

Math417
Mathematical Programming

Homework IV

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1 Completing the proof of the f. theorem

Let $\{J_i\}$ denote the set of subsets of $\{1, \dots, n\}$ with $|J| = m$ and such that a_j with $j \in J_i$ are linearly independent $\forall J_i$. Also let A^i be the matrix

$$\begin{pmatrix} a_j^T \\ \cdot \\ \cdot \\ \cdot \end{pmatrix} \text{ with } j \in J_i$$

Then there exists a unique point, say \bar{x} such that $A^i \bar{x} = b$ because by construction it has full rank. Therefore for all J_i there exists only one associated BFS. Moreover there is only a finite number of J_i so there can only be a finite number of BFS.

2 Derivation of the dual problem

Here we have two cases

i) $v \geq 0$

$$\inf\{v^T x \mid x \geq 0 \quad v \geq 0\} = 0$$

ii) $v_i < 0$ for some i

$$\inf\{v^T x \mid x \geq 0\} = -\infty$$

The first case is quite obvious. To support the second we simply say that we can make the i th components of x go to infinity and keep the others null. So the infimum is unbounded.

3 Dual of the dual is the primal

We start with the dual problem as denoted in the course notes.

$$\max b^T y \quad A^T y \leq c \quad (1)$$

$$\iff -\min -b^T y \quad A^T y + s = c \quad s \geq 0 \quad (2)$$

$$\iff -\min b^T (y^- - y^+) \quad A^T (y^+ - y^-) + s = c \quad (3)$$

Now we have split y to get restrictions on it and introduced a slack variable which must be positive so we have

$$y = y^+ - y^- \quad y^+, y^-, s \geq 0 \quad (4)$$

Now we want to express the last problem in standard form. We construct the new vectors z and \hat{b} and the new matrix \hat{A}

$$z = \begin{pmatrix} y^+ \\ y^- \\ s \end{pmatrix} \quad \hat{A} = (A^T \mid -A^T \mid I_n) \quad \hat{b}^T = (-b^T \mid b^T \mid 0_n) \quad (5)$$

We mention the dimensions of these new objects.

$$\hat{A} \in \mathbb{R}^{n \times 2m+n} \quad s \in \mathbb{R}^{2m+n} \quad z \in \mathbb{R}^{2m+n} \quad \hat{b} \in \mathbb{R}^{2n} \quad (6)$$

We can now reexpress (1) as follows

$$-\min \hat{b}^T z \quad \hat{A}z = c$$

And we can now take its dual

$$\begin{aligned} \mathcal{D} : \quad & -\max c^T \xi \quad \hat{A}^T \xi \leq \hat{b} \\ \iff & \min c^T (-\xi) \quad \hat{A}(-\xi) \geq -\hat{b} \end{aligned}$$

Now let's take a look at the constraints

$$\begin{pmatrix} A \\ -A \\ I_n \end{pmatrix} - \xi = \begin{pmatrix} b \\ -b \\ 0_n \end{pmatrix} \quad (7)$$

Which we can break down into

$$A(-\xi) \geq b \quad -A(-\xi) \geq -b \quad -\xi \geq 0 \quad (8)$$

Now let $x = -\xi \geq 0$. We can see from the above reformulation that the two matrix inequalities yield an equality and so we have the standard form of the primal, namely:

$$\min c^T x \quad Ax = b \quad x \geq 0 \quad (9)$$

And so we have shown the dual of the dual is indeed the primal

4 Generalized duality

4.1 Weak duality

$$q(\lambda, \mu) = \inf_{x \in \mathbb{R}^n} L(x, \lambda, \mu)$$

Let \tilde{x} be a feasible point and $\lambda \geq 0$ then we have

$$\xi(x) = \lambda^\top g(x) + \mu^\top h(x) \leq 0 \quad (10)$$

Since $L(x, \lambda, \mu) = f(x) + \xi(x)$

We immediately have that

$$q(\tilde{x}, \lambda, \mu) = \inf\{L(\tilde{x}, \lambda, \mu) | \lambda \geq 0\} \leq L(\tilde{x}, \lambda, \mu) \leq f(\tilde{x})$$

4.2 Failure of strong duality

We find an example of a convex function where there is a duality gap. Consider the following optimization problem:

$$\min f(x) = \min e^{-x_1}$$

$$g(x) = \frac{x_1^2}{x_2} \leq 0$$

$$h(x) = 0$$

With the restriction that x_2 is strictly positive.
Then let

$$L(x, \lambda, \mu) = f(x) + \lambda g(x)$$

And define

$$q(\lambda, \mu) = \inf_{x \in \mathbb{R}^2} L(x, \lambda, \mu)$$

Then if λ is strictly negative q is unbounded and if it is greater or equal to zero it is null. So its supremum is also null. The dual problem becomes trivial, namely find the maximum of zero. But $p^* = 1$. So we have

$$\sup\{q(\lambda, \mu) | \lambda \geq 0\} = 0 < \inf\{f(x) | g(x) \leq 0, h(x) = 0\} = 1 \quad (11)$$