Faculty of Computer and Information Sciences, Ain Shams University: Too Wrong to Pass Too Correct to Fail

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1	Con	binatorics	1
	1.1	Burnside Lemma	1
	1.2	Catlan Numbers	1
2	Alge	bra	2
	2.1	Primitive Roots	2
	2.2	Discrete Logarithm	2
	2.3	Iteration over submasks	2
	2.4	Totient function	2
	2.5	CRT and EEGCD	2
	2.6	FFT	3
	2.7	Fibonacci	3
	2.8	Gauss Determinant	3
	2.9	GAUSS SLAE	4
	2.10	Matrix Inverse	4
	2.10	NTT	4
	2.11	NTT of KACTL	5
3		Structures	5
	3.1	2D BIT	5
	3.2	2D Sparse table	5
	3.3	hillbert Order	6
	3.4	Merge Sort Bit with updates	6
	3.5	Mo's	7
	3.6	Mo With Updates	7
	3.7	Ordered Set	8
	3.8	Persistent Seg Tree	8
	3.9	Sqrt Decomposition	8
	3.10	Treap	9
	3.11	Wavelet Tree	9
4	\mathbf{DP}		10
	4.1	Dynamic Convex Hull Trick	10
	4.2	Dynamic Connectivety with SegTree	10
	4.3		11
	4.4		12
5	Geo	netry	12
6	Gra	hs	12
7	Mat		12
	7.1	Xor With Gauss	12
	7.2	Josephus	12
	7.3	•	12
	7.4	,	13
_	~ ·		
8	Stri	·O-	13
	8.1		13
	8.2		14
	8.3		14
	8.4		14
	8.5		15
	8.6	Suffix Automaton Anany	15

8.8	Suffix Automaton With Rollback Mostafa	16
8.9	Zalgo Anany	16
Tre	es	17
9.1	Centroid Decomposition	17
9.2	Dsu On Trees	17
9.3	Heavy Light Decomposition (Along with Euler Tour)	17
9.4	LCA	18
9.5	Mo on Trees	18

1 Combinatorics

1.1 Burnside Lemma

```
1
2  // |Classes|=sum (k ^C(pi)) / |G|
3
4  // C(pi) the number of cycles in the permutation pi
5
6  // |G| the number of permutations
```

1.2 Catlan Numbers

```
1 const int MOD = \dots
   const int MAX = ....
 3 int catalan[MAX];
    void init() {
        catalan[0] = catalan[1] = 1;
        for (int i=2; i<=n; i++) {</pre>
            catalan[i] = 0;
            for (int j=0; j < i; j++) {
 9
                catalan[i] += (catalan[j] * catalan[i-j-1]) % MOD;
10
                if (catalan[i] >= MOD) {
11
                    catalan[i] -= MOD;
12
13
14
15
16
   // 1- Number of correct bracket sequence consisting of n opening and n closing
   // 2- The number of rooted full binary trees with n+1 leaves (vertices are not
         numbered).
         A rooted binary tree is full if every vertex has either two children or no
          children.
   // 3- The number of ways to completely parenthesize n+1 factors.
21 // 4- The number of triangulations of a convex polygon with n+2 sides
          (i.e. the number of partitions of polygon into disjoint triangles by using
          the diagonals).
   // 5- The number of ways to connect the 2n points on a circle to form n disjoint
          chords.
   // 6- The number of non-isomorphic full binary trees with n internal nodes (i.e.
          nodes having at least one son).
   // 7- The number of monotonic lattice paths from point (0,0) to point (n,n) in a
          square lattice of size nxn,
          which do not pass above the main diagonal (i.e. connecting (0,0) to (n,n))
   // 8- Number of permutations of length n that can be stack sorted
          (i.e. it can be shown that the rearrangement is stack sorted if and only
         there is no such index i<j<k, such that ak<ai<aj ).
   // 9- The number of non-crossing partitions of a set of n elements.
    // 10- The number of ways to cover the ladder 1..n using n rectangles
   // (The ladder consists of n columns, where ith column has a height i).
```

2 Algebra

2.1 Primitive Roots

```
int powmod (int a, int b, int p) {
2
         int res = 1;
3
         while (b)
            if (b & 1)
                res = int (res * 111 * a % p), --b;
                a = int (a * 111 * a % p), b >>= 1;
8
        return res;
9
10
    int generator (int p) {
11
12
        vector<int> fact;
13
         int phi = p - 1, n = phi;
14
         for (int i = 2; i * i <= n; ++i)</pre>
15
            if (n \% i == 0) {
16
                fact.push_back (i);
17
                 while (n \% i == 0)
18
                    n /= i;
19
20
        if (n > 1)
21
            fact.push_back (n);
22
23
         for (int res = 2; res <= p; ++res) {</pre>
24
             bool ok = true;
25
             for (size_t i = 0; i < fact.size() && ok; ++i)</pre>
26
                ok &= powmod (res, phi / fact[i], p) != 1;
27
             if (ok) return res;
28
29
         return -1;
30
```

2.2 Discrete Logarithm

```
// Returns minimum x for which a \hat{x} \times m = b \times m, a and m are coprime.
    int solve(int a, int b, int m) {
         a %= m, b %= m;
 4
         int n = sqrt(m) + 1;
 5
         int an = 1;
 7
         for (int i = 0; i < n; ++i)
             an = (an * 111 * a) % m;
 Q
10
        unordered_map<int, int> vals;
11
         for (int q = 0, cur = b; q \le n; ++q) {
12
             vals[cur] = q;
13
             cur = (cur * 111 * a) % m;
14
15
16
         for (int p = 1, cur = 1; p \le n; ++p) {
             cur = (cur * 111 * an) % m;
17
             if (vals.count(cur)) {
19
                 int ans = n * p - vals[cur];
20
                 return ans;
21
22
23
         return -1;
\frac{24}{25}
26
    //When a and m are not coprime
    // Returns minimum x for which a ^ x % m = b % m.
    int solve(int a, int b, int m) {
29
        a %= m, b %= m;
30
         int k = 1, add = 0, q;
31
        while ((g = gcd(a, m)) > 1) {
32
             if (b == k)
                 return add;
```

```
if (b % g)
35
                return -1;
36
            b /= g, m /= g, ++add;
37
            k = (k * 111 * a / g) % m;
38
39
40
        int n = sqrt(m) + 1;
41
        int an = 1;
42
        for (int i = 0; i < n; ++i)
43
            an = (an * 111 * a) % m;
44
45
        unordered_map<int, int> vals;
46
        for (int q = 0, cur = b; q \le n; ++q) {
47
            vals[cur] = q;
48
            cur = (cur * 111 * a) % m;
49
50
51
        for (int p = 1, cur = k; p \le n; ++p) {
52
            cur = (cur * 111 * an) % m;
53
            if (vals.count(cur)) {
54
                int ans = n * p - vals[cur] + add;
55
                return ans;
56
57
58
        return -1:
59
```

2.3 Iteration over submasks

```
1 int s = m;
2 while (s > 0) {
3     ... you can use s ...
4     s = (s-1) & m;
5 }
```

2.4 Totient function

```
void phi_1_to_n(int n) {
        vector<int> phi(n + 1);
 3
        phi[0] = 0;
        phi[1] = 1;
        for (int i = 2; i <= n; i++)</pre>
            phi[i] = i;
 8
        for (int i = 2; i <= n; i++) {</pre>
 9
            if (phi[i] == i) {
10
                 for (int j = i; j <= n; j += i)
11
                     phi[j] -= phi[j] / i;
12
13
14
```

2.5 CRT and EEGCD

```
1 11 extended(11 a, 11 b, 11 &x, 11 &y) {
 3
        if(b == 0) {
            x = 1;
 5
            y = 0;
 6
            return a;
        11 x0, y0;
 8
 9
        11 g = extended(b, a % b, x0, y0);
10
        x = y0;
11
        y = x0 - a / b * y0;
12
13
        return g ;
```

```
15 ll de(ll a, ll b, ll c, ll &x, ll &y) {
16
17
         11 g = \text{extended}(abs(a), abs(b), x, y);
18
        if(c % q) return -1;
19
20
        x \star = c / g;
21
        y *= c / g;
22
23
        if(a < 0)x = -x;
24
        if(b < 0)y = -y;
25
         return q;
26
27
    pair<11, 11> CRT(vector<11> r, vector<11> m) {
28
29
        11 r1 = r[0], m1 = m[0];
30
31
         for(int i = 1; i < r.size(); i++) {</pre>
32
33
             11 r2 = r[i], m2 = m[i];
34
             11 x0, y0;
35
             11 q = de(m1, -m2, r2 - r1, x0, y0);
36
37
             if(q == -1) return \{-1, -1\};
38
39
             11 \text{ nr} = x0 * m1 + r1;
40
             11 nm = m1 / g * m2;
42
             r1 = (nr % nm + nm) % nm;
43
             m1 = nm;
44
45
         return {r1, m1};
46
```

