Faculty of Computer and Information Sciences, Ain
Shams University: Too Wrong to Pass Too Correct to
Fail

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2021

1	Com	abinatorics 1	
	1.1	Burnside Lemma	
	1.2	Catlan Numbers	
2	Alge	ebra 2	,
-	2.1	Primitive Roots	
	2.2	Discrete Logarithm	
	2.3	Iteration over submasks	
	2.4	Totient function	
	2.4	CRT and EEGCD	
	2.6	FFT	
	2.7	Fibonacci	
	2.8	Gauss Determinant	
	2.9	GAUSS SLAE	•
	2.10	Matrix Inverse	
	2.11	NTT	
	2.12	NTT of KACTL	1
			Ì
3	Data	a Structures 5	
	3.1	2D BIT	
	3.2	2D Sparse table	
	3.3	hillbert Order	
	3.4	Merge Sort Bit with updates	,
	3.5	Mo's	1
	3.6	Mo With Updates	1
	3.7	Ordered Set	1:
	3.8	Persistent Seg Tree	1
	3.9	Sqrt Decomposition	1
	3.10	Treap	1.
	3.11	Wavelet Tree	
			1
4	\mathbf{DP}	10	
	4.1	Dynamic Convex Hull Trick	1
	4.2	Dynamic Connectivety with SegTree	1:
	4.3	Li Chao Tree	1.
	4.4	CHT Line Container	
			2
5	Geor	metry 12	
U	GCO.	inetry 12	-
6	Graj	phs 12	2
U	Graj	piis	
-	N / L - 4	L 10	2
7	Mat	h 12	
	3.7	1.00	2
8	Nun	nberTheory 12	20
_	a		
9	Stri		
	9.1	Aho-Corasick Mostafa	2
	9.2	Aho-Corasick Anany	i
	9.3	KMP Anany	2
	9.4	Manacher Kactl	3
	9.5	Suffix Array Kactl	-
	9.6	Suffix Automaton Anany	3
	9.7	Suffix Automaton Mostafa	
	9.8	Suffix Automaton With Rollback Mostafa	,

```
      10 Trees
      16

      10.1 Centroid Decomposition
      16

      10.2 Dsu On Trees
      16

      10.3 Heavy Light Decomposition (Along with Euler Tour)
      16

      10.4 LCA
      17

      10.5 Mo on Trees
      17
```

1 Combinatorics

1.1 Burnside Lemma

```
1
2  // |Classes|=sum (k ^C(pi)) / |G|
3
4  // C(pi) the number of cycles in the permutation pi
5  // |G| the number of permutations
```

1.2 Catlan Numbers

```
1 const int MOD = \dots
  const int MAX = ....
  int catalan[MAX];
  void init() {
       catalan[0] = catalan[1] = 1;
       for (int i=2; i<=n; i++) {</pre>
           catalan[i] = 0;
           for (int j=0; j < i; j++) {
               catalan[i] += (catalan[j] * catalan[i-j-1]) % MOD;
               if (catalan[i] >= MOD) {
                   catalan[i] -= MOD;
.5 }
  // 1- Number of correct bracket sequence consisting of n opening and n closing
  // 2- The number of rooted full binary trees with n+1 leaves (vertices are not
        numbered).
        A rooted binary tree is full if every vertex has either two children or no
        children.
^{90} // 3- The number of ways to completely parenthesize n+1 factors.
21 // 4- The number of triangulations of a convex polygon with n+2 sides
         (i.e. the number of partitions of polygon into disjoint triangles by using
         the diagonals).
^{13} ^{\prime\prime} /5- The number of ways to connect the 2n points on a circle to form n disjoint
         chords.
^4 // 6- The number of non-isomorphic full binary trees with n internal nodes (i.e.
        nodes having at least one son).
5 // 7- The number of monotonic lattice paths from point (0,0) to point (n,n) in a
         square lattice of size nxn,
        which do not pass above the main diagonal (i.e. connecting (0,0) to (n,n))
   /\!/ 8- Number of permutations of length n that can be stack sorted
        (i.e. it can be shown that the rearrangement is stack sorted if and only
        if
 // there is no such index i<j<k, such that ak<ai<aj ).</pre>
  // 9- The number of non-crossing partitions of a set of n elements.
  // 10- The number of ways to cover the ladder 1..n using n rectangles
  // (The ladder consists of n columns, where ith column has a height i).
```

2 Algebra

2.1 Primitive Roots

```
int powmod (int a, int b, int p) {
2
         int res = 1;
3
         while (b)
            if (b & 1)
                res = int (res * 111 * a % p), --b;
                a = int (a * 111 * a % p), b >>= 1;
8
        return res;
9
10
    int generator (int p) {
11
12
        vector<int> fact;
13
         int phi = p - 1, n = phi;
14
         for (int i = 2; i * i <= n; ++i)</pre>
15
            if (n \% i == 0) {
16
                fact.push_back (i);
17
                 while (n \% i == 0)
18
                    n /= i;
19
20
        if (n > 1)
21
            fact.push_back (n);
22
23
         for (int res = 2; res <= p; ++res) {</pre>
24
             bool ok = true;
25
             for (size_t i = 0; i < fact.size() && ok; ++i)</pre>
26
                ok &= powmod (res, phi / fact[i], p) != 1;
27
             if (ok) return res;
28
29
         return -1;
30
```

2.2 Discrete Logarithm

```
// Returns minimum x for which a \hat{x} \times m = b \times m, a and m are coprime.
    int solve(int a, int b, int m) {
         a %= m, b %= m;
 4
         int n = sqrt(m) + 1;
 5
         int an = 1;
 7
         for (int i = 0; i < n; ++i)
             an = (an * 111 * a) % m;
 Q
10
        unordered_map<int, int> vals;
11
         for (int q = 0, cur = b; q \le n; ++q) {
12
             vals[cur] = q;
13
             cur = (cur * 111 * a) % m;
14
15
16
         for (int p = 1, cur = 1; p \le n; ++p) {
             cur = (cur * 111 * an) % m;
17
             if (vals.count(cur)) {
19
                 int ans = n * p - vals[cur];
20
                 return ans;
21
22
23
         return -1;
\frac{24}{25}
26
    //When a and m are not coprime
    // Returns minimum x for which a ^ x % m = b % m.
    int solve(int a, int b, int m) {
29
        a %= m, b %= m;
30
         int k = 1, add = 0, q;
31
        while ((g = gcd(a, m)) > 1) {
32
             if (b == k)
                 return add;
```

```
if (b % g)
35
                return -1;
36
            b /= g, m /= g, ++add;
37
            k = (k * 111 * a / g) % m;
38
39
40
        int n = sqrt(m) + 1;
41
        int an = 1;
42
        for (int i = 0; i < n; ++i)
43
            an = (an * 111 * a) % m;
44
45
        unordered_map<int, int> vals;
46
        for (int q = 0, cur = b; q \le n; ++q) {
47
            vals[cur] = q;
48
            cur = (cur * 111 * a) % m;
49
50
51
        for (int p = 1, cur = k; p \le n; ++p) {
52
            cur = (cur * 111 * an) % m;
53
            if (vals.count(cur)) {
54
                int ans = n * p - vals[cur] + add;
55
                return ans;
56
57
58
        return -1:
59
```

2.3 Iteration over submasks

```
1 int s = m;
2 while (s > 0) {
3     ... you can use s ...
4     s = (s-1) & m;
5 }
```

2.4 Totient function

```
void phi_1_to_n(int n) {
        vector<int> phi(n + 1);
 3
        phi[0] = 0;
        phi[1] = 1;
        for (int i = 2; i <= n; i++)</pre>
            phi[i] = i;
 8
        for (int i = 2; i <= n; i++) {</pre>
 9
            if (phi[i] == i) {
10
                 for (int j = i; j <= n; j += i)
11
                     phi[j] -= phi[j] / i;
12
13
14
```

2.5 CRT and EEGCD

```
1 11 extended(11 a, 11 b, 11 &x, 11 &y) {
 3
        if(b == 0) {
            x = 1;
 5
            y = 0;
 6
            return a;
        11 x0, y0;
 8
 9
        11 g = extended(b, a % b, x0, y0);
10
        x = y0;
11
        y = x0 - a / b * y0;
12
13
        return g ;
```

```
15 ll de(ll a, ll b, ll c, ll &x, ll &y) {
16
17
         11 g = \text{extended}(abs(a), abs(b), x, y);
18
        if(c % q) return -1;
19
20
        x \star = c / g;
21
        y *= c / g;
22
23
        if(a < 0)x = -x;
24
        if(b < 0)y = -y;
25
         return q;
26
27
    pair<11, 11> CRT(vector<11> r, vector<11> m) {
28
29
        11 r1 = r[0], m1 = m[0];
30
31
         for(int i = 1; i < r.size(); i++) {</pre>
32
33
             11 r2 = r[i], m2 = m[i];
34
             11 x0, y0;
35
             11 q = de(m1, -m2, r2 - r1, x0, y0);
36
37
             if(q == -1) return \{-1, -1\};
38
39
             11 \text{ nr} = x0 * m1 + r1;
40
             11 nm = m1 / g * m2;
42
             r1 = (nr % nm + nm) % nm;
43
             m1 = nm;
44
45
         return {r1, m1};
46
```

