



structure

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Dynamic Programming : Why the need for optimal sub structure

words dynamic progr

implement on a problem which doesn't have optimal substructure property.

2.Dynamic Programming : Why the need for optimal sub structure

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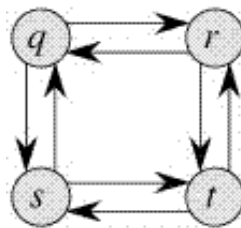
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In dynamic programming agiven problems has Optimal Substructure Proper problem can be obtained by using optimal solutions of its sub problems.

For example the shortest path problem has following optimal substructure p shortest path from a source node U to destination node V then the shortest shortest path from U to X and shortest path from X to V.

But Longest path problem doesn't have the Optimal Substructure property. i.e the longest path between two nodes doesn't have to be the longest path

For example, the longest path $q \rightarrow r \rightarrow t$ is not a combination of longest path t , because the longest path from q to r is $q \rightarrow s \rightarrow t \rightarrow r$.



So here: optimal solution to a problem does not contain an optimal solution

For more details you can read

[Longest path problem from wikipedia](#)[Optimal substructure from wikipedia](#)

3.Dynamic Programming : Why the need for optimal sub structure

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You're perfectly right that the definitions are imprecise. DP is a technique for than an algorithm in itself. The term "optimal substructure" is a vague conc every loop can be expressed as a recursive function: each iteration solves a every algorithm with a loop a DP? Clearly not.

What people actually mean by "optimal substructure" and "overlapping subp are used often enough to decrease the asymptotic complexity of solutions. I useful! In most cases the subtle implication is a decrease from exponential t p

Ex: There is an exponential number of paths between two nodes in a dense looking at only a polynomial number of them because the memos are extre

On the other hand, Traveling salesman can be expressed as a memoized fur the memos cause a factor of $O((1/2)^n)$ time to be saved. But, the numbe $O(n!)$. This is so much bigger that the asymptotic run time is still super-exp an algorithm is generally not called a Dynamic Program even though it's foll the DP for shortest paths. Apparently it's only a DP if it gives a nice result!