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4	4.1 Dynamic Convex Hull Trick	10 10 10	3 4 5 6	<pre>// C(pi) the number of cycles in the permutation pi // G the number of permutations</pre>	
5	5.1 Convex Hull 5.2 Geometry Template 5.3 Half Plane Intersection 5.4 Segments Intersection	13 13 14	2 3	Catlan Numbers const int MOD = const int MAX = int catalan[MAX]; void init() {	
6	6.1 2 SAD 6.2 Ariculation Point 6.3 Bridges Tree and Diameter 6.4 Dinic With Scalling 6.5 Gomory Hu 6.6 HoperaftKarp BPM 6.7 Hungarian	16 16	5 6 7 8 9 0 1 2 3	<pre>catalan[0] = catalan[1] = 1; for (int i=2; i<=n; i++) { catalan[i] = 0; for (int j=0; j < i; j++) { catalan[i] += (catalan[j] * catalan[i-j-1]) % MOD; if (catalan[i] >= MOD) { catalan[i] -= MOD; } } }</pre>	

6.9

6.10

```
17 // 1- Number of correct bracket sequence consisting of n opening and n closing
    // 2- The number of rooted full binary trees with n+1 leaves (vertices are not
        numbered).
         A rooted binary tree is full if every vertex has either two children or no
         children.
   // 3- The number of ways to completely parenthesize n+1 factors.
   // 4- The number of triangulations of a convex polygon with n+2 sides
          (i.e. the number of partitions of polygon into disjoint triangles by using
          the diagonals).
   // 5- The number of ways to connect the 2n points on a circle to form n disjoint
   // 6- The number of non-isomorphic full binary trees with n internal nodes (i.e.
         nodes having at least one son).
   // 7- The number of monotonic lattice paths from point (0,0) to point (n,n) in a
         square lattice of size nxn.
          which do not pass above the main diagonal (i.e. connecting (0,0) to (n,n))
   // 8- Number of permutations of length n that can be stack sorted
         (i.e. it can be shown that the rearrangement is stack sorted if and only
        there is no such index i<j<k, such that ak<ai<aj ).
  // 9- The number of non-crossing partitions of a set of n elements.
  // 10- The number of ways to cover the ladder 1..n using n rectangles
   // (The ladder consists of n columns, where ith column has a height i).
```

2 Algebra

2.1 Primitive Roots

```
int powmod (int a, int b, int p) {
         int res = 1;
3
         while (b)
             if (b & 1)
                 res = int (res * 111 * a % p), --b;
                a = int (a * 111 * a % p), b >>= 1;
8
         return res:
9
10
11
    int generator (int p) {
12
         vector<int> fact;
13
         int phi = p - 1, n = phi;
14
         for (int i = 2; i * i <= n; ++i)
15
             if (n \% i == 0) {
16
                 fact.push_back (i);
17
                 while (n \% i == 0)
18
                     n /= i;
19
20
        if (n > 1)
21
             fact.push_back (n);
22
23
         for (int res = 2; res <= p; ++res) {</pre>
24
             bool ok = true;
25
             for (size_t i = 0; i < fact.size() && ok; ++i)</pre>
26
                ok &= powmod (res, phi / fact[i], p) != 1;
27
             if (ok) return res;
28
29
         return -1;
30
```

2.2 Discrete Logarithm

```
int an = 1;
        for (int i = 0; i < n; ++i)
            an = (an * 111 * a) % m;
 9
        unordered_map<int, int> vals;
11
        for (int q = 0, cur = b; q \le n; ++q) {
            vals[cur] = q;
13
             cur = (cur * 111 * a) % m;
14
        for (int p = 1, cur = 1; p <= n; ++p) {</pre>
16
            cur = (cur * 111 * an) % m;
             if (vals.count(cur)) {
19
                int ans = n * p - vals[cur];
                return ans;
        return -1;
    //When a and m are not coprime
    // Returns minimum x for which a ^ x % m = b % m.
    int solve(int a, int b, int m) {
        a %= m, b %= m;
30
        int k = 1, add = 0, q;
31
        while ((g = gcd(a, m)) > 1) {
32
            if (b == k)
33
                return add;
34
            if (b % g)
35
                return -1;
            b /= g, m /= g, ++add;
36
             k = (k * 111 * a / g) % m;
37
38
39
40
        int n = sqrt(m) + 1;
41
        int an = 1;
42
        for (int i = 0; i < n; ++i)
43
            an = (an * 111 * a) % m;
44
45
        unordered_map<int, int> vals;
46
        for (int q = 0, cur = b; q \le n; ++q) {
47
            vals[cur] = q;
48
            cur = (cur * 111 * a) % m;
49
50
51
        for (int p = 1, cur = k; p \le n; ++p) {
            cur = (cur * 111 * an) % m;
53
            if (vals.count(cur)) {
                int ans = n * p - vals[cur] + add;
55
                return ans;
56
57
58
        return -1:
59
```

2.3 Iteration over submasks

2.4 Totient function

```
1  void phi_1_to_n(int n) {
2     vector<int> phi(n + 1);
3     phi[0] = 0;
```

2.5 CRT and EEGCD

```
11 extended(11 a, 11 b, 11 &x, 11 &y) {
2
3
         if(b == 0) {
4
             x = 1;
             y = 0;
-5
6
             return a;
7
         11 x0, y0;
9
        11 g = \text{extended}(b, a % b, x0, y0);
10
        x = y0;
11
         y = x0 - a / b * y0;
12
13
         return g ;
14
15
    ll de(ll a, ll b, ll c, ll &x, ll &y) {
16
17
         11 g = extended(abs(a), abs(b), x, y);
18
         if(c % g) return -1;
19
20
         x \star = c / q;
21
         y \star = c / g;
22
23
         if(a < 0)x = -x;
24
         if(b < 0)y = -y;
25
26
    pair<11, 11> CRT(vector<11> r, vector<11> m) {
28
29
         11 r1 = r[0], m1 = m[0];
30
31
         for(int i = 1; i < r.size(); i++) {</pre>
32
33
             11 r2 = r[i], m2 = m[i];
34
             11 x0, v0;
35
             11 q = de(m1, -m2, r2 - r1, x0, y0);
36
\frac{37}{38}
             if(q == -1) return \{-1, -1\};
39
             11 \text{ nr} = x0 * m1 + r1;
40
             11 nm = m1 / q * m2;
41
42
             r1 = (nr % nm + nm) % nm;
43
             m1 = nm;
44
45
         return {r1, m1};
```

2.6 FFT

```
1 #include<iostream>
2 #include <bits/stdc++.h>
3 #define l1 long long
4 #define ld long double
5 #define rep(i, a, b) for(int i = a; i < (b); ++i)
6 #define all(x) begin(x), end(x)
7 #define sz(x) (int)(x).size()</pre>
```

```
#define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
   typedef complex<double> C;
11 typedef vector<double> vd;
12 typedef vector<int> vi;
    typedef pair<int, int> pii;
14
    void fft(vector<C>& a) {
15
        int n = sz(a), L = 31 - \underline{builtin_clz(n)};
16
         static vector<complex<long double>> R(2, 1);
17
         static vector<C> rt(2, 1); // (^ 10% fas te r i f double)
18
        for (static int k = 2; k < n; k *= 2) {
19
            R.resize(n);
20
            rt.resize(n);
21
            auto x = polar(1.0L, acos(-1.0L) / k);
22
            rep(i, k, 2 * k) rt[i] = R[i] = i & 1 ? R[i / 2] * x : R[i / 2];
23
24
25
        rep(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
26
        rep(i, 0, n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
         for (int k = 1; k < n; k *= 2)
27
28
             for (int i = 0; i < n; i += 2 * k) rep(j, 0, k) {
29
                Cz = rt[j + k] * a[i + j + k]; //
30
                a[i + j + k] = a[i + j] - z;
31
                a[i + j] += z;
32
33
34 vd conv(const vd& a, const vd& b) {
35
        if (a.empty() || b.empty()) return {};
36
        vd res(sz(a) + sz(b) - 1);
37
        int L = 32 - __builtin_clz(sz(res)), n = 1 << L;</pre>
        vector<C> in(n), out(n);
39
        copy(all(a), begin(in));
40
        rep(i, 0, sz(b)) in[i].imag(b[i]);
41
         fft(in);
42
        for (C\& x : in) x *= x;
43
        rep(i, 0, n) out[i] = in[-i & (n - 1)] - conj(in[i]);
44
45
        rep(i, 0, sz(res)) res[i] = imag(out[i]) / (4 * n);
46
        return res;
47
48
49
    int main() {
50
51
         //Applications
52
         //1-All possible sums
53
54
         //2-All possible scalar products
55
         // We are given two arrays a[] and b[] of length n.
         //We have to compute the products of a with every cyclic shift of b.
56
57
         //We generate two new arrays of size 2n: We reverse a and append n zeros to
58
         //And we just append b to itself. When we multiply these two arrays as
             polynomials,
59
         //and look at the coefficients c[n-1], c[n], ..., c[2n-2] of the product c,
             we get:
60
        //c[k]=sum\ i+j=k\ a[i]b[j]
61
62
         //3-Two stripes
63
         //We are given two Boolean stripes (cyclic arrays of values 0 and 1) a and b
64
         //We want to find all ways to attach the first stripe to the second one,
65
         //such that at no position we have a 1 of the first stripe next to a 1 of
             the second stripe.
```

2.7 Fibonacci

```
1 2 3 // F(n-1) * F(n+1) - F(n)^2 = (-1)^n 4 5 // F(n+k) = F(k) * F(n+1) + F(k-1) * F(n)
```

2.8 Gauss Determinant

```
const double EPS = 1E-9;
    int n;
    vector < vector<double> > a (n, vector<double> (n));
5
    double det = 1;
6
    for (int i=0; i<n; ++i) {</pre>
        int k = i;
         for (int j=i+1; j<n; ++j)</pre>
Q
             if (abs (a[j][i]) > abs (a[k][i]))
10
11
         if (abs (a[k][i]) < EPS) {</pre>
12
             det = 0;
13
             break:
14
15
         swap (a[i], a[k]);
16
         if (i != k)
17
             det = -det;
18
         det *= a[i][i];
19
         for (int j=i+1; j<n; ++j)</pre>
20
             a[i][j] /= a[i][i];
21
         for (int j=0; j < n; ++j)
22
             if (j != i && abs (a[j][i]) > EPS)
23
                 for (int k=i+1; k<n; ++k)</pre>
24
                     a[j][k] = a[i][k] * a[j][i];
25
26
27 cout << det;
```

2.9 GAUSS SLAE

```
const double EPS = 1e-9;
    const int INF = 2; // it doesn't actually have to be infinity or a big number
    int gauss (vector < vector<double> > a, vector<double> & ans) {
 4
         int n = (int) a.size();
 5
         int m = (int) a[0].size() - 1;
         vector<int> where (m, -1);
 8
 Q
         for (int col = 0, row = 0; col < m && row < n; ++col) {
10
             int sel = row;
11
             for (int i = row; i < n; ++i)
12
                 if (abs (a[i][col]) > abs (a[sel][col]))
13
                    sel = i;
14
             if (abs (a[sel][col]) < EPS)</pre>
15
                 continue;
16
             for (int i = col; i <= m; ++i)</pre>
17
                 swap (a[sel][i], a[row][i]);
18
             where [col] = row;
19
20
             for (int i = 0; i < n; ++i)
21
                 if (i != row) {
22
                     double c = a[i][col] / a[row][col];
\frac{23}{24}
                     for (int j = col; j \le m; ++j)
                         a[i][j] = a[row][j] * c;
25
26
             ++row;
27
28
29
         ans.assign (m, 0);
30
         for (int i = 0; i < m; ++i)
31
             if (where[i] != -1)
                 ans[i] = a[where[i]][m] / a[where[i]][i];
```

```
for (int i = 0; i < n; ++i) {
34
            double sum = 0;
35
            for (int j = 0; j < m; ++j)
36
                sum += ans[j] * a[i][j];
37
            if (abs (sum - a[i][m]) > EPS)
38
                return 0;
39
40
41
        for (int i = 0; i < m; ++i)
42
            if (where[i] == -1)
43
                return INF;
44
        return 1;
45
```

2.10 Matrix Inverse

```
\ //\  Sometimes, the questions are complicated - and the answers are simple. \ //\ 
    #pragma GCC optimize ("03")
    #pragma GCC optimize ("unroll-loops")
    #include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    vector < vector<double> > gauss (vector < vector<double> > a) {
11
        int n = (int) a.size();
        vector<vector<double> > ans(n, vector<double>(n, 0));
12
13
14
        for (int i = 0; i < n; i++)
15
            ans[i][i] = 1;
16
        for (int i = 0; i < n; i++) {
17
            for (int j = i + 1; j < n; j++)
18
                if(a[j][i] > a[i][i]) {
19
                    swap(a[j], a[i]);
20
                    swap(ans[j], ans[i]);
21
22
            double val = a[i][i];
23
            for (int j = 0; j < n; j++) {
24
                a[i][j] /= val;
25
                ans[i][j] /= val;
27
            for (int j = 0; j < n; j++) {
28
                if(j == i)continue;
29
                val = a[j][i];
30
                for (int k = 0; k < n; k++) {
                    a[j][k] = val * a[i][k];
                     ans[j][k] = val * ans[i][k];
33
35
36
        return ans;
37
38
    int main() {
39
40
41
        vector<vector<double> > v(3, vector<double> (3) );
42
        for (int i = 0; i < 3; i++)
43
            for (int j = 0; j < 3; j++)
44
                cin >> v[i][j];
45
        for(auto i : gauss(v)) {
46
47
            for(auto j : i)
                cout << j << " ";
48
49
            cout << "\n";
50
```

2.11 NTT

```
struct NTT {
         int mod ;
         int root :
         int root_1 ;
        int root_pw ;
        NTT(int _mod, int primtive_root, int NTT_Len) {
 9
             mod = mod:
10
             root_pw = NTT_Len;
11
             root = fastpower(primtive_root, (mod - 1) / root_pw);
12
             root_1 = fastpower(root, mod - 2);
13
14
        void fft(vector<int> & a, bool invert) {
15
             int n = a.size();
16
17
             for (int i = 1, j = 0; i < n; i++) {
18
                 int bit = n >> 1;
19
                 for (; j & bit; bit >>= 1)
20
                    j ^= bit;
21
                 j ^= bit;
22
23
                 if (i < j)
\frac{24}{25}
                     swap(a[i], a[j]);
26
27
             for (int len = 2; len <= n; len <<= 1) {</pre>
28
                 int wlen = invert ? root_1 : root;
29
                 for (int i = len; i < root pw; i <<= 1)</pre>
30
                     wlen = (int)(1LL * wlen * wlen % mod);
31
32
33
                 for (int i = 0; i < n; i += len) {</pre>
34
                     int w = 1;
35
                     for (int j = 0; j < len / 2; j++) {
36
                         int u = a[i + j], v = (int)(1LL * a[i + j + len / 2] * w %
                              mod);
37
                          a[i + j] = u + v < mod ? u + v : u + v - mod;
38
                         a[i + j + len / 2] = u - v >= 0 ? u - v : u - v + mod;
39
                          w = (int) (1LL * w * wlen % mod);
40
                     }
41
42
43
44
             if (invert) {
45
                 int n_1 = fastpower(n, mod - 2);
46
                 for (int & x : a)
47
                     x = (int) (1LL * x * n_1 % mod);
48
49
50
         vector<int> multiply(vector<int> &a, vector<int> &b) {
51
             vector<int> fa(a.begin(), a.end()), fb(b.begin(), b.end());
52
             int n = 1:
53
             while(n < a.size() + b.size())</pre>
54
                 n \ll 1;
55
56
             fa.resize(n);
57
             fb.resize(n);
58
59
             fft(fa, 0);
60
             fft(fb, 0);
61
62
             for (int i = 0; i < n; i++)
63
                fa[i] = 1LL * fa[i] * fb[i] % mod;
64
             fft(fa, 1);
65
             return fa;
66
67
    };
```

```
for(int i = x; i <= n; i += i & -i)
for(int j = y; j <= m; j += j & -j)

bit[i][j] += val;

}

int get(int x, int y) {
   int ans = 0;
   for(int i = x; i; i -= i & -i)
   for(int j = y; j; j -= j & -j)
   ans += bit[i][j];
}</pre>
```

void upd(int x, int y, int val) {

int md = (s + e) >> 1; return conv(solve(s,md),solve(md+1)) 3 Data Structures

3.1 2D BIT

11 }

```
///If the mod changes don't forget to calculate the primitive root
    using 11 = long long;
    const 11 mod = (119 << 23) + 1, root = 3; // = 998244353
    // For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 << 21
    // and 483 << 21 (same root). The last two are > 10^9.
    typedef vector<ll> v1;
    11 modpow(11 b, 11 e) {
10
        11 \text{ ans} = 1;
11
        for (; e; b = b * b % mod, e /= 2)
12
            if (e & 1) ans = ans * b % mod;
        return ans:
13
14
15
  void ntt(vl &a) {
16
        int n = sz(a), L = 31 - \underline{builtin_clz(n)};
17
        static v1 rt(2, 1);
18
        for (static int k = 2, s = 2; k < n; k *= 2, s++) {
19
            rt.resize(n):
20
             ll z[] = \{1, modpow(root, mod >> s)\};
21
            f(i,k,2*k) rt[i] = rt[i / 2] * z[i & 1] % mod;
22
23
        vector<int> rev(n);
24
        f(i,0,n) \text{ rev}[i] = (\text{rev}[i / 2] | (i \& 1) << L) / 2;
25
        f(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
26
        for (int k = 1; k < n; k *= 2)
27
            for (int i = 0; i < n; i += 2 * k) f(j, 0, k) {
                11 z = rt[j + k] * a[i + j + k] % mod, &ai = a[i + j];
29
                a[i + j + k] = ai - z + (z > ai ? mod : 0);
30
                ai += (ai + z >= mod ? z - mod : z);
31
32
33
   vl conv(const vl &a, const vl &b) {
        if (a.empty() || b.empty()) return {};
34
35
        int s = sz(a) + sz(b) - 1, B = 32 - _builtin_clz(s), n = 1 << B;
        int inv = modpow(n, mod - 2);
37
        vl L(a), R(b), out(n);
38
        L.resize(n), R.resize(n);
39
        ntt(L), ntt(R);
40
        f(i,0,n) out [-i \& (n-1)] = (l1)L[i] * R[i] % mod * inv % mod;
41
        ntt(out);
42
        return {out.begin(), out.begin() + s};
43
44
    vector<int> v;
45
   vector<ll> solve(int s, int e) {
46
        if(s==e) {
47
            vector<11> res(2);
48
            res[0] = 1;
49
            res[1] = v[s];
50
            return res;
51
52
53
        return conv(solve(s, md), solve(md+1, e));
54
```

