

# CSC236 Tutorial Exercises

## Sample Solutions

1. Consider the recurrence relation

$$T(n) = \begin{cases} 1 & n = 1 \\ 1 + T\left(\left\lceil \frac{n}{3} \right\rceil\right) & n > 1 \end{cases}$$

Use complete induction to prove that for every positive natural number  $n$ ,  $T(n) \geq c \lg(n)$ , for some positive real constant  $c$ .

**Sample solution.**

**Basis step:**

$$T(1) = 1 \geq c \lg(1) = 0$$

**Inductive step:**

Assume  $T(i) \geq c \lg i$  for  $1 \leq i < k$  and an arbitrary  $k > 1$ . We must show  $T(k) \geq c \lg k$ .

$$T(k) = 1 + T\left(\left\lceil \frac{k}{3} \right\rceil\right) \quad \text{by definition of } T \text{ when } k > 1$$

$$\Rightarrow T(k) \geq 1 + c \lg \left\lceil \frac{k}{3} \right\rceil \quad \text{by I.H., since } 1 \leq \left\lceil \frac{k}{3} \right\rceil < k \text{ when } k > 1$$

$$\Rightarrow T(k) \geq 1 + c \lg \frac{k}{3} \quad \text{since } \frac{k}{3} \leq \left\lceil \frac{k}{3} \right\rceil \text{ and } \lg \text{ is increasing}$$

$$\Rightarrow T(k) \geq 1 + c \lg k - c \lg 3$$

$$\Rightarrow T(k) \geq c \lg k \quad \text{provided } 1 - c \lg 3 \geq 0 \Rightarrow \frac{1}{\lg 3} \geq c \Rightarrow 0.63 \geq c$$

□

2. Consider the recurrence relation

$$T(n) = \begin{cases} 1 & n = 0 \\ 3 & n = 1 \\ 3T(n-1) - 2T(n-2) & n > 1 \end{cases}$$

Find a closed form for  $T(n)$ , and prove that it is correct using induction.

$$\begin{aligned} T(n) &= 3T(n-1) - 2T(n-2) \\ &= 3(3T(n-2) - 2T(n-3)) - 2T(n-2) = 7T(n-2) - 6T(n-3) \\ &= 7(3T(n-3) - 2T(n-4)) - 6T(n-3) = 15T(n-3) - 14T(n-4) \\ &\quad \dots \\ &= (2^{k+1} - 1)T(n-k) - (2^{k+1} - 2)T(n-k-1) \end{aligned}$$

We must choose (i.e., continue unwinding to) a  $k$  such that  $T(n-k)$  and  $T(n-k-1)$  can be replaced by base cases, i.e.,  $T(1)$  and  $T(0)$ . Hence,  $k = n - 1$ .

$$\begin{aligned} &= (2^n - 1)T(1) - (2^n - 2)T(0) \\ &= (2^n - 1) \cdot 3 - (2^n - 2) \cdot 1 = 3 \cdot 2^n - 3 - 2^n + 2 = 2^{n+1} - 1 \end{aligned}$$

Now, we use complete induction to prove  $T_r(n) = T_c(n)$ , where

$$T_r(n) = \begin{cases} 1 & n = 0 \\ 3 & n = 1 \\ 3T_r(n-1) - 2T_r(n-2) & n > 1 \end{cases} \quad \text{and} \quad T_c(n) = 2^{n+1} - 1$$

**Basis step:**

$$T_r(0) = 1 = T_c(0) = 2^{0+1} - 1 = 1$$

$$T_r(1) = 3 = T_c(1) = 2^{1+1} - 1 = 3$$

**Inductive step:**

Assume  $T_r(i) = T_c(i)$  for  $0 \leq i < k$  and an arbitrary  $k > 1$ . We must show  $T_r(k) = T_c(k)$ .

$$\begin{aligned} T_r(k) &= 3T_r(k-1) - 2T_r(k-2) && \text{by definition of } T_r \text{ when } k > 1 \\ &= 3T_c(k-1) - 2T_c(k-2) && \text{by I.H, since } 0 \leq k-1 < k \text{ when } k > 1 \text{ and} \\ &&& 0 \leq k-2 < k \text{ when } k > 1 \\ &= 3(2^{k-1+1} - 1) - 2(2^{k-2+1} - 1) && \text{by definition of } T_c \\ &= 3 \cdot 2^k - 3 - 2 \cdot 2^{k-1} + 2 \\ &= 3 \cdot 2^k - 2^k - 1 \\ &= 2 \cdot 2^k - 1 \\ &= 2^{k+1} - 1 \\ &= T_c(k) \end{aligned}$$