2.6 FFT

```
#include < iostream >
    #include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define rep(i, a, b) for(int i = a; i < (b); ++i)
    #define all(x) begin(x), end(x)
    #define sz(x) (int)(x).size()
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    typedef complex<double> C;
11
    typedef vector<double> vd;
    typedef vector<int> vi;
13
    typedef pair<int, int> pii;
    void fft(vector<C>& a) {
15
        int n = sz(a), L = 31 - __builtin_clz(n);
16
        static vector<complex<long double>> R(2, 1);
17
        static vector<C> rt(2, 1); // (^ 10% fas te r i f double)
18
        for (static int k = 2; k < n; k \neq 2) {
19
            R.resize(n):
20
21
            auto x = polar(1.0L, acos(-1.0L) / k);
22
            rep(i, k, 2 * k) rt[i] = R[i] = i & 1 ? R[i / 2] * x : R[i / 2];
23
24
25
        rep(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
        rep(i, 0, n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
26
27
        for (int k = 1; k < n; k *= 2)
28
            for (int i = 0; i < n; i += 2 * k) rep(j, 0, k) {
29
                Cz = rt[j + k] * a[i + j + k]; //
30
                a[i + j + k] = a[i + j] - z;
31
                a[i + j] += z;
32
33
34
    vd conv(const vd& a, const vd& b) {
35
        if (a.empty() || b.empty()) return {};
36
        vd res(sz(a) + sz(b) - 1);
        int L = 32 - __builtin_clz(sz(res)), n = 1 << L;</pre>
```

```
vector<C> in(n), out(n);
39
        copy(all(a), begin(in));
40
        rep(i, 0, sz(b)) in[i].imag(b[i]);
41
        fft(in);
42
        for (C& x : in) x *= x;
43
        rep(i, 0, n) out[i] = in[-i & (n - 1)] - conj(in[i]);
44
        fft (out);
45
        rep(i, 0, sz(res)) res[i] = imag(out[i]) / (4 * n);
46
        return res;
47
48
49
    int main() {
50
51
        //Applications
52
        //1-All possible sums
53
54
        //2-All possible scalar products
55
        // We are given two arrays a[] and b[] of length n.
56
        //We have to compute the products of a with every cyclic shift of b.
57
        //We generate two new arrays of size 2n: We reverse a and append n zeros to
58
        //And we just append b to itself. When we multiply these two arrays as
             polynomials,
59
        //and look at the coefficients c[n-1], c[n], ..., c[2n-2] of the product c,
             we aet:
60
        //c[k]=sum i+j=k a[i]b[j]
61
62
        //3-Two stripes
63
        //We are given two Boolean stripes (cyclic arrays of values 0 and 1) a and b
64
        //We want to find all ways to attach the first stripe to the second one,
65
        //such that at no position we have a 1 of the first stripe next to a 1 of
             the second stripe.
66
```

2.7 Fibonacci

2.8 Gauss Determinant

```
1 const double EPS = 1E-9;
    int n:
    vector < vector<double> > a (n, vector<double> (n));
    double det = 1;
    for (int i=0; i<n; ++i) {</pre>
        int k = i;
        for (int j=i+1; j<n; ++j)</pre>
 Q
             if (abs (a[j][i]) > abs (a[k][i]))
10
                 k = i;
11
         if (abs (a[k][i]) < EPS) {
12
             det = 0;
13
             break;
14
15
        swap (a[i], a[k]);
16
        if (i != k)
17
             det = -det;
18
        det *= a[i][i];
19
        for (int j=i+1; j<n; ++j)</pre>
           a[i][j] /= a[i][i];
         for (int j=0; j < n; ++j)
```

2.9 GAUSS SLAE

```
const double EPS = 1e-9;
2
    const int INF = 2; // it doesn't actually have to be infinity or a big number
3
    int gauss (vector < vector<double> > a, vector<double> & ans) {
4
         int n = (int) a.size();
5
         int m = (int) a[0].size() - 1;
6
7
8
         vector<int> where (m, -1);
9
         for (int col = 0, row = 0; col < m && row < n; ++col) {
10
             int sel = row;
11
             for (int i = row; i < n; ++i)
12
                 if (abs (a[i][col]) > abs (a[sel][col]))
13
                     sel = i;
14
             if (abs (a[sel][col]) < EPS)</pre>
15
                 continue;
16
             for (int i = col; i <= m; ++i)</pre>
17
                 swap (a[sel][i], a[row][i]);
18
             where[col] = row;
19
20
             for (int i = 0; i < n; ++i)
21
                 if (i != row) {
22
                      double c = a[i][col] / a[row][col];
23
                      for (int j = col; j <= m; ++j)</pre>
24
                          a[i][j] -= a[row][j] * c;
25
26
             ++row;
27
28
29
         ans.assign (m, 0);
30
         for (int i = 0; i < m; ++i)
31
             if (where[i] != -1)
32
                 ans[i] = a[where[i]][m] / a[where[i]][i];
33
         for (int i = 0; i < n; ++i) {
34
             double sum = 0;
35
             for (int j = 0; j < m; ++j)
    sum += ans[j] * a[i][j];</pre>
36
37
             if (abs (sum - a[i][m]) > EPS)
38
                 return 0;
39
40
41
         for (int i = 0; i < m; ++i)
             if (where [i] == -1)
43
                 return INF;
44
         return 1;
45
```

2.10 Matrix Inverse

```
// Sometimes, the questions are complicated - and the answers are simple. //
pragma GCC optimize ("03")
#pragma GCC optimize ("unroll-loops")
#include <bits/stdc++.h>
#define 11 long long
#define ld long double
#define Io ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
using namespace std;
vector < vector<double> > gauss (vector < vector<double> > a) {

int n = (int) a.size();
vector<vector<double> > ans(n, vector<double>(n, 0));
```

```
14
         for (int i = 0; i < n; i++)
15
             ans[i][i] = 1;
16
         for (int i = 0; i < n; i++) {
17
             for (int j = i + 1; j < n; j++)
18
                 if(a[j][i] > a[i][i]) {
19
                     swap(a[j], a[i]);
20
                     swap(ans[j], ans[i]);
21
22
             double val = a[i][i];
23
             for (int j = 0; j < n; j++) {
24
                 a[i][j] /= val;
25
                 ans[i][j] /= val;
27
             for (int j = 0; j < n; j++) {
                 if(j == i)continue;
28
29
                 val = a[i][i];
30
                 for(int k = 0; k < n; k++) {
    a[j][k] -= val * a[i][k];</pre>
                      ans[j][k] = val * ans[i][k];
33
34
35
36
        return ans;
37
38
    int main() {
39
40
41
         vector<vector<double> > v(3, vector<double> (3) );
42
         for (int i = 0; i < 3; i++)
43
             for (int j = 0; j < 3; j++)
44
                 cin >> v[i][j];
45
46
        for(auto i : gauss(v)) {
47
             for(auto j : i)
                 cout << j << " ";
48
49
             cout << "\n";
50
51
```

2.11 NTT

```
1
    struct NTT {
        int mod ;
 3
        int root ;
        int root 1 ;
        int root_pw ;
        NTT(int _mod, int primtive_root, int NTT_Len) {
 9
            mod = mod;
10
            root_pw = NTT_Len;
11
             root = fastpower(primtive_root, (mod - 1) / root_pw);
12
             root_1 = fastpower(root, mod - 2);
13
14
        void fft(vector<int> & a, bool invert) {
15
            int n = a.size();
16
17
             for (int i = 1, j = 0; i < n; i++) {
                int bit = n >> 1;
18
19
                for (; j & bit; bit >>= 1)
20
                    j ^= bit;
                j ^= bit;
21
22
23
                if (i < j)
                    swap(a[i], a[j]);
25
26
27
            for (int len = 2; len <= n; len <<= 1) {</pre>
                int wlen = invert ? root_1 : root;
                for (int i = len; i < root_pw; i <<= 1)</pre>
                     wlen = (int)(1LL * wlen * wlen % mod);
```

```
32
33
                 for (int i = 0; i < n; i += len) {</pre>
34
                     int w = 1;
35
                     for (int j = 0; j < len / 2; j++) {
36
                         int u = a[i + j], v = (int)(1LL * a[i + j + len / 2] * w %
                              mod);
                         a[i + j] = u + v < mod ? u + v : u + v - mod;
                         a[i + j + len / 2] = u - v >= 0 ? u - v : u - v + mod;
39
                         w = (int) (1LL * w * wlen % mod);
40
41
42
43
44
             if (invert) {
45
                 int n_1 = fastpower(n, mod - 2);
46
                 for (int & x : a)
47
                    x = (int) (1LL * x * n_1 % mod);
48
49
50
         vector<int> multiply(vector<int> &a, vector<int> &b) {
51
             vector<int> fa(a.begin(), a.end()), fb(b.begin(), b.end());
52
             int n = 1:
53
             while(n < a.size() + b.size())</pre>
54
                n <<= 1;
55
56
             fa.resize(n);
57
             fb.resize(n);
58
59
             fft(fa, 0);
60
             fft(fb, 0);
61
             for (int i = 0; i < n; i++)
62
63
                fa[i] = 1LL * fa[i] * fb[i] % mod;
64
             fft(fa, 1);
65
             return fa;
66
67 };
```

2.12 NTT of KACTL

```
1 ///(Note faster than the other NTT)
    ///If the mod changes don't forget to calculate the primitive root
    using 11 = long long;
    const 11 mod = (119 << 23) + 1, root = 3; // = 998244353</pre>
    // For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 << 21
    // and 483 << 21 (same root). The last two are > 10^9.
    typedef vector<ll> vl;
    11 modpow(ll b, ll e) {
10
        11 \text{ ans} = 1;
11
         for (; e; b = b * b % mod, e /= 2)
12
            if (e & 1) ans = ans * b % mod;
13
         return ans:
14
15
    void ntt(vl &a) {
16
        int n = sz(a), L = 31 - \underline{builtin_clz(n)};
17
         static vl rt(2, 1);
18
         for (static int k = 2, s = 2; k < n; k \neq 2, s++) {
19
             rt.resize(n);
20
             ll z[] = \{1, modpow(root, mod >> s)\};
21
             f(i,k,2*k) rt[i] = rt[i / 2] * z[i & 1] % mod;
22
23
         vector<int> rev(n);
24
         f(i,0,n) \text{ rev}[i] = (\text{rev}[i / 2] | (i \& 1) << L) / 2;
         f(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
26
         for (int k = 1; k < n; k *= 2)
27
             for (int i = 0; i < n; i += 2 * k) f(j, 0, k) {
                 11 z = rt[j + k] * a[i + j + k] % mod, &ai = a[i + j];
28
29
                 a[i + j + k] = ai - z + (z > ai ? mod : 0);
30
                 ai += (ai + z >= mod ? z - mod : z);
```

```
32
33
   vl conv(const vl &a, const vl &b) {
        if (a.empty() || b.empty()) return {};
        int s = sz(a) + sz(b) - 1, B = 32 - _builtin_clz(s), n = 1 << B;
35
        int inv = modpow(n, mod - 2);
37
        vl L(a), R(b), out(n);
38
        L.resize(n), R.resize(n);
39
        ntt(L), ntt(R);
40
        f(i,0,n) out [-i \& (n-1)] = (11)L[i] * R[i] % mod * inv % mod;
41
42
        return {out.begin(), out.begin() + s};
43
44
    vector<int> v;
   vector<ll> solve(int s, int e) {
        if(s==e) {
            vector<ll> res(2);
47
48
            res[0] = 1;
49
            res[1] = v[s];
50
            return res;
51
52
        int md = (s + e) >> 1;
53
        return conv(solve(s, md), solve(md+1, e));
```