2.6 FFT

```
#include < iostream >
    #include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define rep(i, a, b) for(int i = a; i < (b); ++i)
    #define all(x) begin(x), end(x)
    #define sz(x) (int)(x).size()
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    typedef complex<double> C;
11
    typedef vector<double> vd;
    typedef vector<int> vi;
13
    typedef pair<int, int> pii;
    void fft(vector<C>& a) {
15
        int n = sz(a), L = 31 - __builtin_clz(n);
16
        static vector<complex<long double>> R(2, 1);
17
        static vector<C> rt(2, 1); // (^ 10% fas te r i f double)
18
        for (static int k = 2; k < n; k \neq 2) {
19
            R.resize(n):
20
21
            auto x = polar(1.0L, acos(-1.0L) / k);
22
            rep(i, k, 2 * k) rt[i] = R[i] = i & 1 ? R[i / 2] * x : R[i / 2];
23
24
25
        rep(i, 0, n) rev[i] = (rev[i / 2] | (i & 1) << L) / 2;
        rep(i, 0, n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
26
27
        for (int k = 1; k < n; k *= 2)
28
            for (int i = 0; i < n; i += 2 * k) rep(j, 0, k) {
29
                Cz = rt[j + k] * a[i + j + k]; //
30
                a[i + j + k] = a[i + j] - z;
31
                a[i + j] += z;
32
33
34
    vd conv(const vd& a, const vd& b) {
35
        if (a.empty() || b.empty()) return {};
36
        vd res(sz(a) + sz(b) - 1);
        int L = 32 - __builtin_clz(sz(res)), n = 1 << L;</pre>
```

```
vector<C> in(n), out(n);
39
        copy(all(a), begin(in));
40
        rep(i, 0, sz(b)) in[i].imag(b[i]);
41
        fft(in);
42
        for (C& x : in) x *= x;
43
        rep(i, 0, n) out[i] = in[-i & (n - 1)] - conj(in[i]);
44
        fft (out);
45
        rep(i, 0, sz(res)) res[i] = imag(out[i]) / (4 * n);
46
        return res;
47
48
49
    int main() {
50
51
        //Applications
52
        //1-All possible sums
53
54
        //2-All possible scalar products
55
        // We are given two arrays a[] and b[] of length n.
56
        //We have to compute the products of a with every cyclic shift of b.
57
        //We generate two new arrays of size 2n: We reverse a and append n zeros to
58
        //And we just append b to itself. When we multiply these two arrays as
             polynomials,
59
        //and look at the coefficients c[n-1], c[n], ..., c[2n-2] of the product c,
             we aet:
60
        //c[k]=sum i+j=k a[i]b[j]
61
62
        //3-Two stripes
63
        //We are given two Boolean stripes (cyclic arrays of values 0 and 1) a and b
64
        //We want to find all ways to attach the first stripe to the second one,
65
        //such that at no position we have a 1 of the first stripe next to a 1 of
             the second stripe.
66
```

2.7 Fibonacci

2.8 Gauss Determinant

```
1 const double EPS = 1E-9;
    int n:
    vector < vector<double> > a (n, vector<double> (n));
    double det = 1;
    for (int i=0; i<n; ++i) {</pre>
        int k = i;
        for (int j=i+1; j<n; ++j)</pre>
 Q
             if (abs (a[j][i]) > abs (a[k][i]))
10
                 k = i;
11
         if (abs (a[k][i]) < EPS) {
12
             det = 0;
13
             break;
14
15
        swap (a[i], a[k]);
16
        if (i != k)
17
             det = -det;
18
        det *= a[i][i];
19
        for (int j=i+1; j<n; ++j)</pre>
           a[i][j] /= a[i][i];
         for (int j=0; j < n; ++j)
```

2.9 GAUSS SLAE

```
const double EPS = 1e-9;
2
    const int INF = 2; // it doesn't actually have to be infinity or a big number
3
    int gauss (vector < vector<double> > a, vector<double> & ans) {
4
         int n = (int) a.size();
5
         int m = (int) a[0].size() - 1;
6
7
8
         vector<int> where (m, -1);
9
         for (int col = 0, row = 0; col < m && row < n; ++col) {
10
             int sel = row;
11
             for (int i = row; i < n; ++i)
12
                 if (abs (a[i][col]) > abs (a[sel][col]))
13
                     sel = i;
14
             if (abs (a[sel][col]) < EPS)</pre>
15
                 continue;
16
             for (int i = col; i <= m; ++i)</pre>
17
                 swap (a[sel][i], a[row][i]);
18
             where[col] = row;
19
20
             for (int i = 0; i < n; ++i)
21
                 if (i != row) {
22
                      double c = a[i][col] / a[row][col];
23
                      for (int j = col; j <= m; ++j)</pre>
24
                          a[i][j] -= a[row][j] * c;
25
26
             ++row;
27
28
29
         ans.assign (m, 0);
30
         for (int i = 0; i < m; ++i)
31
             if (where[i] != -1)
32
                 ans[i] = a[where[i]][m] / a[where[i]][i];
33
         for (int i = 0; i < n; ++i) {
34
             double sum = 0;
35
             for (int j = 0; j < m; ++j)
    sum += ans[j] * a[i][j];</pre>
36
37
             if (abs (sum - a[i][m]) > EPS)
38
                 return 0;
39
40
41
         for (int i = 0; i < m; ++i)
             if (where [i] == -1)
43
                 return INF;
44
         return 1;
45
```

2.10 Matrix Inverse

```
// Sometimes, the questions are complicated - and the answers are simple. //
pragma GCC optimize ("03")
#pragma GCC optimize ("unroll-loops")
#include <bits/stdc++.h>
#define 11 long long
#define ld long double
#define Io ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
using namespace std;
vector < vector<double> > gauss (vector < vector<double> > a) {

int n = (int) a.size();
vector<vector<double> > ans(n, vector<double>(n, 0));
```

```
14
         for (int i = 0; i < n; i++)
15
             ans[i][i] = 1;
16
         for (int i = 0; i < n; i++) {
17
             for (int j = i + 1; j < n; j++)
18
                 if(a[j][i] > a[i][i]) {
19
                     swap(a[j], a[i]);
20
                     swap(ans[j], ans[i]);
21
22
             double val = a[i][i];
23
             for (int j = 0; j < n; j++) {
24
                 a[i][j] /= val;
25
                 ans[i][j] /= val;
27
             for (int j = 0; j < n; j++) {
                 if(j == i)continue;
28
29
                 val = a[i][i];
30
                 for(int k = 0; k < n; k++) {
    a[j][k] -= val * a[i][k];</pre>
                      ans[j][k] = val * ans[i][k];
33
34
35
36
        return ans;
37
38
    int main() {
39
40
41
         vector<vector<double> > v(3, vector<double> (3) );
42
         for (int i = 0; i < 3; i++)
43
             for (int j = 0; j < 3; j++)
44
                 cin >> v[i][j];
45
46
        for(auto i : gauss(v)) {
47
             for(auto j : i)
                 cout << j << " ";
48
49
             cout << "\n";
50
51
```

2.11 NTT

```
1
    struct NTT {
        int mod ;
 3
        int root ;
        int root 1 ;
        int root_pw ;
        NTT(int _mod, int primtive_root, int NTT_Len) {
 9
            mod = mod;
10
            root_pw = NTT_Len;
11
             root = fastpower(primtive_root, (mod - 1) / root_pw);
12
             root_1 = fastpower(root, mod - 2);
13
14
        void fft(vector<int> & a, bool invert) {
15
            int n = a.size();
16
17
             for (int i = 1, j = 0; i < n; i++) {
                int bit = n >> 1;
18
19
                for (; j & bit; bit >>= 1)
20
                    j ^= bit;
                j ^= bit;
21
22
23
                if (i < j)
                    swap(a[i], a[j]);
25
26
27
            for (int len = 2; len <= n; len <<= 1) {</pre>
                int wlen = invert ? root_1 : root;
                for (int i = len; i < root_pw; i <<= 1)</pre>
                     wlen = (int)(1LL * wlen * wlen % mod);
```

```
32
33
                 for (int i = 0; i < n; i += len) {</pre>
34
                     int w = 1;
35
                     for (int j = 0; j < len / 2; j++) {
36
                         int u = a[i + j], v = (int)(1LL * a[i + j + len / 2] * w %
                              mod);
                         a[i + j] = u + v < mod ? u + v : u + v - mod;
                         a[i + j + len / 2] = u - v >= 0 ? u - v : u - v + mod;
39
                         w = (int) (1LL * w * wlen % mod);
40
41
42
43
44
             if (invert) {
45
                 int n_1 = fastpower(n, mod - 2);
46
                 for (int & x : a)
47
                    x = (int) (1LL * x * n_1 % mod);
48
49
50
         vector<int> multiply(vector<int> &a, vector<int> &b) {
51
             vector<int> fa(a.begin(), a.end()), fb(b.begin(), b.end());
52
             int n = 1:
53
             while(n < a.size() + b.size())</pre>
54
                n <<= 1;
55
56
             fa.resize(n);
57
             fb.resize(n);
58
59
             fft(fa, 0);
60
             fft(fb, 0);
61
             for (int i = 0; i < n; i++)
62
63
                fa[i] = 1LL * fa[i] * fb[i] % mod;
64
             fft(fa, 1);
65
             return fa;
66
67 };
```

2.12 NTT of KACTL

```
1 ///(Note faster than the other NTT)
    ///If the mod changes don't forget to calculate the primitive root
    using 11 = long long;
    const 11 mod = (119 << 23) + 1, root = 3; // = 998244353</pre>
    // For p < 2^30 there is also e.g. 5 << 25, 7 << 26, 479 << 21
    // and 483 << 21 (same root). The last two are > 10^9.
    typedef vector<11> v1;
    11 modpow(ll b, ll e) {
10
        11 \text{ ans} = 1;
11
         for (; e; b = b * b % mod, e /= 2)
12
            if (e & 1) ans = ans * b % mod;
13
         return ans:
14
15
    void ntt(vl &a) {
16
        int n = sz(a), L = 31 - \underline{builtin_clz(n)};
17
         static vl rt(2, 1);
18
         for (static int k = 2, s = 2; k < n; k *= 2, s++) {
19
             rt.resize(n);
20
             ll z[] = \{1, modpow(root, mod >> s)\};
21
             f(i,k,2*k) rt[i] = rt[i / 2] * z[i & 1] % mod;
22
23
         vector<int> rev(n);
24
         f(i,0,n) \text{ rev}[i] = (\text{rev}[i / 2] | (i \& 1) << L) / 2;
         f(i,0,n) if (i < rev[i]) swap(a[i], a[rev[i]]);</pre>
26
         for (int k = 1; k < n; k *= 2)
27
             for (int i = 0; i < n; i += 2 * k) f(j, 0, k) {
                 11 z = rt[j + k] * a[i + j + k] % mod, &ai = a[i + j];
28
29
                 a[i + j + k] = ai - z + (z > ai ? mod : 0);
30
                 ai += (ai + z >= mod ? z - mod : z);
```

```
32
33
   vl conv(const vl &a, const vl &b) {
        if (a.empty() || b.empty()) return {};
        int s = sz(a) + sz(b) - 1, B = 32 - _builtin_clz(s), n = 1 << B;
35
        int inv = modpow(n, mod - 2);
37
        vl L(a), R(b), out(n);
38
        L.resize(n), R.resize(n);
39
        ntt(L), ntt(R);
40
        f(i,0,n) out [-i \& (n-1)] = (11)L[i] * R[i] % mod * inv % mod;
41
42
        return {out.begin(), out.begin() + s};
43
44
    vector<int> v;
   vector<ll> solve(int s, int e) {
        if(s==e) {
            vector<ll> res(2);
47
48
            res[0] = 1;
49
            res[1] = v[s];
50
            return res;
51
52
        int md = (s + e) >> 1;
53
        return conv(solve(s, md), solve(md+1, e));
```