NTT of KACTL

6

3.2 2D Sparse table

```
1 /*
        note this isn't the best cache-wise version
        query O(1), Build O(NMlqNlqM)
        be careful when using it and note the he build a dimension above another
        i.e he builds a sparse table for each row
        the build sparse table over each row's sparse table
8
    const int N = 505, LG = 10;
   int st[N][N][LG][LG];
11
    int a[N][N], lg2[N];
12
13
    int yo(int x1, int y1, int x2, int y2) {
14
     x2++:
      y2++;
15
16
      int a = 1g2[x2 - x1], b = 1g2[y2 - y1];
17
      return max (
18
             \max(st[x1][y1][a][b], st[x2 - (1 << a)][y1][a][b]),
19
             \max(st[x1][y2 - (1 << b)][a][b], st[x2 - (1 << a)][y2 - (1 << b)][a][b]
                  ])
20
           );
21
22
23
    void build(int n, int m) { // 0 indexed
24
      for (int i = 2; i < N; i++) lg2[i] = lg2[i >> 1] + 1;
25
      for (int i = 0; i < n; i++) {
26
        for (int j = 0; j < m; j++) {
27
          st[i][j][0][0] = a[i][j];
28
29
30
      for (int a = 0; a < LG; a++) {
31
        for (int b = 0; b < LG; b++) {
32
          if (a + b == 0) continue;
33
          for (int i = 0; i + (1 << a) <= n; i++) {
34
            for (int j = 0; j + (1 << b) <= m; <math>j++) {
35
36
                st[i][j][a][b] = max(st[i][j][a][b-1], st[i][j+(1 << (b-1))][a]
                     ][b - 1]);
              } else {
38
                st[i][j][a][b] = max(st[i][j][a - 1][b], st[i + (1 << (a - 1)))[j][a]
39
40
41
42
43
44
```

3.3 hillbert Order

```
///Faster Sorting MO
3
    const int infinity = (int)1e9 + 42;
    const int64_t llInfinity = (int64_t)1e18 + 256;
5
    const int module = (int)1e9 + 7;
    const long double eps = 1e-8;
    inline int64_t gilbertOrder(int x, int y, int pow, int rotate) {
Q
        if (pow == 0) {
10
             return 0;
11
12
        int hpow = 1 << (pow-1);</pre>
13
        int seq = (x < hpow) ? (
14
             (y < hpow) ? 0 : 3
15
16
             (y < hpow) ? 1 : 2
17
        seg = (seg + rotate) & 3;
```

```
const int rotateDelta[4] = {3, 0, 0, 1};
20
        int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
21
        int nrot = (rotate + rotateDelta[seq]) & 3;
        int64 t subSquareSize = int64 t(1) << (2*pow - 2);
23
        int64_t ans = seg * subSquareSize;
24
        int64_t add = gilbertOrder(nx, ny, pow-1, nrot);
25
        ans += (seg == 1 || seg == 2) ? add : (subSquareSize - add - 1);
26
27
    struct Query {
29
30
        int 1, r, idx;
31
        int64_t ord;
32
33
        inline void calcOrder() {
34
            ord = gilbertOrder(1, r, 21, 0);
35
36
    };
37
38
    inline bool operator<(const Query &a, const Query &b) {</pre>
39
        return a.ord < b.ord;</pre>
40
41
    signed main() {
        #ifndef USE_FILE_IO
44
            ios_base::sync_with_stdio(false);
45
46
47
        mt19937 rnd(42);
48
49
        int n, m, k; cin >> n >> m; k = rnd() % 1048576;
50
        vector<int> p(n+1);
51
        for (int i = 0; i < n; i++) {
            int val = rnd() % 1048576;
53
            p[i+1] = p[i] ^ val;
54
55
56
        vector<Query> gry(m);
57
        for (int i = 0; i < m; i++) {
            int 1 = rnd() % n + 1, r = rnd() % n + 1;
            if (1 > r) {
                swap(1, r);
            qry[i].1 = 1; qry[i].r = r;
            qry[i].idx = i;
64
            qry[i].calcOrder();
65
66
67
        int64_t ans = 0;
68
        vector<int64_t> res(m);
69
        vector<int64_t> cnt((int)2e6, 0);
70
        sort(qry.begin(), qry.end());
71
        int 1 = 0, r = 1;
72
        ans = (p[1] == k);
73
        cnt[p[0]]++; cnt[p[1]]++;
74
75
        for (Query q: qry) {
76
            while (1 > q.1) {
77
78
79
                ans += cnt[p[1] ^ k];
80
                cnt[p[1]]++;
81
82
            while (r < q.r) {
83
84
                ans += cnt[p[r] ^ k];
85
                cnt[p[r]]++;
86
87
             while (1 < q.1) {
88
                cnt[p[1]]--;
89
                ans -= cnt[p[1] ^ k];
90
                1++;
91
             while (r > q.r) {
                cnt[p[r]]--;
```

```
ans -= cnt[p[r] ^{\circ} k];
95
97
              res[q.idx] = ans;
98
99
100
          uint64_t rhsh = 0;
          for (int i = 0; i < m; i++) {
101
102
              rhsh *= (uint64_t)1e9 + 7;
103
              rhsh += (uint64_t)res[i];
104
105
          cout << rhsh << "\n";</pre>
106
107
          return 0;
108
```

3.4 Merge Sort Bit with updates

```
//O(log ^ 2 N) updates and queries
 3
    #include <ext/pb_ds/tree_policy.hpp>
    #include <ext/pb_ds/assoc_container.hpp>
    #include <ext/rope>
    using namespace std;
    using namespace __gnu_pbds;
10
    using namespace __gnu_cxx;
11
12
    template<class T> using Tree = tree<T, null_type, less<T>, rb_tree_tag,
         tree_order_statistics_node_update>;
14
15
    Tree<int> t[N];
16
17
    void add(int idx, int v) {
18
         for (int x = ++idx; x < N; x += x & -x) {
19
             t[x].insert(v);
20
21
    void erase(int idx, int v) {
         for (int x = ++idx; x < N; x += x & -x)
^{24}
            t[x].erase(v);
25
26
    int get(int idx, int limit){
        int ret = 0;
         for (int x = ++idx; x; x -= x & -x)
29
            ret += (t[x].order_of_key(limit+1));
30
31
```

3.5 Mo's

```
#include <bits/stdc++.h>
    int n, qq, arr[N], sz = 1000; // sz is the size of the bucket
    int co[N], ans = 0, ansq[N];
    int cul = 1, cur = 1;
    void add(int x) {
8
        co[arr[x]]++;
Q
        if (co[arr[x]] == 1)
10
            ans++:
        else if (co[arr[x]] == 2)
12
            ans--;
13
14
15
    void remove(int x) {
        co[arr[x]]--;
```

```
17
         if (co[arr[x]] == 1)
18
19
         else if (co[arr[x]] == 0)
20
21
22
23
    void solve(int 1, int r,int ind) {
25
        while (cul < 1) remove(cul++);</pre>
26
        while (cul > 1) add(--cul);
27
        while (cur < r) add(cur++);</pre>
        while (cur > r) remove(--cur);
        ansq[ind] = ans;
30
31
32
33
   int main() {
35
        cin >> qq;
36
                                  \{1/sz,r\},
                                             { 1 , ind}
37
       priority_queue<pair<pair<int, int>, pair<int, int>>, vector<pair<pair<int,</pre>
            int>, pair<int, int>>>, greater<pair<pair<int, int>, pair<int, int>>>> q
38
         for (int i = 0; i < qq; i++) {</pre>
39
             int 1, r;
40
             cin >> 1 >> r;
41
             q.push({{1 / sz, r},{1,i}});
42
43
        while (q.size()) {
44
             int ind=q.top().second.second,l=q.top().second.first,r=q.top().first.
             solve(l, r,ind);
46
             q.pop();
47
48
        for (int i = 0; i < qq; i++)
49
             cout << ansq[i] << endl;</pre>
50
51
52
        return 0;
53
```

3.6 Mo With Updates

```
///O(N^5/3) note that the block size is not a standard size
    #pragma GCC optimize ("03")
    #pragma GCC target ("sse4")
    #include <bits/stdc++.h>
    using namespace std;
10
11
    using 11 = long long;
12
13
    const int N = 1e5 +5;
    const int M = 2 * N;
    const int blk = 2155;
    const int mod = 1e9 + 7;
17
    struct Query{
18
     int 1, r, t, idx;
       Query (int a = 0, int b = 0, int c = 0, int d = 0) \{1=a, r=b, t=c, idx = d\}
      bool operator < (Query o) {</pre>
        if(r / blk == o.r / blk && 1 / blk == o.l / blk)return t < o.t;</pre>
        if(r / blk == o.r / blk)return l < o.l;</pre>
        return r < o.r;
^{24}
25
   } Q[N];
27 int a[N], b[N];
28 int cnt1[M], cnt2[N];
   int L = 0, R = -1, K = -1;
```

```
31
     // cout << x << '\n';
      cnt2[cnt1[x]]--;
33
       cnt1[x]++;
34
       cnt2[cnt1[x]]++;
35
36
     void del(int x){ ///delete item from range
       cnt2[cnt1[x]]--;
38
       cnt1[x]--;
39
       cnt2[cnt1[x]]++;
40
41 map<int,int>id;
42 int cnt;
43 int ans[N];
44
     int p[N], nxt[N];
45
46
     void upd(int idx) { ///update item value
       if(p[idx] >= L && p[idx] <= R)
48
         del(a[p[idx]]), add(nxt[idx]);
49
       a[p[idx]] = nxt[idx];
 50
51
     void err(int idx) {
52
       if(p[idx] >= L \&\& p[idx] <= R)
53
          del(a[p[idx]]), add(prv[idx]);
54
       a[p[idx]] = prv[idx];
55
56
     int main(){
57
58
       int n, q, 1, r, tp;
59
60
       scanf("%d%d", &n, &q);
61
62
       for (int i = 0; i < n; i++) {
63
         scanf("%d", a + i);
64
         if(id.count(a[i]) == 0)
           id[a[i]] = cnt++;
65
66
         a[i] = id[a[i]];
67
         b[i] = a[i];
68
69
       int qIdx = 0;
70
       int ord = 0:
71
       while (q--) {
72
\begin{array}{c} 73 \\ 74 \end{array}
         scanf("%d", &tp);
         if(tp == 1) {
75
76
           /// ADD Query
            scanf("%d%d", &1, &r); --1, --r;
77
           Q[qIdx] = Query(1,r,ord-1,qIdx); qIdx++;
          } else{
 79
           /// ADD Update
 80
           scanf("%d%d",p + ord, nxt + ord); --p[ord];
81
           if(id.count(nxt[ord]) == 0)
            id[nxt[ord]] = cnt++;
 83
           nxt[ord] = id[nxt[ord]];
 84
            prv[ord] = b[p[ord]];
 85
           b[p[ord]] = nxt[ord];
 86
           ++ord;
87
 88
89
90
       sort(Q,Q+qIdx);
91
       for(int i = 0; i < qIdx; i++) {</pre>
         while (L < Q[i].1) del(a[L++]);
92
93
          while (L > Q[i].1) add (a[--L]);
94
         while (R < Q[i].r) add (a[++R]);
 95
         while (R > Q[i].r) del(a[R--]);
96
         while (K < Q[i].t) upd (++K);
97
         while (K > Q[i].t) err(K--);
98
          ///Solve Query I
99
100
       for(int i = 0; i < qIdx; i++)</pre>
101
         printf("%d\n", ans[i]);
102
```

 $\frac{103}{104}$

return 0;

void add(int x) { ///add item to range

105 }

3.7 Ordered Set

3.8 Persistent Seg Tree

```
2
    int val[ N * 60 ], L[ N * 60 ], R[ N * 60 ], ptr, tree[N]; /// N * 1gN
    int upd(int root, int s, int e, int idx) {
        int ret = ++ptr;
 5
        val[ret] = L[ret] = R[ret] = 0;
        if (s == e) {
            val[ret] = val[root] + 1;
            return ret:
 Q
10
        int md = (s + e) >> 1;
11
        if (idx <= md) {
12
            L[ret] = upd(L[root], s, md, idx), R[ret] = R[root];
13
        } else {
14
            R[ret] = upd(R[root], md + 1, e, idx), L[ret] = L[root];
15
16
        val[ret] = max(val[L[ret]], val[R[ret]]);
17
        return ret;
18
19
  int qry(int node, int s, int e, int 1, int r){
20
      if(r < s || e < 1 || !node)return 0; //Punishment Value</pre>
      if(1 <= s && e <= r) {
22
        return val[node];
23
      int md = (s+e) >> 1;
25
      return max(qry(L[node], s, md, l, r), qry(R[node], md+1, e, l, r));
26
27
    int merge(int x, int y, int s, int e) {
28
        if(!x||!y)return x | y;
29
        if(s == e) {
30
            val[x] += val[y];
31
            return x;
32
33
        int md = (s + e) >> 1;
        L[x] = merge(L[x], L[y], s, md);
34
35
        R[x] = merge(R[x], R[y], md+1,e);
        val[x] = val[L[x]] + val[R[x]];
36
37
        return x;
38
```

3.9 Sqrt Decomposition

```
1  // Source: https://cp-algorithms.com/data_structures/sqrt_decomposition.html
2
3  // input data
4  int n;
5  vector<int> a (n);
6
7  // preprocessing
```

```
int len = (int) sqrt (n + .0) + 1; // size of the block and the number of blocks
    vector<int> b (len);
10
    for (int i=0; i<n; ++i)</pre>
        b[i / len] += a[i];
11
12
13
    // answering the gueries
14
    for (;;) {
15
        int 1, r;
16
      // read input data for the next query
17
        int sum = 0;
18
         for (int i=1; i<=r; )</pre>
19
             if (i % len == 0 && i + len - 1 <= r) {
20
                 // if the whole block starting at i belongs to [1, r]
21
                 sum += b[i / len];
22
                 i += len;
23
24
             else {
25
                 sum += a[i];
26
                 ++i;
27
28
29
30
    // If you're getting TLE and can't optimize more, you could reduce the number of
          slow division operations using the following code:
31
32
    int sum = 0;
33
    int c_l = 1 / len,    c_r = r / len;
34
    if (c_l == c_r)
35
        for (int i=1; i<=r; ++i)</pre>
36
             sum += a[i];
37
    else {
38
         for (int i=1, end=(c_1+1)*len-1; i<=end; ++i)</pre>
39
             sum += a[i];
40
         for (int i=c_l+1; i<=c_r-1; ++i)</pre>
41
             sum += b[i];
42
         for (int i=c_r*len; i<=r; ++i)</pre>
43
             sum += a[i];
44
```

3.10 Treap

```
typedef struct item * pitem;
2
    struct item {
         int prior, value, cnt;
4
        bool rev;
5
        pitem 1, r;
6
         item(int x, int y, int z){
7
             value = x;
             prior = y;
8
9
             cnt = z;
             rev = 0;
10
11
             1 = r = NULL;
12
13
    };
14
15
    int cnt (pitem it) {
16
         return it ? it->cnt : 0;
17
18
19
    void upd_cnt (pitem it) {
20
         if (it)
21
             it\rightarrow cnt = cnt(it\rightarrow 1) + cnt(it\rightarrow r) + 1;
22
23
24
    void push (pitem it) {
25
        if (it && it->rev) {
26
             it->rev = false;
27
             swap (it->1, it->r);
             if (it->1) it->1->rev ^= true;
29
             if (it->r) it->r->rev ^= true;
30
31
    }
```

```
void merge (pitem & t, pitem l, pitem r) {
34
        push (1);
35
        push (r);
36
        if (!l || !r)
37
            t = 1 ? 1 : r;
38
        else if (1->prior > r->prior)
39
            merge (1->r, 1->r, r), t = 1;
40
41
            merge (r->1, 1, r->1), t = r;
42
        upd_cnt (t);
43
44
45
    void split (pitem t, pitem & 1, pitem & r, int key, int add = 0) {
46
47
            return void(1 = r = 0);
         push (t);
48
49
        int cur_key = add + cnt(t->1);
50
        if (key <= cur_key)</pre>
51
            split (t->1, 1, t->1, key, add), r = t;
52
53
            split (t->r, t->r, r, key, add + 1 + cnt(t->1)), 1 = t;
54
        upd_cnt (t);
55
56
57
    void reverse (pitem t, int 1, int r) {
58
        pitem t1, t2, t3;
59
        split (t, t1, t2, 1);
        split (t2, t2, t3, r-l+1);
61
        t2->rev ^= true;
62
        merge (t, t1, t2);
63
        merge (t, t, t3);
64
66
   void output (pitem t) {
67
        if (!t) return;
68
        push (t);
69
        output (t->1);
70
        printf ("%c", char(t->value));
71
        output (t->r);
72
73
74
    pitem gettreap(string s) {
75
            pitem ret=NULL;
76
77
           for(i=0;i<s.size();i++)merge(ret,ret,new item(s[i],(rand()<<15)+rand(),</pre>
78
        return ret;
79
```

