

1. Recall the CLIQUE decision problem from last week's tutorial.

**Input:** An undirected graph  $G = (V, E)$  and a positive integer  $k$ .

**Question:** Does  $G$  contain a *clique* of size at least  $k$ , *i.e.*, a subset of  $k$  or more vertices such that  $G$  contains **every** possible edge between the vertices in the clique?

In last week's tutorial, you showed that  $\text{CLIQUE} \in NP$ . Now, show that CLIQUE is NP-hard. Give a detailed reduction and argument of correctness. (HINT: You can use any of the problems you know to be NP-hard from lectures, tutorials, or the textbook — except for CLIQUE itself, of course!)