3 Data Structures

3.1 2D BIT

```
1 void upd(int x, int y, int val) {
2    for(int i = x; i <= n; i += i & -i)
3    for(int j = y; j <= m; j += j & -j)
4    bit[i][j] += val;
5    }
6    int get(int x, int y) {
7       int ans = 0;
8       for(int i = x; i; i -= i & -i)
9       for(int j = y; j; j -= j & -j)
10       ans += bit[i][j];
11 }</pre>
```

3.2 2D Sparse table

```
1 /*
        note this isn't the best cache-wise version
 3
        query O(1), Build O(NMlgNlgM)
        be careful when using it and note the he build a dimension above another
        i.e he builds a sparse table for each row
        the build sparse table over each row's sparse table
    const int N = 505, LG = 10;
10
   int st[N][N][LG][LG];
    int a[N][N], lg2[N];
12
13
    int yo(int x1, int y1, int x2, int y2) {
14
      x2++;
15
      y2++;
16
      int a = \lg 2[x2 - x1], b = \lg 2[y2 - y1];
17
      return max (
18
             \max(st[x1][y1][a][b], st[x2 - (1 << a)][y1][a][b]),
19
             \max(st[x1][y2 - (1 << b)][a][b], st[x2 - (1 << a)][y2 - (1 << b)][a][b]
20
           );
21
    void build(int n, int m) { // 0 indexed
      for (int i = 2; i < N; i++) lg2[i] = lg2[i >> 1] + 1;
```

```
for (int i = 0; i < n; i++) {</pre>
26
         for (int j = 0; j < m; j++) {
27
           st[i][j][0][0] = a[i][j];
28
29
30
      for (int a = 0; a < LG; a++) {</pre>
31
         for (int b = 0; b < LG; b++) {</pre>
          if (a + b == 0) continue;
32
33
           for (int i = 0; i + (1 << a) <= n; i++) {
34
             for (int j = 0; j + (1 << b) <= m; <math>j++) {
35
               if (!a) {
36
                 st[i][j][a][b] = max(st[i][j][a][b - 1], st[i][j + (1 << (b - 1))][a]
37
38
                 st[i][j][a][b] = max(st[i][j][a - 1][b], st[i + (1 << (a - 1))][j][a]
41
42
43
44
```

3.3 hillbert Order

```
///Faster Sorting MO
 2
 3
    const int infinity = (int)1e9 + 42;
    const int64_t llInfinity = (int64_t)1e18 + 256;
    const int module = (int)1e9 + 7;
    const long double eps = 1e-8;
    inline int64_t gilbertOrder(int x, int y, int pow, int rotate) {
 8
 9
10
            return 0;
11
12
         int hpow = 1 << (pow-1);
13
        int seq = (x < hpow) ? (
14
             (y < hpow) ? 0 : 3
15
         ) : (
             (y < hpow) ? 1 : 2
16
17
18
        seg = (seg + rotate) & 3;
19
         const int rotateDelta[4] = {3, 0, 0, 1};
        int nx = x & (x ^ hpow), ny = y & (y ^ hpow);
20
21
        int nrot = (rotate + rotateDelta[seq]) & 3;
22
        int64_t subSquareSize = int64_t(1) << (2*pow - 2);
23
        int64_t ans = seg * subSquareSize;
         int64_t add = gilbertOrder(nx, ny, pow-1, nrot);
24
25
         ans += (seg == 1 || seg == 2) ? add : (subSquareSize - add - 1);
26
        return ans;
27
28
29
    struct Query {
30
        int 1, r, idx;
31
        int64_t ord;
32
33
        inline void calcOrder() {
34
            ord = gilbertOrder(1, r, 21, 0);
35
36
37
38
    inline bool operator<(const Query &a, const Query &b) {</pre>
39
        return a.ord < b.ord;</pre>
40
41
    signed main() {
43
         #ifndef USE_FILE_IO
44
             ios_base::sync_with_stdio(false);
45
         #endif
46
        mt19937 rnd(42);
```

```
49
         int n, m, k; cin >> n >> m; k = rnd() % 1048576;
50
         vector<int> p(n+1);
         for (int i = 0; i < n; i++) {
52
             int val = rnd() % 1048576;
53
             p[i+1] = p[i] ^ val;
54
55
56
         vector<Query> gry(m);
57
         for (int i = 0; i < m; i++) {
             int 1 = rnd() % n + 1, r = rnd() % n + 1;
             if (1 > r) {
                 swap(l, r);
             qry[i].1 = 1; qry[i].r = r;
             qry[i].idx = i;
             gry[i].calcOrder();
65
66
67
         int64_t ans = 0;
68
         vector<int64_t> res(m);
69
         vector<int64_t> cnt((int)2e6, 0);
70
         sort(qry.begin(), qry.end());
 71
         int 1 = 0, r = 1;
72
         ans = (p[1] == k);
73
         cnt[p[0]]++; cnt[p[1]]++;
74
75
         for (Query q: qry) {
76
77
             while (1 > q.1) {
78
                 1--;
79
                 ans += cnt[p[1] ^ k];
80
                 cnt[p[1]]++;
81
82
             while (r < q.r) {
83
                 <u>r</u>++;
84
                 ans += cnt[p[r] ^{\circ} k];
85
                 cnt[p[r]]++;
86
87
             while (1 < q.1) {
                 cnt[p[1]]--;
89
                 ans -= cnt[p[1] ^ k];
91
92
             while (r > q.r) {
93
                 cnt[p[r]]--;
94
                 ans -= cnt[p[r] ^ k];
                 r--;
96
97
             res[q.idx] = ans;
98
99
100
         uint64_t rhsh = 0;
101
         for (int i = 0; i < m; i++) {
102
             rhsh *= (uint64_t)1e9 + 7;
103
             rhsh += (uint64_t)res[i];
104
         cout << rhsh << "\n";</pre>
105
106
107
         return 0;
108
```

3.4 Merge Sort Bit with updates

```
//O(log ^ 2 N) updates and queries

#include <ext/pb_ds/tree_policy.hpp>
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/rope>

using namespace std;
```

7

```
using namespace __gnu_pbds;
10
    using namespace __gnu_cxx;
11
12
    template<class T> using Tree = tree<T, null_type, less<T>, rb_tree_tag,
         tree_order_statistics_node_update>;
13
14
    Tree<int> t[N];
15
16
17
    void add(int idx, int v) {
18
        for (int x = ++idx; x < N; x += x & -x) {
19
            t[x].insert(v);
20
21
22
    void erase(int idx, int v) {
23
        for (int x = ++idx; x < N; x += x & -x)
24
            t[x].erase(v);
25
26
    int get(int idx, int limit) {
27
        int ret = 0;
28
        for (int x = ++idx; x; x -= x & -x)
29
            ret += (t[x].order_of_key(limit+1));
30
        return ret;
31
```

3.5 Mo's

```
#include <bits/stdc++.h>
2
3
    int n, qq, arr[N], sz = 1000; // sz is the size of the bucket
    int co[N], ans = 0, ansq[N];
    int cul = 1, cur = 1;
7
    void add(int x) {
        co[arr[x]]++;
8
Q
         if (co[arr[x]] == 1)
10
             ans++;
11
         else if (co[arr[x]] == 2)
12
            ans--;
13
14
15
    void remove(int x) {
        co[arr[x]]--;
17
        if (co[arr[x]] == 1)
18
            ans++:
19
         else if (co[arr[x]] == 0)
20
            ans--;
21
22
    void solve(int 1, int r, int ind) {
24
25
         while (cul < 1) remove(cul++);</pre>
26
         while (cul > 1) add(--cul);
27
        while (cur < r) add(cur++);</pre>
28
        while (cur > r) remove(--cur);
29
         ansq[ind] = ans;
30
31
32
    int main() {
34
        FIO
35
        cin >> qq;
36
                                  \{1/sz,r\},
                                                { 1 , ind}
37
       priority_queue<pair<int, int>, pair<int, int>>, vector<pair<jair<int,</pre>
            int>, pair<int, int>>>, greater<pair<pair<int, int>, pair<int, int>>>> q
38
         for (int i = 0; i < qq; i++) {</pre>
39
             int 1, r;
40
             cin >> 1 >> r;
41
             q.push(\{\{1 / sz, r\}, \{1,i\}\});
42
         while (q.size()) {
```

```
44
             int ind=q.top().second.second,l=q.top().second.first,r=q.top().first.
45
             solve(1, r,ind);
46
             q.pop();
47
48
         for (int i = 0; i < qq; i++)
49
             cout << ansq[i] << endl;</pre>
50
51
52
        return 0;
53
```

