

- (a) **MINCUT**(N):
- $f \leftarrow \text{FORDFULKERSON}(N)$ # find a maximum flow in N
 - construct the residual network N_f
 - $S \leftarrow \{\text{all vertices reachable from } s \text{ in } N_f\}$ (found using BFS)
 - return** ($S, V - S$)

As explained in lecture, the Ford-Fulkerson algorithm can be implemented to run in worst-case time $\Theta(n^2m)$. Constructing the residual network takes time $\Theta(n+m)$, and so does running BFS on N_f . So the total running time is dominated by the Ford-Fulkerson algorithm at $\Theta(n^2m)$.

- (b) Use the same algorithm as above but run it on the “reverse network” N' , equal to N with all edge directions reversed.

Since the original algorithm guarantees that $|S|$ is minimum, this algorithm guarantees that $|T|$ will be minimum.