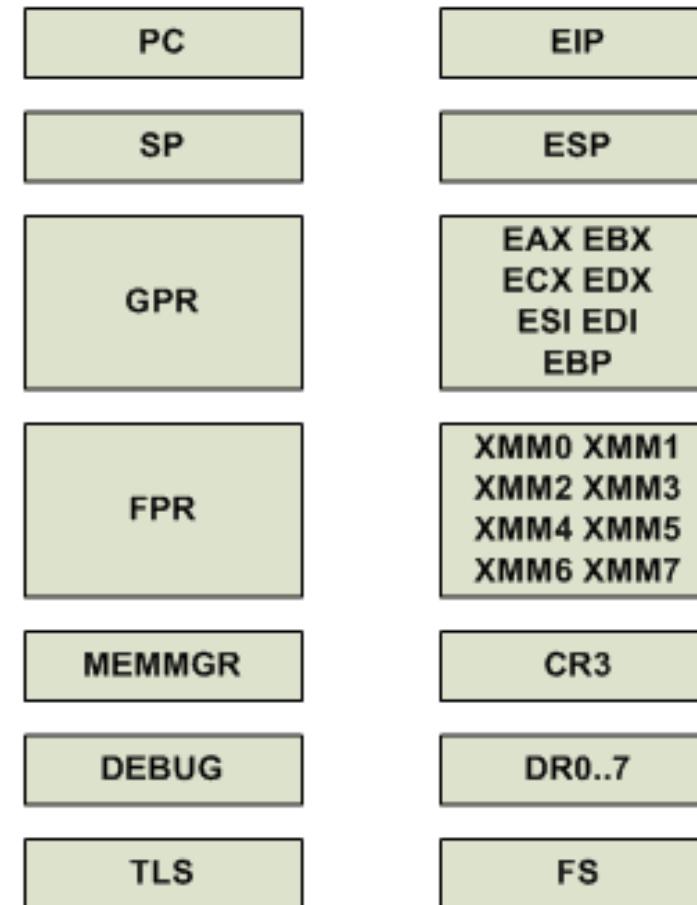


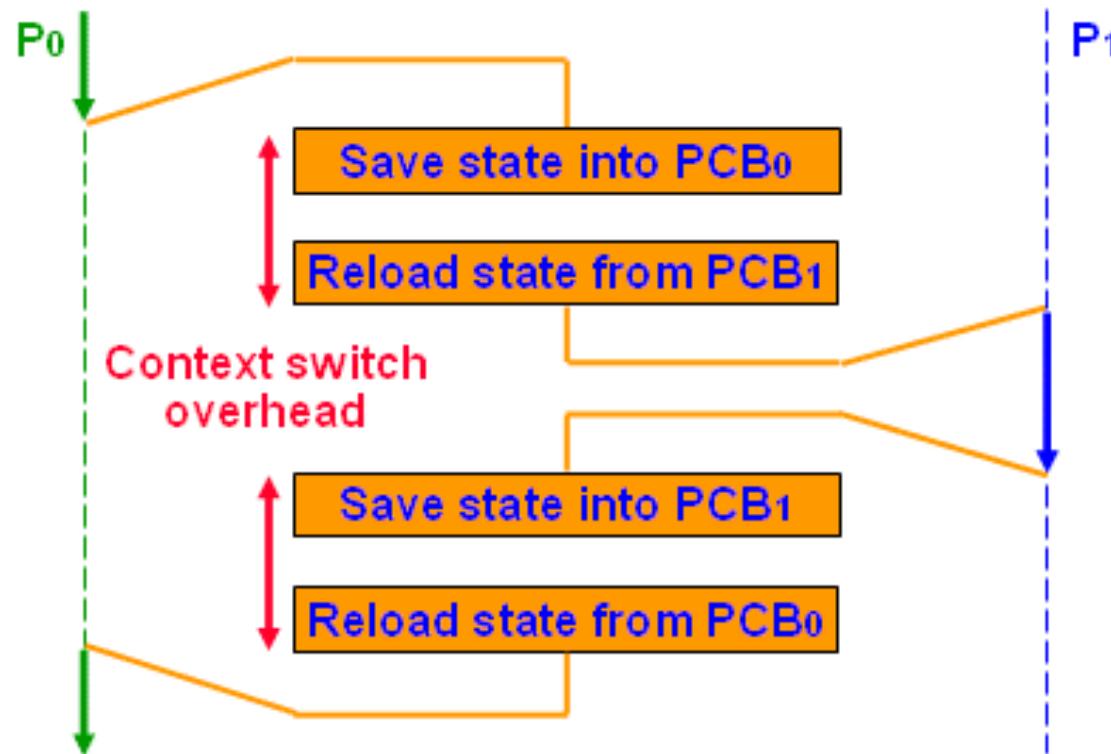
Процессы и потоки. Планирование задач.