3.6 Mo With Updates

```
///O(N^5/3) note that the block size is not a standard size
 3
    #pragma GCC optimize ("03")
    #pragma GCC target ("sse4")
    #include <bits/stdc++.h>
 9
    using namespace std;
10
11
    using 11 = long long;
    const int N = 1e5 +5;
13
    const int M = 2 * N;
   const int blk = 2155;
    const int mod = 1e9 + 7;
    struct Ouerv{
18
     int 1, r, t, idx;
19
      Query (int a = 0, int b = 0, int c = 0, int d = 0) {l=a, r=b, t=c, idx = d;}
20
      bool operator < (Query o) {</pre>
        if(r / blk == o.r / blk && 1 / blk == o.l / blk)return t < o.t;</pre>
21
        if(r / blk == o.r / blk)return 1 < o.1;</pre>
23
        return r < o.r;
24
25
    } Q[N];
27 int a[N], b[N];
28 int cnt1[M], cnt2[N];
   int L = 0, R = -1, K = -1;
    void add(int x) { ///add item to range
31
   // cout << x << '\n';
     cnt2[cnt1[x]]--;
33
      cnt1[x]++;
34
      cnt2[cnt1[x]]++;
35
36 void del(int x) { ///delete item from range
    cnt2[cnt1[x]]--;
38
      cnt1[x]--;
39
      cnt2[cnt1[x]]++;
40
41 map<int,int>id;
42 int cnt;
43 int ans[N];
    int p[N], nxt[N];
44
45
    int prv[N];
46
    void upd(int idx) { ///update item value
      if(p[idx] >= L && p[idx] <= R)
48
        del(a[p[idx]]), add(nxt[idx]);
49
      a[p[idx]] = nxt[idx];
50
51
    void err(int idx) {
      if(p[idx] >= L \&\& p[idx] <= R)
        del(a[p[idx]]), add(prv[idx]);
54
      a[p[idx]] = prv[idx];
55
56
    int main(){
57
      int n, q, 1, r, tp;
```

```
60
       scanf("%d%d", &n, &q);
61
62
       for (int i = 0; i < n; i++) {
63
         scanf("%d", a + i);
 64
         if(id.count(a[i]) == 0)
65
           id[a[i]] = cnt++;
         a[i] = id[a[i]];
66
         b[i] = a[i];
68
69
       int qIdx = 0;
70
       int ord = 0;
71
       while (q--) {
72
73
74
         scanf("%d", &tp);
         if(tp == 1) {
75
           /// ADD Query
76
           scanf("%d%d", &1, &r); --1, --r;
 77
           Q[qIdx] = Query(1,r,ord-1,qIdx); qIdx++;
78
 79
           /// ADD Update
80
           scanf("%d%d",p + ord, nxt + ord); --p[ord];
81
           if(id.count(nxt[ord]) == 0)
 82
              id[nxt[ord]] = cnt++;
 83
           nxt[ord] = id[nxt[ord]];
 84
           prv[ord] = b[p[ord]];
 85
           b[p[ord]] = nxt[ord];
 86
           ++ord;
 87
 88
89
90
       sort(Q,Q+qIdx);
       for(int i = 0; i < qIdx; i++) {</pre>
92
         while (L < Q[i].l) del(a[L++]);
93
         while (L > Q[i].1) add (a[--L]);
94
         while (R < Q[i].r) add (a[++R]);
95
         while (R > Q[i].r) del(a[R--]);
96
         while (K < Q[i].t) upd (++K);
97
         while (K > Q[i].t) err(K--);
98
          ///Solve Query I
99
100
       for(int i = 0; i < qIdx; i++)</pre>
101
         printf("%d\n", ans[i]);
102
103
104
       return 0;
105
```

3.7 Ordered Set

3.8 Persistent Seg Tree

```
1
2  int val[ N * 60 ], L[ N * 60 ], R[ N * 60 ], ptr, tree[N]; /// N * lgN
3  int upd(int root, int s, int e, int idx) {
4    int ret = ++ptr;
5    val[ret] = L[ret] = R[ret] = 0;
```

```
if (s == e) {
            val[ret] = val[root] + 1;
 8
            return ret;
 9
10
        int md = (s + e) \gg 1;
11
        if (idx <= md) {
12
            L[ret] = upd(L[root], s, md, idx), R[ret] = R[root];
13
         } else {
14
            R[ret] = upd(R[root], md + 1, e, idx), L[ret] = L[root];
15
16
        val[ret] = max(val[L[ret]], val[R[ret]]);
17
        return ret;
18
19
   int qry(int node, int s, int e, int 1, int r){
      if(r < s || e < 1 || !node)return 0; //Punishment Value</pre>
21
      if(1 <= s && e <= r) {
22
        return val[node];
23
      int md = (s+e) >> 1;
25
      return max(qry(L[node], s, md, l, r), qry(R[node], md+1, e, l, r));
26
27
    int merge(int x, int y, int s, int e) {
28
        if(!x||!y) return x | y;
29
        if(s == e) {
30
            val[x] += val[y];
31
            return x;
32
33
        int md = (s + e) \gg 1;
        L[x] = merge(L[x], L[y], s, md);
        R[x] = merge(R[x], R[y], md+1,e);
36
        val[x] = val[L[x]] + val[R[x]];
37
        return x;
38
```

3.9 Sqrt Decomposition

```
1 // Source: https://cp-algorithms.com/data_structures/sqrt_decomposition.html
 3
    // input data
    int n;
    vector<int> a (n);
    // preprocessing
    int len = (int) sqrt (n + .0) + 1; // size of the block and the number of blocks
    vector<int> b (len);
    for (int i=0; i<n; ++i)</pre>
11
        b[i / len] += a[i];
12
13
    // answering the queries
14
    for (;;) {
        int 1, r;
16
      // read input data for the next query
17
        int sum = 0;
18
        for (int i=1; i<=r; )</pre>
19
            if (i % len == 0 && i + len - 1 <= r) {</pre>
                // if the whole block starting at i belongs to [1, r]
                 sum += b[i / len];
                 i += len;
            else {
25
                 sum += a[i];
26
                 ++i;
27
28
29
    // If you're getting TLE and can't optimize more, you could reduce the number of
          slow division operations using the following code:
    int sum = 0;
    int c_l = 1 / len,     c_r = r / len;
33
    if (c_l == c_r)
        for (int i=1; i<=r; ++i)</pre>
```

```
sum += a[i];
37
    else {
38
         for (int i=1, end=(c_1+1)*len-1; i<=end; ++i)</pre>
39
              sum += a[i];
40
         for (int i=c_l+1; i<=c_r-1; ++i)</pre>
41
             sum += b[i]:
42
         for (int i=c_r*len; i<=r; ++i)</pre>
43
             sum += a[i];
44
```

3.10 Treap

```
typedef struct item * pitem;
    struct item {
        int prior, value, cnt;
        bool rev;
 -5
        pitem 1, r;
 6
         item(int x, int y, int z){
            value = x;
            prior = y;
 8
            cnt = z;
10
             rev = 0;
11
             1 = r = NULL:
12
13
    };
14
15
    int cnt (pitem it) {
16
        return it ? it->cnt : 0;
17
18
19 void upd_cnt (pitem it) {
20
        if (it)
21
             it->cnt = cnt(it->1) + cnt(it->r) + 1;
22
23
^{24}
    void push (pitem it) {
25
        if (it && it->rev) {
26
             it->rev = false;
27
             swap (it->1, it->r);
             if (it->1) it->1->rev ^= true;
28
29
             if (it->r) it->r->rev ^= true;
30
31
    }
32
33
    void merge (pitem & t, pitem l, pitem r) {
34
        push (1);
35
        push (r);
36
        if (!1 || !r)
37
            t = 1 ? 1 : r;
         else if (l->prior > r->prior)
38
39
            merge (1->r, 1->r, r), t = 1;
40
41
            merge (r->1, 1, r->1), t = r;
42
        upd_cnt (t);
43
44
45
    void split (pitem t, pitem & 1, pitem & r, int key, int add = 0) {
46
47
            return void( 1 = r = 0 );
48
         push (t);
         int cur_key = add + cnt(t->1);
50
         if (key <= cur_key)</pre>
51
            split (t->1, 1, t->1, key, add), r = t;
52
53
            split (t->r, t->r, r, key, add + 1 + cnt(t->1)), 1 = t;
54
        upd_cnt (t);
55
56
57
    void reverse (pitem t, int l, int r) {
58
        pitem t1, t2, t3;
59
         split (t, t1, t2, 1);
60
         split (t2, t2, t3, r-l+1);
```

```
61
        t2->rev ^= true;
62
        merge (t, t1, t2);
63
        merge (t, t, t3);
64
65
66
   void output (pitem t) {
67
        if (!t) return;
        push (t);
69
        output (t->1);
70
        printf ("%c", char(t->value));
71
        output (t->r);
72
73
74
    pitem gettreap(string s) {
75
            pitem ret=NULL;
76
77
           for(i=0;i<s.size();i++)merge(ret,ret,new item(s[i],(rand()<<15)+rand(),</pre>
                1));
78
        return ret;
79
```

