X is in the tree; otherwise, L is an empty list. The code skeleton is given. You only need to fill in the missing predicates, one predicate per blank.

Examples:

```
?- path(c,X).
X = [a, c].
?- path(X,[a,c]).
X = c.
?- path(g,X).
X = [a, b, e, g].
?- path(h,X).
X = [].
```

Grading criteria: 1 pt for each blank Order error of predicates in path(X, L): -1 pt in total

b) Define a predicate height (X,H) such that for a given node X, H is the maximum height of the node in the tree, if the node is in the tree; otherwise, the predicate returns false. The code skeleton is given. You only need to fill in the missing predicates, one predicate per blank.

Examples:

```
?- height(a,X).
X = 4.
?- height(b,X).
X = 3.
?- height(c,X).
X = 1.
?- height(d,X).
X = 1.
?- height(e,X).
X = 2.
?- height(f,X).
X = 1.
?- height(f,X).
```

Grading criteria:

2 pts for each blank Order error of predicates: -1 pt in total

Problem 3. Cuts and Negation in Prolog (10 points)

Given each of the Prolog programs a) - e), write the *first* answer to the query:

```
?- likes(X,a).
a)
bear(a).
animal(a).
likes(X, Y) :- Y=b, !, X=d.
likes(c, Y) :- animal(Y).
 X=c. (2pts)
b)
bear(a).
animal(a).
likes(X, Y) :- Y=a, !, X=d.
likes(X,Y) :- animal(Y).
 X=d. (2pts)
c)
bear(a).
animal(a).
likes(c, X) :- \+ X=a.
likes(c, X) :- \ + \ animal(X).
 false. (2pts)
d)
bear(a).
animal(a).
likes(c, X) :- bear(X), !, fail.
likes(c, X) :- animal(X).
 false. (2pts)
e)
bear(a).
animal(a).
likes(c, X) :- \+ bear(X).
likes(c, X) :- animal(X).
 X=c. (2pts)
```

Problem 4. Prolog Search Tree (10 points)

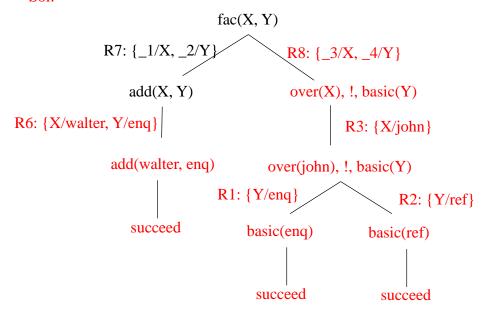
Consider the following program:

```
/*R1*/ basic(enq).
/*R2*/ basic(ref).
/*R3*/ over(john).
/*R4*/ over(walter).
/*R5*/ gen(book).
/*R6*/ add(walter, enq).
/*R7*/ fac(P, F) :- add(P, F).
/*R8*/ fac(P, F) :- over(P), !, basic(F).
/*R9*/ fac(P, F) :- gen(F).
```

Draw the complete Prolog search tree for the query fac(x,y), giving **all** answers. At each **tree edge**, whenever applicable, **i**) indicate the rule number Ri, i=1,...,9, of the rule being applied, and **ii**) the unification(s) being made. At each **tree node** indicate the goal to be satisfied. At each leaf node indicate "succeed" or "fail". The initial step has been done for you.

```
fac(X,Y)
R7: {_1/X, _2/Y}/
add(X,Y)
```

Sol:



Grading criteria:

2 pts for each rule (R6, R8, R3, R1, R2)

-2 pts for every extra answer

Problem 5. Flex and Bison (20 points)

Given the following grammar for expressions on a binary tree where every non-empty tree node has a numeric label:

```
<expression> ::= sum(<tree>) | product(<tree>)
<tree> ::= empty | node(<num>,<tree>,<tree>)
<num> ::= <D1><N> | <N>
<D1> ::= <D1><N> | <D2>
<D2> ::= [1-9]
<N> ::= [0-9]
```

An expression is either a *sum* or a *product*. A *sum* operation adds up the numeric label values of all nodes in the tree. A *product* operation, multiplies the numeric label values of all nodes in the tree. An empty node has a numeric value of 1 in a *product* and 0 in a *sum*. Some examples:

```
sum(node(4,empty,empty)) = 4+0+0 = 4

sum(node(9,node(5,empty,node(2,empty,empty)),node(1,empty,empty)))

= 9+(5+0+(2+0+0))+(1+0+0) = 17

product(node(12,empty,empty)) = 12*1*1 = 12

product(node(20,node(6,empty,node(4,empty,empty)),node(7,empty,empty)))

= 20*(6*1*(4*1*1))*(7*1*1) = 3360
```