Инициализация процессорных ресурсов

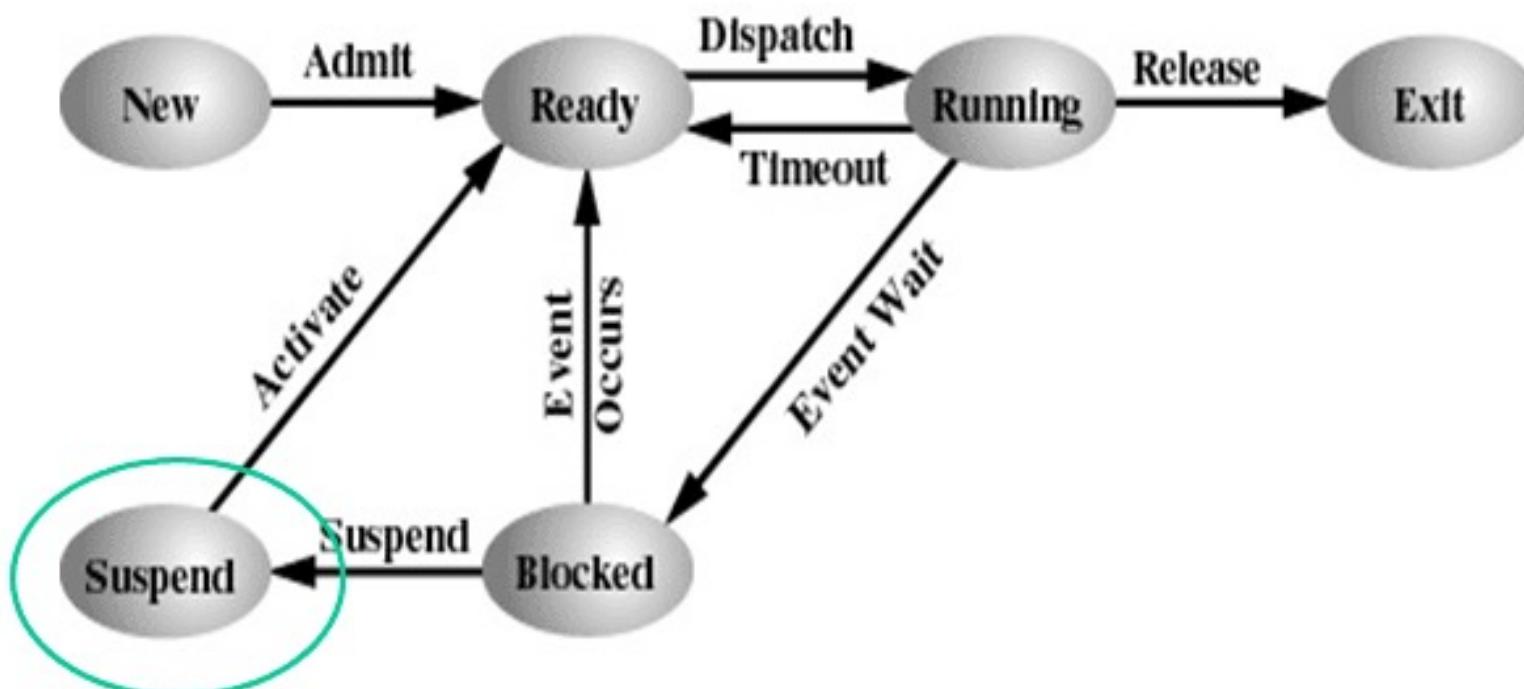
Контекст выполнения процесса



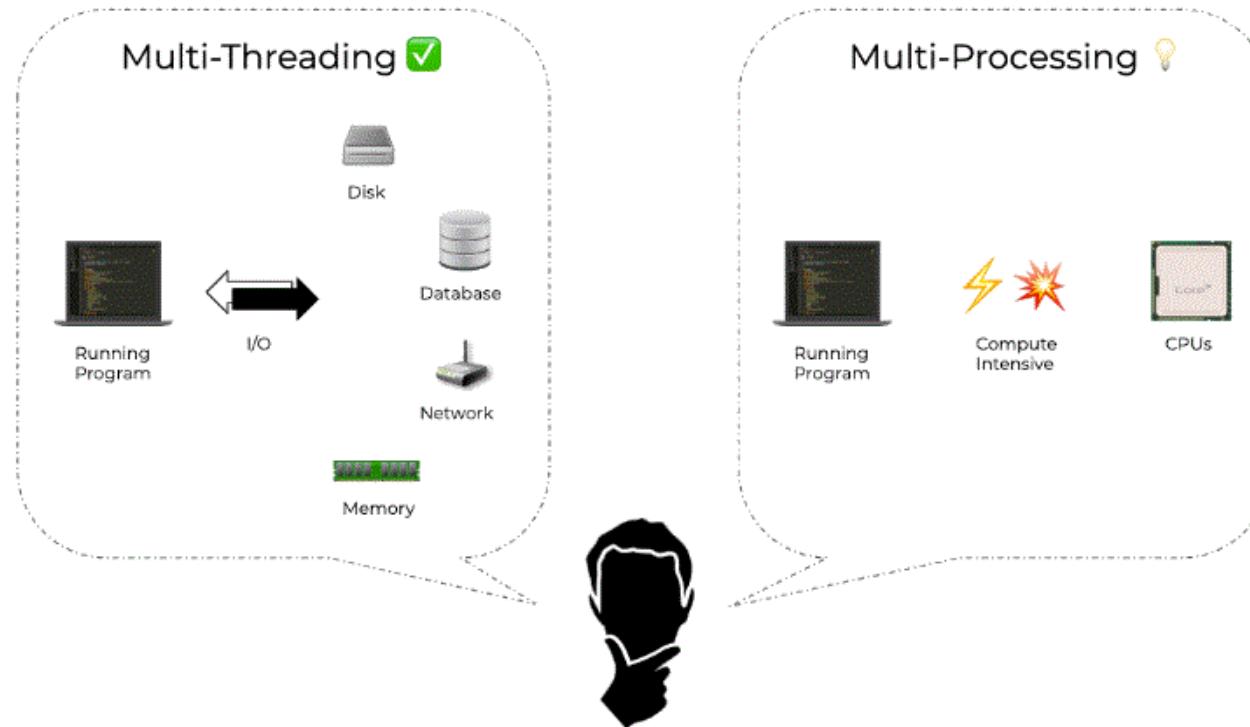
Переключения контекста



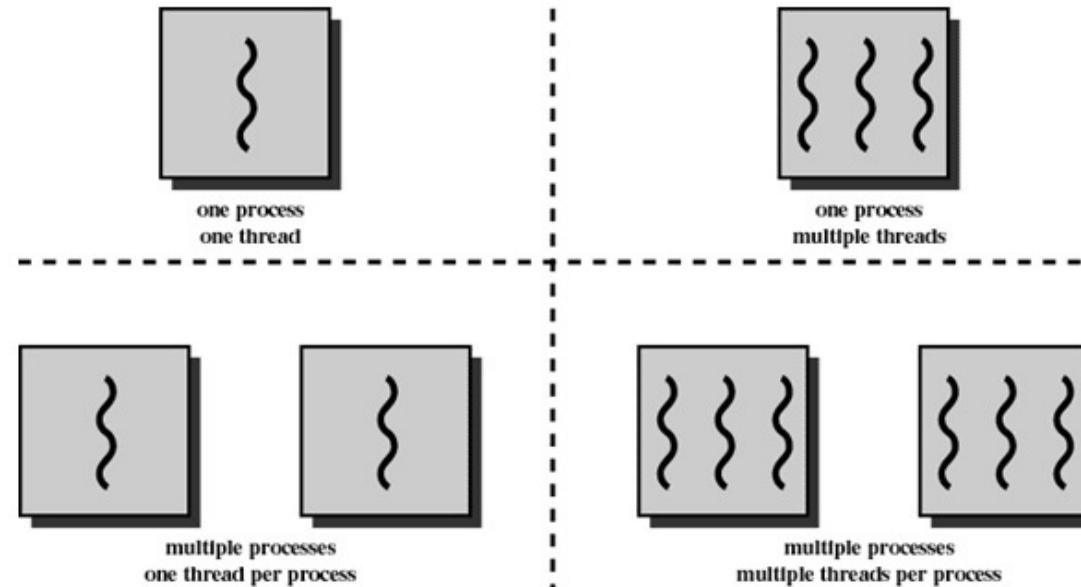
Состояния нити



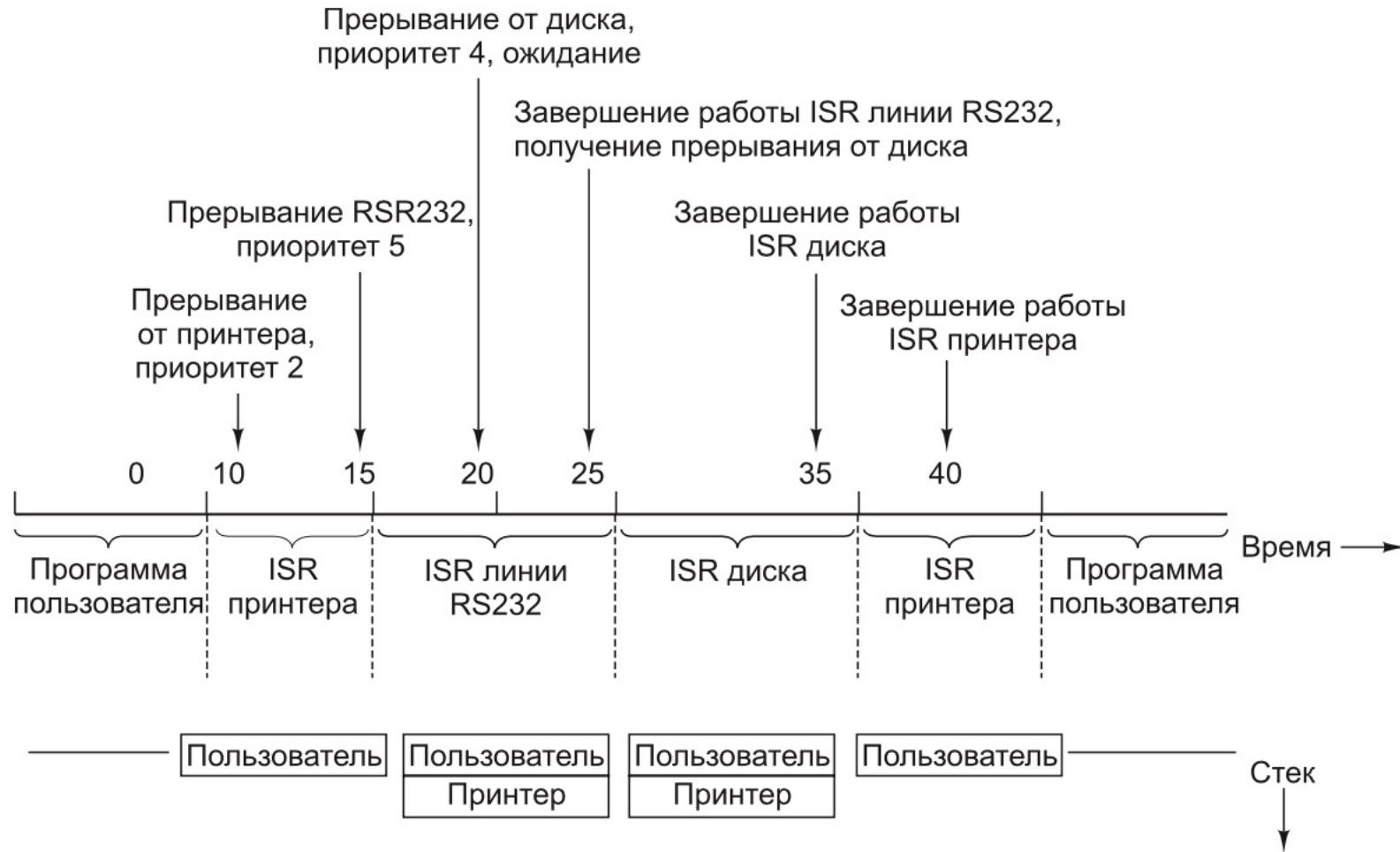
Процессы и потоки



Процессы и потоки

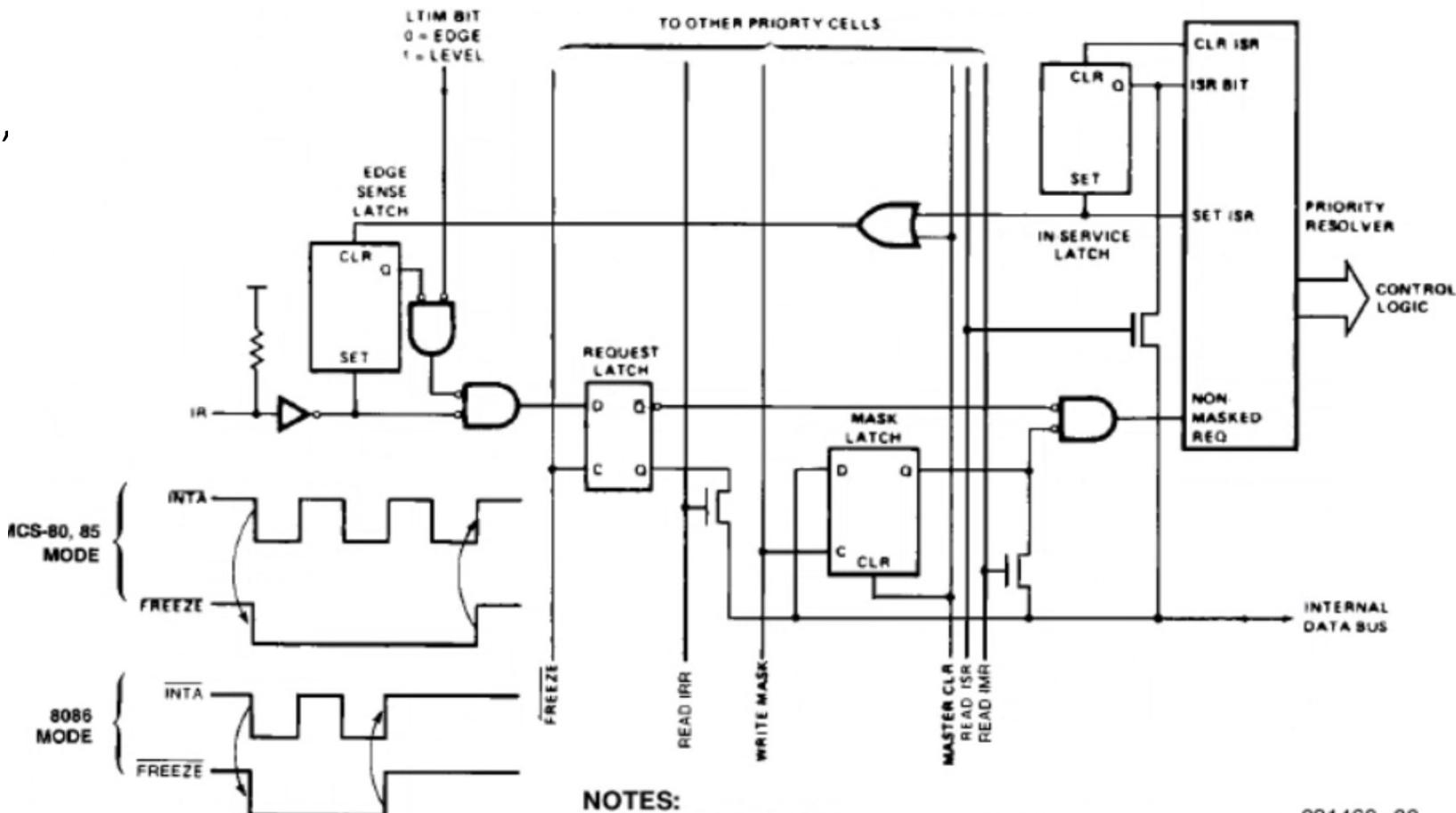


Прерывания



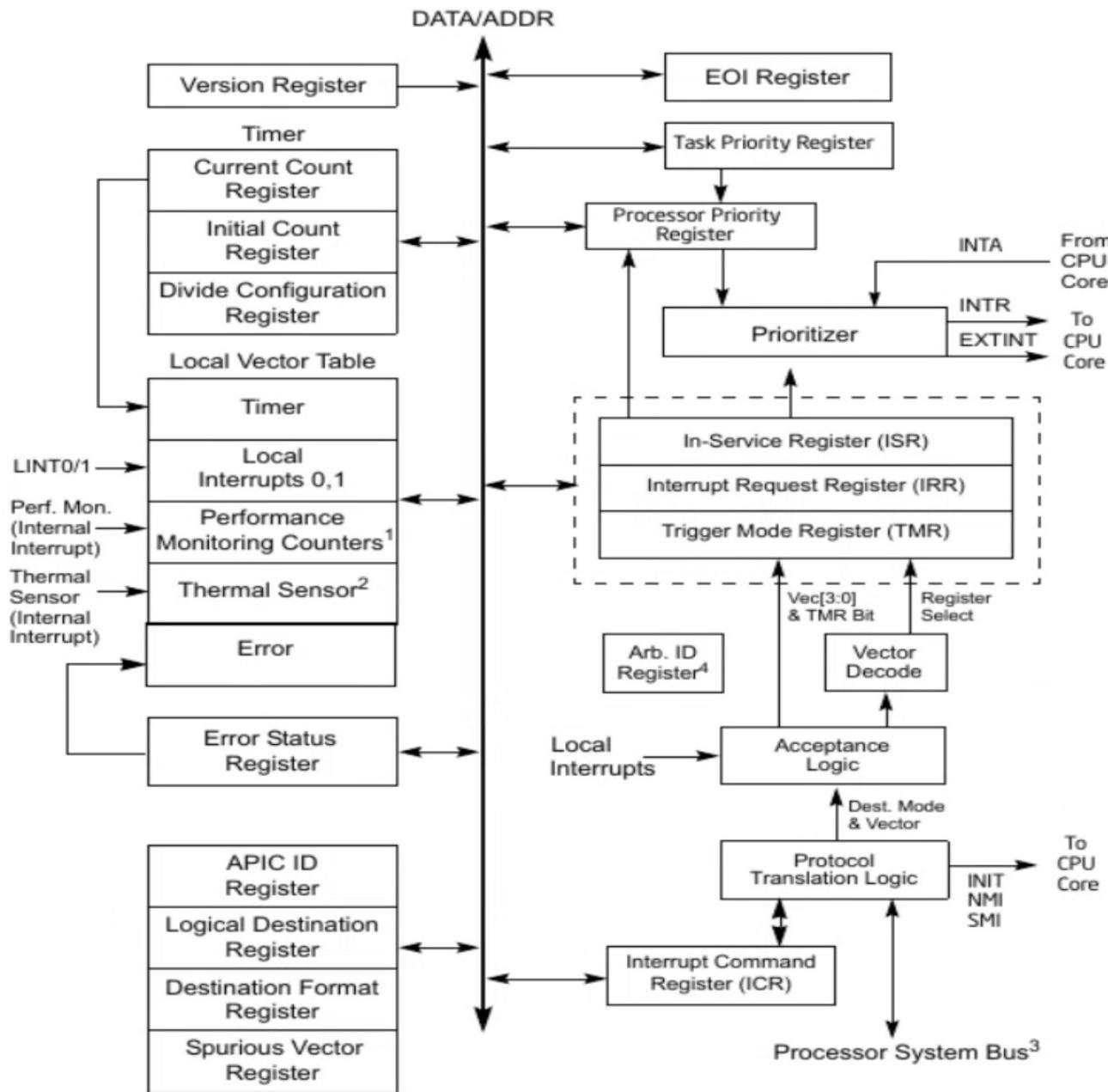
Контролер прерываний

- Если прерывание маскированно, выходим
- Если есть более приоритетные прерывания в обработке ставим очередь.
- Вызываем прерывание у процессора



C	D	Q	Operation
1	Di	Di	Follow
0	X	Qn-1	Hold

APIC



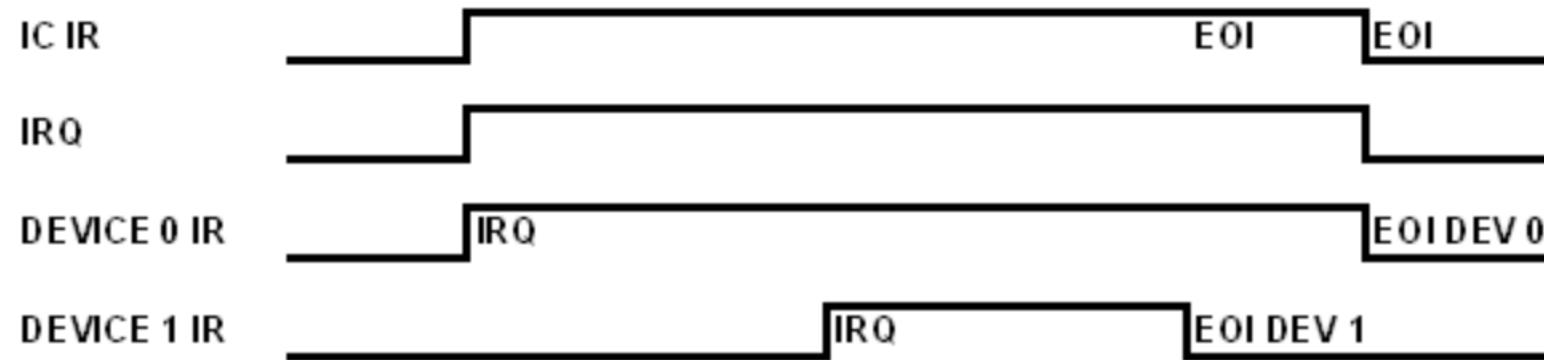
1. Introduced in P6 family processors.

2. Introduced in the Pentium 4 and Intel Xeon processors.

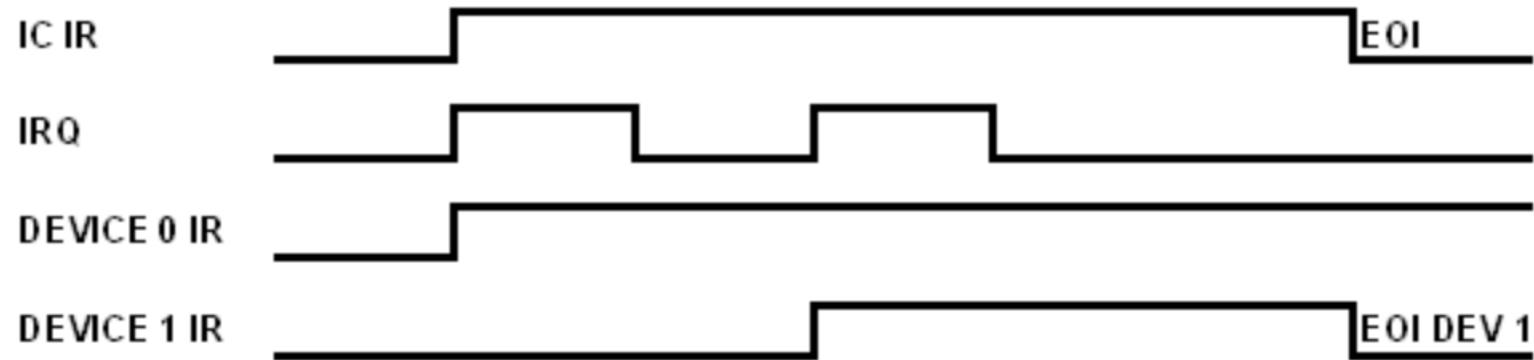
3. Three-wire APIC bus in P6 family and Pentium processors.