3.11 Wavelet Tree

```
// remember your array and values must be 1-based
   struct wavelet tree {
        int lo, hi;
        wavelet_tree *1, *r;
 5
        vector<int> b;
 7
        //nos are in range [x,y]
        //array indices are [from, to)
 9
        wavelet_tree(int *from, int *to, int x, int y) {
10
            lo = x, hi = y;
11
            if (lo == hi or from >= to)
12
                return:
13
             int mid = (lo + hi) / 2;
14
            auto f = [mid] (int x) {
                return x <= mid;</pre>
15
16
17
            b.reserve(to - from + 1);
18
            b.pb(0);
19
            for (auto it = from; it != to; it++)
20
                b.pb(b.back() + f(*it));
21
             //see how lambda function is used here
22
            auto pivot = stable_partition(from, to, f);
23
             l = new wavelet_tree(from, pivot, lo, mid);
24
             r = new wavelet_tree(pivot, to, mid + 1, hi);
25
26
27
        //kth smallest element in [1, r]
28
        int kth(int 1, int r, int k) {
29
            if (1 > r)
30
                return 0;
31
            if (lo == hi)
32
                return lo;
33
            int inLeft = b[r] - b[1 - 1];
34
            int lb = b[1 - 1]; //amt of nos in first (1-1) nos that go in left
35
            int rb = b[r]; //amt of nos in first (r) nos that go in left
36
             if (k <= inLeft)</pre>
37
                return this->l->kth(lb + 1, rb, k);
            return this->r->kth(1 - 1b, r - rb, k - inLeft);
38
39
40
41
         //count of nos in [1, r] Less than or equal to k
42
        int LTE(int 1, int r, int k) {
43
            if (1 > r \text{ or } k < 10)
44
                return 0;
45
            if (hi <= k)
46
                return r - 1 + 1;
47
            int 1b = b[1 - 1], rb = b[r];
48
            return this->1->LTE(lb + 1, rb, k) + this->r->LTE(l - lb, r - rb, k);
49
```

```
__
```

```
51
         //count of nos in [1, r] equal to k
52
         int count(int 1, int r, int k) {
53
             if (1 > r \text{ or } k < 10 \text{ or } k > hi)
54
                 return 0:
55
             if (lo == hi)
56
                 return r - 1 + 1;
             int 1b = b[1 - 1], rb = b[r], mid = (10 + hi) / 2;
57
             if (k <= mid)
59
                 return this->l->count(lb + 1, rb, k);
60
             return this->r->count(1 - 1b, r - rb, k);
61
62
    };
```

4 DP

#include<iostream>

4.1 Dynamic Convex Hull Trick

```
#include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    struct Line
9
         11 m, b;
10
        mutable function<const Line*()> succ;
11
        bool operator<(const Line& other) const
12
13
             return m < other.m;</pre>
14
15
        bool operator<(const 11 &x) const
16
17
             const Line* s = succ();
18
             if (!s)
19
                 return 0;
20
             return b - s \rightarrow b < (s \rightarrow m - m) * x;
21
    // will maintain upper hull for maximum
^{24}
    struct HullDynamic : public multiset<Line, less<>>
25
26
        bool bad(iterator y)
27
28
             auto z = next(y);
29
             if (y == begin())
30
31
                 if (z == end())
32
                     return 0;
33
                 return y->m == z->m && y->b <= z->b;
34
35
             auto x = prev(y);
36
             if (z == end())
37
                 return y->m == x->m && y->b <= x->b;
38
             return (ld) (x-b - y-b)*(z-m - y-m) >= (ld) (y-b - z-b)*(y-m - x-m)
39
40
        void insert_line(ll m, ll b)
41
42
             auto y = insert({ m, b });
43
             y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
44
             if (bad(y))
45
46
                 erase(y);
47
49
             while (next(y) != end() && bad(next(y)))
                 erase(next(y));
             while (y != begin() && bad(prev(y)))
```

```
erase(prev(y));
53
55
        11 query(11 x)
56
57
58
            auto 1 = *lower_bound(x);
            return 1.m * x + 1.b;
60
61
    };
62
    int main()
65
66
```

4.2 Dynamic Connectivety with SegTree

```
#pragma GCC optimize("03")
    #pragma GCC optimize ("unroll-loops")
    #pragma GCC target("avx,avx2,fma")
    using namespace std;
    #include "bits/stdc++.h"
    #define pb push_back
    #define F first
    #define S second
    #define f(i, a, b) for(int i = a; i < b; i++)
    #define all(a) a.begin(),a.end()
    #define rall(a) a.rbegin(),a.rend()
    #define sz(x) (int)(x).size()
    //#define mp make_pair
    #define popCnt(x) (__builtin_popcountll(x))
17
    typedef long long 11;
    typedef pair<int, int> ii;
    using ull = unsigned long long;
    const int N = 1e5+5, LG = 17, MOD = 1e9 + 7;
    const long double PI = acos(-1);
    struct PT{
24
        11 x, y;
25
        PT() {}
26
        PT(11 a, 11 b):x(a), y(b) {}
27
        PT operator - (const PT & o) {return PT{x-o.x,y-o.y};}
28
        bool operator < (const PT & o) const {return make_pair(x,y) < make_pair(o.x,</pre>
             o.y);}
    ll cross(PT x, PT y) {
        return x.x * y.y - x.y * y.x;
33
   PT val[300005];
34 bool in[300005];
   ll qr[300005];
    bool ask[300005];
    11 ans[N];
    vector<PT> t[300005 * 4]; ///segment tree holding points to queries
    void update(int node, int s, int e, int 1, int r, PT x) {
        if(r < s || e < 1) return;
41
        if(1 \le s \&\& e \le r)  { ///add this point to maximize it with queries in
             this range
42
            t[node].pb(x);
43
            return;
44
45
        int md = (s + e) \gg 1;
46
        update(node<<1, s, md, 1, r, x);
47
        update (node <<1|1, md+1, e, l, r, x);
48
49
   vector<PT> stk;
50 inline void addPts(vector<PT> v) {
51
        stk.clear();
                        ///reset the data structure you are using
        sort(all(v));
```

```
53
          ///build upper envelope
54
         for(int i = 0; i < v.size(); i++) {</pre>
55
             while(sz(stk) > 1 && cross(v[i] - stk.back(), stk.back() - stk[stk.size
                   ()-2]) <= 0)
56
                  stk.pop_back();
57
             stk.push_back(v[i]);
58
59
60
     inline 11 calc(PT x, 11 val) {
61
         ///mb+v
62
         return x.x * val + x.y;
63
64
65
    ll query(ll x) {
66
         if(stk.empty())
67
             return LLONG_MIN;
         int lo = 0, hi = stk.size() - 1;
68
69
          while(lo + 10 < hi) {
70
             int md = lo + (hi-lo) / 2;
71
             if(calc(stk[md+1],x) > calc(stk[md],x))
 72
                 lo = md + 1;
73
             else
74
                 hi = md;
75
76
         11 ans = LLONG MIN:
77
         for (int i = lo; i <= hi; i++)</pre>
78
             ans = max(ans, calc(stk[i], x));
 79
          return ans;
 80
81
     void solve(int node, int s, int e) {      ///Solve queries
82
         addPts(t[node]);
                              ///note that there is no need to add/delete just build
              for t[node]
 83
         f(i,s,e+1){
 84
             if(ask[i]) {
 85
                 ans[i] = max(ans[i], query(qr[i]));
 86
87
 88
         if(s==e)return;
 89
         int md = (s + e) \gg 1;
90
         solve(node<<1,s,md);</pre>
91
         solve(node<<1|1,md+1,e);
92
93
     void doWork() {
94
95
         int n;
96
         cin >> n;
97
         stk.reserve(n);
98
         f(i,1,n+1) {
99
             int tp;
100
             cin >> tp;
101
             if(tp == 1) {
                             ///Add Query
                 int x, y;
102
103
                 cin >> x >> y;
104
                 val[i] = PT(x, y);
105
                 in[i] = 1;
106
                 else if(tp == 2) { ///Delete Query
107
                  int x;
108
                  cin >> x;
109
                  if(in[x])update(1, 1, n, x, i - 1, val[x]);
110
                  in[x] = 0;
111
                 else {
112
                  cin >> qr[i];
113
                  ask[i] = true;
114
115
         f(i,1,n+1) ///Finalize Query
116
117
             if(in[i])
                 update(1, 1, n, i, n, val[i]);
118
119
120
         f(i,1,n+1) ans [i] = LLONG_MIN;
121
         solve(1, 1, n);
122
         f(i,1,n+1)
123
         if(ask[i]) {
124
             if(ans[i] == LLONG_MIN)
```

```
125
                  cout << "EMPTY SET\n";</pre>
126
              else
127
                  cout << ans[i] << '\n';
128
129
130
     int32_t main() {
131
     #ifdef ONLINE_JUDGE
132
133
         ios_base::sync_with_stdio(0);
134
          cin.tie(0);
135
     #endif // ONLINE JUDGE
136
         int t = 1;
137
          cin >> t;
138
         while (t--) {
139
              doWork();
140
141
         return 0;
142
```

4.3 Li Chao Tree

```
#include<iostream>
    #include <bits/stdc++.h>
    #define 11 long long
    #define ld long double
    #define IO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    struct Line
 9
10
        Line(11 \, m, 11 \, b) : m(m), b(b) {}
11
         11 operator()(11 x)
12
13
             return m * x + b;
14
15
    };
16
    struct node
17
18
         node * left, * right ;
19
        Line line ;
20
        node(node * left, node *right, Line line):left(left), right(right), line(
             line) {}
21
        node * getLeft()
22
             if(left==NULL)
23
24
                 left= new node (NULL, NULL, Line(0, 1e18));
25
             return left ;
26
27
        node * getright()
28
29
             if(right==NULL)
30
                 right= new node (NULL, NULL, Line (0, 1e18));
31
             return right ;
32
33
        void insert(Line newline, int 1, int r)
34
35
             int m = (1+r)/2;
36
             bool lef=newline(1)<line(1);</pre>
37
             bool mid=newline(m) <line(m);</pre>
38
39
             if (mid)
40
                 swap(line, newline);
41
             if(r-1==1)
42
                 return ;
43
             else if(lef!=mid)
44
                 getLeft()->insert(newline,1,m);
45
             else
46
                 getright()->insert(newline,m,r);
47
48
        11 query(int x, int 1, int r)
49
50
             int m = (1 + r) / 2;
```

```
15
```

```
if(r - 1 == 1)
52
                 return line(x);
53
             else if (x < m)
                 return min(line(x), getLeft()->query(x, 1, m));
55
56
                 return min(line(x), getright()->query(x, m, r));
57
58
        void deletee()
59
             if(left!=NULL)
60
61
                 left->deletee();
62
             if(right!=NULL)
63
                 right->deletee();
             free(this);
65
66
    };
67
    int main()
68
69
70
        node * root = new node(NULL, NULL, Line(0,5));
71
         root->insert(Line(1,-3),1,100);
72
73
         for (int i=1; i<=10; i++)</pre>
74
             cout << root -> query (i, 1, 100) << "\n";
75
```

