

# Map Coloring Theorems

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# Outline

1. Motivation
2. Duality
3. Coloring graphs
4. Coloring algorithm
5. Euler's Formula
6. Art Galleries

# Motivation

Definition.

A  $k$ -coloring of a map is an assignment of  $k$  colors, one to each country, in such a way that no two countries sharing a border have the same color.



A coloring of a map of Europe in thirteen colors

# Motivation

Definition.

A map is *k-colorable* if it has a coloring with at most  $k$  colors.



A coloring of a map of Europe in five colors

# Motivation

De Morgan (1982)

*A student of mine asked me today to give him a reason for a fact which I did not know was a fact - and do not yet. He says that if a figure be anyhow divided and the compartments differently coloured so that figures with any portion of common boundary line are differently coloured - four colours may be wanted, but not more - the following is the case in which four colours are wanted. Query cannot a necessity for five or more be invented.*

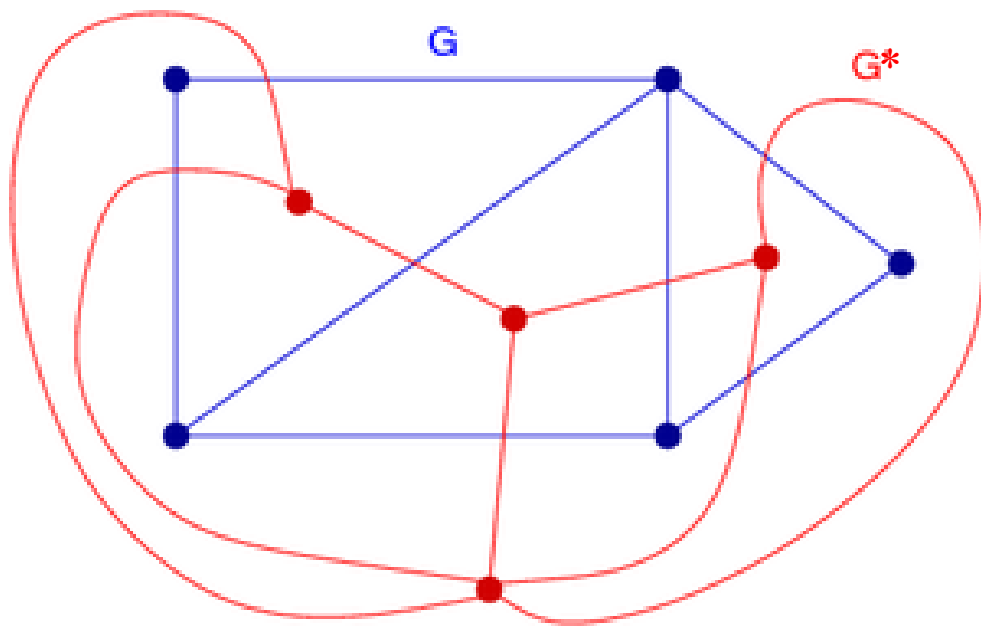
The 4-Color Conjecture (1852)

Is every map 4-colorable?

# Duality

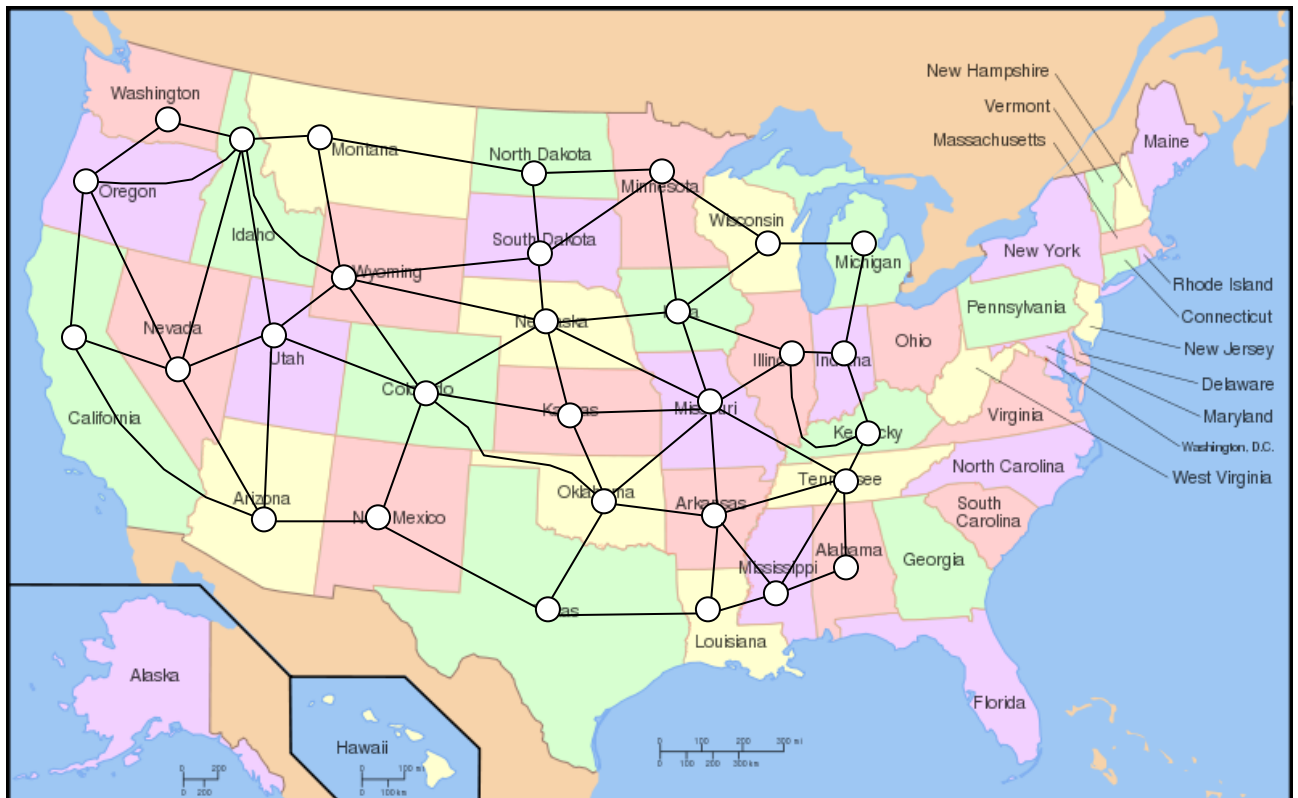
Definition.