3 Data Structures

3.1 2D BIT

```
1 void upd(int x, int y, int val) {
2    for(int i = x; i <= n; i += i & -i)
3    for(int j = y; j <= m; j += j & -j)
4    bit[i][j] += val;
5    }
6    int get(int x, int y) {
7       int ans = 0;
8       for(int i = x; i; i -= i & -i)
9       for(int j = y; j; j -= j & -j)
10       ans += bit[i][j];
11 }</pre>
```

3.2 2D Sparse table

```
1 /*
        note this isn't the best cache-wise version
 3
        query O(1), Build O(NMlgNlgM)
        be careful when using it and note the he build a dimension above another
        i.e he builds a sparse table for each row
        the build sparse table over each row's sparse table
    const int N = 505, LG = 10;
10
   int st[N][N][LG][LG];
    int a[N][N], lg2[N];
12
13
    int yo(int x1, int y1, int x2, int y2) {
14
      x2++;
15
      y2++;
16
      int a = \lg 2[x2 - x1], b = \lg 2[y2 - y1];
17
      return max (
18
             \max(st[x1][y1][a][b], st[x2 - (1 << a)][y1][a][b]),
19
             \max(st[x1][y2 - (1 << b)][a][b], st[x2 - (1 << a)][y2 - (1 << b)][a][b]
20
           );
21
    void build(int n, int m) { // 0 indexed
      for (int i = 2; i < N; i++) lg2[i] = lg2[i >> 1] + 1;
```

```
for (int i = 0; i < n; i++) {</pre>
26
         for (int j = 0; j < m; j++) {
27
           st[i][j][0][0] = a[i][j];
28
29
30
      for (int a = 0; a < LG; a++) {</pre>
31
         for (int b = 0; b < LG; b++) {</pre>
          if (a + b == 0) continue;
32
33
           for (int i = 0; i + (1 << a) <= n; i++) {
34
             for (int j = 0; j + (1 << b) <= m; <math>j++) {
35
               if (!a) {
36
                 st[i][j][a][b] = max(st[i][j][a][b - 1], st[i][j + (1 << (b - 1))][a]
37
38
                 st[i][j][a][b] = max(st[i][j][a - 1][b], st[i + (1 << (a - 1))][j][a]
41
42
43
44
```

3.3 hillbert Order

```
///Faster Sorting MO
 2
 3
    const int infinity = (int)1e9 + 42;
    const int64_t llInfinity = (int64_t)1e18 + 256;
    const int module = (int)1e9 + 7;
    const long double eps = 1e-8;
    inline int64_t gilbertOrder(int x, int y, int pow, int rotate) {
 8
 9
10
            return 0;
11
12
         int hpow = 1 << (pow-1);
13
        int seq = (x < hpow) ? (
14
             (y < hpow) ? 0 : 3
15
         ) : (
             (y < hpow) ? 1 : 2
16
17
18
        seg = (seg + rotate) & 3;
19
         const int rotateDelta[4] = {3, 0, 0, 1};
        int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
20
21
        int nrot = (rotate + rotateDelta[seg]) & 3;
22
        int64_t subSquareSize = int64_t(1) << (2*pow - 2);
23
        int64_t ans = seg * subSquareSize;
         int64_t add = gilbertOrder(nx, ny, pow-1, nrot);
24
25
         ans += (seg == 1 || seg == 2) ? add : (subSquareSize - add - 1);
26
        return ans;
27
28
29
    struct Query {
30
        int 1, r, idx;
31
        int64_t ord;
32
33
        inline void calcOrder() {
34
            ord = gilbertOrder(1, r, 21, 0);
35
36
37
38
    inline bool operator<(const Query &a, const Query &b) {</pre>
39
        return a.ord < b.ord;</pre>
40
41
    signed main() {
43
         #ifndef USE_FILE_IO
44
             ios_base::sync_with_stdio(false);
45
         #endif
46
        mt19937 rnd(42);
```

```
49
         int n, m, k; cin >> n >> m; k = rnd() % 1048576;
50
         vector<int> p(n+1);
         for (int i = 0; i < n; i++) {
52
             int val = rnd() % 1048576;
53
             p[i+1] = p[i] ^ val;
54
55
56
         vector<Query> gry(m);
57
         for (int i = 0; i < m; i++) {
             int 1 = rnd() % n + 1, r = rnd() % n + 1;
             if (1 > r) {
                 swap(l, r);
             qry[i].1 = 1; qry[i].r = r;
             qry[i].idx = i;
             gry[i].calcOrder();
65
66
67
         int64_t ans = 0;
68
         vector<int64_t> res(m);
69
         vector<int64_t> cnt((int)2e6, 0);
70
         sort(qry.begin(), qry.end());
 71
         int 1 = 0, r = 1;
72
         ans = (p[1] == k);
73
         cnt[p[0]]++; cnt[p[1]]++;
74
75
         for (Query q: qry) {
76
77
             while (1 > q.1) {
78
                 1--;
79
                 ans += cnt[p[1] ^ k];
80
                 cnt[p[1]]++;
81
82
             while (r < q.r) {
83
                 <u>r</u>++;
84
                 ans += cnt[p[r] ^{\circ} k];
85
                 cnt[p[r]]++;
86
87
             while (1 < q.1) {
                 cnt[p[1]]--;
89
                 ans -= cnt[p[1] ^ k];
91
92
             while (r > q.r) {
93
                 cnt[p[r]]--;
94
                 ans -= cnt[p[r] ^ k];
                 r--;
96
97
             res[q.idx] = ans;
98
99
100
         uint64_t rhsh = 0;
101
         for (int i = 0; i < m; i++) {
102
             rhsh *= (uint64_t)1e9 + 7;
103
             rhsh += (uint64_t)res[i];
104
         cout << rhsh << "\n";</pre>
105
106
107
         return 0;
108
```

3.4 Merge Sort Bit with updates

```
//O(log ^ 2 N) updates and queries

#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/rope>

using namespace std;
```

7

```
using namespace __gnu_pbds;
10
    using namespace __gnu_cxx;
11
12
    template<class T> using Tree = tree<T, null_type, less<T>, rb_tree_tag,
         tree_order_statistics_node_update>;
13
14
    Tree<int> t[N];
15
16
17
    void add(int idx, int v) {
18
        for (int x = ++idx; x < N; x += x & -x) {
19
            t[x].insert(v);
20
21
22
    void erase(int idx, int v) {
23
        for (int x = ++idx; x < N; x += x & -x)
24
            t[x].erase(v);
25
26
    int get(int idx, int limit) {
27
        int ret = 0;
28
        for (int x = ++idx; x; x -= x & -x)
29
            ret += (t[x].order_of_key(limit+1));
30
        return ret;
31
```

3.5 Mo's

```
#include <bits/stdc++.h>
2
3
    int n, qq, arr[N], sz = 1000; // sz is the size of the bucket
    int co[N], ans = 0, ansq[N];
    int cul = 1, cur = 1;
7
    void add(int x) {
        co[arr[x]]++;
8
Q
         if (co[arr[x]] == 1)
10
             ans++;
11
         else if (co[arr[x]] == 2)
12
            ans--;
13
14
15
    void remove(int x) {
        co[arr[x]]--;
17
        if (co[arr[x]] == 1)
18
            ans++:
19
         else if (co[arr[x]] == 0)
20
            ans--;
21
22
    void solve(int 1, int r, int ind) {
24
25
         while (cul < 1) remove(cul++);</pre>
26
         while (cul > 1) add(--cul);
27
        while (cur < r) add(cur++);</pre>
28
        while (cur > r) remove(--cur);
29
         ansq[ind] = ans;
30
31
32
    int main() {
34
        FIO
35
        cin >> qq;
36
                                  \{1/sz,r\},
                                                { 1 , ind}
37
       priority_queue<pair<int, int>, pair<int, int>>, vector<pair<jair<int,</pre>
            int>, pair<int, int>>>, greater<pair<pair<int, int>, pair<int, int>>>> q
38
         for (int i = 0; i < qq; i++) {</pre>
39
             int 1, r;
40
             cin >> 1 >> r;
41
             q.push(\{\{1 / sz, r\}, \{1,i\}\});
42
         while (q.size()) {
```

```
44
             int ind=q.top().second.second,l=q.top().second.first,r=q.top().first.
45
             solve(1, r,ind);
46
             q.pop();
47
48
         for (int i = 0; i < qq; i++)
49
             cout << ansq[i] << endl;</pre>
50
51
52
        return 0;
53
```

