

AMATH 515: Homework 5
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1. Consider the Euclidean projection operator

$$P_{\Omega}(x) = \arg \min_{z \in \Omega} \|z - x\|^2,$$

for a closed convex set $\Omega \subset \mathbb{R}^n$. Show that P_{Ω} is non-expansive, i.e.,

$$\|P_{\Omega}(x) - P_{\Omega}(x')\| \leq \|x - x'\|, \quad \forall x, x' \in \mathbb{R}^n.$$

We start by defining the projection operator. The Euclidean projection $P_{\Omega}(x)$ of a point $x \in \mathbb{R}^n$ onto the closed convex set Ω is defined as:

$$P_{\Omega}(x) = \arg \min_{z \in \Omega} \|z - x\|^2$$

Since Ω is closed and convex, the projection is well defined and unique for every x . Let $p = P_{\Omega}(x)$ and $p' = P_{\Omega}(x')$ meaning:

$$p = \arg \min_{z \in \Omega} \|z - x\|^2 \text{ and } p' = \arg \min_{z \in \Omega} \|z - x'\|^2$$

A key property of projections onto convex sets is that the projection satisfies the variational inequality:

$$\langle p - x, z - p \rangle \geq 0, \quad \forall z \in \Omega$$

We can apply the same property to p' :

$$\langle p' - x', z - p' \rangle \geq 0, \quad \forall z \in \Omega$$

From $\langle p - x, z - p \rangle \geq 0$, we can apply $z = p'$ to get:

$$\langle p - x, p' - p \rangle \geq 0$$

Similarly, from $\langle p' - x', z - p' \rangle \geq 0$, we can apply $z = p$ to get:

$$\langle p' - x', p - p' \rangle \geq 0$$

Combining $\langle p - x, p' - p \rangle \geq 0$ and $\langle p' - x', p - p' \rangle \geq 0$, we get:

$$\langle p - x, p' - p \rangle + \langle p' - x', p - p' \rangle \geq 0$$

$$\langle p - x, p' - p \rangle - \langle p' - x', p' - p \rangle \geq 0$$

$$\langle (p - x) - (p' - x'), p' - p \rangle \geq 0$$

$$\langle p - p' - x + x', p' - p \rangle \geq 0$$

Using the linearity and symmetry of the inner product, we rewrite:

$$\langle p - p', p' - p \rangle - \langle x - x', p' - p \rangle \geq 0$$

$$\langle p - p', -(p - p') \rangle - \langle x - x', p' - p \rangle \geq 0$$

$$-\|p - p'\|^2 - \langle x - x', p' - p \rangle \geq 0$$

$$-\|p - p'\|^2 \geq \langle x - x', p' - p \rangle$$

$$\begin{aligned}\|p - p'\|^2 &\leq -\langle x - x', p' - p \rangle \\ \|p - p'\|^2 &\leq -\langle x - x', p - p' \rangle\end{aligned}$$

Now, to complete the proof, we apply Cauchy-Schwarz inequality:

$$\langle x - x', p - p' \rangle \leq \|x - x'\| \|p - p'\|$$

Since we established $\|p - p'\|^2 \leq \langle x - x', p - p' \rangle$, using Cauchy-Schwarz gives us:

$$\|p - p'\|^2 \leq \|x - x'\| \|p - p'\|$$

If $p = p'$, then the inequality trivially holds. Thus, we proceed with the non trivial case dividing both sides by $\|p - p'\|$:

$$\|p - p'\| \leq \|x - x'\|$$

Thus, we have that $\|p - p'\| \leq \|x - x'\|$ for both $p = p'$ and $p \neq p'$.

Thus, we conclude:

$$\|P_\Omega(x) - P_\Omega(x')\| \leq \|x - x'\|, \quad \forall x, x' \in \mathbb{R}^n$$

This proves that the projection operator is indeed nonexpansive.

2. Let $\Omega = \{x \in \mathbb{R}^n \mid \|x\| \leq 1\}$ be the unit ball. Show that

$$P_\Omega(x) = \begin{cases} x & \text{if } \|x\| \leq 1 \\ \frac{x}{\|x\|} & \text{if } \|x\| > 1 \end{cases}$$

The projection $P_\Omega(x)$ of a point $x \in \mathbb{R}^n$ onto Ω is defined as the unique point in Ω that minimizes the Euclidean distance to x : $P_\Omega(x) = \arg \min_{y \in \Omega} \|y - x\|$. Since Ω is closed and convex, projections onto Ω are well defined and unique.