3.11 Wavelet Tree

```
1 // remember your array and values must be 1-based
    struct wavelet_tree {
        int lo, hi;
        wavelet_tree *1, *r;
 5
        vector<int> b;
 6
        //nos are in range [x, y]
 8
        //array indices are [from, to)
        wavelet_tree(int *from, int *to, int x, int y) {
 9
10
            lo = x, hi = y;
11
            if (lo == hi or from >= to)
12
                return;
13
            int mid = (lo + hi) / 2;
14
            auto f = [mid] (int x) {
15
                return x <= mid;</pre>
16
17
            b.reserve(to - from + 1);
18
            b.pb(0);
19
            for (auto it = from; it != to; it++)
20
                b.pb(b.back() + f(*it));
```

```
10
```

```
//see how lambda function is used here
22
             auto pivot = stable_partition(from, to, f);
23
             l = new wavelet_tree(from, pivot, lo, mid);
             r = new wavelet_tree(pivot, to, mid + 1, hi);
25
26
27
         //kth smallest element in [1, r]
28
        int kth(int 1, int r, int k) {
29
             if (1 > r)
30
                 return 0;
31
             if (lo == hi)
32
                 return lo;
33
             int inLeft = b[r] - b[1 - 1];
             int lb = b[1 - 1]; //amt of nos in first (1-1) nos that go in left
35
             int rb = b[r]; //amt of nos in first (r) nos that go in left
36
             if (k <= inLeft)</pre>
37
                 return this->l->kth(lb + 1, rb, k);
38
             return this->r->kth(l - lb, r - rb, k - inLeft);
39
40
41
         //count of nos in [1, r] Less than or equal to k
42
         int LTE(int 1, int r, int k) {
43
             if (1 > r \text{ or } k < 10)
44
                return 0;
45
             if (hi <= k)
46
                return r - 1 + 1;
47
             int 1b = b[1 - 1], rb = b[r];
48
             return this->1->LTE(1b + 1, rb, k) + this->r->LTE(1 - 1b, r - rb, k);
49
50
51
         //count of nos in [l, r] equal to k
52
         int count(int 1, int r, int k) {
53
             if (1 > r \text{ or } k < 10 \text{ or } k > hi)
54
                return 0;
55
             if (lo == hi)
56
                return r - 1 + 1;
57
             int lb = b[1 - 1], rb = b[r], mid = (lo + hi) / 2;
58
             if (k <= mid)
                return this->l->count(lb + 1, rb, k);
59
60
             return this->r->count(1 - 1b, r - rb, k);
61
62
```

4 DP

4.1 Dynamic Convex Hull Trick

```
#include<iostream>
    #include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    struct Line
Q
10
         mutable function<const Line*()> succ;
11
         bool operator<(const Line& other) const
12
13
             return m < other.m;</pre>
14
15
         bool operator<(const 11 &x) const
16
17
             const Line* s = succ();
             if (!s)
19
                 return 0;
20
             return b - s \rightarrow b < (s \rightarrow m - m) * x;
21
    // will maintain upper hull for maximum
```

```
struct HullDynamic : public multiset<Line, less<>>
25
26
        bool bad(iterator y)
27
28
            auto z = next(y);
29
            if (y == begin())
30
                if (z == end())
32
                    return 0;
33
                return y->m == z->m && y->b <= z->b;
34
35
            auto x = prev(y);
            if (z == end())
37
                return y->m == x->m && y->b <= x->b;
38
            return (ld) (x-b - y-b)*(z-m - y-m) >= (ld) (y-b - z-b)*(y-m - x-m)
40
        void insert_line(ll m, ll b)
41
42
            auto y = insert({ m, b });
43
            y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
44
            if (bad(y))
45
46
                erase(y);
47
                return;
48
            while (next(y) != end() && bad(next(y)))
                erase(next(y));
51
            while (y != begin() && bad(prev(y)))
                erase(prev(y));
53
55
        11 query(11 x)
56
57
58
            auto 1 = *lower_bound(x);
59
            return 1.m * x + 1.b;
60
61
    };
62
    int main()
63
64
65
66
```

4.2 Dynamic Connectivety with SegTree

```
/// MANGA
    #pragma GCC optimize("03")
    #pragma GCC optimize ("unroll-loops")
    #pragma GCC target("avx,avx2,fma")
    using namespace std;
    #include "bits/stdc++.h"
    #define pb push_back
    #define F first
    #define S second
    #define f(i, a, b) for (int i = a; i < b; i++)
    #define all(a) a.begin(),a.end()
    #define rall(a) a.rbegin(),a.rend()
    #define sz(x) (int)(x).size()
    //#define mp make_pair
17
    #define popCnt(x) (__builtin_popcountll(x))
18
    typedef long long 11;
    typedef pair<int, int> ii;
    using ull = unsigned long long;
    const int N = 1e5+5, LG = 17, MOD = 1e9 + 7;
    const long double PI = acos(-1);
    struct PT{
        11 x, y;
        PT() {}
```

```
26
        PT(11 a, 11 b):x(a), y(b){}
27
         PT operator - (const PT & o) {return PT{x-o.x,y-o.y};}
28
        bool operator < (const PT & o) const {return make_pair(x,y) < make_pair(o.x,</pre>
29
30
    ll cross(PT x, PT y) {
31
        return x.x * y.y - x.y * y.x;
32
33
    PT val[300005];
    bool in[300005];
35
    11 gr[300005];
    bool ask[300005];
                                                                                         107
37
    11 ans[N];
                                                                                         108
38
    vector<PT> t[300005 * 4]; ///segment tree holding points to queries
39
    void update(int node, int s, int e, int 1, int r, PT x) {
                                                                                         110
40
        if(r < s || e < 1)return;</pre>
                                                                                         111
41
        if(1 \le s \&\& e \le r)  { ///add this point to maximize it with queries in
                                                                                         112
             this range
                                                                                         113
42
             t[node].pb(x);
                                                                                         114
43
             return;
                                                                                         115
44
                                                                                         116
45
         int md = (s + e) \gg 1;
                                                                                         117
        update(node<<1,s,md,1,r,x);
                                                                                         118
47
        update(node<<1|1,md+1,e,1,r,x);
                                                                                         119
48
49
   vector<PT> stk;
50
   inline void addPts(vector<PT> v) {
                                                                                         122
51
        stk.clear(); ///reset the data structure you are using
                                                                                         123
52
        sort(all(v));
53
        ///build upper envelope
54
         for(int i = 0; i < v.size(); i++) {</pre>
55
             while(sz(stk) > 1 && cross(v[i] - stk.back(), stk.back() - stk[stk.size
                                                                                         127
                  ()-2]) <= 0)
                 stk.pop_back();
                                                                                         129
57
             stk.push_back(v[i]);
58
                                                                                         131
59
                                                                                         132
60
    inline 11 calc(PT x, 11 val) {
                                                                                         133
61
                                                                                         134
62
         return x.x * val + x.y;
                                                                                         135
63
                                                                                         136
64
                                                                                         137
65
   11 query(11 x) {
                                                                                         138
66
        if(stk.empty())
                                                                                         139
67
             return LLONG_MIN;
                                                                                         140
68
         int lo = 0, hi = stk.size() - 1;
69
        while (lo + 10 < hi) {
            int md = lo + (hi-lo) / 2;
70
71
             if(calc(stk[md+1],x) > calc(stk[md],x))
72
                 10 = md + 1;
73
             else
74
                hi = md;
75
76
         11 ans = LLONG_MIN;
77
         for (int i = lo; i <= hi; i++)</pre>
78
            ans = max(ans, calc(stk[i], x));
79
         return ans;
80
81
    void solve(int node, int s, int e) {      ///Solve queries
82
        addPts(t[node]);
                             ///note that there is no need to add/delete just build
             for t[node]
83
         f(i,s,e+1){
84
             if(ask[i]) {
85
                 ans[i] = max(ans[i], query(qr[i]));
86
87
88
         if(s==e) return;
89
        int md = (s + e) >> 1;
90
         solve(node<<1,s,md);</pre>
91
         solve(node<<1|1,md+1,e);
92
93
    void doWork() {
94
95
        int n:
```

```
cin >> n;
    stk.reserve(n);
    f(i,1,n+1) {
        int tp;
        cin >> tp;
        if(tp == 1) {
                       ///Add Ouerv
            int x, y;
            cin >> x >> y;
            val[i] = PT(x, y);
            in[i] = 1;
           else if(tp == 2) { ///Delete Query
            int x;
            cin >> x;
            if(in[x])update(1, 1, n, x, i - 1, val[x]);
            in[x] = 0;
          else {
            cin >> qr[i];
            ask[i] = true;
    f(i,1,n+1) ///Finalize Query
        if(in[i])
            update(1, 1, n, i, n, val[i]);
    f(i,1,n+1) ans [i] = LLONG_MIN;
    solve(1, 1, n);
    f(i,1,n+1)
    if(ask[i])
        if(ans[i] == LLONG MIN)
            cout << "EMPTY SET\n";
            cout << ans[i] << '\n';
int32_t main() {
#ifdef ONLINE JUDGE
    ios_base::sync_with_stdio(0);
    cin.tie(0);
#endif // ONLINE JUDGE
   int t = 1;
    cin >> t;
    while (t--) {
        doWork();
    return 0;
```

4.3 Li Chao Tree

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130

141

142

```
1 #include<iostream>
    #include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    struct Line
 9
        11 m, b;
10
        Line(ll m, ll b) : m(m), b(b) {}
11
        11 operator()(11 x)
12
13
            return m * x + b;
14
   };
15
16
    struct node
17
18
        node * left,* right ;
19
20
        node(node * left, node *right, Line line):left(left), right(right), line(
             line) {}
        node * getLeft()
```

```
23
             if(left==NULL)
24
                 left= new node (NULL, NULL, Line(0, 1e18));
25
             return left;
26
27
        node * getright()
28
29
             if(right==NULL)
30
                 right= new node (NULL, NULL, Line (0, 1e18));
31
             return right;
32
33
         void insert(Line newline, int 1, int r)
34
35
             int m = (1+r)/2;
36
             bool lef=newline(1)<line(1);</pre>
37
             bool mid=newline(m) <line(m);</pre>
38
39
             if (mid)
40
                 swap(line, newline);
41
             if(r-l==1)
42
                 return ;
43
             else if(lef!=mid)
44
                 getLeft()->insert(newline,1,m);
45
46
                 getright()->insert(newline,m,r);
47
48
         11 query(int x, int 1, int r)
49
50
             int m = (1 + r) / 2;
51
             if(r - 1 == 1)
52
                 return line(x);
53
             else if (x < m)
54
                 return min(line(x), getLeft()->query(x, 1, m));
55
56
                 return min(line(x), getright()->query(x, m, r));
57
58
         void deletee()
59
60
             if(left!=NULL)
61
                 left->deletee();
62
             if(right!=NULL)
63
                 right->deletee();
64
             free(this);
65
66
    };
67
    int main()
68
69
70
        node * root = new node(NULL, NULL, Line(0,5));
71
         root->insert(Line(1,-3),1,100);
72
73
         for (int i=1; i<=10; i++)</pre>
74
             cout<<root->query(i,1,100)<<"\n";</pre>
75
```

4.4 CHT Line Container

```
struct Line
2
3
         mutable 11 m, b, p;
        bool operator<(const Line& o) const
4
5
6
             return m < o.m;</pre>
        bool operator<(11 x) const
10
             return p < x;
11
12
    };
13
14
    struct LineContainer : multiset<Line, less<>>
15
```

```
16
         // (for doubles, use inf = 1/.0, div(a,b) = a/b)
17
        static const 11 inf = LLONG_MAX;
18
        ll div(ll db, ll dm) // floored division
19
20
            return db / dm - ((db ^ dm) < 0 && db % dm);
21
22
        bool isect(iterator x, iterator y)
23
24
            if (y == end())
25
26
                x->p = inf:
27
                return false;
29
            if (x->m == y->m)
30
                x->p = x->b > y->b ? inf : -inf;
31
32
                x->p = div(y->b - x->b, x->m - y->m);
33
            return x->p >= y->p;
34
35
        void add(ll m, ll b)
36
37
            auto z = insert(\{m, b, 0\}), y = z++, x = y;
38
            while (isect(y, z))
39
                z = erase(z);
40
            if (x != begin() && isect(--x, y))
41
                isect(x, y = erase(y));
            while ((y = x) != begin() \&\& (--x)->p >= y->p)
42
43
                isect(x, erase(y));
44
45
        11 query(11 x)
46
47
            assert(!empty());
48
            auto 1 = *lower_bound(x);
49
            return 1.m * x + 1.b;
50
51
```

5 Geometry

5.1 Convex Hull

```
1
    struct point {
        11 x, y;
 3
        point(11 x, 11 y) : x(x), y(y) {}
        point operator -(point other) {
 5
             return point(x - other.x, y - other.y);
 6
        bool operator <(const point &other) const {</pre>
 8
             return x != other.x ? x < other.x : y < other.y;</pre>
 9
10
    11 cross(point a, point b) {
11
12
        return a.x * b.y - a.y * b.x;
13
    11 dot(point a, point b) {
15
        return a.x * b.x + a.y * b.y;
16
17
    struct sortCCW {
18
        point center;
19
20
        sortCCW(point center) : center(center) {}
21
^{22}
        bool operator()(point a, point b) {
23
            11 res = cross(a - center, b - center);
24
            if(res)
25
                 return res > 0;
26
            return dot(a - center, a - center) < dot(b - center, b - center);</pre>
27
28
    vector<point> hull(vector<point> v) {
```

```
_
```

```
30
        sort(v.begin(), v.end());
31
        sort(v.begin() + 1, v.end(), sortCCW(v[0]));
32
        v.push back(v[0]);
33
        vector<point> ans ;
34
        for(auto i : v) {
35
            int sz = ans.size();
36
            while (sz > 1 \&\& cross(i - ans[sz - 1], ans[sz - 2] - ans[sz - 1]) \le 0)
37
               ans.pop_back(), sz--;
            ans.push back(i);
39
40
        ans.pop back();
41
        return ans;
42
```

5.2 Geometry Template

```
using ptype = double edit this first ;
    double EPS = 1e-9;
    struct point {
        ptype x, y;
        point(ptype x, ptype y) : x(x), y(y) {}
        point operator - (const point & other) const {
            return point(x - other.x, y - other.y);
10
11
12
        point operator + (const point & other) const {
13
            return point(x + other.x, y + other.y);
14
15
16
        point operator *(ptype c) const {
17
            return point(x * c, y * c);
18
19
20
        point operator / (ptype c) const {
21
            return point(x / c, y / c);
22
23
        point prep() {
24
            return point(-y, x);
25
26
27
    ptype cross(point a, point b) {
29
        return a.x * b.y - a.y * b.x;
30
31
32
    ptype dot(point a, point b) {
33
        return a.x * b.x + a.y * b.y;
34
35
    double abs(point a) {
36
        return sqrt (dot (a, a));
37
38
    // angle between [0 , pi]
39
    double angle (point a, point b) {
        return acos(dot(a, b) / abs(a) / abs(b));
41
42
   // a : point in Line
   // d : Line direction
    point LineLineIntersect(point a1, point d1, point a2, point d2) {
45
        return a1 + d1 * cross(a2 - a1, d2) / cross(d1, d2);
46
47 // Line a---b
   // point C
49
   point ProjectPointLine(point a, point b, point c) {
        return a + (b - a) * 1.0 * dot(c - a, b - a) / dot(b - a, b - a);
50
51
   // segment a---b
   // point C
54
    point ProjectPointSegment(point a, point b, point c) {
55
        double r = dot(c - a, b - a) / dot(b - a, b - a);
        if(r < 0)
```

```
return a;
58
         if(r > 1)
 59
             return b;
         return a + (b - a) * r;
61
    // Line a---b
63
    // point p
    point reflectAroundLine(point a, point b, point p) {
         //(proj-p) *2 + p
66
         return ProjectPointLine(a, b, p) * 2 - p;
67
     // Around origin
68
    point RotateCCW(point p, double t) {
70
         return point(p.x * cos(t) - p.y * sin(t),
71
                      p.x * sin(t) + p.y * cos(t));
72
 73
     // Line a---b
74 vector<point> CircleLineIntersect(point a, point b, point center, double r) {
        a = a - center;
76
        b = b - center;
         point p = ProjectPointLine(a, b, point(0, 0)); // project point from center
              to the Line
         if(dot(p, p) > r * r)
 79
             return {};
80
         double len = sqrt(r * r - dot(p, p));
81
         if(len < EPS)</pre>
82
            return {center + p};
83
         point d = (a - b) / abs(a - b);
85
         return {center + p + d * len, center + p - d * len};
86
87
     vector<point> CircleCircleIntersect(point c1, double r1, point c2, double r2) {
88
89
         if(r1 < r2) {
90
             swap(r1, r2);
91
             swap(c1, c2);
92
93
         double d = abs(c1 - c2); // distance between c1,c2
         if(d > r1 + r2 || d < r1 - r2)
95
             return {};
96
97
         double angle = acos(min((d * d + r1 * r1 - r2 * r2) / (2 * r1 * d), 1.0));
98
         point p = (c2 - c1) / d * r1;
100
         if(angle < EPS)</pre>
101
             return {p};
102
103
         return {RotateCCW(p, angle), RotateCCW(p, -angle)};
104
105
106 point circumcircle(point p1, point p2, point p3) {
107
108
         return LineLineIntersect((p1 + p2) / 2, (p1 - p2).prep(),
109
                                  (p1 + p3) / 2, (p1 - p3).prep());
110
111 //S : Area.
    //I : number points with integer coordinates lying strictly inside the polygon.
113 //B : number of points lying on polygon sides by B.
114 	 //S = I + B/2 - 1
```