The dual  $G^*$  of a map  $G$  is obtained by replacing each state with its capital and then putting an edge between two capitals if the states share a border.



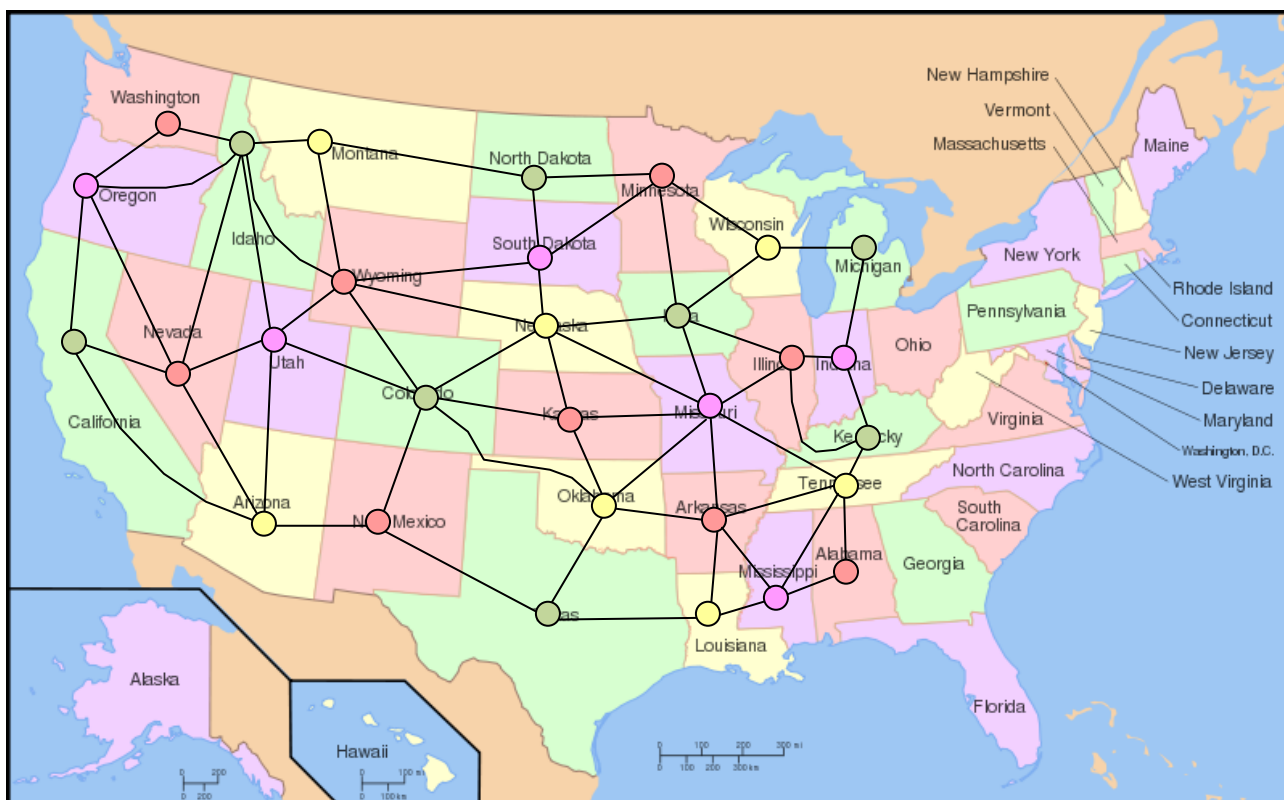
A simple example of duality

# Duality



The dual of the USA

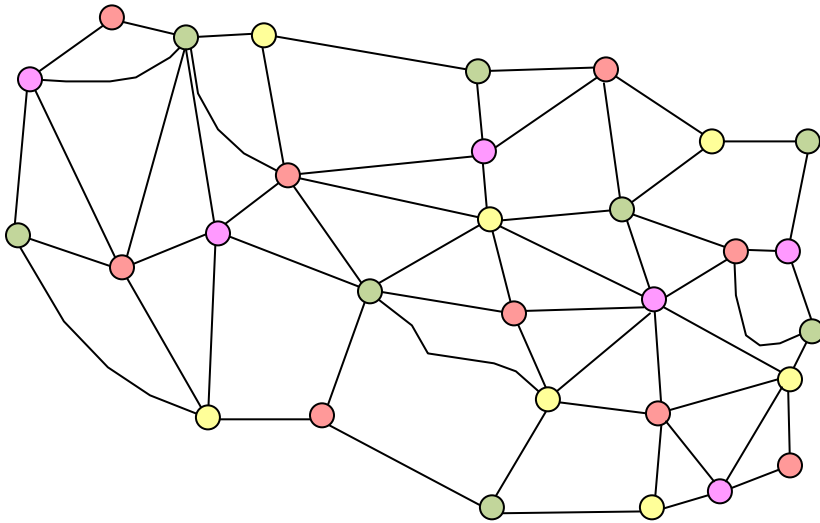
# Duality



The dual of the USA



# Duality

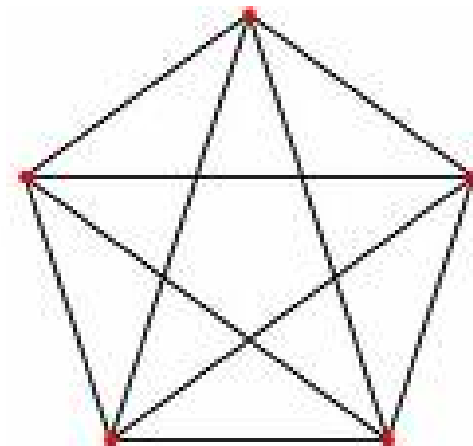
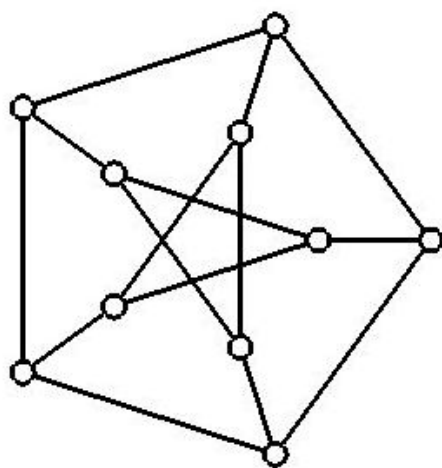


Coloring the dual of the USA

# Coloring graphs

Definition.