Case 1: $\|x\| \leq 1$:

If x is already in Ω , then clearly x is its own projection since the function $\|y - x\|$ attains its minimum at $y = x$ as when $y = x$, $\|y - x\| = 0$. Hence, $P_\Omega(x) = x$ when $\|x\| \leq 1$.

Case 2: $\|x\| > 1$:

If $\|x\| > 1$, then $x \notin \Omega$. The closest point in Ω to x must be on the boundary $\partial\Omega = \{y \in \mathbb{R}^n \mid \|y\| = 1\}$. We now solve the constrained optimization problem: $\min_{y \in \Omega} \|y - x\|^2$ subject to $\|y\| \leq 1$. We define the

Lagrangian: $L(y, \lambda) = \|y - x\|^2 + \lambda(\|y\|^2 - 1)$. We find $\nabla_y L = 2(y - x) + 2\lambda y$. We set $\nabla_y L$ to 0 in order to find the critical points of the Lagrangian, which correspond to the optimal solutions that

satisfy the constraint $\|y\| = 1$ for our projection. From setting $2(y - x) + 2\lambda y$ to 0, we get:

$$\begin{aligned}
2(y - x) + 2\lambda y &= 0 && \text{(We set this to 0.)} \\
y - x + \lambda y &= 0 \\
y + \lambda y &= x \\
y(1 + \lambda) &= x \\
y &= \frac{x}{1 + \lambda} && \text{(Not needed now, but will be useful later)} \\
\|y\|(1 + \lambda) &= \|x\| && \text{(Taking norms on both sides as } \|y\| = 1) \\
1 + \lambda &= \|x\| && \text{(As } \|y\| = 1) \\
\lambda &= \|x\| - 1 && \text{(Solving for } \lambda) \\
y &= \frac{x}{1 + \lambda} && \text{(Stated before)} \\
y &= \frac{x}{1 + \|x\| - 1} && \text{(Substituting } \lambda = \|x\| - 1) \\
y &= \frac{x}{\|x\|}
\end{aligned}$$

Because $y = \frac{x}{\|x\|}$, the Euclidean projection of x onto Ω for $\|x\| > 1$ is $P_\Omega(x) = \frac{x}{\|x\|}$.

From our cases of $\|x\| \leq 1$ and $\|x\| > 1$, and we found using them, we have shown that

$$P_\Omega(x) = \begin{cases} x & \text{if } \|x\| \leq 1 \\ \frac{x}{\|x\|} & \text{if } \|x\| > 1 \end{cases}$$

3. Given scalars $L \geq m > 0$ show that $\max\{|1 - \alpha m|, |1 - \alpha L|\}$ is minimized when $\alpha = \frac{2}{L+m}$.

We will consider two cases: $L = m$ and $L > m$.

Case 1: $L = m$

When $L = m$, $\alpha = \frac{2}{2L}$, or more simply, $\alpha = \frac{1}{L}$. Furthermore, $\max\{|1 - \alpha m|, |1 - \alpha L|\} = \max\{|1 - \alpha L|, |1 - \alpha L|\}$. This means that $\max\{|1 - \alpha m|, |1 - \alpha L|\} = |1 - \alpha L|$ because the max of two identical values is simply that value itself. Furthermore, $|1 - \alpha L| = |1 - \frac{L}{L}| = |1 - 1| = 0$. Thus, in this trivial case, $\max\{|1 - \alpha m|, |1 - \alpha L|\}$ simply equals 0. We conclude that the expression $\max\{|1 - \alpha m|, |1 - \alpha L|\}$ is minimized at 0 in this trivial case.