5.3 Half Plane Intersection

```
// Redefine epsilon and infinity as necessary. Be mindful of precision errors.
const long double eps = le-9, inf = le9;

// Basic point/vector struct.
struct Point {

long double x, y;
explicit Point(long double x = 0, long double y = 0) : x(x), y(y) {}

// Addition, substraction, multiply by constant, cross product.
```

```
12
        friend Point operator + (const Point& p, const Point& q) {
13
             return Point(p.x + q.x, p.y + q.y);
14
15
16
         friend Point operator - (const Point& p, const Point& q) {
17
             return Point(p.x - q.x, p.y - q.y);
18
19
20
        friend Point operator * (const Point& p, const long double& k) {
21
             return Point(p.x * k, p.y * k);
22
23
^{24}
        friend long double cross(const Point& p, const Point& q) {
25
            return p.x * q.y - p.y * q.x;
26
27
    };
    // Basic half-plane struct.
    struct Halfplane {
31
         // 'p' is a passing point of the line and 'pq' is the direction vector of
32
             the line.
33
        Point p, pq;
34
        long double angle;
35
36
        Halfplane() {}
        Halfplane (const Point& a, const Point& b) : p(a), pq(b - a) {
38
             angle = atan21(pq.y, pq.x);
39
40
41
        // Check if point 'r' is outside this half-plane.
        // Every half-plane allows the region to the LEFT of its line.
43
        bool out(const Point& r) {
44
            return cross(pq, r - p) < -eps;</pre>
45
46
47
        // Comparator for sorting.
48
         // If the angle of both \hat{h}alf-planes is equal, the leftmost one should go
49
        bool operator < (const Halfplane& e) const {</pre>
50
             if (fabsl(angle - e.angle) < eps) return cross(pg, e.p - p) < 0;</pre>
51
             return angle < e.angle;</pre>
52
53
54
        // We use equal comparator for std::unique to easily remove parallel half-
        bool operator == (const Halfplane& e) const {
56
             return fabsl(angle - e.angle) < eps;</pre>
57
58
59
        // Intersection point of the lines of two half-planes. It is assumed they're
              never parallel.
60
        friend Point inter(const Halfplane& s, const Halfplane& t) {
61
             long double alpha = cross((t.p - s.p), t.pq) / cross(s.pq, t.pq);
62
             return s.p + (s.pq * alpha);
63
64
    };
65
66
67
    // Actual algorithm
    vector<Point> hp_intersect(vector<Halfplane>& H) {
70
71
        Point box[4] = { // Bounding box in CCW order
72
            Point(inf, inf),
73
             Point (-inf, inf),
74
             Point (-inf, -inf),
75
            Point(inf, -inf)
76
77
78
        for(int i = 0; i<4; i++) { // Add bounding box half-planes.</pre>
79
             Halfplane aux(box[i], box[(i+1) % 4]);
80
             H.push_back(aux);
```

```
// Sort and remove duplicates
sort(H.begin(), H.end());
H.erase(unique(H.begin(), H.end()), H.end());
deque<Halfplane> dq;
int len = 0;
for(int i = 0; i < int(H.size()); i++) {</pre>
    // Remove from the back of the deque while last half-plane is redundant
    while (len > 1 && H[i].out(inter(dg[len-1], dg[len-2]))) {
        dq.pop_back();
        --len;
    // Remove from the front of the deque while first half-plane is
    while (len > 1 && H[i].out(inter(dq[0], dq[1]))) {
        dq.pop_front();
        --len;
    // Add new half-plane
    dq.push_back(H[i]);
    ++len;
// Final cleanup: Check half-planes at the front against the back and vice-
while (len > 2 && dq[0].out(inter(dq[len-1], dq[len-2]))) {
    dq.pop_back();
    --len;
while (len > 2 && dq[len-1].out(inter(dq[0], dq[1]))) {
    dq.pop_front();
    --len;
// Report empty intersection if necessary
if (len < 3) return vector<Point>();
// Reconstruct the convex polygon from the remaining half-planes.
vector<Point> ret(len);
for (int i = 0; i+1 < len; i++) {
    ret[i] = inter(dq[i], dq[i+1]);
ret.back() = inter(dq[len-1], dq[0]);
return ret;
```

5.4 Segments Intersection

83

85

86

87

90

91

95

96

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99

100

101

 $\frac{102}{103}$

104

105

106

107

108

109

110

111

112

113

114

115

116

117

118

119

120

121

122

123

124

125

 $\frac{126}{127}$

128

129

```
const double EPS = 1E-9;
 3
    struct pt {
        double x, y;
    struct seq {
        pt p, q;
        int id;
10
11
        double get_y (double x) const {
12
            if (abs(p.x - q.x) < EPS)
13
               return p.y;
14
            return p.y + (q.y - p.y) * (x - p.x) / (q.x - p.x);
15
16
17
18
    bool intersect1d(double 11, double r1, double 12, double r2) {
        if (11 > r1)
```

```
swap(11, r1);
21
         if (12 > r2)
22
             swap(12, r2);
23
         return max(11, 12) <= min(r1, r2) + EPS;</pre>
24
25
26
    int vec(const pt& a, const pt& b, const pt& c) {
         double s = (b.x - a.x) * (c.y - a.y) - (b.y - a.y) * (c.x - a.x);
         return abs(s) < EPS ? 0 : s > 0 ? +1 : -1;
28
29
30
31
    bool intersect (const seq& a, const seq& b)
\frac{32}{33}
         return intersect1d(a.p.x, a.q.x, b.p.x, b.q.x) &&
34
                intersect1d(a.p.y, a.q.y, b.p.y, b.q.y) &&
35
                vec(a.p, a.q, b.p) * vec(a.p, a.q, b.q) <= 0 &&
36
                vec(b.p, b.q, a.p) * vec(b.p, b.q, a.q) <= 0;
37
38
39
    bool operator<(const seq& a, const seq& b)
40
41
         double x = max(min(a.p.x, a.q.x), min(b.p.x, b.q.x));
         return a.get_y(x) < b.get_y(x) - EPS;</pre>
43
44
45
    struct event {
46
         double x;
47
         int tp, id;
48
49
         event() {}
50
        event(double x, int tp, int id) : x(x), tp(tp), id(id) {}
51
52
        bool operator<(const event& e) const {</pre>
53
             if (abs(x - e.x) > EPS)
54
                 return x < e.x;
55
             return tp > e.tp;
56
57
    };
58
59
60
    vector<set<seg>::iterator> where;
62
    set<seg>::iterator prev(set<seg>::iterator it) {
63
         return it == s.begin() ? s.end() : --it;
64
65
    set<seg>::iterator next(set<seg>::iterator it) {
67
         return ++it;
68
69
70
    pair<int, int> solve(const vector<seq>& a) {
        int n = (int)a.size();
71
72
         vector<event> e;
73
         for (int i = 0; i < n; ++i) {
74
             e.push_back(event(min(a[i].p.x, a[i].q.x), +1, i));
75
             e.push_back(event(max(a[i].p.x, a[i].q.x), -1, i));
76
77
        sort(e.begin(), e.end());
78
79
80
         where.resize(a.size());
81
         for (size_t i = 0; i < e.size(); ++i) {</pre>
82
             int id = e[i].id;
83
             if (e[i].tp == +1) {
84
                 set<seg>::iterator nxt = s.lower_bound(a[id]), prv = prev(nxt);
85
                 if (nxt != s.end() && intersect(*nxt, a[id]))
86
                     return make_pair(nxt->id, id);
87
                 if (prv != s.end() && intersect(*prv, a[id]))
88
                     return make_pair(prv->id, id);
89
                 where[id] = s.insert(nxt, a[id]);
90
91
                 set<seg>::iterator nxt = next(where[id]), prv = prev(where[id]);
                 if (nxt != s.end() && prv != s.end() && intersect(*nxt, *prv))
93
                     return make_pair(prv->id, nxt->id);
94
                 s.erase(where[id]);
```

20

```
95
96
97
98
        return make_pair(-1, -1);
99
```

5.5 Rectangles Union

```
#include<bits/stdc++.h>
    #define P(x,y) make_pair(x,y)
    using namespace std;
    class Rectangle {
    public:
        int x1, y1, x2, y2;
        static Rectangle empt;
        Rectangle() {
            x1 = y1 = x2 = y2 = 0;
10
11
        Rectangle (int X1, int Y1, int X2, int Y2) {
12
13
            y1 = Y1;
            x2 = X2;
14
15
            y2 = Y2;
16
    };
17
18
    struct Event {
19
        int x, y1, y2, type;
20
        Event() {}
        Event (int x, int y1, int y2, int type): x(x), y1(y1), y2(y2), type(type) {}
   bool operator < (const Event&A, const Event&B) {</pre>
   //if(A.x != B.x)
        return A.x < B.x;</pre>
    //if(A.y1 != B.y1) return A.y1 < B.y1;
    //if(A.y2 != B.y2()) A.y2 < B.y2;
    const int MX = (1 << 17);
30
    struct Node {
        int prob, sum, ans;
        Node() {}
33
        Node (int prob, int sum, int ans): prob(prob), sum(sum), ans(ans) {}
35
   Node tree[MX * 4];
36
    int interval[MX];
37
    void build(int x, int a, int b) {
        tree[x] = Node(0, 0, 0);
39
        if(a == b) {
40
            tree[x].sum += interval[a];
41
            return;
42
43
        build(x * 2, a, (a + b) / 2);
44
        build(x * 2 + 1, (a + b) / 2 + 1, b);
45
        tree[x].sum = tree[x * 2].sum + tree[x * 2 + 1].sum;
46
47
    int ask(int x) {
48
        if(tree[x].prob)
49
            return tree[x].sum;
50
        return tree[x].ans;
51
52
    int st, en, V;
    void update(int x, int a, int b) {
        if(st > b \mid \mid en < a)
55
            return;
56
        if(a >= st && b <= en) {
57
            tree[x].prob += V;
58
            return;
59
60
        update(x * 2, a, (a + b) / 2);
61
        update(x * 2 + 1, (a + b) / 2 + 1, b);
62
        tree[x].ans = ask(x * 2) + ask(x * 2 + 1);
63
    Rectangle Rectangle::empt = Rectangle();
```

```
vector < Rectangle > Rect;
     vector < int > sorted;
     vector < Event > sweep;
     void compressncalc() {
69
         sweep.clear();
 70
         sorted.clear();
71
         for(auto R : Rect) {
72
             sorted.push_back(R.y1);
73
             sorted.push_back(R.y2);
74
 75
         sort(sorted.begin(), sorted.end());
76
         sorted.erase(unique(sorted.begin(), sorted.end()), sorted.end());
77
         int sz = sorted.size();
 78
         for (int j = 0; j < sorted.size() - 1; <math>j++)
79
             interval[j + 1] = sorted[j + 1] - sorted[j];
 80
         for(auto R : Rect) {
81
             sweep.push_back(Event(R.x1, R.y1, R.y2, 1));
82
             sweep.push_back(Event(R.x2, R.y1, R.y2, -1));
83
84
         sort(sweep.begin(), sweep.end());
 85
         build(1, 1, sz - 1);
86
87 long long ans;
     void Sweep() {
89
         ans = 0:
90
         if(sorted.empty() || sweep.empty())
91
             return;
92
         int last = 0, sz_ = sorted.size();
93
         for(int j = 0; j < sweep.size(); j++) {</pre>
94
             ans += 111 * (sweep[j].x - last) * ask(1);
95
             last = sweep[j].x;
96
             V = sweep[j].type;
97
             st = lower_bound(sorted.begin(), sorted.end(), sweep[j].yl) - sorted.
                  begin() + 1;
 98
             en = lower_bound(sorted.begin(), sorted.end(), sweep[j].y2) - sorted.
                  begin();
99
             update(1, 1, sz_ - 1);
100
101
    int main() {
          freopen("in.in", "r", stdin);
103
104
         int n;
         scanf("%d", &n);
105
106
         for (int j = 1; j \le n; j++) {
107
             int a, b, c, d;
108
             scanf("%d %d %d %d", &a, &b, &c, &d);
109
             Rect.push_back(Rectangle(a, b, c, d));
110
111
         compressncalc();
112
         Sweep();
113
         cout << ans << endl;
114
```

6 Graphs

6.1 2 SAD

```
* ts.atMostOne({0,\tilde1,2}); // <= 1 of vars 0, \tilde1 and 2 are true
13
    * ts.solve(); // Returns true iff it is solvable
    * ts.values[0..N-1] holds the assigned values to the vars
     * Time: O(N+E), where N is the number of boolean variables, and E is the number
           of clauses.
     * Status: stress-tested
17
18
    #pragma once
19
20
    struct TwoSat {
21
        int N;
22
        vector<vi> gr;
23
        vi values; // 0 = false, 1 = true
24
25
        TwoSat(int n = 0) : N(n), gr(2*n) {}
26
27
        int addVar() { // (optional)
28
            gr.emplace_back();
29
            gr.emplace_back();
30
            return N++;
31
32
33
        void either(int f, int i) {
34
            f = \max(2*f, -1-2*f);
35
            j = \max(2*j, -1-2*j);
36
            gr[f].push_back(j^1);
37
            gr[j].push_back(f^1);
38
39
        void setValue(int x) { either(x, x); }
40
41
        void atMostOne(const vi& li) { // (optional)
42
            if (sz(li) <= 1) return;</pre>
            int cur = ~1i[0];
43
44
            rep(i,2,sz(li)) {
45
                int next = addVar();
46
                either(cur, ~li[i]);
47
                either(cur, next);
48
                either(~li[i], next);
49
                cur = "next;
50
51
            either(cur, ~li[1]);
52
53
54
        vi val, comp, z; int time = 0;
55
        int dfs(int i) {
            int low = val[i] = ++time, x; z.push_back(i);
56
57
            for(int e : gr[i]) if (!comp[e])
58
                low = min(low, val[e] ?: dfs(e));
59
            if (low == val[i]) do {
60
                x = z.back(); z.pop_back();
61
                comp[x] = low;
62
                if (values[x>>1] == -1)
63
                    values[x>>1] = x&1;
            } while (x != i);
64
65
            return val[i] = low;
66
67
68
        bool solve() {
69
            values.assign(N, -1);
70
            val.assign(2*N, 0); comp = val;
71
            rep(i,0,2*N) if (!comp[i]) dfs(i);
72
            rep(i,0,N) if (comp[2*i] == comp[2*i+1]) return 0;
73
            return 1;
74
   };
75
```

6.2 Ariculation Point

```
1 vector<int> adj[N];
2 int dfsn[N], low[N], instack[N], ar_point[N], timer;
3 stack<int> st;
```

```
void dfs(int node, int par) {
         dfsn[node] = low[node] = ++timer;
         int kam = 0;
          for(auto i: adj[node]) {
 9
              if(i == par) continue;
10
              if(dfsn[i] == 0){
11
                  kam++;
12
                  dfs(i, node);
13
                  low[node] = min(low[node], low[i]);
14
                  if(dfsn[node] <= low[i] && par != 0) ar_point[node] = 1;</pre>
15
16
              else low[node] = min(low[node], dfsn[i]);
17
18
         if(par == 0 && kam > 1) ar point[node] = 1;
19
20
21
     void init(int n){
         for(int i = 1; i <= n; i++) {</pre>
22
23
             adi[i].clear();
24
              low[i] = dfsn[i] = 0;
25
              instack[i] = 0;
26
              ar_point[i] = 0;
27
28
         timer = 0:
29
30
31
     int main(){
32
         int tt;
33
         cin >> tt:
34
         while (tt--) {
35
              // Input
36
              init(n);
37
              for(int i = 1; i <= n; i++) {</pre>
38
                  if(dfsn[i] == 0) dfs(i, 0);
39
40
              int c = 0;
41
              for(int i = 1; i <= n; i++) {</pre>
42
                  if(ar_point[i]) c++;
43
44
              cout << c << '\n';
45
46
          return 0;
47
```

