# COMP4109 Midterm 1 General Notes

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# 1 Types of Cryptography

- symmetric key
  - ▶ shift ciphers
  - ▶ block ciphers
  - ▶ stream ciphers
- asymmetric key (public-private key)
- hashing
- protocols

## 2 Security Notions Models

- three components of a model
  - 1. attack model
    - ► ciphertext only attack (P) (COA)
      - attacker attempts to decrypt ciphertext to plaintext
    - ▶ known plaintext attack (P) (KPA)
      - attacker knows one or more plaintext-ciphertext pairs
    - ► chosen plaintext attack (A) (CPA)
      - attacker chooses a plaintext and encrypts it to receive ciphertext
    - ► chosen ciphertext attack (A) (CCA)
      - attacker chooses a ciphertext and decrypts it to receive plaintext
  - 2. security goal
    - ▶ (IND) indistinguishability
      - ciphertext should be indistinguishable from random string
    - ► (NM) non-malleability
      - cannot modify ciphertext so it decrypts to another plaintext that makes sense
  - 3. level of security
    - ▶ information theoretic
      - attacker has unlimited resources at their disposal
    - ► complexity theoretic
      - $\circ$  attacker has resources bounded O(p) where p is the security parameter
    - ► computational (realistic)
      - $\circ~$  attacker has the resources of n computers
- two components of a notion
  - ▶ goal + attack model
  - ▶ e.g. IND-COA or NM-COA or IND-KPA, etc.

# 3 Unicity Distance

- expected minimum length of ciphertext needed to uniquely compute a secret key
- $\bullet \quad \frac{\log_2|K|}{R_L\log_2|P|}$ 
  - $\blacktriangleright$  where  $R_L$  is redundancy of the language
  - ►  $R_{\text{English}}$  is about 0.75

# 4 Shift Ciphers

## 4.1 Caesar Cipher

- choose a key from  $\mathbb{Z}_{|P|}$
- $c_i = p_i + k \mod |P|$

#### 4.1.1 Strengths

• none really, this sucks

#### 4.1.2 Weaknesses

- easy to brute force
- weak to frequency analysis

## 4.2 Affine Cipher

- choose any a and  $b \mod 26$ 
  - ightharpoonup except  $a \gcd(a,26)$  must be 1
- k = (a, b) where
  - $E_k(m) = (am + b) \bmod 26$
  - $D_k(c) = a^{-1}(c b) \mod 26$

### 4.2.1 Strengths

- better than caesar cipher
- two unknowns

### 4.2.2 Weaknesses

- use frequency analysis to solve for a and b
- not much better than Caesar really

## 4.3 Substitution Cipher

- permute P to get A
- sub  $P_i$  fo  $A_i$

#### 4.3.1 Strengths

• no strengths, don't use this

#### 4.3.2 Weaknesses

- weak to CPA
- · weak to KPA
- · weak to COA
  - ► frequency analysis
  - ▶ exhaustive search won't work though

## 4.4 Vigenère Cipher

- choose some  $k_l$  as a plaintext string of length l
- encrypt  $c_i = p_i + k_{i \mod l} \mod |P|$

### 4.4.1 Strengths

- much better than what we've seen so far
- if the length of the key is equal to the length of the message, very strong

#### 4.4.2 Weaknesses

- can find candidate key lengths by factoring
- weak to frequency analysis
- multiple encryptions with same key opens up attacks

#### 4.5 One-Time Pad

- like Vigenère except:
  - ▶ change key each time
  - ▶ perfect security if key length is equal to message length

### 4.5.1 Strengths

- perfect security for key length = message length
  - ▶ semantically secure in information theoretic security against COA

#### 4.5.2 Weaknesses

- key can only be used one time
- key length the same as message length is kind of silly
  - ▶ why not just send the message over the secure channel in the first place
  - ▶ very long keys are impractical
- each key needs to be truly random
- · has malleability
  - ▶ no authentication, only confidentiality
- 5 Block Ciphers
- 6 Stream Ciphers
- 7 Hashing Functions
- 8 MACs