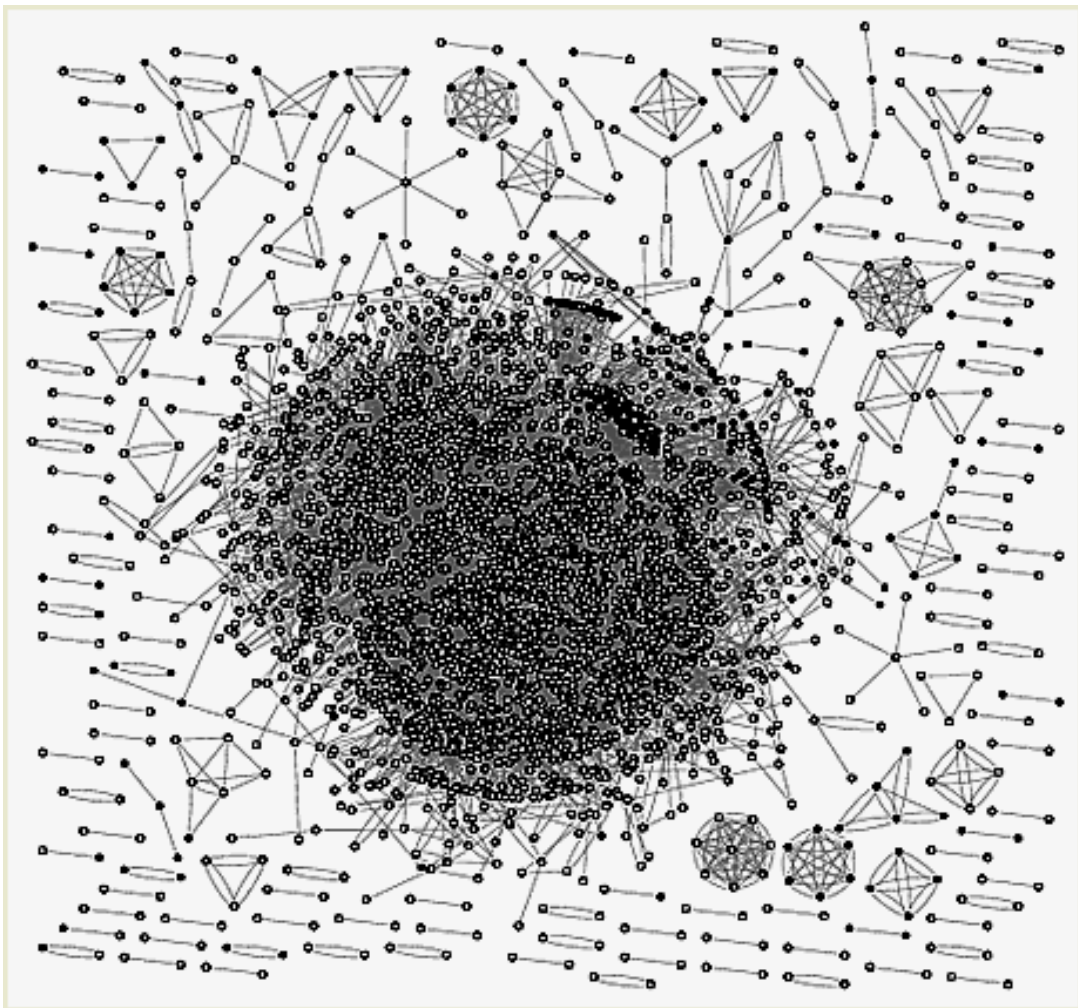
A **graph** consists of vertices together with edges joining pairs of vertices.



# Coloring graphs

Definition.

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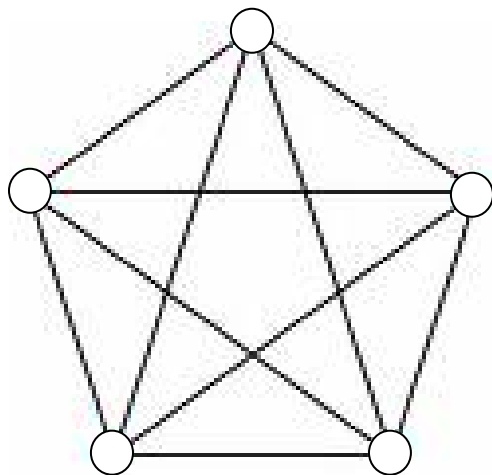


# Coloring graphs

Definition.

A  *$k$ -coloring* of a graph is an assignment of  $k$  colors, one to each vertex, in such a way that no two adjacent vertices share the same color.

The *chromatic number* of a graph  $G$  is the minimum  $k$  for which a  $k$ -coloring exists. It is written  $\chi(G)$ .