6.3 Bridges Tree and Diameter

```
#include <bits/stdc++.h>
    #define 11 long long
    using namespace std;
    const int N = 3e5 + 5, mod = 1e9 + 7;
    vector<int> adj[N], bridge_tree[N];
    int dfsn[N], low[N], cost[N], timer, cnt, comp_id[N], kam[N], ans;
    stack<int> st;
10
11
    void dfs(int node, int par) {
12
         dfsn[node] = low[node] = ++timer;
13
         st.push(node);
14
         for(auto i: adj[node]) {
15
             if(i == par) continue;
16
             if(dfsn[i] == 0){
17
                 dfs(i, node);
18
                 low[node] = min(low[node], low[i]);
19
20
             else low[node] = min(low[node], dfsn[i]);
21
\frac{22}{23}
         if(dfsn[node] == low[node]){
24
             while (1) {
25
                 int cur = st.top();
                 st.pop();
```

```
comp_id[cur] = cnt;
28
                 if(cur == node) break;
29
30
31
   }
32
   void dfs2(int node, int par) {
33
        kam[node] = 0;
35
        int mx = 0, second_mx = 0;
36
        for(auto i: bridge_tree[node]) {
37
            if(i == par) continue;
38
            dfs2(i, node);
            kam[node] = max(kam[node], 1 + kam[i]);
40
             if(kam[i] > mx) {
41
                 second_mx = mx;
42
                 mx = kam[i];
43
44
             else second_mx = max(second_mx, kam[i]);
45
46
        ans = max(ans, kam[node]);
47
        if(second mx) ans = max(ans, 2 + mx + second mx);
48
49
50 int main(){
51
        ios_base::sync_with_stdio(0);cin.tie(0);cout.tie(0);
        int n, m;
53
        cin >> n >> m;
54
        while (m--) {
55
            int u, v;
56
             cin >> u >> v;
57
            adj[u].push_back(v);
58
             adj[v].push_back(u);
59
60
        dfs(1, 0);
61
        for(int i = 1; i <= n; i++) {</pre>
62
            for(auto j: adj[i]){
63
                if(comp_id[i] != comp_id[j]){
64
                    bridge_tree[comp_id[i]].push_back(comp_id[j]);
65
66
67
68
        dfs2(1, 0);
        cout << ans;
70
71
        return 0:
72
```

6.4 Dinic With Scalling

```
1 ///O(ElqFlow) on Bipratite Graphs and O(EVlqFlow) on other graphs (I think)
    struct Dinic {
        #define vi vector<int>
        #define rep(i,a,b) f(i,a,b)
        struct Edge {
            int to, rev;
            11 c, oc;
 8
 9
            11 flow() { return max(oc - c, OLL); } // if you need flows
10
        vi lvl, ptr, q;
11
12
        vector<vector<Edge>> adj;
13
        Dinic(int n) : lvl(n), ptr(n), q(n), adj(n) {}
14
        void addEdge(int a, int b, ll c, int id, ll rcap = 0) {
15
            adj[a].push_back({b, sz(adj[b]), c, c, id});
            adj[b].push_back({a, sz(adj[a]) - 1, rcap, rcap,id});
16
17
18
        11 dfs(int v, int t, 11 f) {
            if (v == t || !f) return f;
19
20
            for (int& i = ptr[v]; i < sz(adj[v]); i++) {</pre>
21
                Edge& e = adj[v][i];
                if (lvl[e.to] == lvl[v] + 1)
                     if (ll p = dfs(e.to, t, min(f, e.c))) {
```

```
e.c -= p, adj[e.to][e.rev].c += p;
25
                         return p;
26
27
28
             return 0:
29
30
        11 calc(int s, int t) {
31
             11 \text{ flow} = 0; q[0] = s;
             rep(L,0,31) do { // 'int L=30' maybe faster for random data
32
33
                 lvl = ptr = vi(sz(q));
34
                 int qi = 0, qe = lvl[s] = 1;
35
                 while (qi < qe && !lvl[t]) {
36
                     int v = q[qi++];
37
                     for (Edge e : adj[v])
38
                         if (!lvl[e.to] && e.c >> (30 - L))
39
                             q[qe++] = e.to, lvl[e.to] = lvl[v] + 1;
40
41
                 while (ll p = dfs(s, t, LLONG_MAX)) flow += p;
42
             } while (lvl[t]);
43
             return flow;
44
45
         bool leftOfMinCut(int a) { return lvl[a] != 0; }
46
    };
```

6.5 Gomory Hu

```
1 /**
     * Author: chilli, Takanori MAEHARA
     * Date: 2020-04-03
     * License: CC0
     * Source: https://github.com/spaghetti-source/algorithm/blob/master/graph/
          gomory_hu_tree.cc#L102
     * Description: Given a list of edges representing an undirected flow graph,
     * returns edges of the Gomory-Hu tree. The max flow between any pair of
     * vertices is given by minimum edge weight along the Gomory-Hu tree path.
     * Time: $0(V)$ Flow Computations
10
     * Status: Tested on CERC 2015 J, stress-tested
11
12
     * Details: The implementation used here is not actually the original
     * Gomory-Hu, but Gusfield's simplified version: "Very simple methods for all
     * pairs network flow analysis". PushRelabel is used here, but any flow
15
     * implementation that supports 'leftOfMinCut' also works.
16
17
    #pragma once
18
    #include "PushRelabel.h"
19
20
21
    typedef array<11, 3> Edge;
    vector<Edge> gomoryHu(int N, vector<Edge> ed) {
23
        vector<Edge> tree;
24
        vi par(N);
25
26
            PushRelabel D(N); // Dinic also works
27
            for (Edge t : ed) D.addEdge(t[0], t[1], t[2], t[2]);
28
            tree.push_back({i, par[i], D.calc(i, par[i])});
            rep(j,i+1,N)
30
                if (par[j] == par[i] && D.leftOfMinCut(j)) par[j] = i;
31
32
        return tree;
33
```

6.6 HopcraftKarp BPM

```
1  /**
2  * Author: Chen Xing
3  * Date: 2009-10-13
4  * License: CC0
5  * Source: N/A
6  * Description: Fast bipartite matching algorithm. Graph $g$ should be a list
```

```
* of neighbors of the left partition, and $btoa$ should be a vector full of
     * -1's of the same size as the right partition. Returns the size of
    * the matching. $btoa[i]$ will be the match for vertex $i$ on the right side,
    * or $-1$ if it's not matched.
11
    * Usage: vi btoa(m, -1); hopcroftKarp(g, btoa);
     * Time: O(\sqrt{V}E)
13
     * Status: stress-tested by MinimumVertexCover, and tested on oldkattis.
          adkbipmatch and SPOJ:MATCHING
15
    #pragma once
16
17
    bool dfs(int a, int L, vector<vi>& q, vi& btoa, vi& A, vi& B) {
18
        if (A[a] != L) return 0;
19
        A[a] = -1;
20
        for (int b : g[a]) if (B[b] == L + 1) {
21
            B[b] = 0;
            if (btoa[b] == -1 || dfs(btoa[b], L + 1, g, btoa, A, B))
23
                return btoa[b] = a, 1;
24
25
26
27
    int hopcroftKarp(vector<vi>& g, vi& btoa) {
        int res = 0;
30
        vi A(g.size()), B(btoa.size()), cur, next;
31
        for (;;) {
32
            fill(all(A), 0);
33
            fill(all(B), 0);
34
            /// Find the starting nodes for BFS (i.e. layer 0).
35
            cur.clear();
36
            for (int a : btoa) if (a != -1) A[a] = -1;
37
            rep(a,0,sz(g)) if(A[a] == 0) cur.push_back(a);
38
            /// Find all layers using bfs.
39
            for (int lay = 1;; lay++) {
40
                bool islast = 0;
41
                next.clear();
42
                for (int a : cur) for (int b : g[a]) {
43
                    if (btoa[b] == -1) {
44
                        B[b] = lay;
45
                        islast = 1;
46
47
                    else if (btoa[b] != a && !B[b]) {
48
                        B[b] = lay;
49
                        next.push_back(btoa[b]);
50
51
52
                if (islast) break;
53
                if (next.empty()) return res;
54
                for (int a : next) A[a] = lay;
55
                cur.swap(next);
56
57
            /// Use DFS to scan for augmenting paths.
58
            rep(a, 0, sz(g))
59
                res += dfs(a, 0, g, btoa, A, B);
61
```

6.7 Hungarian

```
1  /*
2     Notes:
3     note that n must be <= m
4     so in case in your problem n >= m, just swap
5     also note this
6     void set(int x, int y, ll v) {a[x+1][y+1]=v;}
7     the algorithim assumes you're using 0-index
8     but it's using 1-based
9     */
10     struct Hungarian {
11         const ll INF = 10000000000000000; ///10^18
12         int n,m;
13     vector<vector<ll> > a;
```

```
14
         vector<11> u, v; vector<int> p, way;
15
         Hungarian(int n, int m):
16
         n(n), m(m), a(n+1, vector<11>(m+1, INF-1)), u(n+1), v(m+1), p(m+1), way (m+1) {}
17
         void set(int x, int y, ll v) {a[x+1][y+1]=v;}
18
         ll assign(){
19
             for(int i = 1; i <= n; i++) {
20
                 int j0=0;p[0]=i;
21
                  vector<11> minv(m+1, INF);
22
                  vector<char> used(m+1, false);
23
                 do {
24
                      used[j0]=true;
25
                      int i0=p[j0], j1; l1 delta=INF;
26
                      for(int j = 1; j <= m; j++)if(!used[j]){</pre>
                          11 cur=a[i0][j]-u[i0]-v[j];
27
28
                          if (cur<minv[j])minv[j]=cur, way[j]=j0;</pre>
29
                          if (minv[j] < delta) delta = minv[j], j1 = j;</pre>
30
31
                      for (int j = 0; j \le m; j++)
                          if(used[j])u[p[j]]+=delta,v[j]-=delta;
33
                          else minv[j]-=delta;
34
                      i0=i1;
35
                  } while(p[j0]);
36
                  do {
37
                      int j1=way[j0];p[j0]=p[j1];j0=j1;
38
                 } while(j0);
39
40
             return -v[0];
41
42
         vector<int> restoreAnswer() { ///run it after assign
43
             vector<int> ans (n+1);
44
             for (int j=1; j<=m; ++j)</pre>
45
                 ans[p[j]] = j;
46
             return ans;
47
48
   };
```

6.8 Kosaraju

```
g: Adjacency List of the original graph
      rg : Reversed Adjacency List
      vis : A bitset to mark visited nodes
      adj : Adjacency List of the super graph
      stk : holds dfs ordered elements
      cmp[i] : holds the component of node i
      go[i] : holds the nodes inside the strongly connected component i
 Q
10
    #define FOR(i,a,b) for(int i = a; i < b; i++)
12
    #define pb push_back
13
14
    const int N = 1e5+5;
15
16
    vector<vector<int>>q, rq;
17
    vector<vector<int>>go;
18
    bitset<N>vis;
19
   vector<vector<int>>adj;
20
   stack<int>stk;
21
    int n, m, cmp[N];
    void add_edge(int u, int v){
      g[u] push_back(v);
23
24
      rg[v].push_back(u);
25
    void dfs(int u) {
27
      vis[u]=1;
28
      for(auto v : g[u])if(!vis[v])dfs(v);
29
      stk.push(u);
30
31
    void rdfs(int u,int c) {
32
      vis[u] = 1;
33
      cmp[u] = c;
      go[c].push_back(u);
```

```
for(auto v : rg[u])if(!vis[v])rdfs(v,c);
36
37
38
      vis.reset();
39
      for(int i = 0; i < n; i++)if(!vis[i])</pre>
40
41
      vis.reset();
42
      int c = 0:
43
      while(stk.size()){
44
        auto cur = stk.top();
45
        stk.pop();
46
        if(!vis[cur])
47
          rdfs(cur,c++);
48
49
50
      return c;
51
```

6.9 Krichoff

```
2
         Count number of spanning trees in a graph
 3
 4
    int power(long long n, long long k) {
      int ans = 1;
      while (k) {
        if (k \& 1) ans = (long long) ans * n % mod;
 8
        n = (long long) n * n % mod;
 9
        k >>= 1;
10
11
      return ans;
12
13 int det(vector<vector<int>> a) {
14
      int n = a.size(), m = (int)a[0].size();
15
      int free_var = 0;
      const long long MODSQ = (long long) mod * mod;
16
17
      int det = 1, rank = 0;
18
       for (int col = 0, row = 0; col < m && row < n; col++) {</pre>
        int mx = row;
19
20
        for (int k = row; k < n; k++) if (a[k][col] > a[mx][col]) mx = k;
21
         if (a[mx][col] == 0) {
22
          det = 0;
23
           continue;
24
25
         for (int j = col; j < m; j++) swap(a[mx][j], a[row][j]);</pre>
         if (row != mx) det = det == 0 ? 0 : mod - det;
27
         det = 1LL * det * a[row][col] % mod;
28
         int inv = power(a[row][col], mod - 2);
         for (int i = 0; i < n && inv; i++) {</pre>
29
30
          if (i != row && a[i][col]) {
31
             int x = ((long long)a[i][col] * inv) % mod;
32
             for (int j = col; j < m && x; j++) {
33
                \textbf{if} \ (a[row][j]) \ a[i][j] = (MODSQ + a[i][j] - ((\textbf{long long})a[row][j] \ \star \ x) 
34
35
36
37
        row++:
38
         ++rank;
39
40
      return det;
41
```

6.10 Manhattan MST

```
#include<bits/stdc++.h>
using namespace std;

const int N = 2e5 + 9;
```

```
int n;
         vector<pair<int, int>> g[N];
        struct PT {
            int x, y, id;
            bool operator < (const PT &p) const {</pre>
11
                return x == p.x ? y < p.y : x < p.x;
12
13
        struct node {
14
15
            int val, id;
        } t[N];
16
17
        struct DSU {
19
             void init(int n) { for (int i = 1; i <= n; i++) p[i] = i; }</pre>
             int find(int u) { return p[u] == u ? u : p[u] = find(p[u]); }
21
             void merge(int u, int v) { p[find(u)] = find(v); }
         struct edge {
24
            int u, v, w;
            bool operator < (const edge &p) const { return w < p.w; }</pre>
26
        vector<edge> edges;
        int query(int x) {
            int r = 2e9 + 10, id = -1;
            for (; x \le n; x += (x \& -x)) if (t[x].val < r) r = t[x].val, id = t[x].id;
31
32
         void modify(int x, int w, int id) {
34
            for (; x > 0; x -= (x \& -x)) if (t[x].val > w) t[x].val = w, t[x].id = id;
35
36
         int dist(PT &a, PT &b) {
37
            return abs(a.x - b.x) + abs(a.y - b.y);
38
39
        void add(int u, int v, int w) {
40
            edges.push_back({u, v, w});
41
       long long Kruskal() {
44
            sort(edges.begin(), edges.end());
45
             long long ans = 0;
             for (edge e : edges) {
                int u = e.u, v = e.v, w = e.w;
48
                 if (dsu.find(u) != dsu.find(v)) {
49
                     ans += w;
50
                     g[u].push_back({v, w});
51
                      //g[v].push_back({u, w});
                     dsu.merge(u, v);
53
54
55
             return ans;
56
        void Manhattan() {
58
             for (int i = 1; i <= n; ++i) p[i].id = i;
59
            for (int dir = 1; dir <= 4; ++dir) {</pre>
60
                 if (dir == 2 || dir == 4) {
61
                     for (int i = 1; i <= n; ++i) swap(p[i].x, p[i].y);</pre>
62
                  else if (dir == 3) {
64
                     for (int i = 1; i \le n; ++i) p[i].x = -p[i].x;
65
66
                  sort(p + 1, p + 1 + n);
67
                 vector<int> v;
                 static int a[N];
                  for (int i = 1; i <= n; ++i) a[i] = p[i].y - p[i].x, v.push_back(a[i]);</pre>
70
                  sort(v.begin(), v.end());
 71
                  v.erase(unique(v.begin(), v.end()), v.end());
72
                  for (int i = 1; i \le n; i \ge n; i
                           v.begin() + 1;
                  for (int i = 1; i <= n; ++i) t[i].val = 2e9 + 10, t[i].id = -1;</pre>
74
                  for (int i = n; i >= 1; --i) {
 75
                     int pos = query(a[i]);
76
                     if (pos != -1) add(p[i].id, p[pos].id, dist(p[i], p[pos]));
77
                     modify(a[i], p[i].x + p[i].y, i);
```

```
78     }
79     }
80   }
81     int32_t main() {
82         ios_base::sync_with_stdio(0);
83         cin.tie(0);
84         cin >> n;
85         for (int i = 1; i <= n; i++) cin >> p[i].x >> p[i].y;
86         Manhattan();
87         cout << Kruskal() << '\n';
88         for (int u = 1; u <= n; u++) {
89              for (auto x: g[u]) cout << u - 1 << ' ' << x.first - 1 << '\n';
90         }
91         return 0;
92     }
</pre>
```