4.4 CHT Line Container

```
1
    struct Line
 2
 3
         mutable 11 m, b, p;
 4
         bool operator < (const Line& o) const
 5
 6
             return m < o.m;</pre>
 7
 8
         bool operator<(11 x) const
 Q
10
             return p < x;
11
12
13
14
    struct LineContainer : multiset<Line, less<>>
15
16
         // (for doubles, use inf = 1/.0, div(a,b) = a/b)
         static const 11 inf = LLONG_MAX;
17
         ll div(ll db, ll dm) // floored division
18
19
             return db / dm - ((db ^ dm) < 0 && db % dm);</pre>
20
21
22
         bool isect(iterator x, iterator y)
23
24
             if (y == end())
\frac{25}{26}
                 x->p = inf;
27
                 return false;
29
             if (x->m == y->m)
30
                 x->p = x->b > y->b ? inf : -inf;
31
32
                 x->p = div(y->b - x->b, x->m - y->m);
\frac{33}{34}
             return \bar{x}->p >= y->p;
35
         void add(ll m, ll b)
36
37
             auto z = insert(\{m, b, 0\}), y = z++, x = y;
             while (isect(y, z))
39
                 z = erase(z);
40
             if (x != begin() && isect(--x, y))
41
                 isect(x, y = erase(y));
42
             while ((y = x) != begin() && (--x)->p >= y->p)
43
                 isect(x, erase(y));
```

5 Geometry

6 Graphs

7 Math

7.1 Xor With Gauss

```
1 /*
 2
        Some applications
 3
        If you want to find the maximum in xor subset
        just ans = max(ans, ans ^p[i]) for all i
        if you want to count the number of subsets with a certain value
        check all different subsets of p
   ll p[66];
 9
   bool add(ll x) {
        for(int i = 60; (~i) && x; --i) {
10
            if(x >> i & 1) {
11
                if(!p[i]) {
12
13
                    p[i] = x;
14
                    return true;
15
                    else {
16
                    x ^= p[i];
17
18
19
20
        return false;
21
```

7.2 Josephus

```
1 // n = total person
   // will kill every kth person, if k = 2, 2, 4, 6, ...
   // returns the mth killed person
    11 josephus(11 n, 11 k, 11 m) {
     m = n - m;
      if (k <= 1) return n - m;</pre>
      11 i = m;
      while (i < n) {
        11 r = (i - m + k - 2) / (k - 1);
10
        if ((i + r) > n) r = n - i;
11
        else if (!r) r = 1;
        i += r;
13
        m = (m + (r * k)) % i;
14
      } return m + 1;
15
```

7.3 Matrix Power/Multiplication

```
1  struct Matrix {
2
3      const static int D = 100;
4      int a[D][D];
```

```
2
3
3
3
3
3
```

```
6
        Matrix(int val) {
             for (int i = 0; i < D; i++)
                 for (int j = 0; j < D; j++)
9
                    a[i][j] = val;
10
11
12
            memset(a, 0, sizeof a);
13
        void initIdentity() {
14
15
            clear();
16
             for (int i = 0; i < D; i++)
17
                a[i][i] = 1;
18
19
        int * operator [](int r) {
20
            return a[r];
21
22
        const int * operator [](int r) const{
23
            return a[r];
24
25
26
        friend Matrix operator * (const Matrix & a, const Matrix & b) {
27
            Matrix ret(0);
28
             for (int k = 0; k < D; k++)
29
                 for(int i = 0; i < D; i++)if(a[i][k])</pre>
30
                     for (int j = 0; j < D; j++)
31
                         ret[i][j] = (ret[i][j] + 111 * a[i][k] * b[k][j]) % MOD;
32
             return ret;
33
34
35
36
    Matrix raiseMatrix (Matrix trans, 11 k) {
37
        Matrix res(0);
38
        res.initIdentity();
39
        for(;k;k>>=1,trans = trans * trans)
40
            if(k & 1)
41
                res = res * trans;
42
        return res;
43
```

7.4 Rabin Miller Primality check

```
2
    // n < 4,759,123,141
                                         3: 2, 7, 61
    // n < 1,122,004,669,633
                                         4: 2, 13, 23, 1662803
    // n < 3,474,749,660,383
                                         6 : pirmes <= 13
   // n < 3,825,123,056,546,413,051
                                       9 : primes <= 23
    int testPrimes[] = {2,3,5,7,11,13,17,19,23};
7
    struct MillerRabin{
10
     ///change K according to n
11
      const int K = 9;
12
      11 mult(11 s, 11 m, 11 mod){
13
        if(!m) return 0;
14
        11 \text{ ret} = \text{mult}(s, m/2, mod);
15
        ret = (ret + ret) % mod;
16
        if (m & 1) ret = (ret + s) % mod;
17
        return ret;
18
19
      11 power(ll x, ll p, ll mod){
21
        11 s = 1, m = x;
22
        while (p) {
23
          if(p&1) s = mult(s, m, mod);
          p >>= 1;
          m = mult(m, m, mod);
26
27
        return s;
28
29
      bool witness(ll a, ll n, ll u, int t) {
```

```
11 x = power(a, u, n), nx;
32
         for (int i = 0; i < t; i++) {
          nx = mult(x, x, n);
33
34
           if (nx == 1 \text{ and } x != 1 \text{ and } x != n-1) return 1;
35
36
37
        return x != 1;
38
39
40
      bool isPrime(ll n){ // return 1 if prime, 0 otherwise
41
        if(n < 2) return 0;
42
        if(!(n&1)) return n == 2;
        for(int i = 0; i < K; i++)if(n == testPrimes[i])return 1;</pre>
43
44
        11 u = n-1; int t = 0;
45
46
        while (u&1) u >>= 1, t++; // n-1 = u*2^t
47
48
        for(int i = 0; i < K; i++) if(witness(testPrimes[i], n, u, t)) return 0;</pre>
49
50
51 }tester;
```

8 Strings

8.1 Aho-Corasick Mostafa

```
struct AC FSM {
    #define ALPHABET SIZE 26
         struct Node {
             int child[ALPHABET_SIZE], failure = 0, match_parent = -1;
             vector<int> match;
 8
            Node() {
 9
                 for (int i = 0; i < ALPHABET_SIZE; ++i)child[i] = -1;</pre>
10
11
        };
12
13
        vector<Node> a;
14
15
        AC_FSM() {
16
            a.push_back(Node());
17
18
19
        void construct_automaton(vector<string> &words) {
20
             for (int w = 0, n = 0; w < words.size(); ++w, n = 0) {
21
                 for (int i = 0; i < words[w].size(); ++i) {</pre>
22
                     if (a[n].child[words[w][i] - 'a'] == -1) {
                         a[n].child[words[w][i] - 'a'] = a.size();
24
                         a.push_back(Node());
25
26
                     n = a[n].child[words[w][i] - 'a'];
27
                 a[n].match.push_back(w);
29
30
             queue<int> q;
31
             for (int k = 0; k < ALPHABET_SIZE; ++k) {</pre>
32
                 if (a[0].child[k] == -1) a[0].child[k] = 0;
33
                 else if (a[0].child[k] > 0) {
34
                     a[a[0].child[k]].failure = 0;
35
                     q.push(a[0].child[k]);
36
37
38
             while (!q.empty()) {
39
                 int r = q.front();
40
41
                 for (int k = 0, arck; k < ALPHABET_SIZE; ++k) {</pre>
42
                     if ((arck = a[r].child[k]) != -1) {
43
                         q.push(arck);
                         int v = a[r].failure;
```

```
while (a[v].child[k] == -1) v = a[v].failure;
                         a[arck].failure = a[v].child[k];
                         a[arck].match_parent = a[v].child[k];
                         while (a[arck].match parent != -1 &&
49
                               a[a[arck].match_parent].match.empty())
50
                             a[arck].match_parent =
51
                                     a[a[arck].match_parent].match_parent;
53
                }
54
55
56
57
        void aho_corasick(string &sentence, vector<string> &words,
58
                          vector<vector<int> > &matches) {
59
            matches.assign(words.size(), vector<int>());
60
            int state = 0, ss = 0;
            for (int i = 0; i < sentence.length(); ++i, ss = state) {</pre>
61
62
                while (a[ss].child[sentence[i] - 'a'] == -1)
63
                    ss = a[ss].failure;
64
                state = a[state].child[sentence[i] - 'a'] = a[ss].child[sentence[i]
65
                for (ss = state; ss != -1; ss = a[ss].match_parent)
                    for (int w: a[ss].match)
67
                        matches[w].push_back(i + 1 - words[w].length());
68
69
70
    };
```

8.2 Aho-Corasick Anany

```
int trie[N][A];
   int go[N][A]; //holds the node that you will go to after failure and stuff
    int ptr;
    11 ans[N]; //this node is a string terminator;
    int fail[N]: ///the failure function for each
    void BFS() {
7
        queue<int> q;
8
        f(i,0,A) {
Q
            if(trie[0][i]) {
10
                q.push(trie[0][i]);
11
                fail[trie[0][i]] = 0;
12
13
            go[0][i] = trie[0][i];
14
15
16
        while(q.size()) {
            auto node = q.front();
17
18
19
            ans[node] += ans[fail[node]]; //propagate fail[i] to ans[i]
20
            for (int i = 0; i < A; i++) {
21
                if(trie[node][i]) { ///calculate failure for you child
22
                    int to = trie[node][i];
23
                    int cur = fail[node]; ///int g = pi[i-1]
24
                    while(cur && !trie[cur][i]) ///while(q && s[q] != s[i])
25
                        cur = fail[cur]; ///g = pi[g-1]
26
                    if(trie[cur][i])cur = trie[cur][i]; ///g += s[i] == s[g]
                    fail[to] = cur; //pi[i] = g
28
                    q.push(to);
                    go[node][i] = trie[node][i];
30
                    else {
31
                    go[node][i] = go[fail[node]][i];
32
33
34
35
    void ins(string s, ll val) {
37
        int cur = 0;
38
        string sx = "";
39
        for(char c : s) {
40
            sx.push_back(c);
            if(!trie[cur][c - 'a']) {
41
                trie[cur][c - 'a'] = ++ptr;
```