Complete the Flex and Bison files so that when they are compiled and run, it will give the following output on the **input**:

```
sum(node(4,empty,empty))
4
sum(node(9,node(5,empty,node(2,empty,empty)),node(1,empty,empty)))
17
product(node(12,empty,empty))
12
product(node(20,node(6,empty,node(4,empty,empty)),node(7,empty,empty)))
3360
```

Flex file "tree.lex":

```
%option noyywrap
% {
struct treenode
 int sum;
int product;
};
#define YYSTYPE treenode
#include "tree.tab.h"
용}
num [0-9][1-9][0-9]+ (4pts)
op [(),\n]
ws [ \t]+
응응
           { <u>yylval.sum = atoi(yytext); yylval.product = atoi(yytext);</u> (4pts)
{num}
              return NUM; }
{qo}
              return *yytext;
\{\mathtt{ws}\}
empty return EMPTY;
sum return SUM;
product return PRODUCT;
node
응응
```

Bison file "tree.y":

```
#include <iostream>
using namespace std;
struct treenode
 int sum;
int product;
};
#define YYSTYPE treenode
int yylex(void);
int yyerror(const char*);
왕}
%token NUM
%token EMPTY
%token SUM
%token PRODUCT
응응
/* Fill in the blanks in the grammar rules and actions*/
input: /* empty */ | input line ;
line: '\n'
|SUM '(' tree ')' '\n' { cout << $3.sum << endl;} (1pt)
| PRODUCT '(' tree ')' '\n' { cout << $3.product << endl;}; (1pt)
tree:
EMPTY { \$\$.SUM = 0; \$\$.PRODUCT = 1;} (4pts, 1pt for grammar, 3pts for
| '(' NUM ',' tree ',' tree ')' { $$.sum = $2.sum + $4.sum + $6.sum; $$.product =
$2.product * $4.product * $6.product;}; (6pts, 3pts for grammar, 3pts for action)
응응
int main() { return yyparse();}
int yyerror(const char* s) {
      cout << "error" << endl;</pre>
      return 0;
}
```

Problem 6. Parameter Passing Methods (10 Points)

The following program is in an imaginary D language, which has a syntax similar to the C language, but can apply static or dynamic scoping as well as various parameter passing methods as we specify. Determine the output of the following D program with each specified scoping and parameter passing method:

Static scoping, call by value:

```
(3,2)(2,2)(4,2) (2.5pts)
```

Static scoping, call by reference:

```
(3,3)(2,2)(1,2) (2.5pts)
```

Static scoping, call by value-result:

```
(3,2)(2,2)(1,3) (2.5pts)
```

Dynamic scoping, call by name:

```
(2,3)(1,2)(1,2) (2.5pts)
```

(* 0.5 pts are deducted for each incorrect value, at most 2.5 pts are deducted for each blank. *)

(* 0.5 pts in total are deducted for missing parentheses and comma. *)

Problem 7. Activation Records (10 points)

Recall that the C language by default uses static scoping on variable names and passing-by-value for parameter passing in procedure calls. Complete the activation records, including the variables and their values if known, the parameters and their values if known, and the control links for the following C program at specified point in time:

- (i) right before calling main;
- (ii) right before calling mul;
- (iii) right before exiting the call of mul;
- (iv) right before exiting main;
- (v) right after exiting main and before the program terminates.

```
#include <stdio.h>
int x = 3;
int y = 4;
int a[2] = \{5, 6\};
void mul(int x, int b, int z[ ])
      int i=0, size=2;
      for (i=0; i < size; i++) {
            a[i] *= x++;
            z[i] *= b++;
      }
      y=x*b;
}
int main()
{
      int y = 5;
      int b[2] = \{2, 3\};
      mul(y, x, b);
```

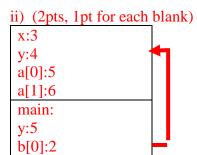
```
i) (1pt)

x:3

y:4

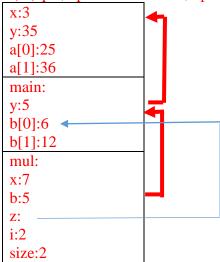
a[0]:5

a[1]:6
```

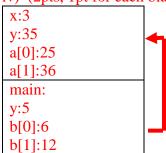


b[1]:3

iii) (4pts, 1pt for each blank, 1pt for the control link)



iv) (2pts, 1pt for each blank)



v) (1pt) x:3 y:35 a[0]:25 a[1]:36