```
rep(i,0,n) if (!lfound[i]) cover.push back(i);
rep(i,0,m) if (seen[i]) cover.push_back(n+i);
assert(sz(cover) == res);
return cover;
```

GeneralMatching.h

**Description:** Matching for general graphs. Fails with probability N/mod.

Time:  $\mathcal{O}(N^3)$ "../numerical/MatrixInverse-mod.h" vector<pii> generalMatching(int N, vector<pii>& ed) { vector<vector<ll>> mat(N, vector<ll>(N)), A; for (pii pa : ed) { int a = pa.first, b = pa.second, r = rand() % mod; mat[a][b] = r, mat[b][a] = (mod - r) % mod;int r = matInv(A = mat), M = 2\*N - r, fi, f; assert(r % 2 == 0);**if** (M != N) **do** { mat.resize(M, vector<ll>(M)); rep(i,0,N) { mat[i].resize(M); rep(j,N,M) { int r = rand() % mod; mat[i][j] = r, mat[j][i] = (mod - r) % mod;} while (matInv(A = mat) != M); vi has(M, 1); vector<pii> ret; rep(it, 0, M/2) { rep(i,0,M) **if** (has[i]) rep(j,i+1,M) if (A[i][j] && mat[i][j]) { fi = i; fj = j; goto done; } assert(0); done: if (fj < N) ret.emplace\_back(fi, fj);</pre> has[fi] = has[fj] = 0;rep(sw,0,2) { 11 a = modpow(A[fi][fi], mod-2); rep(i,0,M) if (has[i] && A[i][fj]) { ll b = A[i][fj] \* a % mod;rep(j, 0, M) A[i][j] = (A[i][j] - A[fi][j] \* b) % mod;swap(fi,fj); return ret;

# DFS algorithms

Description: Calculates a valid assignment to boolean variables a, b, c,... to a 2-SAT problem, so that an expression of the type (a||b)&&(!a||c)&&(d||!b)&&... becomes true, or reports that it is unsatisfiable. Negated variables are represented by bit-inversions ( $\sim x$ ).

```
Usage: TwoSat ts(number of boolean variables);
ts.either(0, \sim3); // Var 0 is true or var 3 is false
ts.setValue(2); // Var 2 is true
ts.atMostOne(\{0, \sim 1, 2\}); // <= 1 of vars 0, \sim 1 and 2 are true
ts.solve(); // Returns true iff it is solvable
ts.values[0..N-1] holds the assigned values to the vars
```

**Time:**  $\mathcal{O}(N+E)$ , where N is the number of boolean variables, and E is the number of clauses. 5f9706, 56 lines

```
struct TwoSat {
 int N;
 vector<vi> gr;
 vi values; // 0 = false, 1 = true
 TwoSat(int n = 0) : N(n), gr(2*n) {}
 int addVar() { // (optional)
   gr.emplace_back();
   gr.emplace back();
    return N++;
 void either(int f, int j) {
   f = \max(2*f, -1-2*f);
   j = \max(2 * j, -1 - 2 * j);
   gr[f].push_back(j^1);
    gr[j].push_back(f^1);
 void setValue(int x) { either(x, x); }
 void atMostOne(const vi& li) { // (optional)
   if (sz(li) <= 1) return;</pre>
   int cur = ~li[0];
   rep(i,2,sz(li)) {
     int next = addVar();
     either(cur, ~li[i]);
     either(cur, next);
     either(~li[i], next);
     cur = ~next;
    either(cur, ~li[1]);
 vi val, comp, z; int time = 0;
 int dfs(int i) {
   int low = val[i] = ++time, x; z.push_back(i);
   for(int e : gr[i]) if (!comp[e])
     low = min(low, val[e] ?: dfs(e));
   if (low == val[i]) do {
     x = z.back(); z.pop_back();
     comp[x] = low;
     if (values[x>>1] == -1)
       values[x>>1] = x&1;
    } while (x != i);
   return val[i] = low;
 bool solve() {
   values.assign(N, -1);
   val.assign(2*N, 0); comp = val;
   rep(i,0,2*N) if (!comp[i]) dfs(i);
   rep(i,0,N) if (comp[2*i] == comp[2*i+1]) return 0;
   return 1:
};
```

