Developer Manual for TreeScaper

Zhifeng Deng

July 30, 2020

1 Installation

CLVTreeScaper requires a CLAPACK properly installed and linked on your machine. CLAPACK-3.2.1 has been attached to this repository. You may also download here.

See detailed instruction on using BLAS library optimized for your machine in CLA-PACK/README.install at step (4).

For a fast default installation, you will need to

- Clone TreeScaper repository from GitHub	(see step 1 below)
- Relocate CLAPACK-3.2.1 and modify CLAPACK make.inc file	(see step 2 below)
- Modify TreeScaper makeCLVTreeScaper.inc file	(see step 2 below)
- Make CLAPACK library	(see step 3 below)
- Make CLVTreeScaper binary	(see step 3 below)

Procedure for installing CLAPACK.

(1) git clone -b zdver https://github.com/TreeScaper/TreeScaper.git to build the following directory structure:

TreeScaper/README.install this file

TreeScaper/makeCLVTreeScaper.inc compiler, compile flags and library

definitions for TreeScaper.

 $\label{eq:clapack-3.2.1} TreeScaper/CLAPACK-3.2.1/ \qquad \qquad CLAPACK \ attached \ in \ TreeScaper.$

TreeScaper/CLAPACK-3.2.1/make.inc compiler, compile flags and library

definitions, for TreeScaper.

(2) Move /CLAPACK-3.2.1 outside TreeScaper and modify /CLAPACK-3.2.1/make.inc. For default installation, you need to only modify the OS postfix name PLAT in /CLAPACK-3.2.1/make.inc. For advanced installation, please refer to /CLAPACK-3.2.1/README.install

Update the path of CLAPACK: CLAPPATH in makeCLVTreeScaper.inc and make sure the OS postfix name is consistent with CLAPACK setting, i.e. PLAT in /CLAPACK-3.2.1/make.inc and in makeCLVTreeScaper.inc must be the same.

(2)' If there is a CLAPACK already built in your machine. Make sure it has the following directory structure:

CLAPACK/BLAS/ C source for BLAS

CLAPACK/F2CLIBS/ f2c I/O functions (libI77) and math functions (libF77)

CLAPACK/INSTALL/ Testing functions and pre-tested make.inc files

for various platforms.

CLAPACK/INCLUDE/ header files - clapack.h is including C prototypes

of all the CLAPACK routines.

CLAPACK/SRC/ C source of LAPACK routines

Update the path of CLAPACK: CLAPPATH in makeCLVTreeScaper.inc and check the OS postfix name of lapack_XXX.a and blas_XXX.a and modify PLAT in makeCLVTreeScaper.inc

For example, if the naming is lapack_MAC.a and lapack_MAC.a then, modify

PLAT = _LINUX

in makeCLVTreeScaper.inc. If the naming is lapack.a and blas.a, modify

PLAT =

in makeCLVTreeScaper.inc.

(3) Go to TreeScaper directory. To install the CLAPACK, run make CLAPACK

To compile the TreeScaper, run make or make CLVTreeScaper.

You may move the binary CLVTreeScaper to other location for your, convenience. Make sure you also move the default parameters files nldr_parameters.csv and dimest_parameters.csv to maintain the structure:

/CLVTreeScaper the CLVTreeScaper binary parameters for nldr routines

/dimest_parameters.csv parameters for dimension estimation routines

To update the 'zdver' GitHub branch,

- 1) Keep your customized makeCLVTreeScaper.inc file.
- 2) Run git pull
- 3) If the makeCLVTreeScaper.inc got overwritten, restore your customized version.
- 4) If there is no change on CLAPACK side, which is usually the case, run make CLVTreeScaper to get the new binary.

Warning: you are not suggested to comment any local modification on this branch.

2 Command line arguments list

The command line version binary file accept a long list of arguments for particular tasks, especially for complicated tasks on trees. Figure 1 shows the structure of keywords <code>-dimest</code>, <code>-nldr</code> and <code>-trees</code> with their outputs and possible dependency between these output. In this figure, solid arrowed line represents output file of routine and dashed arrowed line represents possible dependency between files and routines.

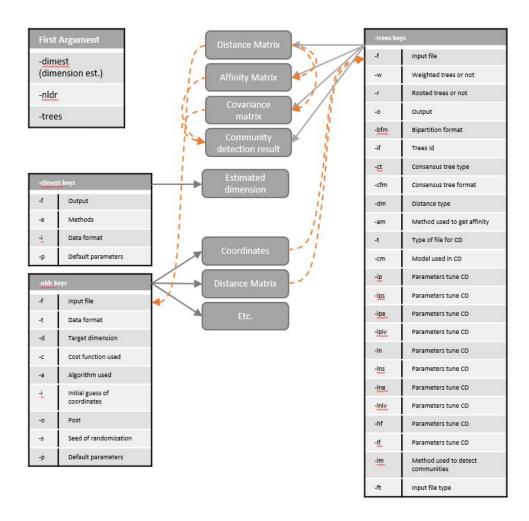


Figure 1: Argument List and routine structure of the current version.

TreeScaper binary takes a long command to execute a complicated task like performing community detection on a set of trees. The typical procedures TreeScaper is going to perform are

- 1. compute bipartition matrix of trees;
- 2. compute trees distance matrix and affinity matrix or compute bipartition covariance matrix;
- 3. perform community detection methods.

These procedures are specified with one command like

./CLVTreeScaper -trees -f test.tre -ft trees -r 1 -w 1 -o Community -t Affinity -dm RF -cm CNM -lm auto

and produce bipartition matrix, affinity/covariance matrix and CD results, which is what the dashed line between output files from keyword -trees. Since these dependency happen internally in one command, we have no control over the intermediate results, i.e., modifications like manually setting threshold for affinity before CD is not possible.

