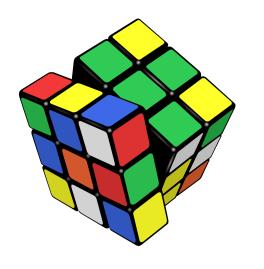
Quiz 1

Section 1 (LEC0101, LEC2003)





Computer Science

Fall 2017 contest winner

CSC263 Sep 27 – 2017

Assume that algorithm A is asymptotically more efficient than B. Which statement is true?

Players correct: 58%

- A) A will be a better choice for all inputs
- B) B will be a better choice for small inputs
- C) A will be a better choice for all inputs except large inputs
- D) A will be a better choice for all inputs except small inputs

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Let $T_W(n)$ be the worst-case and $T_A(n)$ be the average-case running time of an algorithm. Which one is always true?

Players correct: 40%

A)
$$T_W(n) = O(T_A(n))$$

B)
$$T_W(n) = \Omega(T_A(n))$$

C)
$$T_W(n) = \Theta(T_A(n))$$

D)
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There are three possible procedures A, B, C for running an application. Procedure A takes 2 seconds, procedure B takes 3 seconds and procedure C takes 4 seconds to run. This application to be run chooses algorithm A with the probabilty of 0.30, the algorithm B with the probabilty of 0.20 and C with probabilty of 0.50. What is the average case running time of the application?

Players correct: 84%

A) 2

B) 4.5

C) 3.2

D)3

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- A) 2
- B) 4.5 C) 3.2

D)3

Average-case running time = $\sum x p(x) = 2 * 0.3 + 4 * 0.2 + 4 * 0.5 = 3.2$

What is the minimum number of comparisons required to find a minimum value stored in a max heap? Players correct: 24%

A) n

B) n/2 C) $O(\log n)$

$$(D)^{\frac{n}{2}} - 1$$

What is the minimum number of comparisons required to find a minimum value stored in a max heap? Players correct: 24%

- A) n B) n/2 C) $O(\log n)$ D) $\frac{n}{2} 1$

$$(D)^{\frac{n}{2}} - 1$$

Answer: The minimum element in a max-heap can be anywhere in the leaves of the heap. (why?)

There are up to n/2 elements in the leaves, so finding the minimum needs n/2 times.

What is time complexity of increasing and decreasing a value of an element in a max heap?

Players correct: 71%

$$B)O(\log n), O(n) \quad C) O(\log n), O(\log n)$$

What is time complexity of increasing and decreasing a value of an element in a max heap?

Players correct: 71%

Answer:

- Increasing a value need a max-heapify-up operation
- Decreasing a value needs a max-heapify-down operation

Both operations are in $O(\log n)$ since needs to traverse the height of the tree in the worst case.

In a heap of height *h*, what is the minimum number of elements?

(Assume the height of the root is zero).

A)
$$2^{(h+1)} - 1$$

B)
$$2^{h}$$

$$C) 2^{(h-1)}$$

$$(D) 2^{(h+1)}$$

Players correct: 25%

In a heap of height *h*, what is the minimum number of elements?

(Assume the height of the root is zero).

Players correct: 25%

A)
$$2^{(h+1)} - 1$$

B)
$$2^{h}$$

$$C) 2^{(h-1)}$$

$$(D) 2^{(h+1)}$$

Heap is a complete binary tree, so it is full at all level bottom level which has at least 1 node. The number of nodes at level 0 is $1 = 2^0$, at level 1 is $2 = 2^1$, ..., at level h - 1 is 2^{h-1} and at level h is 1.

Total number of nodes: $(1 + 2 + 2^2 ... + 2^{h-1}) + 1 = (2^h - 1) + 1 = 2^h$.

Given four numeric values a < b < c < d as the elements' priorities, how many possible max heaps would be possible to build?

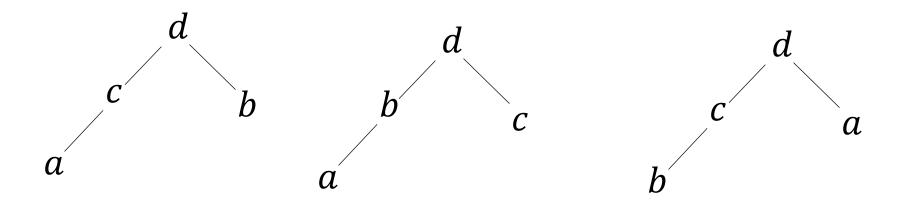
Players correct: 40%

A)2 B)3 C)4 D)5

Given four numeric values a < b < c < d as the elements' priorities, how many possible max heaps would be possible to build?

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A)2 B)3 C) 4 D) 5



Assume that the root of a heap lies on level 1, the immediate children of the root lies on level 2. Then, k^{th} largest element in the max-heap lies always on the k level of the heap.

A) True

B) False

Players correct: 79%

Assume that the root of a heap lies on level 1, the immediate children of the root lies on level 2. Then, k^{th} largest element in the max-heap lies always on the k level of the heap.

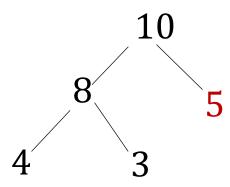
A) True

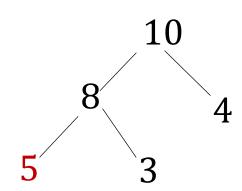
B) False

Players correct: 79%

For values of k > 2 it is not always true.

Counter example:





A **ternary** max heap is defined similar to a **binary** max heap but each node has at most 3 **children** instead of 2 **children**. Let $A = \{15, 6, 8, 14, 3, 2\}$ be a ternary max heap with starting index at zero. We insert a new element 11. What is the index of 11 after being inserted to A. (A must remain a ternary max heap.)

A)1

B)2

C) 3

D) 4

Players correct: 44%

A **ternary** max heap is defined similar to a **binary** max heap but each node has at most **3 children** instead of **2 children**.

Let $A = \{20, 16, 17, 19, 8, 2, 14\}$ be a ternary max heap with starting index at zero. We insert a new element 18. What is the index of 18 after being inserted to A. (A must remain a ternary max heap.)