## EulerWalk.h

Description: Eulerian undirected/directed path/cycle algorithm. Input should be a vector of (dest, global edge index), where for undirected graphs, forward/backward edges have the same index. Returns a list of nodes in the Eulerian path/cycle with src at both start and end, or empty list if no cycle/path exists. To get edge indices back, add .second to s and ret. Time:  $\mathcal{O}(V+E)$ 

```
780b64, 15 lines
vi eulerWalk(vector<vector<pii>>& gr, int nedges, int src=0) {
 int n = sz(qr);
 vi D(n), its(n), eu(nedges), ret, s = {src};
 D[src]++; // to allow Euler paths, not just cycles
```

```
while (!s.emptv()) {
  int x = s.back(), y, e, &it = its[x], end = sz(gr[x]);
  if (it == end) { ret.push_back(x); s.pop_back(); continue; }
  tie(y, e) = gr[x][it++];
  if (!eu[e]) {
    D[x] --, D[y] ++;
    eu[e] = 1; s.push_back(y);
for (int x : D) if (x < 0 \mid \mid sz(ret) != nedges+1) return {};
return {ret.rbegin(), ret.rend()};
```

### 6.5 Heuristics

#### MaximalCliques.h

Description: Runs a callback for all maximal cliques in a graph (given as a symmetric bitset matrix; self-edges not allowed). Callback is given a bitset representing the maximal clique.

```
Time: \mathcal{O}\left(3^{n/3}\right), much faster for sparse graphs
```

b0d5b1, 12 lines

```
typedef bitset<128> B;
template<class F>
void cliques(vector<B>& eds, F f, B P = \simB(), B X={}, B R={}) {
 if (!P.any()) { if (!X.any()) f(R); return; }
  auto q = (P | X)._Find_first();
  auto cands = P & ~eds[q];
 rep(i,0,sz(eds)) if (cands[i]) {
   R[i] = 1;
    cliques (eds, f, P & eds[i], X & eds[i], R);
    R[i] = P[i] = 0; X[i] = 1;
```

#### MaximumClique.h

Description: Quickly finds a maximum clique of a graph (given as symmetric bitset matrix; self-edges not allowed). Can be used to find a maximum independent set by finding a clique of the complement graph.

Time: Runs in about 1s for n=155 and worst case random graphs (p=.90). Runs faster for sparse graphs.

f7c0bc, 49 lines

```
typedef vector<bitset<200>> vb;
struct Maxclique {
 double limit=0.025, pk=0;
 struct Vertex { int i, d=0; };
 typedef vector<Vertex> vv;
 wh e:
 vv V;
 vector<vi> C;
 vi qmax, q, S, old;
  void init(vv& r) {
    for (auto \& v : r) v.d = 0;
    for (auto& v : r) for (auto j : r) v.d += e[v.i][j.i];
    sort(all(r), [](auto a, auto b) { return a.d > b.d; });
    int mxD = r[0].d;
    rep(i, 0, sz(r)) r[i].d = min(i, mxD) + 1;
 void expand(vv& R, int lev = 1) {
    S[lev] += S[lev - 1] - old[lev];
    old[lev] = S[lev - 1];
    while (sz(R)) {
      if (sz(q) + R.back().d <= sz(qmax)) return;</pre>
      g.push_back(R.back().i);
      vv T;
      for(auto v:R) if (e[R.back().i][v.i]) T.push_back({v.i});
      if (sz(T)) {
        if (S[lev]++ / ++pk < limit) init(T);</pre>
        int \dot{\eta} = 0, mxk = 1, mnk = max(sz(qmax) - sz(q) + 1, 1);
        C[1].clear(), C[2].clear();
```

for (auto v : T) {