Such commands are inconvenient and also difficult to cooperate with the GUI system we are building on CloudForest. Therefore, keyword and command structure shown in Fig.2 is proposed, while the old structure will remain in the TreeScaper binary until the new system become reliable.

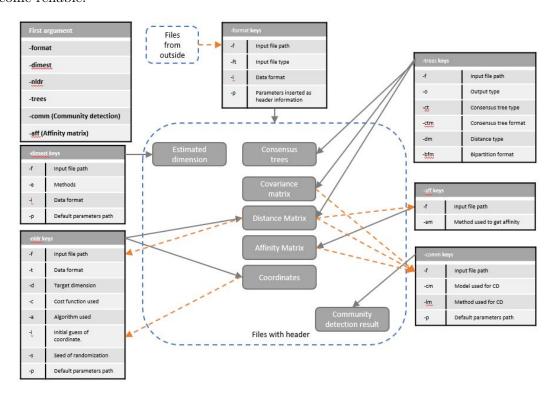


Figure 2: Argument List and routine structure of the new version.

New key arguments <code>-format</code>, <code>-comm</code> and <code>-aff</code> are added to TreeScaper. The new structure is centered at files with header information (in order to simplify file names and shorten argument list). <code>-format</code> takes care of converting files from other source with different format to files consistent with TreeScaper.

-aff and -comm which stand for affinity and community detection are separated from -tree, which now only takes care of computing bipartition matrix, distance matrix and covariance matrix from tree file. Users are now able to perform any modification on output files from -tree, for example, send them to -nldr. Then -comm will take these (modified) files and other necessary information to perform CD methods.

2.1 -dimest arguments

2.2 -nldr arguments

The key -nldr takes in a distance matrix of n points $p_i, i=1,\cdots,n$ on certain metric space and return coordinates of p_i , lies on Euclidean space \mathbb{R}^k along with other useful information. In particular, the coordinates matrix \tilde{C} of \tilde{p}_i stores in NLDR_COR.out and the Euclidean distance matrix \tilde{D} stores in NLDR_DIS.out.

Arguments	Description	options
-f	Source file	
-t	Data type	COR, DIS
-d	Target dimension, k	
-с	Cost function	CCA, SAMMON, CLASSIC_MDS
		KRUSKAL1, NORMALIZED
-a	Algorithm	STOCHASTIC, METROPOLIS, GAUSS_SEIDEL
		LINEAR_ITERATION, MAJORIZATION
-i	Initial guess	RAND, CLASSIC_MDS
-0	Output postfix	
	Parameter file	

2.3 -trees arguments

The key -trees is able to accomplish multiple tasks:

- 1. Distance: it takes in n trees and return a distance matrix $D \in \mathbb{R}^{n \times n}$
- 2. Affinity: it takes in n trees or a distance matrix D and return an affinity matrix A computed from D or certain distance matrix from the input trees set.
- 3. Consensus tree: it takes in n trees and return the consensus tree from input trees set.
- 4. Covariance: it takes in n trees and return the covariance matrix $Cov \in \mathbb{R}^{m \times m}$ of m bipartition appeared in the trees set.
- 5. Community: it takes in affinity matrix or covariance matrix and perform the community detection method.

Note that all tasks are coupled with each other if there is mathematical dependence. For example, when -o Affinity appears, -trees computes and outputs bipartition matrix, distance matrix and affinity matrix; when -o Community -t Covariance appears, -trees computes and outputs bipartition matrix, covariance matrix and community detection results.

ì
ity, Consensus
e, Community
ĹX
Strict
ick
t, SPR
RNM, NNM
Covariance
i e

2.4 -aff arguments

The key -aff takes in any distance matrix D and return affinity matrix A computed from D.

Arguments	Description	options
-f	Source file	
-am	Affinity type	Rec, Exp

2.5 -comm arguments

2.6 -format arguments

3 Basic data structures

Most of the basic data structures are constructed by Wen Huang. They includes basic array, matrix, string, mapping and file stream. They are, basically the c++ built-in structure warped up with convenient functions and operators. For example, the matrix class integrates singular value decomposition from CLAPACK. These data structures' header file and implementation files are prefixed with "w".

There are other more complicated data structures for specific algorithm and mathematical objects such as trees and community. They will be addressed in the next section.

3.1 Array

Members of array of type T are consisted of a pointer of T^* vec and a static integer length that indicates the length. The member functions and operators are given below.

1. friend std::istream &operator>

This operator does nothing and will not assign value to the array from the istream. According to particular needs of reading data, this may later be implemented with actual reading functionality.

2. friend std::ostream &operator«

This operator will output the length and its components separated by ":" and ",".

Example: an array of char[] from "a" to "e" is outputted in the format of

$$\{5: a, b, c, d, e\}$$

3. const Array& Array::operator=(const Array &right)

The assignment operator will free the pointer vec on the left and allocate a new vec. Then it assigns values from right hand side to left hand side component-wise.

The operator returns a pointer of the array on the left.

4. const Array& Array::operator+=(const Array &right)

This operator creates a new array with the right hand side attached to left hand side. It allocates Array of the correct length and then assigns values accordingly. Then call the assignment operator = to overwrite the current array.

Warning: calling = costs repeated and unnecessary copy-pasting.

5. const Array& Array::operator-=(const Array &right)

This operator creates a new array from the left hand side, with every component presented in the right removed and calls assignment operator to overwrite the current Array.

Warning: calling = costs repeated and unnecessary copy-pasting.