6.11 Maximum Clique

```
///Complexity O(3 ^ (N/3)) i.e works for 50
    ///you can change it to maximum independent set by flipping the edges 0->1, 1->0
    ///if you want to extract the nodes they are 1-bits in R
    int g[60][60];
     int res;
     long long edges[60];
     \textbf{void} \ \texttt{BronKerbosch}(\textbf{int} \ \texttt{n}, \ \textbf{long} \ \textbf{long} \ \texttt{R}, \ \textbf{long} \ \textbf{long} \ \texttt{P}, \ \textbf{long} \ \textbf{long} \ \texttt{X}) \ \ \{
       if (P == 0LL && X == 0LL) { //here we will find all possible maximal cliques (
            not maximum) i.e. there is no node which can be included in this set
         int t = __builtin_popcount11(R);
10
         res = max(res, t);
11
         return;
13
       int u = 0;
14
       while (!((1LL << u) & (P | X))) u ++;
15
       for (int v = 0; v < n; v++) {
16
         if (((1LL << v) & P) && !((1LL << v) & edges[u])) {</pre>
17
            BronKerbosch(n, R | (1LL << v), P & edges[v], X & edges[v]);</pre>
18
            P -= (1LL << v);
19
            X \mid = (1LL << v);
20
21
      }
22
    int max_clique (int n) {
       for (int i = 1; i <= n; i++) {
26
         edges[i - 1] = 0;
         for (int j = 1; j \le n; j++) if (g[i][j]) edges[i-1] = (1LL \le (j-1)
       BronKerbosch (n, 0, (1LL \ll n) - 1, 0);
30
       return res;
```

6.12 MCMF

```
struct MCMF {
16
17
        const int inf = 1000000010;
18
19
        vector<vector<Edge>> g;
20
21
        MCMF(int _n) {
            n = _n + 1;
23
             g.resize(n);
24
25
26
        void addEdge(int u, int v, int cap, int cost) {
             Edge e1 = \{v, cost, cap, 0, (int) g[v].size()\};
28
             Edge e2 = \{u, -\cos t, 0, 0, (int) g[u].size()\};
29
             g[u].push_back(e1);
30
             q[v].push back(e2);
31
32
33
        pair<int, int> minCostMaxFlow(int s, int t) {
34
             int flow = 0;
35
             int cost = 0;
36
             vector<int> state(n), from(n), from_edge(n);
37
            vector<int> d(n);
38
             deque<int> q;
39
             while (true) {
40
                 for (int i = 0; i < n; i++)
41
                    state[i] = 2, d[i] = inf, from[i] = -1;
42
                 state[s] = 1;
43
                q.clear();
44
                q.push_back(s);
45
                d[s] = 0;
46
                 while (!q.empty()) {
                     int v = q.front();
48
                     q.pop_front();
49
                     state[v] = 0;
50
                     for (int i = 0; i < (int) g[v].size(); i++) {
51
                         Edge e = g[v][i];
52
                         if (e.flow >= e.cap \mid \mid (d[e.to] \leq d[v] + e.cost))
53
                            continue;
54
                         int to = e.to;
55
                         d[to] = d[v] + e.cost;
56
                         from[to] = v;
57
                         from_edge[to] = i;
58
                         if (state[to] == 1) continue;
                         if (!state[to] || (!q.empty() && d[q.front()] > d[to]))
60
                             q.push_front(to);
61
                         else q.push_back(to);
62
                         state[to] = 1;
63
                     }
64
65
                if (d[t] == inf) break;
66
                int it = t, addflow = inf;
67
                 while (it != s) {
68
                     addflow = min(addflow,
                                   g[from[it]][from_edge[it]].cap
70
                                    - g[from[it]][from_edge[it]].flow);
71
                     it = from[it];
73
                it = t;
74
                 while (it != s) {
75
                    g[from[it]][from_edge[it]].flow += addflow;
76
                     g[it][g[from[it]][from_edge[it]].backEdge].flow -= addflow;
77
                     cost += g[from[it]][from_edge[it]].cost * addflow;
78
                     it = from[it];
79
                 flow += addflow;
81
82
             return {cost, flow};
83
   };
```

6.13 Minimum Arbroscene in a Graph

```
1 const int maxn = 2510, maxm = 7000000;
   const 11 maxint = 0x3f3f3f3f3f3f3f3f3fLL;
    int n, ec, ID[maxn], pre[maxn], vis[maxn];
   11 in[maxn];
    struct edge_t {
        int u, v;
 Q
        11 w;
10
    } edge[maxm];
11
    void add(int u, int v, ll w) {
        edge[++ec].u = u, edge[ec].v = v, edge[ec].w = w;
13
14
15
   11 arborescence(int n, int root) {
16
        11 \text{ res} = 0, \text{ index};
17
        while (true) {
18
             for (int i = 1; i \le n; ++i) {
19
                 in[i] = maxint, vis[i] = -1, ID[i] = -1;
20
21
            for (int i = 1; i <= ec; ++i) {</pre>
22
                int u = edge[i].u, v = edge[i].v;
                 if (u == v || in[v] <= edge[i].w) continue;</pre>
24
                 in[v] = edge[i].w, pre[v] = u;
25
26
            pre[root] = root, in[root] = 0;
27
            for (int i = 1; i <= n; ++i) {</pre>
28
                res += in[i];
29
                 if (in[i] == maxint) return -1;
30
31
            index = 0;
32
            for (int i = 1; i \le n; ++i) {
33
                if (vis[i] != -1) continue;
                 int u = i, v;
35
                 while (vis[u] == -1) {
36
                    vis[u] = i;
37
                     u = pre[u];
38
39
                if (vis[u] != i || u == root) continue;
40
                 for (v = u, u = pre[u], ++index; u != v; u = pre[u]) ID[u] = index;
41
                 ID[v] = index;
42
43
            if (index == 0) return res;
44
            for (int i = 1; i <= n; ++i) if (ID[i] == -1) ID[i] = ++index;</pre>
            for (int i = 1; i <= ec; ++i) {
45
46
                 int u = edge[i].u, v = edge[i].v;
47
                 edge[i].u = ID[u], edge[i].v = ID[v];
48
                 edge[i].w -= in[v];
49
50
            n = index, root = ID[root];
51
52
        return res;
53 }
```

6.14 Minmimum Vertex Cover (Bipartite)

```
int myrandom (int i) { return std::rand()%i;}

struct MinimumVertexCover {
    int n, id;
    vector<vector<int> > g;
    vector<int> color, m, seen;
    vector<int> comp[2];
    MinimumVertexCover() {}
    MinimumVertexCover(int n, vector<vector<int> > g) {
    this->n = n;
    this->g = g;
    color = m = vector<int>(n, -1);
```

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73 74 75

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77 78

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86

87

seen = vector<int>(n, 0);

```
makeBipartite();
void dfsBipartite(int node, int col) {
    if (color[node] != -1) {
        assert(color[node] == col); /* MSH BIPARTITE YA BASHMOHANDES */
    color[node] = col;
    comp[col].push back(node);
    for (int i = 0; i < int(g[node].size()); i++)</pre>
        dfsBipartite(g[node][i], 1 - col);
void makeBipartite() {
    for (int i = 0; i < n; i++)
        if (color[i] == -1)
            dfsBipartite(i, 0);
// match a node
bool dfs(int node) {
  random_shuffle(g[node].begin(),g[node].end());
    for (int i = 0; i < g[node].size(); i++) {</pre>
        int child = g[node][i];
        if (m[child] == -1) {
            m[node] = child;
            m[child] = node;
            return true;
        if (seen[child] == id)
            continue:
        seen[child] = id;
        int enemy = m[child];
        m[node] = child;
        m[child] = node;
        m[enemy] = -1;
        if (dfs(enemy))
            return true;
        m[node] = -1;
        m[child] = enemy;
        m[enemy] = child;
    return false:
void makeMatching() {
for (int j = 0; j < 5; j++)
  random_shuffle(comp[0].begin(),comp[0].end(),myrandom );
    for (int i = 0; i < int(comp[0].size()); i++) {</pre>
        id++:
        if(m[comp[0][i]] == -1)
            dfs(comp[0][i]);
}
void recurse(int node, int x, vector<int> &minCover, vector<int> &done) {
    if (m[node] != -1)
        return;
    if (done[node])return;
    done[node] = 1;
    for (int i = 0; i < int(g[node].size()); i++) {</pre>
        int child = g[node][i];
        int newnode = m[child];
        if (done[child]) continue;
        if(newnode == -1) {
            continue;
        done[child] = 2;
        minCover.push_back(child);
        m[newnode] = -1;
        recurse(newnode, x, minCover, done);
```

```
89
90
91
         vector<int> getAnswer() {
92
             vector<int> minCover, maxIndep;
93
              vector<int> done(n, 0);
94
             makeMatching();
95
             for (int x = 0; x < 2; x++)
96
                  for (int i = 0; i < int(comp[x].size()); i++) {</pre>
97
                      int node = comp[x][i];
98
                      if (m[node] == -1)
99
                          recurse (node, x, minCover, done);
100
101
102
              for (int i = 0; i < int(comp[0].size()); i++)</pre>
103
                  if (!done[comp[0][i]]) {
104
                      minCover.push_back(comp[0][i]);
105
106
              return minCover;
107
108
    };
```

6.15 Prufer Code

```
#include<bits/stdc++.h>
 1
    using namespace std;
    const int N = 3e5 + 9;
 6
    prufer code is a sequence of length n-2 to uniquely determine a labeled tree
         with n vertices
    Each time take the leaf with the lowest number and add the node number the leaf
         is connected to
    the sequence and remove the leaf. Then break the algo after n-2 iterations
11
   //0-indexed
12 int n;
    vector<int> g[N];
14
    int parent[N], degree[N];
15
16
  void dfs (int v) {
      for (size_t i = 0; i < q[v].size(); ++i) {</pre>
17
        int to = g[v][i];
18
19
        if (to != parent[v]) {
20
          parent[to] = v;
21
          dfs (to);
22
23
      }
^{24}
    vector<int> prufer_code() {
27
      parent[n-1] = -1;
      dfs (n - 1);
29
      int ptr = -1;
      for (int i = 0; i < n; ++i) {
31
        degree[i] = (int) g[i].size();
32
        if (degree[i] == 1 && ptr == -1) ptr = i;
33
34
      vector<int> result;
35
      int leaf = ptr;
      for (int iter = 0; iter < n - 2; ++iter) {</pre>
37
        int next = parent[leaf];
38
        result.push_back (next);
39
        --degree[next];
40
        if (degree[next] == 1 && next < ptr) leaf = next;</pre>
41
42
43
          while (ptr < n && degree[ptr] != 1) ++ptr;</pre>
44
          leaf = ptr;
45
46
      }
```

```
22
```

```
47
      return result;
48
49
    vector < pair<int, int> > prufer_to_tree(const vector<int> & prufer_code) {
      int n = (int) prufer_code.size() + 2;
51
      vector<int> degree (n, 1);
52
      for (int i = 0; i < n - 2; ++i) ++degree[prufer_code[i]];</pre>
53
54
55
      while (ptr < n && degree[ptr] != 1) ++ptr;</pre>
56
      int leaf = ptr;
57
      vector < pair<int, int> > result;
58
      for (int i = 0; i < n - 2; ++i) {
59
        int v = prufer_code[i];
60
        result.push_back (make_pair (leaf, v));
61
        --degree[leaf];
62
        if (--degree[v] == 1 && v < ptr) leaf = v;</pre>
63
64
          ++ptr;
65
           while (ptr < n && degree[ptr] != 1) ++ptr;</pre>
66
           leaf = ptr;
67
68
69
      for (int v = 0; v < n - 1; ++v) if (degree[v] == 1) result.push_back (
           make_pair (v, n - 1));
70
      return result:
71
72
73
    int32_t main() {
74
75
      return 0;
76
```

6.16 Push Relabel Max Flow

```
struct edge
2
3
        int from, to, cap, flow, index;
4
        edge(int from, int to, int cap, int flow, int index):
5
            from(from), to(to), cap(cap), flow(flow), index(index) {}
6
    struct PushRelabel
9
10
        int n;
11
        vector<vector<edge> > q;
12
        vector<long long> excess;
13
        vector<int> height, active, count;
14
        queue<int> Q;
15
16
        PushRelabel(int n):
17
            n(n), g(n), excess(n), height(n), active(n), count(2*n) {}
18
19
        void addEdge(int from, int to, int cap)
20
21
             g[from].push_back(edge(from, to, cap, 0, g[to].size()));
22
            if(from==to)
23
                g[from].back().index++;
24
            g[to].push_back(edge(to, from, 0, 0, g[from].size()-1));
25
26
27
        void enqueue(int v)
28
29
            if(!active[v] && excess[v] > 0)
30
31
                 active[v]=true;
32
                 Q.push(v);
33
34
35
36
        void push (edge &e)
37
38
            int amt=(int)min(excess[e.from], (long long)e.cap - e.flow);
```

```
return:
    e.flow += amt;
    g[e.to][e.index].flow -= amt;
    excess[e.to] += amt;
    excess[e.from] -= amt;
    enqueue (e.to);
void relabel(int v)
    count[height[v]]--;
    int d=2*n;
    for(auto &it:g[v])
        if(it.cap-it.flow>0)
            d=min(d, height[it.to]+1);
    height[v]=d;
    count[height[v]]++;
    enqueue (v);
void gap(int k)
    for (int v=0; v<n; v++)
        if(height[v]<k)</pre>
            continue;
        count[height[v]]--;
        height[v]=max(height[v], n+1);
        count[height[v]]++;
        enqueue (v);
void discharge(int v)
    for(int i=0; excess[v]>0 && i<g[v].size(); i++)</pre>
        push(g[v][i]);
    if(excess[v]>0)
        if(count[height[v]]==1)
            gap(height[v]);
        else
            relabel(v);
long long max_flow(int source, int dest)
    count[0] = n-1;
    count[n] = 1;
    height[source] = n;
    active[source] = active[dest] = 1;
    for(auto &it:g[source])
        excess[source] += it.cap;
        push(it);
    while(!Q.empty())
        int v=Q.front();
        Q.pop();
        active[v]=false;
        discharge(v);
    long long max_flow=0;
    for(auto &e:g[source])
        max_flow+=e.flow;
    return max_flow;
```

if(height[e.from] <=height[e.to] || amt==0)</pre>

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6.17 Tarjan Algo

114 };

```
vector< vector<int> > scc;
    vector<int> adj[N];
    int dfsn[N], low[N], cost[N], timer, in_stack[N];
    stack<int> st;
    // to detect all the components (cycles) in a directed graph
    void tarjan(int node) {
         dfsn[node] = low[node] = ++timer;
         in_stack[node] = 1;
10
         st.push (node);
11
         for(auto i: adj[node]){
12
            if(dfsn[i] == 0){
13
                 tarjan(i);
14
                 low[node] = min(low[node], low[i]);
15
16
             else if(in_stack[i]) low[node] = min(low[node], dfsn[i]);
17
18
         if(dfsn[node] == low[node]){
19
             scc.push_back(vector<int>());
20
             while(1){
21
                 int cur = st.top();
                 st.pop();
23
                 in_stack[cur] = 0;
24
                 scc.back().push_back(cur);
25
                 if(cur == node) break;
26
27
28
29
    int main(){
30
        int m;
31
         cin >> m;
32
         while (m--) {
33
             int u, v;
34
             cin >> u >> v;
35
             adj[u].push_back(v);
36
37
        for(int i = 1; i <= n; i++) {</pre>
38
             if(dfsn[i] == 0) {
39
                 tarjan(i);
40
41
42
43
         return 0;
```

6.18 Bipartite Matching

```
#include < iostream >
    #include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    struct graph
8
9
        int L, R;
10
        vector<vector<int> > adj;
11
        graph(int 1, int r) : L(1), R(r), adj(1+1) {}
12
        void add_edge(int u, int v)
13
14
            adj[u].push_back(v+L);
15
16
        int maximum_matching()
17
18
            vector<int> mate(L+R+1,-1), level(L+1);
```

```
19
             function<bool (void) > levelize = [&]()
20
21
                 queue<int> q;
22
                 for (int i=1; i<=L; i++)</pre>
23
24
                      level[i]=-1;
25
                     if(mate[i]<0)
26
                          q.push(i), level[i]=0;
27
28
                 while(!q.empty())
29
30
                      int node=q.front();
31
                      q.pop();
32
                      for(auto i : adj[node])
33
34
                          int v=mate[i];
35
                          if(v<0)
36
                              return true;
37
                          if(level[v]<0)
38
39
                              level[v]=level[node]+1;
40
                              q.push(v);
41
42
43
44
                 return false;
45
46
             function<bool (int)> augment =[&] (int node)
47
48
                 for(auto i : adj[node])
49
50
                      int v=mate[i];
51
                     if(v<0 || (level[v]>level[node] && augment(v)))
52
53
                          mate[node]=i;
54
                          mate[i]=node;
55
                          return true;
56
57
58
                 return false;
59
             };
60
             int match=0;
61
             while(levelize())
                 for(int i=1; i<=L; i++)</pre>
63
                      if(mate[i] < 0 && augment(i))</pre>
64
                          match++;
65
             return match;
66
67
    };
68
69
    int main()
70
71
72
        int L, R, m;
73
         cin>>L>>R>>m;
74
         graph g(L, R);
75
         for (int i = 0; i < m; ++i)
76
77
             int u, v;
78
             cin>>u>>v;
79
             g.add_edge(u, v);
80
81
         cout<<q.maximum_matching();</pre>
```

7 Math

7.1 Xor With Gauss

1 /*

```
၁
π
```

```
Some applications
        If you want to find the maximum in xor subset
        just ans = max(ans, ans ^p[i]) for all i
        if you want to count the number of subsets with a certain value
        check all different subsets of p
7
    11 p[66];
    bool add(ll x) {
10
        for (int i = 60; (~i) && x; --i) {
11
            if(x >> i & 1) {
12
                if(!p[i]) {
13
                    p[i] = x;
                    return true;
14
15
                    else {
16
                    x ^= p[i];
17
18
19
20
        return false;
21
```