4. Not implemented in Pentium 4 and Intel Xeon processors.

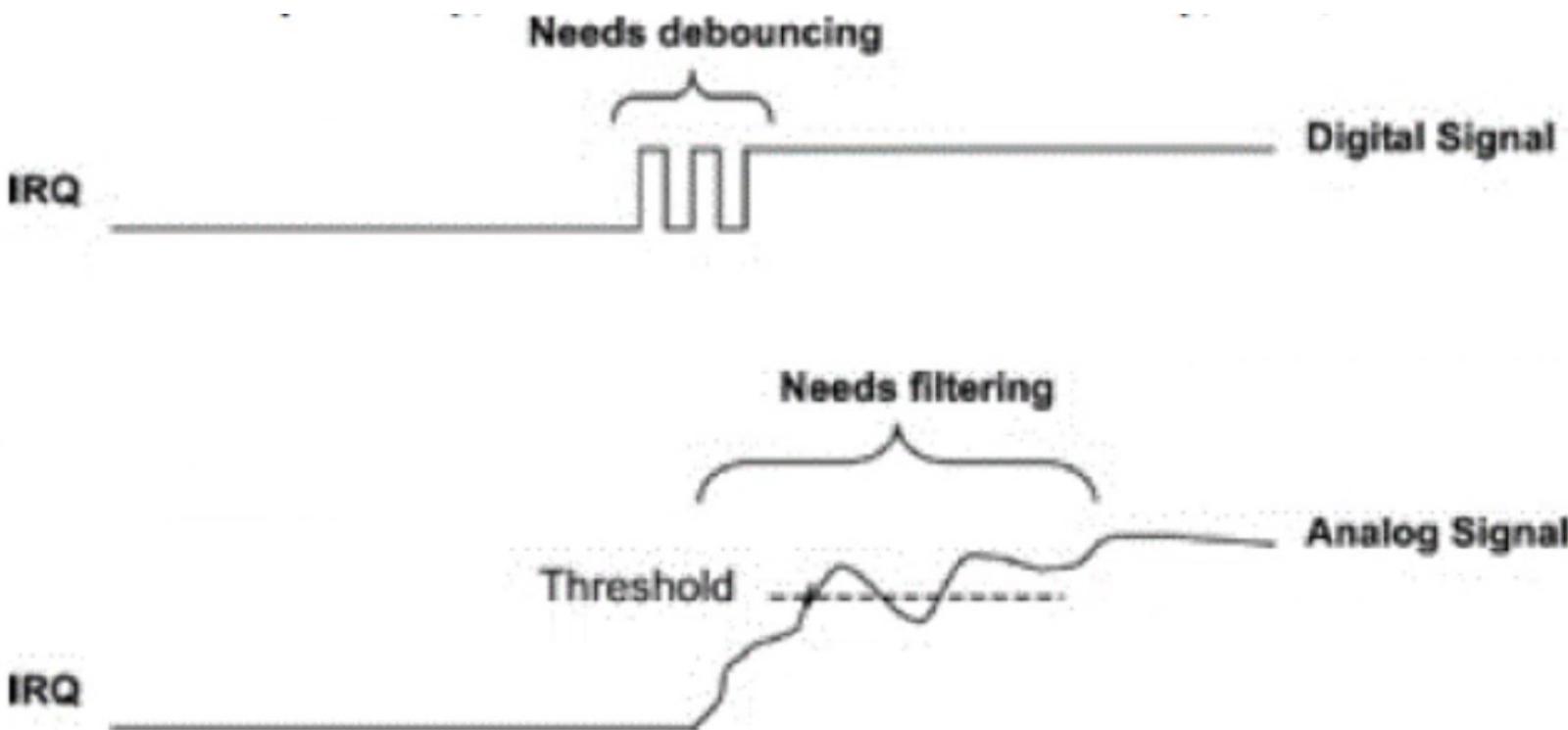
Прерывания по уровню



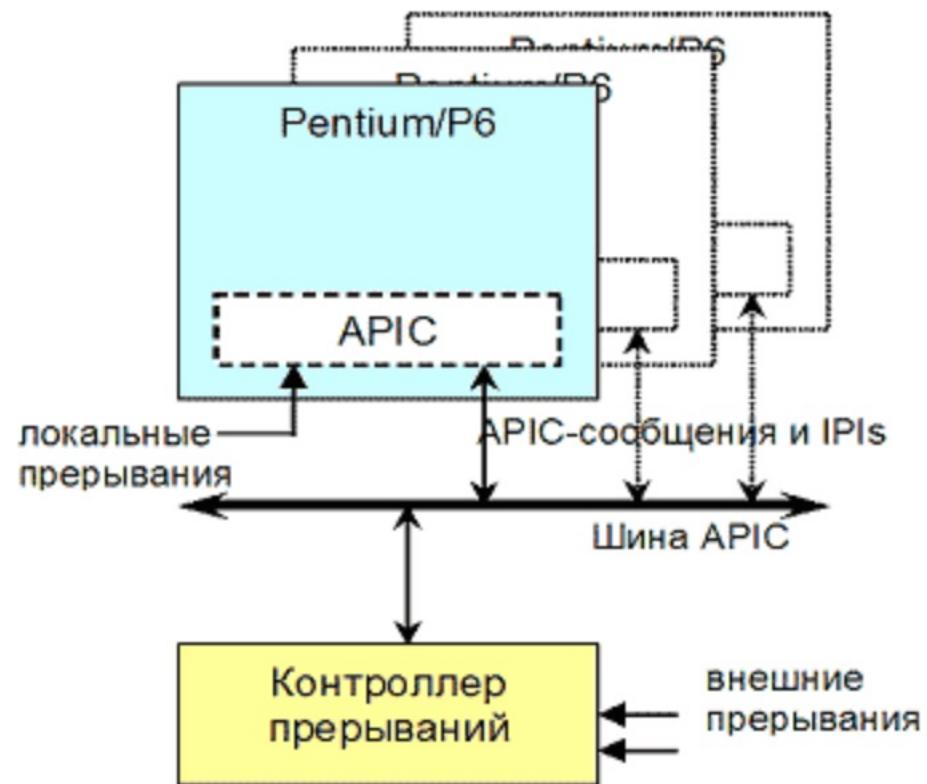
Прерывания по фронту



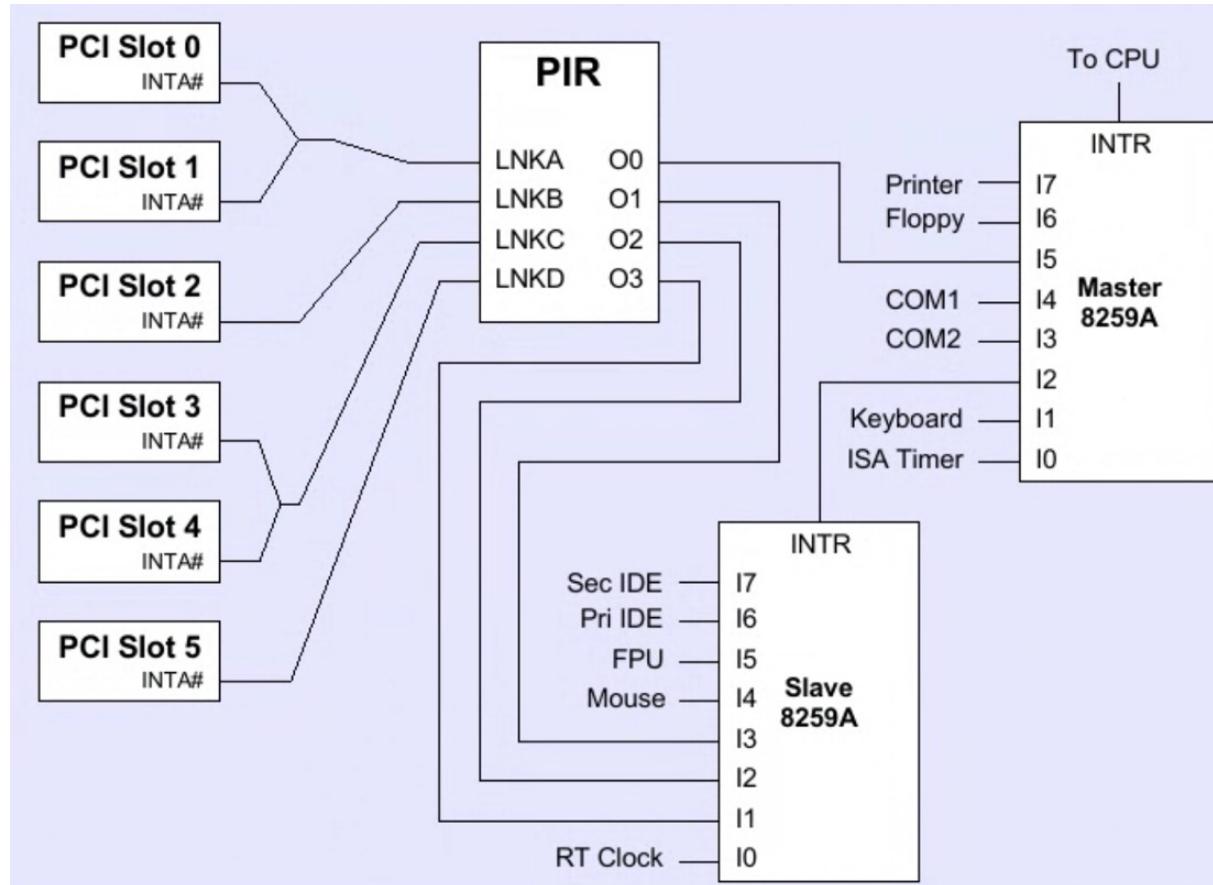
Spurious interrupt



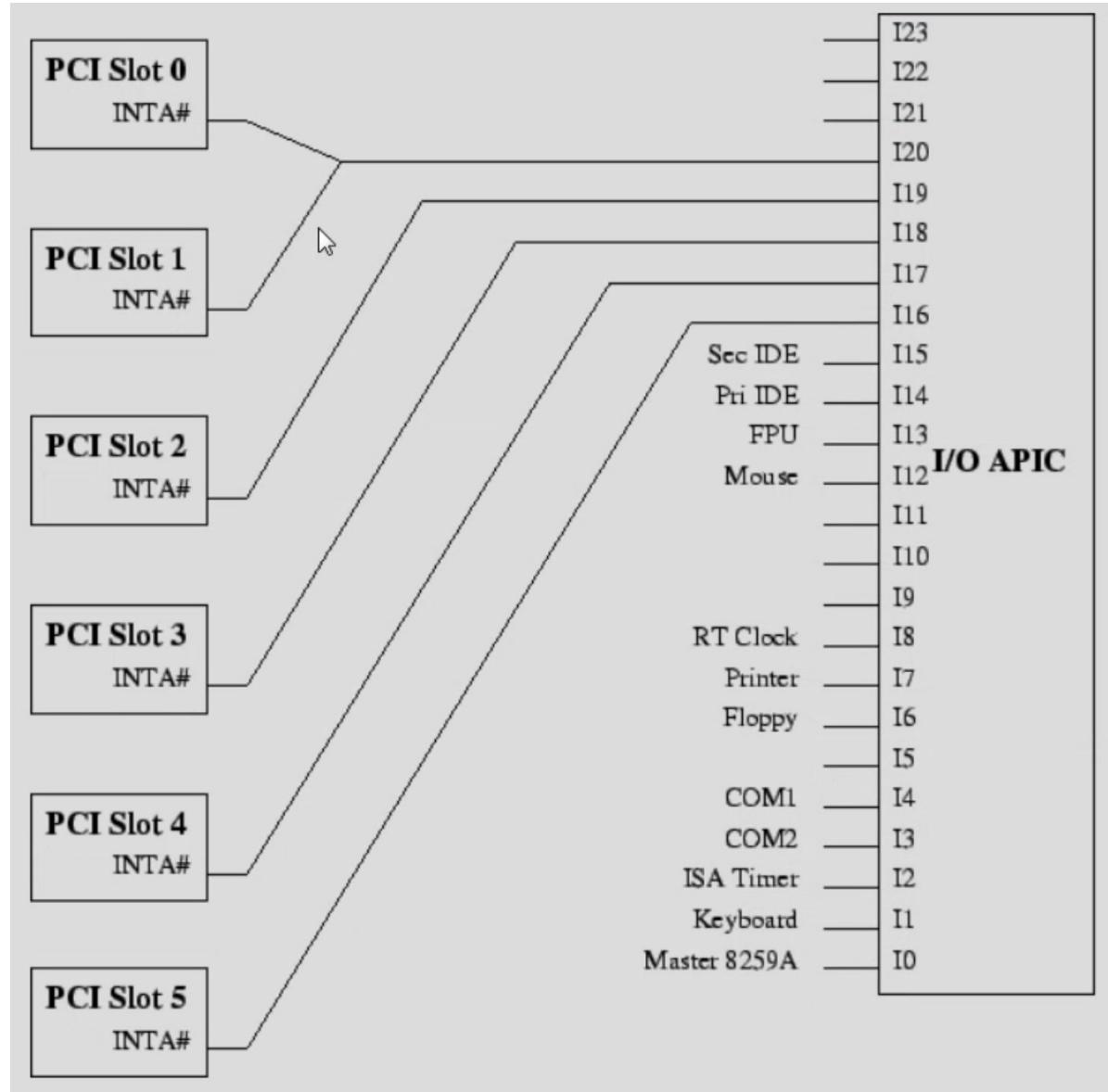
Общая схема обработки прерываний



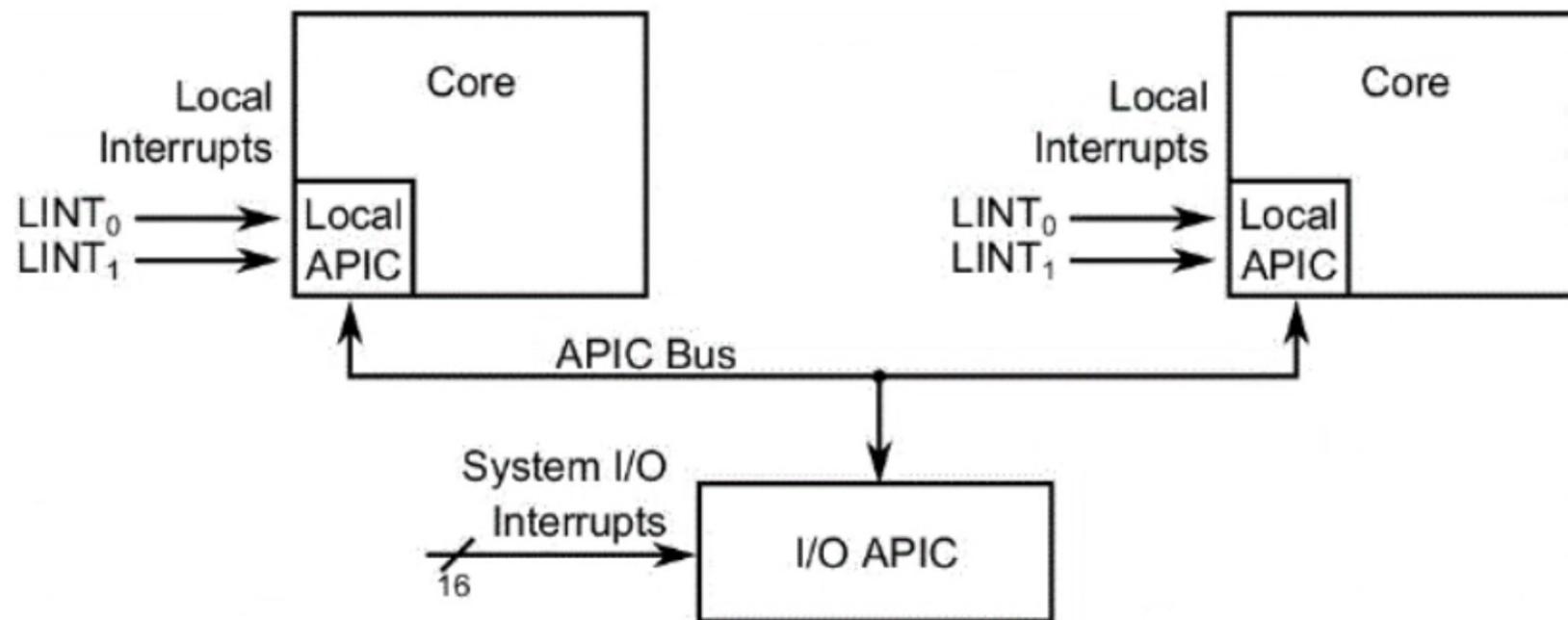
PIR – PCI Interrupt routing



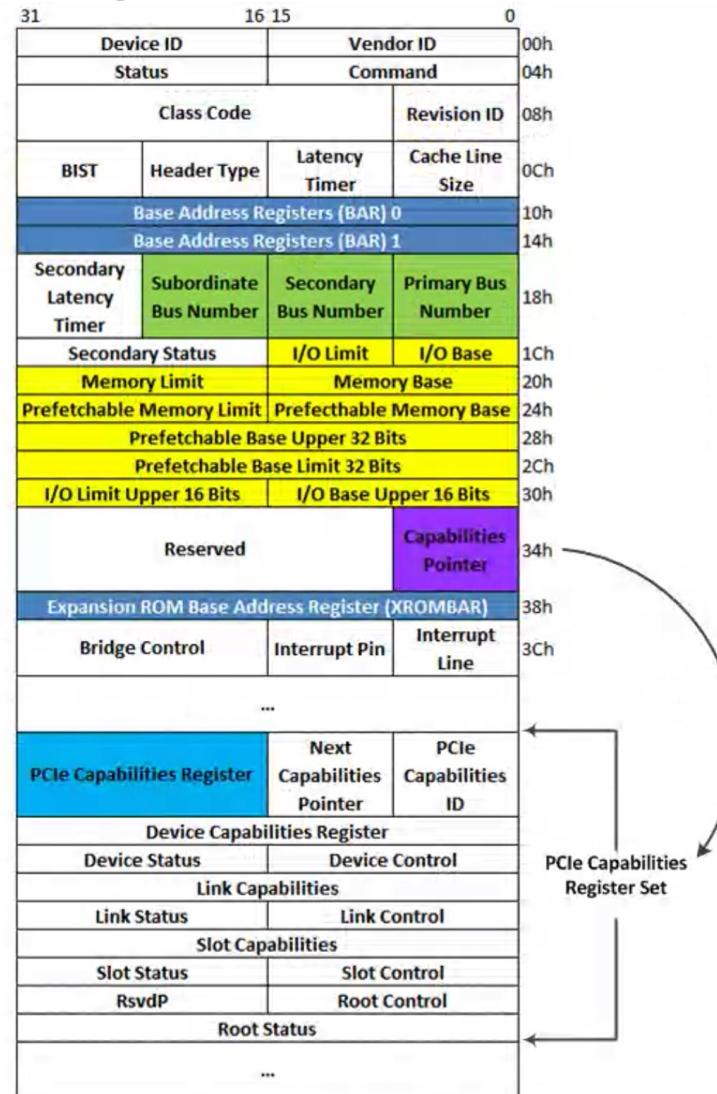
i/o apic



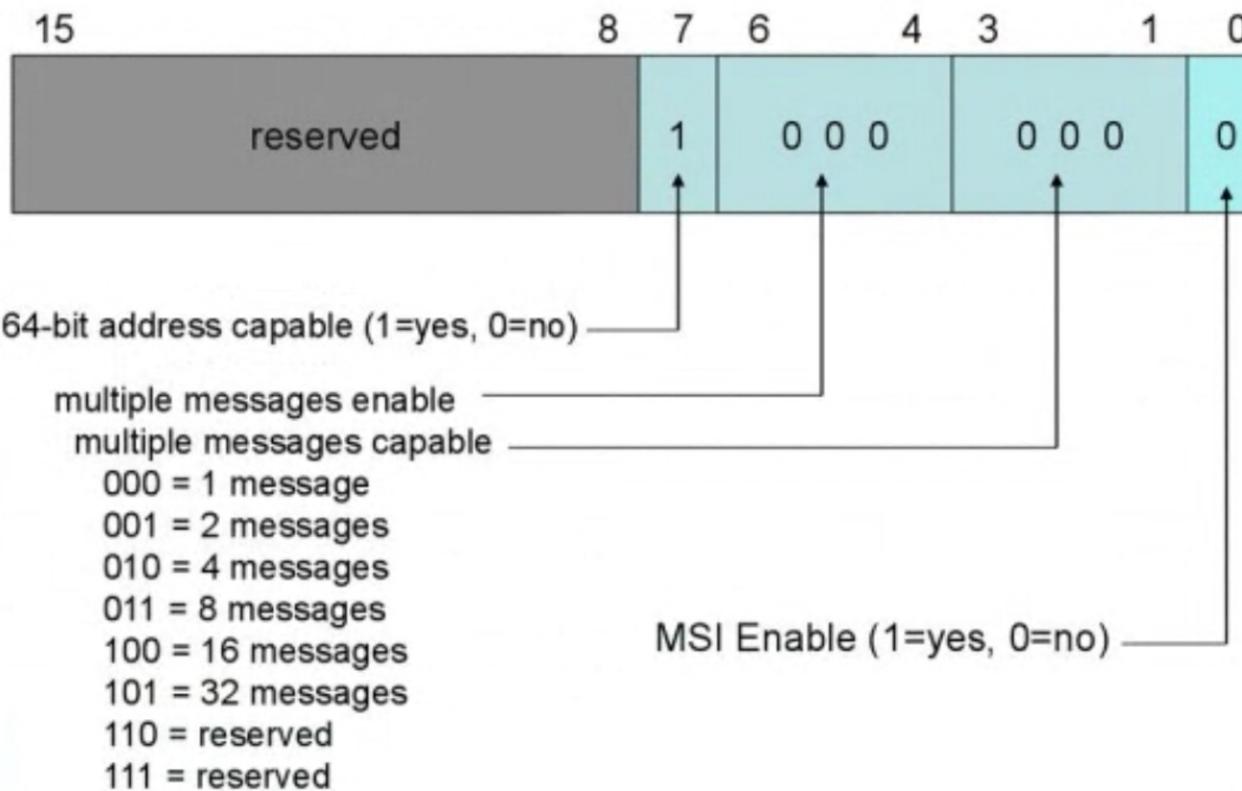
i/o apic



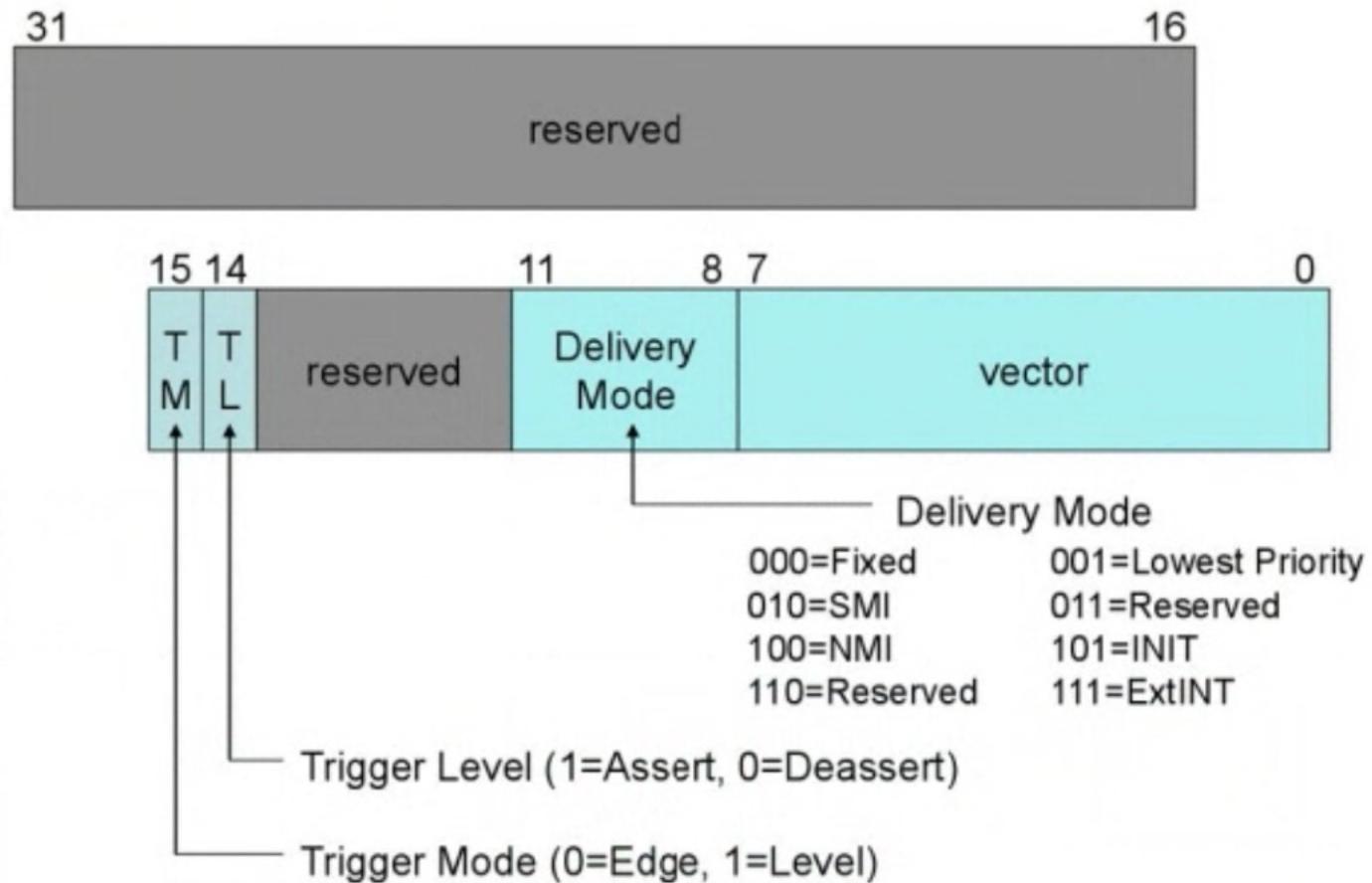
MSI – Message-signal interrupt



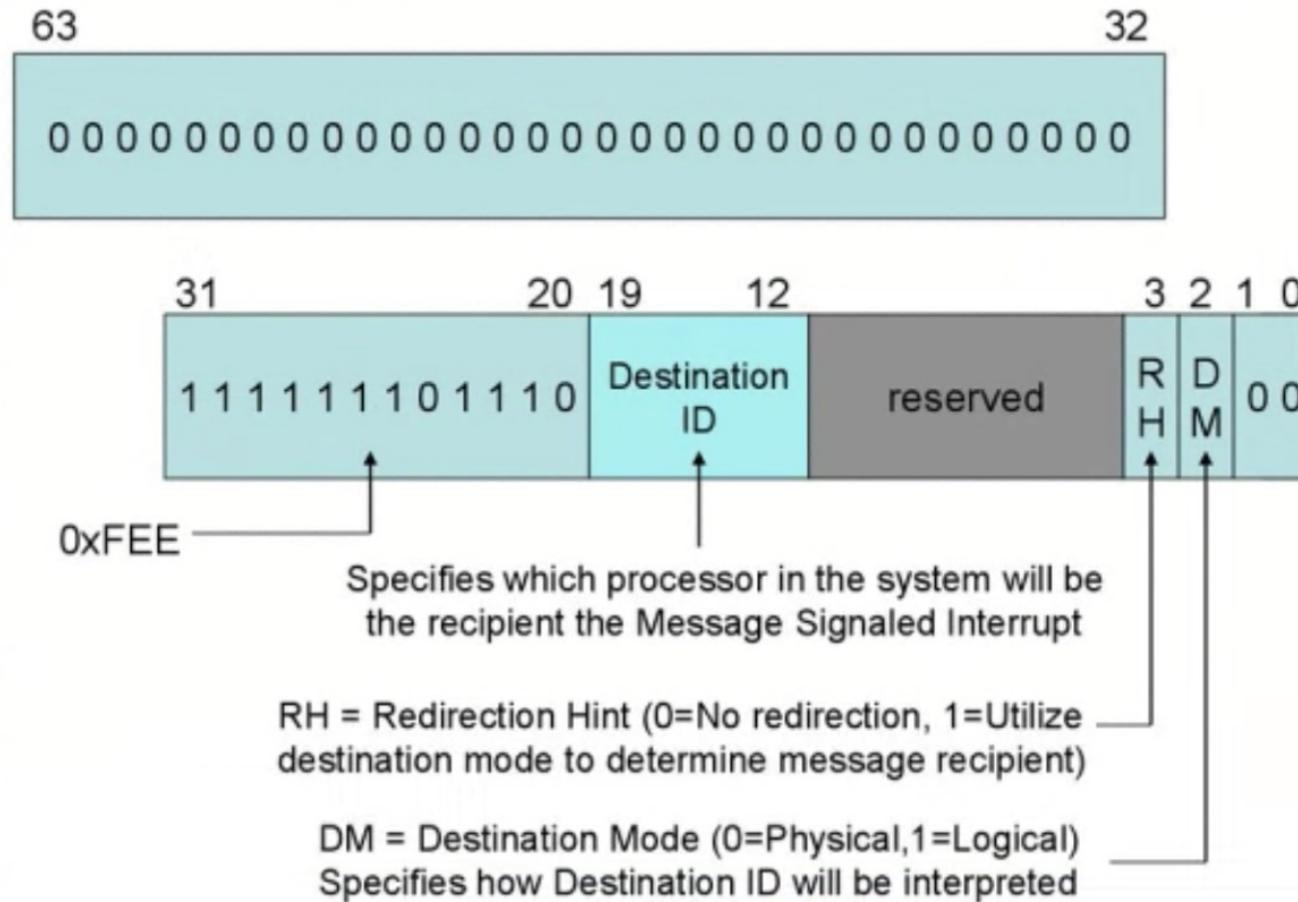
MSI control register



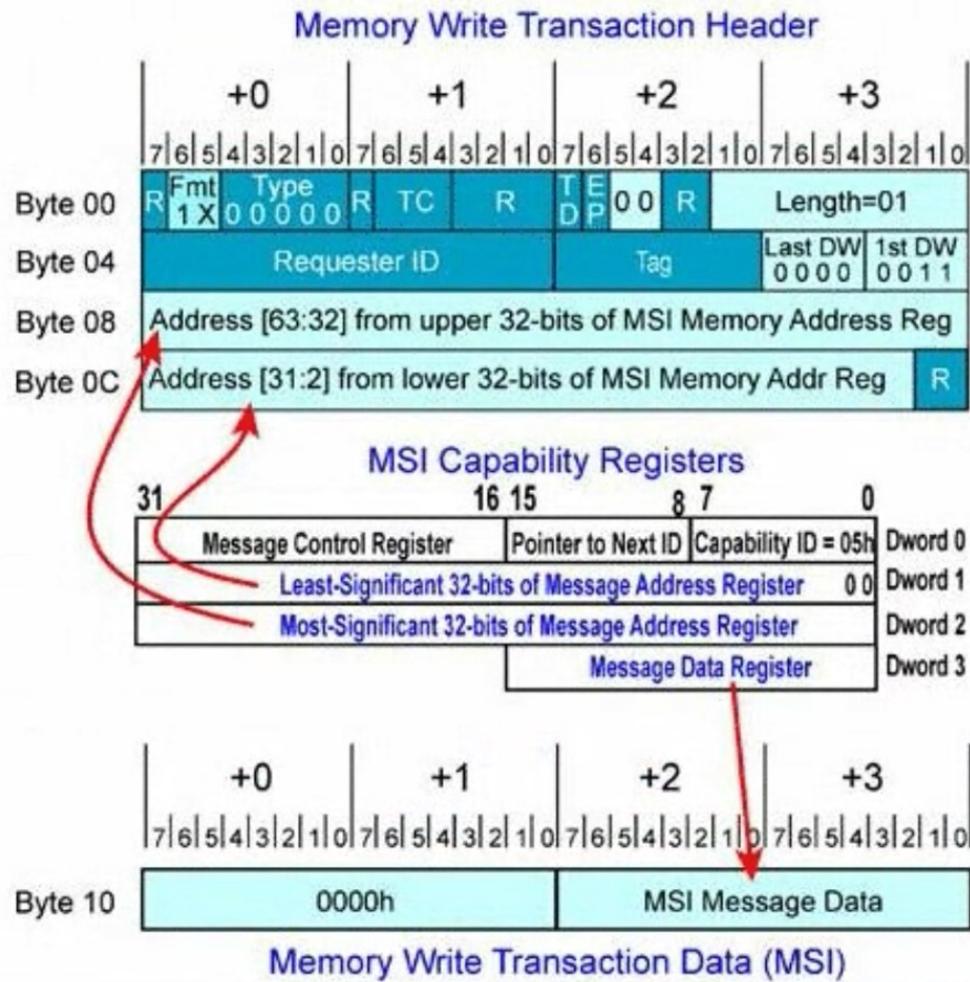
MSI data register



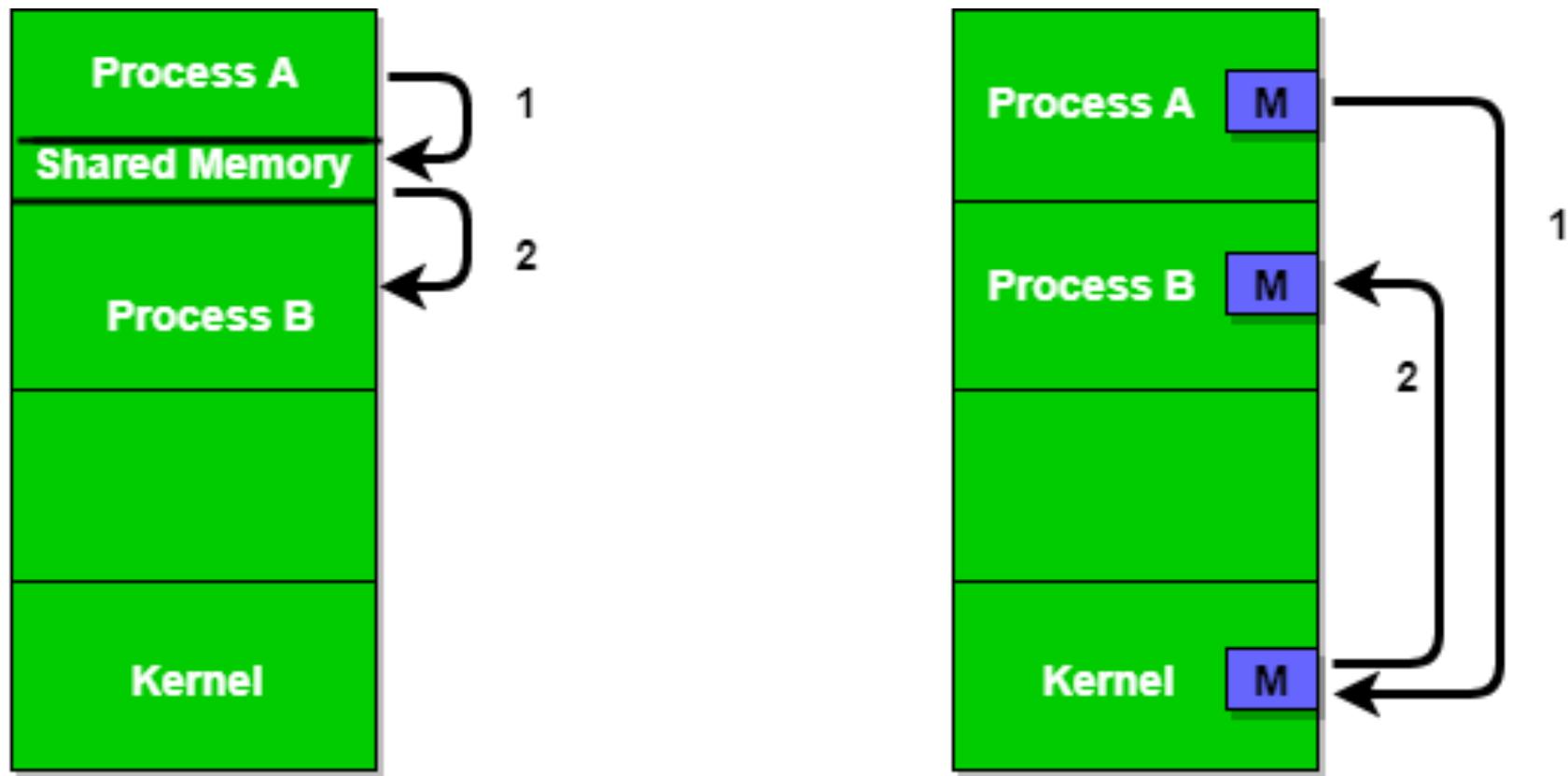