3.6 Mo With Updates

```
///O(N^5/3) note that the block size is not a standard size
 3
    #pragma GCC optimize ("03")
    #pragma GCC target ("sse4")
    #include <bits/stdc++.h>
 9
    using namespace std;
10
11
    using 11 = long long;
    const int N = 1e5 +5;
13
    const int M = 2 * N;
   const int blk = 2155;
    const int mod = 1e9 + 7;
    struct Ouerv{
18
     int 1, r, t, idx;
19
      Query (int a = 0, int b = 0, int c = 0, int d = 0) {l=a, r=b, t=c, idx = d;}
20
      bool operator < (Query o) {</pre>
        if(r / blk == o.r / blk && 1 / blk == o.l / blk)return t < o.t;</pre>
21
        if(r / blk == o.r / blk)return 1 < o.1;</pre>
23
        return r < o.r;
24
25
    } Q[N];
27 int a[N], b[N];
28 int cnt1[M], cnt2[N];
   int L = 0, R = -1, K = -1;
    void add(int x) { ///add item to range
31
   // cout << x << '\n';
     cnt2[cnt1[x]]--;
33
      cnt1[x]++;
34
      cnt2[cnt1[x]]++;
35
36 void del(int x) { ///delete item from range
    cnt2[cnt1[x]]--;
38
      cnt1[x]--;
39
      cnt2[cnt1[x]]++;
40
41 map<int,int>id;
42 int cnt;
43 int ans[N];
    int p[N], nxt[N];
44
45
    int prv[N];
46
    void upd(int idx) { ///update item value
      if(p[idx] >= L && p[idx] <= R)
48
        del(a[p[idx]]), add(nxt[idx]);
49
      a[p[idx]] = nxt[idx];
50
51
    void err(int idx) {
      if(p[idx] >= L \&\& p[idx] <= R)
        del(a[p[idx]]), add(prv[idx]);
54
      a[p[idx]] = prv[idx];
55
56
    int main(){
57
      int n, q, 1, r, tp;
```

```
60
       scanf("%d%d", &n, &q);
61
62
       for (int i = 0; i < n; i++) {
63
         scanf("%d", a + i);
 64
         if(id.count(a[i]) == 0)
65
           id[a[i]] = cnt++;
         a[i] = id[a[i]];
66
         b[i] = a[i];
68
69
       int qIdx = 0;
70
       int ord = 0;
71
       while (q--) {
72
73
74
         scanf("%d", &tp);
         if(tp == 1) {
75
           /// ADD Query
76
           scanf("%d%d", &1, &r); --1, --r;
 77
           Q[qIdx] = Query(1,r,ord-1,qIdx); qIdx++;
78
 79
           /// ADD Update
80
           scanf("%d%d",p + ord, nxt + ord); --p[ord];
81
           if(id.count(nxt[ord]) == 0)
 82
              id[nxt[ord]] = cnt++;
 83
           nxt[ord] = id[nxt[ord]];
 84
           prv[ord] = b[p[ord]];
 85
           b[p[ord]] = nxt[ord];
 86
           ++ord;
 87
 88
89
90
       sort(Q,Q+qIdx);
       for(int i = 0; i < qIdx; i++) {</pre>
92
         while (L < Q[i].l) del(a[L++]);
93
         while (L > Q[i].1) add (a[--L]);
94
         while (R < Q[i].r) add (a[++R]);
95
         while (R > Q[i].r) del(a[R--]);
96
         while (K < Q[i].t) upd (++K);
97
         while (K > Q[i].t) err(K--);
98
          ///Solve Query I
99
100
       for(int i = 0; i < qIdx; i++)</pre>
101
         printf("%d\n", ans[i]);
102
103
104
       return 0;
105
```

3.7 Ordered Set

3.8 Persistent Seg Tree

```
1
2  int val[ N * 60 ], L[ N * 60 ], R[ N * 60 ], ptr, tree[N]; /// N * lgN
3  int upd(int root, int s, int e, int idx) {
4    int ret = ++ptr;
5    val[ret] = L[ret] = R[ret] = 0;
```

```
if (s == e) {
            val[ret] = val[root] + 1;
 8
            return ret;
 9
10
        int md = (s + e) \gg 1;
11
        if (idx <= md) {
12
            L[ret] = upd(L[root], s, md, idx), R[ret] = R[root];
13
         } else {
14
            R[ret] = upd(R[root], md + 1, e, idx), L[ret] = L[root];
15
16
        val[ret] = max(val[L[ret]], val[R[ret]]);
17
        return ret;
18
19
   int qry(int node, int s, int e, int 1, int r){
      if(r < s || e < 1 || !node)return 0; //Punishment Value</pre>
21
      if(1 <= s && e <= r) {
22
        return val[node];
23
      int md = (s+e) >> 1;
25
      return max(qry(L[node], s, md, l, r), qry(R[node], md+1, e, l, r));
26
27
    int merge(int x, int y, int s, int e) {
28
        if(!x||!y) return x | y;
29
        if(s == e) {
30
            val[x] += val[y];
31
            return x;
32
33
        int md = (s + e) \gg 1;
        L[x] = merge(L[x], L[y], s, md);
        R[x] = merge(R[x], R[y], md+1,e);
36
        val[x] = val[L[x]] + val[R[x]];
37
        return x;
38
```

3.9 Sqrt Decomposition

```
1 // Source: https://cp-algorithms.com/data_structures/sqrt_decomposition.html
 3
    // input data
    int n;
    vector<int> a (n);
    // preprocessing
    int len = (int) sqrt (n + .0) + 1; // size of the block and the number of blocks
    vector<int> b (len);
    for (int i=0; i<n; ++i)</pre>
11
        b[i / len] += a[i];
12
13
    // answering the queries
14
    for (;;) {
        int 1, r;
16
      // read input data for the next query
17
        int sum = 0;
18
        for (int i=1; i<=r; )</pre>
19
            if (i % len == 0 && i + len - 1 <= r) {</pre>
                // if the whole block starting at i belongs to [1, r]
                 sum += b[i / len];
                 i += len;
            else {
25
                 sum += a[i];
26
                 ++i;
27
28
29
    // If you're getting TLE and can't optimize more, you could reduce the number of
          slow division operations using the following code:
    int sum = 0;
    int c_l = 1 / len,     c_r = r / len;
33
    if (c_l == c_r)
        for (int i=1; i<=r; ++i)</pre>
```

```
sum += a[i];
37
    else {
38
         for (int i=1, end=(c_1+1)*len-1; i<=end; ++i)</pre>
39
              sum += a[i];
40
         for (int i=c_l+1; i<=c_r-1; ++i)</pre>
41
             sum += b[i]:
42
         for (int i=c_r*len; i<=r; ++i)</pre>
43
             sum += a[i];
44
```

3.10 Treap

```
typedef struct item * pitem;
    struct item {
        int prior, value, cnt;
        bool rev;
 -5
        pitem 1, r;
 6
         item(int x, int y, int z){
            value = x;
            prior = y;
 8
            cnt = z;
10
             rev = 0;
11
             1 = r = NULL:
12
13
    };
14
15
    int cnt (pitem it) {
16
        return it ? it->cnt : 0;
17
18
19 void upd_cnt (pitem it) {
20
        if (it)
21
             it->cnt = cnt(it->1) + cnt(it->r) + 1;
22
23
^{24}
    void push (pitem it) {
25
        if (it && it->rev) {
26
             it->rev = false;
27
             swap (it->1, it->r);
             if (it->1) it->1->rev ^= true;
28
29
             if (it->r) it->r->rev ^= true;
30
31
    }
32
33
    void merge (pitem & t, pitem l, pitem r) {
34
        push (1);
35
        push (r);
36
        if (!1 || !r)
37
            t = 1 ? 1 : r;
         else if (l->prior > r->prior)
38
39
            merge (1->r, 1->r, r), t = 1;
40
41
            merge (r->1, 1, r->1), t = r;
42
        upd_cnt (t);
43
44
45
    void split (pitem t, pitem & 1, pitem & r, int key, int add = 0) {
46
47
            return void( 1 = r = 0 );
48
         push (t);
         int cur_key = add + cnt(t->1);
50
         if (key <= cur_key)</pre>
51
            split (t->1, 1, t->1, key, add), r = t;
52
53
            split (t->r, t->r, r, key, add + 1 + cnt(t->1)), 1 = t;
54
        upd_cnt (t);
55
56
57
    void reverse (pitem t, int l, int r) {
58
        pitem t1, t2, t3;
59
         split (t, t1, t2, 1);
60
         split (t2, t2, t3, r-l+1);
```

```
61
        t2->rev ^= true;
62
        merge (t, t1, t2);
63
        merge (t, t, t3);
64
65
66
   void output (pitem t) {
67
        if (!t) return;
        push (t);
69
        output (t->1);
70
        printf ("%c", char(t->value));
71
        output (t->r);
72
73
74
    pitem gettreap(string s) {
75
            pitem ret=NULL;
76
77
           for(i=0;i<s.size();i++)merge(ret,ret,new item(s[i],(rand()<<15)+rand(),</pre>
                1));
78
        return ret;
79
```