7.2 Josephus

```
1 // n = total person
   // will kill every kth person, if k = 2, 2, 4, 6, ...
   // returns the mth killed person
   ll josephus(ll n, ll k, ll m) {
      m = n - m;
      if (k <= 1) return n - m;
      11 i = m:
      while (i < n) {
9
       11 r = (i - m + k - 2) / (k - 1);
10
       if ((i + r) > n) r = n - i;
11
        else if (!r) r = 1;
       i += r;
12
       m = (m + (r * k)) % i;
13
14
      } return m + 1;
15
```

7.3 Matrix Power/Multiplication

```
struct Matrix {
         const static int D = 100;
         int a[D][D];
6
        Matrix(int val) {
             for (int i = 0; i < D; i++)
                 for (int j = 0; j < D; j++)
9
                     a[i][j] = val;
10
11
        void clear() {
12
            memset(a, 0, sizeof a);
13
14
        void initIdentity() {
15
16
             for (int i = 0; i < D; i++)
                a[i][i] = 1;
17
18
19
         int * operator [](int r) {
20
             return a[r];
21
        const int * operator [](int r) const{
23
             return a[r];
\frac{24}{25}
26
         friend Matrix operator * (const Matrix & a, const Matrix & b) {
27
             Matrix ret(0);
             for (int k = 0; k < D; k++)
```

```
for(int i = 0; i < D; i++)if(a[i][k])</pre>
30
                     for (int j = 0; j < D; j++)
31
                         ret[i][j] = (ret[i][j] + 111 * a[i][k] * b[k][j]) % MOD;
32
             return ret;
33
34
35
    Matrix raiseMatrix (Matrix trans, 11 k) {
37
        Matrix res(0);
        res.initIdentity();
38
39
        for(;k;k>>=1,trans = trans * trans)
40
            if(k & 1)
                res = res * trans;
42
        return res;
43
```

7.4 Rabin Miller Primality check

```
// n < 4,759,123,141
                                          3: 2, 7, 61
   // n < 1,122,004,669,633
                                          4 : 2, 13, 23, 1662803
    // n < 3,474,749,660,383
                                          6 : pirmes <= 13
    // n < 3,825,123,056,546,413,051
                                         9 : primes <= 23
    int testPrimes[] = {2,3,5,7,11,13,17,19,23};
 Q
    struct MillerRabin{
     ///change K according to n
11
      const int K = 9;
12
      11 mult(11 s, 11 m, 11 mod) {
13
        if(!m) return 0;
14
        11 \text{ ret} = \text{mult}(s, m/2, mod);
15
        ret = (ret + ret) % mod;
16
        if (m & 1) ret = (ret + s) % mod;
17
        return ret;
18
19
20
      11 power(11 x, 11 p, 11 mod) {
21
        11 s = 1, m = x;
22
        while(p){
23
          if(p&1) s = mult(s, m, mod);
24
          p >>= 1;
25
          m = mult(m, m, mod);
26
27
        return s;
28
29
30
      bool witness(ll a, ll n, ll u, int t) {
31
        11 x = power(a, u, n), nx;
        for (int i = 0; i < t; i++) {
33
          nx = mult(x, x, n);
34
          if (nx == 1 \text{ and } x != 1 \text{ and } x != n-1) return 1;
35
          x = nx;
36
37
        return x != 1;
38
39
40
      bool isPrime(ll n) { // return 1 if prime, 0 otherwise
41
        if(n < 2) return 0;
42
        if(!(n&1)) return n == 2;
43
        for(int i = 0; i < K; i++)if(n == testPrimes[i])return 1;</pre>
44
        11 u = n-1; int t = 0;
45
        while (u&1) u >>= 1, t++; // n-1 = u*2^t
47
        for(int i = 0; i < K; i++) if(witness(testPrimes[i], n, u, t)) return 0;</pre>
49
        return 1;
50
51 }tester;
```

8.1 Aho-Corasick Mostafa

```
struct AC FSM {
    #define ALPHABET_SIZE 26
         struct Node {
5
            int child[ALPHABET_SIZE], failure = 0, match_parent = -1;
6
            vector<int> match;
                 for (int i = 0; i < ALPHABET_SIZE; ++i)child[i] = -1;</pre>
10
11
        };
12
13
        vector<Node> a:
14
15
        AC_FSM() {
16
            a.push_back(Node());
17
18
19
        void construct_automaton(vector<string> &words) {
20
            for (int w = 0, n = 0; w < words.size(); ++w, n = 0) {</pre>
\frac{1}{21}
                 for (int i = 0; i < words[w].size(); ++i) {</pre>
22
                    if (a[n].child[words[w][i] - 'a'] == -1) {
                         a[n].child[words[w][i] - 'a'] = a.size();
                         a.push_back(Node());
\frac{25}{26}
                     n = a[n].child[words[w][i] - 'a'];
                 a[n].match.push_back(w);
29
30
             queue<int> q;
31
             for (int k = 0; k < ALPHABET_SIZE; ++k) {</pre>
32
                if (a[0].child[k] == -1) a[0].child[k] = 0;
                 else if (a[0].child[k] > 0) {
33
34
                    a[a[0].child[k]].failure = 0;
35
                     q.push(a[0].child[k]);
             while (!q.empty()) {
                 int r = q.front();
40
                 q.pop();
                 for (int k = 0, arck; k < ALPHABET_SIZE; ++k) {</pre>
                    if ((arck = a[r].child[k]) != -1) {
                         q.push(arck);
                         int v = a[r].failure;
45
                         while (a[v].child[k] == -1) v = a[v].failure;
                         a[arck].failure = a[v].child[k];
                         a[arck].match_parent = a[v].child[k];
                         while (a[arck].match_parent != -1 &&
                                a[a[arck].match_parent].match.empty())
50
                             a[arck].match_parent =
51
                                     a[a[arck].match_parent].match_parent;
52
53
54
55
57
        void aho_corasick(string &sentence, vector<string> &words,
                           vector<vector<int> > &matches) {
59
             matches.assign(words.size(), vector<int>());
60
             int state = 0, ss = 0;
61
             for (int i = 0; i < sentence.length(); ++i, ss = state) {</pre>
                 while (a[ss].child[sentence[i] - 'a'] == -1)
62
63
                    ss = a[ss].failure;
                 state = a[state].child[sentence[i] - 'a'] = a[ss].child[sentence[i]
                 for (ss = state; ss != -1; ss = a[ss].match_parent)
66
                     for (int w: a[ss].match)
67
                         matches[w].push_back(i + 1 - words[w].length());
```

```
69
70 };
```

8.2 Aho-Corasick Anany

```
1 int trie[N][A];
   int go[N][A]; //holds the node that you will go to after failure and stuff
    11 ans[N]; //this node is a string terminator;
   int fail[N]; ///the failure function for each
    void BFS() {
        queue<int> q;
        f(i,0,A) {
            if(trie[0][i]) {
10
                q.push(trie[0][i]);
                fail[trie[0][i]] = 0;
11
12
13
            go[0][i] = trie[0][i];
14
15
        while(g.size()) {
17
            auto node = q.front();
18
            q.pop();
19
            ans[node] += ans[fail[node]]; ///propagate fail[i] to ans[i]
20
            for (int i = 0; i < A; i++) {</pre>
21
                if(trie[node][i]) { ///calculate failure for you child
                    int to = trie[node][i];
                    int cur = fail[node]; ///int g = pi[i-1]
                    while(cur && !trie[cur][i]) ///while(q && s[q] != s[i])
                        cur = fail[cur]; ///g = pi[g-1]
                    if(trie[cur][i])cur = trie[cur][i]; ///g += s[i] == s[g]
                    fail[to] = cur; ///pi[i] = g
                    q.push(to);
29
                    go[node][i] = trie[node][i];
30
                    else (
                    go[node][i] = go[fail[node]][i];
32
33
34
35
   void ins(string s, ll val) {
        int cur = 0;
38
        string sx = "";
39
        for(char c : s) {
            sx.push_back(c);
            if(!trie[cur][c - 'a']) {
                trie[cur][c - 'a'] = ++ptr;
            cur = trie[cur][c - 'a'];
45
        ans[cur] += val;
```

8.3 KMP Anany

```
1  vector<int> fail(string s) {
2    int n = s.size();
3    vector<int> pi(n);
4    for(int i = 1; i < n; i++) {
5       int g = pi[i-1];
6       while(g && s[i] != s[g])
7       g = pi[g-1];
8       g += s[i] == s[g];
9       pi[i] = g;
10    }
11    return pi;
12  }
13  vector<int> KMP(string s, string t) {
14    vector<int> pi = fail(t);
```

```
57
```

```
15
        vector<int> ret;
16
        for(int i = 0, g = 0; i < s.size(); i++) {
17
            while (g && s[i] != t[g])
                q = pi[q-1];
19
            g += s[i] == t[g];
20
            if(g == t.size()) { ///occurrence found
21
                ret.push_back(i-t.size()+1);
22
                g = pi[g-1];
23
24
25
        return ret:
26
```

8.4 Manacher Kactl

```
1 // If the size of palindrome centered at i is x, then dl[i] stores (x+1)/2.
 3
    vector<int> d1(n);
 4
    for (int i = 0, l = 0, r = -1; i < n; i++) {
         int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
         while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k]) {
 7
 8
 9
        d1[i] = k--;
10
        if (i + k > r) {
11
            1 = i - k;
12
            r = i + k;
13
14
15
17
    // If the size of palindrome centered at i is x, then d2[i] stores x/2
18
19
   vector<int> d2(n);
    for (int i = 0, l = 0, r = -1; i < n; i++) {
         int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1);
         while (0 \le i - k - 1 \&\& i + k \le n \&\& s[i - k - 1] == s[i + k]) {
22
23
24
25
        d2[i] = k--;
26
        if (i + k > r) {
27
            1 = i - k - 1;
28
            r = i + k;
29
30 }
```

8.5 Suffix Array Kactl

```
struct SuffixArray {
        using vi = vector<int>;
3
        #define rep(i,a,b) for(int i = a; i < b; i++)
            Note this code is considers also the empty suffix
-5
            so hear sa[0] = n and sa[1] is the smallest non empty suffix
            and sa[n] is the largest non empty suffix
            also LCP[i] = LCP(sa[i-1], sa[i]), meanining LCP[0] = LCP[1] = 0
            if you want to get LCP(i..j) you need to build a mapping between
10
            sa[i] and i, and build a min sparse table to calculate the minimum
11
            note that this minimum should consider sa[i+1...j] since you don't want
12
            to consider LCP(sa[i], sa[i-1])
13
14
            you should also print the suffix array and lcp at the beginning of the
                contest
15
            to clarify this stuff
16
17
        vi sa, lcp;
18
        SuffixArray(string& s, int lim=256) { // or basic_string<int>
19
            int n = sz(s) + 1, k = 0, a, b;
20
            vi x(all(s)+1), y(n), ws(max(n, lim)), rank(n);
```

```
sa = lcp = y, iota(all(sa), 0);
22
            for (int j = 0, p = 0; p < n; j = max(1, j * 2), lim = p) {
23
                p = j, iota(all(y), n - j);
24
                rep(i,0,n) if (sa[i] >= j) y[p++] = sa[i] - j;
25
                fill(all(ws), 0);
26
                rep(i,0,n) ws[x[i]]++;
27
                rep(i, 1, lim) ws[i] += ws[i - 1];
28
                for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
29
                swap(x, y), p = 1, x[sa[0]] = 0;
30
                rep(i,1,n) = sa[i-1], b = sa[i], x[b] =
31
                     (y[a] == y[b] && y[a + j] == y[b + j]) ? p - 1 : p++;
32
33
            rep(i,1,n) rank[sa[i]] = i;
            for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)</pre>
34
35
                for (k \&\& k--, j = sa[rank[i] - 1];
36
                        s[i + k] == s[j + k]; k++);
37
38
   };
```

8.6 Suffix Automaton Anany

```
///Note it's better to use addNode to clear a node before using it
    ///at the start of each test case use initAutomaton
    int last = 0, cntState = 1;
 4
    int nxt[N * 2][26];
    int len[N * 2], link[N * 2], firstPos[N * 2], cnt[N * 2];
    void addNode(int i) {
        memset(nxt[i], 0, sizeof nxt[i]);
        link[i] = -1;
11
        cnt[i] = 0;
12
13
14
    void initAutomaton() {
15
        cntState = 1:
16
        last = 0;
17
        addNode(last);
18
19
    int addChar(char c) {
21
        c -= 'a'; ///note this offset
23
        int p = last;
24
        int cur = cntState++;
        addNode(cur);
26
        cnt[cur] = 1; ///extra
27
        len[cur] = len[last] + 1;
28
        firstPos[cur] = len[cur] - 1; ///extra
29
        while(p != -1 && nxt[p][c] == 0) {
            nxt[p][c] = cur;
31
            p = link[p];
32
33
34
        if(p == -1) {
35
            link[cur] = 0;
36
           else {
37
            int q = nxt[p][c];
            if(len[q] == len[p] + 1) {
38
39
                link[cur] = q;
40
               else {
                int clone = cntState++;
42
                link[clone] = link[q];
43
                firstPos[clone] = firstPos[q]; ///extra
                len[clone] = len[p] + 1;
44
45
                link[q] = link[cur] = clone;
                memcpy(nxt[clone], nxt[q], sizeof nxt[q]);
47
                cnt[clone] = 0; ///extra
                f(i,0,26)nxt[clone][i] = nxt[q][i];
48
49
                while (p != -1 && nxt[p][c] == q) {
                    nxt[p][c] = clone;
                    p = link[p];
```

```
0
```

8.7 Suffix Automaton Mostafa

#include <bits/stdc++.h>

```
#define FIO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    typedef long long 11;
    typedef long double ld;
    const int N = 2e6 + 9, M = 5e5 + 9;
Q
    struct SA {
10
        struct node {
11
            int to[26];
12
            int link, len, co = 0;
13
14
15
                memset(to, 0, sizeof to);
16
                 co = 0, link = 0, len = 0;
17
18
        };
19
20
        int last, sz;
21
        vector<node> v;
22
23
\frac{24}{25}
             v = vector<node>(1);
             last = 0, sz = 1;
26
27
28
        void add_letter(int c) {
29
            int p = last;
30
             last = sz++;
31
             v.push_back({});
32
            v[last].len = v[p].len + 1;
33
             v[last].co = 1;
34
             for (; v[p].to[c] == 0; p = v[p].link)
35
                v[p].to[c] = last;
36
             if (v[p].to[c] == last) {
37
                v[last].link = 0;
38
                 return;
39
40
             int q = v[p].to[c];
41
             if (v[q].len == v[p].len + 1) {
                v[last].link = q;
43
                 return;
44
45
             int cl = sz++;
46
             v.push_back(v[q]);
47
             v.back().co = 0;
48
             v.back().len = v[p].len + 1;
49
            v[last].link = v[q].link = cl;
50
51
             for (; v[p].to[c] == q; p = v[p].link)
                v[p].to[c] = cl;
53
54
55
        void build_co() {
56
             priority_queue<pair<int, int>> q;
57
             for (int i = sz - 1; i > 0; i--)
58
                 q.push({v[i].len, i});
59
             while (q.size()) {
60
                 int i = q.top().second;
61
62
                 v[v[i].link].co += v[i].co;
```

```
64 }
65 };
66 67 int main() {
68 FIO
70 return 0;
71 }
```

8.8 Suffix Automaton With Rollback Mostafa

```
#include <bits/stdc++.h>
    #define FIO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    typedef long long 11;
    typedef long double ld;
    const int N = 2e6 + 9, M = 5e5 + 9;
    struct SA {
10
        struct node {
11
            int to[26];
12
            int link, len, co = 0;
13
14
            node() {
15
                memset(to, 0, sizeof to);
16
                co = 0, link = 0, len = 0;
17
18
        };
19
20
        struct LogNode {
21
            int last, sz;
22
            vector<pair<int, int>, int>> edges;
23
            pair<int, int> LinksUpdate = {0, 0};
24
        };
25
26
        int last, sz;
27
        vector<node> v;
28
        vector<LogNode> logs;
29
30
31
            v = vector<node>(1);
32
            last = 0, sz = 1;
33
34
35
        void add_letter(int c) {
            logs.push_back({});
37
            logs.back().last = last;
38
            logs.back().sz = sz;
39
40
            int p = last;
41
            last = sz++;
42
            v.push_back({});
43
            v[last].len = v[p].len + 1;
44
            v[last].co = 1;
45
            for (; v[p].to[c] == 0; p = v[p].link) {
46
                logs.back().edges.push_back({{p, c}, 0});
47
                v[p].to[c] = last;
48
49
            if (v[p].to[c] == last) {
50
                v[last].link = 0;
51
                return;
52
53
            int q = v[p].to[c];
54
            if (v[q].len == v[p].len + 1) {
55
                v[last].link = q;
56
                return;
57
            int cl = sz++;
59
            v.push_back(v[q]);
60
            v.back().co = 0;
            v.back().len = v[p].len + 1;
```

```
logs.back().LinksUpdate = {q, v[q].link};
63
             v[last].link = v[q].link = cl;
64
             for (; v[p].to[c] == q; p = v[p].link) {
65
                 logs.back().edges.push_back({{p, c}, q});
66
                 v[p].to[c] = c1;
67
68
69
        void rollback() {
70
             assert(logs.size());
71
             auto log = logs.back();
72
             while (v.size() > log.sz)
73
74
75
                v.pop_back();
             for (auto edge: log.edges)
                 v[edge.first.first].to[edge.first.second] = edge.second;
76
77
             if (log.LinksUpdate.first != 0)
                v[log.LinksUpdate.first].link = log.LinksUpdate.second;
78
             last = log.last;
79
             sz = log.sz:
80
             logs.pop_back();
81
82
    };
83
^{84}_{85}
    int main() {
        FIO
86
87
         return 0;
88
```