Warning: the new array is built incrementally and calls **resize** everytime, will brings the complexity to $O(length^2 \times right.length)$ other than $O(length \times right.length)$.

6. bool Array::operator==

Compares two array component-wise after comparing the length.

7. bool Array::operator<

The logic of the compare operator is to set the array with component-wise greater component to be the greater one. If the lengths are different, only compare the first k components where k is the smaller length. If the first k components happens to be the same (component-wise), the longer Array is greater than the shorter one.

8. Array Array::operator(const int index, const int end)

This operator extract sub-array from the current array. It takes two indices as parameters and return a new array that has the values from index to end.

Warning: this implementation is different than the implementation of String::operator(), which also takes two integers as parameters but the first one is the starting index and the second one is the length of the sub-string, instead of the ending index.

3.2 Matrix

Members of a matrix of type T are consisted of a pointer of pointers T** implemented inrow-major and two static integers row and col which indicate the dimensions of the matrix. This class also calls classic linear algebra algorithms from CLAPACK.

Overloaded operators are given below.

1. friend istream & operator ».

Warning: This operator does nothing, i.e., it does not assign values from the input stream.

2. friend ostream & operator «.

This operator output the matrix in the format of

```
{ (3,2) a, b, c d, e, f }
```

3. friend Matrix operator+(Matrix<T> left, Matrix<T> right).

This operator resize the left and right matrices to the lager dimension by calling member function <u>resize</u> and then create a new matrix and assign values from entry-wise addition.

Warning: the resize of left and right is silent here, which will permanently change the matrix being summed.

4. friend Matrix operator+(Matrix<T> left, S right).

This operator accepting a number right of type S on the right will add right entry-wise to each element in Matrix left, i.e., shift the matrix by right.

5. friend Matrix operator+(S left, Matrix<T> right)

Shift the matrix by left entry-wise.

6. friend Matrix operator-

See friend Matrix operator+.

- 7. Matrix operator*(const Matrix<T> &left, const Matrix<T> &right)
 Implement matrix multiplication.
- 8. friend Matrix operator*(S value, Matrix<T> mat) or (Matrix<T> mat, S value).

This operator return a new matrix entry-wise scaled by value. Note that this operator does not change the original matrix.

Warning: the matrix getting rescaled should be passed by reference in order to avoid construction/destruction computation.

9. friend Matrix operator/(Matrix<T> mat, S value)

This operator return a new matrix entry-wise divided by value. Note that this operator does not change the original matrix. Also note that this is not the syntax used in some advanced language where A/B means $B^{-1}A$.

Warning: the matrix getting rescaled should be passed by reference in order to avoid construction/destruction computation.

10. T &operator()(const int r, const int c = 0)

This operator returns the entry at r-row and c-column.

Important member functions are given below.

1. Matrix<double> compute_scalar_product_matrix()

This function return a **double** type matrix $S \in \mathbb{R}^{n \times n}$ from the current matrix $D^{(2)} \in \mathbb{R}^{n \times n}$. S is the centering-scaled $D^{(2)}$, which is assumed to be a squared distance matrix of n points $\{p_i\}_{i=1,\dots,n}$, $D^{(2)}_{ij} = D^{(2)}_{ji} = d^2(p_i, p_j)$.

Note that the classical multidimensional scaling method, MDS assumes these n points lie on some Euclidean space \mathbb{R}^k equipped with classic 2-norm distance. And S is a squared distance matrix of transformed n points such that the arithmetic mean of new

points is $0^k \in \mathbb{R}^k$. Also note that the arithmetic mean in Euclidean space is also the Karcher mean defined by

$$\arg\min_{m\in\mathbb{R}^k}\sum_{i=1}^n d^2(p_i,m).$$

The formula of computing S is given by

$$S = -\frac{1}{2}JD^{(2)}J$$

where $J = I - \frac{1}{n} \mathbf{1} \mathbf{1}^T$ and $\mathbf{1} \mathbf{1}^T$ is all-1 matrix. Also note that the transformation $D^{(2)} \to S$ preserves the solution of MDS, i.e.,

$$\arg\min_{B \in \mathbb{R}^{n \times k}} \left\| BB^T - D^{(2)} \right\|_F^2 = \arg\min_{B \in \mathbb{R}^{n \times k}} \left\| -\frac{1}{2} JBB^T J + \frac{1}{2} JD^{(2)} J \right\|_F^2$$

where B is the Euclidean coordinate matrix of n points in \mathbb{R}^k which generates (approximately) the squared distance matrix $D^{(2)}$.

Error: The invariance of transformation seems to be only true for Euclidean space. For more abstract $D^{(2)}$ generated from more general metric on Riemannian manifold, the scaling $-\frac{1}{2}JD^{(2)}J$ does not preserves positive definiteness of the squared distance matrix, which further causes negative eigenvalues in the following PCA, the eigendecomposition, process.

Warning: The squared distance matrix $D^{(2)}$ used in here is inconsistent with the distance matrix D computed in compute_Distance_Matrix in difference of entry-wise squared or not.

2. Matrix<double> compute_Distance_Matrix()

This function return a **double** type matrix $D \in \mathbb{R}^{n \times n}$ computed from the current matrix $M \in \mathbb{R}^{n \times k}$. M is consider as coordinate matrix of n points in k-dimensional Euclidean space, i-th row represents the k-tuple Euclidean coordinates of a point p_i . And D is the distance matrix where $D_{ij} = D_{ji} = d(p_i, p_j) = ||p_i - p_2 j||_2$ is the 2-norm distance between points p_i, p_j .

Warning: the resulting distance matrix D is dense, the symmetric structure is not exploited here.