□

3. Consider another recurrence relation

$$T(n) = \begin{cases} 1 & n = 0 \\ T(n-1) + n - 2 & n > 0 \end{cases}$$

Unwind the recurrence **carefully**, following the pattern below, for some  $n$  that is comfortably greater than 1:

$$\begin{aligned} T(n) &= T(n-1) + n - 2 \\ &= T(n-2) + n - 1 - 2 + n - 2 = T(n-2) + 2n - 5 \\ &= T(n-3) + n - 2 - 2 + 2n - 5 = T(n-3) + 3n - 9 \end{aligned}$$

Continue to see a pattern that leads to a guess at a closed form for  $T(n)$ .

**Sample solution 1.**

$$\begin{aligned} T(n) &= T(n-1) + n - 2 \\ &= T(n-2) + n - 1 - 2 + n - 2 = T(n-2) + 2n - 5 \\ &= T(n-3) + n - 2 - 2 + 2n - 5 = T(n-3) + 3n - 9 \\ &= T(n-4) + n - 3 - 2 + 3n - 9 = T(n-4) + 4n - 14 \end{aligned}$$

So far, we see the pattern, in a general case, is like

$$T(n) = T(n-k) + kn - \dots \quad (*)$$

The question is, in (\*), how to express the last term (i.e., 2, 5, 9, 14, etc.) in terms of  $k$ . See the following table:

$k$	value
1	2
2	5
3	9
4	14
...	...

We need to manipulate the **value**, rewrite it using the corresponding  $k$ . There are different ways to do so. One is as follows:

$k$	value	
1	2	=1+1
2	5	=2+3
3	9	=3+6
4	14	=4+10
...	...	...

Now the question is what is the relationship of 1,3, 6, 10, etc., with their corresponding  $k$  (i.e., 1, 2, 3, 4, etc.)

$k$	value		
1	2	=1+1	=1+(1*2)/2
2	5	=2+3	=2+(2*3)/2
3	9	=3+6	=3+(3*4)/2
4	14	=4+10	=4+(4*5)/2
...	...	..	...

We found the pattern for the last term:  $k+(k*(k+1))/2$ . So, we can write the (\*) as follows:

$$T(n) = T(n-k) + kn - (k + k * (k+1)/2)$$

We must choose (*i.e.*, continue unwinding to) a  $k$  such that  $T(n - k)$  can be replaced by the base cases,  $T(0)$ . Hence,  $k = n$ . Therefore,

$$\begin{aligned}
 T(n) &= T(n - n) + n \cdot n - \left(n + \frac{n(n+1)}{2}\right) \\
 T(n) &= T(0) + n^2 - \left(n + \frac{n^2+n}{2}\right) \\
 T(n) &= 1 + n^2 - n - \frac{n^2+n}{2} \\
 T(n) &= \frac{2+2n^2-2n-n^2-n}{2} \\
 T(n) &= \frac{n^2-3n+2}{2} = \frac{(n-1)(n-2)}{2}
 \end{aligned}$$

### Sample solution 2.

Sometimes, we could faster find the pattern if we do not simplify the unwinding relations a lot:

$$\begin{aligned}
 T(n) &= T(n-1) + n - 2 \\
 &= T(n-2) + n - 1 - 2 + n - 2 = T(n-2) + 2n - 2 - 3 \\
 &= T(n-3) + n - 2 - 2 + 2n - 2 - 3 = T(n-3) + 3n - 2 - 3 - 4 \\
 &= T(n-4) + n - 3 - 2 + 3n - 2 - 3 - 4 = T(n-4) + 4n - 2 - 3 - 4 - 5 \\
 &\dots \\
 &= T(n-k) + kn - 2 - 3 - 4 - 5 - \dots - (k+1) \\
 &= T(n-k) + kn - \frac{(k+1)(k+2)}{2} + 1
 \end{aligned}$$

We must choose (*i.e.*, continue unwinding to) a  $k$  such that  $T(n - k)$  can be replaced by the base cases,  $T(0)$ . Hence,  $k = n$ .

$$\begin{aligned}
 &= T(n-n) + nn - \frac{(n+1)(n+2)}{2} + 1 \\
 &= T(0) + n^2 - \frac{(n+1)(n+2)}{2} + 1 \\
 &= 1 + n^2 - \frac{n^2 + 3n + 2}{2} + 1 \\
 &= \frac{2n^2 - n^2 - 3n - 2 + 4}{2} \\
 &= \frac{n^2 - 3n + 2}{2} = \frac{(n-1)(n-2)}{2}
 \end{aligned}$$