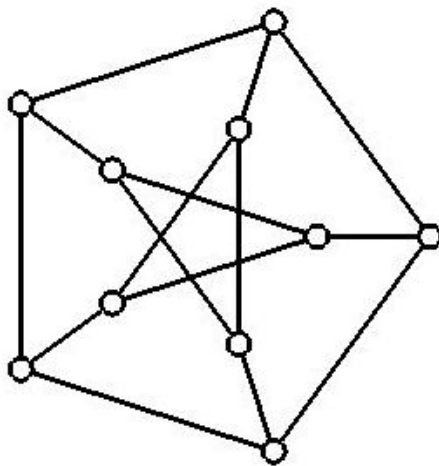


# Coloring graphs

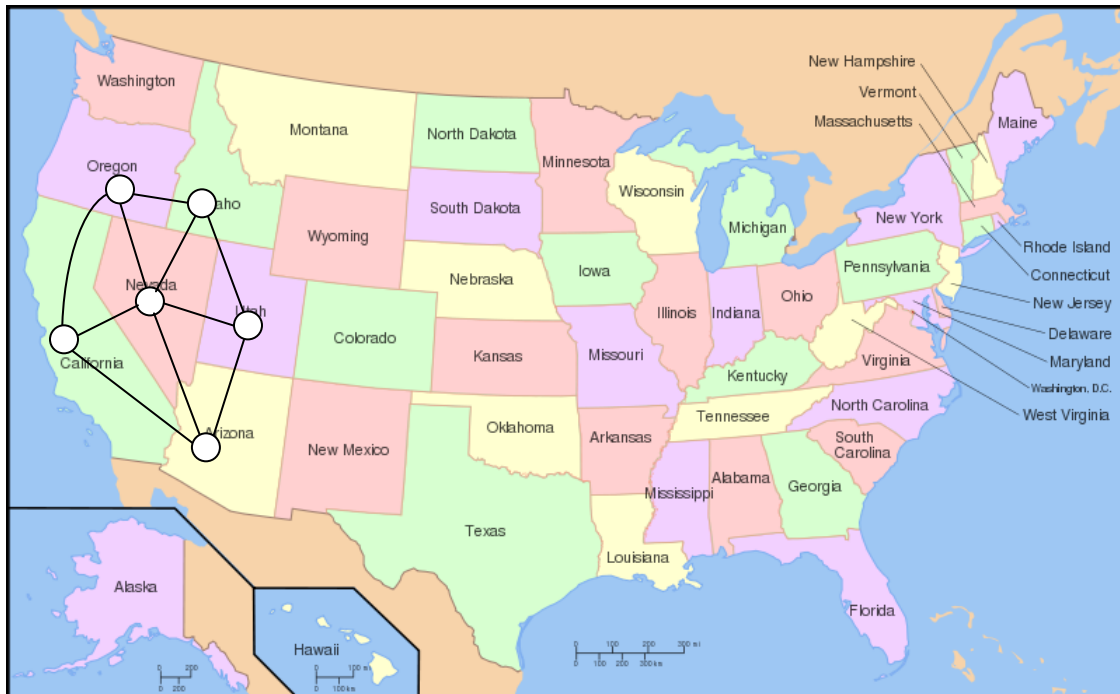
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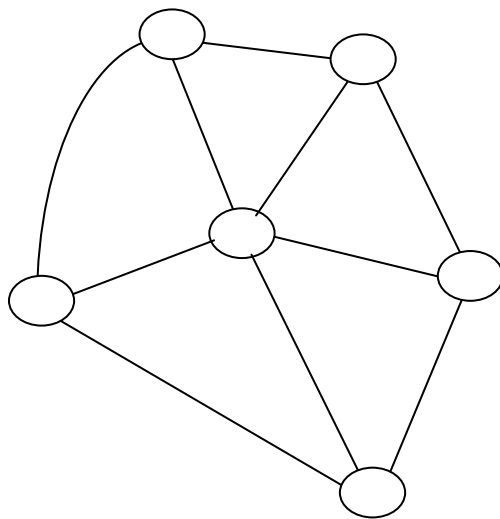
The *chromatic number* of a graph  $G$  is the minimum  $k$  for which a  $k$ -coloring exists. It is written  $\chi(G)$ .



# Coloring graphs



## A 4-coloring of the USA



# Coloring graphs

Observation 1.

*If a graph  $G$  contains a graph  $H$ , then*

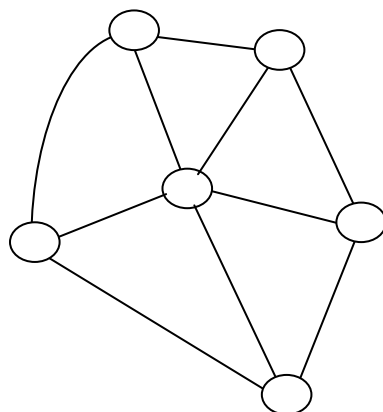
$$\chi(G) \geq \chi(H)$$

Therefore the chromatic number of the USA is 4.

# Coloring algorithm

Definition.

The **degree of a vertex** in a graph is the number of edges out of that vertex.



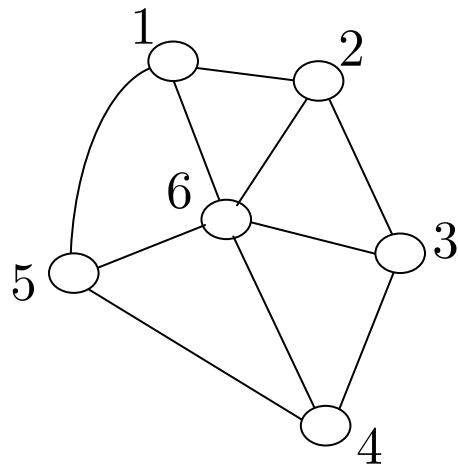
The minimum degree of a graph is the minimum of all the degrees of the vertices in the graph.