MSI address register



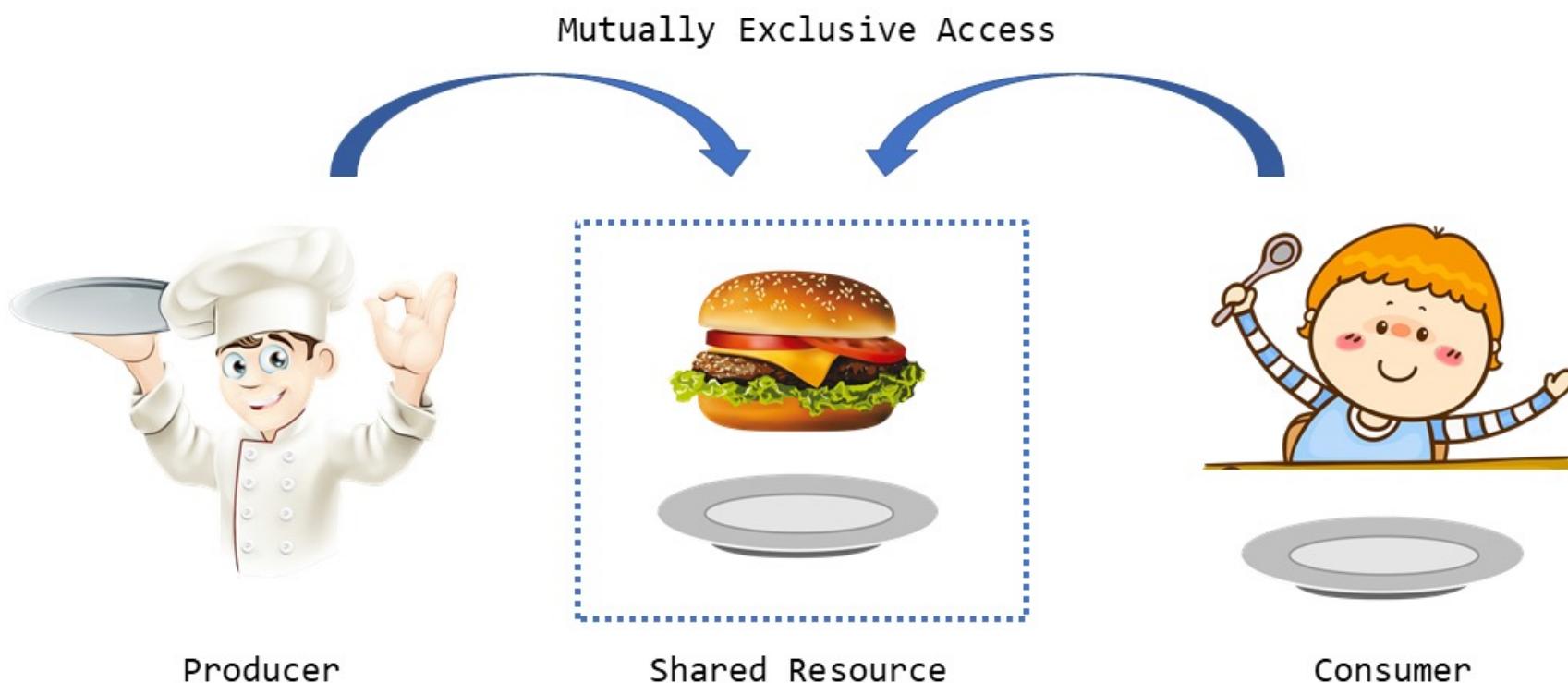
MSI memory transaction



IPC



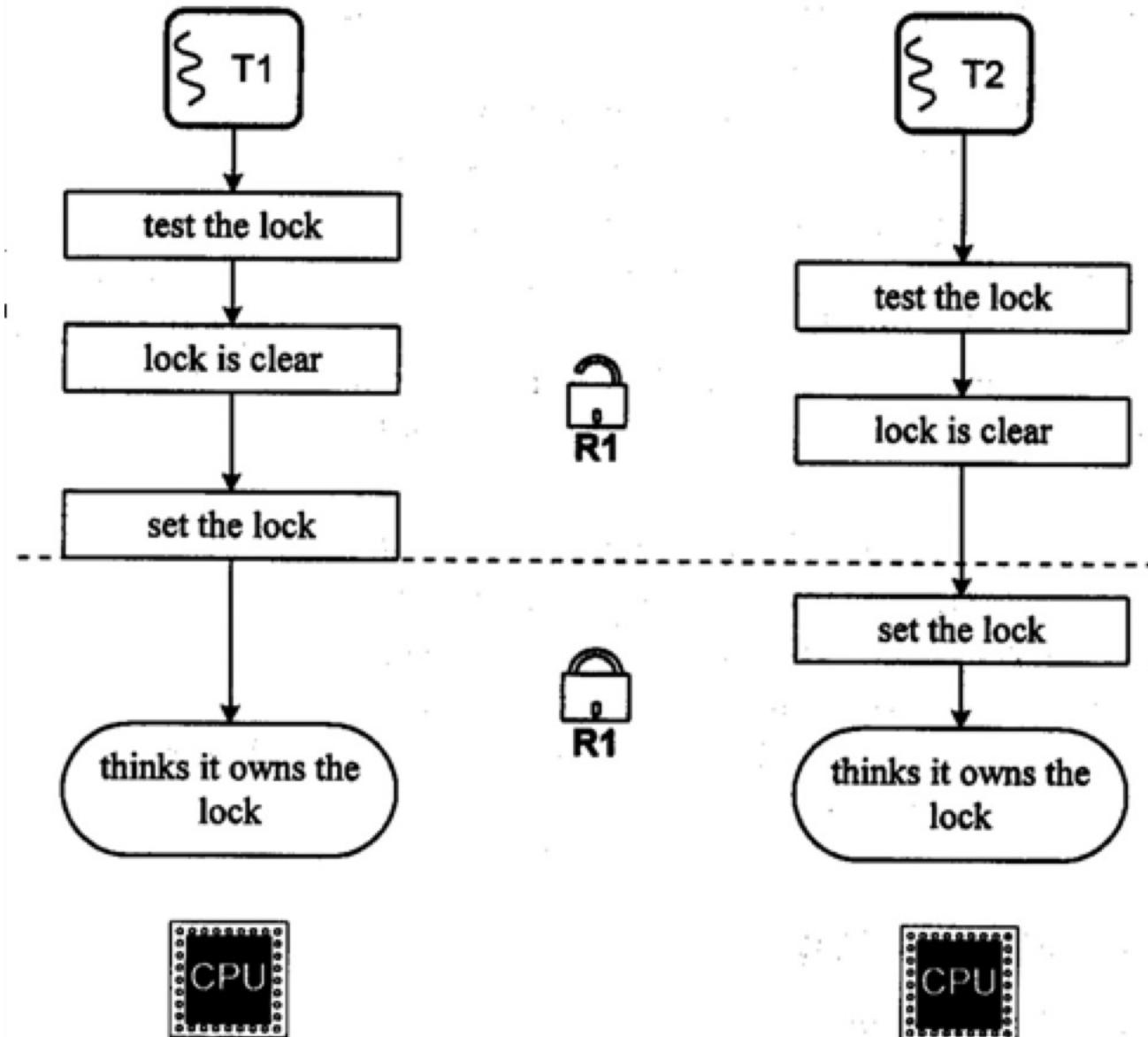
Producer-Consumer problem



spin/wait

spin mutex

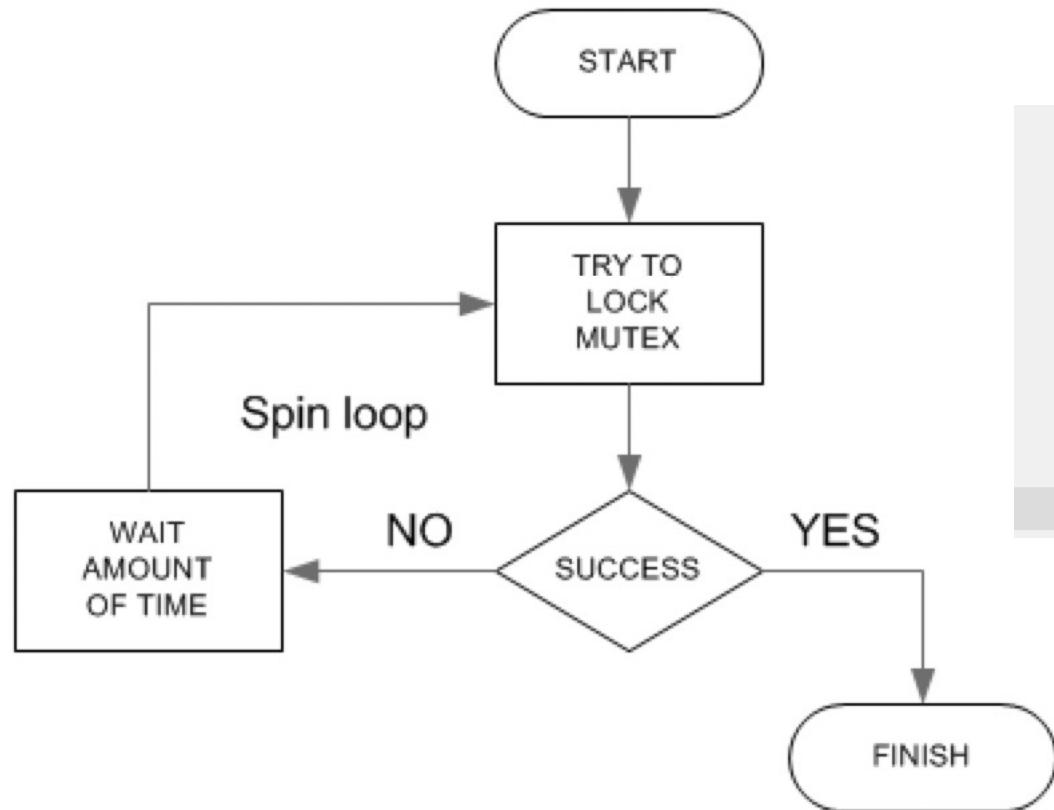
- init
- trylock
- lock
- unlock



spin mutex

```
1 |-----  
2 mutex_init(&mutex, MUTEX_UNLOCKED)      mov [mutex], 0  
3 mutex_init(&mutex, MUTEX_LOCKED)        mov [mutex], 1  
4 -----  
5 bool    mutex_trylock(&mutex)  
6     .....  
7         mov eax, 1  
8         lock xchg  [mutex],eax  
9         or  eax,eax  
10        mov eax, FALSE  
11        mov ebx, TRUE  
12        cmovz  eax, ebx  
13        ret  
14 -----  
15 mutex_unlock(&mutex)                 lock mov  [mutex],0
```