3.11 Wavelet Tree

```
// remember your array and values must be 1-based
   struct wavelet tree {
        int lo, hi;
        wavelet_tree *1, *r;
 5
        vector<int> b;
 7
        //nos are in range [x,y]
        //array indices are [from, to)
 9
        wavelet_tree(int *from, int *to, int x, int y) {
10
            lo = x, hi = y;
11
            if (lo == hi or from >= to)
12
                return:
13
             int mid = (lo + hi) / 2;
14
            auto f = [mid] (int x) {
                return x <= mid;</pre>
15
16
17
            b.reserve(to - from + 1);
18
            b.pb(0);
19
            for (auto it = from; it != to; it++)
20
                b.pb(b.back() + f(*it));
21
             //see how lambda function is used here
22
            auto pivot = stable_partition(from, to, f);
23
             l = new wavelet_tree(from, pivot, lo, mid);
24
             r = new wavelet_tree(pivot, to, mid + 1, hi);
25
26
27
        //kth smallest element in [1, r]
28
        int kth(int 1, int r, int k) {
29
            if (1 > r)
30
                return 0;
31
            if (lo == hi)
32
                return lo;
33
            int inLeft = b[r] - b[1 - 1];
34
            int lb = b[1 - 1]; //amt of nos in first (1-1) nos that go in left
35
            int rb = b[r]; //amt of nos in first (r) nos that go in left
36
             if (k <= inLeft)</pre>
37
                return this->l->kth(lb + 1, rb, k);
            return this->r->kth(1 - 1b, r - rb, k - inLeft);
38
39
40
41
         //count of nos in [1, r] Less than or equal to k
42
        int LTE(int 1, int r, int k) {
43
            if (1 > r \text{ or } k < 10)
44
                return 0;
45
            if (hi <= k)
46
                return r - 1 + 1;
47
            int 1b = b[1 - 1], rb = b[r];
48
            return this->1->LTE(lb + 1, rb, k) + this->r->LTE(l - lb, r - rb, k);
49
```

```
__
```

```
51
         //count of nos in [1, r] equal to k
52
         int count(int 1, int r, int k) {
53
             if (1 > r \text{ or } k < 10 \text{ or } k > hi)
54
                 return 0:
55
             if (lo == hi)
56
                 return r - 1 + 1;
             int 1b = b[1 - 1], rb = b[r], mid = (10 + hi) / 2;
57
             if (k <= mid)
59
                 return this->l->count(lb + 1, rb, k);
60
             return this->r->count(1 - 1b, r - rb, k);
61
62
    };
```

4 DP

#include<iostream>

4.1 Dynamic Convex Hull Trick

```
#include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    struct Line
9
         11 m, b;
10
        mutable function<const Line*()> succ;
11
        bool operator<(const Line& other) const
12
13
             return m < other.m;</pre>
14
15
        bool operator<(const 11 &x) const
16
17
             const Line* s = succ();
18
             if (!s)
19
                 return 0;
20
             return b - s \rightarrow b < (s \rightarrow m - m) * x;
21
    // will maintain upper hull for maximum
^{24}
    struct HullDynamic : public multiset<Line, less<>>
25
26
        bool bad(iterator y)
27
28
             auto z = next(y);
29
             if (y == begin())
30
31
                 if (z == end())
32
                     return 0;
33
                 return y->m == z->m && y->b <= z->b;
34
35
             auto x = prev(y);
36
             if (z == end())
37
                 return y->m == x->m && y->b <= x->b;
38
             return (ld) (x-b - y-b)*(z-m - y-m) >= (ld) (y-b - z-b)*(y-m - x-m)
39
40
        void insert_line(ll m, ll b)
41
42
             auto y = insert({ m, b });
43
             y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
44
             if (bad(y))
45
46
                 erase(y);
47
49
             while (next(y) != end() && bad(next(y)))
                 erase(next(y));
             while (y != begin() && bad(prev(y)))
```

```
erase(prev(y));
53
55
        11 query(11 x)
56
57
58
            auto 1 = *lower_bound(x);
            return 1.m * x + 1.b;
60
61
    };
62
    int main()
65
66
```

4.2 Dynamic Connectivety with SegTree

```
#pragma GCC optimize("03")
    #pragma GCC optimize ("unroll-loops")
    #pragma GCC target("avx,avx2,fma")
    using namespace std;
    #include "bits/stdc++.h"
    #define pb push_back
    #define F first
    #define S second
    #define f(i, a, b) for(int i = a; i < b; i++)
    #define all(a) a.begin(),a.end()
    #define rall(a) a.rbegin(),a.rend()
    #define sz(x) (int)(x).size()
    //#define mp make_pair
    #define popCnt(x) (__builtin_popcountll(x))
17
    typedef long long 11;
    typedef pair<int, int> ii;
    using ull = unsigned long long;
    const int N = 1e5+5, LG = 17, MOD = 1e9 + 7;
    const long double PI = acos(-1);
    struct PT{
24
        11 x, y;
25
        PT() {}
26
        PT(11 a, 11 b):x(a), y(b) {}
27
        PT operator - (const PT & o) {return PT{x-o.x,y-o.y};}
28
        bool operator < (const PT & o) const {return make_pair(x,y) < make_pair(o.x,</pre>
             o.y);}
    ll cross(PT x, PT y) {
        return x.x * y.y - x.y * y.x;
33
   PT val[300005];
34 bool in[300005];
   ll qr[300005];
    bool ask[300005];
    11 ans[N];
    vector<PT> t[300005 * 4]; ///segment tree holding points to queries
    void update(int node, int s, int e, int 1, int r, PT x) {
        if(r < s || e < 1) return;
41
        if(1 \le s \&\& e \le r)  { ///add this point to maximize it with queries in
             this range
42
            t[node].pb(x);
43
            return;
44
45
        int md = (s + e) \gg 1;
46
        update(node<<1, s, md, 1, r, x);
47
        update (node <<1|1, md+1, e, l, r, x);
48
49
   vector<PT> stk;
50 inline void addPts(vector<PT> v) {
51
        stk.clear();
                        ///reset the data structure you are using
        sort(all(v));
```

```
53
          ///build upper envelope
54
         for(int i = 0; i < v.size(); i++) {</pre>
55
             while(sz(stk) > 1 && cross(v[i] - stk.back(), stk.back() - stk[stk.size
                   ()-2]) <= 0)
56
                  stk.pop_back();
57
             stk.push_back(v[i]);
58
59
60
     inline 11 calc(PT x, 11 val) {
61
         ///mb+v
62
         return x.x * val + x.y;
63
64
65
    ll query(ll x) {
66
         if(stk.empty())
67
             return LLONG_MIN;
         int lo = 0, hi = stk.size() - 1;
68
69
          while(lo + 10 < hi) {
70
             int md = lo + (hi-lo) / 2;
71
             if(calc(stk[md+1],x) > calc(stk[md],x))
 72
                 lo = md + 1;
73
             else
74
                 hi = md;
75
76
         11 ans = LLONG MIN:
77
         for (int i = lo; i <= hi; i++)</pre>
78
             ans = max(ans, calc(stk[i], x));
 79
          return ans;
 80
81
     void solve(int node, int s, int e) {      ///Solve queries
82
         addPts(t[node]);
                              ///note that there is no need to add/delete just build
              for t[node]
 83
         f(i,s,e+1){
 84
             if(ask[i]) {
 85
                 ans[i] = max(ans[i], query(qr[i]));
 86
87
 88
         if(s==e)return;
 89
         int md = (s + e) \gg 1;
90
         solve(node<<1,s,md);</pre>
91
         solve(node<<1|1,md+1,e);
92
93
     void doWork() {
94
95
         int n;
96
         cin >> n;
97
         stk.reserve(n);
98
         f(i,1,n+1) {
99
             int tp;
100
             cin >> tp;
101
             if(tp == 1) {
                             ///Add Query
                 int x, y;
102
103
                 cin >> x >> y;
104
                 val[i] = PT(x, y);
105
                 in[i] = 1;
106
                 else if(tp == 2) { ///Delete Query
107
                  int x;
108
                  cin >> x;
109
                  if(in[x])update(1, 1, n, x, i - 1, val[x]);
110
                  in[x] = 0;
111
                 else {
112
                  cin >> qr[i];
113
                  ask[i] = true;
114
115
         f(i,1,n+1) ///Finalize Query
116
117
             if(in[i])
                 update(1, 1, n, i, n, val[i]);
118
119
120
         f(i,1,n+1) ans [i] = LLONG_MIN;
121
         solve(1, 1, n);
122
         f(i,1,n+1)
123
         if(ask[i]) {
124
             if(ans[i] == LLONG_MIN)
```

```
125
                  cout << "EMPTY SET\n";</pre>
126
              else
127
                  cout << ans[i] << '\n';
128
129
130
     int32_t main() {
131
     #ifdef ONLINE_JUDGE
132
133
         ios_base::sync_with_stdio(0);
134
          cin.tie(0);
135
     #endif // ONLINE JUDGE
136
         int t = 1;
137
          cin >> t;
138
         while (t--) {
139
              doWork();
140
141
         return 0;
142
```

4.3 Li Chao Tree

```
#include<iostream>
    #include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    struct Line
 9
10
        Line(11 \, m, 11 \, b) : m(m), b(b) {}
11
         11 operator()(11 x)
12
13
             return m * x + b;
14
15
    };
16
    struct node
17
18
         node * left, * right ;
19
        Line line ;
20
        node(node * left, node *right, Line line):left(left), right(right), line(
             line) {}
21
        node * getLeft()
22
             if(left==NULL)
23
24
                 left= new node (NULL, NULL, Line(0, 1e18));
25
             return left ;
26
27
        node * getright()
28
29
             if(right==NULL)
30
                 right= new node (NULL, NULL, Line (0, 1e18));
31
             return right ;
32
33
        void insert(Line newline, int 1, int r)
34
35
             int m = (1+r)/2;
36
             bool lef=newline(1)<line(1);</pre>
37
             bool mid=newline(m) <line(m);</pre>
38
39
             if (mid)
40
                 swap(line, newline);
41
             if(r-1==1)
42
                 return ;
43
             else if(lef!=mid)
44
                 getLeft()->insert(newline,1,m);
45
             else
46
                 getright() ->insert(newline, m, r);
47
48
        11 query(int x, int 1, int r)
49
50
             int m = (1 + r) / 2;
```

```
15
```

```
if(r - 1 == 1)
52
                 return line(x);
53
             else if (x < m)
                 return min(line(x), getLeft()->query(x, 1, m));
55
56
                 return min(line(x), getright()->query(x, m, r));
57
58
        void deletee()
59
             if(left!=NULL)
60
61
                 left->deletee();
62
             if(right!=NULL)
63
                 right->deletee();
             free(this);
65
66
    };
67
    int main()
68
69
70
        node * root = new node(NULL, NULL, Line(0,5));
71
         root->insert(Line(1,-3),1,100);
72
73
         for (int i=1; i<=10; i++)</pre>
74
             cout << root -> query (i, 1, 100) << "\n";
75
```