# Coloring algorithm

```
"
Uvctv"ykvj"c"itcrj"I3"ykvj"p"xgtvkegu0"
"
"    Hqt"k"? "3"vq"p"fq"
"
        30  Hkpf"c"xgtvgz"xk"qh"uocnnguv"
            fgitgg"kp"Ik0"
        40  Tgrnceg"Ik"ykvj"Ik"/"xk0"
"
Eqnqt"xp"ykvj"eqnqt"30" "
"
"    Hqt"k"? "p/3"vq"3"fq"
"
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```

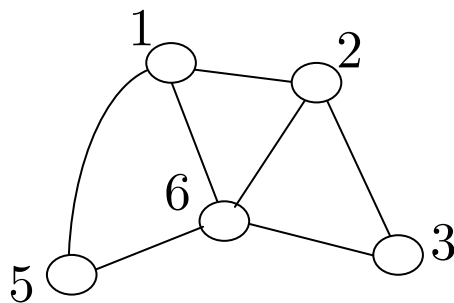
# Coloring algorithm



# Coloring algorithm

Pick a vertex of smallest degree, say 4.

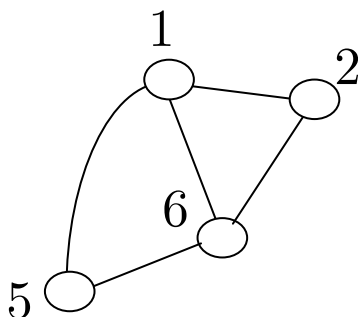
Delete it



# Coloring algorithm

Pick vertex of smallest degree, say 3.

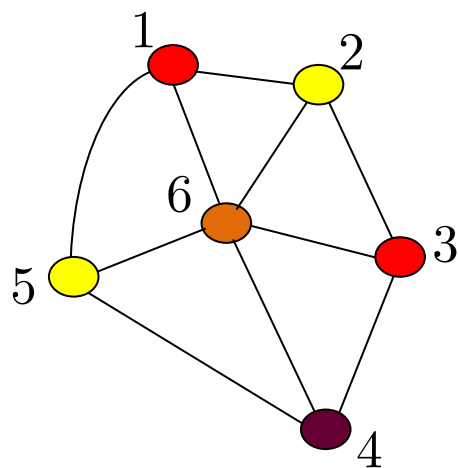
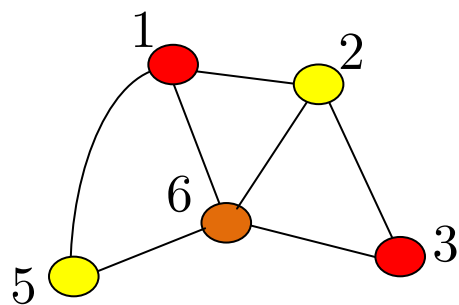
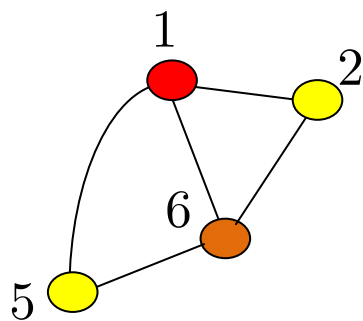
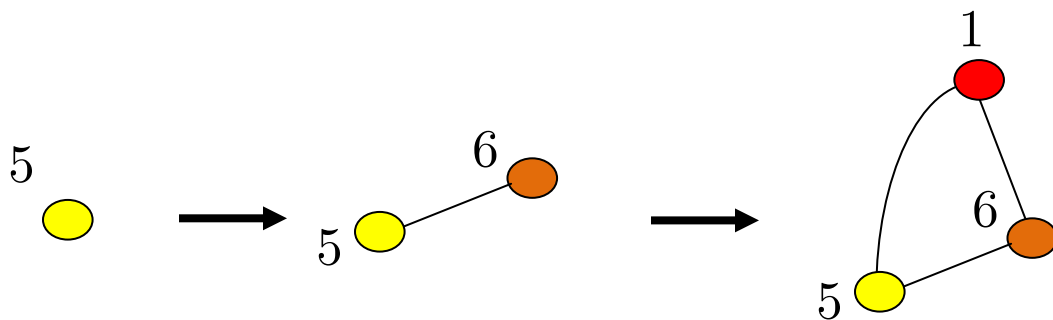
Delete it.



Pick vertex of smallest degree, say 2.

Delete it.

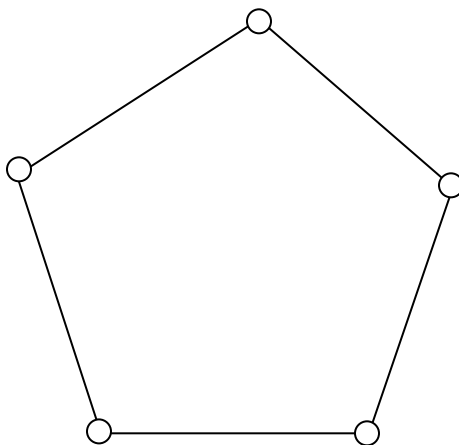
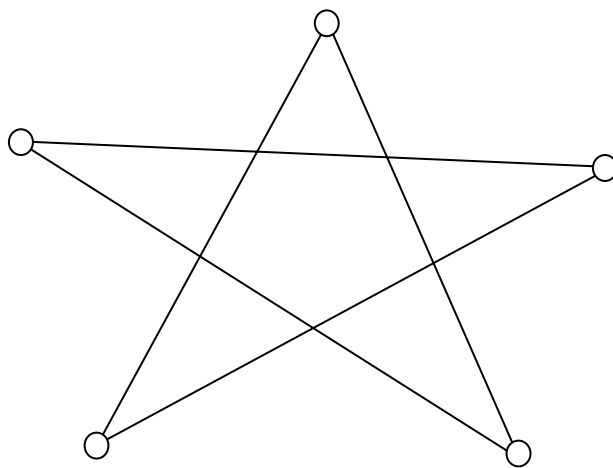
# Coloring algorithm



# Euler's Formula

Definition.

A **planar graph** is a graph which can be drawn in the plane without any two edges crossing.



# Euler's Formula

## Definition.

A face of a graph drawn without crossings is one of the lakes that you get when you pour water on the picture and treat the edges as walls.