spin lock



```
1 spinmutex_lock(&mutex)
2
3
4
5
6
7
8
9
10
```

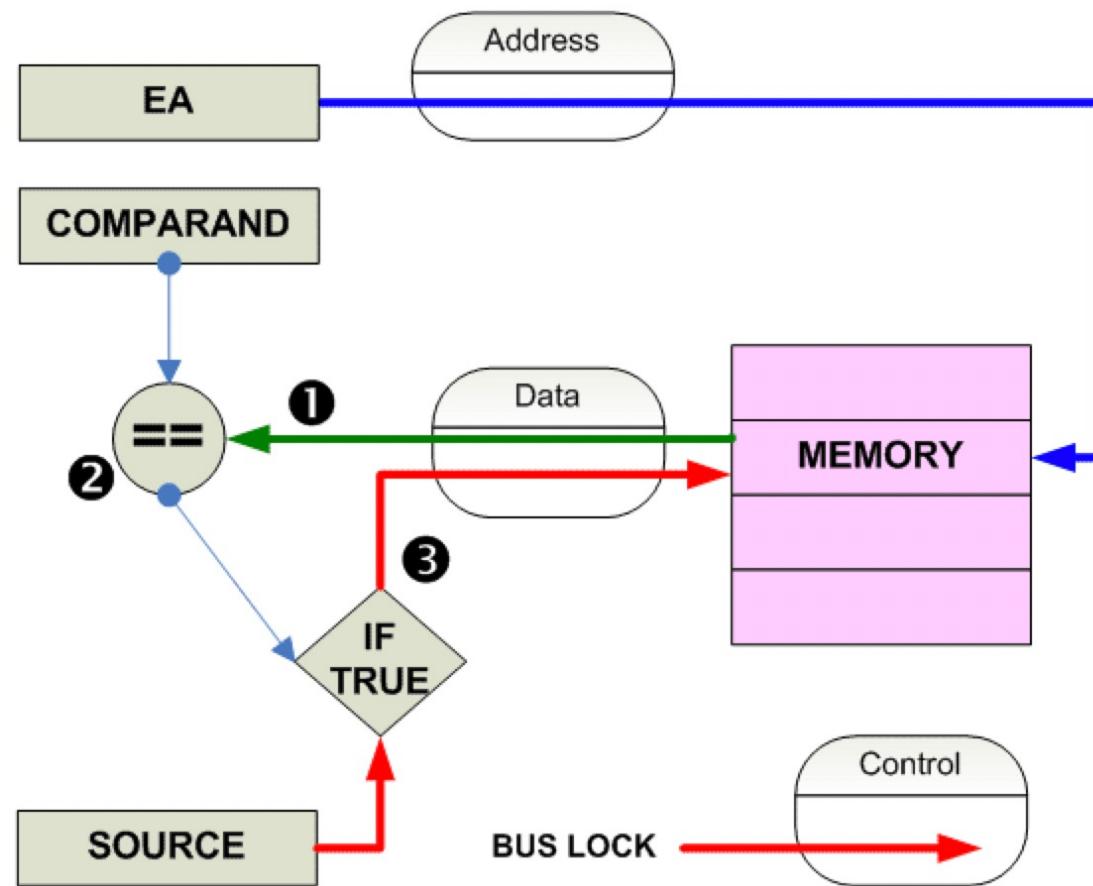
?spin_loop:
mov eax,1
lock xchg [mutex],eax
or eax,eax
jz ?lock_success
;; wait
jmp ?spin_loop
?lock_success:
ret

spin lock

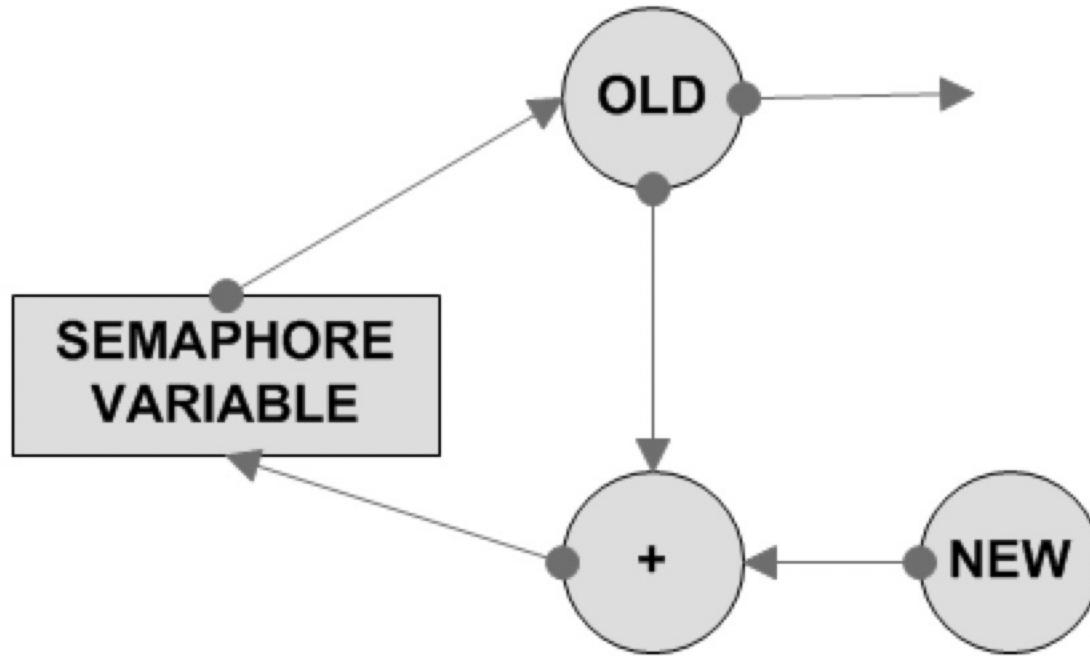
```
1 spinmutex_lock(&mutex)
2
3         ?spin_loop:
4             mov eax,1
5             lock xchg [mutex],eax
6             or eax,eax
7             jz ?lock_success
8             ;; wait
9             jmp ?spin_loop
10            ?lock_success:
11                ret
```

```
1 spinmutex_lock(&mutex)
2
3         ?spin_loop:
4             mov eax,1
5             lock xchg [mutex],eax
6             or eax,eax
7             jz ?lock_success
8             xor ecx,ecx
9             xor edx,edx
10            lea eax,[mutex]
11            monitor
12            mwait
13            jmp ?spin_loop
14            ?lock_success:
15                ret
```

Atomics



semaphore



semaphore

```
1-----  
2 semaphore_init(&semaphore)          mov [semaphore], 0  
3-----  
4 bool    semaphore_trylock(&semaphore)  
5           mov eax, 1  
6           lock xadd    [semaphore], eax  
7           cmp eax, MAX_VALUE  
8           mov eax, TRUE  
9           jnc ?success  
10          mov eax, -1  
11          lock xadd    [semaphore],eax  
12          mov eax, FALSE  
13 ?success:  
14          ret  
15-----  
16 semaphore_unlock(&semaphore)      lock dec [semaphore]  
17-----|
```

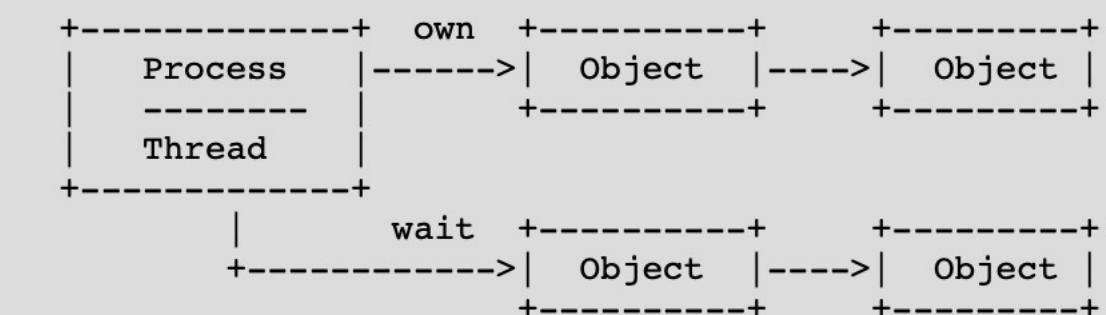
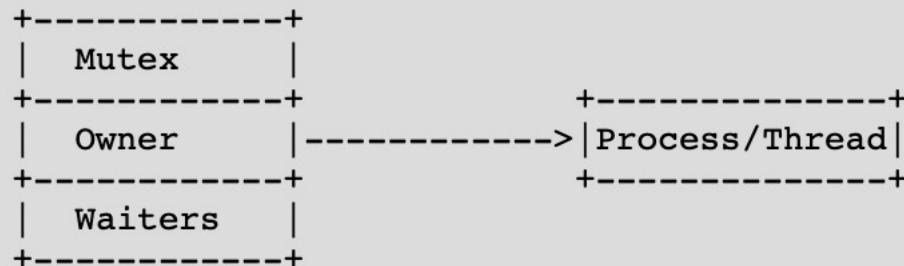
semaphore

```
1 |-----  
2 bool semaphore_trylock(&semaphore)      spinmutex_lock(&mutex)  
3 |-----  
4     if (*semaphore < MAX_VALUE)  
5     {  
6         (*semaphore)++;  
7         mutex_unlock(&mutex);  
8         return TRUE;  
9     }  
10    mutex_unlock(&mutex);  
11    return FALSE;  
12 -----  
12 semaphore_unlock(&semaphore)      spinmutex_lock(&mutex)  
13 |-----  
14     (*semaphore)--;  
15     mutex_unlock(&mutex)  
16 -----  
16 spinsemaphore_lock(&semaphore)      while(semaphore_trylock(&semaphore) != TRUE) ;  
17 -----
```

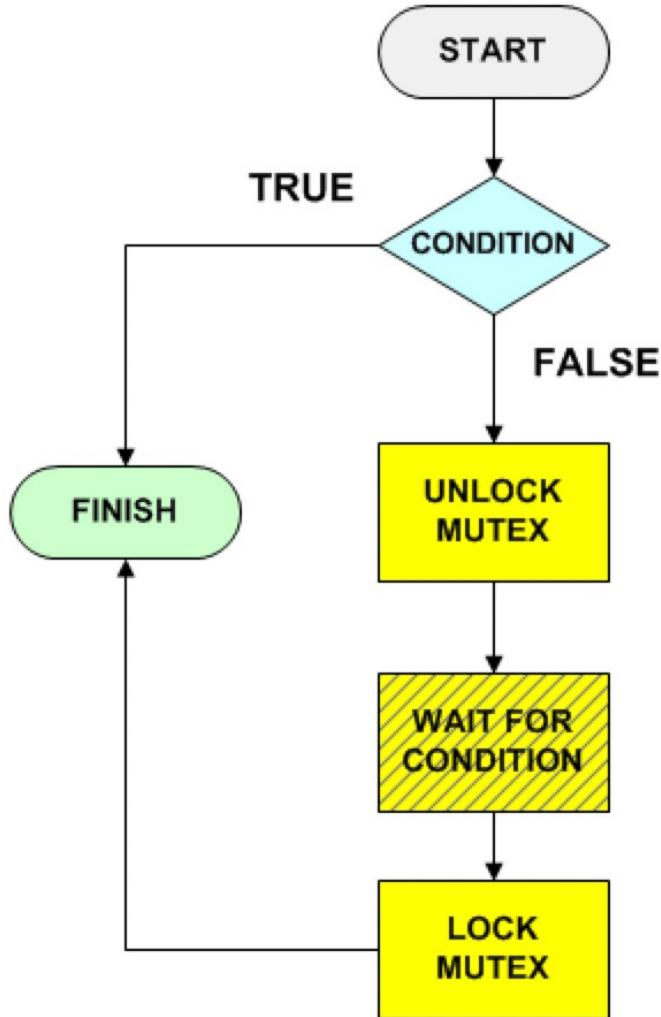
1 proc

```
1
2     bool    sp_mutex_trylock(&mutex)
3             cli
4             mov eax,[mutex]
5             mov [mutex], 1
6             sti
7             ret
8
9     bool    sp_semaphore_trylock(&mutex)
10            cli
11            mov eax, [sem]
12            add eax, 1
13            mov [sem], eax
14            cmp eax, MAX_VALUE
15            jnc ?success
16            sub eax, 1
17            mov [sem], eax
18            mov eax, FALSE
19            sti
20            ret
21        ?success:
22            mov eax, TRUE
23            sti
24            ret
25
```

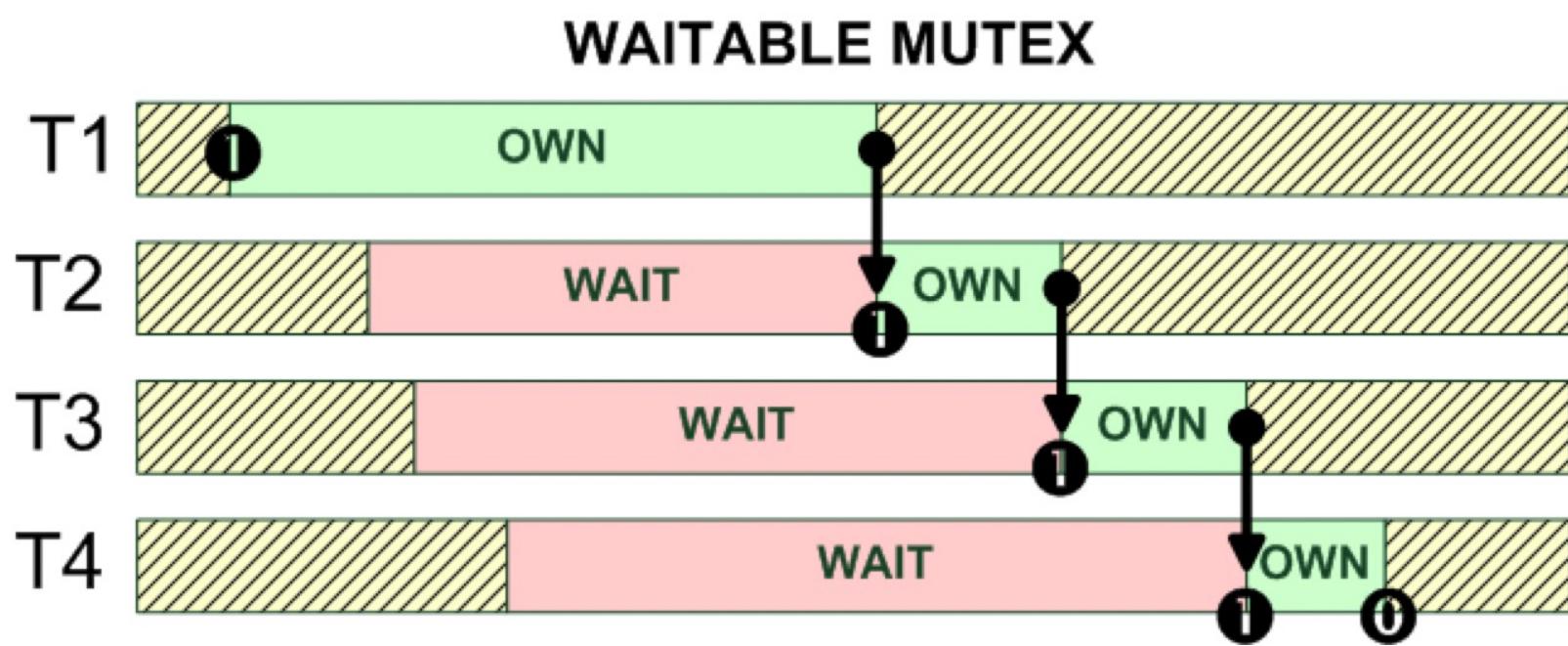
writable objects



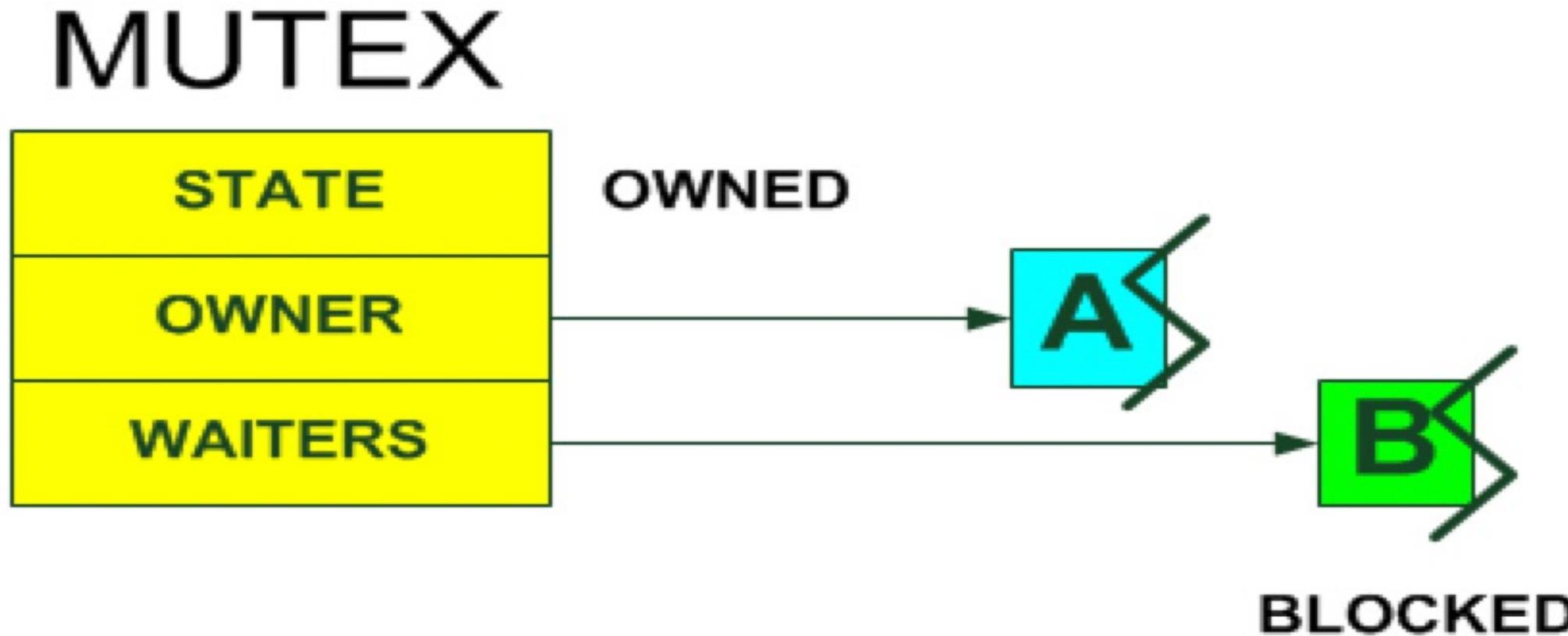
waitable condvar



waitable mutex



waitable mutex



APC (APPLICATION SPECIFIC CALLBACKS)