Case 2: $L > m$

To achieve minimization, we should make both terms equal: $|1 - \alpha m| = |1 - \alpha L|$. Since in this case, $L > m > 0$, we have that $\alpha > 0$. We analyze when both absolute values are equal while still considering $L > m$ as this means $L \neq m$. This means we consider $1 - \alpha m = -(1 - \alpha L)$ instead of $1 - \alpha m = 1 - \alpha L$ as once again, $L \neq m$. So, we proceed with $1 - \alpha m = -(1 - \alpha L)$.

$$1 - \alpha m = -(1 - \alpha L)$$

$$1 - \alpha m = -1 + \alpha L$$

$$2 = \alpha m + \alpha L$$

$$\frac{2}{\alpha} = m + L$$

$$\alpha = \frac{2}{m + L}$$

We have that in both cases $\max\{|1 - \alpha m|, |1 - \alpha L|\}$ is minimized when $\alpha = \frac{2}{L+m}$. Thus, with our scalars $L \geq m > 0$, $\max\{|1 - \alpha m|, |1 - \alpha L|\}$ is minimized when $\alpha = \frac{2}{L+m}$.

4. For a constant vector $c \in \mathbb{R}^n$ and variable $x \in \mathbb{R}^n$ find the minimizer of $c^T x$ over Ω , where Ω is each of the following sets:

- (a) The unit ball $\{x \mid \|x\| \leq 1\}$

Since $c^T x$ represents the dot product of c and x , the function $c^T x$ is minimized when x is in the opposite direction of c . This is because the dot product is maximized when x is aligned with c and minimized when x is in the opposite direction. The constraint requires that x lies within or on the boundary of the unit ball meaning $\|x\| \leq 1$. To achieve the smallest possible value of $c^T x$, we should take x to be in the direction of $-c$ and have the largest possible norm, where $\|x^*\| = 1$ because $\left\| -\frac{c}{\|c\|} \right\| = \frac{\|c\|}{\|c\|}$, and $\frac{\|c\|}{\|c\|}$ is simply 1. Thus, for the unit ball, the optimal choice is $x^* = -\frac{c}{\|c\|}$, which satisfies $\|x^*\| = 1$. With our minimizer $x^* = -\frac{c}{\|c\|}$, we have that our minimum value is calculated from $c^T x^*$, which is $c^T \left(-\frac{c}{\|c\|} \right)$, which simplifies to $-\|c\|$. Thus, the minimum value is $-\|c\|$ with minimizer $x^* = -\frac{c}{\|c\|}$.

- (b) A box $\{x \mid 0 \leq x_i \leq 1, \quad i = 1, \dots, n\}$

The constraint means that each component x_i is bounded between 0 and 1. The function $f(x) = c^T x$ is linear, so it is minimized at a vertex of the box since a linear function attains its extreme values at the boundary of a convex set because a linear function's minimum will occur at an extreme point of the feasible region as a linear function attains its extreme values at the boundary of a convex set since it has no curvature and always increases and decreases in a fixed direction

and in a monotone manner. The function $c^T x$ is the sum: $f(x) = \sum_{i=1}^n c_i x_i$. Each x_i should be set to 0 if $c_i > 0$ to minimize the contribution to $c^T x$, and each x_i should be set to 1 if $c_i < 0$ to minimize the contribution to $c^T x$. Thus, the optimal solution is: $x_i^* = \begin{cases} 0 & \text{if } c_i > 0 \\ 1 & \text{if } c_i < 0 \end{cases}$. Thus,

our minimizer is $x_i^* = \begin{cases} 0 & \text{if } c_i > 0 \\ 1 & \text{if } c_i < 0 \end{cases}$ with the minimum value $-\sum_{i=1}^n c_i^{\ominus}$, where c_i^{\ominus} is defined as

$c_i^{\ominus} = \begin{cases} -c_i & \text{when } c_i < 0 \\ 0 & \text{when } c_i > 0 \end{cases}$. Note that $c_i = 0$ the term $c_i x_i = 0$ regardless of the choice of x_i , so any value of $x_i \in [0, 1]$ is optimal.

Hint: try to visualize the problems in 2D before attempting a proof.

5. Consider the projected steepest descent update formula

$$x^{(k+1)} = P_{\Omega}(x^{(k)} - \alpha^{(k)} \nabla f(x^{(k)}))$$

with step size $\alpha^{(k)} > 0$, a closed convex set $\Omega \subset \mathbb{R}^n$, and a continuously differentiable f . Show that

$$x^{(k+1)} = \arg \min_{x \in \Omega} f(x^{(k)}) + \nabla f(x^{(k)})^T (x - x^{(k)}) + \frac{1}{2\alpha^{(k)}} \|x - x^{(k)}\|^2$$

The projection operator P_{Ω} projects a given point onto the closed convex set Ω by solving the minimization problem: $P_{\Omega}(y) = \arg \min_{x \in \Omega} \frac{1}{2} \|x - y\|^2$. For our given update rule, we set: $y = x^{(k)} -$