4.4 CHT Line Container

```
1
    struct Line
 2
 3
         mutable 11 m, b, p;
 4
         bool operator<(const Line& o) const
 5
 6
             return m < o.m;</pre>
 7
 8
         bool operator<(11 x) const
 Q
10
             return p < x;
11
12
13
14
    struct LineContainer : multiset<Line, less<>>
15
16
         // (for doubles, use inf = 1/.0, div(a,b) = a/b)
         static const 11 inf = LLONG_MAX;
17
         ll div(ll db, ll dm) // floored division
18
19
             return db / dm - ((db ^ dm) < 0 && db % dm);</pre>
20
21
22
         bool isect(iterator x, iterator y)
23
24
             if (y == end())
\frac{25}{26}
                 x->p = inf;
27
                 return false;
29
             if (x->m == y->m)
30
                 x->p = x->b > y->b ? inf : -inf;
31
32
                 x->p = div(y->b - x->b, x->m - y->m);
\frac{33}{34}
             return \bar{x}->p >= y->p;
35
         void add(ll m, ll b)
36
37
             auto z = insert(\{m, b, 0\}), y = z++, x = y;
             while (isect(y, z))
39
                 z = erase(z);
40
             if (x != begin() && isect(--x, y))
41
                 isect(x, y = erase(y));
42
             while ((y = x) != begin() && (--x)->p >= y->p)
43
                 isect(x, erase(y));
```

5 Geometry

6 Graphs

7 Math

7.1 Xor With Gauss

```
1 /*
 2
        Some applications
 3
        If you want to find the maximum in xor subset
        just ans = max(ans, ans ^p[i]) for all i
        if you want to count the number of subsets with a certain value
        check all different subsets of p
   ll p[66];
 9
   bool add(ll x) {
        for(int i = 60; (~i) && x; --i) {
10
            if(x >> i & 1) {
11
                if(!p[i]) {
12
13
                    p[i] = x;
14
                    return true;
15
                    else {
16
                    x ^= p[i];
17
18
19
20
        return false;
21
```

7.2 Josephus

```
1 // n = total person
   // will kill every kth person, if k = 2, 2, 4, 6, ...
   // returns the mth killed person
    11 josephus(11 n, 11 k, 11 m) {
     m = n - m;
      if (k <= 1) return n - m;</pre>
      11 i = m;
      while (i < n) {
        11 r = (i - m + k - 2) / (k - 1);
10
        if ((i + r) > n) r = n - i;
11
        else if (!r) r = 1;
        i += r;
13
        m = (m + (r * k)) % i;
14
      } return m + 1;
15
```

7.3 Matrix Power/Multiplication

```
1  struct Matrix {
2
3      const static int D = 100;
4      int a[D][D];
```

```
2
3
3
3
3
3
```

```
6
        Matrix(int val) {
             for (int i = 0; i < D; i++)
                 for (int j = 0; j < D; j++)
9
                    a[i][j] = val;
10
11
12
            memset(a, 0, sizeof a);
13
        void initIdentity() {
14
15
            clear();
16
             for (int i = 0; i < D; i++)
17
                a[i][i] = 1;
18
19
        int * operator [](int r) {
20
            return a[r];
21
22
        const int * operator [](int r) const{
23
            return a[r];
24
25
26
        friend Matrix operator * (const Matrix & a, const Matrix & b) {
27
            Matrix ret(0);
28
             for (int k = 0; k < D; k++)
29
                 for(int i = 0; i < D; i++)if(a[i][k])</pre>
30
                     for (int j = 0; j < D; j++)
31
                         ret[i][j] = (ret[i][j] + 111 * a[i][k] * b[k][j]) % MOD;
32
             return ret;
33
34
35
36
    Matrix raiseMatrix (Matrix trans, 11 k) {
37
        Matrix res(0);
38
        res.initIdentity();
39
        for(;k;k>>=1,trans = trans * trans)
40
            if(k & 1)
41
                res = res * trans;
42
        return res;
43
```

7.4 Rabin Miller Primality check

```
2
    // n < 4,759,123,141
                                         3: 2, 7, 61
    // n < 1,122,004,669,633
                                         4: 2, 13, 23, 1662803
    // n < 3,474,749,660,383
                                         6 : pirmes <= 13
   // n < 3,825,123,056,546,413,051
                                       9 : primes <= 23
    int testPrimes[] = {2,3,5,7,11,13,17,19,23};
7
    struct MillerRabin{
10
     ///change K according to n
11
      const int K = 9;
12
      11 mult(11 s, 11 m, 11 mod){
13
        if(!m) return 0;
14
        11 \text{ ret} = \text{mult}(s, m/2, mod);
15
        ret = (ret + ret) % mod;
16
        if (m & 1) ret = (ret + s) % mod;
17
        return ret;
18
19
      11 power(ll x, ll p, ll mod){
21
        11 s = 1, m = x;
22
        while (p) {
23
          if(p&1) s = mult(s, m, mod);
          p >>= 1;
          m = mult(m, m, mod);
26
27
        return s;
28
29
      bool witness(ll a, ll n, ll u, int t) {
```

```
11 x = power(a, u, n), nx;
32
         for (int i = 0; i < t; i++) {
          nx = mult(x, x, n);
33
34
           if (nx == 1 \text{ and } x != 1 \text{ and } x != n-1) return 1;
35
36
37
        return x != 1;
38
39
40
      bool isPrime(ll n){ // return 1 if prime, 0 otherwise
41
        if(n < 2) return 0;
42
        if(!(n&1)) return n == 2;
        for(int i = 0; i < K; i++)if(n == testPrimes[i])return 1;</pre>
43
44
        11 u = n-1; int t = 0;
45
46
        while (u&1) u >>= 1, t++; // n-1 = u*2^t
47
48
        for(int i = 0; i < K; i++) if(witness(testPrimes[i], n, u, t)) return 0;</pre>
49
50
51 }tester;
```

8 Strings

8.1 Aho-Corasick Mostafa

```
struct AC FSM {
    #define ALPHABET SIZE 26
         struct Node {
             int child[ALPHABET_SIZE], failure = 0, match_parent = -1;
             vector<int> match;
 8
            Node() {
 9
                 for (int i = 0; i < ALPHABET_SIZE; ++i)child[i] = -1;</pre>
10
11
        };
12
13
        vector<Node> a;
14
15
        AC_FSM() {
16
            a.push_back(Node());
17
18
19
        void construct_automaton(vector<string> &words) {
20
             for (int w = 0, n = 0; w < words.size(); ++w, n = 0) {
21
                 for (int i = 0; i < words[w].size(); ++i) {</pre>
22
                     if (a[n].child[words[w][i] - 'a'] == -1) {
                         a[n].child[words[w][i] - 'a'] = a.size();
24
                         a.push_back(Node());
25
26
                     n = a[n].child[words[w][i] - 'a'];
27
                 a[n].match.push_back(w);
29
30
             queue<int> q;
31
             for (int k = 0; k < ALPHABET_SIZE; ++k) {</pre>
32
                 if (a[0].child[k] == -1) a[0].child[k] = 0;
33
                 else if (a[0].child[k] > 0) {
34
                     a[a[0].child[k]].failure = 0;
35
                     q.push(a[0].child[k]);
36
37
38
             while (!q.empty()) {
39
                 int r = q.front();
40
41
                 for (int k = 0, arck; k < ALPHABET_SIZE; ++k) {</pre>
42
                     if ((arck = a[r].child[k]) != -1) {
43
                         q.push(arck);
                         int v = a[r].failure;
```

```
while (a[v].child[k] == -1) v = a[v].failure;
                         a[arck].failure = a[v].child[k];
                         a[arck].match_parent = a[v].child[k];
                         while (a[arck].match parent != -1 &&
49
                               a[a[arck].match_parent].match.empty())
50
                             a[arck].match_parent =
51
                                     a[a[arck].match_parent].match_parent;
53
                }
54
55
56
57
        void aho_corasick(string &sentence, vector<string> &words,
58
                          vector<vector<int> > &matches) {
59
            matches.assign(words.size(), vector<int>());
60
            int state = 0, ss = 0;
            for (int i = 0; i < sentence.length(); ++i, ss = state) {</pre>
61
62
                while (a[ss].child[sentence[i] - 'a'] == -1)
63
                    ss = a[ss].failure;
64
                state = a[state].child[sentence[i] - 'a'] = a[ss].child[sentence[i]
65
                for (ss = state; ss != -1; ss = a[ss].match_parent)
                    for (int w: a[ss].match)
67
                        matches[w].push_back(i + 1 - words[w].length());
68
69
70
    };
```

8.2 Aho-Corasick Anany

```
int trie[N][A];
   int go[N][A]; //holds the node that you will go to after failure and stuff
    int ptr;
    11 ans[N]; //this node is a string terminator;
    int fail[N]: ///the failure function for each
    void BFS() {
7
        queue<int> q;
8
        f(i,0,A) {
Q
            if(trie[0][i]) {
10
                q.push(trie[0][i]);
11
                fail[trie[0][i]] = 0;
12
13
            go[0][i] = trie[0][i];
14
15
16
        while(q.size()) {
            auto node = q.front();
17
18
19
            ans[node] += ans[fail[node]]; //propagate fail[i] to ans[i]
20
            for (int i = 0; i < A; i++) {
21
                if(trie[node][i]) { ///calculate failure for you child
22
                    int to = trie[node][i];
23
                    int cur = fail[node]; ///int g = pi[i-1]
24
                    while(cur && !trie[cur][i]) ///while(q && s[q] != s[i])
25
                        cur = fail[cur]; ///g = pi[g-1]
26
                    if(trie[cur][i])cur = trie[cur][i]; ///g += s[i] == s[g]
                    fail[to] = cur; //pi[i] = g
28
                    q.push(to);
                    go[node][i] = trie[node][i];
30
                    else {
31
                    go[node][i] = go[fail[node]][i];
32
33
34
35
    void ins(string s, ll val) {
37
        int cur = 0;
38
        string sx = "";
39
        for(char c : s) {
40
            sx.push_back(c);
            if(!trie[cur][c - 'a']) {
41
                trie[cur][c - 'a'] = ++ptr;
```