8.3 KMP Anany

```
1 vector<int> fail(string s) {
        int n = s.size();
 3
        vector<int> pi(n);
        for(int i = 1; i < n; i++) {</pre>
            int q = pi[i-1];
            while (g \&\& s[i] != s[g])
              g = pi[g-1];
            q += s[i] == s[q];
            pi[i] = g;
11
        return pi;
13
    vector<int> KMP(string s, string t) {
14
        vector<int> pi = fail(t);
        vector<int> ret;
15
16
        for(int i = 0, g = 0; i < s.size(); i++) {</pre>
17
            while (g && s[i] != t[g])
               g = pi[g-1];
19
            q += s[i] == t[q];
20
            if(g == t.size()) { ///occurrence found
21
                ret.push_back(i-t.size()+1);
22
                q = pi[q-1];
25
        return ret;
26
```

8.4 Manacher Kactl

```
1 // If the size of palindrome centered at i is x, then d1[i] stores (x+1)/2.
 3
    vector<int> d1(n);
    for (int i = 0, l = 0, r = -1; i < n; i++) {
        int k = (i > r) ? 1 : min(d1[1 + r - i], r - i + 1);
        while (0 \le i - k \&\& i + k \le n \&\& s[i - k] == s[i + k]) {
 9
        d1[i] = k--;
        if(i + k > r) {
10
           1 = i - k;
11
            r = i + k;
12
13
14
15
16
17
    // If the size of palindrome centered at i is x, then d2[i] stores x/2
19
   vector<int> d2(n);
    for (int i = 0, l = 0, r = -1; i < n; i++) {
        int k = (i > r) ? 0 : min(d2[1 + r - i + 1], r - i + 1);
        while (0 \le i - k - 1 \&\& i + k \le n \&\& s[i - k - 1] == s[i + k]) {
24
25
        d2[i] = k--;
        if (i + k > r) {
            1 = i - k - 1;
            r = i + k;
29
30
```

```
struct SuffixArray {
         using vi = vector<int>;
         #define rep(i,a,b) for(int i = a; i < b; i++)
 3
 5
             Note this code is considers also the empty suffix
 6
             so hear sa[0] = n and sa[1] is the smallest non empty suffix
             and sa[n] is the largest non empty suffix
             also LCP[i] = LCP(sa[i-1], sa[i]), meanining LCP[0] = LCP[1] = 0
             if you want to get LCP(i..j) you need to build a mapping between
10
             sa[i] and i, and build a min sparse table to calculate the minimum
             note that this minimum should consider sa[i+1...j] since you don't want
11
12
             to consider LCP(sa[i], sa[i-1])
13
14
             you should also print the suffix array and lcp at the beginning of the
15
             to clarify this stuff
16
17
         vi sa, lcp;
18
        SuffixArray(string& s, int lim=256) { // or basic_string<int>
            int n = sz(s) + 1, k = 0, a, b;
19
20
             vi x(all(s)+1), y(n), ws(max(n, lim)), rank(n);
21
             sa = lcp = y, iota(all(sa), 0);
             for (int j = 0, p = 0; p < n; j = max(1, j * 2), lim = p) {
                p = j, iota(all(y), n - j);
rep(i,0,n) if (sa[i] >= j) y[p++] = sa[i] - j;
23
24
25
                 fill(all(ws), 0);
26
                 rep(i,0,n) ws[x[i]]++;
                 rep(i,1,lim) ws[i] += ws[i - 1];
27
                 for (int i = n; i--;) sa[--ws[x[y[i]]]] = y[i];
29
                 swap(x, y), p = 1, x[sa[0]] = 0;
30
                 rep(i,1,n) = sa[i-1], b = sa[i], x[b] =
31
                     (y[a] == y[b] && y[a + j] == y[b + j]) ? p - 1 : p++;
32
33
             rep(i,1,n) rank[sa[i]] = i;
34
             for (int i = 0, j; i < n - 1; lcp[rank[i++]] = k)
35
                 for (k \&\& k--, j = sa[rank[i] - 1];
36
                        s[i + k] == s[j + k]; k++);
37
    };
```

8.6 Suffix Automaton Anany

```
///Note it's better to use addNode to clear a node before using it
   ///at the start of each test case use initAutomaton
3
    int last = 0, cntState = 1;
    int nxt[N * 2][26];
    int len[N * 2], link[N * 2], firstPos[N * 2], cnt[N * 2];
    void addNode(int i) {
        memset(nxt[i], 0, sizeof nxt[i]);
10
        link[i] = -1;
11
        cnt[i] = 0;
12
13
14
    void initAutomaton() {
15
        cntState = 1;
16
        last = 0:
17
        addNode(last);
18
19
20
   int addChar(char c) {
        c -= 'a'; ///note this offset
23
        int p = last;
24
        int cur = cntState++;
25
        addNode(cur);
26
        cnt[cur] = 1; ///extra
        len[cur] = len[last] + 1;
```

```
firstPos[cur] = len[cur] - 1; ///extra
29
        while (p != -1 && nxt[p][c] == 0) {
30
            nxt[p][c] = cur;
31
            p = link[p];
32
33
34
        if(p == -1) {
35
            link[cur] = 0;
36
37
            int q = nxt[p][c];
38
            if(len[q] == len[p] + 1) {
39
                link[cur] = q;
40
               else {
                int clone = cntState++;
42
                link[clone] = link[q];
                firstPos[clone] = firstPos[q]; ///extra
                len[clone] = len[p] + 1;
                link[q] = link[cur] = clone;
45
                memcpy(nxt[clone], nxt[q], sizeof nxt[q]);
47
                cnt[clone] = 0; ///extra
                f(i,0,26)nxt[clone][i] = nxt[q][i];
49
                while (p != -1 && nxt[p][c] == q) {
                    nxt[p][c] = clone;
                    p = link[p];
52
53
55
        last = cur:
56
        return cur;
```

8.7 Suffix Automaton Mostafa

```
#include <bits/stdc++.h>
    #define FIO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    typedef long long 11;
    typedef long double ld;
    const int N = 2e6 + 9, M = 5e5 + 9;
    struct SA {
10
        struct node {
11
            int to[26];
12
            int link, len, co = 0;
13
14
            node() {
15
                memset(to, 0, sizeof to);
16
                co = 0, link = 0, len = 0;
17
18
        };
19
20
        int last, sz;
21
        vector<node> v;
22
23
            v = vector<node>(1);
25
            last = 0, sz = 1;
26
27
28
        void add_letter(int c) {
29
            int p = last;
30
            last = sz++;
31
            v.push_back({});
32
            v[last].len = v[p].len + 1;
33
            v[last].co = 1;
34
            for (; v[p].to[c] == 0; p = v[p].link)
35
                v[p].to[c] = last;
36
            if (v[p].to[c] == last) {
37
                v[last].link = 0;
                return;
```

```
int q = v[p].to[c];
41
             if (v[q].len == v[p].len + 1) {
                v[last].link = q;
43
                 return;
44
45
             int cl = sz++;
46
             v.push_back(v[q]);
47
             v.back().co = 0;
48
             v.back().len = v[p].len + 1;
49
             v[last].link = v[q].link = cl;
50
51
             for (; v[p].to[c] == q; p = v[p].link)
52
                v[p].to[c] = cl;
53
54
55
        void build co() {
56
             priority_queue<pair<int, int>> q;
57
             for (int i = sz - 1; i > 0; i--)
58
                 q.push({v[i].len, i});
59
             while (q.size()) {
60
                 int i = q.top().second;
61
                q.pop();
62
                v[v[i].link].co += v[i].co;
63
64
65
    };
66
67
    int main() {
68
69
70
         return 0;
71
```

8.8 Suffix Automaton With Rollback Mostafa

```
#include <bits/stdc++.h>
    #define FIO ios_base::sync_with_stdio(0); cin.tie(0); cout.tie(0);
    using namespace std;
    typedef long long 11;
    typedef long double 1d;
    const int N = 2e6 + 9, M = 5e5 + 9;
 9
    struct SA {
10
         struct node {
11
             int to[26];
12
             int link, len, co = 0;
13
14
             node() {
15
                 memset(to, 0, sizeof to);
16
                 co = 0, link = 0, len = 0;
17
18
        };
19
20
         struct LogNode {
21
             int last, sz;
22
             vector<pair<int, int>, int>> edges;
23
             pair<int, int> LinksUpdate = {0, 0};
24
25
26
         int last, sz;
\frac{27}{28}
        vector<node> v;
        vector<LogNode> logs;
29
30
31
             v = vector<node>(1);
32
             last = 0, sz = 1;
33
34
35
         void add_letter(int c) {
36
             logs.push_back({});
             logs.back().last = last;
```

```
logs.back().sz = sz;
39
40
             int p = last;
41
             last = sz++;
42
             v.push_back({});
43
             v[last].len = v[p].len + 1;
44
             v[last].co = 1;
45
             for (; v[p].to[c] == 0; p = v[p].link) {
46
                 logs.back().edges.push_back({{p, c}, 0});
47
                 v[p].to[c] = last;
48
49
             if (v[p].to[c] == last) {
50
                 v[last].link = 0;
51
                 return;
52
53
             int q = v[p].to[c];
54
             if (v[q].len == v[p].len + 1) {
55
                 v[last].link = q;
56
                 return;
57
             int c1 = sz++;
58
59
             v.push_back(v[q]);
60
             v.back().co = 0;
61
             v.back().len = v[p].len + 1;
62
             logs.back().LinksUpdate = {q, v[q].link};
v[last].link = v[q].link = cl;
63
64
             for (; v[p].to[c] == q; p = v[p].link) {
                 logs.back().edges.push_back({{p, c}, q});
65
66
                 v[p].to[c] = cl;
67
68
69
        void rollback() {
70
             assert(logs.size());
71
             auto log = logs.back();
72
             while (v.size() > log.sz)
73
                 v.pop_back();
             for (auto edge: log.edges)
74
                 v[edge.first.first].to[edge.first.second] = edge.second;
75
76
             if (log.LinksUpdate.first != 0)
                 v[log.LinksUpdate.first].link = log.LinksUpdate.second;
78
             last = log.last;
79
             sz = log.sz;
80
             logs.pop_back();
81
82
    };
83
84
    int main() {
85
        FIO
86
87
         return 0;
```