$\alpha^{(k)} \nabla f(x^{(k)})$. Thus, we have $x^{(k+1)} = \arg \min_{x \in \Omega} \frac{1}{2} \|x - (x^{(k)} - \alpha^{(k)} \nabla f(x^{(k)}))\|^2$ from $x^{(k+1)} = P_{\Omega}(x^{(k)} - \alpha^{(k)} \nabla f(x^{(k)}))$.

$$\begin{aligned} x^{(k+1)} &= P_{\Omega}(x^{(k)} - \alpha^{(k)} \nabla f(x^{(k)})) \\ x^{(k+1)} &= \arg \min_{x \in \Omega} \frac{1}{2} \|x - (x^{(k)} - \alpha^{(k)} \nabla f(x^{(k)}))\|^2 \\ x^{(k+1)} &= \arg \min_{x \in \Omega} \frac{1}{2} \|x - x^{(k)} + \alpha^{(k)} \nabla f(x^{(k)})\|^2 \\ x^{(k+1)} &= \arg \min_{x \in \Omega} \frac{1}{2} \left(\|x - x^{(k)}\|^2 + 2\alpha^{(k)} \nabla f(x^{(k)})^T (x - x^{(k)}) + (\alpha^{(k)})^2 \|\nabla f(x^{(k)})\|^2 \right) \\ x^{(k+1)} &= \arg \min_{x \in \Omega} \frac{1}{2} \left(\|x - x^{(k)}\|^2 \right) + \left(\alpha^{(k)} \nabla f(x^{(k)})^T (x - x^{(k)}) \right) + \frac{1}{2} \left((\alpha^{(k)})^2 \|\nabla f(x^{(k)})\|^2 \right) \end{aligned}$$

Notice how the term $\frac{1}{2} (\alpha^{(k)})^2 \|\nabla f(x^{(k)})\|^2$ does not influence which x minimizes the expression as when minimizing over x , $f(x^{(k)})$ does not influence the choice of x and $f(x^{(k)})$ only shifts the objective function by a constant amount because $f(x^{(k)})$ is independent of x and does not contribute to the gradient or curvature of the function being minimized. Since optimization depends only on the terms that vary with x , constants can be ignored without affecting the minimization. Thus, $f(x^{(k)})$ and resultantly, $\frac{1}{2} ((\alpha^{(k)})^2 \|\nabla f(x^{(k)})\|^2)$ do not alter the minimization of the argument of our given objective function that we are set to minimize. Thus, we proceed by minimizing $\frac{1}{2} (\|x - x^{(k)}\|^2) + (\alpha^{(k)} \nabla f(x^{(k)})^T (x - x^{(k)}))$ and proceed with $x^{(k+1)} = \arg \min_{x \in \Omega} \frac{1}{2} (\|x - x^{(k)}\|^2) + (\alpha^{(k)} \nabla f(x^{(k)})^T (x - x^{(k)}))$.

Now, we proceed with the minimization problem:

$$x^{(k+1)} = \arg \min_{x \in \Omega} \frac{1}{2} (\|x - x^{(k)}\|^2) + (\alpha^{(k)} \nabla f(x^{(k)})^T (x - x^{(k)}))$$

Rearranging, we get:

$$x^{(k+1)} = \arg \min_{x \in \Omega} \alpha^{(k)} \nabla f(x^{(k)})^T (x - x^{(k)}) + \frac{1}{2} \|x - x^{(k)}\|^2$$

Now, we divide by our step size, $\alpha^{(k)}$ to get:

$$x^{(k+1)} = \arg \min_{x \in \Omega} f(x^{(k)}) + \nabla f(x^{(k)})^T (x - x^{(k)}) + \frac{1}{2\alpha^{(k)}} \|x - x^{(k)}\|^2$$

Thus, we have shown from our projected steepest descent update formula: $x^{(k+1)} = P_{\Omega}(x^{(k)} - \alpha^{(k)} \nabla f(x^{(k)}))$, we can get $x^{(k+1)} = \arg \min_{x \in \Omega} f(x^{(k)}) + \nabla f(x^{(k)})^T (x - x^{(k)}) + \frac{1}{2\alpha^{(k)}} \|x - x^{(k)}\|^2$.