8.3 KMP Anany

```
1 vector<int> fail(string s) {
        int n = s.size();
 3
        vector<int> pi(n);
        for(int i = 1; i < n; i++) {</pre>
            int q = pi[i-1];
            while (g \&\& s[i] != s[g])
              g = pi[g-1];
            q += s[i] == s[q];
            pi[i] = g;
11
        return pi;
13
    vector<int> KMP(string s, string t) {
14
        vector<int> pi = fail(t);
        vector<int> ret;
15
16
        for(int i = 0, g = 0; i < s.size(); i++) {</pre>
17
            while (g && s[i] != t[g])
               g = pi[g-1];
19
            q += s[i] == t[q];
20
            if(g == t.size()) { ///occurrence found
21
                ret.push_back(i-t.size()+1);
22
                q = pi[q-1];
25
        return ret;
26
```

8.4 Manacher Kactl

```
1 // If the size of palindrome centered at i is x, then d1[i] stores (x+1)/2.
 3
    vector<int> d1(n);
    for (int i = 0, l = 0, r = -1; i < n; i++) {
        int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
        while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k]) {
 9
        d1[i] = k--;
        if(i + k > r) {
10
           1 = i - k;
11
            r = i + k;
12
13
14
15
16
17
    // If the size of palindrome centered at i is x, then d2[i] stores x/2
19
   vector<int> d2(n);
    for (int i = 0, l = 0, r = -1; i < n; i++) {
        int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1);
        while (0 \le i - k - 1 \&\& i + k \le n \&\& s[i - k - 1] == s[i + k]) {
24
25
        d2[i] = k--;
        if (i + k > r) {
            1 = i - k - 1;
            r = i + k;
29
30
```

```
struct SuffixArray {
         using vi = vector<int>;
         #define rep(i,a,b) for(int i = a; i < b; i++)
 3
 5
             Note this code is considers also the empty suffix
 6
             so hear sa[0] = n and sa[1] is the smallest non empty suffix
             and sa[n] is the largest non empty suffix
             also LCP[i] = LCP(sa[i-1], sa[i]), meanining LCP[0] = LCP[1] = 0
             if you want to get LCP(i..j) you need to build a mapping between
10
             sa[i] and i, and build a min sparse table to calculate the minimum
             note that this minimum should consider sa[i+1...j] since you don't want
11
12
             to consider LCP(sa[i], sa[i-1])
13
14
             you should also print the suffix array and lcp at the beginning of the
15
             to clarify this stuff
16
17
         vi sa, lcp;
18
        SuffixArray(string& s, int lim=256) { // or basic_string<int>
            int n = sz(s) + 1, k = 0, a, b;
19
20
             vi x(all(s)+1), y(n), ws(max(n, lim)), rank(n);
21
             sa = lcp = y, iota(all(sa), 0);
             for (int j = 0, p = 0; p < n; j = max(1, j * 2), lim = p) {
                p = j, iota(all(y), n - j);
rep(i,0,n) if (sa[i] >= j) y[p++] = sa[i] - j;
23
24
25
                 fill(all(ws), 0);
26
                 rep(i,0,n) ws[x[i]]++;
                 rep(i,1,lim) ws[i] += ws[i - 1];
27
                 for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
29
                 swap(x, y), p = 1, x[sa[0]] = 0;
30
                 rep(i,1,n) = sa[i-1], b = sa[i], x[b] =
31
                     (y[a] == y[b] && y[a + j] == y[b + j]) ? p - 1 : p++;
32
33
             rep(i,1,n) rank[sa[i]] = i;
34
             for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)
35
                 for (k \&\& k--, j = sa[rank[i] - 1];
36
                        s[i + k] == s[j + k]; k++);
37
    };
```

8.6 Suffix Automaton Anany

```
///Note it's better to use addNode to clear a node before using it
   ///at the start of each test case use initAutomaton
3
    int last = 0, cntState = 1;
    int nxt[N * 2][26];
    int len[N * 2], link[N * 2], firstPos[N * 2], cnt[N * 2];
    void addNode(int i) {
        memset(nxt[i], 0, sizeof nxt[i]);
10
        link[i] = -1;
11
        cnt[i] = 0;
12
13
14
    void initAutomaton() {
15
        cntState = 1;
16
        last = 0:
17
        addNode(last);
18
19
20
   int addChar(char c) {
        c -= 'a'; ///note this offset
23
        int p = last;
24
        int cur = cntState++;
25
        addNode(cur);
26
        cnt[cur] = 1; ///extra
        len[cur] = len[last] + 1;
```

```
firstPos[cur] = len[cur] - 1; ///extra
29
        while (p != -1 && nxt[p][c] == 0) {
30
            nxt[p][c] = cur;
31
            p = link[p];
32
33
34
        if(p == -1) {
35
            link[cur] = 0;
36
37
            int q = nxt[p][c];
38
            if(len[q] == len[p] + 1) {
39
                link[cur] = q;
40
               else {
                int clone = cntState++;
42
                link[clone] = link[q];
                firstPos[clone] = firstPos[q]; ///extra
                len[clone] = len[p] + 1;
                link[q] = link[cur] = clone;
45
                memcpy(nxt[clone], nxt[q], sizeof nxt[q]);
47
                cnt[clone] = 0; ///extra
                f(i,0,26)nxt[clone][i] = nxt[q][i];
49
                while (p != -1 && nxt[p][c] == q) {
                    nxt[p][c] = clone;
                    p = link[p];
52
53
55
        last = cur:
56
        return cur;
```

8.7 Suffix Automaton Mostafa

```
#include <bits/stdc++.h>
    #define FIO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    typedef long long 11;
    typedef long double ld;
    const int N = 2e6 + 9, M = 5e5 + 9;
    struct SA {
10
        struct node {
11
            int to[26];
12
            int link, len, co = 0;
13
14
            node() {
15
                memset(to, 0, sizeof to);
16
                co = 0, link = 0, len = 0;
17
18
        };
19
20
        int last, sz;
21
        vector<node> v;
22
23
            v = vector<node>(1);
25
            last = 0, sz = 1;
26
27
28
        void add_letter(int c) {
29
            int p = last;
30
            last = sz++;
31
            v.push_back({});
32
            v[last].len = v[p].len + 1;
33
            v[last].co = 1;
34
            for (; v[p].to[c] == 0; p = v[p].link)
35
                v[p].to[c] = last;
36
            if (v[p].to[c] == last) {
37
                v[last].link = 0;
                return;
```

```
int q = v[p].to[c];
41
             if (v[q].len == v[p].len + 1) {
                v[last].link = q;
43
                 return;
44
45
             int cl = sz++;
46
             v.push_back(v[q]);
47
             v.back().co = 0;
48
             v.back().len = v[p].len + 1;
49
             v[last].link = v[q].link = cl;
50
51
             for (; v[p].to[c] == q; p = v[p].link)
52
                v[p].to[c] = cl;
53
54
55
        void build co() {
56
             priority_queue<pair<int, int>> q;
57
             for (int i = sz - 1; i > 0; i--)
58
                 q.push({v[i].len, i});
59
             while (q.size()) {
60
                 int i = q.top().second;
61
                q.pop();
62
                v[v[i].link].co += v[i].co;
63
64
65
    };
66
67
    int main() {
68
69
70
         return 0;
71
```

8.8 Suffix Automaton With Rollback Mostafa

```
#include <bits/stdc++.h>
    #define FIO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    typedef long long 11;
    typedef long double 1d;
    const int N = 2e6 + 9, M = 5e5 + 9;
 9
    struct SA {
10
         struct node {
11
             int to[26];
12
             int link, len, co = 0;
13
14
             node() {
15
                 memset(to, 0, sizeof to);
16
                 co = 0, link = 0, len = 0;
17
18
        };
19
20
         struct LogNode {
21
             int last, sz;
22
             vector<pair<int, int>, int>> edges;
23
             pair<int, int> LinksUpdate = {0, 0};
24
25
26
         int last, sz;
\frac{27}{28}
        vector<node> v;
        vector<LogNode> logs;
29
30
31
             v = vector<node>(1);
32
             last = 0, sz = 1;
33
34
35
         void add_letter(int c) {
36
             logs.push_back({});
             logs.back().last = last;
```

```
logs.back().sz = sz;
39
40
             int p = last;
41
             last = sz++;
42
             v.push_back({});
43
             v[last].len = v[p].len + 1;
44
             v[last].co = 1;
45
             for (; v[p].to[c] == 0; p = v[p].link) {
46
                 logs.back().edges.push_back({{p, c}, 0});
47
                 v[p].to[c] = last;
48
49
             if (v[p].to[c] == last) {
50
                 v[last].link = 0;
51
                 return;
52
53
             int q = v[p].to[c];
54
             if (v[q].len == v[p].len + 1) {
55
                 v[last].link = q;
56
                 return;
57
             int c1 = sz++;
58
59
             v.push_back(v[q]);
60
             v.back().co = 0;
61
             v.back().len = v[p].len + 1;
62
             logs.back().LinksUpdate = {q, v[q].link};
v[last].link = v[q].link = cl;
63
64
             for (; v[p].to[c] == q; p = v[p].link) {
                 logs.back().edges.push_back({{p, c}, q});
65
66
                 v[p].to[c] = cl;
67
68
69
        void rollback() {
70
             assert(logs.size());
71
             auto log = logs.back();
72
             while (v.size() > log.sz)
73
                 v.pop_back();
             for (auto edge: log.edges)
74
                 v[edge.first.first].to[edge.first.second] = edge.second;
75
76
             if (log.LinksUpdate.first != 0)
                 v[log.LinksUpdate.first].link = log.LinksUpdate.second;
78
             last = log.last;
79
             sz = log.sz;
80
             logs.pop_back();
81
82
    };
83
84
    int main() {
85
        FIO
86
87
         return 0;
```