4 Algorithms

4.1 Nonlinear dimensional reduction(NLDR)

This part collects algorithms and their important subroutines implemented in TreeScaper. The main goal in these algorithms is that given a squared distance matrix $D^{(2)}$ or distance matrix D of n points $\{p_i\}_{i=1,\dots,n}$ on some metric space, find n points $\{p_i'\}_{i=1,\dots,n} \subset \mathbb{R}^k$, such that the distance matrix D' for \mathbf{p}' approximates D the best, under the cost functions defined in different algorithm. Note that these n new points \mathbf{p}' can be represents by the coordinate matrix $B \in \mathbb{R}^{n \times k}$ which is often used as the output of these NLDR algorithms.

4.1.1 Classical Multidimensional scaling(MDS)

.

Classical Multidimensional scaling assumes $D^{(2)}$ is generated from Euclidean space with typical vector 2-norm as distance. The implemented algorithm NLDR::CLASSIC_MDS contains 4 parts:

- Compute centered matrix S from the given tree distance matrix D by calling <u>Matrix</u>::compute_Scalar_Matrix.
- 2. Perform singular value decompositions(SVD) to S by calling Matrix::SVD_LIB to obtain

$$S = U \Sigma V^T$$
.

Note that since S is symmetric, there exist eigen-decomposition $S = Q\Lambda Q^T$, i.e., there exists a signature matrix E, which has only 1 or -1 in diagonal and 0 elsewhere, such that $U\Sigma V^T = Q\Lambda EEQ^T = Q(\Lambda E)(QE)^T$ and U = Q, $\Sigma = \Lambda E$ and V = QE. This implies eigen-decomposition from CLAPACK is more efficient.

Also note that if $D^{(2)}$ uses vector 2-norm in Euclidean space, S is positive definite and SVD coincides with eigen-decomposition.

Warning: when there exist files named consistently that indicates SVD has been done and U, V, Σ has been stored, the routine will not do it again but simply read them from files. This is silent and could cause problem if those files are not actually inconsistent.

3. In case of performing MDS for $D^{(2)}$ from other metric space, which cause the presence of negative eigenvalues, it selects the k eigenvectors Q_{i_j} , $j=1,\dots,k$, where λ_{i_j} are the k most largest positive eigenvalues.

Error: memory leakage happens in this process whenever a negative eigenvalues encountered in the k most largest in magnitude eigenvalues. This problem is temporarily fixed but the theoretical explanation and necessity of this process is still needed. the classical MDS may not be suitable at all for Tree subjects.

4. Produce the coordinates matrix B for n points $\mathbf{p}' \subset \mathbb{R}^k$ by

$$B = \left[\sqrt{\lambda_{i_1}} Q_{i_1} \quad \cdots \quad \sqrt{\lambda_{i_k}} Q_{i_k} \right] \in \mathbb{R}^{n \times k}.$$

5. Compute the stress that estimate how good BB^T approximate $D^{(2)}$ by calling NLDR::CLASSIC_MDS_stress. Note that classical MDS do not need to compute the stress since B already minimized the stress function. However, since Tree space is not a Euclidean space with appropriate distance, the output B does not minimize the stress function.

For more information of MDS, see here.

Implementations of some routines

Data structure

1. Matrix

Description Row-major 2-dimensional array.

Member row Number of rows.

col Number of columns.

**matrix Pointers to each row.

Member function <u>resize</u> Change the dimensions.

2. Ptree

Description Index base array-type unweighted tree with adjacency matrix.

Member leaf_number

*parent Array of indices of the parent.

*lchild Array of indices of the right child.

*rchild Array of indices of the left child.

**edge Adjacency matrix.

Member function none

3. NEWICKNODE

Description Linked node pointed to its children and parent.

Member Nchildren Number of children.

label weight

*child List of children.

hv1 Hash value for unknown use.

hv2 Hash value that identifies the bipar-

tition.

bitstr Bit string that represents the leaves

contained in the (sub-)tree.

parent

Member function none

 $4. \ {\tt NEWICKTREE}$

Description A NEWICKNODE that represents the root.

Member root A NEWICKNODE.

Member function none

 $5. \ {\tt TreeOPE}$

Description Operation associated to one <u>NEWICKTREE</u>. Note that

most of the method are implemented in recursive pre-

order.

Member

Member function <u>loadnewicktree</u> Read <u>NEWICKTREE</u>.

loadnewicktree2Read NEWICKTREE.floadnewicktreeRead NEWICKTREE.loadnodeRead NEWICKTREE.loadleafRead NEWICKTREE.parsetreeRead NEWICKTREE.parsenodeRead NEWICKTREE.parseleafRead NEWICKTREE.

addchild Link child to the parent.

dfs compute hash Assigned hash values to all (sub-

)tree which identifies the structure

and therefore the bipartition.

<u>bipart</u> Store hash values in one big array for

computing RF distance.

<u>findleaf</u> Find a leaf by the

NEWICKNODE::label.

<u>normalizedTree</u> Lift a unrooted tree to a rooted tree.

<u>newick2lcbb</u> Convert <u>NEWICKTREE</u> to <u>Ptree</u> for

computing matching distance.

<u>newick2ptree</u> Implementation of <u>newick2lcbb</u>.

sumofdegree

<u>bipartcount</u> Count the occurrence of particular bi-

partition.

Addbipart Insert nodes to the current tree so

that there exist a (sub-)tree that con-

tains only a given set of leaves.

6. Trees

Description Multiple <u>NEWICKTREE</u>s with member function that com-

putes different distances.

 ${\bf Member}$

Member functioninitialTreesRead trees from file.ReadTreesRead trees from file.

compute_numofbipart

<u>Compute_Hash</u> Generate hash table for

computing hash values in a

tree.

