

COMP 9334

Capacity Planning

Assignment, Session 1, 2017

Assignment 1

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Question 1

Answer:

$$(a) D(j) = V(j) \cdot S(j) = \frac{C(j)}{C(0)} \times \frac{B(j)}{C(j)} = \frac{B(j)}{C(0)}$$

$B(j)$ = Busy time for the device j .

$C(0)$ = Number of completions for the system.

$$\therefore D(\text{CPU}) = \frac{B(\text{CPU})}{C(0)} = \frac{1102}{1237} \approx 0.895 \text{ s} = 895 \text{ ms}$$

$$D(\text{Disk1}) = \frac{B(\text{Disk1})}{C(0)} = \frac{929}{1231} \approx 0.755 \text{ s} = 755 \text{ ms}$$

$$D(\text{Disk2}) = \frac{B(\text{Disk2})}{C(0)} = \frac{1017}{1231} \approx 0.826 \text{ s} = 826 \text{ ms}$$

$$D(\text{Disk3}) = \frac{B(\text{Disk3})}{C(0)} = \frac{1265}{1231} \approx 1.028 \text{ s} = 1028 \text{ ms}$$

(b)

$$X(0) \leq \min \left[\frac{1}{\max D_i}, \frac{N}{\sum_{i=1}^K D_i} \right]$$

$$\text{Bond 1} = \frac{1}{\max D_i} = \frac{1}{D(\text{Disk3})} = \frac{1}{1.028} \text{ jobs/s} \approx 0.973 \text{ job/s}$$

$$\text{Bond 2} = \frac{N}{\sum_{i=1}^K D_i} = \frac{40}{0.895 + 0.755 + 0.826 + 1.028} \approx 11.42 \text{ jobs/s}$$

(c)

$$M_{\max} = X_0 \cdot (Z + R) \quad \dots \dots \dots \left. \begin{array}{l} M_{\max} = M, \\ R \text{ is the } \text{min} \text{ minimum} \\ \text{time bond.} \end{array} \right\}$$

$$\therefore \frac{M_{\max}}{X_0} - Z = R$$

$$\therefore R = \frac{40}{0.973} - 27 = 14.1 \text{ s.}$$

Hence, The minimum possible response time is 14.1 s.

Question 2

Answer :

- (a) Probability that an arrival is blocked
 = Probability that there are m customers in the system.

$$P_m = \frac{\rho^m}{m!} \bigg/ \sum_{k=0}^m \frac{\rho^k}{k!}, \text{ where } \rho = \frac{\lambda}{\mu}, \text{ according to the question,}$$

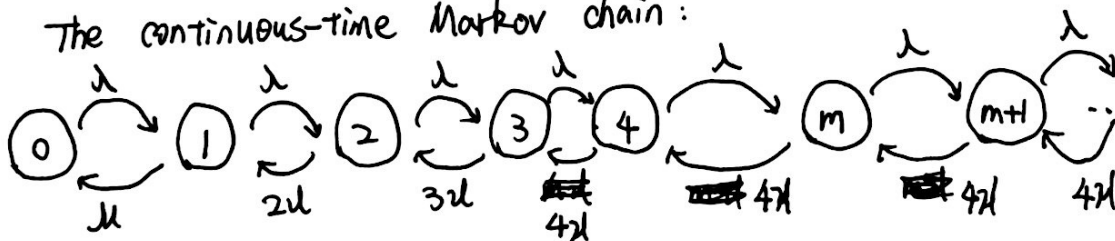
$$m=4, \quad \rho = \frac{\lambda}{\mu} = \frac{\frac{1}{3}}{\frac{1}{10}} = \frac{10}{3} \approx 3.33$$

$$\therefore P_4 = \frac{\frac{(10/3)^4}{4!}}{\frac{(10/3)^0}{0!} + \frac{(10/3)^1}{1!} + \frac{(10/3)^2}{2!} + \frac{(10/3)^3}{3!} + \frac{(10/3)^4}{4!}} \approx 0.243$$

(b)

(i)

The continuous-time Markov chain:



state 0 = There is no call in the system.

state 1 = There is 1 call in the system (= 1 call at ^{one of the} operators, no call in the slots)

state 2 = There is 2 calls in the system (= 2 calls at the two of the operators, no call in the slots.)