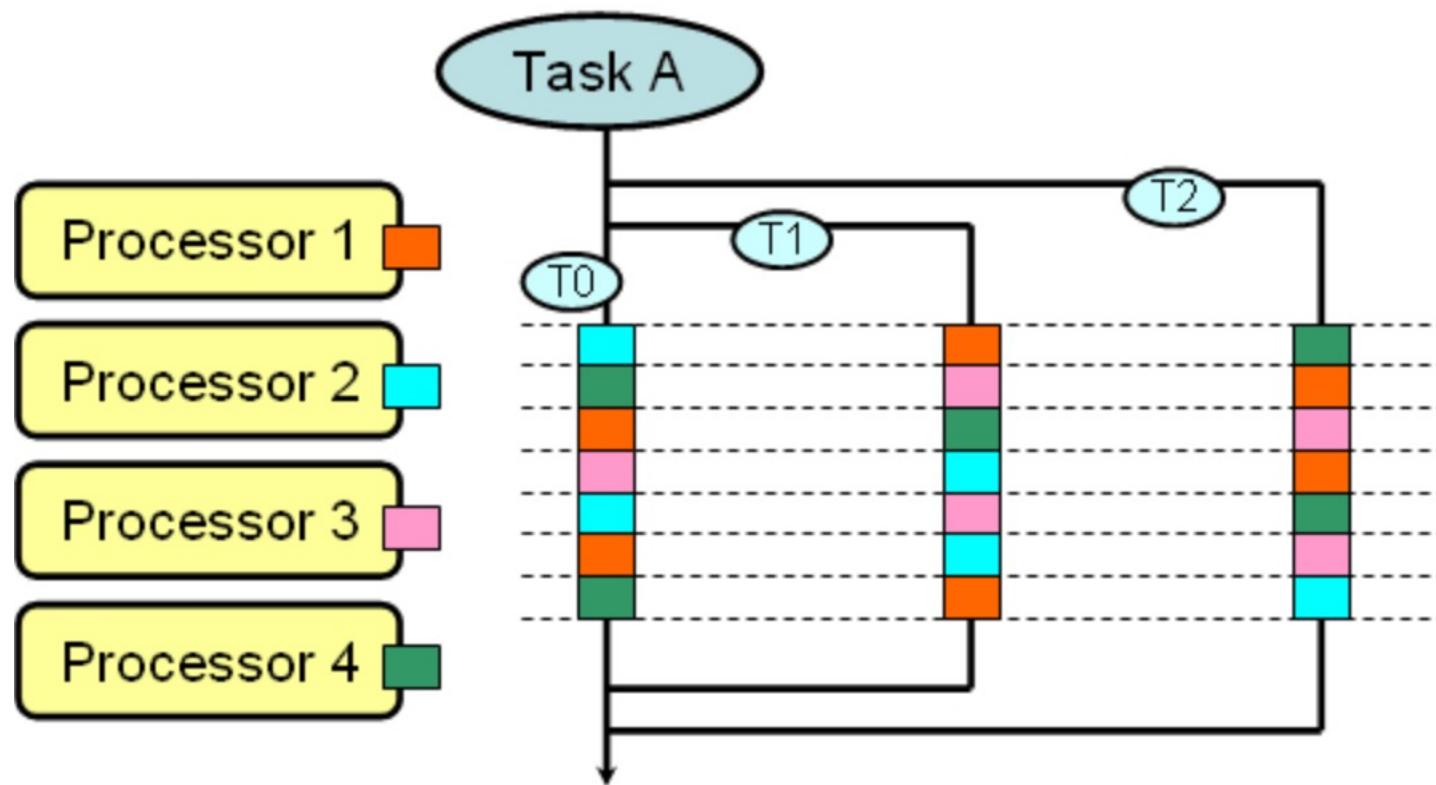
```
1 // POSIX
2
3 result = sigaction(int signalnum,
4                     const struct sigaction* action,
5                     struct sigaction* prevaction);
6
7
8     struct sigaction
9     {
10         void (*sa_handler)();
11         void (*sa_sigaction)(int, siginfo_t*, void*);
12         sigset_t sa_mask;
13         int sa_flags;
14             SA_RESETHAND
15             SA_NODEFER
16             SA_RESTART
17             SA_SIGINFO
18
19     };
```

```
1 //| signal.h
2 void* signal(int signum, void* handler);
3             SIG_IGN          // ignore
4             SIG_DFL          // default action
5
6     int    raise(int signum);
7
8     int    raise(int signum)
9     {
10         return kill(getpid(), signum);
11     }
12
13     char*  strsignal(int signum)
14
15
16     union sigval
17     {
18         int sival_int;
19         void*   sival_ptr;
20     };
21
22 void* sigqueue(pid_t pid, int signum, const union sigval value);
```

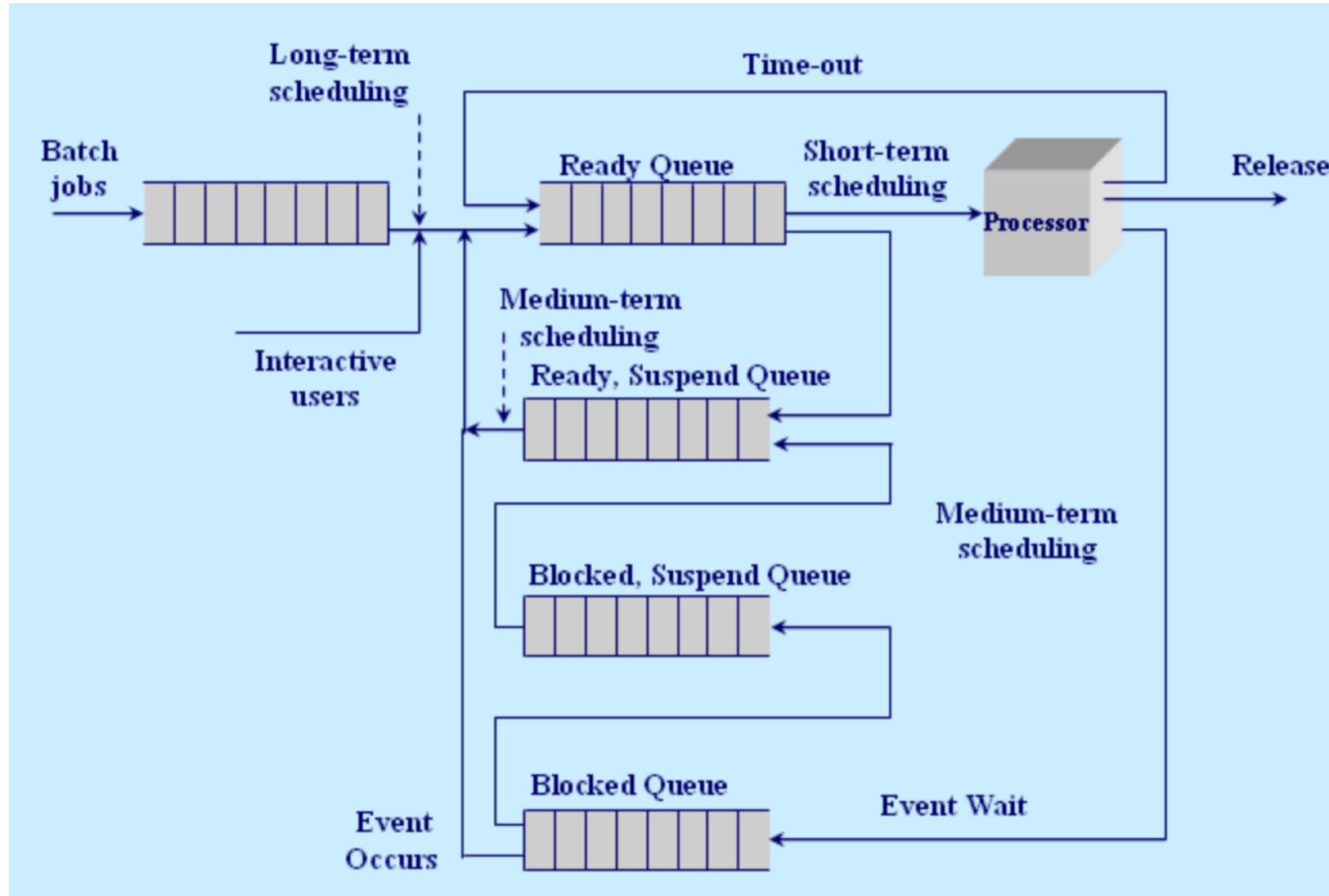
APC (APPLICATION SPECIFIC CALLBACKS)

```
1 NTSTATUS NTAPI NtQueueApcThread(IN HANDLE ThreadHandle,
2                               IN PUSER_APIC_ROUTINE ApcRoutine,
3                               IN PVOID FirstArg, // QueryUserAPC() APC address
4                               IN PVOID SecondArg, // QueryUserAPC() APC argument
5                               IN PVOID ThirdArg);
6
7
8 VOID NTAPI _NativeUserRoutine(PVOID First Arg, PVOID SecondArg, PVOID ThirdArg)
9 {
10    ((PAPCFUNC) FirstArg, ( (DWORD) SecondArg);
11 }
12
13 NTSTATUS NtAlertThread(HANDLE ThreadHandle);
14 NTSTATUS NtTestAlert(VOID);
```