8.9 Zalgo Anany

```
1 int z[N], n;
    void Zalgo(string s) {
         int L = 0, R = 0;
         for (int i = 1; i < n; i++) {
 5
             if(i<=R&&z[i-L] < R - i + 1)z[i] = z[i-L];</pre>
 6
             else {
                 L = i;
                 R = max(R, i);
 Q
                 while (R < n \&\& s[R-L] == s[R])R++;
10
                 z[i] = R-L; --R;
11
12
13
```

9 Trees

9.1 Centroid Decomposition

```
1
2
            1. consider path(a,b) can be decomposed to path(a,lca(a,b)) and path(b,
            where lca(a,b) is the lca on the centroid tree
            2. Each one of the n^2 paths is the concatenation of two paths in a set
            paths from a node to all its ancestors in the centroid decomposition.
            3. Ancestor of a node in the original tree is either an ancestor in the
            a descendadnt
   vector<int> adj[N]; ///adjacency list of original graph
11
    int sz[N];
   bool used[N];
    int centPar[N]; //parent in centroid
    void init(int node, int par) { ///initialize size
16
        sz[node] = 1;
17
        for(auto p : adj[node])
18
            if(p != par && !used[p]) {
19
               init(p, node);
20
                sz[node] += sz[p];
21
    ^{24}
        for(int p : adj[node])
25
            if(!used[p] && p != par && sz[p] * 2 > limit)
26
            return centroid(p, node, limit);
27
        return node:
28
29
    int decompose(int node) {
        init(node, node); ///calculate size
31
        int c = centroid(node, node, sz[node]); //get centroid
32
        used[c] = true;
33
        for(auto p : adj[c])if(!used[p.F]) {      ///initialize parent for others and
            decompose
            centPar[decompose(p.F)] = c;
35
36
        return c;
37
    void update(int node, int distance, int col) {
        int centroid = node;
40
        while (centroid) {
41
            ///solve
42
            centroid = centPar[centroid];
45
    int query(int node) {
47
        int ans = 0;
48
49
        int centroid = node;
        while(centroid) {
51
            ///solve
52
            centroid = centPar[centroid];
53
55
        return ans;
```

9.2 Dsu On Trees

```
1 const int N = 1e5 + 9;
2 vector<int> adj[N];
3 int bigChild[N], sz[N];
```

```
void dfs(int node, int par) {
        for(auto v : adj[node]) if(v != par){
            dfs(v, node);
            sz[node] += sz[v];
            if(!bigChild[node] || sz[v] > sz[bigChild[node]]) {
                bigChild[node] = v;
10
11
    void add(int node, int par, int bigChild, int delta) {
14
        ///modify node to data structure
        for(auto v : adj[node])
        if(v != par && v != bigChild)
19
            add(v, node, bigChild, delta);
20
    void dfs2(int node, int par, bool keep) {
        for(auto v : adj[node])if(v != par && v != bigChild[node]) {
            dfs2(v, node, 0);
        if(bigChild[node]) {
            dfs2(bigChild[node], node, true);
29
        add(node, par, bigChild[node], 1);
30
        ///process queries
31
        if(!keep) {
            add(node, par, -1, -1);
33
34
```

9.3 Heavy Light Decomposition (Along with Euler Tour)

```
Notes:
            2. solve function iterates over segments and handles them seperatly
            if you're gonna use it make sure you know what you're doing
            3. to update/query segment in[node], out[node]
            4. to update/query chain in[nxt[node]], in[node]
            nxt[node]: is the head of the chain so to go to the next chain node =
 9
10
    int sz[mxN], nxt[mxN];
    int in[N], out[N], rin[N];
11
    vector<int> g[mxN];
    int par[mxN];
15
    void dfs_sz(int v = 0, int p = -1) {
16
        sz[v] = 1;
        par[v] = p;
17
18
        for (auto &u : g[v]) {
19
            if (u == p) {
20
                swap(u, g[v].back());
21
            if(u == p) continue;
23
            dfs_sz(u,v);
            sz[v] += sz[u];
            if (sz[u] > sz[g[v][0]])
                swap(u, g[v][0]);
28
        if(v != 0)
29
            g[v].pop_back();
30
    void dfs_hld(int v = 0) {
33
        in[v] = t++;
        rin[in[v]] = v;
35
        for (auto u : g[v]) {
            nxt[u] = (u == g[v][0] ? nxt[v] : u);
            dfs_hld(u);
```

```
15
```

```
39
        out[v] = t;
41
42
    int n;
43
    bool isChild(int p, int u) {
44
      return in[p] <= in[u] && out[u] <= out[p];</pre>
45
46
    int solve(int u,int v) {
47
        vector<pair<int,int> > segu;
48
         vector<pair<int,int> > seqv;
49
         if(isChild(u,v)){
50
          while(nxt[u] != nxt[v]){
51
             seqv.push_back(make_pair(in[nxt[v]], in[v]));
52
             v = par[nxt[v]];
53
54
           segv.push_back({in[u], in[v]});
         } else if(isChild(v,u)){
55
           while(nxt[u] != nxt[v]){
57
           segu.push_back(make_pair(in[nxt[u]], in[u]));
58
          u = par[nxt[u]];
59
60
          segu.push_back({in[v], in[u]});
61
      } else {
62
          while (u != v) {
63
            if(nxt[u] == nxt[v]) {
64
               if(in[u] < in[v]) seqv.push_back({in[u],in[v]}), R.push_back({u+1,v}</pre>
65
               else sequ.push_back({in[v],in[u]}), L.push_back({v+1,u+1});
66
              u = v:
67
               break;
68
             } else if(in[u] > in[v]) {
69
               segu.push_back({in[nxt[u]],in[u]}), L.push_back({nxt[u]+1, u+1});
70
               u = par[nxt[u]];
71
72
               segv.push_back({in[nxt[v]],in[v]}), R.push_back({nxt[v]+1, v+1});
73
               v = par[nxt[v]];
74
75
76
77
         reverse(segv.begin(), segv.end());
78
         int res = 0, state = 0;
79
         for(auto p : sequ) {
80
             gry(1,1,0,n-1,p.first,p.second,state,res);
81
82
         for(auto p : segv) {
83
             qry(0,1,0,n-1,p.first,p.second,state,res);
84
85
         return res;
86
```

9.4 LCA

```
const int N = 1e5 + 5;
    const int LG = 18;
4
    vector<int> adj[N];
5
    int pa[N][LG], lvl[N];
    int in[N], out[N], timer;
    void dfs(int u, int p) {
      in[u] = ++timer;
      for (int k = 1; k < LG; k++)
10
        pa[u][k] = pa[pa[u][k-1]][k-1];
11
      for(auto v : adj[u])
12
        if(v != p){
           lvl[v] = lvl[u] + 1;
13
14
           pa[v][0] = u;
15
           dfs(v, u);
16
17
      out[u] = timer;
18
    int LCA(int u, int v) {
```

```
if(lvl[u] > lvl[v])
21
        swap(u,v);
22
       int d = lvl[v] - lvl[u];
       for (int k = 0; k < LG; k++)
24
        if(d >> k \& 1)
25
          v = pa[v][k];
26
       if (u == v) return u;
       for(int i = LG - 1; i >= 0; --i)
28
        if(pa[u][i] != pa[v][i]){
29
          u = pa[u][i];
30
          v = pa[v][i];
31
32
      return pa[u][0];
33
```

9.5 Mo on Trees

```
1 int BL[N << 1], ID[N << 1];</pre>
  int lvl[N], par[17][N];
    int ans[N];
    vector<ii> adj[N];
    struct query{
      int id, 1, r, 1c;
      bool operator < (const query & rhs) {</pre>
        return (BL[1] == BL[rhs.1]) ? (r < rhs.r) : (BL[1] < BL[rhs.1]);</pre>
 9
10
    } Q [ N ] ;
11
    int in[N], out[N], val[N], timer;
    void dfs(int node, int p) {
      in[node] = ++timer; ID[timer] = node;
14
      for(int i = 1; i < 17; i++)par[i][node] = par[i-1][par[i-1][node]];</pre>
15
      for(auto child : adj[node])if(child.F != p){
16
        lvl[child.F] = lvl[node] + 1;
17
        par[0][child.F] = node;
18
        val[child.F] = child.S;
19
        dfs(child.F, node);
20
21
      out[node] = ++timer; ID[timer] = node;
22
   int LCA(int u, int v) {
      if(lvl[u] > lvl[v])swap(u,v);
25
      for (int k = 0; k < 17; k++)
        if((lvl[v] - lvl[u]) >> k & 1)
27
         v = par[k][v];
      if(u == v)
29
        return u;
30
      for (int i = 16; i >= 0; --i)
31
        if(par[i][u] != par[i][v])
32
          u = par[i][u], v = par[i][v];
33
      return par[0][u];
34
35 bool vis[N];
    int inSet[N];
37
    void add(int node, int & res){
      if(val[node] > N) return;
39
      if(!vis[node]){
40
        inSet[val[node]]++;
        while(inSet[res])res++;
41
42
43
        inSet[val[node]]--;
44
        if(!inSet[val[node]] && val[node] < res)</pre>
45
          res = val[node];
46
47
      vis[node] ^= 1;
48
49
    //-----Adding Queries-----/
50
51
        int u, v;
        cin >> u >> v; if(lvl[u] > lvl[v])swap(u, v);
52
53
        int lca = LCA(u, v);
        Q[i].id = i;
        Q[i].lc = lca;
```

```
Faculty of Computer and Information Sciences,
Ain Shams University: Too Wrong to Pass Too Correct to Fail
```