# Euler's Formula

If  $G$  is a graph drawn without crossings in the plane, let  $f(G), n(G), e(G)$  be the number of faces, vertices and edges in the graph.

## Euler's Formula (1752)

*For any connected plane graph  $G$*

$$n(G) - f(G) + e(G) = 2$$



# Euler's Formula

How to prove Euler's Formula:

Basic idea

If you delete one edge from a cycle in the graph, then the number of faces drops by 1.

So continuing deleting edges in cycles,

$$n(G) - f(G) + e(G)$$

never changes until there are no cycles left.

But if there are no cycles left, then the graph is a tree. For a tree we have

$$e(G) = n(G) - 1 \qquad f(G) = 1$$

and so the value of the formula is 2.

# Euler's Formula

Definition.

The degree of a face surrounded by a cycle is the number of edges in that cycle.

Observation.

If  $d(F)$  is the degree of face  $F$  in a graph  $G$ , then

$$\sum_F d(F) = 2e(G)$$

This follows from the fact that if we add up the number of edges on the boundary of each face, then each edge of the graph is counted twice.

# Euler's Formula

## Theorem (Plato, <360bc)

*There is only one connected planar graph where all faces are quadrilaterals and all vertices have degree three.*

How to prove Plato's Theorem:

Euler's Formula says

$$n(G) - e(G) + f(G) = 2$$

Adding up degrees of faces says

$$4f(G) = 2e(G)$$

Adding up degrees of vertices says

$$3n(G) = 2e(G)$$

# Euler's Formula

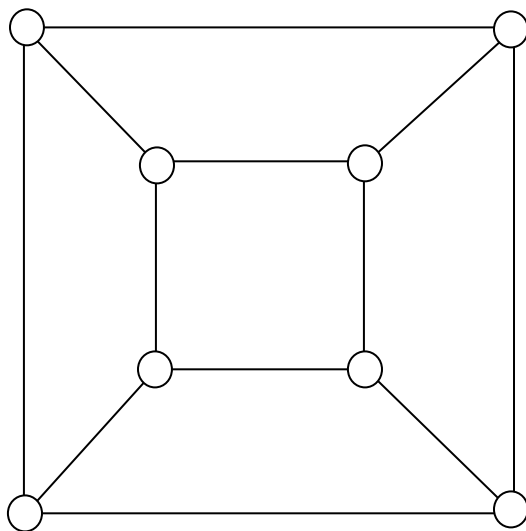
The only solution to the equations

$$n - e + f = 2$$

$$4f = 2e$$

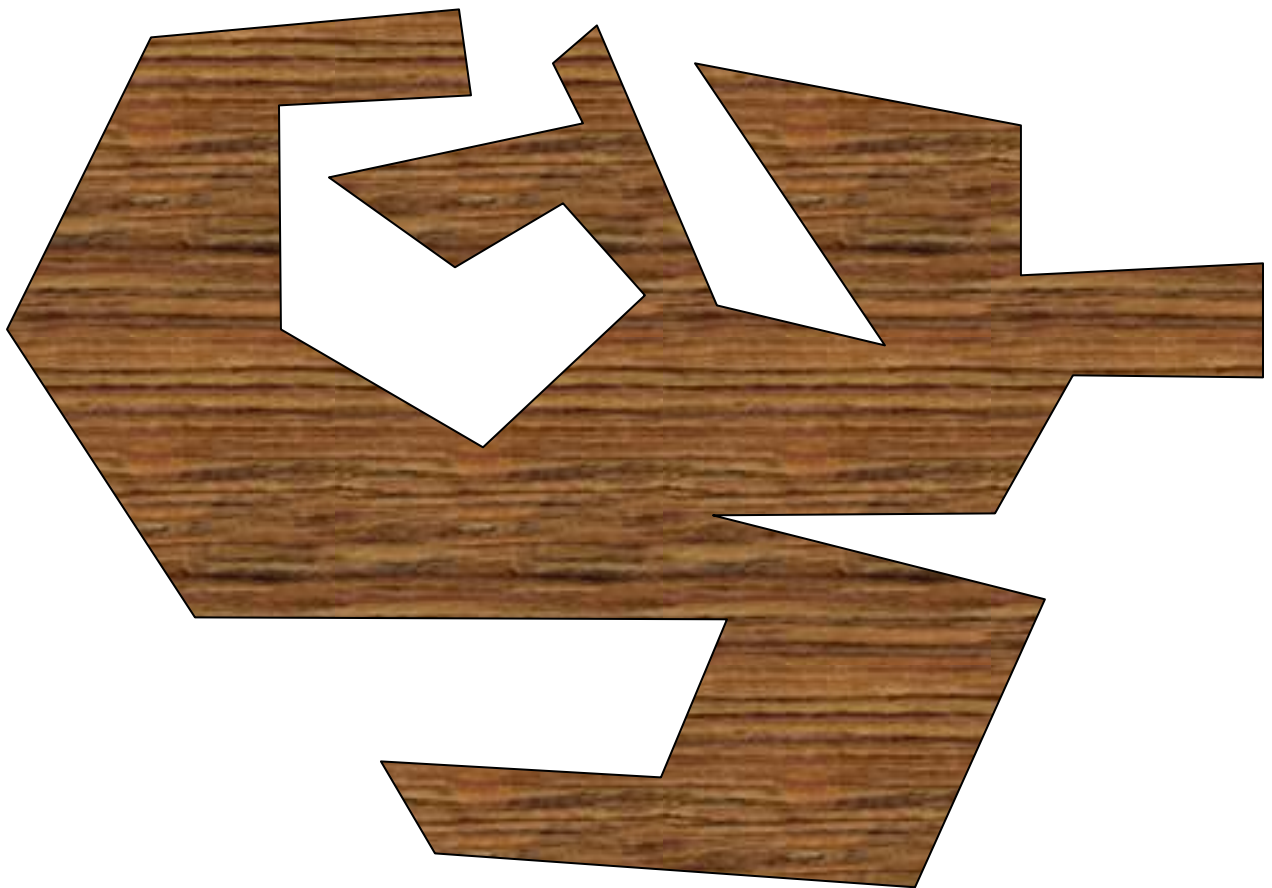
$$3n = 2e$$

is  $n = 8, e = 12, f = 6$ . We recognize the cube:



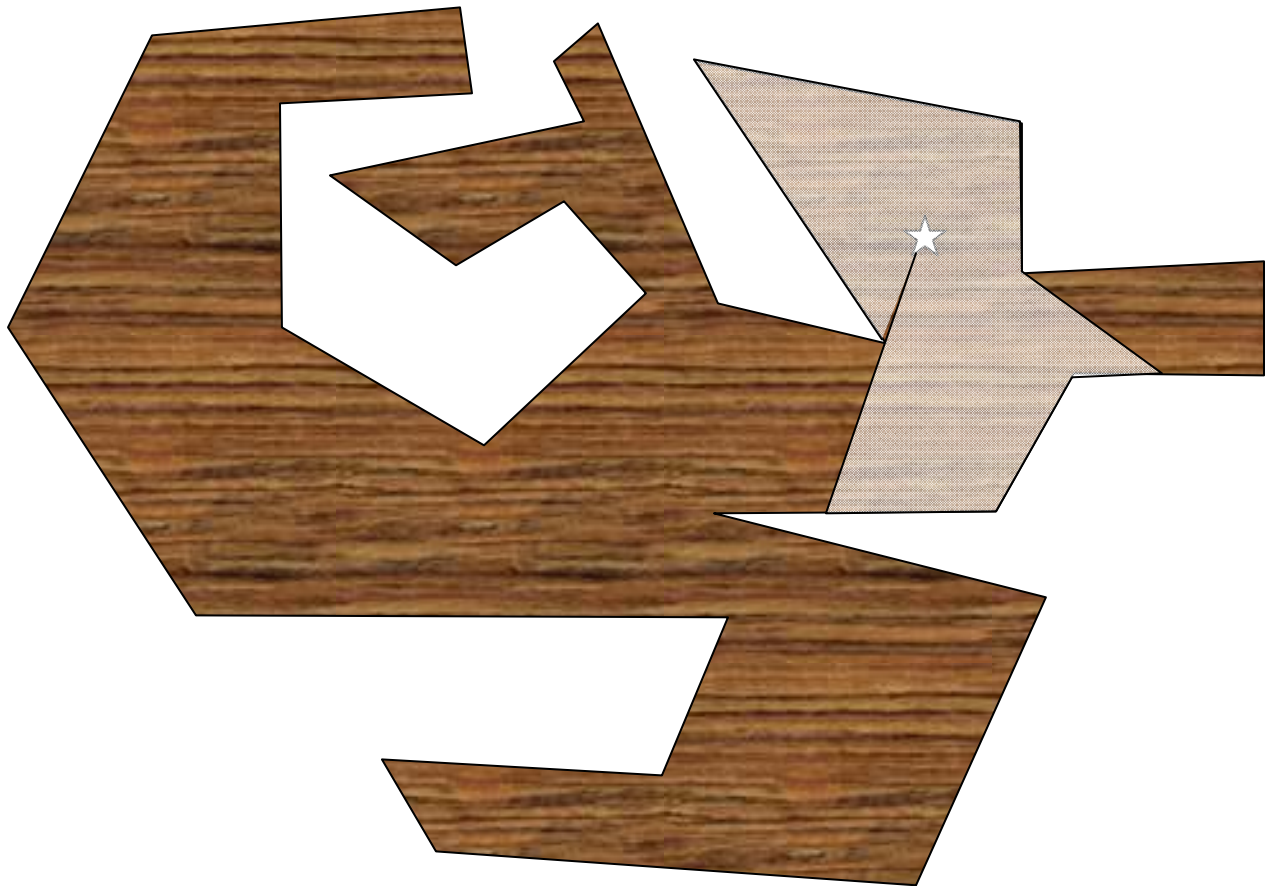
# Art Galleries

An art gallery is a region in the plane bounded by a simple polygon.



# Art Galleries

A guard is a point and the visibility of a guard is the set of points which can be joined to the guard by a straight line that does not cross through a wall of the art gallery.



# Art Galleries

For each art gallery, what is the smallest number of guards required so that every point of the art gallery is visible to at least one guard?

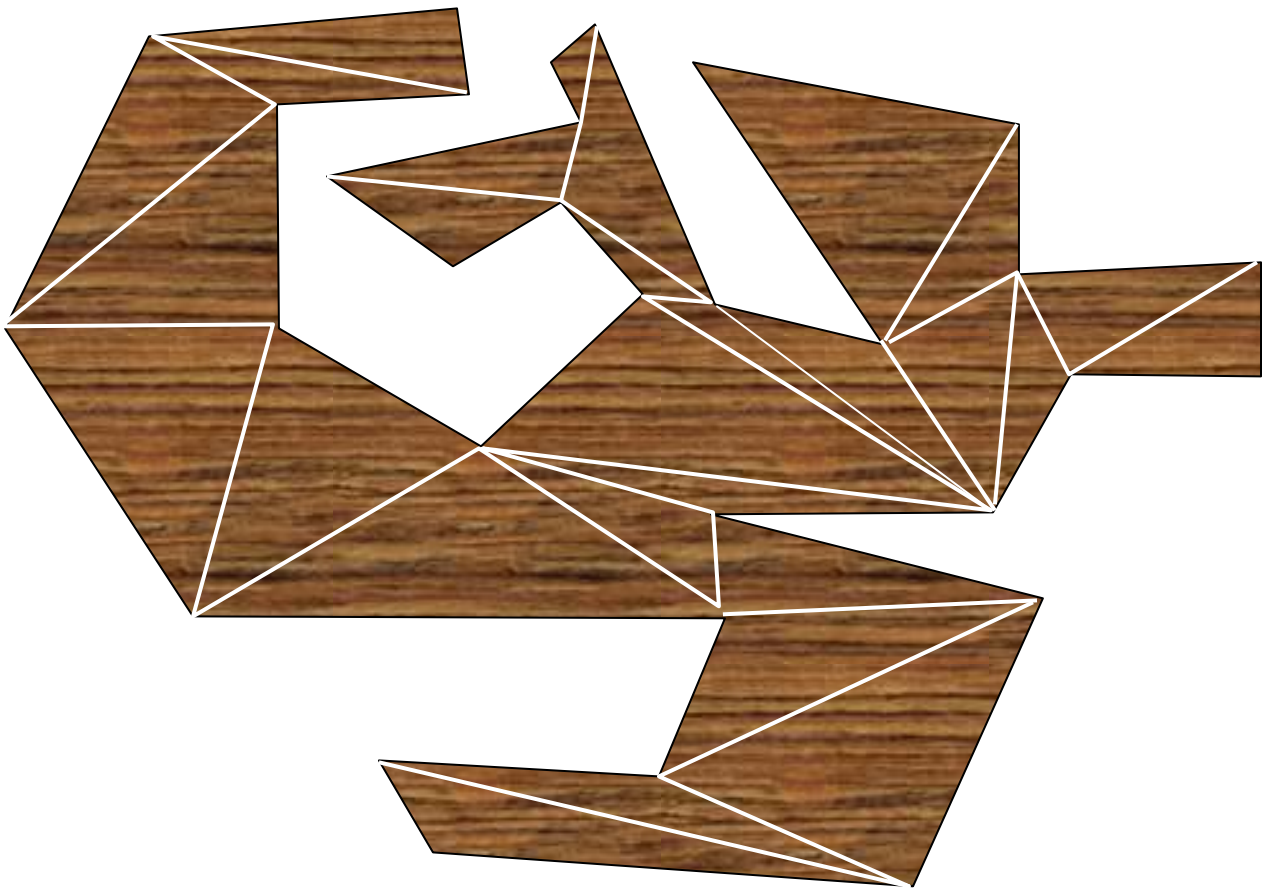
## Art Gallery Theorem (Chvátal, 1975)

*For any art gallery with  $n$  corners,  $n/3$  guards are enough to guard the whole gallery.*

# Art Galleries

How to prove the art gallery theorem

Triangulate the polygon