state 3 = There is 3 calls in the system (= 3 calls at the operators, no calls in the slots)

state 4 = There is 4 calls in the system (= 4 calls at the operators, no calls in the slots)

state m = There is m calls in the system (= 4 calls at the operators, $(m-4)$ calls in the slots).

state $m+1$ = There is $m+1$ calls in the system (= 4 calls at the operators, $(m-3)$ calls in the slots)

(b) (i) the transition rates are shown ~~at~~^{at} the chain, where
 $\lambda = \frac{1}{3} \approx 0.333$ calls/min ~~It is 15 calls per hour~~
 $\mu = \frac{1}{\text{mean service rate}} = \frac{1}{10} = 0.1$ service/min.

~~(b)~~

(ii)

$$\lambda P_0 = \mu P_1$$

$$2\mu P_2 = \lambda P_1$$

$$3\mu P_3 = \lambda P_2$$

$$4\mu P_4 = \lambda P_3$$

$$\vdots$$

$$4\mu P_m = \lambda P_{m-1}$$

$$4\mu P_{m+1} = \lambda P_m$$

\vdots

$$P_0 + P_1 + P_2 + P_3 + P_4 + \dots + P_m + P_{m+1} \dots = 1$$

where $\lambda = 0.333$ calls/min

$\mu = 0.1$ service/min.

(iii) The steady state probabilities:

I don't make expressions for them, because I made the calculation jobs by using a python program [which I submitted as "solve-equations.py"]. Therefore, I could calculate $P_0, P_1, P_2 \dots P_m, P_{m+1}$ without express them by using only λ and μ .

Please check my codes attached, thank you.

(iv) From the result of Part (a)

the rejection rate with no slots is 0.243, which 50% is about 0.121, hence we need to fetch a $P(M+4) < 0.121$, where we have $m=4$ slots.

According to the calculation with "solve-equation.py" I made, ~~the~~ the smallest M could be 3, i.e. there are 4 operators and 3 holding slots. In the situation,

$P(M+4) = P(7) = 0.0929$, which is the ~~smallest M~~ first time (the smallest M), $P(M+4) < 0.121$.

Therefore, the smallest M is 3.

(v) According to the Little's Law:

$$N_{avg} = X(0) \times R_{avg}, \text{ where.}$$

$$X(0) = \lambda$$

$$R_{avg} = T = \frac{CP.m}{m \mu (1-p)} + \frac{1}{\mu}. \quad (\lambda = \frac{1}{3}, \mu = \frac{1}{10}, \therefore p = \frac{5}{6})$$

$$= \frac{0.0514}{4 \times \frac{1}{10} \times (\frac{1}{6})} + \frac{1}{\frac{1}{10}}.$$

$$\Rightarrow N_{avg} = \frac{1}{3} \times 10.772 = 3.591 \text{ calls.}$$

$$\therefore \text{The waiting time} = \frac{N_{avg} \times P(\text{slots})}{X(0)}$$

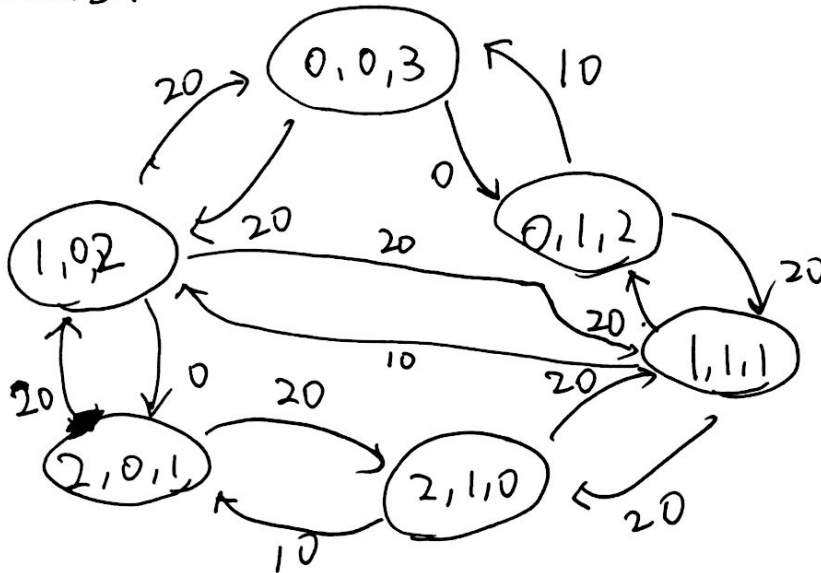
$$= \frac{N_{avg} \times [P(5) + P(6) + P(7)]}{X(0)}$$

$$= \frac{3.591 \times [0.1338 + 0.1115 + 0.0929]}{\frac{1}{3}}$$

$$\rightarrow \underline{\underline{3.6429 \text{ (mins)}}}$$

Question 3.