<u>Compute_Bipart_Matrix</u> Generate a sparse matrix

that stores the weight of bipartition, its frequency of

occurrence.

 $\underline{ {\tt Compute_Bipart_Covariance}} \\ {\tt Generate} \quad {\tt the} \quad {\tt covariance}$

matrix according to the

formula.

Compute RF dist by hash Generate the RF-distance

matrix according to the for-

mula.

<u>pttree</u> Construct the adjacency

matrix of a Ptree.

<u>compute_matrix</u> Generate matrix for com-

puting matching distance by accumulating common edges from two <u>Ptree</u>s.

Compute Matching dist Compute the matching dis-

tance between two trees by the XOR table created from all possible bipartitions.

Compute Affinity dist

tance from the given dis-

tance matrix.

TreeOPE related routines.

1. TreeOPE::loadnewicktree.

${f Argument}$	(char *fname, int *error)
Description	Read tree from formatted string that stores biparti-
	tion. The implementation is given in floadnewicktree .
	Same level of the node is paired by "()" and separated
	by ",".
Complexity	
Memory space	
Associated routine	<u>floadnewicktree</u> Implementation by recursive process-
	ing the string in preorder.
Comments	This routine is better implemented by stack structure.
	It can only process unweighted tree. Also this routine
	takes the file name as input while the duplication version
	<u>loadnewicktree2</u> takes FILE type, customized fstream
	type. This routine seems to be insecure and redundant.
Error code	-1 Out of memory.
	-2 Parse error, the parentheses in string
	does not match.

2. <u>TreeOPE</u>::loadnewicktree2.

${f Argument}$	(FILE *fp, int *error)	
Description	Duplication version of los	adnewicktree but with cus-
	tomized fstream. Actual in	nplementation is not given in
	here, but in floadnewickt	ree
Complexity		
Memory space		
Associated routines	<u>floadnewicktree</u>	Implementation by recur-
		sive processing the string in
		preorder.
Comments	This routine also seems to	be redundant since the main
	thread of TreeScaper never	called it. There is another in-
	put routine parsetree, wh	ich can handle both weighted
	and unweighted tree, is use	ed in TreeScaper.
Error code	-1	Out of memory.
	-2	Parse error, the parentheses
		in string does not match.

$3. \ \underline{{\tt TreeOPE}}{\tt ::floadnewicktree}.$

	Argument	(FILE *fp, int *er	ror)
	Description	A pair of nodes are	created by <u>loadnode</u> when "(" is
	-	encountered.	,
	Complexity		
	Memory space		
	Associated routine	loadnode	
	Comments		as to be redundant since the main
	Comments		
		_	ever called it. There is another in-
		-	e, which can handle both weighted
		and unweighted tree,	_
	Error code		Out of memory.
			Parse error, the parentheses in string
		C	oes not match.
4.	TreeOPE::loadnode.		
	Argument	(FILE *fp, int *er	ror)
	Description	Create internal nodes	. When this function is called, a
		"(" has been read, if :	fp continue to read "(", next pair
		of nodes should be g	enerated, i.e., <u>loadnode</u> is called
		_	is encountered and <u>loadleaf</u> will
		· ,	s encountered, it is at the end of
		· · · · · · · · · · · · · · · · · · ·	des and should exit the routine to
		returned to previous l	
	Complexity	returned to previous i	ever of flode.
	Memory space		
	Associated routine	loadleaf	Add a loof and nature to provious
	Associated routine	loadleal	Add a leaf and return to previous level.
		11 1 1 1	
		addchild	Add the new pair of nodes to
			their parent.
		readlabelandweight	Read additional information
			from string.
	Comments	This is better impleme	nted by stack structure. Also note
		that this method read	leaves in preorder traversal.
	Error code	-1	Out of memory.
		-2	Parse error, the parentheses in
			string does not match.
5.	TreeOPE::parsetree.		
	Argument	(char *str. int *e	rror, NEWICKTREE *testtree)
	Description	Duplicate version of <u>f</u>	
	Complexity	Duplicate version of 1	Todanewickuice.
	Memory space		
	Associated routine	nargonodo	
		<u>parsenode</u> This is the routine use	nd in TracSanger
	Comments	This is the routine use	
	Error code		Out of memory.
		-2 H	Parse error, the parentheses in string
			loes not match.

6. <u>TreeOPE</u>::parsenode.

Argument	(FILE *fp, int *erro	r)
Description	Duplicated version load	lnode.
Complexity		
Memory space		
Associated routine	parseleaf	Add a leaf and return to previous
		level.
	addchild	Add the new pair of nodes to
		their parent.
	parselabelandweight	Read additional information
		from string.
Error code	-1	Out of memory.
	-2	Parse error, the parentheses in
		string does not match.