# Art Galleries

Then we get a planar graph  $\mathcal{A}$  with  $n$  vertices.  $\mathcal{A}$  has a vertex of degree two. Remove the vertex of degree two from  $\mathcal{A}$ . We get another triangulation of an art gallery. Again find a vertex of degree two and remove it.

So we conclude that  $\mathcal{A}$  is built up from a single vertex by repeatedly adding a vertex of degree at most two.

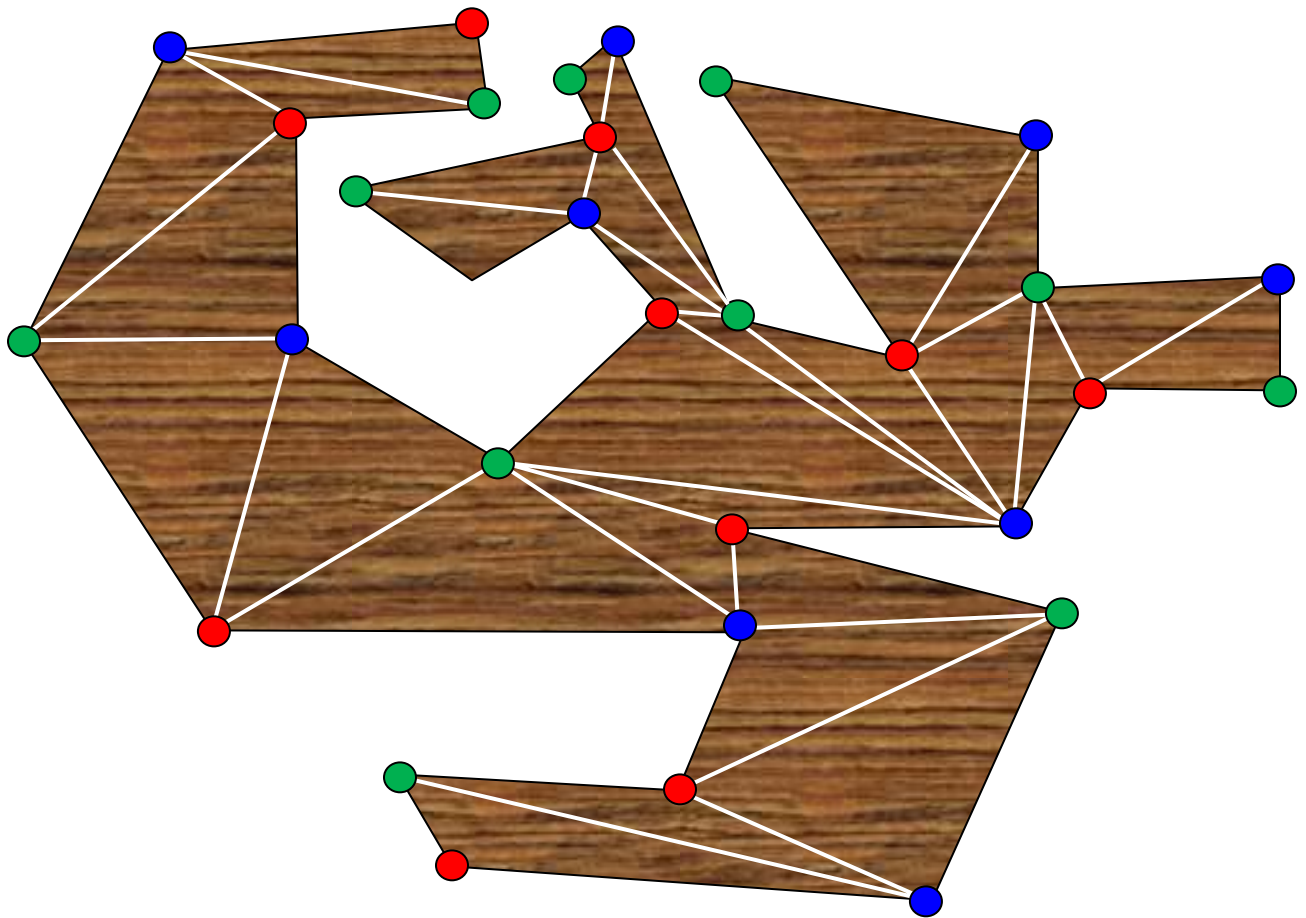
By the coloring algorithm it is 3-colorable.

Pick color 1. Then every triangle has color 1 on it since three colors are needed to color a triangle.

Then place guards at vertices of color 1 and they see the whole art gallery.

# Art Galleries

How to prove the art gallery theorem



- (1) Can you draw an art gallery with
- (a) 6 corners and cannot be guarded by 1 guard?
  - (b) 7 corners and cannot be guarded by 2 guards?
  - (c)  $3n$  corners that cannot be guarded by  $n - 1$  guards?

(3)(a) Color the map below with four colors. Can it be colored with three colors?



(b) Find the chromatic number of each continent shown below.



Is it possible to hang a picture on a wall using a string and two nails in such a way that the picture does not fall but if either nail is removed then the picture falls down.