8.9 Zalgo Anany

```
1 int z[N], n;
    void Zalgo(string s) {
         int L = 0, R = 0;
         for (int i = 1; i < n; i++) {
 5
             if(i<=R&&z[i-L] < R - i + 1)z[i] = z[i-L];</pre>
 6
             else {
                 L = i;
                 R = max(R, i);
 Q
                 while (R < n \&\& s[R-L] == s[R])R++;
10
                 z[i] = R-L; --R;
11
12
13
```

9 Trees

9.1 Centroid Decomposition

```
1
2
            1. consider path(a,b) can be decomposed to path(a,lca(a,b)) and path(b,
            where lca(a,b) is the lca on the centroid tree
            2. Each one of the n^2 paths is the concatenation of two paths in a set
            paths from a node to all its ancestors in the centroid decomposition.
            3. Ancestor of a node in the original tree is either an ancestor in the
            a descendadnt
   vector<int> adj[N]; ///adjacency list of original graph
11
    int sz[N];
   bool used[N];
    int centPar[N]; //parent in centroid
    void init(int node, int par) { ///initialize size
16
        sz[node] = 1;
17
        for(auto p : adj[node])
18
            if(p != par && !used[p]) {
19
               init(p, node);
20
                sz[node] += sz[p];
21
    ^{24}
        for(int p : adj[node])
25
            if(!used[p] && p != par && sz[p] * 2 > limit)
26
            return centroid(p, node, limit);
27
        return node:
28
29
    int decompose(int node) {
        init(node, node); ///calculate size
31
        int c = centroid(node, node, sz[node]); //get centroid
32
        used[c] = true;
33
        for(auto p : adj[c])if(!used[p.F]) {      ///initialize parent for others and
            decompose
            centPar[decompose(p.F)] = c;
35
36
        return c;
37
    void update(int node, int distance, int col) {
        int centroid = node;
40
        while (centroid) {
41
            ///solve
42
            centroid = centPar[centroid];
45
    int query(int node) {
47
        int ans = 0;
48
49
        int centroid = node;
        while(centroid) {
51
            ///solve
52
            centroid = centPar[centroid];
53
55
        return ans;
```

9.2 Dsu On Trees

```
1 const int N = 1e5 + 9;
2 vector<int> adj[N];
3 int bigChild[N], sz[N];
```

```
void dfs(int node, int par) {
        for(auto v : adj[node]) if(v != par){
            dfs(v, node);
            sz[node] += sz[v];
            if(!bigChild[node] || sz[v] > sz[bigChild[node]]) {
                bigChild[node] = v;
10
11
    void add(int node, int par, int bigChild, int delta) {
14
        ///modify node to data structure
        for(auto v : adj[node])
        if(v != par && v != bigChild)
19
            add(v, node, bigChild, delta);
20
    void dfs2(int node, int par, bool keep) {
        for(auto v : adj[node])if(v != par && v != bigChild[node]) {
            dfs2(v, node, 0);
        if(bigChild[node]) {
            dfs2(bigChild[node], node, true);
29
        add(node, par, bigChild[node], 1);
30
        ///process queries
31
        if(!keep) {
            add(node, par, -1, -1);
33
34
```

9.3 Heavy Light Decomposition (Along with Euler Tour)

```
Notes:
            2. solve function iterates over segments and handles them seperatly
            if you're gonna use it make sure you know what you're doing
            3. to update/query segment in[node], out[node]
            4. to update/query chain in[nxt[node]], in[node]
            nxt[node]: is the head of the chain so to go to the next chain node =
 9
10
    int sz[mxN], nxt[mxN];
    int in[N], out[N], rin[N];
11
    vector<int> g[mxN];
    int par[mxN];
15
    void dfs_sz(int v = 0, int p = -1) {
16
        sz[v] = 1;
        par[v] = p;
17
18
        for (auto &u : g[v]) {
19
            if (u == p) {
20
                swap(u, g[v].back());
21
            if(u == p) continue;
23
            dfs_sz(u,v);
            sz[v] += sz[u];
            if (sz[u] > sz[g[v][0]])
                swap(u, g[v][0]);
28
        if(v != 0)
29
            g[v].pop_back();
30
    void dfs_hld(int v = 0) {
33
        in[v] = t++;
        rin[in[v]] = v;
35
        for (auto u : g[v]) {
            nxt[u] = (u == g[v][0] ? nxt[v] : u);
            dfs_hld(u);
```

```
15
```

```
39
        out[v] = t;
41
42
    int n;
43
    bool isChild(int p, int u) {
44
      return in[p] <= in[u] && out[u] <= out[p];</pre>
45
46
    int solve(int u,int v) {
47
        vector<pair<int,int> > segu;
48
         vector<pair<int,int> > seqv;
49
         if(isChild(u,v)){
50
          while(nxt[u] != nxt[v]){
51
             seqv.push_back(make_pair(in[nxt[v]], in[v]));
52
             v = par[nxt[v]];
53
54
           segv.push_back({in[u], in[v]});
         } else if(isChild(v,u)){
55
           while(nxt[u] != nxt[v]){
57
           segu.push_back(make_pair(in[nxt[u]], in[u]));
58
          u = par[nxt[u]];
59
60
          segu.push_back({in[v], in[u]});
61
      } else {
62
          while (u != v) {
63
            if(nxt[u] == nxt[v]) {
64
               if(in[u] < in[v]) seqv.push_back({in[u],in[v]}), R.push_back({u+1,v})</pre>
65
               else sequ.push_back({in[v],in[u]}), L.push_back({v+1,u+1});
66
              u = v:
67
               break;
68
             } else if(in[u] > in[v]) {
69
               segu.push_back({in[nxt[u]],in[u]}), L.push_back({nxt[u]+1, u+1});
70
               u = par[nxt[u]];
71
72
               segv.push_back({in[nxt[v]],in[v]}), R.push_back({nxt[v]+1, v+1});
73
               v = par[nxt[v]];
74
75
76
77
         reverse(segv.begin(), segv.end());
78
         int res = 0, state = 0;
79
         for(auto p : sequ) {
80
             gry(1,1,0,n-1,p.first,p.second,state,res);
81
82
         for(auto p : segv) {
83
             qry(0,1,0,n-1,p.first,p.second,state,res);
84
85
         return res;
86
```

9.4 LCA

```
const int N = 1e5 + 5;
    const int LG = 18;
4
    vector<int> adj[N];
5
    int pa[N][LG], lvl[N];
    int in[N], out[N], timer;
    void dfs(int u, int p) {
      in[u] = ++timer;
      for (int k = 1; k < LG; k++)
10
        pa[u][k] = pa[pa[u][k-1]][k-1];
11
      for(auto v : adj[u])
12
        if(v != p){
           lvl[v] = lvl[u] + 1;
13
14
           pa[v][0] = u;
15
           dfs(v, u);
16
17
      out[u] = timer;
18
    int LCA(int u, int v) {
```

```
if(lvl[u] > lvl[v])
21
        swap(u,v);
22
       int d = lvl[v] - lvl[u];
       for (int k = 0; k < LG; k++)
24
        if(d >> k \& 1)
25
          v = pa[v][k];
26
       if (u == v) return u;
       for(int i = LG - 1; i >= 0; --i)
28
        if(pa[u][i] != pa[v][i]){
29
          u = pa[u][i];
30
          v = pa[v][i];
31
32
      return pa[u][0];
33
```

9.5 Mo on Trees

```
1 int BL[N << 1], ID[N << 1];</pre>
  int lvl[N], par[17][N];
    int ans[N];
    vector<ii> adj[N];
    struct query{
      int id, 1, r, 1c;
      bool operator < (const query & rhs) {</pre>
        return (BL[1] == BL[rhs.1]) ? (r < rhs.r) : (BL[1] < BL[rhs.1]);</pre>
 9
10
    } Q [ N ] ;
11
    int in[N], out[N], val[N], timer;
    void dfs(int node, int p) {
      in[node] = ++timer; ID[timer] = node;
14
      for(int i = 1; i < 17; i++)par[i][node] = par[i-1][par[i-1][node]];</pre>
15
      for(auto child : adj[node])if(child.F != p){
16
        lvl[child.F] = lvl[node] + 1;
17
        par[0][child.F] = node;
18
        val[child.F] = child.S;
19
        dfs(child.F, node);
20
21
      out[node] = ++timer; ID[timer] = node;
22
   int LCA(int u, int v) {
      if(lvl[u] > lvl[v])swap(u,v);
25
      for (int k = 0; k < 17; k++)
        if((lvl[v] - lvl[u]) >> k & 1)
27
         v = par[k][v];
      if(u == v)
29
        return u;
30
      for (int i = 16; i >= 0; --i)
31
        if(par[i][u] != par[i][v])
32
          u = par[i][u], v = par[i][v];
33
      return par[0][u];
34
35 bool vis[N];
    int inSet[N];
37
    void add(int node, int & res){
      if(val[node] > N) return;
39
      if(!vis[node]){
40
        inSet[val[node]]++;
        while(inSet[res])res++;
41
42
43
        inSet[val[node]]--;
44
        if(!inSet[val[node]] && val[node] < res)</pre>
45
          res = val[node];
46
47
      vis[node] ^= 1;
48
49
    //-----Adding Queries-----/
50
51
        int u, v;
        cin >> u >> v; if(lvl[u] > lvl[v])swap(u, v);
52
53
        int lca = LCA(u, v);
        Q[i].id = i;
        Q[i].lc = lca;
```

```
Faculty of Computer and Information Sciences,
Ain Shams University: Too Wrong to Pass Too Correct to Fail
```