8.9 Zalgo Anany

```
int z[N], n;
    void Zalgo(string s) {
         int L = 0, R = 0;
         for(int i = 1; i < n; i++) {</pre>
4
             if(i \le R\&\&z[i-L] < R - i + 1)z[i] = z[i-L];
6
             else {
7
                 L = i;
                 R = max(R, i);
9
                 while (R < n \&\& s[R-L] == s[R])R++;
                 z[i] = R-L; --R;
11
12
```

9 Trees

9.1 Centroid Decomposition

```
1
2
        Properties:
            1. consider path(a,b) can be decomposed to path(a,lca(a,b)) and path(b,
                 1ca(a,b))
            where lca(a,b) is the lca on the centroid tree
5
            2. Each one of the n^2 paths is the concatenation of two paths in a set
                 of O(n lg(n))
            paths from a node to all its ancestors in the centroid decomposition.
            3. Ancestor of a node in the original tree is either an ancestor in the
                 CD tree or
            a descendadnt
    vector<int> adj[N]; ///adjacency list of original graph
    int sz[N];
    bool used[N];
14
    int centPar[N]; //parent in centroid
    void init(int node, int par) { ///initialize size
15
        sz[node] = 1;
```

```
for(auto p : adj[node])
18
            if(p != par && !used[p]) {
19
                init(p, node);
20
                sz[node] += sz[p];
21
22
    int centroid(int node, int par, int limit) {      ///get centroid
        for(int p : adj[node])
            if(!used[p] && p != par && sz[p] * 2 > limit)
26
            return centroid(p, node, limit);
27
        return node:
28
    int decompose(int node) {
30
        init(node, node);
                           ///calculate size
31
        int c = centroid(node, node, sz[node]); ///get centroid
        used[c] = true;
33
        for(auto p : adj[c])if(!used[p.F]) {      ///initialize parent for others and
             decompose
34
            centPar[decompose(p.F)] = c;
35
36
        return c;
37
    void update(int node, int distance, int col) {
        int centroid = node;
40
        while (centroid) {
41
            ///solve
42
            centroid = centPar[centroid];
43
44
    int query(int node) {
45
46
47
        int ans = 0;
48
49
        int centroid = node;
50
        while(centroid) {
51
            ///solve
52
            centroid = centPar[centroid];
53
55
        return ans;
56
```

9.2 Dsu On Trees

```
1 const int N = 1e5 + 9;
   vector<int> adj[N];
    int bigChild[N], sz[N];
    void dfs(int node, int par) {
        for(auto v : adj[node]) if(v != par){
            dfs(v, node);
            sz[node] += sz[v];
            if(!bigChild[node] || sz[v] > sz[bigChild[node]]) {
 9
                bigChild[node] = v;
10
11
    void add(int node, int par, int bigChild, int delta) {
14
15
        ///modify node to data structure
16
17
        for(auto v : adj[node])
        if(v != par && v != bigChild)
19
            add(v, node, bigChild, delta);
20
21
    void dfs2(int node, int par, bool keep) {
        for(auto v : adj[node])if(v != par && v != bigChild[node]) {
24
            dfs2(v, node, 0);
26
        if(bigChild[node]) {
            dfs2(bigChild[node], node, true);
```

```
ی
```

```
29     add(node, par, bigChild[node], 1);
30     ///process queries
31     if(!keep) {
32         add(node, par, -1, -1);
33     }
34 }
```

9.3 Heavy Light Decomposition (Along with Euler Tour)

```
1 /*
 2
         Notes:
 3
             1. 0-based
 4
             2. solve function iterates over segments and handles them seperatly
             if you're gonna use it make sure you know what you're doing
             3. to update/query segment in[node], out[node]
             4. to update/query chain in[nxt[node]], in[node]
             nxt[node]: is the head of the chain so to go to the next chain node =
                  par[nxt[node]]
 9
    int sz[mxN], nxt[mxN];
    int in[N], out[N], rin[N];
    vector<int> g[mxN];
    int par[mxN];
14
15
    void dfs_sz(int v = 0, int p = -1) {
16
        sz[v] = 1;
17
         par[v] = p;
18
         for (auto &u : g[v]) {
19
             if (u == p) {
20
                 swap(u, g[v].back());
21
22
             if(u == p) continue;
\begin{array}{c} 23 \\ 24 \end{array}
             dfs_sz(u,v);
             sz[v] += sz[u];
25
             if (sz[u] > sz[g[v][0]])
26
                 swap(u, g[v][0]);
27
\frac{28}{29}
         if(v != 0)
             g[v].pop_back();
30
31
32
    void dfs_hld(int v = 0) {
33
        in[v] = t++;
         rin[in[v]] = v;
34
35
         for (auto u : g[v]) {
36
             nxt[u] = (u == g[v][0] ? nxt[v] : u);
37
             dfs_hld(u);
38
39
         out[v] = t;
40
41
42 int n;
    bool isChild(int p, int u) {
44
      return in[p] <= in[u] && out[u] <= out[p];</pre>
45
46
    int solve(int u,int v) {
47
         vector<pair<int,int> > segu;
48
         vector<pair<int,int> > seqv;
49
         if(isChild(u,v)){
50
           while(nxt[u] != nxt[v]){
             segv.push_back(make_pair(in[nxt[v]], in[v]));
51
52
             v = par[nxt[v]];
53
54
           segv.push_back({in[u], in[v]});
55
         } else if(isChild(v,u)){
56
           while(nxt[u] != nxt[v]){
57
           segu.push_back(make_pair(in[nxt[u]], in[u]));
58
           u = par[nxt[u]];
59
60
           segu.push_back({in[v], in[u]});
61
       } else {
           while (u != v) {
```

```
if(nxt[u] == nxt[v]) {
64
              if(in[u] < in[v]) segv.push_back({in[u],in[v]}), R.push_back({u+1,v})
65
              else sequ.push_back({in[v],in[u]}), L.push_back({v+1,u+1});
66
              u = v:
67
              break;
68
            } else if(in[u] > in[v]) {
               segu.push_back({in[nxt[u]],in[u]}), L.push_back({nxt[u]+1, u+1});
70
            } else {
               segv.push_back({in[nxt[v]],in[v]}), R.push_back({nxt[v]+1, v+1});
73
               v = par[nxt[v]];
74
75
76
77
        reverse(seqv.begin(), seqv.end());
78
        int res = 0,state = 0;
79
        for(auto p : sequ) {
80
            qry(1,1,0,n-1,p.first,p.second,state,res);
81
82
        for(auto p : seqv) {
83
            qry(0,1,0,n-1,p.first,p.second,state,res);
85
        return res;
86
```

9.4 LCA

```
const int N = 1e5 + 5;
    const int LG = 18;
    vector<int> adj[N];
    int pa[N][LG], lvl[N];
    int in[N], out[N], timer;
    void dfs(int u, int p) {
      in[u] = ++timer;
      for (int k = 1; k < LG; k++)
        pa[u][k] = pa[pa[u][k-1]][k-1];
11
      for(auto v : adj[u])
12
        if(v != p) {
13
           lvl[v] = lvl[u] + 1;
14
           pa[v][0] = u;
15
           dfs(v, u);
16
17
      out[u] = timer;
18
19
    int LCA(int u, int v) {
      if(lvl[u] > lvl[v])
21
        swap(u,v);
       int d = lvl[v] - lvl[u];
       for (int k = 0; k < LG; k++)
        if(d >> k \& 1)
          v = pa[v][k];
26
       if (u == v) return u;
27
       for(int i = LG - 1; i >= 0; --i)
        if(pa[u][i] != pa[v][i]){
29
          u = pa[u][i];
30
          v = pa[v][i];
31
32
      return pa[u][0];
```

9.5 Mo on Trees

```
1 int BL[N << 1], ID[N << 1];
2 int lvl[N], par[17][N];
3 int ans[N];
4 vector<ii> adj[N];
5 struct query{
```

```
int id, 1, r, lc;
      bool operator < (const query & rhs) {</pre>
        return (BL[1] == BL[rhs.1]) ? (r < rhs.r) : (BL[1] < BL[rhs.1]);
10
    } Q [ N ] ;
11
    int in[N], out[N], val[N], timer;
12
    void dfs(int node, int p){
13
      in[node] = ++timer; ID[timer] = node;
14
      for(int i = 1; i < 17; i++)par[i][node] = par[i-1][par[i-1][node]];</pre>
15
      for(auto child : adj[node])if(child.F != p){
        lvl[child.F] = lvl[node] + 1;
16
17
        par[0][child.F] = node;
18
        val[child.F] = child.S;
19
        dfs(child.F, node);
20
21
      out[node] = ++timer; ID[timer] = node;
22
23
    int LCA(int u, int v) {
      if(lvl[u] > lvl[v])swap(u,v);
      for (int k = 0; k < 17; k++)
        if((lvl[v] - lvl[u]) >> k & 1)
27
          v = par[k][v];
28
      if(u == v)
29
       return u;
30
      for(int i = 16; i >= 0; --i)
31
       if(par[i][u] != par[i][v])
32
          u = par[i][u], v = par[i][v];
33
      return par[0][u];
34
35
    bool vis[N];
    int inSet[N];
    void add(int node, int & res){
38
      if(val[node] > N)return;
      if(!vis[node]){
40
        inSet[val[node]]++;
41
        while(inSet[res])res++;
42
43
        inSet[val[node]]--;
44
        if(!inSet[val[node]] && val[node] < res)</pre>
45
          res = val[node];
46
47
      vis[node] ^= 1;
48
49
            -----Adding Queries-----/
   f(i,0,q){
51
        int u, v;
52
        cin >> u >> v; if(lvl[u] > lvl[v]) swap(u, v);
53
        int lca = LCA(u, v);
54
        Q[i].id = i;
        Q[i].lc = lca;
        if(lca == u)Q[i].l = in[u], Q[i].r = in[v];
56
57
58
          Q[i].1 = out[u];
59
          Q[i].r = in[v];
60
61
62
                ------Processing Queries-----/
63
    f(i,0,q){
64
            while (curL < Q[i].1) add(ID[curL++], res);</pre>
65
            while (curL > Q[i].1) add(ID[--curL], res);
66
            while (curR < Q[i].r) add(ID[++curR], res);</pre>
67
            while (curR > Q[i].r) add(ID[curR--], res);
            int u = ID[Q[i].1];
69
            int v = ID[Q[i].r];
70
            if(Q[i].lc == u)add(Q[i].lc, res);
            ans[Q[i].id] = res;
72
            if(Q[i].lc == u)add(Q[i].lc, res);
```

10 Numerical

10.1 Lagrange Polynomial

```
class LagrangePoly {
    public:
        LagrangePoly(std::vector<long long> _a) {
            //f(i) = \_a[i]
            //interpola o vetor em um polinomio de grau y.size() - 1
            den.resize(y.size());
            int n = (int) y.size();
            for(int i = 0; i < n; i++) {
 9
10
                y[i] = (y[i] % MOD + MOD) % MOD;
                den[i] = ifat[n - i - 1] * ifat[i] % MOD;
11
                if((n-i-1) % 2 == 1) {
13
                    den[i] = (MOD - den[i]) % MOD;
14
15
16
17
18
        long long getVal(long long x) {
19
            int n = (int) y.size();
20
            x = (x % MOD + MOD) % MOD;
21
            if(x < n) {
                //return y[(int) x];
23
            std::vector<long long> 1, r;
25
            1.resize(n);
26
            1[0] = 1;
27
            for(int i = 1; i < n; i++) {</pre>
28
                1[i] = 1[i - 1] * (x - (i - 1) + MOD) % MOD;
29
30
            r.resize(n);
            r[n - 1] = 1;
31
            for (int i = n - 2; i >= 0; i--) {
33
                r[i] = r[i + 1] * (x - (i + 1) + MOD) % MOD;
34
35
            long long ans = 0;
36
            for (int i = 0; i < n; i++) {
37
                long long coef = l[i] * r[i] % MOD;
38
                ans = (ans + coef * y[i] % MOD * den[i]) % MOD;
39
40
            return ans:
41
    private:
        std::vector<long long> y, den;
45
```

11 formulas

11.1 Geometry

- Lattice points on line: gcd(dx,dy) + 1
- Pick's theorem: $A = I + \frac{B}{2} 1$
- sine rule: $\frac{A}{\sin(a)} = \frac{B}{\sin(b)} = \frac{C}{\sin(c)}$
- cosine rule: $C^2 = A^2 + B^2 2AB \times cos(c)$
- Dot product = $|A||B| \times cos(a)$
- Cross product = $|A||B| \times sin(a)$

- Rotation around axis: $R = (cos(a) \times Id + sin(a) \times crossU + (1 cos(a)) \times 11.4$ outerU)
- Angle of regular polygon = $\frac{180 \times (n-2)}{n}$
- # Diagonals of regular polygon = $\frac{n(n-3)}{n}$
- Triangulation of n-gon = Catalan (n-2)

11.2 Area

- triangle = $\frac{B \times H}{2}$
- triangle = $\sqrt{(S \times (S A) \times (S B) \times (S C))}$, S = PERIMETER/2
- triangle = $r \times S$, r = radius of inscribed circle
- circle = $R^2 \times \pi$
- ellipse = $\pi \times r_1 \times r_2$
- sector = $\frac{(r^2 \times a)}{2}$
- circular cap = $\frac{R^2 \times (a \sin(a))}{2}$
- trapzoid = $\frac{(B1+B2)}{2} \times H$
- prsim = perimeter(B)L + 2area(B)
- sphere = $4\pi r^2$

11.3 Volume

- Right circular cylinder = $\pi r^2 h$
- Pyramid = $\frac{Bh}{3}$
- Right circular cone = $\frac{\pi r^2 h}{3}$
- Sphere = $\frac{4}{3}\pi r^2 h$
- Sphere sector= $\frac{2}{3}\pi r^2 h = \frac{2}{3}\pi r^3 (1 \cos(a))$
- Sphere cap = $\frac{\pi h^2(3r-h)}{3}$

11.4 Combinatorics

- Cayley formula: number of forest with k trees where first k nodes belongs to different trees = kn^{n-k-1} . Multinomial theorem for trees of given degree sequence $\binom{n}{d_i}$
- Prufer sequence (M5da calls it parent array)
- K-Cyclic permutation = $\binom{n}{k} \times (k-1)!$
- Stirling numbers $S(n,k) = k \times S(n-1,k) + S(n,k-1)$ number of way to partition n in k sets.
- Bell number $B_n = \sum_{1}^{n} (n-1, k) B_k$
- Arithmetic-geometric-progression $S_n = \frac{A_1 \times G_1 A_{n+1} \times G_{n+1}}{1-r} + \frac{dr}{(1-r)^2} \times (G_1 G_{n+1})$

11.5 Graph Theory

- Graph realization problem: sorted decreasing degrees: $\sum_{1}^{k} d_i = k(k-1) + sum_(k+1)^n \min(d_i, k)$ (first k form clique and all other nodes are connected to them).
- Euler formula: v + f = e + c + 1
- # perfect matching in bipartite graph, DP[S][j] = DP[S][j-1] + DP[S/v][j-1] for all v connected to the j node.

11.6 Max flow with lower bound

- feasible flow in a network with both upper and lower capacity constraints, no source or sink: capacities are changed to upper bound lower bound. Add a new source and a sink. let M[v] = (sum of lower bounds of ingoing edges to v) (sum of lower bounds of outgoing edges from v). For all v, if $M[v]_{\dot{c}}0$ then add edge (S,v) with capacity M, otherwise add (v,T) with capacity -M. If all outgoing edges from S are full, then a feasible flow exists, it is the flow plus the original lower bounds.
- maximum flow in a network with both upper and lower capacity constraints, with source s and sink t: add edge (t,s) with capacity infinity. Binary search for the lower bound, check whether a feasible exists for a network WITH-OUT source or sink (B).

11.7 Sum of floor function

```
Algorithm:
    t = GCD(p, q)
    p = p/t
    q = q/t
    s = 0
    z = 1
    while (q > 0) and (n > 0)
        (point A)
    t = [p/q]
    s = s + ztn(n+1)/2
    p = p - qt
        (point B)
    t = [n/q]
    s = s + zp(n+1) - zt(pqt + p+q-1)/2
    n = n - qt
        (point C)
    t = [np/q]
```

$$s = s + ztn$$

 $n = t$
 $swap p and q$
 $z = -z$

11.8 Joseph problem

$$g(n,k) = \begin{cases} 0 & \text{if } n = 1\\ (g(n-1,k) + k) \bmod n & \text{if } 1 < n < k\\ \left\lfloor \frac{k((g(n',k) - n \bmod k) \bmod n')}{k-1} \right\rfloor \text{ where } n' = n - \left\lfloor \frac{n}{k} \right\rfloor & \text{if } k \le n \end{cases}$$