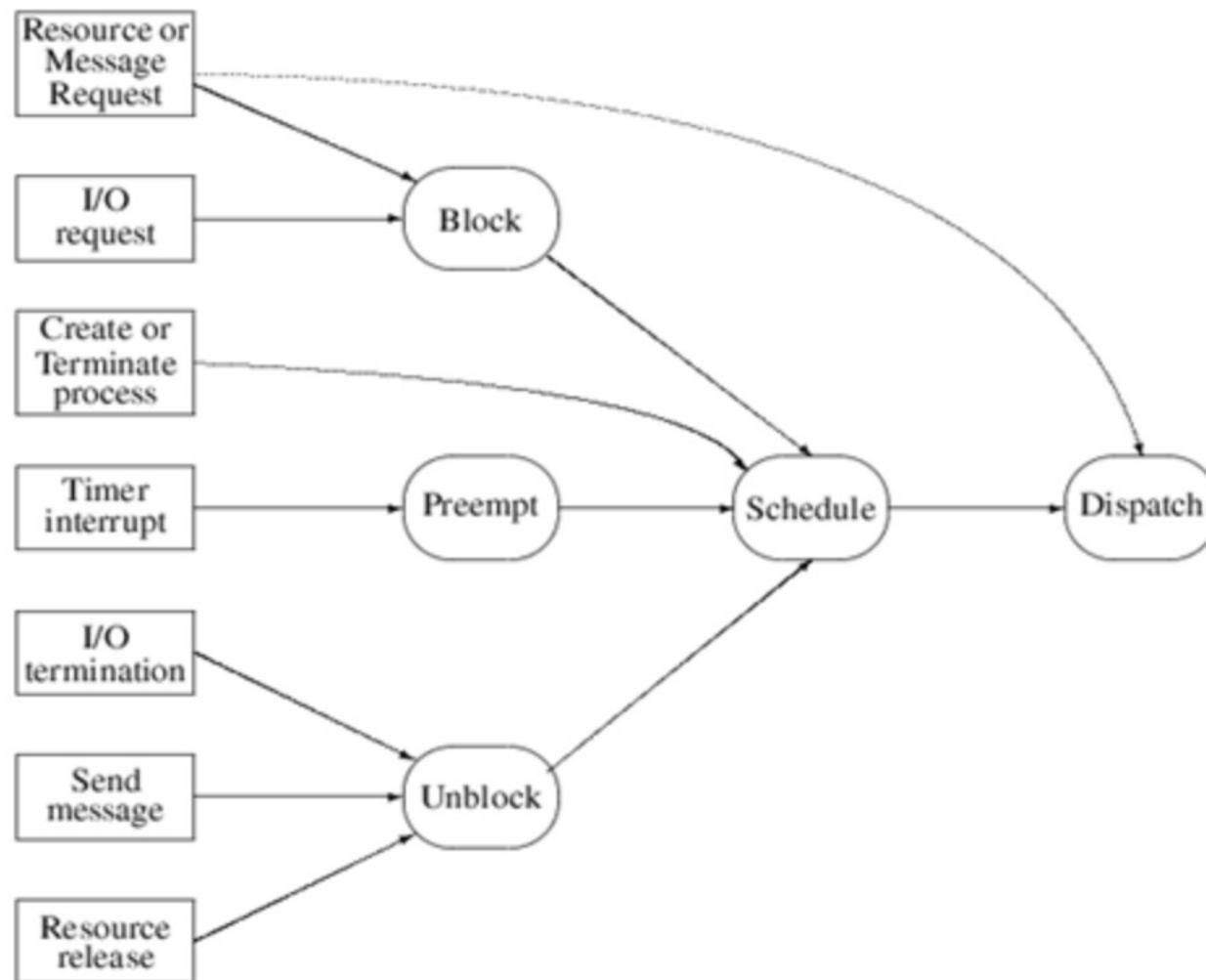
Планировка задач



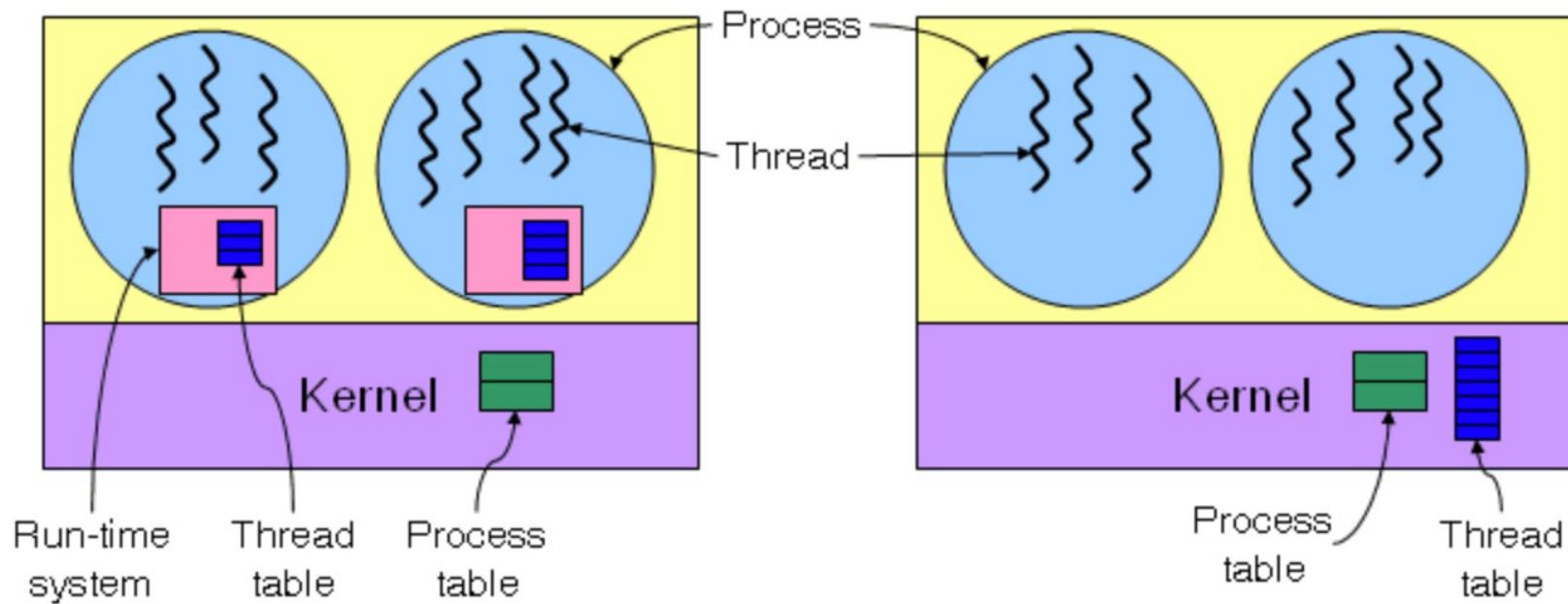
Планировщик



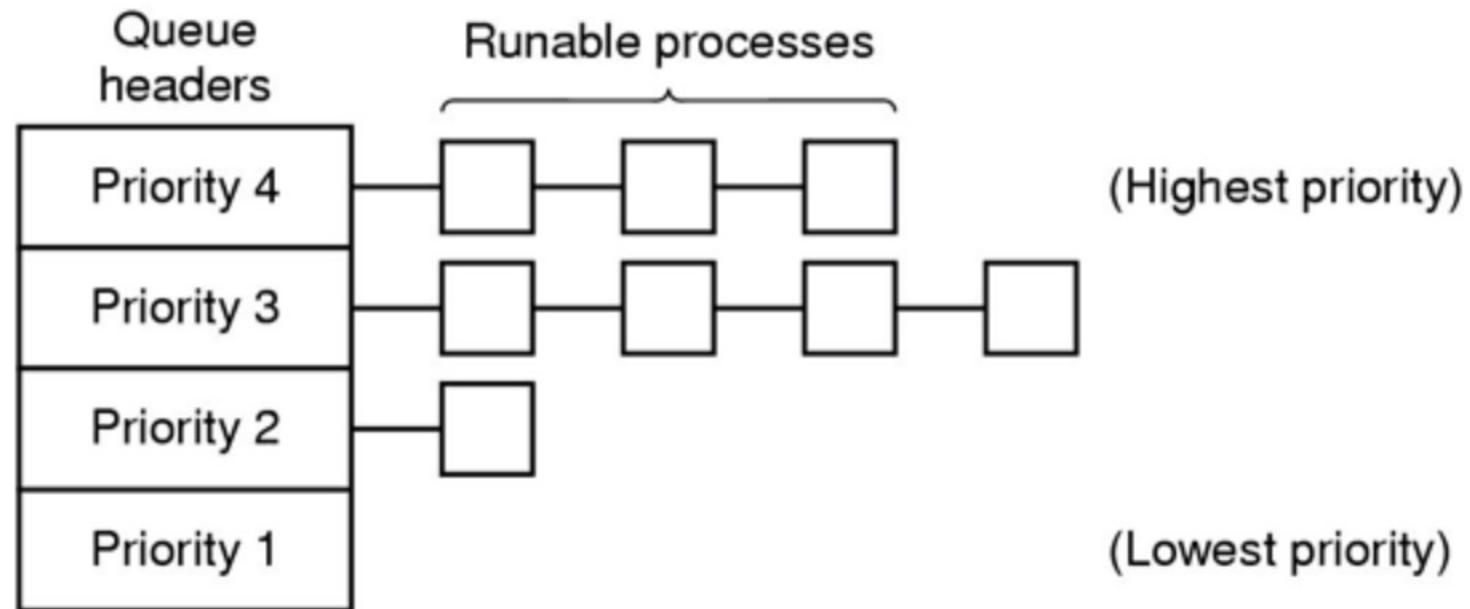
Перепланировка



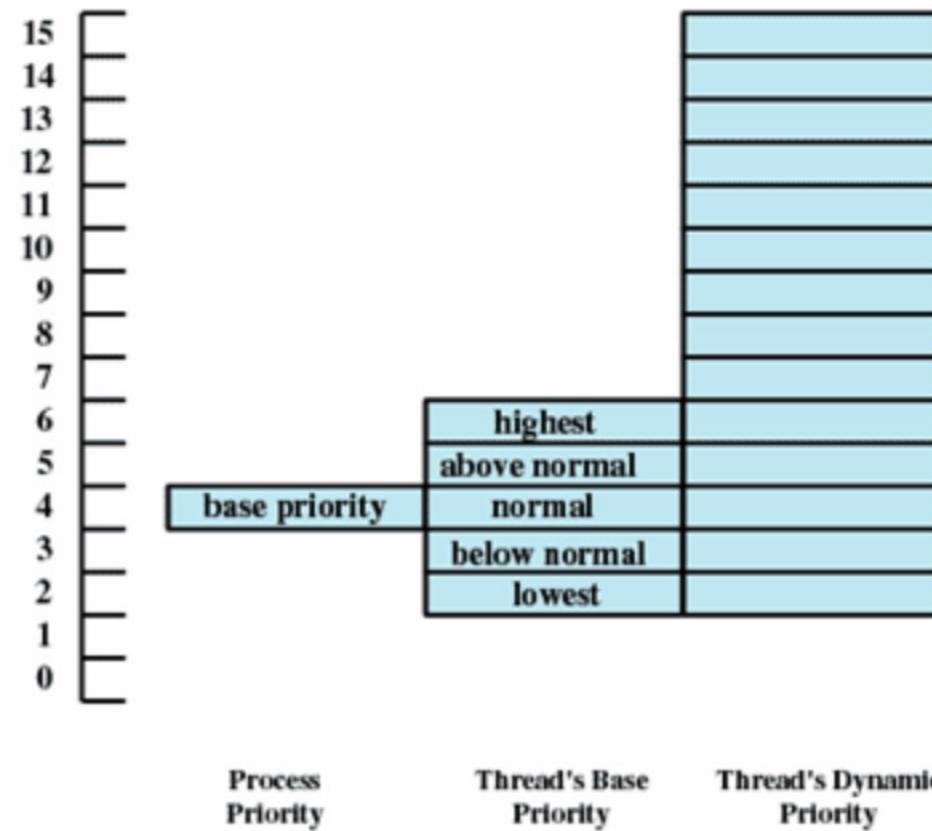
Стратегии планирования



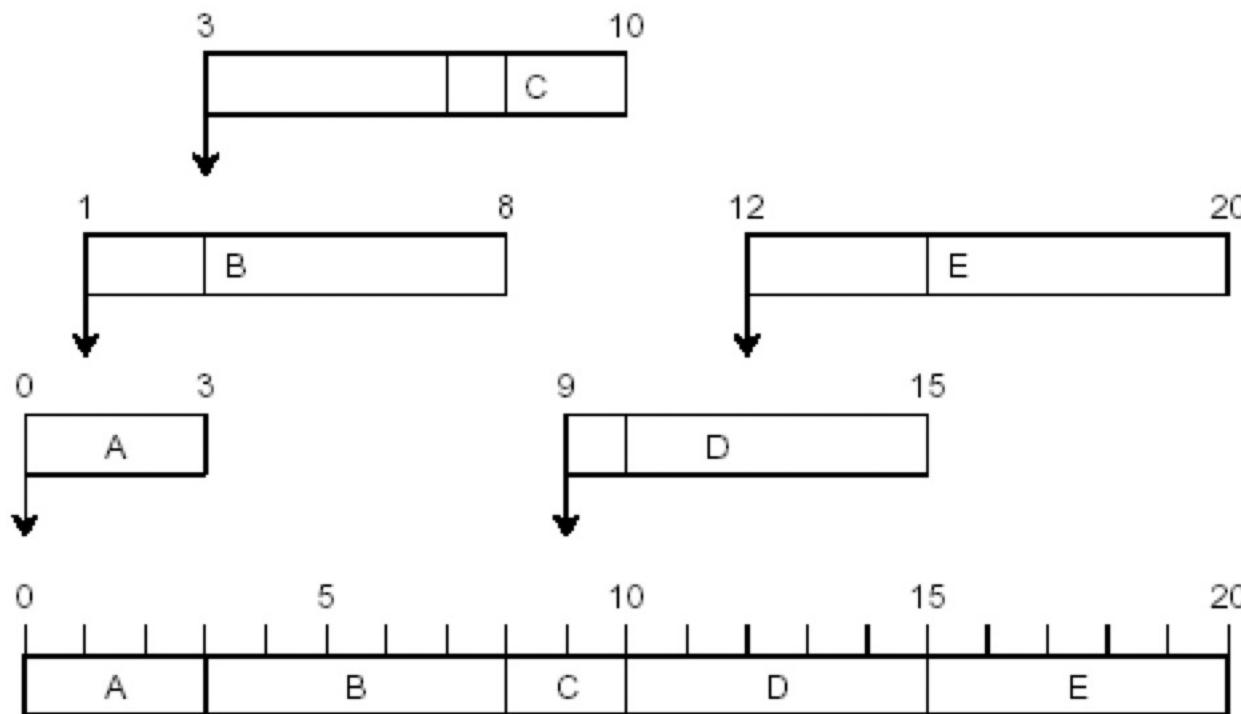
Планировка с приоритетами



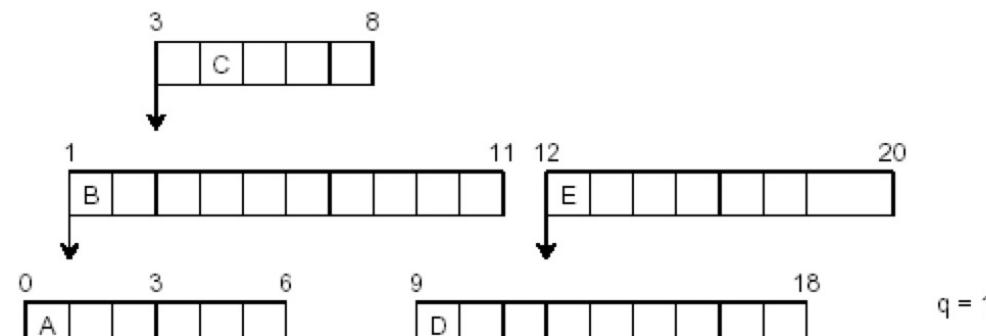
Планировка с приоритетами



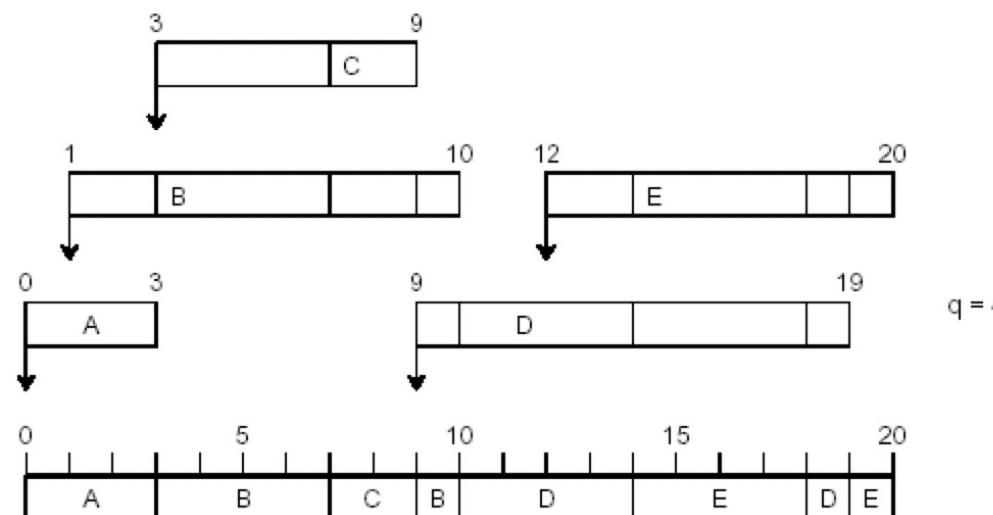
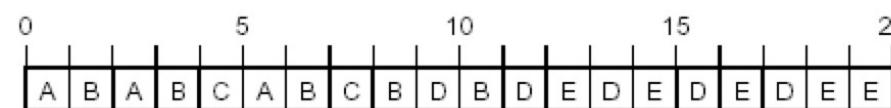
FCFS (First Come First Serving)



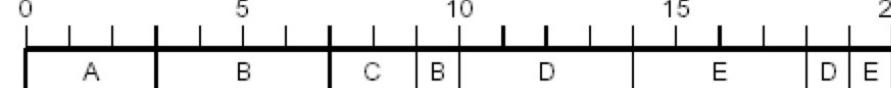
RR (Round Robin)



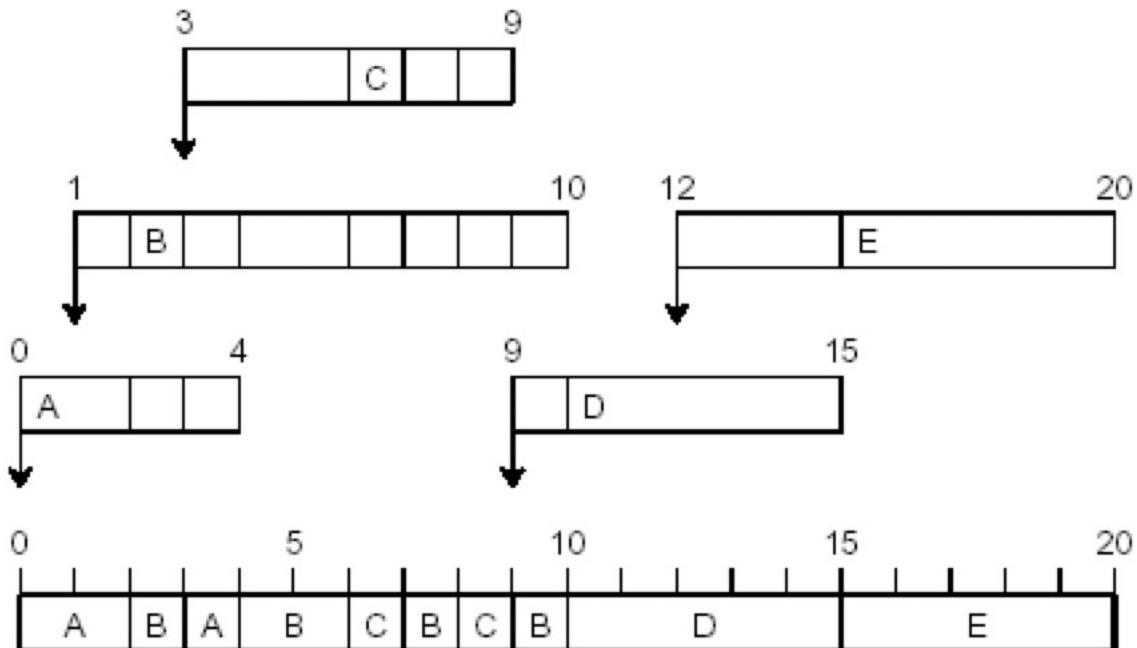
$q = 1$



$q = 4$



SRR (Selfish Round Robin)



0 1 2 3 d

0 2 3 4 5 6 7 8 9 d

0 2 4 6 7 8 d

0 2 3 4 5 6 d

0 2 4 6 7 8 9 10 d

A

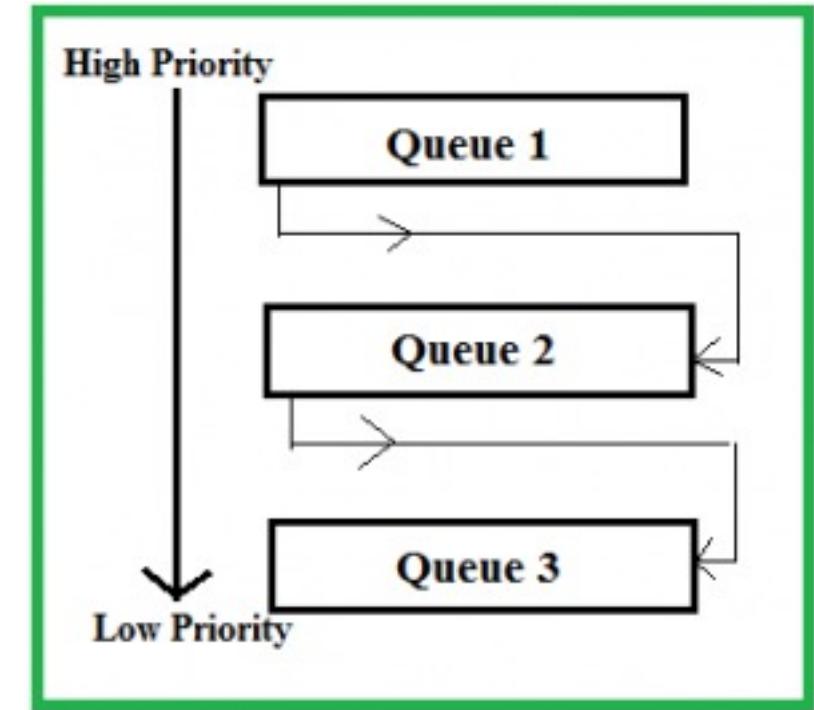
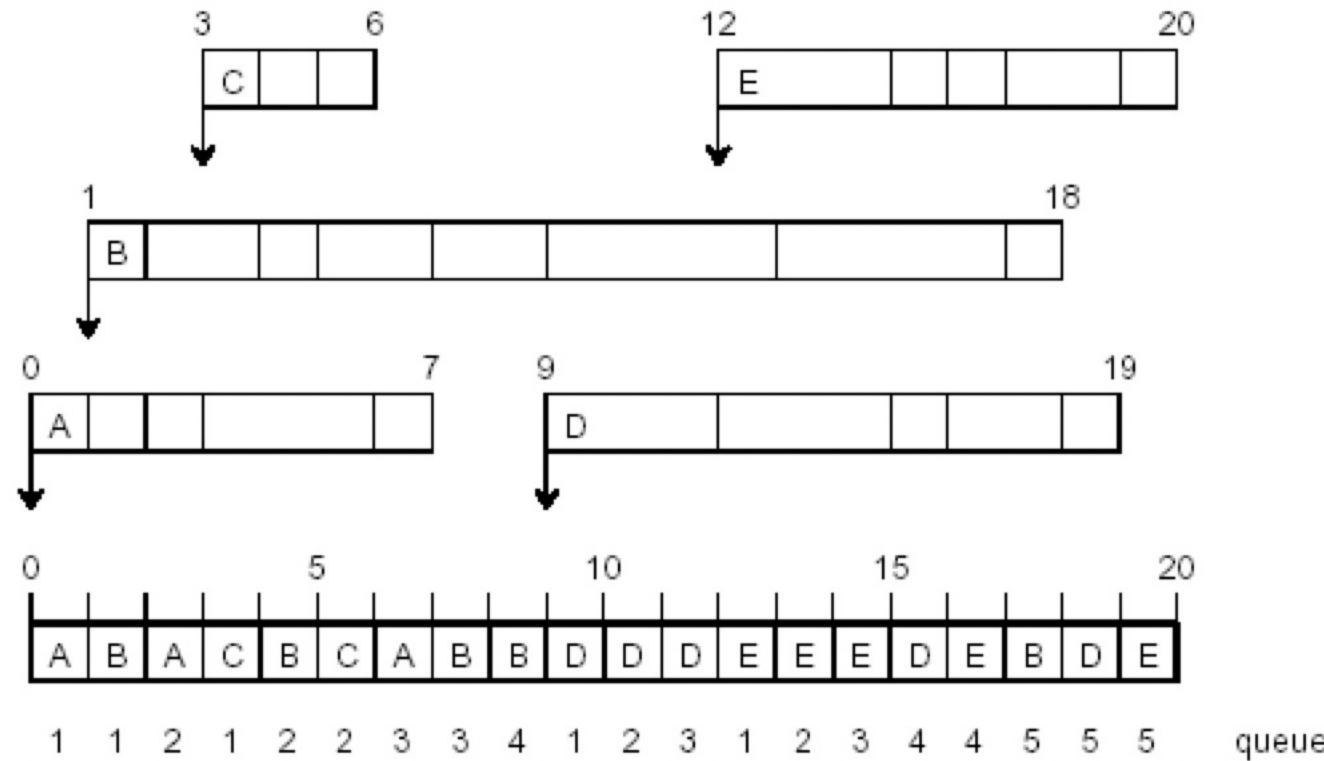
B

C

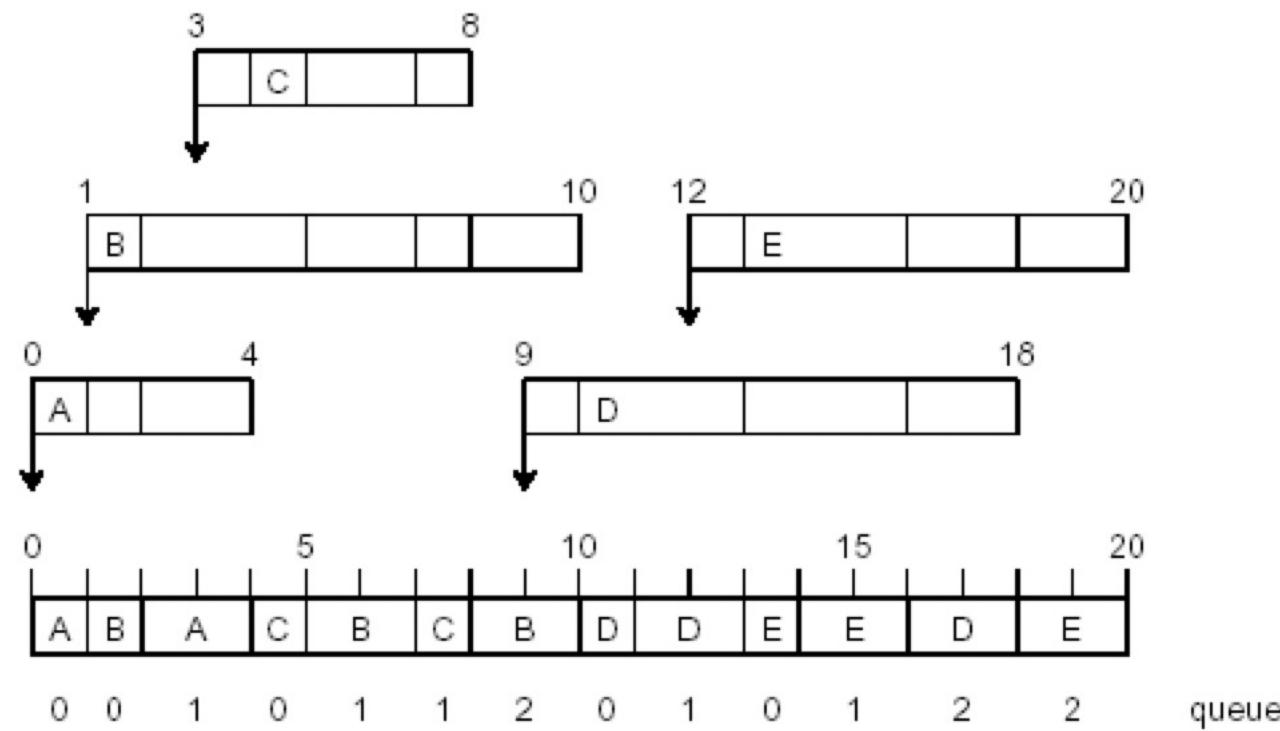
D

E

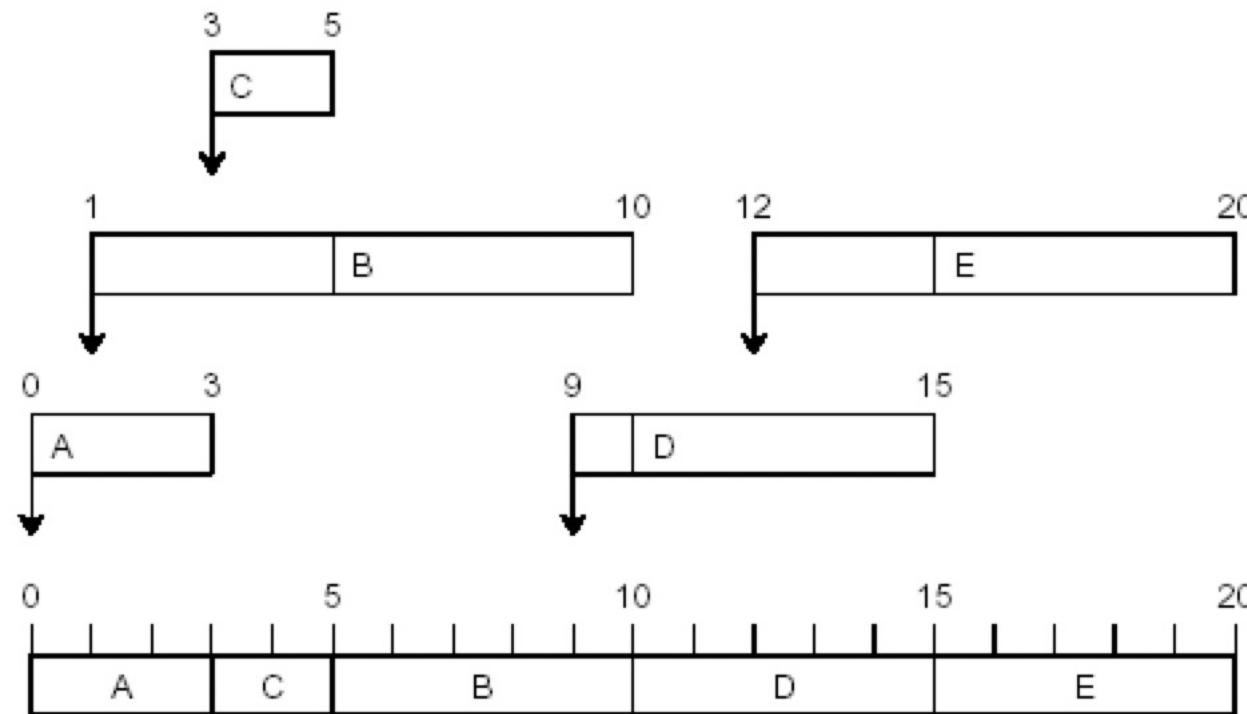
Multilevel Feedback Queue Scheduling (MLFQ)