 $^{7. \ \}underline{\texttt{TreeOPE}}{::} \texttt{dfs_compute_hash}.$

Argument	(NEWICKNODE* st	artNode, LabelMap &lm,
-		ashrf, unsigned treeIdx,
	• –	str, unsigned long long
	•	g long m2, bool WEIGHTED,
	•	[_Taxa, map <unsigned long<="" th=""></unsigned>
	•	> *> &hash2bitstr, int
	numofbipartions)	
Description	It assigned hash val	ue to all leaves set, for internal node,
	the hash values are	computed by the sum of its children's
	hash values (and m	od m1 or m2). For each internal node,
	it determines a sub-	-tree rooted by itself from the current
	tree.	
		quely represented by the hash value
		ves contained in the subtree are also
	represented by the	bit string. For example, 01001100
	represents that the	subtree contains leaf 2, 5 and 6. The
	mapping from has	h values to the leaves it contain is
	stored in hash2bit	str.
Complexity		
Memory space		
Associated routine	<u>Array</u> ::SetBitArra	y Set the some positions, the index of
		leaves, of a bit array to 1.
	<pre>Array::OrbitOPE</pre>	OR operation of bit array, it realizes
		the functionality of making the bit
		string of the root having 1 in every
		leaf's index that the subtree has.
	add_of	Bit-wise addition for hash values.
Comments	Note that hash va	lue to subtree is bijection and sub-
	tree to leaves it co	ntains is subjection. Therefore, the
	mapping hash2bit	str is subjection. Also note that the
	operations, additio	n and modulus, on hash values are
	done in bit-wise ma	anner.
Error code	none	Terminate with specific error message

8. <u>TreeOPE</u>::bipart.

	Argument	(NEWICKNODE *const startnode, unsigned int
	_	&treeIdx, unsigned long long *matrix_hv,
		unsigned int *matrix_treeIdx, double
		*matrix_weight, int &idx, int depth, bool
		isrooted)
	Description	Store hash values, TreeIdx and weights in the given ar-
	Description	rays.
	Complexity	Tays.
	Memory space	
	Associated routine	
		NT
	Comments	Note that the "TreeIdx" is an identical array. Each tree
		will generate one set of such arrays and these arrays
		from different trees are pasted together and sorted by
		the hash values. By comparing hash values, identical
		bipartitions among different trees can be easily found.
	Error code	-1 Out of memory.
		-2 Parse error, the parentheses in string
		does not match.
9.	<pre>TreeOPE::findleaf.</pre>	
	Argument	(std::string leafname, NEWICKNODE
	_	*currentnode, NEWICKNODE *parent, int *icpt)
	Description	Find leaf leafname and return it. icpt also record which
	-	subtree under root the leaf lies in.
	Complexity	
	Memory space	
	Associated routine	none
10.	TreeOPE::normalizedTr	·ee.
	Argument	(NEWICKNODE *lrpt, NEWICKTREE *newickTree, int
		indexchild)
	Description	Lift a unrooted tree to a rooted tree.
	Complexity	
	Memory space	
	Associated routine	normalizedNode It's implementation.
11		

11. TreeOPE::newick2lcbb.

	Argument	<pre>(const NEWICKTREE *nwtree, int num_leaves, struct Ptree *tree)</pre>
	Description	Convert <u>NEWICKTREE</u> to <u>Ptree</u> , which is used to compute matching distance.
	Complexity	
	Memory space	
	Associated routine	newick2ptree Implementation of newick2lcbb .
	Comments	Note that <u>Ptree</u> does not stored hash values and weights, i.e., the bipartition and weight information are lost. Also note that the edges matrix of <u>Ptree</u> is not computed here.
12.	<pre>TreeOPE::sumofdegree.</pre>	
	Argument	(NEWICKNODE *node, bool isrooted, int depth)
	Description	Return the sum of degrees of all nodes.
	Complexity	
	Memory space	
	Associated routine	
	Comments	
	Error code	-1 Out of memory.
		-2 Parse error, the parentheses in string does not match.
13.	<pre>TreeOPE::bipartcount.</pre>	
	Argument	(NEWICKNODE *node, bool isrooted, map <unsigned long="" long,="" unsigned=""> &bipcount, int depth)</unsigned>
	Description	Count the occurrence of particular subtree, bipartition, by its hash value and store the result in the external mapping bipcount
	Complexity	
	Memory space	
	Associated routine	
	Comments	
14.	TreeOPE::Addbipart.	
	Argument	(NEWICKNODE* startNode, double freq, unsigned
		long long hash, Array <char> &bitstr, int</char>
		NumTaxa, bool &iscontained)
	Description	Given bitstr that represents a set of leaves. Insert
		internal nodes from leaf-set to root that collects those
		leaves lie in bitstr so that there is a subtree containing
	Complexity	exactly the same set of leaves in the resulting new tree.
	Memory space	
	Associated routine	none
	Comments	There is a better way to implement this functionality.
		There is a second way to implement this functionantly.

<u>Trees</u> related routines.

1. <u>Trees</u>::initialTrees.