(a)



STATE	Job amount in CPU1	Job amount in CPU2	Job amount in disk
state (0,0,3)	0	0	3
state (1,0,2)	1	0	2
state (0,1,2)	0	1	2
state (2,0,1)	2	0	1
state (1,1,1)	1	1	1
state (2,1,0)	2	1	0

The transition rates are shown in the graph, for:

- ① Jobs complete at CPU1 at a rate of $\frac{1000}{50} = 20$ jobs/s.
- ② Jobs complete at CPU2 at a rate of $\frac{1000}{100} = 10$ jobs/s.
- ③ Jobs complete at Disk at a rate of $\frac{1000}{50} = 20$ jobs/s.

Question 3

(b) Equations could be fetched:

$$\textcircled{1} 20P(0,0,3) - 20P(1,0,2) - 10P(0,1,2) + 0P(2,0,1) - P(1,1,1) + 0P(2,1,0) = 0$$

$$\textcircled{2} 30P(0,1,2) - 20P(1,1,1) = 0$$

$$\textcircled{3} -20P(2,0,1) - 20P(1,1,1) + 30P(2,1,0) = 0$$

$$\textcircled{4} 40P(2,1,0) - 10P(2,1,0) = 0$$

$$\textcircled{5} -20P(1,0,2) - 20P(0,1,2) + 50P(1,1,1) - 20P(2,1,0) = 0$$

$$\textcircled{6} P(0,0,3) + P(1,0,2) + P(0,1,2) + P(2,0,1) + P(1,1,1) + P(2,1,0) = 1$$

(c) By using "q3.py" (attached in the zip file);

$$P(0,0,3) = 0.2697$$

$$P(1,0,2) = 0.2039$$

$$P(0,1,2) = 0.1316$$

$$P(2,0,1) = 0.0395$$

$$P(1,1,1) = 0.1974$$

$$P(2,1,0) = 0.1579$$

Question 3:

(d) For this system.

$$\begin{aligned} \text{The system throughput} &= \text{CPU1 throughput} + \text{CPU2 throughput} \\ &= \text{Utilization of CPU1} \times \text{Service rate of CPU1} \\ &\quad + U(\text{CPU2}) \times S(\text{CPU2}) \\ &= [P(1,0,2) + P(2,0,1) + P(1,1,1) + P(2,1,0)] \times \\ &\quad 20 \text{ jobs/s} + [P(0,1,2) + P(1,1,1) + P(2,1,0)] \\ &\quad \times 10 \text{ jobs/s} = 11.974 + 4.869 \\ &\quad = 16.843 \text{ jobs/s} \end{aligned}$$

(e).

$$\begin{aligned}\text{Throughput of CPU} &= C(\text{CPU}) \times S(\text{CPU}) \\ &= [P(1,0,2) + P(2,0,1) + P(1,1,1) + \\ &\quad P(2,1,0)] \times 20 \text{ jobs/s} \\ &= (0.2039 + 0.0395 + 0.1974 + 0.1579) \times 20 \\ &= 11.974 \text{ jobs/s.}\end{aligned}$$

(f)

The throughput of the disk = $U(\text{Disk}) \times S(\text{Disk})$
 $= [1 - P(2, 1, 0)] \times 20 \text{ jobs/s}$
 $= 16.843 \text{ jobs/s.}$

$$\text{Response time of the disk} = \frac{\text{user amount in the system}}{\text{the throughput of the disk}}$$
$$= \frac{3}{16.843} \text{ s.}$$

Service time of the disk = $\frac{\text{Utilization (disk)}}{\text{throughput (disk)}} = \frac{0.842}{16.843} = 0.05 \text{ s} = 50 \text{ ms}$

\therefore the waiting time = Response time of the disk -
the service time of the disk
 $= 178 - 50 \text{ ms} = 128 \text{ ms}$.