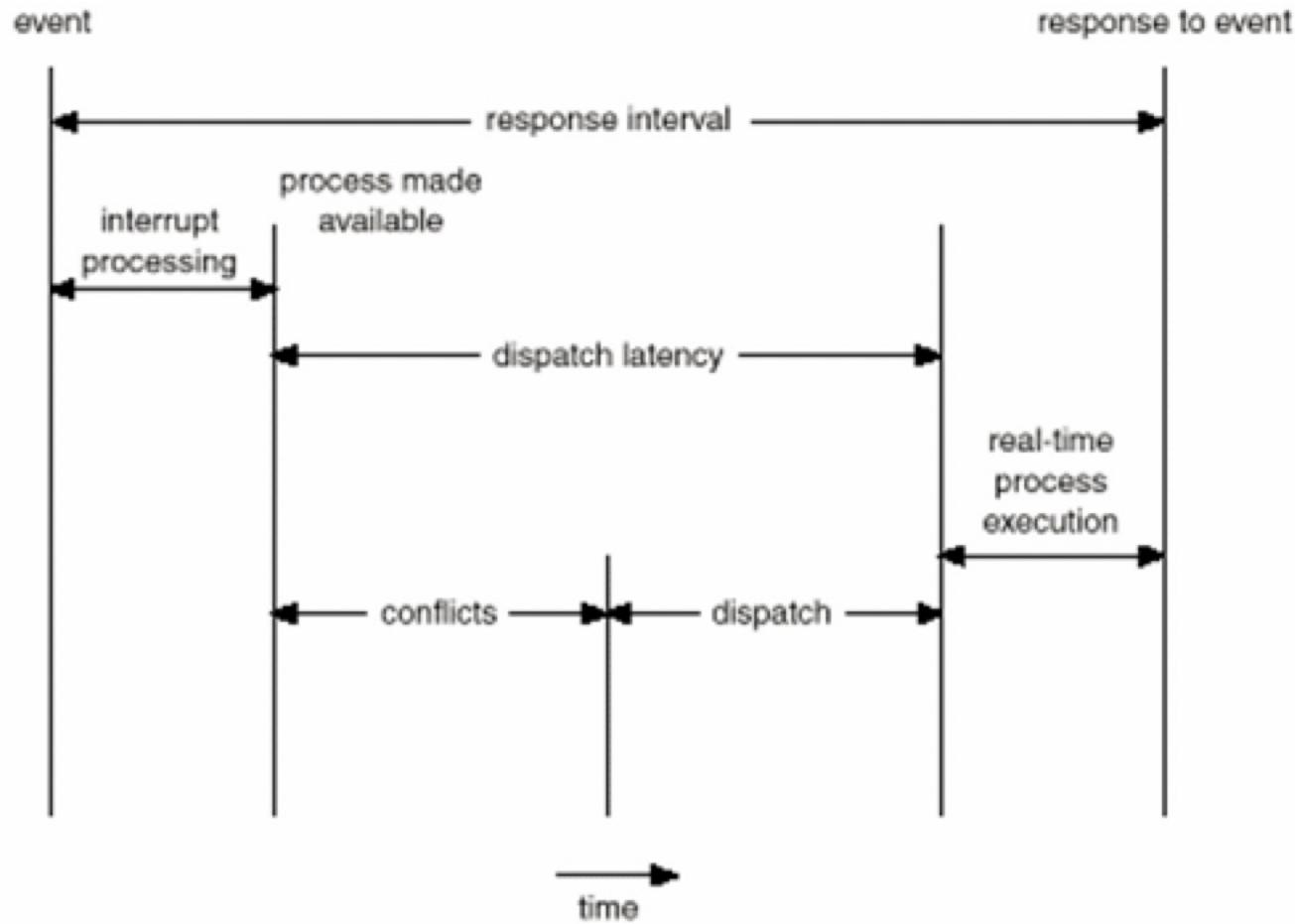
$$\text{expFB } t = b^c \ f = \frac{1}{b^c}$$



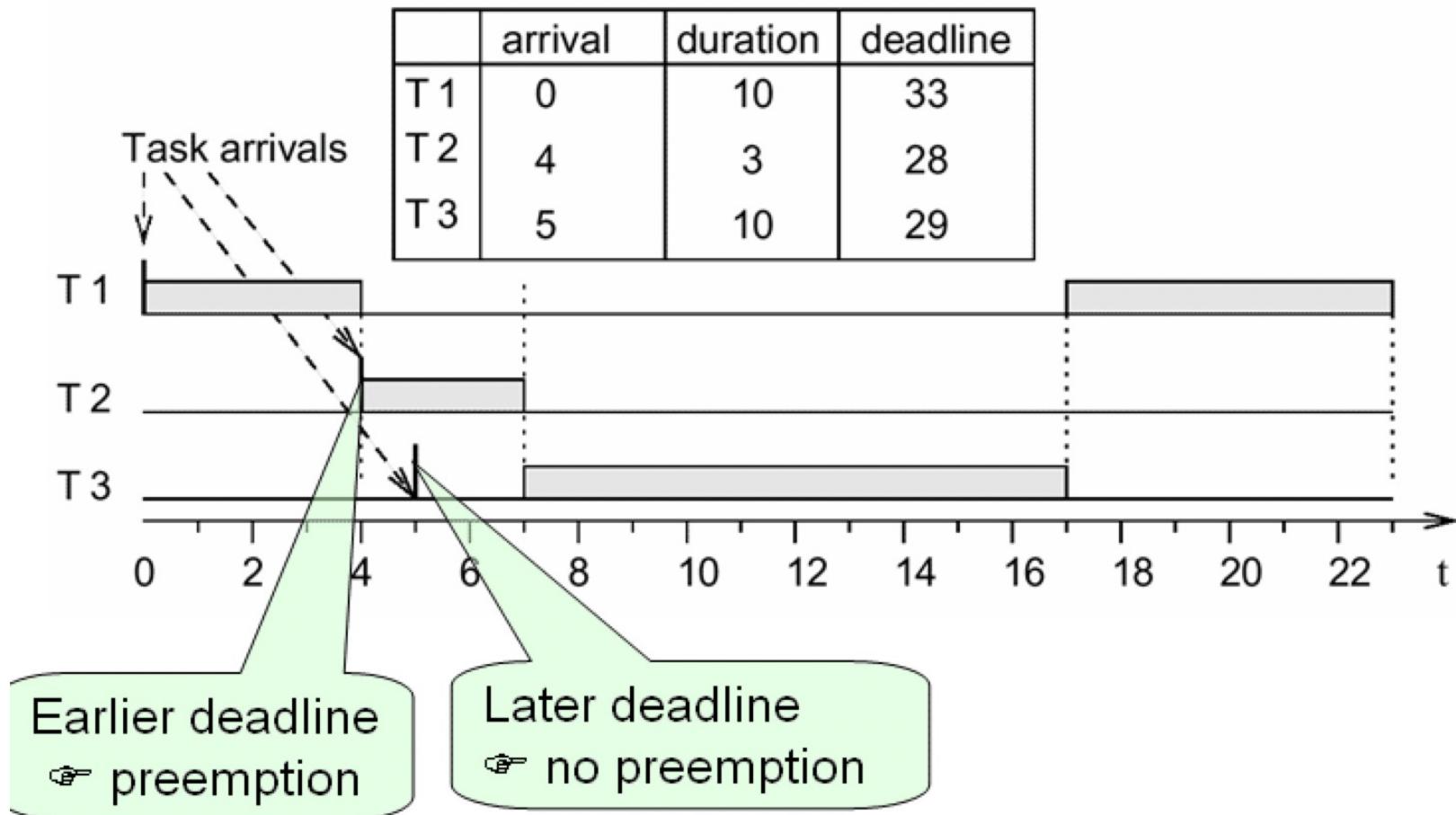
SPN (Shortest Process Next):



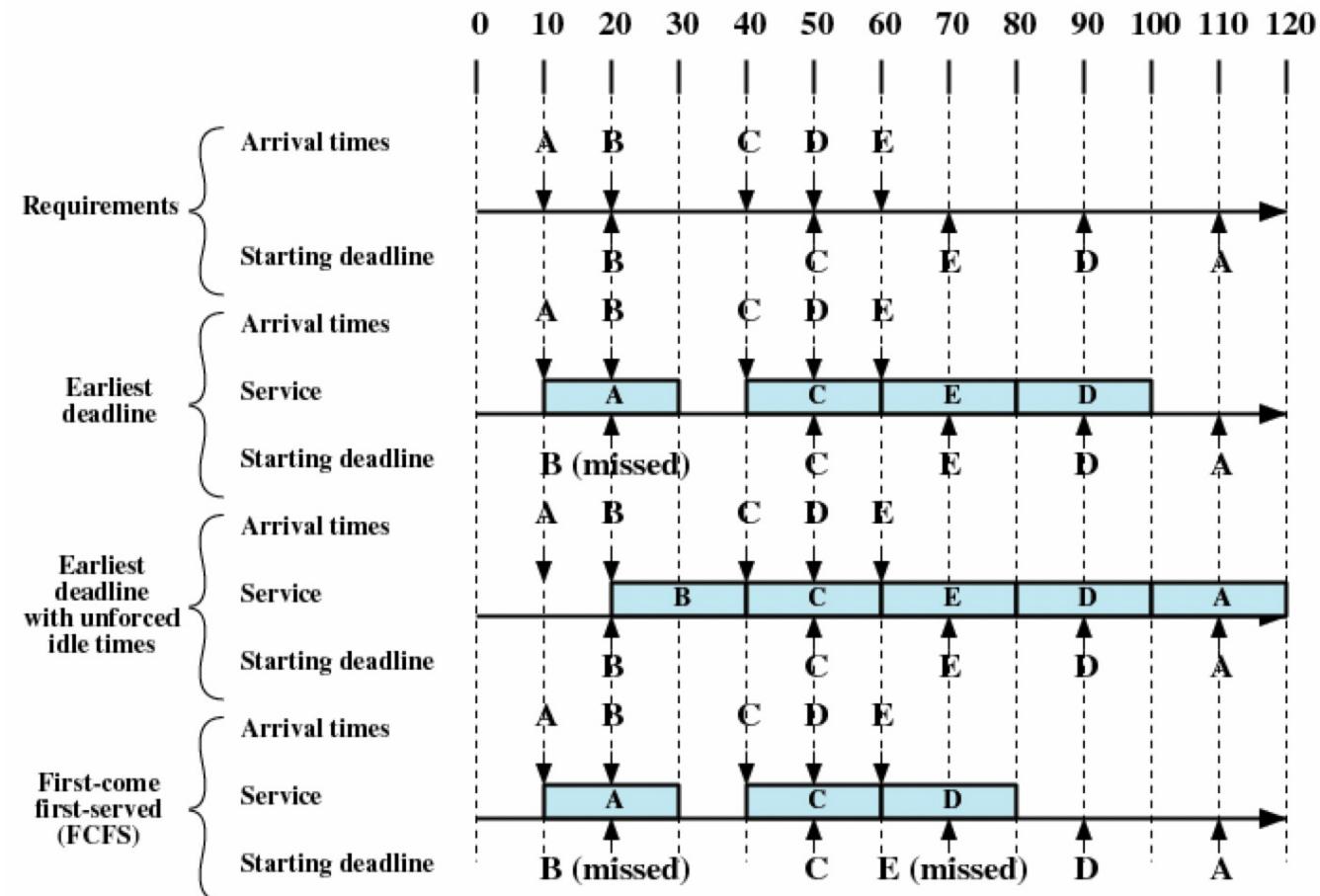
Real time Scheduling



EDF (Earliest Deadline first)



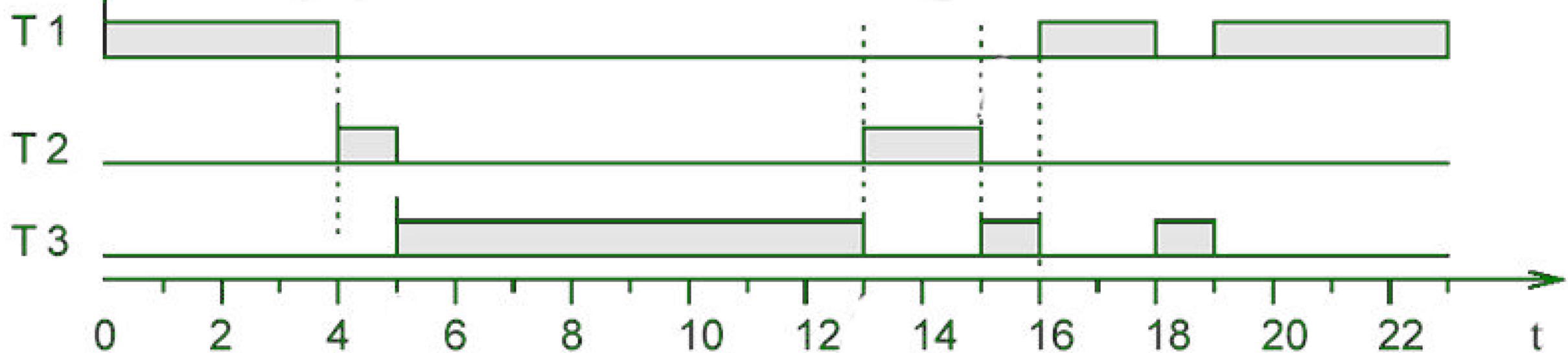
EDF (Earliest Deadline first)



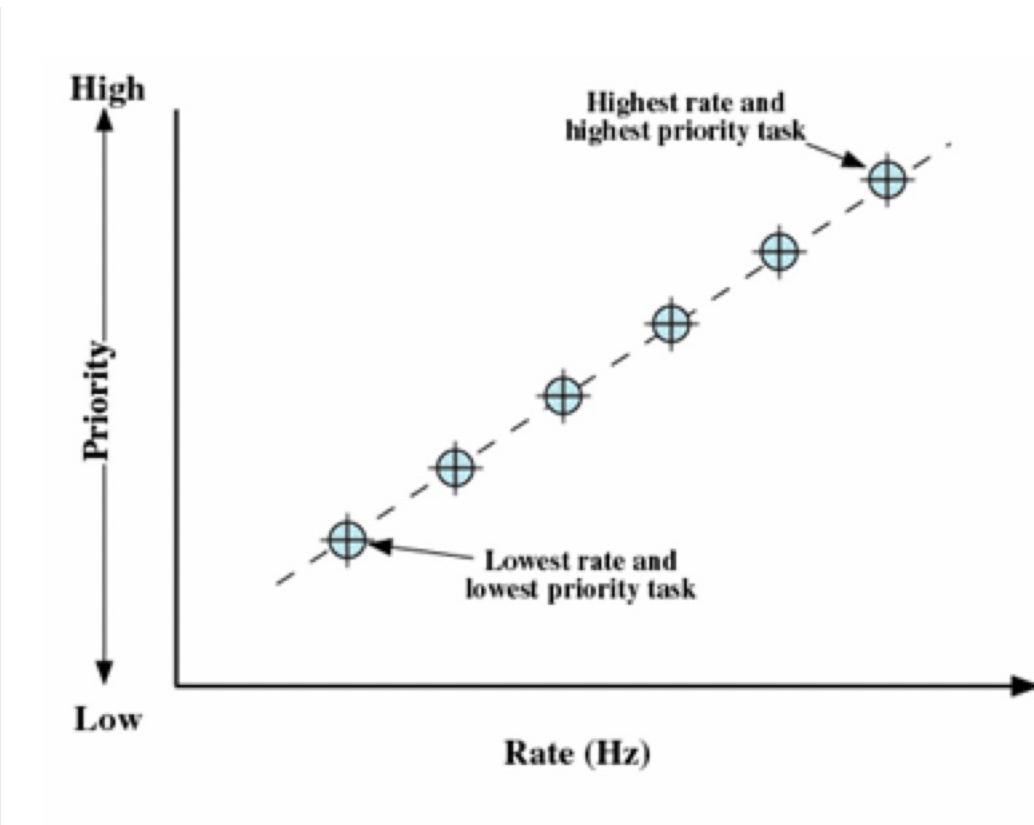
LL/LST (Least Laxity, Least Slack Time First)

	arrival	duration	deadline
T 1	0	10	33
T 2	4	3	28
T 3	5	10	29

$$\text{slack_time} = (D - t - e')$$



RMS (Rate Monotonic Scheduling)



$$\sum_{k=1}^n \frac{C_i}{T_i} \leq U = n(2^{1/n} - 1)$$