Argument	(string fname)	
Description	Initialize a set	of <u>NEWICKEDTREE</u> s by calling
	<u>loadnewickedtree</u>	2. For Nexus trees, it only cre-
	ate a leaveslabel	smaps that stores the labels of leaf
	set.	•
Complexity		
Memory space		
Associated routine	<u>loadnewicktree2</u>	Create each tree.
Comments	Complicated string	operations are done here, which is
	unnecessary.	
Error code	-1	Out of memory.
	-2	Parse error, the parentheses in string
		does not match.
	-3	Failure of opening file.
Trees::ReadTrees.	-3	Failure of opening file.
Trees::ReadTrees. Argument	none	Failure of opening file.
	none	Failure of opening file. on of <u>initialTrees</u> except it calls
Argument	none A duplicated version	
Argument	none A duplicated version	on of <u>initialTrees</u> except it calls Newicked and NEXUS type of tree.
Argument	none A duplicated version parsetree for both	on of <u>initialTrees</u> except it calls Newicked and NEXUS type of tree.
Argument Description	none A duplicated version parsetree for both	on of <u>initialTrees</u> except it calls Newicked and NEXUS type of tree.
Argument Description Complexity	none A duplicated version parsetree for both	on of <u>initialTrees</u> except it calls Newicked and NEXUS type of tree.
Argument Description Complexity Memory space	none A duplicated versic parsetree for both Also lifted the tree	on of <u>initialTrees</u> except it calls Newicked and NEXUS type of tree. if it is unrooted.
Argument Description Complexity Memory space	none A duplicated version parsetree for both Also lifted the tree parsetree normalizedTree	on of <u>initialTrees</u> except it calls Newicked and NEXUS type of tree. if it is unrooted. Create each tree.
Argument Description Complexity Memory space Associated routine	none A duplicated version parsetree for both Also lifted the tree parsetree normalizedTree	on of <u>initialTrees</u> except it calls Newicked and NEXUS type of tree. if it is unrooted. Create each tree. Lift a unrooted tree. ring operations are done here, which
Argument Description Complexity Memory space Associated routine	none A duplicated versic parsetree for both Also lifted the tree parsetree normalizedTree Very complicated st	on of <u>initialTrees</u> except it calls Newicked and NEXUS type of tree. if it is unrooted. Create each tree. Lift a unrooted tree. ring operations are done here, which
Argument Description Complexity Memory space Associated routine Comments	none A duplicated version parsetree for both Also lifted the tree parsetree normalizedTree Very complicated strip is really unnecessar.	on of initialTrees except it calls Newicked and NEXUS type of tree. if it is unrooted. Create each tree. Lift a unrooted tree. ring operations are done here, which y. Out of memory.
Argument Description Complexity Memory space Associated routine Comments	none A duplicated versic parsetree for both Also lifted the tree parsetree normalizedTree Very complicated st is really unnecessary	on of initialTrees except it calls Newicked and NEXUS type of tree. if it is unrooted. Create each tree. Lift a unrooted tree. ring operations are done here, which y.

3. Trees::compute_numofbipart.

Argument	none		
Description	It computes the numbers of bipartition for all trees and		
	stores them in the array <u>number of bipartition</u> . The		
	formula is given by		
	s/2-n		
	where s is the sum of degrees and n is the number of leaf.		
Complexity			
Memory space			
Associated routine	ciated routine <u>sumofdegree</u>		

4. <u>Trees</u>::Compute_Hash.

Argument none					
_	Description				
		in a tree.			
_	Complexity				
	Memory space				
	Associated routine	dfs_compute_hash			
5. <u>]</u>	<u>'rees</u> ::Compute_Bipart_Matrix.				
-	Argument	none			
_	Description	The arrays of indivial tree's hashvalue, tree index and weight created from <u>bipart</u> were combined and sorted. Since the hash value represents the unique subtree structure, i.e a bipartition, the number of unique bipartion can be counted via checking the hash value. As a result, a sparse bipartition matrix that stores weight of unique bipartition versus trees is created.			
-	Complexity	Sipervisor versus v	1000 10 01000004.		
	Memory space				
-	Associated routine	<u>bipart</u>	Create arrays of hash values, weights		
	Associated Toutille	<u>bipart</u>	with tree index of one tree.		
		Sort	Sort the 3 arrays attached from all trees by the hash values, so that we can easily count the occurrence for each hash value, i.e., bipartition.		
		sort	Seems to be built-in sort for array that sort a temperate hash value array for certain later operation.		
_	Comments	The sort which is different then Sort is confusing here. Is it the default sort in c++?			
6. <u>]</u>	<u>Γrees</u> ::Vec_multiply.				
-		rgument (const double* Vec1, const do			
	Argument	(const double* V	ec1, const double* Vec2, int		
	Argument	(const double* Vousingue_idx)	ec1, const double* Vec2, int		
-	Argument Description				
-		Unique_idx)	natrix		
		Unique_idx)			
-		Unique_idx)	natrix		
-		Unique_idx)	natrix		
-	Description	Unique_idx)	natrix		
-	Description Complexity	Unique_idx)	natrix		
-	Description Complexity Memory space	Unique_idx) It return a rank-1 r	natrix		

 $7. \ \underline{\texttt{Trees}}{::} \texttt{Compute_Bipart_Covariance}.$

Argument	(bool ISWEIGHTED)				
Description	Compute the bipartition covariance matrix from the matrix, C, created by <u>Compute Bipart Matrix</u> , M. Let $M_1 = MM^T$, $v_1 = mean(M)$, $v_2 = sum(M)$, $M_2 = v2v1^T$ and $M_3 = v1v1^T$, then				
	$C = (M_1 - M_2 - M_2^T + n * M_3)/(n-1).$				
Complexity					
Memory space					
Associated routine	SparseMatrix::transpose				
	SparseMatrix::Multiply	Matrix-Matrix multiplication.			
	SparseMatrix::Mean	Matrix mean.			
	SparseMatrix::Multiply_ve				
		tion.			
	Trees::Vec_Multiply	Rank-1 matrix.			
Comments	Note that it is implemented				
	multiplication.				
Trees::Compute_RF_dis	·				
Argument	(bool ISWEIGHTED)				
Description	Compute the unweighted/weighted RF distance. For the unweighted distance, accumulate the number of each unique bipartition's occurrencein each tree, f_{ij} , and the number of bipartitions, n_i , then				
	unique bipartition's occurrer	ncein each tree, f_{ij} , and the			
	unique bipartition's occurrer	ncein each tree, f_{ij} , and the hen			
	unique bipartition's occurrer number of bipartitions, n_i , t	ncein each tree, f_{ij} , and the hen $\frac{n_j - 2f_{ij}}{2}.$ e complicated. The result is			
Complexity	unique bipartition's occurrent number of bipartitions, n_i , to $d_{ij} = \frac{n_i + 1}{2}$ For weighted case, it is more	ncein each tree, f_{ij} , and the hen $\frac{n_j - 2f_{ij}}{2}.$ e complicated. The result is			
Complexity Memory space	unique bipartition's occurrent number of bipartitions, n_i , to $d_{ij} = \frac{n_i + 1}{2}$ For weighted case, it is more	ncein each tree, f_{ij} , and the hen $\frac{n_j - 2f_{ij}}{2}.$ e complicated. The result is			
	unique bipartition's occurrent number of bipartitions, n_i , to $d_{ij} = \frac{n_i + 1}{2}$ For weighted case, it is more	ncein each tree, f_{ij} , and the hen $\frac{n_j - 2f_{ij}}{2}.$ e complicated. The result is			
Memory space	unique bipartition's occurrer number of bipartitions, n_i , t $d_{ij} = \frac{n_i + 1}{m_i}$ For weighted case, it is more stored in the matrix $dist_U$	ncein each tree, f_{ij} , and the hen $\frac{n_j - 2f_{ij}}{2}.$ e complicated. The result is			
Memory space Associated routine	unique bipartition's occurrer number of bipartitions, n_i , t $d_{ij} = \frac{n_i + 1}{m_i}$ For weighted case, it is more stored in the matrix ${\tt dist_U}$	ncein each tree, f_{ij} , and the hen $\frac{n_j - 2f_{ij}}{2}.$ e complicated. The result is			
Memory space Associated routine Comments Trees::pttree. Argument	unique bipartition's occurrent number of bipartitions, n_i , the $d_{ij}=\frac{n_i+1}{2}$. For weighted case, it is more stored in the matrix ${\tt dist_U}$ none none (struct Ptree *treeA, in	ncein each tree, f_{ij} , and the hen $\frac{n_j-2f_{ij}}{2}$. The complicated. The result is RF or dist_RF.			
Memory space Associated routine Comments Trees::pttree.	unique bipartition's occurrer number of bipartitions, n_i , t $d_{ij} = \frac{n_i + 1}{m_i}$ For weighted case, it is more stored in the matrix $dist_U$ none none	ncein each tree, f_{ij} , and the hen $\frac{n_j-2f_{ij}}{2}$. The complicated. The result is RF or dist_RF.			
Memory space Associated routine Comments Trees::pttree. Argument Description	unique bipartition's occurrent number of bipartitions, n_i , the $d_{ij}=\frac{n_i+1}{2}$. For weighted case, it is more stored in the matrix ${\tt dist_U}$ none none (struct Ptree *treeA, in	ncein each tree, f_{ij} , and the hen $\frac{n_j-2f_{ij}}{2}$. The complicated. The result is RF or dist_RF.			
Memory space Associated routine Comments Trees::pttree. Argument Description Complexity	unique bipartition's occurrent number of bipartitions, n_i , to $d_{ij} = \frac{n_i + 1}{n_i}$ For weighted case, it is more stored in the matrix $dist_U$ none none $(struct\ Ptree\ *treeA,\ information It constructs the edge matrix of the struct of the edge matrix of the edge mat$	ncein each tree, f_{ij} , and the hen $\frac{n_j-2f_{ij}}{2}$. The complicated. The result is RF or dist_RF.			
Memory space Associated routine Comments Trees::pttree. Argument Description	unique bipartition's occurrent number of bipartitions, n_i , to $d_{ij} = \frac{n_i + 1}{n_i}$ For weighted case, it is more stored in the matrix $dist_U$ none none $(struct\ Ptree\ *treeA,\ information It constructs the edge matrix of the struct of the edge matrix of the edge mat$	ncein each tree, f_{ij} , and the hen $\frac{n_j-2f_{ij}}{2}$. The complicated. The result is RF or dist_RF.			

10. <u>Trees</u>::compute_matrix.

8.

9.

Argument	<pre>(int *r, int range, struct Ptree *tree1, struct Ptree *tree2)</pre>			
Description	It accumulates the number common edges from two trees			
Description	and store in a vectorized matrix, r.			
Complexity	,			
Memory space				
Associated routine	none			
Comments	For n trees, there are $\binom{n}{2} = n(n-1)$ comparisons and this function will be called $n(n-1)$ times.			
<u>Trees</u> ::tree_mmdis.				
Argument	none			
Description	This distance is given by the solution of Hungarian algorithm of the cost matrix, r, given by compute matrix.			
Complexity				
Memory space				
Associated routine	array_to_matrix Recover r to a matrix.			
Comments	r is an $(k-3) \times (k-3)$ matrix where k is the number			
	of leaves. The main complexity goes into generating			
	distance matrix and running Hungarian algorithm.			
Trees::Compute_Matchi	ng_dist.			
Argument	none			
Description	The matching distance is given by the solution to Hungarian algorithm on the table with entries of number of XOR element in bitstrofatree, which are all possible bipartitions of one tree.			
Complexity	bipartitions of one tree.			
Memory space				
Associated routine	Get_bipartitionofonetree			
Comments	Line 1415 may have a bug.			
Trees::Compute_Affinity_dist.				
Argument	(String str_matrix, int type)			
Description	This routine compute the affinity distance, d_a , from the			
_ 02011P 01011	given distance d . The formula is either			
	$d_a = \frac{1}{\varepsilon_{rel} + d}$			
	or			
	$d_a = e^{-d},$			
	depending on the flag type. It accepts unweighted/weighted RF-distance, Matching-distance,			
	SPR-distance or distance given in file.			
Complexity	SPR-distance or distance given in file.			
Complexity Memory space	SPR-distance or distance given in file.			

14. <u>Trees</u>::temp.

Argument	none	
Description		
Complexity		
Memory space		
Associated routing	ne	
Comments		
Error code	-1	Out of memory.
	-2	Parse error, the parentheses in string
		does not match.